

# Foundations of Software

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Slides in part adapted from:

University of Pennsylvania CIS 500: Software Foundations - Fall 2006

by Benjamin Pierce

# Course Overview

# What is “software foundations” ?

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Software foundations (or “theory of programming languages”) is the mathematical study of the **meaning** of programs.

The goal is finding ways to describe program behaviors that are both **precise** and **abstract**.

- ▶ **precise** so that we can use mathematical tools to formalize and check interesting properties
- ▶ **abstract** so that properties of interest can be discussed clearly, without getting bogged down in low-level details

# Why study software foundations?

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- ▶ To prove specific properties of particular programs (i.e., program verification)
  - ▷ Important in some domains (safety-critical systems, hardware design, security protocols, inner loops of key algorithms, ...), but still quite difficult and expensive
- ▶ To develop intuitions for *informal* reasoning about programs
- ▶ To prove general facts about all the programs in a given programming language (e.g., safety or isolation properties)
- ▶ To understand language features (and their interactions) deeply and develop principles for better language design  
(PL is the "materials science" of computer science...)

# What you can expect to get out of the course

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- ▶ A more sophisticated perspective on programs, programming languages, and the activity of programming
  - ▷ See programs and whole languages as formal, mathematical objects
  - ▷ Make and prove rigorous claims about them
  - ▷ Detailed knowledge of a variety of core language features
- ▶ Deep intuitions about key language properties such as type safety
- ▶ Powerful tools for language design, description, and analysis

Most software designers are language designers!

# Greenspun's Tenth Rule Of Programming

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Any sufficiently complicated C or Fortran program contains an ad-hoc, informally-specified, bug-ridden, slow implementation of half of Common Lisp.

– Philip Greenspun

# What this course is not

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- ▶ An introduction to programming
- ▶ A course on functional programming (though we'll be doing some functional programming along the way)
- ▶ A course on compilers (you should already have basic concepts such as lexical analysis, parsing, abstract syntax, and scope under your belt)
- ▶ A comparative survey of many different programming languages and styles

# Approaches to Program Meaning

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- ▶ **Denotational semantics** and **domain theory** view programs as simple mathematical objects, abstracting away their flow of control and concentrating on their input-output behavior.
- ▶ **Program logics** such as **Hoare logic** and **dependent type theories** focus on logical rules for reasoning about programs.
- ▶ **Operational semantics** describes program behaviors by means of abstract machines. This approach is somewhat lower-level than the others, but is extremely flexible.
- ▶ **Process calculi** focus on the communication and synchronization behaviors of complex concurrent systems.
- ▶ **Type systems** describe approximations of program behaviors, concentrating on the shapes of the values passed between different parts of the program.



# Overview

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This course will concentrate on operational techniques and type systems.

- ▶ Part I: Modeling programming languages
  - ▷ Syntax and parsing
  - ▷ Operational semantics
  - ▷ Inductive proof techniques
  - ▷ The lambda-calculus
  - ▷ Syntactic sugar; fully abstract translations
- ▶ Part II: Type systems
  - ▷ Simple types
  - ▷ Type safety
  - ▷ References
  - ▷ Subtyping

# Overview

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- ▶ Part III: Object-oriented features (case study)
  - ▷ A simple imperative object model
  - ▷ An analysis of core Java
  - ▷ An analysis of core Scala

# Organization of the Course

# Staff

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Instructor:	Martin Odersky <martin.odersky@epfl.ch>
Teaching Assistants:	Eugene Burmako Contact him about the mailing list! <eugene.burmako@epfl.ch>
	Denys Shabalin <denys.shabalin@epfl.ch>
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# Information

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Textbook: Types and Programming Languages,  
Benjamin C. Pierce, MIT Press, 2002

Webpage: <https://fos2015.github.io>

# Elements of the Course

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- ▶ The Foundations of Software course consists of
  - ▷ lectures (Tue 10:15-12:00, INM 200)
  - ▷ exercises and project work (Wed 10:15-12:00, CO 020)
- ▶ The lecture will follow in large parts the textbook.
- ▶ For lack of time, we cannot treat all essential parts of the book in the lectures, that's why the [textbook is required reading](#) for participants of the course.

# Homework and Projects

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You will be asked to

- ▶ solve and hand in some written exercise sheets,
- ▶ do a number of programming assignments, including
  - ▷ parsers,
  - ▷ interpreters and reduction engines,
  - ▷ type checkersfor a variety of small languages.
- ▶ The recommended implementation language for these assignments is [Scala](#).

