

## 五、目标代码生成 (16. 指令选择)

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2024 年 06 月 12 日



## Instruction Selection

Instruction Selection is the process of translating LLVM code presented to the code generator into target-specific machine instructions. There are several well-known ways to do this in the literature. LLVM uses a SelectionDAG based instruction selector.

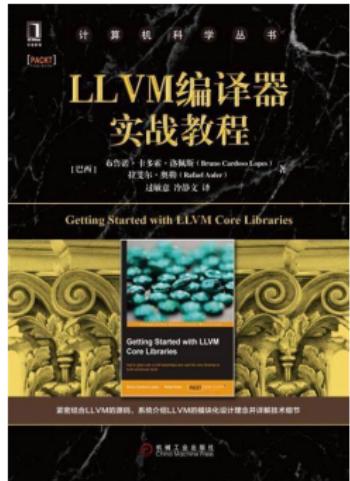
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“略有耳闻”

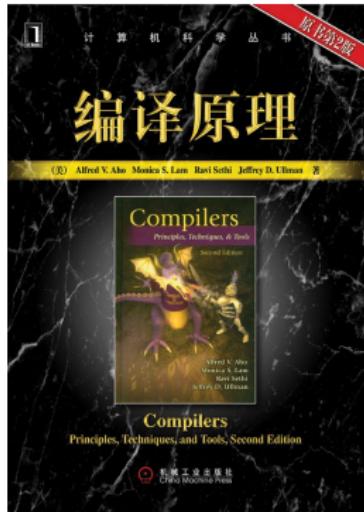
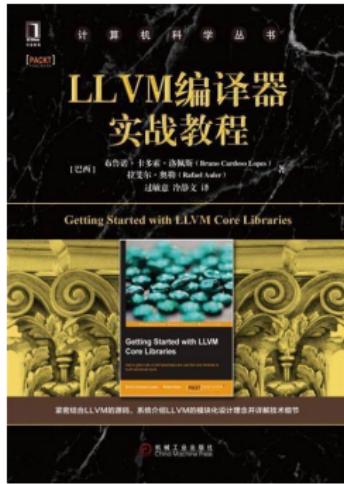
遗憾的是，没有“合适的”教材

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[https://llvm.org/docs/  
CodeGenerator.html](https://llvm.org/docs/CodeGenerator.html)

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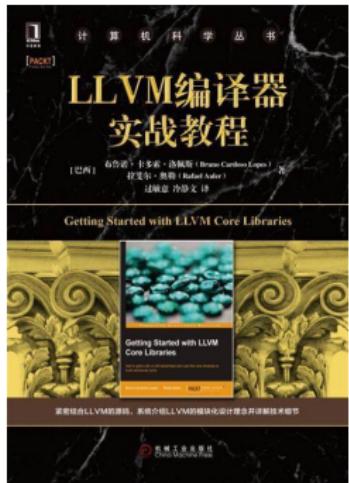


<https://llvm.org/docs/>

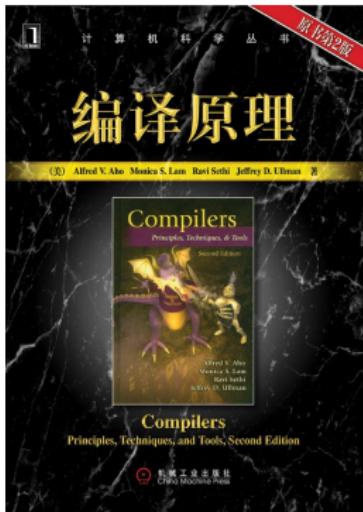
CodeGenerator.html

树翻译方案

遗憾的是，没有“合适的”教材



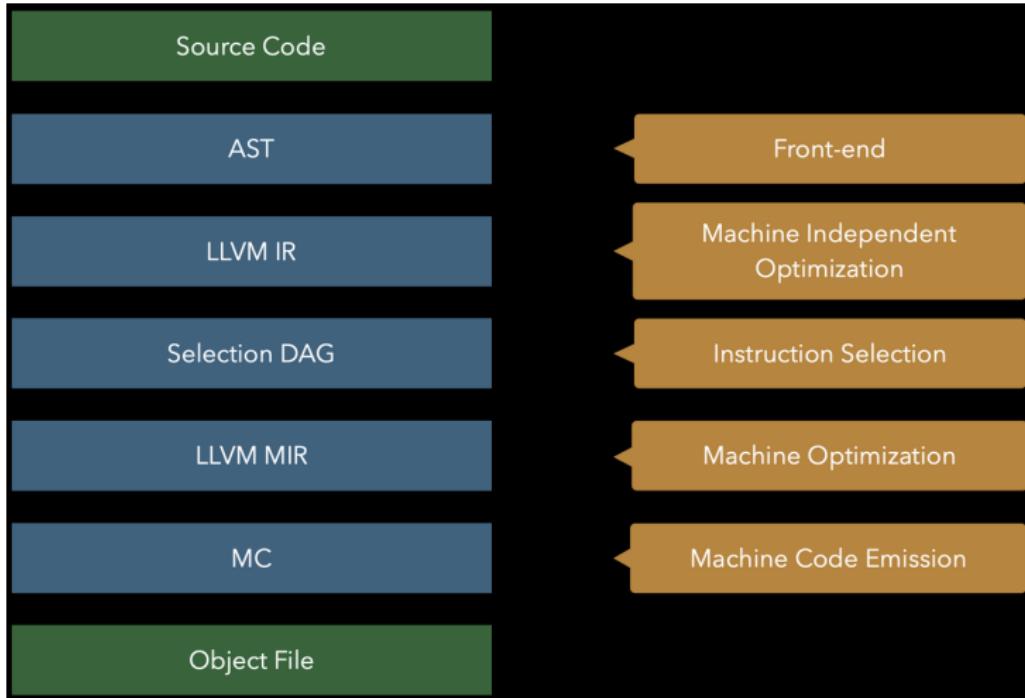
[https://llvm.org/docs/  
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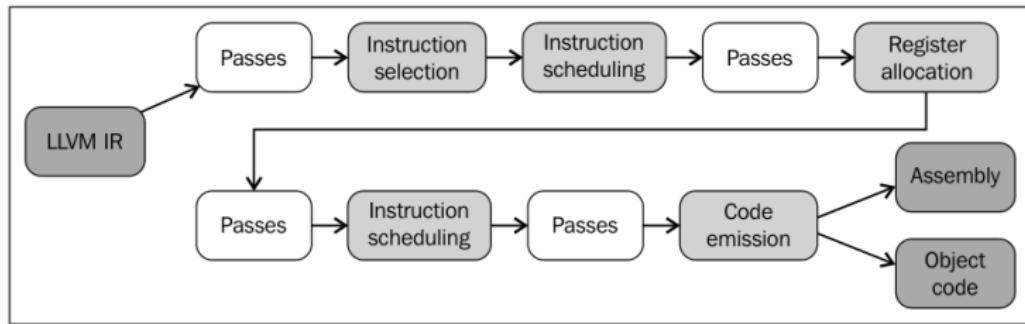
树翻译方案



