四、中间代码生成(2. 表达式的翻译)

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| 产生式 | 语义规则 | |
|---------------------------|--|--|
| $S \rightarrow id = E$; | $S.code = E.code \mid \mid$ gen(top.get(id.lexeme) '=' E.addr) | |
| $E \rightarrow E_1 + E_2$ | $\begin{split} E. & addr = \mathbf{new} \ Temp () \\ & E. code = E_1. code \mid \mid E_2. code \mid \mid \\ & gen(E. addr' =' E_1. addr' +' E_2. addr) \end{split}$ | |
| - E ₁ | $ E.addr = \mathbf{new} \ Temp () $ $E.code = E_1.code \mid $ $gen(E.addr'=' \text{ 'minus' } E_1.addr) $ | |
| (E ₁) | $E.addr = E_1.addr$ $E.code = E_1.code$ | |
| id | E.addr = top.get(id.lexeme) E.code = '' | |

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|---------------------------|--|
| $S \rightarrow id = E$; | $S.code = E.code \mid \mid$ gen(top.get(id.lexeme)'=' E.addr) |
| $E \rightarrow E_1 + E_2$ | $ E.addr = \mathbf{new} \ Temp () $ $E.code = E_1.code \parallel E_2.code \parallel \\ gen(E.addr'=' E_1.addr'+' E_2.addr) $ |
| - E ₁ | $ E.addr = \mathbf{new} \ Temp () $ $E.code = E_1.code \parallel $ $gen(E.addr'=' 'minus' E_1.addr) $ |
| (E ₁) | $E.addr = E_1.addr$ $E.code = E_1.code$ |
| id | E.addr = top.get(id.lexeme) E.code = '' |

```
{ gen( top.get(id.lexeme) '=' E.addr); }
                  { gen(L.array.base '[' L.addr ']' '=' E.addr); }
E \rightarrow E_1 + E_2
                \{ E.addr = new Temp() \}
                   gen(E.addr'='E_1.addr'+'E_2.addr);
                  { E.addr = top.get(id.lexeme); }
   \mid L
                  \{E.addr = new Temp();
                   gen(E,addr'=' L.array.base'[' L.addr']'); }
L \rightarrow id [E] \{ L.array = top.get(id.lexeme);
                   L.type = L.array.type.elem;
                   L.addr = new Temp();
                   gen(L.addr'='E.addr'*'L.tupe.width): 
   L_1 [E] \{L.array = L_1.array\}
                   L.type = L_1.type.elem;
                   t = new Temp();
                   L.addr = new Temp();
                   gen(t'='E.addr'*'L.type.width);
                   gen(L.addr'='L_1.addr'+'t);
```

表达式的中间代码翻译

综合属性 E.code: 中间代码

| 产生式 | 语义规则 |
|---------------------------|--|
| $S \rightarrow id = E$; | $S.code = E.code \mid \mid$ $gen(top.get(id.lexeme)'='E.addr)$ |
| $E \rightarrow E_1 + E_2$ | $E.addr = \mathbf{new} \ Temp ()$ $E.code = E_1.code \mid\mid E_2.code \mid\mid$ $gen(E.addr'='E_1.addr'+'E_2.addr)$ |
| - E ₁ | $E.addr = new \ Temp()$ $E.code = E_1.code \mid \mid gen(E.addr'=' 'minus' \ E_1.addr)$ |
| (E ₁) | $E.addr = E_1.addr$ $E.code = E_1.code$ |
| id | E.addr = top.get(id.lexeme) $E.code = ''$ |

综合属性 E.addr: 变量名 (包括临时变量)、常量

表达式的中间代码翻译

综合属性 E.code: 中间代码

| 产生式 | 语义规则 | |
|---------------------------|--|--|
| $S \rightarrow id = E$; | $S.code = E.code \mid \mid$ gen(top.get(id.lexeme)' = 'E.addr) | |
| $E \rightarrow E_1 + E_2$ | $E.addr = \mathbf{new} \ Temp ()$ $E.code = E_1.code \mid\mid E_2.code \mid\mid$ $gen(E.addr'=' E_1.addr'+' E_2.addr)$ | |
| - E ₁ | $E.addr = new \ Temp()$ $E.code = E_1.code \mid gen(E.addr'=' 'minus' \ E_1.addr)$ | |
| (E ₁) | $E.addr = E_1.addr$ $E.code = E_1.code$ | |
| id | E.addr = top.get(id.lexeme) E.code = '' | |

$$a = b + -c$$

$$t_1 = \min c$$

$$t_2 = b + t_1$$

$$a = t_2$$

综合属性 E.addr: 变量名 (包括临时变量)、常量

```
int main() {
  int a = 0, b = 1, c = 2;

a = b + -c;
  return 0;
}
```

```
int main() {
  int a = 0, b = 1, c = 2;
  a = b + -c;
  return 0;
}
```

```
%7 = sub nsw i32 0, %6
%8 = add nsw i32 %5, %7
store i32 %8, i32* %2, align 4
```

