

On the Median and Small Parsimony Problems in some Genome Rearrangement Models

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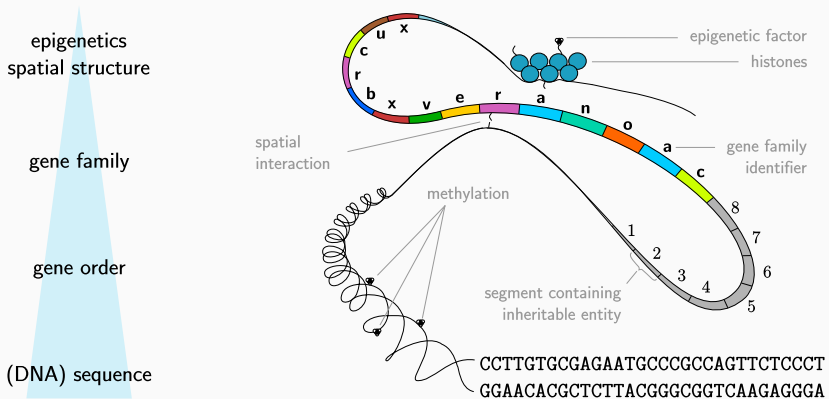
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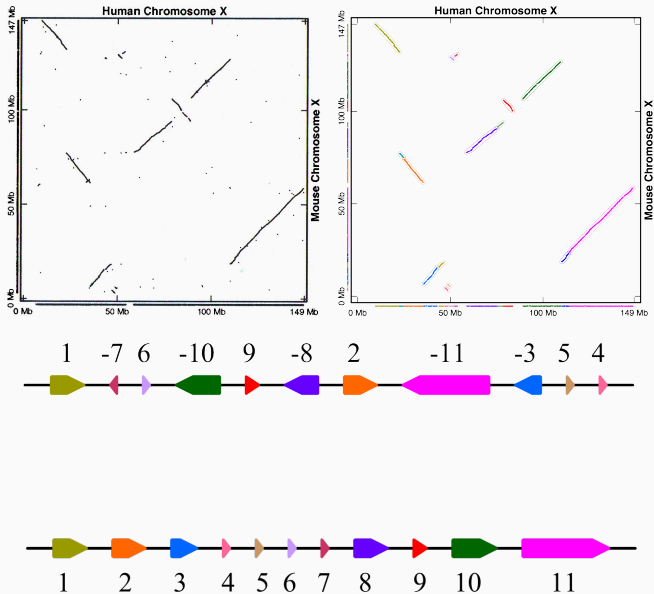
Bielefeld Block Seminar on Comparative Applied Genomics



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Introduction





A gene order with a single (resp. multiple) chromosome(s) is *unichromosomal* (resp. *multichromosomal*).

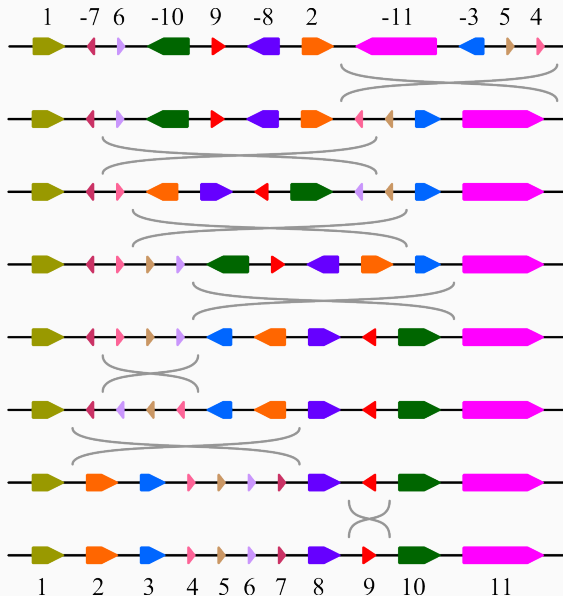
A gene order having only *linear* (resp. *circular*) chromosomes is *linear* (resp. *circular*).

A gene order having both *linear circular chromosomes* is *mixed*.