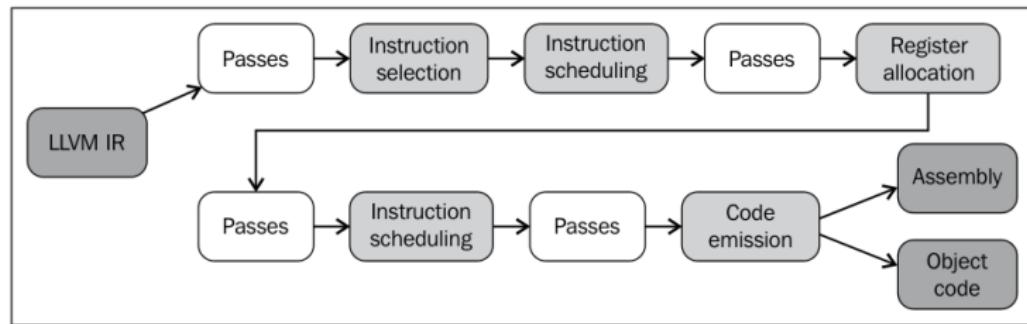
最佳覆盖与最优覆盖算法



(in-memory) LLVM IR    SelectionDAG    MachineInstr    MCInst



(in-memory) LLVM IR    SelectionDAG    MachineInstr    MCInst



Where is “Prologue/Epilogue Insertion”?

# f0-O0 @ Compiler Explorer

The screenshot shows the Compiler Explorer interface with three main panes:

- Left Pane:** C source code for a factorial function:

```
1 int facto
2
3 int main()
4     return
5 }
```
- Middle Pane:** LLVM IR Viewer showing the generated RISC-V LLVM IR:

```
1 define dso_local i32 @main()
2 entry:
3     %retval = add i32 1, 0
4     %argc.addr = getelementptr i32, i32* %argc, i32 0
5     %argv.addr = getelementptr i32*, i32* %argv, i32 0
6     store i32 0, i32 %argc.addr
7     store i32 %argc, i32 %argc.addr
8     store ptr %argc.addr, i32* %argv.addr
9     %call = call i32 @facto(i32 %argc)
10    %mul = mul i32 %call, 1
11    %cmp = icmp eq i32 %mul, 1
12    %conv = zext i32 %cmp, i1
13    ret i32 %conv
14 }
15
16 declare void @facto(i32)
17
18 declare dso_local i32 @facto(i32)
```
- Right Pane:** Opt Pipeline Viewer showing the compilation passes:

Passes:	Function:
RISC-V DAG->DAG Pattern	main
Instruction Selection (riscv-isel)	
Eliminate PHI nodes for register allocation (phi-node-elimination)	
Two-Address instruction pass (twoaddressinstruction)	
Fast Register Allocator (regallocfast)	
Slot index numbering (slotindexes)	
Live Interval Analysis (liveintervals)	
Fast Register Allocator (regallocfast)	
Prologue/Epilogue Insertion & Frame Finalization (prologueepilog)	
Post-RA pseudo instruction expansion pass (postrapseudos)	
SW killed \$x1, \$x2, 28 :: (stor	14
SW killed \$x8, \$x2, 24 :: (stor	15
frame-setup CFI_INSTRUCTION off	16
frame-setup CFI_INSTRUCTION off	17
\$x8 = frame-setup ADDI \$x2, 32	18
frame-setup CFI_INSTRUCTION def	19
renamable \$x12 = COPY \$x0	20
SW killed renamable \$x12, %stack	1
SW killed renamable \$x10, %stack	1
SW killed renamable \$x11, %stack	1
ADDCALLSTACKDOWN 0, 0, implicit	1
SW killed renamable \$x12, \$x8,	21
SW killed renamable \$x10, \$x8,	22
SW killed renamable \$x11, \$x8,	23
renamable \$x10 = ADDI \$x0, 2	24
PseudoCALL target-flags(riscv-c	25
ADDCALLSTACKUP 0, 0, implicit-c	1
renamable \$x11 = COPY killed \$x	26
renamable \$x10 = SLLI renamable	1
renamable \$x10 = SUB killed ren	27
renamable \$x10 = ADDI killed re	28
renamable \$x10 = SLTIU killed r	29
\$x1 = LW \$x2, 28 :: (load (s32)	30
\$x8 = LW \$x2, 24 :: (load (s32)	31
\$x2 = frame-destroy ADDI \$x2, 3	32
PseudoRET implicit killed \$x10;	33
	34
	35

# f0-O1 @ Compiler Explorer

The screenshot shows the Compiler Explorer interface with two main panes. The left pane displays the LLVM IR Viewer for RISC-V code, and the right pane displays the Opt Pipeline Viewer.