# Scala

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- ▶ Scala is a functional and object-oriented language that is closely interoperable with Java.
- ▶ It is very well suited as an implementation language for type-checkers, in particular because it supports:
  - ▷ pattern matching,
  - ▷ higher-order functions,
  - ▷ inheritance and mixins.



# Learning Scala

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If you don't know Scala yet, there's help:

- ▶ The Scala web site:

[www.scala-lang.org](http://www.scala-lang.org)

- ▶ On this site, the documents:

- ▷ *A Brief Scala Tutorial* - an introduction to Scala for Java programmers. (short and basic).
- ▷ *An Introduction to Scala* (longer and more comprehensive).
- ▷ *An Overview of the Scala Programming Language* (high-level).
- ▷ *Scala By Example* (long, comprehensive, tutorial style).

- ▶ The assistants.

# Grading and Exams

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Final course grades will be computed as follows:

- ▶ Homework and project: 30%
- ▶ Mid-term exam: 30%
- ▶ Final exam: 40%

## Exams:

1. Mid-term: first week of November (exact date TBA)
2. Final exam: to be announced

# Collaboration

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- ▶ Collaboration on homework is **strongly encouraged**.
- ▶ Studying with other people is the best way to internalize the material
- ▶ Form pair programming and study groups!  
2-3 people is a good size. 4 is too many for all to have equal input.

"You never really misunderstand something  
until you try to teach it...  
" – Anon.

# Plagiarism

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- ▶ A single group will of course share code.
- ▶ But plagiarizing `code` by `other groups` as part of a project is unethical and will not be tolerated, whatever the source.

# Part I

## Modelling programming languages

# Syntax and Parsing

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- ▶ The first-level of modeling a programming language concerns its **context-free syntax**.
- ▶ Context free syntax determines a set of legal **phrases** and determines the **(tree-)structure** of each of them.
- ▶ It is often given on two levels:
  - ▷ **concrete**: determines the exact (character-by-character) set of legal phrases
  - ▷ **abstract**: concentrates on the tree-structure of legal phrases.
- ▶ We will be mostly concerned with abstract syntax in this course.
- ▶ But to be able to write complete programming tools, we need a convenient way to map character sequences to trees.

# Approaches to Parsing

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There are two ways to construct a parser:

- ▶ **By hand** Derive a parser program from a grammar.
- ▶ **Automatic** Submit a grammar to a tool which generates the parser program.

In the second approach, one uses a special **grammar description language** to describe the input grammar.

# Domain-Specific Languages

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- ▶ The grammar description language is an example of a domain-specific language (DSL).
- ▶ The parser generator acts as a processor (“compiler”) for this language — that’s why it’s sometimes called grandly a “compiler-compiler”.
- ▶ Example of a “program” in the grammar description DSL:

```
Expr    ::= Term { '+' Term | '-' Term }.  
Term    ::= Factor { '*' Factor | '/' Factor }.  
Factor  ::= Number | '(' Expr ')'
```



# Embedded Domain Specific Languages

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- ▶ An alternative to a stand-alone DSL is an **Embedded DSL**.
- ▶ Here, the DSL does not exist as a separate language but as an API in a **host language**.
- ▶ The host language is usually a general purpose programming language.

We will now develop this approach for grammar description languages.

# An EDSL for Parsing in Scala

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```
def expr : Parser[Any] = term ~ rep("+" ~ term | "-" ~ term)
def term : Parser[Any] = factor ~ rep("*" ~ factor | "/" ~ factor)
def factor: Parser[Any] = "(" ~ expr ~ ")" | numericLit
```

A small Perl hack could derive the above code from the informal grammar definition above:

- ▶ Insert a `def` at the beginning of each production.
- ▶ The “`::=`” becomes “`: Parser[Any] =`”.
- ▶ Sequential composition is now expressed by a `~`.
- ▶ Repetition `{...}` is now expressed by `rep(...)`.
- ▶ Option `[...]` is now expressed by `opt(...)`.

# Parser Combinators

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- ▶ The differences between Grammar A and Grammar B are fairly minor.  
  
(Note in particular that existing DSLs for grammar descriptions also tend to add syntactic complications to the idealized Grammar A we have seen).
- ▶ The important difference is that Grammar B is a valid Scala program, when combined with an API that defines the necessary primitives.
- ▶ These primitives are called [parser combinators](#).

# The Basic Idea

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For each language (identified by grammar symbol  $S$ ), define a function  $f_S$  that, given an input stream  $i$ ,

- ▶ if a prefix of  $i$  is in  $S$ , return  $\text{Success}(\text{Pair}(x, i'))$  where  $x$  is a result for  $S$  and  $i'$  is the rest of the input.
- ▶ otherwise, return  $\text{Failure}(\text{msg}, i)$  where  $\text{msg}$  is an error message string.

The first behavior is called **success**, the second **failure**.

# The Basic Idea in Code

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Assume:

```
class StandardTokenParsers {  
  type Parser = Input ⇒ ParseResult
```

where

```
type Input = Reader[Token]  // a stream of Tokens with positions.
```

and we assume a class `Token` with subclasses

- ▷ case class `Keyword(chars: String)` for keywords,
- ▷ case class `NumericLit(chars: String)` for numbers,
- ▷ case class `StringLit(chars: String)` for strings,
- ▷ case class `Identifier(chars: String)` for identifiers.

Also assume a class `ParseResult[T]` with subclasses

```
case class Success[T](result: T, in: Input)
extends ParseResult[T]

case class Failure(msg: String, in: Input)
extends ParseResult[Nothing]
```

# Object-Oriented Parser Combinators

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- ▶ In fact, we will also need to express `|` and `~` as methods of parsers.
- ▶ That's why we extend the function type of parsers as follows:

```
abstract class Parser[T] extends (Input ⇒ ParseResult[T]) {  
  // An unspecified method that defines the parser function.  
  def apply(in : Input): ParseResult  
  
  // A parser combinator for sequential composition  
  def ~ ...  
  
  // A parser combinator for alternative composition  
  def | ...  
}
```

It remains to define concrete combinators that implement this class (see below).

# A Generic Single-Token Parser

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- ▶ The following parser succeeds if the first token in the input satisfies a given predicate `p`.
- ▶ If it succeeds, it reads the token string and returns it as a result.

```
def token(kind: String)(p: Token ⇒ boolean) = new Parser[String] {  
  def apply(in: Input) =  
    if (p(in.head)) Success(in.head.chars, in.tail)  
    else Failure(kind+" expected", in)  
}
```



## Specific Single-Token Parsers

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- ▶ The following parser succeeds if the first token in the input is the given keyword `chars`:
- ▶ If it succeeds, it returns a keyword token as a result.

```
implicit def keyword(chars: String) = token("'" + chars + "'") {  
  case Keyword(chars1) ⇒ chars == chars1  
  case _ ⇒ false  
}
```

- ▶ Note that `keyword` is marked `implicit`, so that you can usually write `"if"` instead of `keyword("if")`.

- ▶ The following parsers succeed if, respectively, the first token in the input is a numeric or string literal, or an identifier.

```
def numericLit = token("number")(_._1.isInstanceOf[NumericLit])  
def stringLit = token("string literal")(_._1.isInstanceOf[StringLit])  
def ident = token("identifier")(_._1.isInstanceOf[Identifier])
```

# The Sequence Combinator

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- ▶ The sequence combinator  $P \sim Q$  succeeds if  $P$  and  $Q$  succeed sequentially, and returns the concatenation of their results.
- ▶  $\sim$  is implemented as a method of class `Parser`.

```
abstract class Parser[T] {  
  def ~ [U](q: Parser[U]) = new Parser[T ~ U] {  
    def apply(in: Input) = Parser.this(in) match {  
      case Success(x, in1) =>  
        q(in1) match {  
          case Success(y, in2) => Success(new ~(x, y), in2)  
          case failure => failure  
        }  
      case failure => failure  
    }  
  }  
}
```

## Concatenating Results

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Normally, the `~` combinator returns the results of the two parsers that are run sequentially wrapped in a `~`-object.

The case class `~` is similar to `Pair` (but `~` turns out to be nicer – see below).

```
case class ~[T, U](_1: T, _2: U) {  
  override def toString = "(" + _1 + " ~ " + _2 + ")"  
}
```

There are also two variants of `~` which return only the left or only the right operand:

```
P <~ Q    // returns only P's result  
P ~> Q    // returns only Q's result
```

# The Alternative Combinator

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- ▶ The alternative combinator  $P \mid Q$  succeeds if either  $P$  or  $Q$  succeeds.
- ▶ It returns the result of  $P$  if  $P$  succeeds, or the result of  $Q$ , if  $Q$  succeeds.
- ▶ The alternative combinator is implemented as a method of class `Parser`.

