

MODULE *Pactus*

The specification of the Pactus consensus algorithm :  
<https://docs.pactus.org/protocol/consensus/protocol/>

EXTENDS *Integers, Sequences, FiniteSets, TLC*

CONSTANT

The maximum number of rounds, limiting the range of behaviors evaluated by TLC.

*MaxRound*,

The maximum number of change – proposer(CP)rounds, limiting the range of behaviors evaluated by TLC.

*MaxCPRound*,

The total number of nodes in the network, denoted as *N* in the protocol.

*N*,

The maximum number of faulty nodes in the network, denoted as *F* in the protocol.

*F*,

The indices of faulty nodes.

*FaultyNodes*

VARIABLES

The set of messages received by the network.

*network*,

The set of messages delivered to each replica.

*logs*,

The state of each replica in the consensus protocol.

*states*

Helper expressions for common values.

*ThreeFPlusOne*  $\triangleq (3 * F) + 1$

*TwoFPlusOne*  $\triangleq (2 * F) + 1$

*OneFPlusOne*  $\triangleq (1 * F) + 1$

A tuple containing all variables in the spec for ease of use in temporal conditions.

*vars*  $\triangleq \langle \textit{network}, \textit{logs}, \textit{states} \rangle$

ASSUME

Ensure the number of nodes is sufficient to tolerate the specified number of faults.

$\wedge N \geq \textit{ThreeFPlusOne}$

Ensure that *FaultyNodes* is a valid subset of node indices.

$\wedge \textit{FaultyNodes} \subseteq 0 .. N - 1$

Helper functions

Check if the replica is the proposer for this round.

The proposer starts with the first replica and moves to the next in the change – proposer phase.

*IsProposer(index)*  $\triangleq$   
 $\textit{states}[\textit{index}].\textit{round} \% N = \textit{index}$

*Check if a node is faulty.*  
 $IsFaulty(index) \triangleq index \in FaultyNodes$

*Returns a subset of bag where each element matches all criteria specified in params*  
 $SubsetOfMsgs(bag, params) \triangleq$   
 $\{i \in bag : \forall field \in DOMAIN \ params : i[field] = params[field]\}$

*Check if the node has received  $3f + 1$  PRECOMMIT votes for a proposal in the current round.*  
 $HasPreCommitAbsolute(index) \triangleq$   
 $Cardinality(SubsetOfMsgs(logs[index], [$   
 $type \mapsto "PRECOMMIT",$   
 $round \mapsto states[index].round])) \geq ThreeFPlusOne$

*Check if the node has received  $2f + 1$  PRECOMMIT votes for a proposal in the current round.*  
 $HasPreCommitQuorum(index) \triangleq$   
 $Cardinality(SubsetOfMsgs(logs[index], [$   
 $type \mapsto "PRECOMMIT",$   
 $round \mapsto states[index].round])) \geq TwoFPlusOne$

$CPHasPreVotesQuorum(index) \triangleq$   
 $Cardinality(SubsetOfMsgs(logs[index], [$   
 $type \mapsto "CP:PRE-VOTE",$   
 $round \mapsto states[index].round,$   
 $cp\_round \mapsto states[index].cp\_round])) \geq TwoFPlusOne$

$CPHasPreVotesQuorumForYes(index) \triangleq$   
 $Cardinality(SubsetOfMsgs(logs[index], [$   
 $type \mapsto "CP:PRE-VOTE",$   
 $round \mapsto states[index].round,$   
 $cp\_round \mapsto states[index].cp\_round,$   
 $cp\_val \mapsto 1])) \geq TwoFPlusOne$

$CPHasPreVotesQuorumForNo(index) \triangleq$   
 $Cardinality(SubsetOfMsgs(logs[index], [$   
 $type \mapsto "CP:PRE-VOTE",$   
 $round \mapsto states[index].round,$   
 $cp\_round \mapsto states[index].cp\_round,$   
 $cp\_val \mapsto 0])) \geq TwoFPlusOne$

$CPHasPreVotesMinorityForNo(index) \triangleq$   
 $Cardinality(SubsetOfMsgs(logs[index], [$   
 $type \mapsto "CP:PRE-VOTE",$   
 $round \mapsto states[index].round,$   
 $cp\_round \mapsto states[index].cp\_round,$   
 $cp\_val \mapsto 0])) \geq OneFPlusOne$