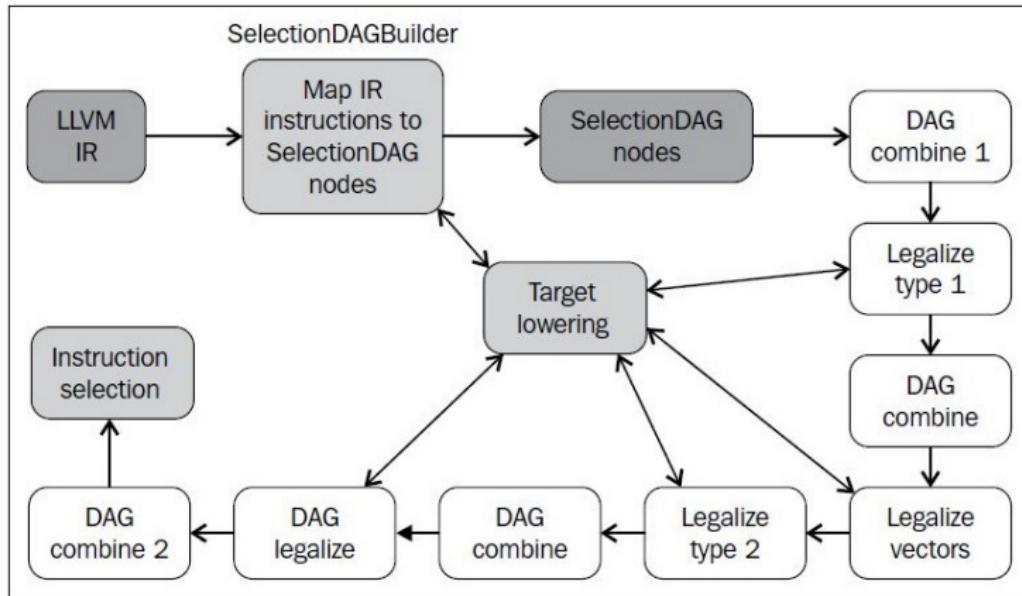
**LLVM IR Viewer (RISC-V rv32gc clang (trunk)):**

```
define dso_local range(i32 0, 2)
entry:
%call = tail call i32 @_factorial(%call)
%cmp = icmp eq i32 %call, 6
%conv = zext i32 %cmp to i32
ret i32 %conv
}
# Machine code for function main:
bb.0.entry:
ADJCALLSTACKDOWN 0, 0, implicit
%2:gpr = ADDI $x0, 2
$x10 = COPY %2:gpr; example.c:4
PseudoCALL target-flags(riscv-c)
ADJCALLSTACKUP 0, 0, implicit-c
%3:gpr = COPY $x10; example.c:4
%4:gpr = ADDI %3:gpr, -6; example.c:4
%5:gpr = SLTIU killed %4:gpr, 1
$x10 = COPY %5:gpr; example.c:4
PseudoRET implicit $x10; example.c:4
# End machine code for function main:
```

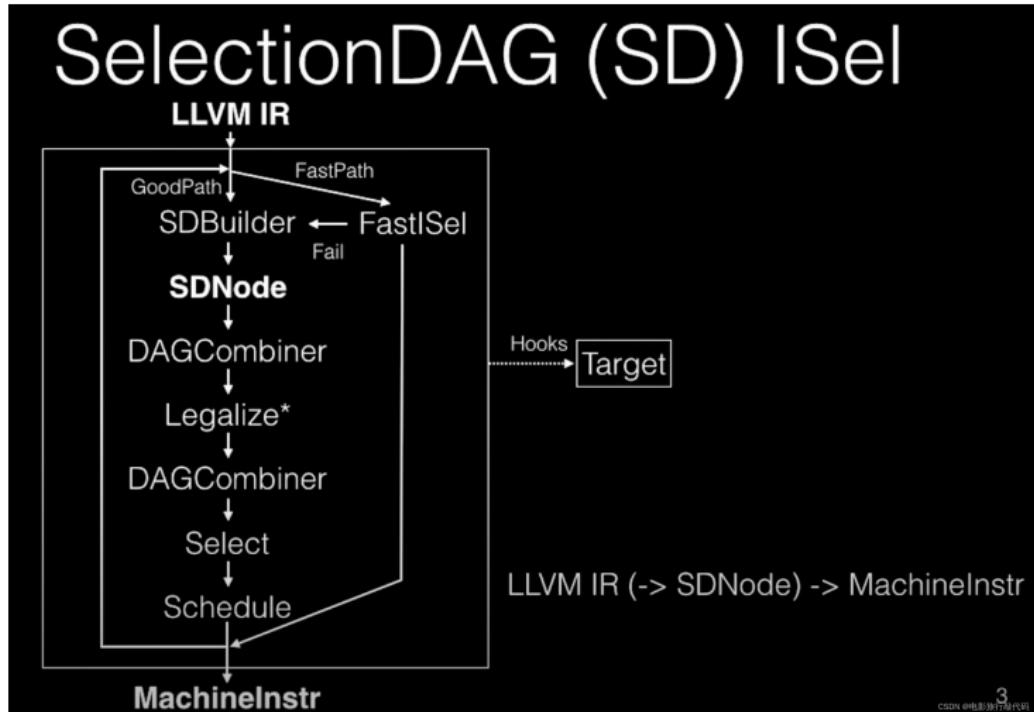
**Opt Pipeline Viewer (RISC-V rv32gc clang (trunk)):**

The pipeline viewer shows the sequence of optimization passes applied to the LLVM IR:

- AssignmentTrackingPass on [module]
- SROAPass on main
- IPSCCPPass on [module]
- GlobalOptPass on [module]
- InstCombinePass on main
- PostOrderFunctionAttrsPass on (main)
- TailCallElimPass on main
- RISC-V DAG->DAG Pattern
- Instruction Selection (riscv-isel)
- Slot index numbering (slotindexes)
- Merge disjoint stack slots (stack-coloring)
- Machine code sinking (machine-sink)
- Live Variable Analysis (livevars)
- Eliminate PHI nodes for register allocation (phi-node-elimination)
- Two-Address instruction pass (twoaddressinstruction)
- Slot index numbering (slotindexes)
- Live Interval Analysis



