

$termvar, x, y, z, f$   
 $typevar, X, Y, Z$   
 $index, i, j, k$   
 $t, c, v, s, n$

$::=$		term
	$x$	variable
	$triv$	unit
	$squash_S$	injection of the retract
	$split_S$	surjection of the retract
	$box_C$	generalize to the untyped universe
	$unbox_C$	specialize the untyped universe to a specific type
	$\Lambda X.t$	
	$[A]t$	
	$\langle A \rangle t$	type cast
	$\lambda x : A. t$	$\lambda$ -abstraction
	$t_1 t_2$	function application
	$(t_1, t_2)$	pair constructor
	$fst\ t$	first projection
	$snd\ t$	second projection
	$succ\ t$	successor function
	$0$	zero
	$case\ t\ of\ t_1 \parallel x.t_2$	
	$(t)$	S

$h$

$::=$		head-normal forms
	$triv$	
	$split_S$	
	$squash_S$	
	$box_C$	
	$unbox_C$	
	$\lambda x : A. t$	
	$(t_1, t_2)$	
	$fst\ t$	
	$snd\ t$	
	$succ\ t$	
	$0$	

$T$

$::=$		terminating types
	$Unit$	unit type
	$Nat$	natural number type
	$T_1 \rightarrow T_2$	function type
	$T_1 \times T_2$	cartesian product type
	$(T)$	S

$K$

$::=$	
	$\star$

$A, B, C, D, E, S, U$

$::=$		type
	$X$	
	$\forall X. A$	
	$Unit$	unit type

	1	
	<b>Nat</b>	natural number type
	?	untyped universe
	$A_1 \rightarrow A_2$	function type
	$A_1 \times A_2$	cartesian product type
	$(A)$	S

$\Gamma$	::=	typing context
	.	empty context
	$\Gamma, X : \star$	
	$\Gamma, x : A$	cons

$vd$	::=	
	$\vdash$	
	$\nvdash$	

$\boxed{\Gamma \vdash A : \star}$

$\frac{X : \star \in \Gamma}{\Gamma \vdash X : \star}$	K_VAR
$\frac{}{\Gamma \vdash \mathbf{Unit} : \star}$	K_UNIT
$\frac{}{\Gamma \vdash \mathbf{Nat} : \star}$	K_NAT
$\frac{}{\Gamma \vdash ? : \star}$	K_UNITYPE
$\frac{\Gamma \vdash A : \star \quad \Gamma \vdash B : \star}{\Gamma \vdash A \rightarrow B : \star}$	K_ARROW
$\frac{\Gamma \vdash A : \star \quad \Gamma \vdash B : \star}{\Gamma \vdash A \times B : \star}$	K_PROD
$\frac{\Gamma, X : \star \vdash A : \star}{\Gamma \vdash \forall X. A : \star}$	K_FORALL

$\boxed{\Gamma \text{ Ok}}$

$\frac{}{\cdot \text{Ok}}$	OK_EMPTY
$\frac{\Gamma \text{ Ok}}{(\Gamma, X : \star) \text{ Ok}}$	OK_TYPEVAR
$\frac{\Gamma \text{ Ok} \quad \Gamma \vdash A : \star}{(\Gamma, x : A) \text{ Ok}}$	OK_VAR

$\boxed{A \sim B}$   $A$  is consistent with  $B$

$\frac{}{A \sim A}$	REFL
$\frac{}{A \sim ?}$	BOX
$\frac{}{? \sim A}$	UNBOX
$\frac{A_1 \sim A_2 \quad B_1 \sim B_2}{A_1 \rightarrow B_1 \sim A_2 \rightarrow B_2}$	ARROW

$$\frac{A_1 \sim A_2 \quad B_1 \sim B_2}{A_1 \times B_1 \sim A_2 \times B_2} \quad \text{PROD}$$

