

DBMS :-

→ Issues with file system :-

1. Data mapping & access
2. Redundancy of data
3. Inconsistency of data
4. Unauthorised access of security
5. No concurrent access
6. No backup & recovery.

→ Why DBMS is required?

→ Overcomes drawbacks of file system.

optimizes :-

→ storage of large data

→ Faster retrieval of data

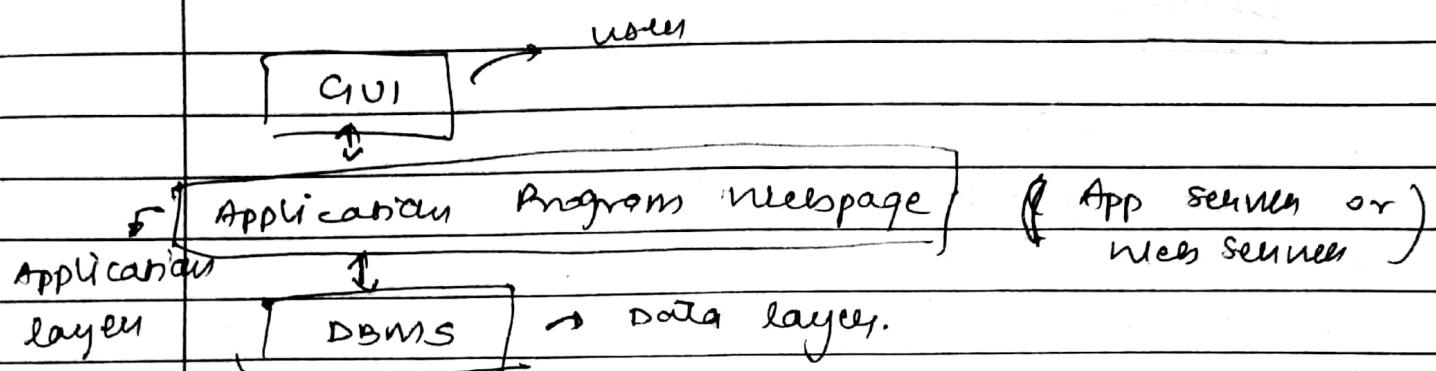
Characteristics :-

- Real world entity
- Relation based tables
- Less Redundancy
- Query language
- Consistency
- Concurrent Access
- Security

Architecture of DBMS :-

→ major advantage → we can access different modules independently.