$CPHasPreVotesMinorityForYes(index) \triangleq$

$$\begin{aligned}
& \text{Cardinality}(\text{SubsetOfMsgs}(\text{logs}[\text{index}], [ \\
& \quad \text{type} \quad \mapsto \text{"CP:PRE-VOTE"}, \\
& \quad \text{round} \quad \mapsto \text{states}[\text{index}].\text{round}, \\
& \quad \text{cp\_round} \mapsto \text{states}[\text{index}].\text{cp\_round}, \\
& \quad \text{cp\_val} \quad \mapsto 1])) \geq \text{OneFPlusOne} \\
\text{CPHasPreVotesForYesAndNo}(\text{index}) & \triangleq \\
& \wedge \text{Cardinality}(\text{SubsetOfMsgs}(\text{logs}[\text{index}], [ \\
& \quad \text{type} \quad \mapsto \text{"CP:PRE-VOTE"}, \\
& \quad \text{round} \quad \mapsto \text{states}[\text{index}].\text{round}, \\
& \quad \text{cp\_round} \mapsto \text{states}[\text{index}].\text{cp\_round}, \\
& \quad \text{cp\_val} \quad \mapsto 0])) \geq 1 \\
& \wedge \text{Cardinality}(\text{SubsetOfMsgs}(\text{logs}[\text{index}], [ \\
& \quad \text{type} \quad \mapsto \text{"CP:PRE-VOTE"}, \\
& \quad \text{round} \quad \mapsto \text{states}[\text{index}].\text{round}, \\
& \quad \text{cp\_round} \mapsto \text{states}[\text{index}].\text{cp\_round}, \\
& \quad \text{cp\_val} \quad \mapsto 1])) \geq 1 \\
\text{CPHasOneMainVotesNoInPrvRound}(\text{index}) & \triangleq \\
& \text{Cardinality}(\text{SubsetOfMsgs}(\text{logs}[\text{index}], [ \\
& \quad \text{type} \quad \mapsto \text{"CP:MAIN-VOTE"}, \\
& \quad \text{round} \quad \mapsto \text{states}[\text{index}].\text{round}, \\
& \quad \text{cp\_round} \mapsto \text{states}[\text{index}].\text{cp\_round} - 1, \\
& \quad \text{cp\_val} \quad \mapsto 0])) > 0 \\
\text{CPHasOneMainVotesYesInPrvRound}(\text{index}) & \triangleq \\
& \text{Cardinality}(\text{SubsetOfMsgs}(\text{logs}[\text{index}], [ \\
& \quad \text{type} \quad \mapsto \text{"CP:MAIN-VOTE"}, \\
& \quad \text{round} \quad \mapsto \text{states}[\text{index}].\text{round}, \\
& \quad \text{cp\_round} \mapsto \text{states}[\text{index}].\text{cp\_round} - 1, \\
& \quad \text{cp\_val} \quad \mapsto 1])) > 0 \\
\text{CPAllMainVotesAbstainInPrvRound}(\text{index}) & \triangleq \\
& \text{Cardinality}(\text{SubsetOfMsgs}(\text{logs}[\text{index}], [ \\
& \quad \text{type} \quad \mapsto \text{"CP:MAIN-VOTE"}, \\
& \quad \text{round} \quad \mapsto \text{states}[\text{index}].\text{round}, \\
& \quad \text{cp\_round} \mapsto \text{states}[\text{index}].\text{cp\_round} - 1, \\
& \quad \text{cp\_val} \quad \mapsto 2])) \geq \text{TwoFPlusOne} \\
\text{CPOneFPlusOneMainVotesAbstainInPrvRound}(\text{index}) & \triangleq \\
& \text{Cardinality}(\text{SubsetOfMsgs}(\text{logs}[\text{index}], [ \\
& \quad \text{type} \quad \mapsto \text{"CP:MAIN-VOTE"}, \\
& \quad \text{round} \quad \mapsto \text{states}[\text{index}].\text{round}, \\
& \quad \text{cp\_round} \mapsto \text{states}[\text{index}].\text{cp\_round} - 1, \\
& \quad \text{cp\_val} \quad \mapsto 2])) \geq \text{OneFPlusOne} \\
\text{CPHasMainVotesQuorum}(\text{index}) & \triangleq
\end{aligned}$$