SDISel    FastISel (per basic block)    GlobalISel (per function)



3  
CSDN @电影旅行时代机

GlobalISel (per function)

```
unsigned int MUL(unsigned long long int x, unsigned int y)
{
    return x * y;
}
```

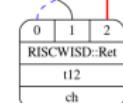
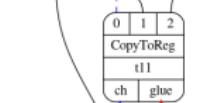
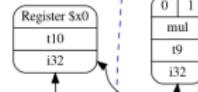
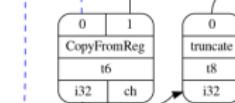
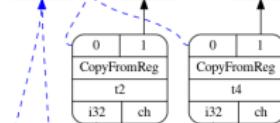
```
unsigned int MUL(unsigned long long int x, unsigned int y)
{
    return x * y;
}
```

```
define dso_local i32 @MUL(i64 %x, i32 %y) local_unnamed_addr #0 {
entry:
    %0 = trunc i64 %x to i32
    %conv1 = mul i32 %0, %y
    ret i32 %conv1
}
```

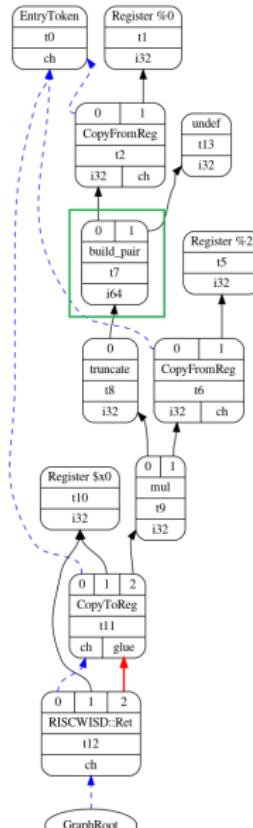
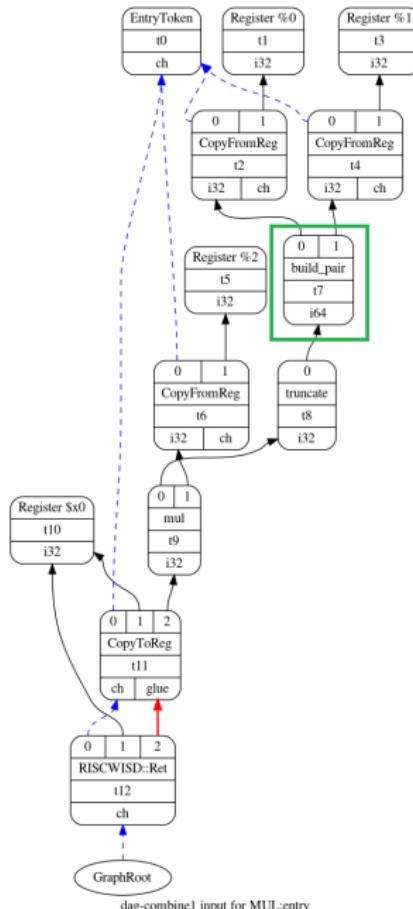
```

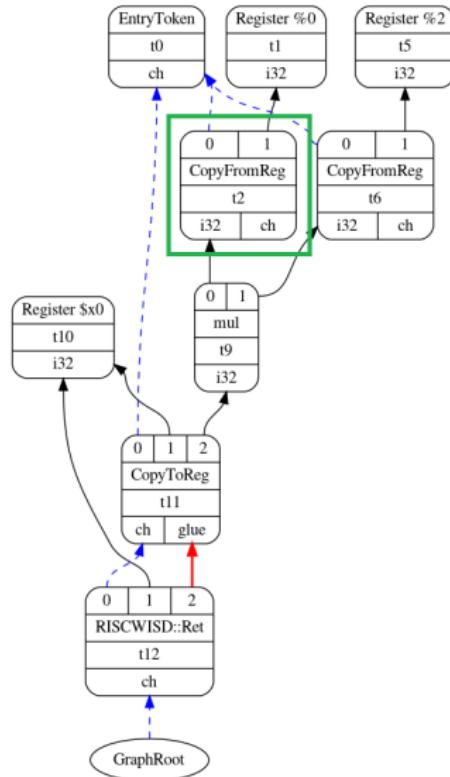
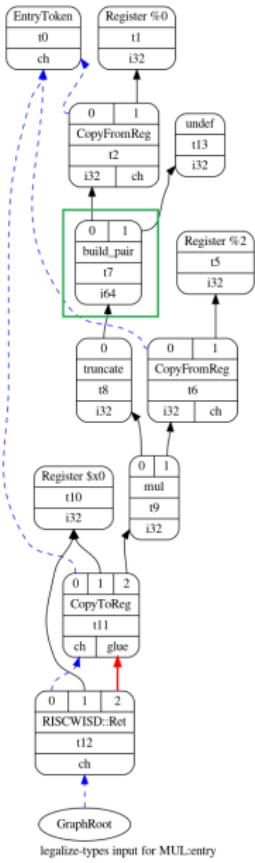
define dso_local i32 @MUL(i64 %x, i32 %y) local_unnamed_addr #0 {
entry:
    %0 = trunc i64 %x to i32
    %conv1 = mul i32 %0, %y
    ret i32 %conv1
}

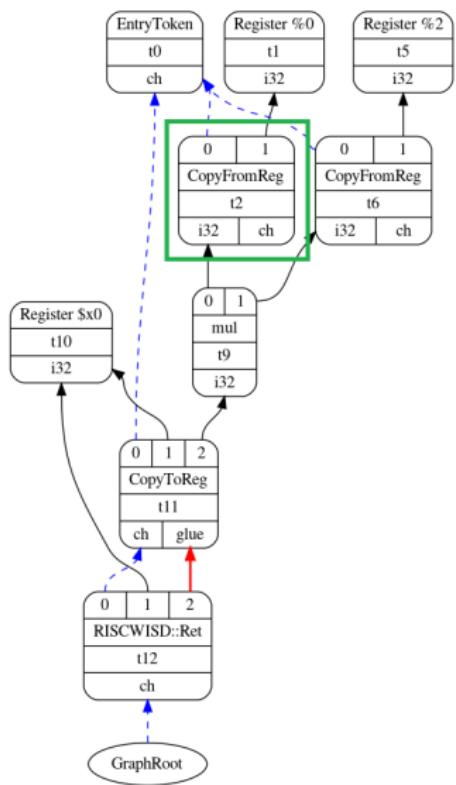
```



