

CS 162a, Home work 1

Brian
Lester

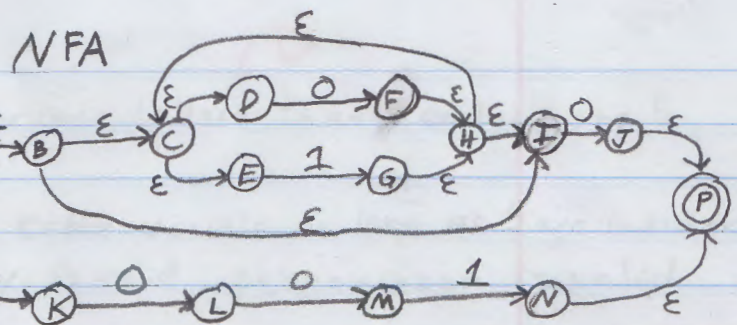
1) $((011)^*0)1(001)$

NFA State
Table

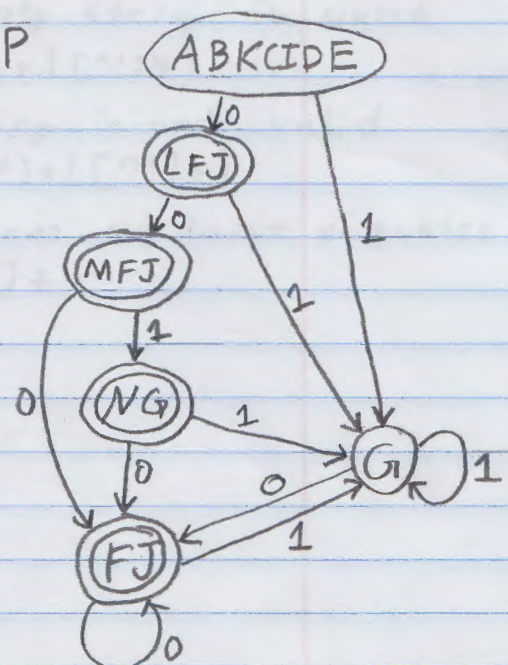
	0	1	ϵ
A	LFJ	G	BKIDE
B	FJ	G	CIDE
C	F	G	DE
D	F		
E		G	
F	FJ	G	HCIDE
G	FJ	G	HCIDE
H	FJ	G	CIDE
I	J		
J			P
K	L		
L	M		
M		N	
N			P

DFA STATE TABLE

	0	1
ABKIDE	LFJ	G
(LFJ)	MFJ	G
G	FJ	G
(MFJ)	FJ	NG
(FJ)	FJ	G
(NG)	FJ	G



DFA



2) a) All strings of lowercase letters that contain the 5 vowels in order

- assuming there can be extra vowels as long as 5 are in order
ie zaaueiouy is valid zabeauiozuv is valid

$$L = [a-z]^*$$

$$RE = \{L\}a\{L\}e\{L\}i\{L\}o\{L\}u\{L\}$$

- assuming there can be only one instance of each vowel

$$L = [bcdfghjklmnpqrstvwxyz]$$

$$RE = \{L\}a\{L\}e\{L\}i\{L\}o\{L\}u\{L\}$$

- assuming vowels can appear multiple times but must be ordered
ie aaeiou is valid aeaiou is not

$$L_a = [bcdfghjklmnpqrstvwxyz] \quad L_e = [bcdfghjklmnpqrstvwxyz]$$

$$L_i = [bcdfghjklmnpqrstvwxyz] \quad L_o = [bcdfghjklmnpqrstvwxyz]$$

$$L_u = [bcdfghjklmnpqrstvwxyz]$$

$$RE = \{L_a\}a\{L_e\}e\{L_i\}i\{L_o\}o\{L_u\}u\{L_u\}$$

b) Strings with an even number of quotes in them.

- assuming a quote with zero quotes is valid, assuming the empty string is valid

$$RE = ([^']*[^']*[^']*)+|[^']*$$

- assuming the empty string is not valid

$$RE = ([^']*[^']*[^']*)+|[^']*$$

- assuming the string requires at least 2 quotes

$$RE = ([^']*[^']*[^']*)+$$

3)a) Set of strings with an equal number of 1's and 0's

$$S \rightarrow 1S0 \mid 0S1 \mid SS \mid \epsilon$$

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This grammar works because any time a 1 is added a 0 is also added