$$\begin{aligned} & \text{Cardinality}(\text{SubsetOfMsgs}(\text{logs}[\text{index}], [ \\ & \quad \text{type} \mapsto \text{"CP:MAIN-VOTE"}, \\ & \quad \text{round} \mapsto \text{states}[\text{index}].\text{round}, \\ & \quad \text{cp\_round} \mapsto \text{states}[\text{index}].\text{cp\_round}])) \geq \text{TwoFPlusOne} \end{aligned}$$

$$\begin{aligned} \text{CPHasMainVotesQuorumForNo}(\text{index}) &\triangleq \\ & \text{Cardinality}(\text{SubsetOfMsgs}(\text{logs}[\text{index}], [ \\ & \quad \text{type} \mapsto \text{"CP:MAIN-VOTE"}, \\ & \quad \text{round} \mapsto \text{states}[\text{index}].\text{round}, \\ & \quad \text{cp\_round} \mapsto \text{states}[\text{index}].\text{cp\_round}, \\ & \quad \text{cp\_val} \mapsto 0])) \geq \text{TwoFPlusOne} \end{aligned}$$

$$\begin{aligned} \text{CPHasMainVotesQuorumForYes}(\text{index}) &\triangleq \\ & \text{Cardinality}(\text{SubsetOfMsgs}(\text{logs}[\text{index}], [ \\ & \quad \text{type} \mapsto \text{"CP:MAIN-VOTE"}, \\ & \quad \text{round} \mapsto \text{states}[\text{index}].\text{round}, \\ & \quad \text{cp\_round} \mapsto \text{states}[\text{index}].\text{cp\_round}, \\ & \quad \text{cp\_val} \mapsto 1])) \geq \text{TwoFPlusOne} \end{aligned}$$

$$\begin{aligned} \text{CPHasDecideVotesForNo}(\text{index}) &\triangleq \\ & \text{Cardinality}(\text{SubsetOfMsgs}(\text{logs}[\text{index}], [ \\ & \quad \text{type} \mapsto \text{"CP:DECIDED"}, \\ & \quad \text{round} \mapsto \text{states}[\text{index}].\text{round}, \\ & \quad \text{cp\_val} \mapsto 0])) > 0 \end{aligned}$$

$$\begin{aligned} \text{CPHasDecideVotesForYes}(\text{index}) &\triangleq \\ & \text{Cardinality}(\text{SubsetOfMsgs}(\text{logs}[\text{index}], [ \\ & \quad \text{type} \mapsto \text{"CP:DECIDED"}, \\ & \quad \text{round} \mapsto \text{states}[\text{index}].\text{round}, \\ & \quad \text{cp\_val} \mapsto 1])) > 0 \end{aligned}$$

$$\begin{aligned} & \text{Check if the node has received a proposal in the current round.} \\ \text{HasProposal}(\text{index}) &\triangleq \\ & \text{Cardinality}(\text{SubsetOfMsgs}(\text{logs}[\text{index}], [ \\ & \quad \text{type} \mapsto \text{"PROPOSAL"}, \\ & \quad \text{round} \mapsto \text{states}[\text{index}].\text{round}])) > 0 \end{aligned}$$

$$\begin{aligned} \text{HasPrecommitted}(\text{index}) &\triangleq \\ & \text{Cardinality}(\text{SubsetOfMsgs}(\text{logs}[\text{index}], [ \\ & \quad \text{type} \mapsto \text{"PRECOMMIT"}, \\ & \quad \text{round} \mapsto \text{states}[\text{index}].\text{round}, \\ & \quad \text{index} \mapsto \text{index}])) = 1 \end{aligned}$$

$$\begin{aligned} & \text{Check if the node has received a block announce message in the current round.} \\ \text{HasBlockAnnounce}(\text{index}) &\triangleq \\ & \text{Cardinality}(\text{SubsetOfMsgs}(\text{logs}[\text{index}], [ \end{aligned}$$

$type \mapsto \text{"BLOCK-ANNOUNCE"},$   
 $round \mapsto states[index].round)) > 0$

Check if the block is committed.

A block is considered committed if a supermajority of non – faulty replicas announce the same block.

$IsCommitted \triangleq$   
 LET  $subset \triangleq SubsetOfMsgs(network, [type \mapsto \text{"BLOCK-ANNOUNCE"}])$   
 IN  $\wedge Cardinality(subset) \geq TwoFPlusOne$   
 $\wedge \forall m1, m2 \in subset : m1.round = m2.round$

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#### Network functions

Simulate a replica sending a message by appending it to the network

The message is delivered to the sender's log immediately.