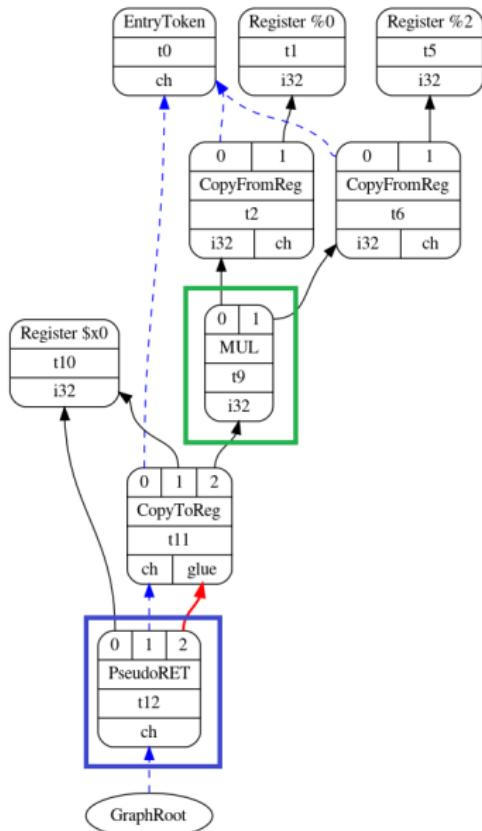
dag-combine1 input for MUL:entry



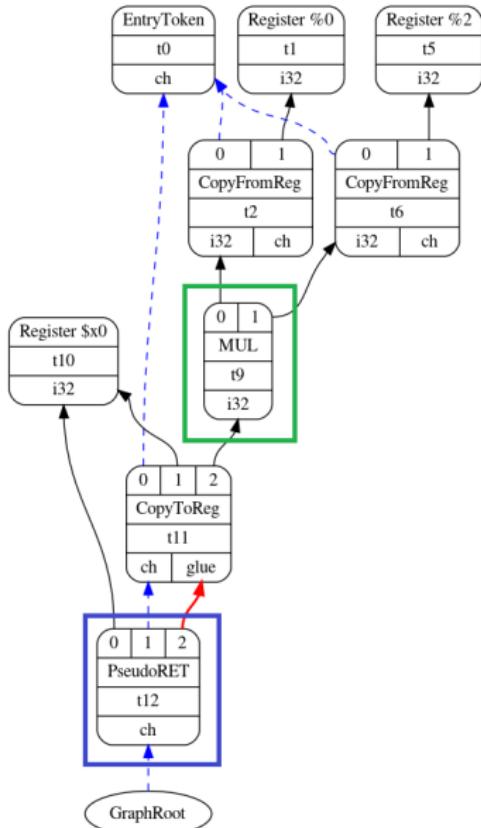




### **dag-combine2 input for MUL:entry**



### scheduler input for MUL:entry



scheduler input for MUL:entry

bb. 0. entry:

liveins: \$x0, \$x2

%2:gpr = COPY \$x2

%0:gpr = COPY \$x0

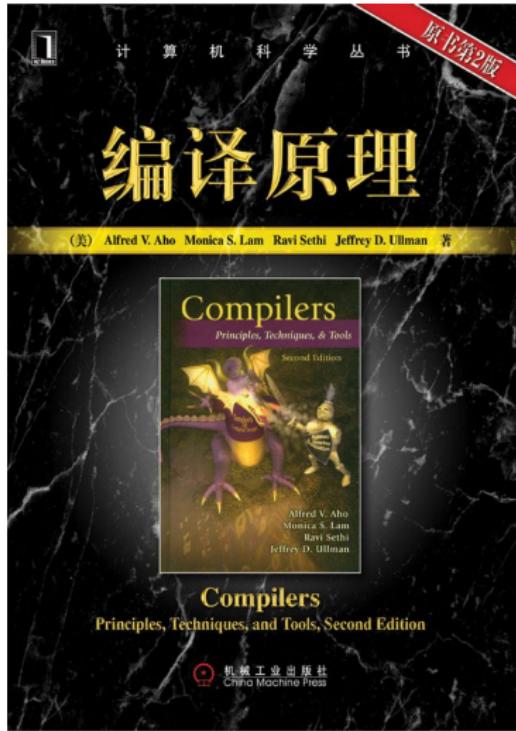
%3:gpr = MUL %0:gpr, %2:gpr

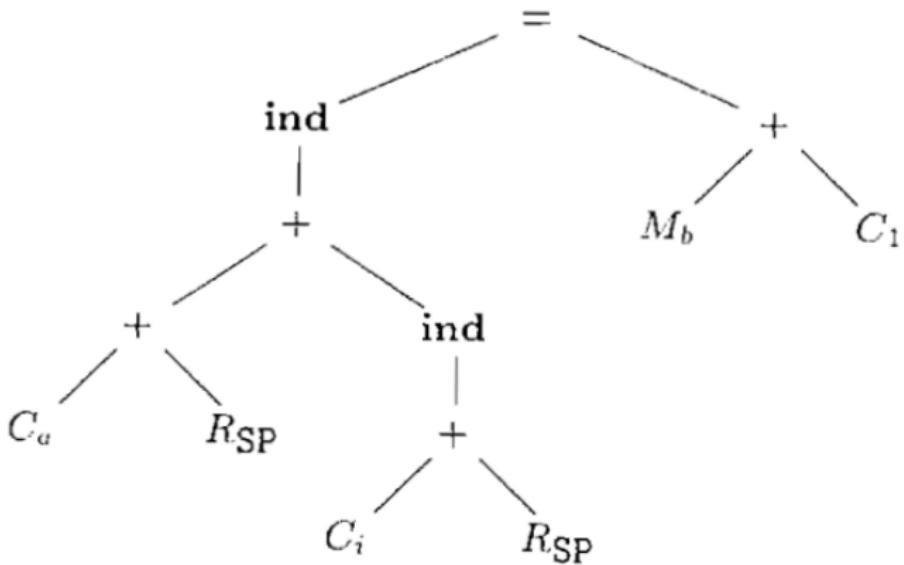
\$x0 = COPY %3:gpr

PseudoRET implicit \$x0

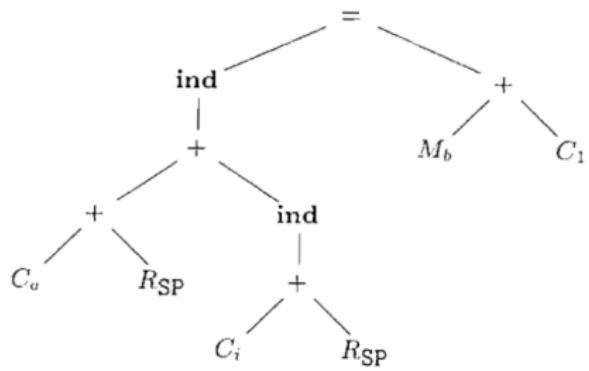
## SelectionDAG Select Phase

The Select phase is the bulk of the target-specific code for instruction selection. This phase takes a legal SelectionDAG as input, pattern matches the instructions supported by the target to this DAG, and produces a new DAG of target code. For example, consider the following LLVM fragment:



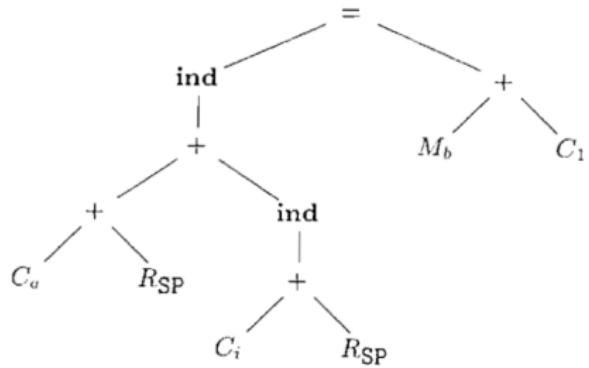


局部变量 a 与 i 的地址: 相对于 SP (栈指针) 的常数偏移量  $C_a$  和  $C_i$



```
LD  R0, #a
ADD R0, R0, SP
ADD R0, R0, i(SP)
LD  R1, b
INC R1
ST  *R0, R1
```

$$R_i \leftarrow \begin{array}{c} + \\ R_i \quad C_a \end{array} \quad \left\{ \begin{array}{l} \text{if } (a = 1) \\ \quad \text{INC } R_i \\ \text{else} \\ \quad \text{ADD } R_i, R_i, \#a \end{array} \right. \}$$



$= \text{ind} + + C_a R_{SP} \text{ ind} + C_i R_{SP} + M_b C_1$

$= \text{ind} + + c_a \text{ sp ind} + c_i \text{ sp} + m_b c_1$

1)	$R_i \rightarrow c_a$	{ LD Ri, #a }
2)	$R_i \rightarrow M_x$	{ LD Ri, x }
3)	$M \rightarrow = M_x R_i$	{ ST x, Ri }
4)	$M \rightarrow = \text{ind} R_i R_j$	{ ST *Ri, Rj }
5)	$R_i \rightarrow \text{ind} + c_a R_j$	{ LD Ri, a(Rj) }
6)	$R_i \rightarrow + R_i \text{ ind} + c_a R_j$	{ ADD Ri, Ri, a(Rj) }
7)	$R_i \rightarrow + R_i R_j$	{ ADD Ri, Ri, Rj }
8)	$R_i \rightarrow + R_i c_1$	{ INC Ri }
9)	$R \rightarrow \text{sp}$	
10)	$M \rightarrow \text{m}$	