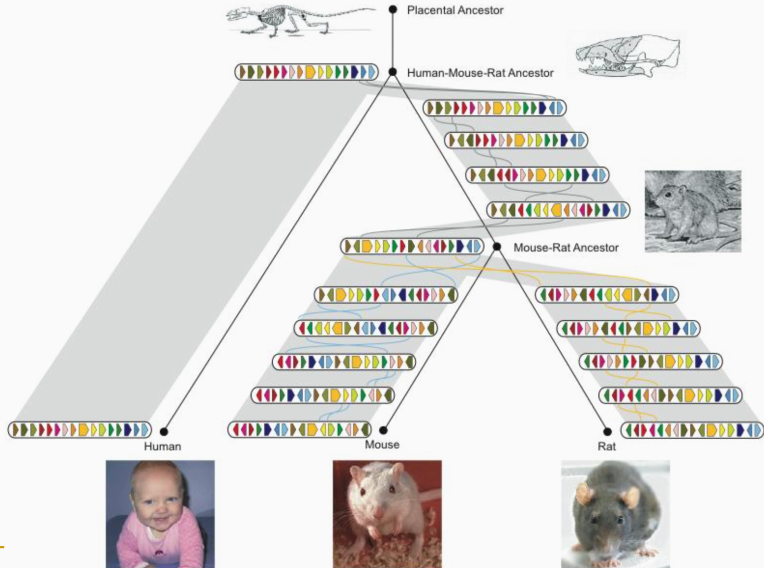
A gene order where some gene appears more than once is said to contain *duplications*.

A set of gene orders where each gene appears once and exactly once per gene order is said to have *equal gene content*.

Rearrangement Scenario



Rearrangement Evolution



Given two gene orders A and B

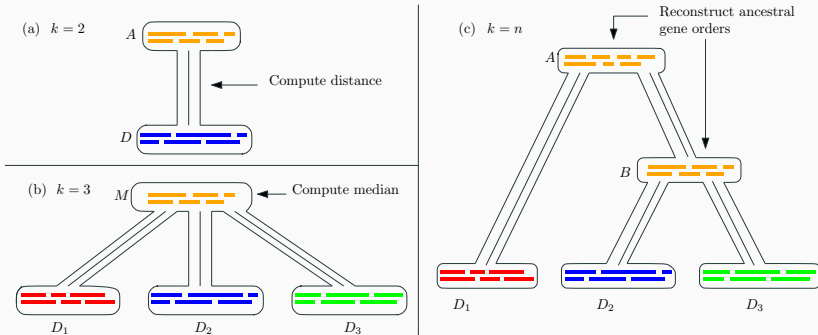
- **Parsimonious Distance:** Compute the minimum # of rearrangements to transform gene order A into gene order B .
- **Parsimonious Scenarios:** Counting, sampling, computing parsimonious scenarios transforming A into B .

Given $k \geq 3$ gene orders

Median of k : Compute a gene order M that minimizes the sum of the pairwise distances to k given gene orders.

Given a species phylogeny and the leaves gene orders

Small Parsimony (SPP): Compute ancestral (internal nodes) gene orders that minimize the sum of per-edge distances.



Genome rearrangement problems are defined by

- **Gene orders:** The structure of the considered gene orders (chromosomal structure, gene content structure).
- **Genome Rearrangement Model:** The considered genome rearrangements operators.

The Pairwise Distance Problem

Linear Unichromosomal Gene Orders, Equal Gene content, no Duplication, Reversal

Gene order structure: signed permutations

All gene orders have the same gene content with each gene appearing once and exactly once per gene order, and each gene order is composed of a single chromosome.

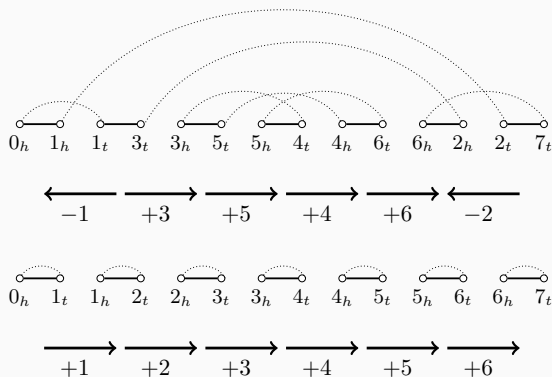
$$\pi = (-1 \quad 3 \quad 5 \quad 4 \quad 6 \quad -2)$$

Evolution model

The only allowed operation to transform a gene order is the reversal.

