

Model-theoretic embeddings

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- what do we actually mean by “semantics”?
 - formal definition of “truth” relies on “models”
 - universal algebra over formal languages (with signature Σ)

Name	Syntax	Semantics
top	\top	$\Delta^{\mathcal{I}}$
bottom	\perp	\emptyset
nominal	$\{a\}$	$\{a^{\mathcal{I}}\}$
conjunction	$C \sqcap D$	$C^{\mathcal{I}} \cap D^{\mathcal{I}}$
existential restriction	$\exists r.C$	$\{x \in \Delta^{\mathcal{I}} \mid \exists y \in \Delta^{\mathcal{I}} : (x, y) \in r^{\mathcal{I}} \wedge y \in C^{\mathcal{I}}\}$
generalized concept inclusion	$C \sqsubseteq D$	$C^{\mathcal{I}} \subseteq D^{\mathcal{I}}$
role inclusion	$r_1 \circ \dots \circ r_n \sqsubseteq r$	$r_1^{\mathcal{I}} \circ \dots \circ r_n^{\mathcal{I}} \subseteq r^{\mathcal{I}}$

- Interpretations and Σ -structures
- Model \mathfrak{A} of a formula ϕ : ϕ is true in \mathfrak{A} ($\mathfrak{A} \models \phi$)
- Theory T : set of formulas
- \mathfrak{A} is a model of T if \mathfrak{A} is a model of all formulas in T
- Ontologies are (special kinds of) theories

- given a theory/ontology T with signature $\Sigma(T)$
- aim: find $f_e : \Sigma(T) \mapsto \mathbb{R}^n$ s.t. $f_e(\Sigma(T))$ is a model of T
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- more general: find an algorithm that maps symbols (signatures) into \mathbb{R}^n so that the *semantics* of the symbol (expressed through axioms and explicit in model structures) is preserved
 - or: the embedding function *is* an interpretation function

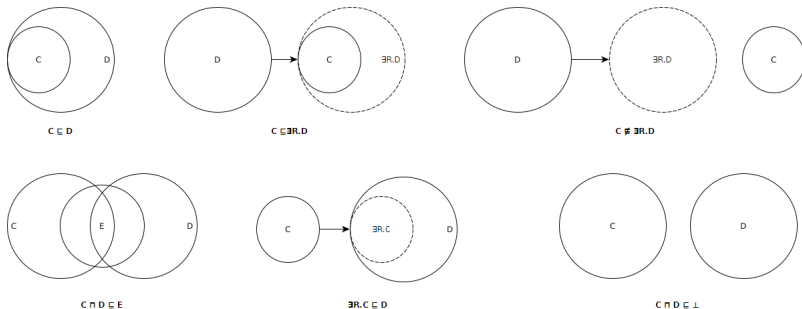
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- any consistent \mathcal{EL}^{++} theory has infinite models
- any consistent \mathcal{EL}^{++} theory has models in \mathbb{R}^n (Loewenheim-Skolem, upwards; compactness)

- for all $r \in \Sigma(T)$ and $C \in \Sigma(T)$, define $f_e(r)$ and $f_e(C)$
- $f_e(C)$ maps to points in an open n -ball such that $f_e(C) = C^{\mathcal{I}}$:
 $C^{\mathcal{I}} = \{x \in \mathbb{R}^n \mid \|f_e(C) - x\| < r_e(C)\}$
 - these are the *extension* of a class in \mathbb{R}^n
- $f_e(r)$ maps a binary relation r to a vector such that
 $r^{\mathcal{I}} = \{(x, y) \mid x + f_e(r) = y\}$
 - that's the TransE property for *individuals*
- use the axioms in T as constraints

- normalize the theory:
 - every \mathcal{EL}^{++} theory can be expressed using four normal forms (Baader et al., 2005)
- eliminate the ABox: replace each individual symbol with a singleton class: a becomes $\{a\}$
- rewrite relation assertions $r(a, b)$ and class assertions $C(a)$ as $\{a\} \sqsubseteq \exists r. \{b\}$ and $\{a\} \sqsubseteq C$
 - something to remember for the next class-vs-instance discussion?
- normalization rules to generate:
 - $C \sqsubseteq D$
 - $C \sqcap D \sqsubseteq E$
 - $C \sqsubseteq \exists R.D$
 - $\exists R.C \sqsubseteq D$

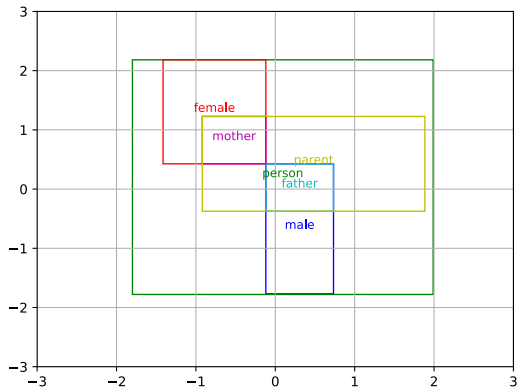
Algorithm: loss functions



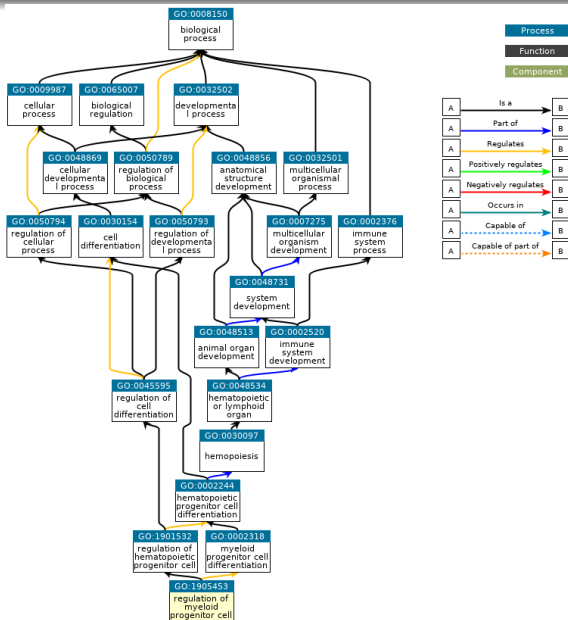
<i>Male</i>	\sqsubseteq <i>Person</i>	(1)
<i>Female</i>	\sqsubseteq <i>Person</i>	(2)
<i>Father</i>	\sqsubseteq <i>Male</i>	(3)
<i>Mother</i>	\sqsubseteq <i>Female</i>	(4)
<i>Father</i>	\sqsubseteq <i>Parent</i>	(5)
<i>Mother</i>	\sqsubseteq <i>Parent</i>	(6)
<i>Female</i> \sqcap <i>Male</i>	$\sqsubseteq \perp$	(7)
<i>Female</i> \sqcap <i>Parent</i>	\sqsubseteq <i>Mother</i>	(8)
<i>Male</i> \sqcap <i>Parent</i>	\sqsubseteq <i>Father</i>	(9)
$\exists hasChild. Person$	\sqsubseteq <i>Parent</i>	(10)
<i>Parent</i>	\sqsubseteq <i>Person</i>	(11)
<i>Parent</i>	$\sqsubseteq \exists hasChild. \top$	(12)

- model with $\Delta = R^n$
- support quantifiers, negation, conjunction,...

EL BOX Embeddings



Gene Ontology



QuickGO - <https://www.ebi.ac.uk/QuickGO>