- 3 tier Adv :-



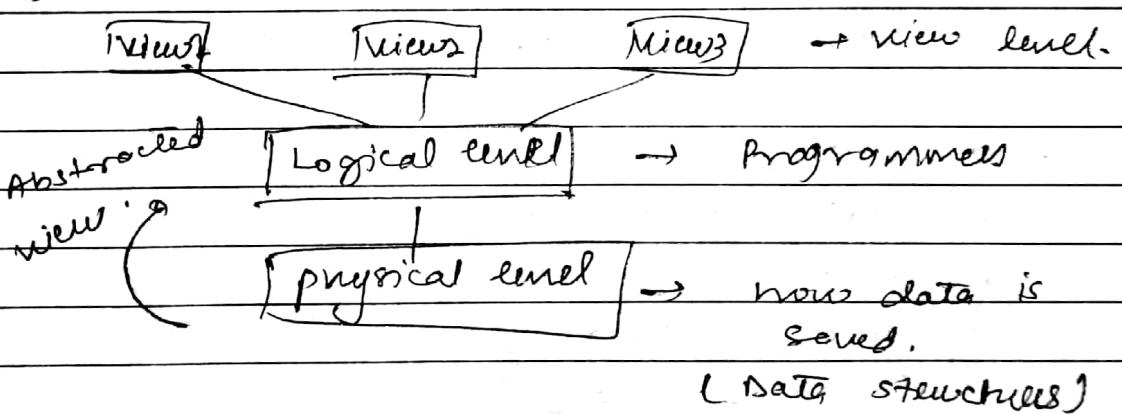
Adv :-

1. Easy to maintain & modify
2. Security
3. Performance.

Disadv :- complex

- Data Abstraction :-

Levels :-



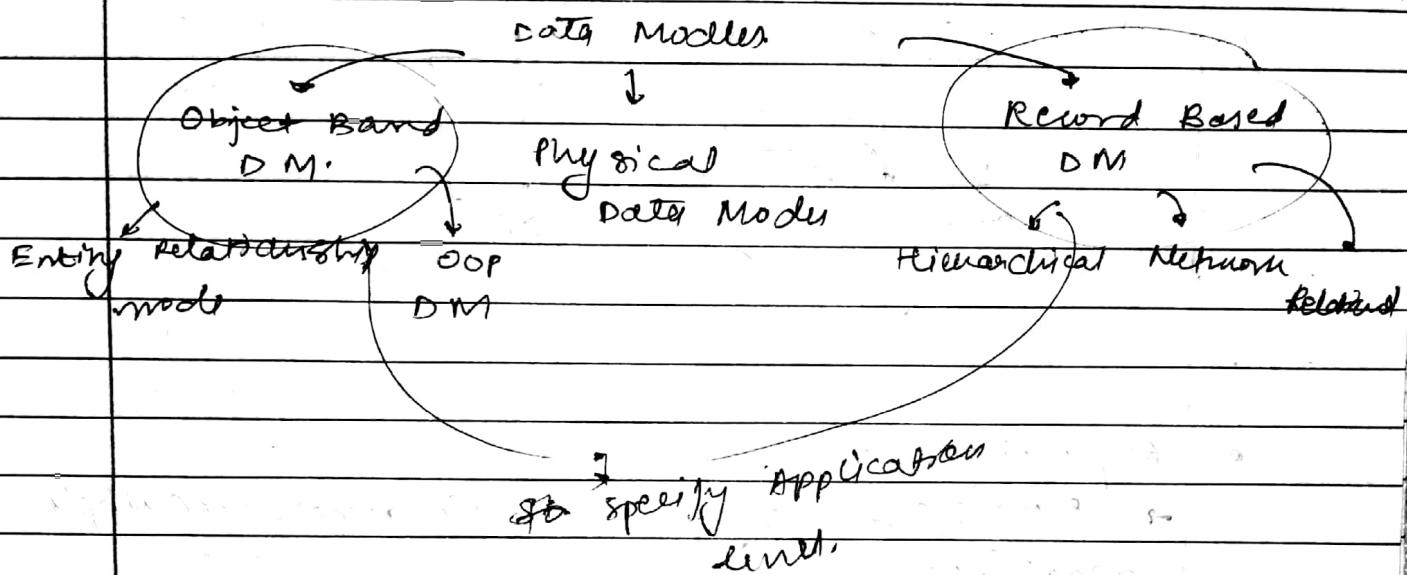
- Instance : Saving data in database at any point of time is called instance.

- Schema :-

- Physical schema : Design of physical env.
- logical schema : defines logical env.
- View schema : how my data is showed to the end user.

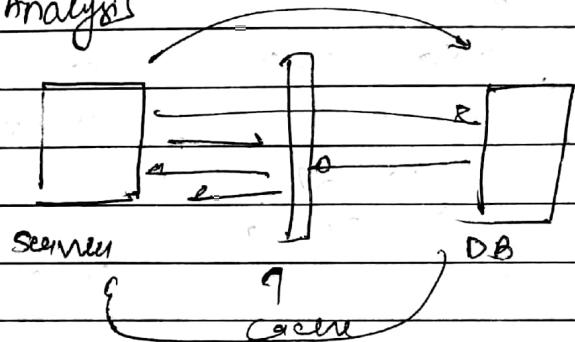
→ Physical Data Independence :- changing in physical level will not effect the logical one.

Data Models :- Describes the database.



System Design :-

1. Requirements collection.
2. Architecture Design → what all components.
3. Final Analysis



- Replication :-
→ Frequently copying the data across multiple machines.
→ Post replication, multiple copies of the data exists across machines.
- Vertical scaling → Adding more servers.
→ increasing resources,
(RAM, CPU, storage),
(horizontal scaling)
- Sharding : sharding refers to splitting the very large database into smaller, faster & more manageable parts called data shards.

Availability
↑
CAP Theorem. Not possible to
consistency partition simultaneously
To tolerate. guarantee
all of the following

Steps to approach :-

1. Feature expectations
2. Estimations
3. Design Goals
4. Skeleton of the design
5. Deep dive.

constraints :-

- Primary key → unique & not null.
- Composite key
- Unique key → can be null. (but only 1).
- Foreign key
- check constraint
- Not Null.