Main result

Computing the parsimonious/estimated distance is tractable.



$$d_{\text{reversal}}(A, B) = n - c + \text{something}$$

Linear Multichromosomal Gene Orders, Eq. Gene Content, no Duplication, Reversals/Translocations/Fusions/Fissions

Genome structure: “broken” signed permutations

$$A = (-1 \ 3)(5 \ 4 \ 6 \ -2), \ B = Id_6$$

Evolution model

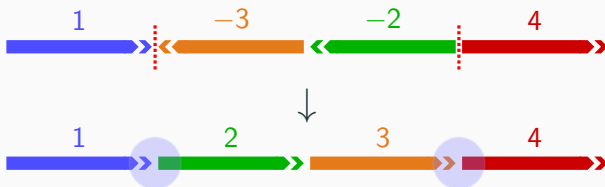
Reversal, translocation, chromosome fusion and fission.

Main result

- Computing the parsimonious distance is tractable.
- Idea: there is a way to order chromosomes together in order to reduce the problem to the reversal distance problem. The proof is highly non-trivial and required several corrections.

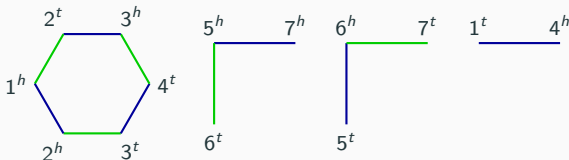
Double-Cut-and-Join

Cut two adjacencies (including possibly a terminal one) and reconnect the four created extremities by creating two adjacencies.



$$A = (1\ 2\ 3\ 4\ 5\ 6\ 7) = \{1^h 2^t, 2^h 3^t, 3^h 4^t, 4^h 5^t, 5^h 6^t, 6^h 7^t\}$$

$$B = [1\ -2\ -3\ -4](6\ 5\ -7) = \{1^h 2^h, 2^t 3^h, 3^t 4^t, 4^h 1^t, 6^h 5^t, 5^h 7^h\}$$



$$d_{\text{DCJ}}(A, B) = n - c - e/2 = 5$$

Bergeron, Mixtacki and Stoye (2006-2013)

The distance formula works for all gene order structures

(uni/multi-chromosomal, linear/circular), and provides a unifying framework allowing to retrieve easily other distance formulas.

SCJ. Meidanis and Feijao (2009-2011)

Either cut one adjacency or create one adjacency.

Breakpoint

A breakpoint is an adjacency or telomere (chromosome extremity) present in a gene order and not in the other.

$$A = (A_a, A_t) = \left(\{1^h 2^t, 2^h 3^t, 3^h 4^t, 5^t 7^h, 5^h 6^t, 6^h 7^t\}, \{1^t, 4^h\} \right)$$

$$B = (B_a, B_t) = \left(\{1^t 3^t, 1^h 2^t, 2^h 4^h, 3^h 4^t, 5^t 7^h, 5^h 6^h\}, \{6^t, 7^t\} \right)$$

$$d_{\text{SCJ}}(A, B) = |A_a \Delta B_a| = 6$$

$$d_{\text{BP}}(A, B) = n - |A_a \cap B_a| - \frac{|A_t \cap B_t|}{2} = 7 - 3 - 0/2 = 4$$

Inversions

There is no exact polynomial time to count the number of inversion scenarios between two unichromosomal genomes. Some approximations method do exist that work well on small instances (Braga et al (2008)). Sampling can be approximated through an MCMC (Miklós, Tannier (2010)).

DCJ

Counting and sampling (under the uniform distribution) DCJ scenarios is tractable for circular genomes (Braga, Stoye (2010)). In the general case it can be approximated with an FPRAS (Miklós, Tannier (2012)).

SCJ

Counting and sampling (under the uniform distribution) SCJ scenarios is tractable for circular genomes (Miklós, Kiss, Tannier (2014)).