$SendMsg(msg) \triangleq$   
 IF  $msg.cp\_round < MaxCPRound$  THEN  
 $\wedge network' = network \cup \{msg\}$   
 $\wedge logs' = [logs \text{ EXCEPT } ![msg.index] = logs[msg.index] \cup \{msg\}]$   
 ELSE  
 UNCHANGED  $\langle network, logs \rangle$

Deliver a message to the specified replica's log.

$DeliverMsg(index) \triangleq$   
 LET  $undeliveredMsgs \triangleq network \setminus logs[index]$   
 IN IF  $Cardinality(undeliveredMsgs) = 0$  THEN  
 UNCHANGED  $\langle vars \rangle$   
 ELSE  
 LET  $msg \triangleq \text{CHOOSE } x \in undeliveredMsgs : \text{TRUE}$   
 IN  
 $\wedge logs' = [logs \text{ EXCEPT } ![index] = logs[index] \cup \{msg\}]$   
 $\wedge \text{UNCHANGED } \langle states, network \rangle$

Broadcast a *PROPOSAL* message into the network.

$SendProposal(index) \triangleq$   
 $SendMsg([$   
 $type \mapsto \text{"PROPOSAL"},$   
 $round \mapsto states[index].round,$   
 $index \mapsto index,$   
 $cp\_round \mapsto 0,$   
 $cp\_val \mapsto 0])$

Broadcast *PRECOMMIT* votes into the network.

$SendPreCommitVote(index) \triangleq$   
 $SendMsg([$   
 $type \mapsto \text{"PRECOMMIT"},$   
 $round \mapsto states[index].round,$

$$\begin{aligned}
index &\mapsto index, \\
cp\_round &\mapsto 0, \\
cp\_val &\mapsto 0])
\end{aligned}$$

Broadcast *CP:PRE-VOTE* votes into the network.

$$\begin{aligned}
SendCPPreVote(index, cp\_val) &\triangleq \\
SendMsg([ & \\
type &\mapsto \text{"CP:PRE-VOTE"}, \\
round &\mapsto states[index].round, \\
index &\mapsto index, \\
cp\_round &\mapsto states[index].cp\_round, \\
cp\_val &\mapsto cp\_val])
\end{aligned}$$

Broadcast *CP:MAIN-VOTE* votes into the network.

$$\begin{aligned}
SendCPMainVote(index, cp\_val) &\triangleq \\
SendMsg([ & \\
type &\mapsto \text{"CP:MAIN-VOTE"}, \\
round &\mapsto states[index].round, \\
index &\mapsto index, \\
cp\_round &\mapsto states[index].cp\_round, \\
cp\_val &\mapsto cp\_val])
\end{aligned}$$

Broadcast *CP:DECIDED* votes into the network.

$$\begin{aligned}
SendCPDecideVote(index, cp\_val) &\triangleq \\
SendMsg([ & \\
type &\mapsto \text{"CP:DECIDED"}, \\
round &\mapsto states[index].round, \\
cp\_round &\mapsto states[index].cp\_round, \\
index &\mapsto 0, \text{ reduce the model size} \\
cp\_val &\mapsto cp\_val])
\end{aligned}$$

Broadcast *BLOCK-ANNOUNCE* messages into the network.

$$\begin{aligned}
AnnounceBlock(index) &\triangleq \\
SendMsg([ & \\
type &\mapsto \text{"BLOCK-ANNOUNCE"}, \\
round &\mapsto states[index].round, \\
index &\mapsto index, \\
cp\_round &\mapsto 0, \\
cp\_val &\mapsto 0])
\end{aligned}$$


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State transition functions

Transition to the propose state.

$$\begin{aligned}
Propose(index) &\triangleq \\
&\wedge \neg IsFaulty(index) \\
&\wedge states[index].name = \text{"propose"}
\end{aligned}$$

$\wedge$   
 IF  $IsProposer(index)$  THEN  
      $SendProposal(index)$   
 ELSE  
     UNCHANGED  $\langle logs, network \rangle$   
 $\wedge states' = [states \text{ EXCEPT } ![index].name = \text{"precommit"}]$

Transition to the *precommit* state.

$PreCommit(index) \triangleq$   
 $\wedge \neg IsFaulty(index)$   
 $\wedge states[index].name = \text{"precommit"}$   
 $\wedge HasProposal(index)$   
 $\wedge SendPreCommitVote(index)$   
 $\wedge states' = states$

Transition to the fast commit state.