```
def | (q: => Parser[T]) = new Parser[T] {  
  def apply(in: Input) = Parser.this(in) match {  
    case s1 @ Success(_, _) => s1  
    case failure => q(in)  
  }  
}
```

# Failure And Success Parsers

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- ▶ The parser `failure(msg)` always fails with the given error message. It is implemented as follows:

```
def failure(msg: String) = new Parser[Nothing] {  
  def apply(in: Input) = Failure(msg, in)  
}
```

- ▶ The parser `success(result)` always succeeds with the given result. It does not consume any input. It is implemented as follows:

```
def success[T](result: T) = new Parser[T] {  
  def apply(in: Input) = Success(result, in)  
}
```

## Result Conversion

---

The parser  $P \hat{\hat{}} f$  succeeds iff  $P$  succeeds. In that case it returns the result of applying  $f$  to the result of  $P$ .

```
def ^^ [U](f: T => U) = new Parser[U] {  
  def apply(in: Input) = Parser.this(in) match {  
    case Success(x, in1) => Success(f(x), in1)  
    case f => f  
  }  
}
```

A variant  $\hat{\hat{\hat{}}}$  takes a value  $V$  as right hand side argument.

It returns  $V$  if the left hand parser succeeds:

```
def ^^^ [U](r: U): Parser[U] = ^^ (x => r)
```

# Option and Repetition Combinators

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- ▶ The `opt(P)` combinator always succeeds and returns an `Option` result.  
It returns `Some(R)` iff *P* succeeds with *R*. In case of failure, the result is `None`.
- ▶ The `rep(P)` combinator applies *P* zero or more times until it fails, the result is a list of *P*'s results.

The two combinators are implemented as follows:

```
def opt[T](p: Parser[T]): Parser[Option[T]] =  
  p ^^ Some | success(None)  
def rep[T](p: Parser[T]): Parser[List[T]] =  
  p ~ rep(p) ^^ { case x ~ xs => x :: xs } | success(List())
```

Note that neither of these combinators can fail!



# The Interleaved Repetition Combinator

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The `repsep(P, Q)` parser parses a (possibly empty) sequence

*P Q P ... Q P*

It returns a list of all results returned by *P*.

**Example:** `repsep(ident, ",")` parses a list of identifiers separated by commas.

The `repsep` combinator is implemented as follows:

```
def repsep[T, U](p: Parser[T], q: Parser[U]): Parser[List[T]] =  
  p ~ rep(q ~> p) ^^ { case r ~ rs => r :: rs } | success(List())
```

## Other Combinators

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More combinators can be defined if necessary.

**Exercise:** Implement the  $\text{rep1}(P)$  parser combinator, which applies  $P$  one or more times.

**Exercise:** Define  $\text{opt}$  and  $\text{rep}$  directly, without making use of  $\sim$ ,  $|$ , and  $\text{empty}$ .

## An Example: JSON

---

JSON, or JavaScript Object Notation, is a popular data interchange format.

JSON data essentially consists of objects `{...}`, arrays `[...]`, numbers, and strings.

Here is an example of a JSON value:

```
{ "address book": {  
    "name": "John Smith",  
    "address": { "street": "10 Market Street",  
                 "city"  : "San Francisco, CA",  
                 "zip"   : 94111 },  
    "phone numbers": ["408 338-4238", "408 111-6892"]  
  }  
}
```

# A JSON parser

---

```
package examples.parsing

// import the standard parser class
import scala.util.parsing.combinator.syntactical.StandardTokenParsers

object JSON extends StandardTokenParsers {

  // fix some delimiter symbols ...
  lexical.delimiters += ("{" , "}" , "[" , "]" , ":" )
  // ... and some reserved words
  lexical.reserved += ("null" , "true" , "false")

  // here are the four productions making up the JSON grammar
  def obj  : Parser[Any] = "{" ~ repsep(member, ",") ~ "}"
  def arr  : Parser[Any] = "[" ~ repsep(value, ",") ~ "]"
  def member : Parser[Any] = stringLit ~ ":" ~ value
  def value : Parser[Any] = stringLit | numericLit | obj | arr |
    "null" | "true" | "false"
```

# Testing the JSON Parser

---

Add a method `main` that can be used to test the parser.

```
def main(args: Array[String]) {  
    val tokens = new lexical.Scanner(args(0))  
    println(args(0))  
    println(phrase(value)(tokens))  
}
```

Here are two test runs:

```
>java examples.parsing.JSON "{ \"x\": true, \"y\": [1, 2, 3] }"  
{ x: true, y: [1, 2, 3] }  
[1.26] parsed: ((({ ~ List(((x ~ :) ~ true), ((y ~ :) ~ ([ ~ List(1,  
2, 3)) ~ ])))) ~ })
```

```
>java examples.parsing.JSON "{ \"x\": true \"y\": [1, 2] }"  
{ x: true y: [1, 2] }  
[1.13] failure: unexpected token string literal y  
{ \"x\": true ^ \"y\": [1, 2] }
```

## Getting Better Output

---

- ▶ The result of the previous JSON parser was a tree containing all input tokens (slightly chaotically).
- ▶ We transform the productions using the `^^` combinator to get more structured output:

```
def obj  : Parser[Any] = // return a Map
  "{" ~> rep(member) <~ "}" ^^ (ms => Map() ++ ms)

def arr  : Parser[Any] = // return a List
  "[" ~> rep{value} <~ "]"

def member : Parser[Any] = // return a name/value pair
  stringLit ~ ":" ~ value ^^
  { case name ~ ":" ~ value => (name, value) }
```

**Question:** What are more precise types for `obj`, `arr`, and `member`?