The Median and Small Parsimony Problems

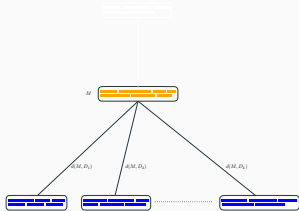
Early Results

Discovering Tractability Islands: SCJ and Breakpoint Distances

The SCJ/Breakpoint Small Parsimony Problem

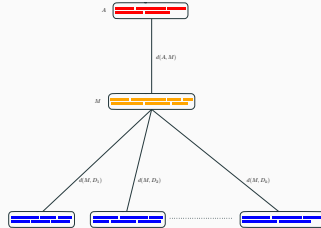
Early Results

Directed median



Minimize: $\sum_{i=1}^k d(M, D_i)$

Rooted median



Minimize: $d(A, M) + \sum_{i=1}^k d(M, D_i)$

Remark. In models with equal gene content, both problems are identical.

Small Parsimony (SPP)

Compute ancestral gene orders that minimize the sum of per-edge distances along a given unrooted species tree S .

Median-based heuristic for the SPP

- Assign an initial gene order to each internal node of S .
- *Repeat*
 - Chose an internal node U .
 - Let A, B, C be the three neighbours of U in S .
 - Assign to U the gene order of a parsimonious median of A, B, C .
- *Until* a convergence criterion is met.

Bryant (1998); Pe'er and Shamir (1998)

The breakpoint median of three is NP-hard if the median is constrained to be unichromosomal, linear or circular.

Caprara (2003)

The reversal median of three is APX-hard if the median is constrained to be unichromosomal linear.

Reduction idea: alternating cycles decomposition, that implies also NP-hardness of the unichromosomal DCJ median.

Tannier, Zheng and Sankoff (2009)

The DCJ median is NP-hard for multichromosomal genomes.

Discovering Tractability Islands: SCJ and Breakpoint Distances

Tannier, Zheng and Sankoff (2009)

The circular and mixed multichromosomal breakpoint median problems can be solved through a reduction to a Maximum-Weight Matching (MWM) problem.

- For every pair of gene extremities x and y draw an edge weighted by the number of gene orders having xy as an adjacency.
- For every gene extremity x draw an edge xt_x weighted by half the number of gene orders having x as a telomere.

The linear case

If the median is constrained to be linear, the breakpoint median is NP-hard.

Feijao and Meidanis (2011)

The multichromosomal linear, circular and mixed SCJ median problems are tractable.

- **A weighted graph:** For an adjacency xy , $f(xy)$ is the difference between the number of gene orders that do not contain xy and the number of gene orders containing xy .
- **Mixed case:** Pick all adjacencies xy s.t. $f(xy) \leq 0$.
- **Circular case:** Compute a Min-Weight Perfect Matching.
- **Linear case:** Solve the mixed case and remove an adjacency of maximum weight per circular chromosome.

The SCJ/Breakpoint Small Parsimony Problem

Question

Some median problems are tractable (breakpoint distance for circular and mixed genomes, SCJ distance for multichromosomal genomes). Does the tractability extend to the SPP ?

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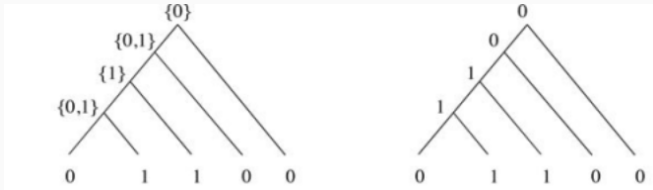
Feijao and Meidanis (2011)

The SCJ SPP is tractable for multichromosomal mixed gene orders.

It is easy for a single adjacency

Consider each adjacency xy as a binary (0/1) character.

Computing a parsimonious evolutionary scenario for a given adjacency can be done easily using Dynamic Programming (DP): Fitch or Sankoff-Rousseau algorithms.



Pitfall

Can we process each adjacency separately to obtain a global solution?

Not immediately as it can happen that a gene extremity is assigned to more than two adjacencies at some internal node of S .

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Solution

Use Fitch algorithm and, if, at some node, both presence (1) or absence (0) of an adjacency are parsimonious, chose absence (0).

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Solution

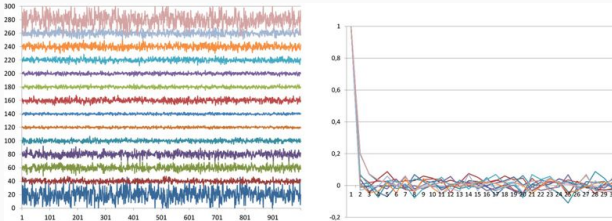
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Consequence

This excludes valid optimal solutions and results in solutions that might be (artificially) highly fragmented.