$FastCommit(index) \triangleq$   
 $\wedge \neg IsFaulty(index)$   
 $\wedge states[index].name \neq \text{"commit"}$  to prevent shuttering  
 $\wedge HasPreCommitAbsolute(index)$   
 $\wedge states' = [states \text{ EXCEPT } ![index].name = \text{"commit"}]$   
 $\wedge$  UNCHANGED  $\langle network, logs \rangle$

Transition to the commit state.

$Commit(index) \triangleq$   
 $\wedge \neg IsFaulty(index)$   
 $\wedge states[index].name = \text{"commit"}$   
 $\wedge HasProposal(index)$   
 $\wedge HasPreCommitQuorum(index)$   
 $\wedge AnnounceBlock(index)$   
 $\wedge$  UNCHANGED  $\langle states \rangle$

Transition for timeout: a non-faulty replica changes the proposer if its timer expires.

$Timeout(index) \triangleq$   
 $\wedge \neg IsFaulty(index)$   
 $\wedge states[index].name = \text{"precommit"}$   
 $\wedge$   
     To limit the the behaviours.  
      $\vee states[index].round < MaxRound$   
      $\vee HasPreCommitQuorum(index)$   
 $\wedge states' = [states \text{ EXCEPT } ![index].name = \text{"cp:pre-vote"}]$   
 $\wedge$  UNCHANGED  $\langle network, logs \rangle$

Transition to the *CP* pre-vote state.

$CPPreVote(index) \triangleq$   
 $\wedge \neg IsFaulty(index)$

$$\begin{aligned}
& \wedge \text{states}[index].name = \text{"cp:pre-vote"} \\
& \wedge \\
& \quad \text{IF } \text{states}[index].cp\_round = 0 \text{ THEN} \\
& \quad \quad \text{IF } \text{HasPreCommitQuorum}(index) \text{ THEN} \\
& \quad \quad \quad \wedge \text{SendCPPreVote}(index, 0) \\
& \quad \quad \quad \wedge \text{states}' = [\text{states} \text{ EXCEPT } ![index].name = \text{"cp:main-vote"}] \\
& \quad \quad \text{ELSE IF } \neg \text{HasPrecommitted}(index) \text{ THEN} \\
& \quad \quad \quad \wedge \text{SendCPPreVote}(index, 1) \\
& \quad \quad \quad \wedge \text{states}' = [\text{states} \text{ EXCEPT } ![index].name = \text{"cp:main-vote"}] \\
& \quad \quad \text{ELSE IF } \vee \text{CPHasPreVotesMinorityForYes}(index) \\
& \quad \quad \quad \vee \text{Cardinality}( \\
& \quad \quad \quad \quad \text{SubsetOfMsgs}(\text{logs}[index], [type \mapsto \text{"PRECOMMIT"}, round \mapsto \text{states}[index].round} \\
& \quad \quad \quad \quad \text{SubsetOfMsgs}(\text{logs}[index], [type \mapsto \text{"CP:PRE-VOTE"}, round \mapsto \text{states}[index].round} \\
& \quad \quad \quad \quad ) \geq \text{TwoFPlusOne} \text{ THEN} \\
& \quad \quad \quad \wedge \text{SendCPPreVote}(index, 1) \\
& \quad \quad \quad \wedge \text{states}' = [\text{states} \text{ EXCEPT } ![index].name = \text{"cp:main-vote"}] \\
& \quad \quad \text{ELSE} \\
& \quad \quad \quad \wedge \text{UNCHANGED } \langle vars \rangle \\
& \quad \text{ELSE} \\
& \quad \quad \wedge \\
& \quad \quad \quad \vee \\
& \quad \quad \quad \quad \wedge \text{CPHasOneMainVotesNoInPrvRound}(index) \\
& \quad \quad \quad \quad \wedge \text{SendCPPreVote}(index, 0) \\
& \quad \quad \quad \quad \vee \\
& \quad \quad \quad \quad \wedge \text{CPHasOneMainVotesYesInPrvRound}(index) \\
& \quad \quad \quad \quad \wedge \text{SendCPPreVote}(index, 1) \\
& \quad \quad \quad \quad \vee \\
& \quad \quad \quad \quad \wedge \text{CPAllMainVotesAbstainInPrvRound}(index) \\
& \quad \quad \quad \quad \wedge \text{SendCPPreVote}(index, 0) \text{ biased to zero} \\
& \quad \quad \wedge \text{states}' = [\text{states} \text{ EXCEPT } ![index].name = \text{"cp:main-vote"}]
\end{aligned}$$