NORMALIZATION :-

- Insertion Anomaly
- Deletion Anomaly
- updation Anomaly.

Functional Dependency:

- It is a relationship that exists between multiple attributes of a relationship.
- Dependent → when one attribute uniquely identifies another attribute of the same table..

Axioms :

$$ABC \rightarrow AB \rightarrow \cancel{WZ} \rightarrow W \rightarrow Z \quad X \rightarrow Y \rightarrow WX \rightarrow YZ$$

Reflexive.

C1804ME 2 -

→ candidate key should be minimal,

1NF :-

- There cannot be multivalued or composite attributes
- Each record needs to be unique.

2NF :-

- Table should be in 1NF $(P \rightarrow N.P)$
- There should be no Partial dependency.

Partial dependency

$X \rightarrow a$

- 1) X is a proper subset of some candidate key & a is non-prime or non-key attribute.
↓ part of any candidate key.

3NF :-

$(N.P \rightarrow N.P)$

- Table should be in 2NF

- No transitive dependency

$X \rightarrow a$

both a and a' are prime

BCNF :-

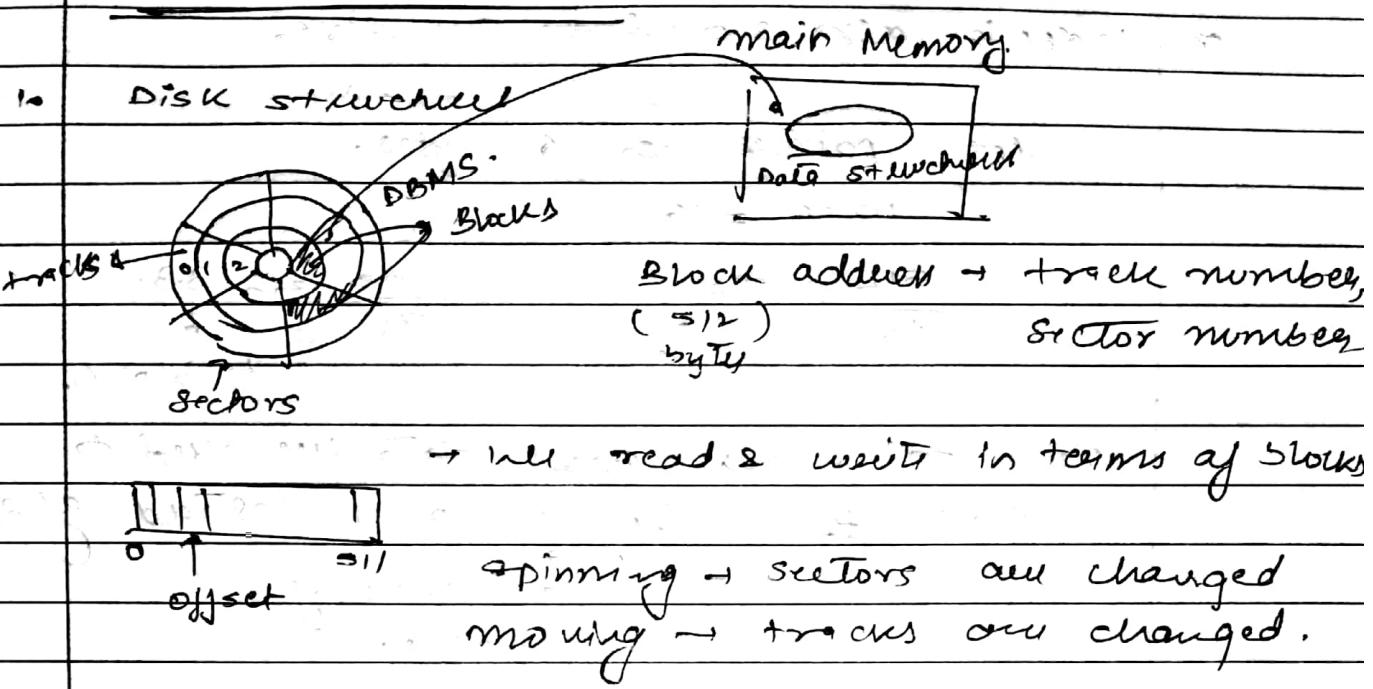
$X \rightarrow Y$

\uparrow
superkey.