Miklós and Smith (2015)

A Markov Chain that uses the Sankoff-Rousseau DP algorithm to propose local moves. This does not exclude any optimal solutions.



Idea

Weight ancestral adjacencies with some kind of prior and penalize discarding adjacencies.

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Weight: Boltzmann Sampling Single Adjacency Scenarios

If we use the Sankoff-Rousseau DP algorithm, we can compute, for each adjacency xy and node U of S , the probability (in the Boltzmann framework) of observing xy at U : $w_U(xy)$.

Luhmann, Lafond et al (2016): The Weighted SCJ SPP

Find ancestral gene orders that minimize

$$\alpha \left(\sum_{U \in S, xy \notin G_U} w_U(xy) \right) + (1 - \alpha) \left(\sum_{(U,V) \in S} d_{SCJ}(U, V) \right)$$

Complexity Results

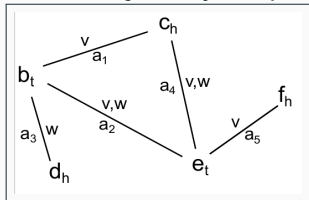
- Tractable for $\alpha = 1$ (no SCJ cost, MWM at each node U) and $\alpha = 0$ (only SCJ cost, SCJ SPP)
- NP-hard for $1 > \alpha > 33/34$.

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Sankoff-Rousseau type DP Algorithm

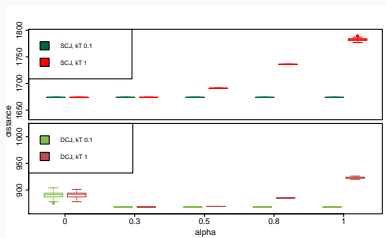
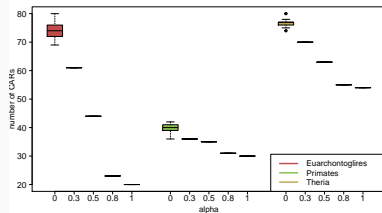
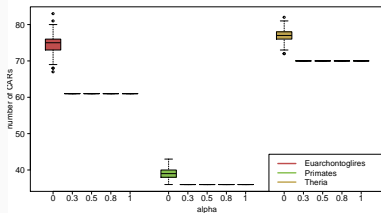
Global Adjacency Graph



n species, N genes, c conn. comp., D max. vertex degree, M max. size of a conn. comp.

Sampling an opt. solution can be done in time $O(nN(1 + D)^{2M})$ and space $O(nN(1 + D)^M)$.

The Weighted SPP (3)



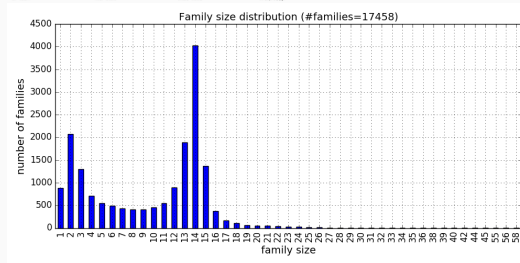
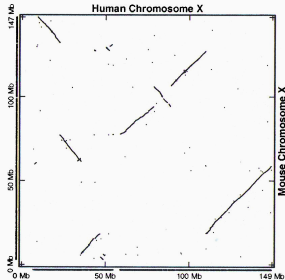
Handling Duplications

The Pairwise Distance Problem

The Median Problem

The Small Parsimony Problem

Gene Duplication is a Frequent Evolutionary Event



Genomes are then represented by “sequences”, not “permutations” 35

The Pairwise Distance Problem

The Exemplar Approach, Sankoff (1999)

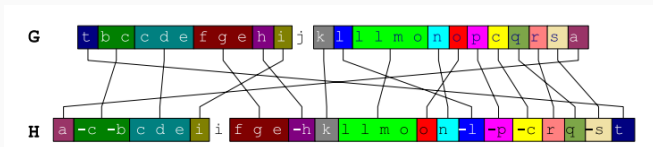
Keep one (*exemplar*) copy per gene family in each genome s.t. the distance between the resulting “permutations” is minimized.

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The Maximal Matching Approach

Find a maximal (A, B) matching respecting gene families s.t. the distance between the resulting “permutations” is minimized.