```

def value : Parser[Any] = (
  obj
| arr
| stringLit
| numericLit ^^ (_.toInt) // return an Int
| "null" ^^^ null         // return 'null'
| "true"  ^^^ true        // return 'true'
| "false" ^^^ false       // return 'false'
)

```

If we run the test now, we get:

```

>java examples.parsing.JSON1 "{ \"x\": true, \"y\": [1, 2, 3] }"
{ x: true, y: [1, 2, 3] }
[1.30] parsed: Map(x → true, y → List(1, 2, 3))

```



## Table of Parser Combinators

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ident	identifier
keyword(...)	keyword or special symbol (implicit)
numericLit	integer number
stringLit	string literal
$P \sim Q$	sequential composition
$P <\sim Q, P \sim> Q$	sequential composition; keep left/right only
$P \mid Q$	alternative
opt(P)	option
rep(P)	repetition
repsep(P, Q)	interleaved repetition
$P \hat{\hat{}} f$	result conversion
$P \hat{\hat{\hat{}}} v$	constant result

# Arithmetic Expressions Again

---

Here is the full parser for arithmetic expressions:

```
object Arithmetic extends StandardTokenParsers {  
  lexical.delimiters += List("(", ")", "+", "-", "*", "/")  
  def expr: Parser[Any] = term ~ rep("+ ~ term | - ~ term)  
  def term = factor ~ rep("* ~ factor | / ~ factor)  
  def factor: Parser[Any] = "(" ~ expr ~ ")" | numericLit
```

**Question:** How can we make it evaluate the parsed expression?

# A Problem with Top-Down Parsing

---

As parser combinators work top-down, they don't allow left-recursion.

A production like

```
def expr = expr ~ "—" ~ term
```

would go into an infinite recursion when executed.

The alternative:

```
def expr = term ~ rep("—" ~ term)
```

produces a “right-leaning” tree:  $X - Y - Z$  parses

```
X ~ List("—" ~ Y, "—" ~ Z)
```

But the correct reduction/evaluation of  $+, -, *, /$  is left-leaning!

## Evaluate by FoldLeft

---

We can solve this problem by delaying reduction until all elements of a repetition have been parsed and then performing a fold left on the list:

```
def expr : Parser[Int] =  
  term ~ rep ("+" ~ term | "-" ~ term) ^^ reduceList  
def term : Parser[Int] =  
  factor ~ rep ("*" ~ factor | "/" ~ factor) ^^ reduceList  
def factor: Parser[Int] =  
  "(" ~> expr <~ ")" | numericLit ^^ (_.toInt)
```

Here, `reduceList` is defined in terms of the fold-left operation `/:` ...

```
val reduceList: Expr ~ List[String ~ Expr] => Expr = {  
  case i ~ ps => (i /: ps)(reduce)  
}
```

... and `reduce` is defined as follows:

```
def reduce(x: Int, r: String ~ Int) = r match {  
  case "+" ~ y => x + y  
  case "-" ~ y => x - y  
  case "*" ~ y => x * y  
  case "/" ~ y => x / y  
  case _ => throw new MatchError("illegal case: " + r)  
}
```

With this, we get:

```
java examples.parsing.ArithmeticParsers1 "2 * (3 + 7)"  
2 * (3 + 7)  
[1.12] parsed: 20
```

(Note: the `chain1` combinator in the standard library does this for you.)

# Conclusion

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- ▶ Combinator parsers give a provide a concise, flexible, and high-level way to construct parsers.
- ▶ The token classes of a context free grammar are modeled as primitive parsers.
- ▶ The combination forms are modeled as higher-order parsers.
- ▶ Combinator parsers are an example of an embedded DSL.
- ▶ By contrast, classical parser generators are stand-alone DSLs.
- ▶ Advantage of an embedded DSL over a parser generator: It's easier to connect the results of combinator parsers with the environment.
- ▶ Disadvantage: lower efficiency – but this can be overcome (using packrat parsing)