Search engine :-

- Indexes of other websites / webpages.
- crawling / spiders. → web scraping.

B-Trees & B+-Trees :-



2. How data is organised.

Employee

eid - 10	size = 4
name - 50	128 \times
dep - 10	4 records
section - 6	per block.
addr - 50	$100 \times 128 = 25$ blocks in total.
	128 \times

- so the data is stored in 25 blocks

Index

eid	Pointers	Pointers to the data on the block.
1		
2		
3		

* Index is also stored on the disk.

let pointer is 6 bytes.

$$10 + 6 = 16$$

$$\frac{512}{16} = 32$$

$$\frac{100}{16} = 3 \text{ remainder } 4$$

4 blocks.

So to access some data we need to search for only 4 blocks instead of 25.
+
1 block of table.

Now we can have answers index-pointer
for those 4 blocks.
↓
sparse index

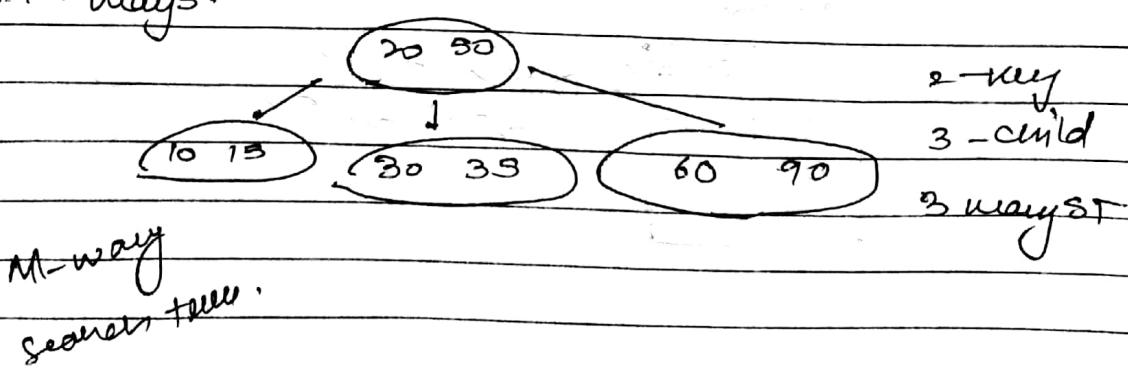
3. Multilevel Index

→ It looks like a tree if we turn it upside down.

4. M-way Search Tree.

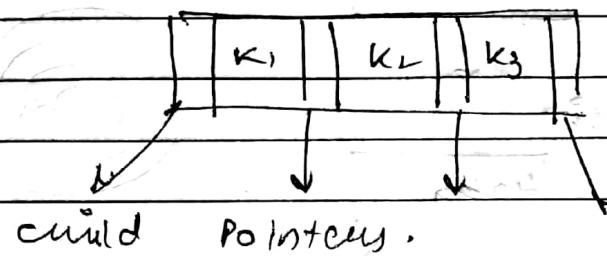
Binary Trees → 2 keys & children.

M-way.

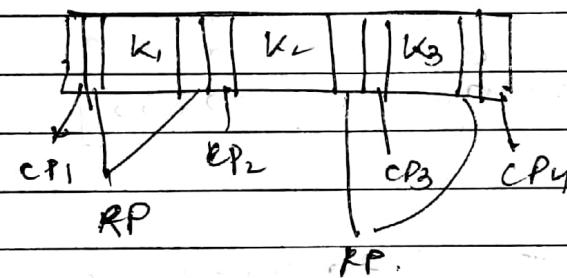


* m-way \rightarrow m-1 keys.

4-way tree :-



\rightarrow M-wayst for data base indexing



Problem with m-way search tree :-

- \rightarrow No control on insertion, we can insert as we like.
- \rightarrow height of m-way search tree can be n .
- \rightarrow it will take a lot of time in searching.

B-Trees :-

\rightarrow M-way search trees with some rules.