Bryant (2000); Bulteau and Jiang (2012)

The exemplar distance problem is NP-hard, in the breakpoint and reversal models, even if each gene family has one copy in a genome and at most two copies in the second genome.

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Blin, C. Fertin (2004)

The maximal matching distance is NP-hard, in the breakpoint model, even if a single gene family contains duplicates.

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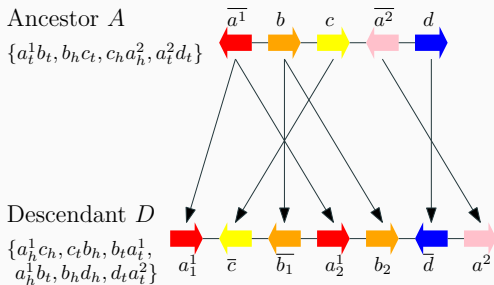
Zeira and Shamir (2016)

If duplicated genes appear through Whole-Chromosome Dup. (WCD) events, the SCJ-WCD distance problem is tractable.

Definition

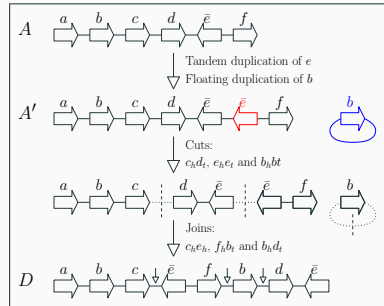
The distance between A and B is *directed* if A is an ancestor of B .

A is *trivial* for B if, for each gene of B , we know its parent gene in A (*orthology relations*).



Feijao, Mane, C. (2017)

If duplicated genes appear through single-gene Tandem Duplications (TD) or Floating Duplications (FD), the directed SCJ-TD-FD distance problem is tractable.

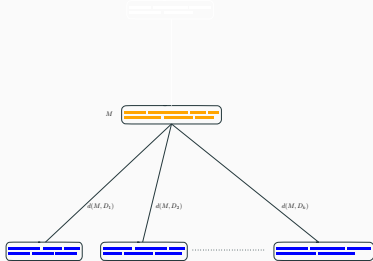


$$d_{SCJ-TD-FD}(A, D) = |A - D| + |D - A| + 2n_d$$

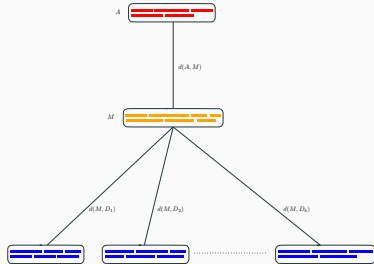
where n_d is the number of duplications in D . Moreover a parsimonious scenario can be computed in linear time.

The Median Problem

The Directed Median



The Rooted Median



Remark

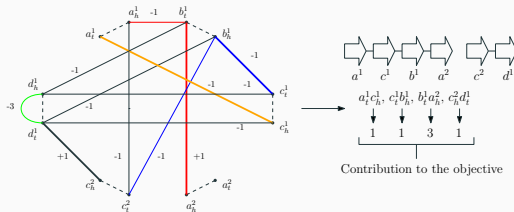
In both cases we assume the *gene content* of the median gene order M is given, as well as the orthology relations between genes of M and genes of A , and genes of the D_i 's and genes of M .

Feijao, Mane and C. (2017)

The Directed SCJ-TD-FD Median problem is tractable (reduction to a MWM problem).

Mane, Lafond, Feijao and C. (2018)

The Rooted SCJ-TD-FD Median problem is NP-hard, even with $k = 2$ and $D_1 = D_2$. It can be described as a simple Integer Linear Program (ILP) solving efficiently a colored MWM problem.



Candidate Adjacencies

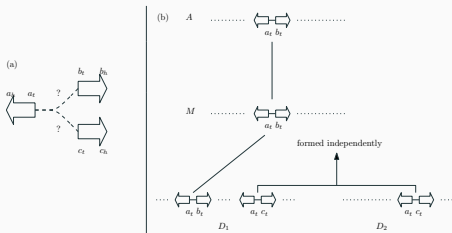
Adjacencies seen in majority of the genomes A , D_1, D_2 .

An optimal median is composed only of candidate adjacencies.

Cand. adj. are *conflicting* if they contain the same gene extremity.