二义性: 倾向于执行较大的归约, 而不是较小的归约

归约/归约冲突

移入/归约冲突

优先选择较长的归约

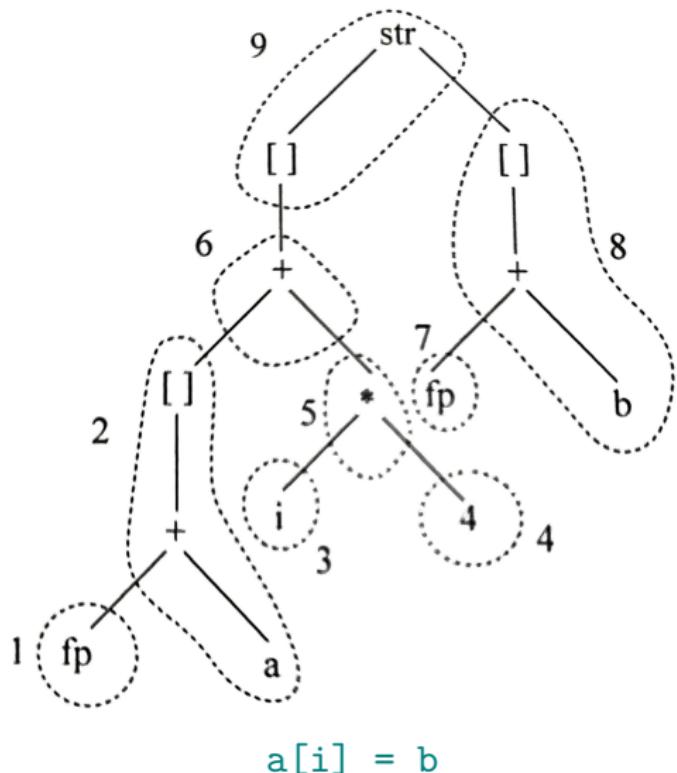
优先选择移入动作



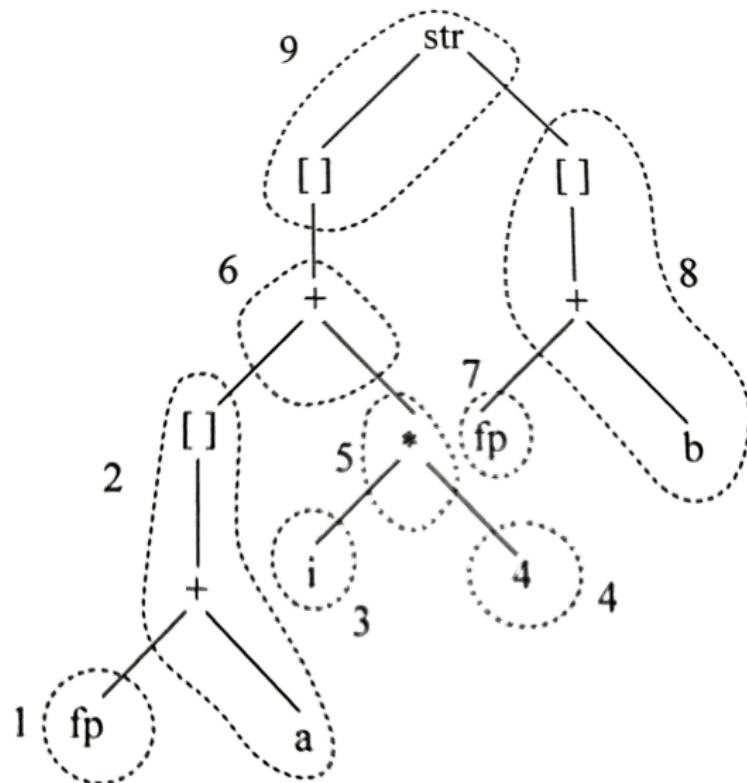
# 算术与存储指令及树型 (鲲鹏处理器)

指    令	树    型
$r_i$	x
add $r_i, r_j, r_k$	$+ \diagup \diagdown$
mul $r_i, r_j, r_k$	$* \diagup \diagdown$
sub $r_i, r_j, r_k$	$- \diagup \diagdown$
div $r_i, r_j, r_k$	$/ \diagup \diagdown$
addi $r_i, r_j, c$	$+ \diagup \quad + \diagup \quad c$ $c \diagdown \quad c \diagdown$
subi $r_i, r_j, c$	$- \diagup \quad - \diagup \quad c$ $c \diagdown \quad c \diagdown$
ldr $r_i, [r_j, c]$	$[] \quad [] \quad [] \quad []$ $c \quad + \quad c \quad + \quad c \quad  $
str $[r_j, c], r_i$	$str \quad str \quad str \quad str$ $[] \quad [] \quad [] \quad []$ $c \quad + \quad c \quad  $

## 瓦片覆盖 (tiling)



```
| ldr r1, [fp, a]  
| addi r2, r0, 4  
| mul r2, ri, r2  
| add r1, r1, r2  
| ldr r2, [fp, b]  
| str [r1, 0], r2
```



`a[i] = b`

目标代码生成

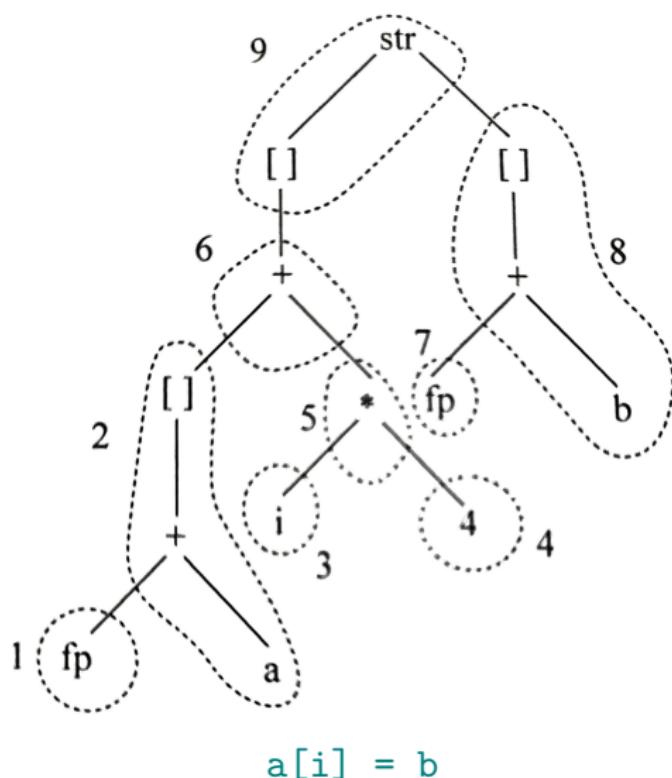
魏恒峰 (hfwei@nju.edu.cn)

A set of small, light-blue navigation icons typically found in presentation software like Beamer. The icons include symbols for back, forward, search, and other document-related functions.