b) Set of strings with an unequal number of 1's and 0's

- language with more 1's than 0's

$$S_1 \rightarrow A1S_1 \mid A1A$$

$$A \rightarrow 0A1 \mid 1A0 \mid AA \mid \epsilon$$

- language with more 0's than 1's

$$S_0 \rightarrow A0S_0 \mid A0A$$

$$A \rightarrow 0A1 \mid 1A0 \mid AA \mid \epsilon$$

- Combine these two

$$S \rightarrow S_1 \mid S_0$$

$$S_1 \rightarrow A1S_1 \mid A1A$$

$$A \rightarrow 0A1 \mid 1A0 \mid AA \mid \epsilon$$

$$S_0 \rightarrow B0S_0 \mid B0B$$

$$B \rightarrow 0B1 \mid 1B0 \mid BB \mid \epsilon$$

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4) Construct the First and Follow sets for

$$A \rightarrow BAC|FE$$

$$B \rightarrow bEF|g$$

$$E \rightarrow e|\epsilon$$

$$F \rightarrow f|EH$$

$$H \rightarrow h$$

$$\begin{aligned} \text{First}(A) &= \text{First}(B) + \text{First}(F) = \{b, g\} + \{f, e, h\} \\ &= \{b, g, f, e, h\} \end{aligned}$$

$$\text{First}(B) = \{b, g\}$$

$$\text{First}(E) = \{e, \epsilon\}$$

$$\begin{aligned} \text{First}(F) &= \{f\} + \text{First}(E) - \epsilon + \text{First}(H) = \{f\} + \{e\} + \{h\} \\ &= \{f, e, h\} \end{aligned}$$

$$\text{First}(H) = \{h\}$$

$$\text{Follow}(A) = \{\$ \} + \text{First}(c) = \{\$ \} + \{c\} = \{\$, c\}$$

$$\text{Follow}(B) = \text{First}(A) = \{b, g, f, e, h\}$$

$$\begin{aligned} \text{Follow}(E) &= \text{First}(F) + \text{Follow}(A) + \text{First}(H) = \{f, e, h\} + \{\$, c\} \\ &\quad + \{h\} = \{f, e, h, \$, c\} \end{aligned}$$

$$\begin{aligned} \text{Follow}(F) &= \text{First}(E) - \epsilon + \text{Follow}(A) + \text{Follow}(B) \\ &= \{e\} + \{\$, c\} + \{b, g, f, e, h\} = \{b, g, f, e, h, \$, c\} \end{aligned}$$

$$\text{Follow}(H) = \text{Follow}(F) = \{b, g, f, e, h, \$, c\}$$

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5) Determine if the grammars are LL(1)

a) $S \rightarrow [S] \mid A$ $\text{First}^+([S]) = \{ [\}$
 $A \rightarrow [A] \mid \epsilon$ $\text{First}^+(A) = \{ [\epsilon \}$
 Not LL1 $\text{First}^+([S]) \cap \text{First}^+(A) = \{ [\} \neq \emptyset$

b) $S \rightarrow ABc$ LL(1) because for A + B when
 $A \rightarrow a | \epsilon$ the derive ϵ the First(~~a~~) \cap Follow(A)
 $B \rightarrow b | \epsilon$ $= \emptyset$ ($\{a\} \cap \text{First}(b) \cap \{b\} = \emptyset$)
 $(\{b\} \cap \{c\} = \emptyset)$

c) $S \rightarrow ABBA$ $\text{First}(a) = \{a\}$ $\text{First}(b) = \{b\}$ $\text{First}(\epsilon) = \{\epsilon\}$
 $A \rightarrow a | \epsilon$ A) ✓ $\text{First}(a) \cap \text{First}(\epsilon) = \emptyset$
 $B \rightarrow b | \epsilon$ ✓ only ϵ derives ϵ
 ✓ B derives $\epsilon + \text{First}(a) \cap \text{Follow}(A)$
 $\text{Follow}(A) = \text{First}(B) + \text{Follow}(S)$
 $= \{b, \$ \}$
 $\{a\} \cap \{b, \$ \} = \emptyset$

B) Follow(B) = First(B) + First(A)
 $= \{b, a, \epsilon\}$
 $\{b\} \cap \{b, a, \epsilon\} = \{b\} \neq \emptyset$
 Not LL(1)

D) $S \rightarrow aAbc \mid bAc$ ✓ $\text{First}(aAbc) \cap \text{First}(bAc) = \emptyset$
 $A \rightarrow b \mid \epsilon$ $\{a\} \cap \{b\} = \emptyset$
 ✓ neither derives ϵ
 A) ✓ $\text{First}(b) \cap \text{first}(\epsilon) = \emptyset$
 ✓ only ϵ derives ϵ
 $\text{Follow}(A) = \{b, c\}$
 $\text{First}(b) \cap \{b, c\}$
 $\{b\} \cap \{b, c\} = \{b\} \neq \emptyset$
 Not LL(1)