1. Every node must have $\lceil n/2 \rceil$ (ceil) value.
2. Root can have minimum 2 children.
3. All leaf nodes must be at same level.
4. Process is bottom up.

B-Tree 6-

$$m = 4$$

Key's 10, 20, 40, 50, 60

70, 80, 130, 135

10 20 40

↓

split

40 ~~40~~

10 20

50 60 70

split

40 70

10 20 30

50, 60

80

split 15 | 30 40 70 |

~~10 20 35 50 60 80~~

root

5 10 15 20

5 10 20

1. 40

15. 30

20

5 10

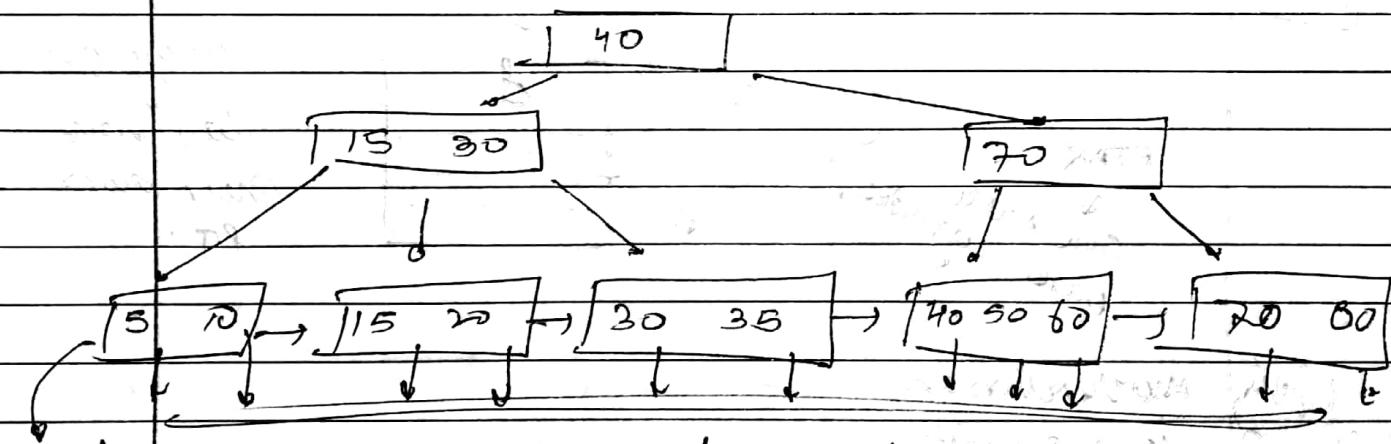
20

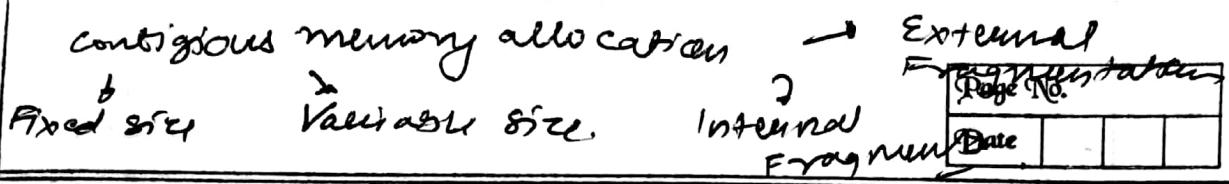
35

50 60

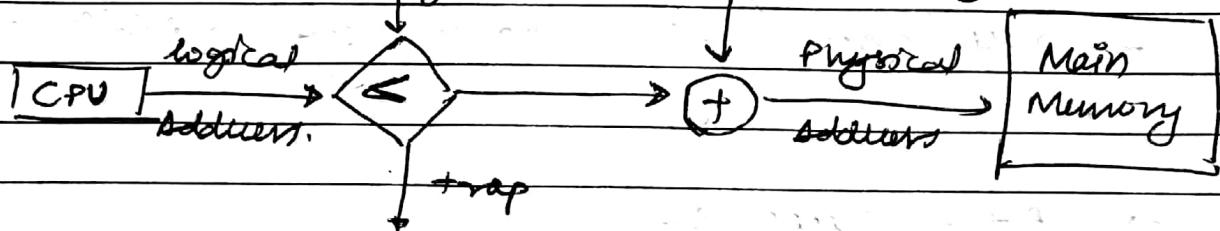
180

- As we are taking the right middle element it is right biased b-tree.
- we can also take left middle elements.
- B+ → Tree :-
- we don't have record pointers from every node.
- we only have record pointers from leaf nodes.
- every parent will have a copy of its node in the child.