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## 瓦片覆盖 (tiling)



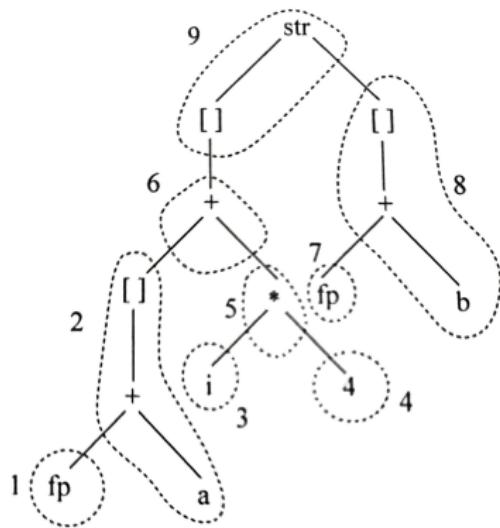
最大吞进算法  
(maximal munch)

```
ldr r1, [fp, a]
addi r2, r0, 4
mul r2, ri, r2
add r1, r1, r2
ldr r2, [fp, b]
str [r1, 0], r2
```

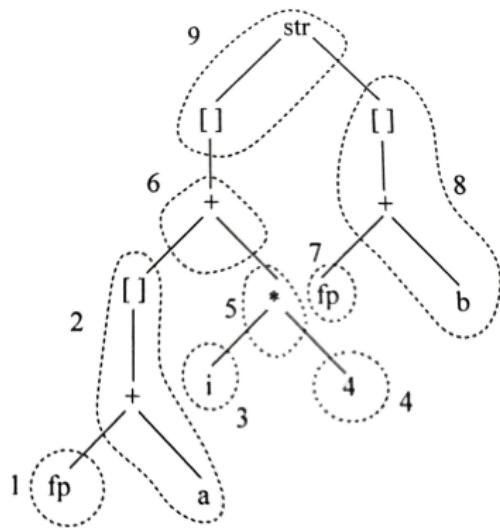
最佳覆盖方案

(不存在可以合并成更小代价树型的相邻树型)

## 使用动态规划思想计算最优覆盖方案

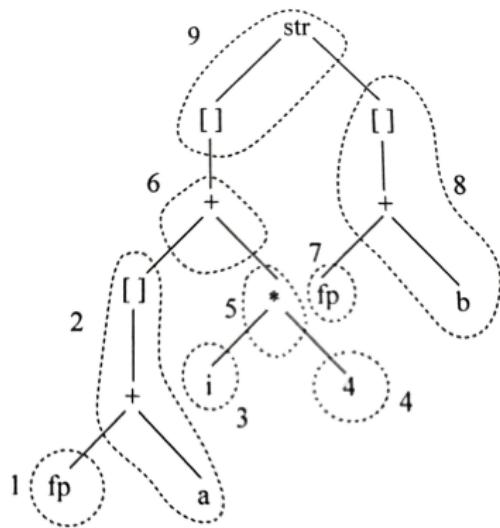


## 使用动态规划思想计算最优覆盖方案



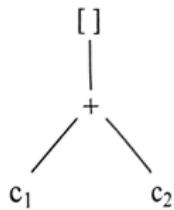
$$\forall n \in Node. c(n) = \min_{t \in tiles(n)} \left( c(t) + \sum_{n_i \in descendants(n)} c(n_i) \right)$$

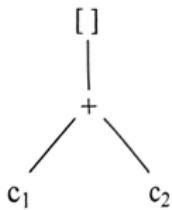
## 使用动态规划思想计算最优覆盖方案



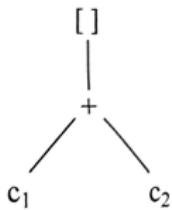
$$\forall n \in Node. c(n) = \min_{t \in tiles(n)} \left( c(t) + \sum_{n_i \in descendants(n)} c(n_i) \right)$$

$$\text{Opt} = c(r) \quad (r \text{ is the root})$$



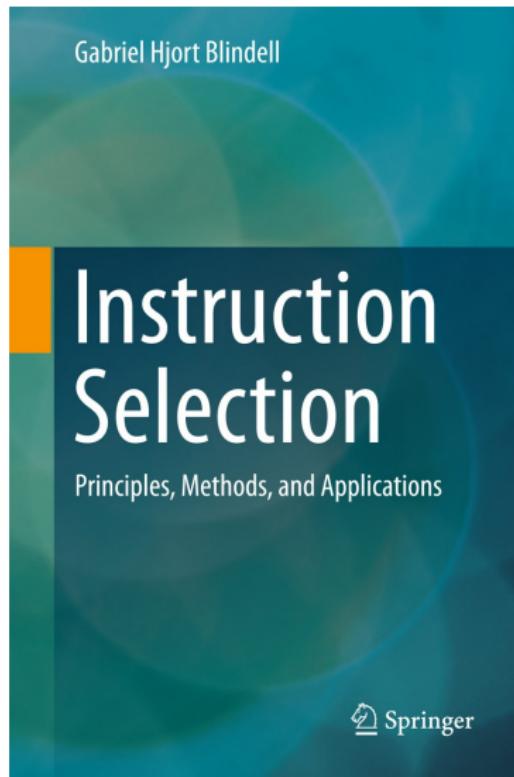


瓦片	指令	瓦片代价	叶子节点代价	总代价
	add	1	$1+1$	3
	addi	1	1	2
	addi	1	1	2



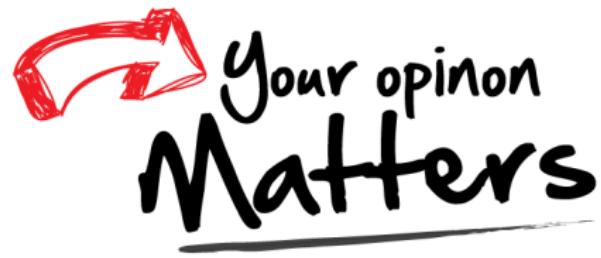
瓦片	指令	瓦片代价	叶子节点代价	总代价
	add	1	1+1	3
	addi	1	1	2
	addi	1	1	2

瓦片	指令	瓦片代价	叶子节点代价	总代价
	ldr	1	2	3
	ldr	1	1	2
	ldr	1	1	2



Tree Covering, DAG Covering, Graph Covering

# Thank You!



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