$\boxed{\Gamma vdt : A}$      $t$  has type  $A$  in context  $\Gamma$

$$\begin{array}{c} \frac{x : A \in \Gamma}{\Gamma \vdash x : A} \quad \text{VAR} \\[10pt] \frac{x : A \in \Gamma \quad \Gamma \text{Ok}}{\Gamma \vdash x : A} \quad \text{VARP} \\[10pt] \overline{\Gamma \vdash \text{box}_C : C \rightarrow ?} \quad \text{BOX} \\[10pt] \frac{\Gamma \text{Ok}}{\Gamma \vdash \text{box}_C : C \rightarrow ?} \quad \text{BoxP} \\[10pt] \overline{\Gamma \vdash \text{unbox}_C : ? \rightarrow C} \quad \text{UNBOX} \\[10pt] \frac{\Gamma \text{Ok}}{\Gamma \vdash \text{unbox}_C : ? \rightarrow C} \quad \text{UNBOXP} \\[10pt] \overline{\Gamma \vdash \text{Box}_A : A \rightarrow ?} \quad \text{BOXG} \\[10pt] \overline{\Gamma \vdash \text{Unbox}_A : ? \rightarrow A} \quad \text{UNBOXG} \\[10pt] \overline{\Gamma \vdash \text{squash}_U : U \rightarrow ?} \quad \text{SQUASH} \\[10pt] \frac{\Gamma \text{Ok}}{\Gamma \vdash \text{squash}_U : U \rightarrow ?} \quad \text{SQUASHP} \\[10pt] \overline{\Gamma \vdash \text{split}_U : ? \rightarrow U} \quad \text{SPLIT} \\[10pt] \frac{\Gamma \text{Ok}}{\Gamma \vdash \text{split}_U : ? \rightarrow U} \quad \text{SPLITP} \\[10pt] \overline{\Gamma \vdash \text{triv} : \text{Unit}} \quad \text{UNIT} \\[10pt] \frac{\Gamma \text{Ok}}{\Gamma \vdash \text{triv} : \text{Unit}} \quad \text{UNITP} \\[10pt] \overline{\Gamma \vdash 0 : \text{Nat}} \quad \text{ZERO} \\[10pt] \frac{\Gamma \text{Ok}}{\Gamma \vdash 0 : \text{Nat}} \quad \text{ZEROP} \\[10pt] \frac{\Gamma \vdash t : \text{Nat}}{\Gamma \vdash \text{succ } t : \text{Nat}} \quad \text{SUCC} \\[10pt] \frac{\Gamma \vdash t : \text{Nat} \quad \Gamma \vdash t_1 : A \quad \Gamma, x : \text{Nat} \vdash t_2 : A}{\Gamma \vdash \text{case } t \text{ of } t_1 \parallel x. t_2 : A} \quad \text{CASE} \\[10pt] \frac{\Gamma \vdash t_1 : A_1 \quad \Gamma \vdash t_2 : A_2}{\Gamma \vdash (t_1, t_2) : A_1 \times A_2} \quad \text{PAIR} \\[10pt] \frac{\Gamma \vdash t : A_1 \times A_2}{\Gamma \vdash \text{fst } t : A_1} \quad \text{FST} \end{array}$$

$$\begin{array}{c}
\frac{\Gamma \vdash t : A_1 \times A_2}{\Gamma \vdash \text{snd } t : A_2} \quad \text{SND} \\
\\
\frac{\Gamma, x : A \vdash t : B}{\Gamma \vdash \lambda x : A_1. t : A \rightarrow B} \quad \text{LAM} \\
\\
\frac{\Gamma \vdash t_1 : A \rightarrow B \quad \Gamma \vdash t_2 : A}{\Gamma \vdash t_1 t_2 : B} \quad \text{APP} \\
\\
\frac{\Gamma, X : \star \vdash t : A}{\Gamma \vdash \Lambda X. t : \forall X. A} \quad \text{LAM} \\
\\
\frac{\Gamma \vdash t : \forall X. B \quad \Gamma \vdash A : \star}{\Gamma \vdash [A]t : [A/X]B} \quad \text{TYPEAPP} \\
\\
\frac{\Gamma \vdash t : ?}{\Gamma \vdash \text{succ } t : ?} \quad \text{SUCCU} \\
\\
\frac{\Gamma \vdash t : ?}{\Gamma \vdash \text{fst } t : ?} \quad \text{FSTU} \\
\\
\frac{\Gamma \vdash t : ?}{\Gamma \vdash \text{snd } t : ?} \quad \text{SNDU} \\
\\
\frac{\Gamma \vdash t_1 : ? \quad \Gamma \vdash t_2 : A}{\Gamma \vdash t_1 t_2 : ?} \quad \text{APPU} \\
\\
\frac{\Gamma \vdash t_1 : A_1 \rightarrow B \quad \Gamma \vdash t_2 : A_2 \quad A_1 \sim A_2}{\Gamma \vdash t_1 t_2 : B} \quad \text{APPC} \\
\\
\frac{\Gamma \vdash t : A_1 \times B \quad A_1 \sim A_2}{\Gamma \vdash \text{fst } t : A_2} \quad \text{FSTC} \\
\\
\frac{\Gamma \vdash t : A \times B_1 \quad B_1 \sim B_2}{\Gamma \vdash \text{snd } t : B_2} \quad \text{SND C} \\
\\
\frac{\Gamma \vdash t : A \quad A \sim B}{\Gamma \vdash \langle B \rangle t : B} \quad \text{CAST}
\end{array}$$

