

§20 Square roots in modular arithmetics.

§20.1. The case of odd prime p .

Let p be an odd prime.

Definition. Let $a \in \mathbb{Z}$, $a \not\equiv 0 \pmod{p}$. a is a quadratic residue (QR) modulo p if the equation

$$x^2 \equiv a \pmod{p}$$

has solutions. Otherwise it is called quadratic non-residue (NR) modulo p .

Note: If $a \equiv 0 \pmod{p}$ then it is neither QR nor NR mod p .

If a is a QR mod p then $x^2 \equiv a \pmod{p}$ has exactly two solutions:

$$x \equiv \pm b \pmod{p}.$$

Q: How to check whether $a \in \{1, 2, \dots, p-1\}$ is a QR or not?

We can solve it with help of prim. roots and discrete logs.

Let g be a prim. root mod p .

$$x \equiv g^i \pmod{p}, \quad a \equiv g^k \pmod{p}$$

$$x^2 \equiv a \pmod{p} \Leftrightarrow \underset{\substack{\uparrow \\ \text{even}}}{2i} \equiv \underset{\substack{\uparrow \\ \text{even}}}{k} \pmod{p-1}$$

If k is odd the equation does not have solutions $\Rightarrow a \equiv g^k \pmod{p}$ is NR.

If $k=2m$ is even then

$$2i \equiv k \pmod{p-1} \Leftrightarrow i \equiv m \pmod{\frac{p-1}{2}}$$

$\Rightarrow x \equiv g^m$ or $g^{m+\frac{p-1}{2}} \pmod{p}$ are solutions

$\Rightarrow a \equiv g^{2m} \pmod{p}$ is a ~~prim. root~~. QR.

Problem: this method involves finding a prim. root and solving the DLP - very hard and practically useless for big p .

Lemma: a is a QR mod $p \Leftrightarrow a^{\frac{p-1}{2}} \equiv 1 \pmod{p}$

a is a NR mod $p \Leftrightarrow a^{\frac{p-1}{2}} \equiv -1 \pmod{p}$.

Proof.

a is a QR mod $p \Leftrightarrow a \equiv g^{2m} \pmod{p}$ where g is a prim. root.

$$\Rightarrow a^{\frac{p-1}{2}} \equiv (g^{2m})^{\frac{p-1}{2}} \equiv g^{m(p-1)} \equiv 1 \pmod{p}.$$

a is NR mod $p \Leftrightarrow a \equiv g^k \pmod{p}$, k is odd

$$\Rightarrow a^{\frac{p-1}{2}} \equiv (g^k)^{\frac{p-1}{2}} \equiv (g^{\frac{p-1}{2}})^k \equiv (-1)^k \equiv -1 \pmod{p}.$$

-1 (check-Ex).



Example: $p=7$. Is 3 a QR?

$$3^{\frac{p-1}{2}} = 3^3 = 6 \equiv -1 \pmod{7} \Rightarrow 3 \text{ is a NR.}$$

Q (Square Root Problem): Given $a \in \{1, \dots, p-1\}$ is a QR, solve the equation $x^2 \equiv a \pmod{p}$.

As before we can solve it with help of prim. roots and discrete logs. - too long. We want something faster.

Easy case: $p \equiv 3 \pmod{4}$. Then $\frac{p+1}{4} \in \mathbb{Z}$.

Lemma: Let p be prime, $p \equiv 3 \pmod{4}$ and a is a QR mod p . Then $a^{\frac{p+1}{4}}$ is a solution of $x^2 \equiv a \pmod{p}$.

Proof: $(a^{\frac{p+1}{4}})^2 \equiv a^{\frac{p+1}{2}} \equiv a^{\frac{p-1}{2}} \cdot a \equiv [a \text{ is a QR}] \equiv a \pmod{p}$. ☑

Example: $p=11$. Solve $x^2 \equiv 3 \pmod{11}$. $11 \equiv 3 \pmod{4}$

Check whether 3 is a QR:

$$3^{\frac{p-1}{2}} \equiv 3^5 \equiv 1 \pmod{11} \Rightarrow 3 \text{ is QR.}$$

~~THE~~ solution is: $3^{\frac{p+1}{4}} \equiv 3^3 \equiv 5 \pmod{11}$
one

$\Rightarrow 5$ is a square root of 3 mod 11.

The general solution is $x \equiv \pm 5 \pmod{11}$.

General case: Let $p-1 = 2^k \cdot m$ where $k \in \mathbb{Z}^+$, $m \in \mathbb{Z}$ is odd

(Easy case is when $k=1$).

Algorithm for finding of a square root of a modulo p :

Step 0: Check if a is QR by checking if $a^{\frac{p-1}{2}} \equiv 1 \pmod{p}$.

Step 1: Find $b \in \{1, 3, \dots, p-1\}$ such that $\text{ord}_p(b) = 2^k$.

Method: find a NR r modulo p by checking $r^{2^{k-1}m} \equiv -1 \pmod{p}$. Do this by random search.

There are $\frac{p-1}{2}$ NR's so we should find such r quickly.

Take $b \equiv r^m \pmod{p}$.

Check: $b^{2^k} \equiv r^{2^k m} \equiv r^{p-1} \equiv 1 \pmod{p}$

$$b^{2^{k-1}} \equiv r^{\frac{p-1}{2}} \equiv -1 \pmod{p}$$

$$\Rightarrow \text{ord}_p(b) \mid 2^k \text{ but } \text{ord}_p(b) \nmid 2^{k-1}$$

$$\Rightarrow \text{ord}_p(b) = 2^k.$$

Step 2: We have that the numbers

$b^0, b^2, b^4, \dots, b^{2^{k-2}}$ are (2^{k-1}) roots of 1.

\Rightarrow they are all roots of 1 of degree 2^{k-1} .

On the other hand a^m is also (2^{k-1}) th root of 1

$$\Rightarrow \exists j \text{ s.t. } b^{2^j} \equiv a^m \pmod{p}, \quad j \in \{0, 1, \dots, 2^{k-1}-1\}$$

In other words $j = \frac{1}{2} \log_{b,p}(a^m)$.

Find this j (By Pohling-Hellman).

This is easy because the order is 2^{k+m} .

Step 3:

$$x \equiv \pm b^j \cdot a^{-(m-1)} \pmod{p}.$$