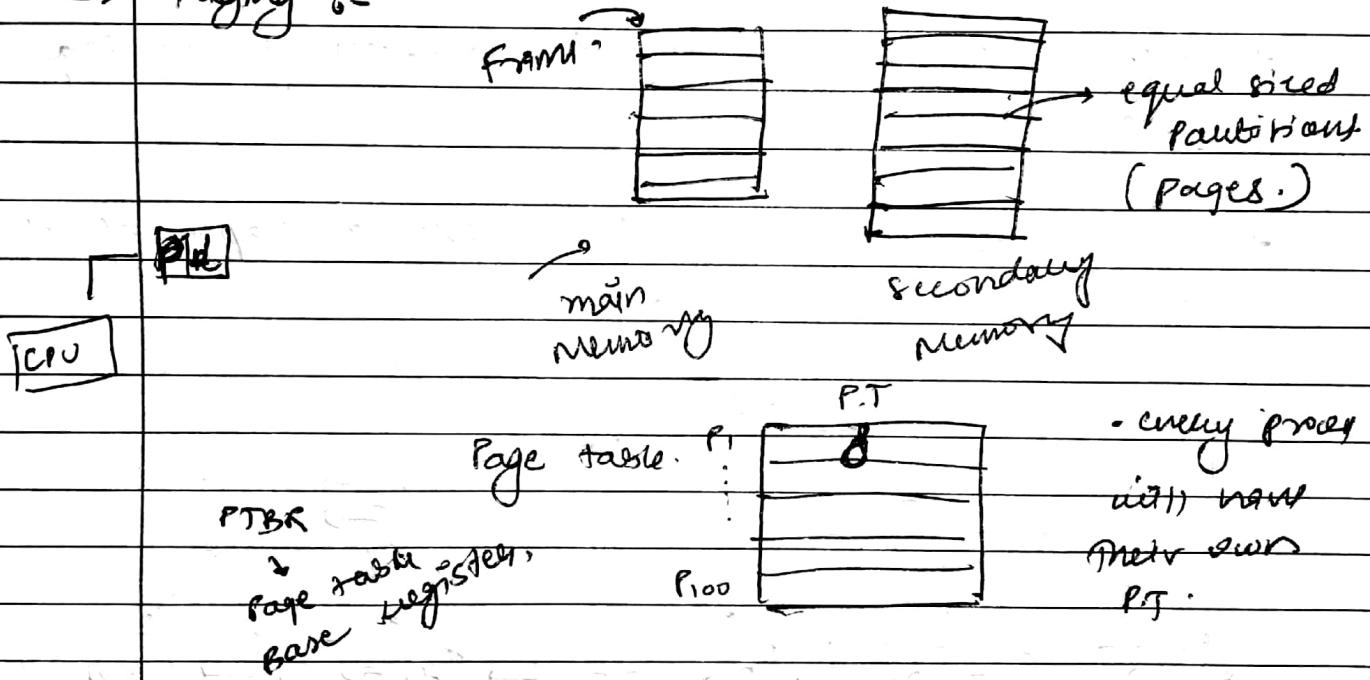


→ Contiguous Memory Allocation :-



Non contiguous. (Fixed size Partitioning)

→ Paging :-



Read of

Mutual Exclusion.

Hold & wait

No preemption

Circular wait

REACT :-

- Redux a layer top on React.
 - helps with the state management of the app.
 - data in the app
 - UI states of the app.

- ## → Component & templates :-

- Components look like HTML templates.
 - They can contain static
 - They also can contain JavaScript for functionality

- component state :- → Data

- JavaScript object.
 - describes the current state of the component
i.e data, UI-state
 - The state can be updated over time
 - Modal could close, the data we output could change.

e.g. state = q

name = 'Ryu',

age : 30

3

- ## DOM Events :-

Event handlers eg's onClick = {function} }

this purchase.

change state using \exists

~~setString~~ SetString()