Transition to the *CP* main-vote state.

$$\begin{aligned}
& \text{CPMainVote}(index) \triangleq \\
& \quad \wedge \neg \text{IsFaulty}(index) \\
& \quad \wedge \text{states}[index].name = \text{"cp:main-vote"} \\
& \quad \wedge \text{CPHasPreVotesQuorum}(index) \\
& \quad \wedge \\
& \quad \quad \vee \\
& \quad \quad \quad \text{all votes for 0} \\
& \quad \quad \quad \wedge \text{CPHasPreVotesQuorumForNo}(index) \\
& \quad \quad \wedge \text{CPHasPreVotesMinorityForNo}(index) \text{ To reduce the behaviours.} \\
& \quad \quad \wedge \text{states}' = [\text{states} \text{ EXCEPT } ![index].name = \text{"commit"}] \\
& \quad \quad \wedge \text{UNCHANGED } \langle network, logs \rangle \\
& \quad \quad \vee \\
& \quad \quad \quad \text{all votes for 1}
\end{aligned}$$





$\wedge \text{UNCHANGED } \langle network, logs \rangle$

**Initial state**

$Init \triangleq$   
 $\wedge network = \{\}$   
 $\wedge logs = [index \in 0 \dots N - 1 \mapsto \{\}]$   
 $\wedge states = [index \in 0 \dots N - 1 \mapsto [$   
 $\quad name \mapsto \text{"propose"},$   
 $\quad round \mapsto 0,$   
 $\quad cp\_round \mapsto 0]]$

**State transition relation**

$Next \triangleq$   
 $\exists index \in 0 \dots N - 1 :$   
 $\vee Propose(index)$   
 $\vee PreCommit(index)$   
 $\vee Timeout(index)$   
 $\vee Commit(index)$   
 $\vee FastCommit(index)$   
 $\vee CPPreVote(index)$   
 $\vee CPMainVote(index)$   
 $\vee CPDecide(index)$   
 $\vee CPStrongTerminate(index)$   
 $\vee DeliverMsg(index)$

**Specification**

$Spec \triangleq$   
 $Init \wedge \Box [Next]_{vars} \wedge WF_{vars}(Next)$

Success: All non-faulty nodes eventually commit.

$Success \triangleq \Diamond(IsCommitted)$

TypeOK is the type-correctness invariant.

$TypeOK \triangleq$   
 $\wedge \forall index \in 0 \dots N - 1 :$   
 $\wedge states[index].round \leq MaxRound$   
 $\wedge states[index].cp\_round \leq MaxCPRound$   
 $\wedge states[index].name = \text{"propose"} \wedge states[index].round > 0 \Rightarrow$   
 $\wedge Cardinality(SubsetOfMsgs(network, [$   
 $\quad type \mapsto \text{"CP:DECIDED"},$   
 $\quad round \mapsto states[index].round - 1,$   
 $\quad cp\_val \mapsto 1])) = 1$   
 $\wedge Cardinality(SubsetOfMsgs(network, [$   
 $\quad type \mapsto \text{"BLOCK-ANNOUNCE"},$

$$\begin{aligned}
& \text{round} \mapsto \text{states}[\text{index}].\text{round} - 1]) = 0 \\
\wedge \text{states}[\text{index}].\text{name} = \text{"commit"} \Rightarrow \\
& \wedge \text{Cardinality}(\text{SubsetOfMsgs}(\text{network}, [ \\
& \quad \text{type} \mapsto \text{"PRECOMMIT"}, \\
& \quad \text{round} \mapsto \text{states}[\text{index}].\text{round}])) \geq \text{TwoFPlusOne} \\
& \wedge \text{Cardinality}(\text{SubsetOfMsgs}(\text{network}, [ \\
& \quad \text{type} \mapsto \text{"PROPOSAL"}, \\
& \quad \text{round} \mapsto \text{states}[\text{index}].\text{round}])) = 1 \\
& \wedge \text{LET } \text{subset} \triangleq \text{SubsetOfMsgs}(\text{network}, [\text{type} \mapsto \text{"BLOCK-ANNOUNCE"}]) \\
& \text{IN } \wedge \forall m1, m2 \in \text{subset} : m1.\text{round} = m2.\text{round}
\end{aligned}$$

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