Convergent Evolution

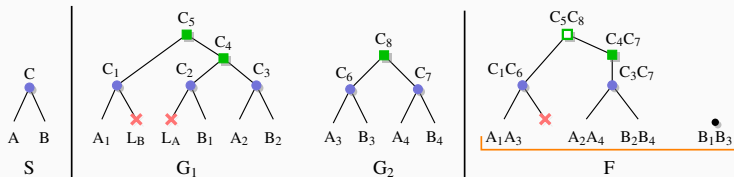
Hardness of the median problem results from the presence of conflicting candidate adjacencies created by duplication and convergent evolution.



The Small Parsimony Problem

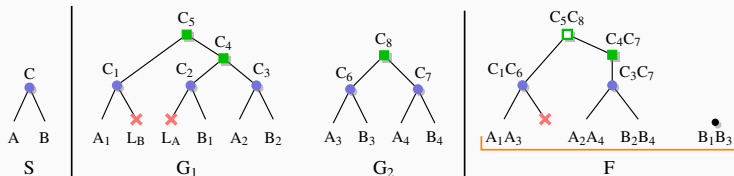
DeCoStar: Bérard et al (2012-2017)

Given two gene families x and y , whose reconciled gene trees are provided, one can solve the SPP for adjacencies between genes of x and y using Sankoff-Rousseau-like DP.



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C. and Ponty (2015)

The DeCo DP scheme is complete and unambiguous \rightarrow Boltzmann sampling can be done efficiently.

Applications: adjacency probabilities, MCMC SPP sampling.

A Negative Point of View

The SCJ-TD-FD SPP and Weighted SPP are NP-hard.

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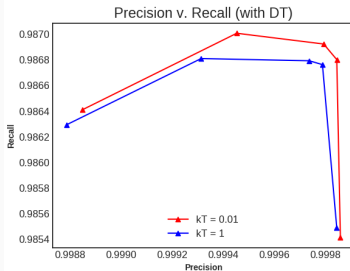
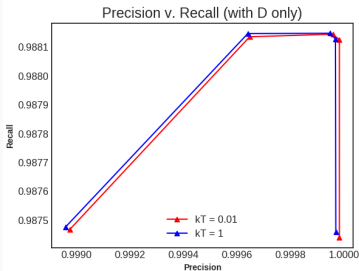
A Positive Point of View

- If reconciled gene trees are given, it is tractable for a single adjacency and provide ancestral adjacency (prior) probabilities.
- The SCJ-TD-FD distance is tractable and the distance formula is amenable to being implemented in an ILP.

So we have all the elements to design an ILP for the generalization of the Weighted SCJ SPP to the SCJ-TD-FD model.

Feijao, Mane and C. (submitted)

The Weighted SCJ-TD-FD SPP (given ancestral gene content and weight on ancestral adjacencies) can be solved efficiently by an ILP, even in the linear case; accurate results on simulated data.



Conclusion

No duplicated genes

- The pairwise distance problem is well understood.
- DCJ and SCJ are the most natural frameworks.
- The median problems is only tractable in the SCJ/Breakpoint models; there are subtle differences between both models.
- The SPP is only tractable in the SCJ model, with restricting assumptions on the gene order model.

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Handling duplicated genes

- The pairwise distance problem is hard generally with exceptions in the SCJ model.
- The median and SPP problems are hard.
- ILP works well in the SCJ model for the median and SPP.

In most cases, the available algorithms provide a single optimal solution, chosen somewhat arbitrarily within a huge space of co-optimal or slightly sub-optimal solutions.

There is a need to develop approaches that

- estimate the size of the space of co-optimal solutions;
- sample within the space of solutions (co-optimal and sub-optimal);
- integrate biological constraints in the genome rearrangement models (contact data, repeats mediating rearrangements, replication-centered rearrangements, ...).

- Interpolating between the breakpoint and DCJ distances ($n - c_2$ versus $n - c$): when does the median become hard?
- MCMC sampling of SCJ SPP solutions with duplicated genes.
- ILP for the DCJ median and SPP with duplicated genes and prior on adjacencies (Shao, Lin, Moret (2014); Avdeyev et al (2017)).
- Integrating segmental duplications in the SCJ-TD-FD or SCJ-WD models.
- SCJ+WD SPP (cancer evolution).
- FPT results for the median and SPP.