

Recall: we compute

$$2^{2016} \equiv 2^{360} \pmod{1739}$$

Approach 3 (Successive squaring):

Write 360 as a sum of powers of 2:

$$360 = 2^8 + 2^6 + 2^5 + 2^3$$

256    64    32    8

Then compute the sequence  $a_n \equiv 2^{2^n} \pmod{1739}$

Notice  $a_{n+1} \equiv a_n^2 \pmod{1739}$ .

$n$	0	1	2	3	4	5	6	7	8
$2^{2^n} \pmod{1739}$	2	4	16	256	1193	747	1529	625	1089

Compute

$$2^{360} = 2^{2^8} \cdot 2^{2^6} \cdot 2^{2^5} \cdot 2^{2^3} \equiv 1089 \cdot 1529 \cdot 747 \cdot 256$$

$$\equiv 858 \cdot 1681 \equiv 667 \pmod{1739}.$$

In total this method requires  $8+3=11$  (heavy) multiplications.

Approach 4: We work with residues modulo 37 and 47 separately and then use CRT.

$$\text{By FLT, } 2^{36} \equiv 1 \pmod{37} \Rightarrow 2^{360} = (2^{36})^{10} \equiv 1 \pmod{37}$$

$$2^{46} \equiv 1 \pmod{47} \Rightarrow 2^{360} = 2^{7 \cdot 46 + 38} \equiv 2^{38} \pmod{47}$$

We can use successive squarings. That will use  $5+2=7$  (light) multiplications.

$$2^{38} \equiv 9 \pmod{47} \quad (\text{Check!})$$

Now we solve the system

$$\begin{cases} x \equiv 1 \pmod{37} \\ x \equiv 9 \pmod{47} \end{cases}$$

Use EEA:

$$47 = 1 \cdot 37 + 10$$

$$37 = 3 \cdot 10 + 7$$

$$10 = 1 \cdot 7 + 3$$

$$7 = 2 \cdot 3 + 1$$

47	37	10	7	3	1
-	-	1	3	1	2
0	1	-1	4	5	14
1	0	1	-3	4	11

Finally,  $t = -11 \cdot 47 + 14 \cdot 37$ .

The answer then is  $2^{360} \equiv 1 \cdot (-11) \cdot 47 + 9 \cdot 14 \cdot 37$   
 $\equiv 4145 \equiv 667 \pmod{1739}$

## §9 Computation of $k$ 'th roots modulo a number.

Task: Given  $a \in \mathbb{Z}$ ,  $k, m \in \mathbb{Z}^+$ , find an integer  $x$  such that  $x^k \equiv a \pmod{m}$ .

Important restrictions:

(a)  $\gcd(k, \varphi(m)) = 1$ .

(b)  $\gcd(a, m) = 1$

(In some cases we can drop this restriction, in particular for  $m=p$  or  $m=pq$  where  $p \neq q$  are prime, see "RSA Theorem").

Example: Find cubic root of 18 (mod 11).

$$x^3 \equiv 18 \pmod{11} \equiv 7 \pmod{11}$$

Compose a table of cubes mod 11:

$x$	0	1	2	3	4	5	6	7	8	9	10
$x^3 \pmod{11}$	0	1	8	5	9	4	(7)	2	6	3	10

(note that  $(-x)^3 \equiv -x^3$ ).

From the table we see that there is the unique solution  $x \equiv 6 \pmod{11}$ .

Much more efficient method for computing  $k$ 'th roots:

(1) Compute  $\varphi(m)$

(2) Find integer  $s, t, s > 0$  such that  
 $1 = s \cdot k + t \cdot \varphi(m)$  (EEA)

(3) Compute

$x \equiv a^s \pmod{m}$  (successive squaring)

$$\begin{aligned} \text{(Indeed } a^1 &= a^{sk + t\varphi(m)} = (a^s)^k \cdot (a^{\varphi(m)})^t \\ &\equiv [E \cdot FT] \equiv (a^s)^k \pmod{m} \end{aligned}$$

In other words  $x \equiv a^s \pmod{m}$  satisfies  
 $x^k \equiv a \pmod{m}$ )

Remark: We just provide one solution, we do not show that it is unique. In fact, it is unique, but we will not show that.

Example:  $x^{101} \equiv 262 \pmod{667}$ .

(1) One can find that  $667 = 23 \cdot 29$ .

$$\text{Then } \varphi(667) = 22 \cdot 28 = 616$$

(2) Apply EEA,

$$616 = 6 \cdot 101 + 10$$

$$101 = 10 \cdot 10 + 1$$

$$\begin{aligned} \text{We can write } 1 &= 101 - 10 \cdot 10 = 101 - 10(616 - 6 \cdot 101) \\ &= 61 \cdot 101 - 10 \cdot 616 \end{aligned}$$

(3) Compute  $x \equiv 262^{61} \pmod{667}$ .

We apply successive squarings (or approach 4) to get

$$x \equiv 262^{61} \equiv 233 \pmod{667}. \quad (\text{Check!})$$