A wonderful tool: Compiler Explorer

数组引用的中间代码翻译

声明: int a[2][3]

数组引用: x = a[1][2]; a[1][2] = x

数组引用的中间代码翻译

声明: int a[2][3]

数组引用: x = a[1][2]; a[1][2] = x

需要计算 a[1][2] 相对于数组基地址 a 的偏移地址

$$addr(a[1][2]) = base + 1 \times 12 + 2 \times 4$$

| | 类型 | 宽度 |
|---------|-----------------------------|----|
| a | array(2, array(3, integer)) | 24 |
| a[i] | array(3, integer) | 12 |
| a[i][j] | integer | 4 |

int a[2][3]

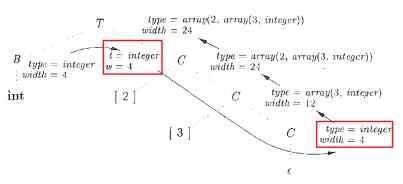
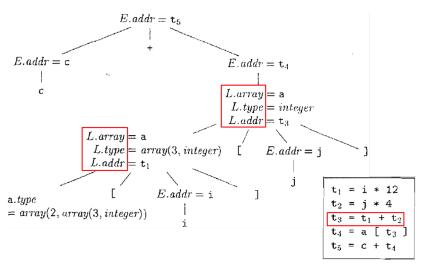


图 6-16 数组类型的语法制导翻译

综合属性 L.array(.base): 数组基地址 (即,数组名)

```
S \rightarrow id = E; { gen(top.get(id.lexeme)' = 'E.addr); }
     L = E; { gen(L.array.base'['|L.addr']'' = 'E.addr); }
E \rightarrow E_1 + E_2 + E.addr = new Temp();
                    gen(E.addr'='E_1.addr'+'E_2.addr); 
      id
                  \{E.addr = top.get(id.lexeme);\}
    \mid L \mid
                  \{E.addr = new\ Temp();
                    gen(E.addr'='L.array base'['L.addr']'); \}
                  \{ L.array = top.get(id.lexeme); 
                    L.type = L.array.type.elem;
                    L.addr = new Temp():
                    gen(L.addr'='E.addr'*'L.type.width); 
      L_i [ E ]
                  \{ L.array = L_1.array; \}
                    L.type = L_1.type.elem;
                    t = new Temp():
                    L.addr = new Temp();
                    gen(t'='E.addr'*'L.type.width);
                    qen(L.addr'='L_1.addr'+'|t);
```



$$c + a[i][j]$$

```
int main() {
   int a[2][3] = { 0 };

int i = 1, j = 2;
   int c = 10, d = 20;

   d = c + a[i][j];
   return 0;
}
```

%2 = alloca [2 x [3 x i32]], align 16

```
int main() {
   int a[2][3] = { 0 };

int i = 1, j = 2;
   int c = 10, d = 20;

   d = c + a[i][j];

   return 0;
}
```

%2 = alloca [2 x [3 x i32]], align 16

int $a[2][3] = { 0 };$

int i = 1, j = 2; int c = 10, d = 20:

d = c + a[i][i];

int main() {

```
return 0;
}

%8 = load i32, i32* %5, align 4 %8:c
%9 = load i32, i32* %3, align 4 %9:i
%10 = sext i32 %9 to i64
%11 = [getelementptr] inbounds [2 x [3 x i32]], [2 x [3 x i32]]* %2, i64 0, i64 %10
%12 = load i32, i32* %4, align 4 %12:j
%13 = sext i32 %12 to i64
%14 = [getelementptr] inbounds [3 x i32], [3 x i32]* %11, i64 0, i64 %13
%15 = load i32, i32* %14, align 4 %15:a[i][]
%16 = add nsw i32 %8, %15
store i32 %16, i32* %6, align 4
```

GEP provides a way to access arrays and manipulate pointers.

GEP abstract away details like size of types.

getelementptr

<base-type>, <base-type>* <base-addr>, [i32 <index>]+

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<base-type>: base type used for the first index

- ▶ It does *not* change the pointer type.
- ► It offsets by the <base-type>.

getelementptr

<base-type>, <base-type>* <base-addr>, [i32 <index>]+

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- ▶ It does *not* change the pointer type.
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Further indices:

- ► Offset inside arrays (aggregate types)
- ► Change the pointer type by removing one layer of "aggregation"

The Often Misunderstood GEP Instruction @ LLVM Docs

The Often Misunderstood GEP Instruction @ LLVM Docs



LLVM IR Tutorial: Phis, GEPs and Other Things, Oh My! @ bilibili

Thank You!



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