$\boxed{\Gamma \vdash t_1 \rightsquigarrow t_2 : A}$   $t_1$  reduces to  $t_2$  with type  $A$  in context  $\Gamma$

$$\begin{array}{c}
\frac{\Gamma \vdash s : A}{\Gamma \vdash s \rightsquigarrow s : A} \quad \text{RD\_VALUES} \\
\\
\frac{\Gamma \vdash t : C}{\Gamma \vdash \text{unbox}_C(\text{box}_C t) \rightsquigarrow t : C} \quad \text{RD\_RETRACT} \\
\\
\frac{\Gamma \vdash t : A}{\Gamma \vdash \text{Unbox}_A(\text{Box}_A t) \rightsquigarrow t : A} \quad \text{RD\_RETRACTG} \\
\\
\frac{\Gamma \vdash t : U}{\Gamma \vdash \text{split}_U(\text{squash}_U t) \rightsquigarrow t : U} \quad \text{RD\_RETRACTU} \\
\\
\frac{\Gamma \vdash t \rightsquigarrow t' : \text{Nat}}{\Gamma \vdash \text{succ } t \rightsquigarrow \text{succ } t' : \text{Nat}} \quad \text{RD\_SUCC} \\
\\
\frac{\Gamma \vdash t_1 : A \quad \Gamma, x : \text{Nat} \vdash t_2 : A}{\Gamma \vdash \text{case } 0 \text{ of } t_1 \parallel x. t_2 \rightsquigarrow t_1 : A} \quad \text{RD\_CASE0}
\end{array}$$

$$\begin{array}{c}
\frac{\Gamma \vdash t : \text{Nat} \quad \Gamma \vdash t_1 : A \quad \Gamma, x : \text{Nat} \vdash t_2 : A}{\Gamma \vdash \text{case}(\text{succ } t) \text{ of } t_1 \parallel x.t_2 \rightsquigarrow [t/x]t_2 : A} \text{RD\_CASESUCC} \\
\\
\frac{\Gamma \vdash t \rightsquigarrow t' : \text{Nat} \quad \Gamma \vdash t_1 : A \quad \Gamma, x : \text{Nat} \vdash t_2 : A}{\Gamma \vdash \text{case } t \text{ of } t_1 \parallel x.t_2 \rightsquigarrow \text{case } t' \text{ of } t_1 \parallel x.t_2 : A} \text{RD\_CASE1} \\
\\
\frac{\Gamma \vdash t : \text{Nat} \quad \Gamma \vdash t_1 \rightsquigarrow t'_1 : A \quad \Gamma, x : \text{Nat} \vdash t_2 : A}{\Gamma \vdash \text{case } t \text{ of } t_1 \parallel x.t_2 \rightsquigarrow \text{case } t \text{ of } t'_1 \parallel x.t_2 : A} \text{RD\_CASE2} \\
\\
\frac{\Gamma \vdash t : \text{Nat} \quad \Gamma \vdash t : A \quad \Gamma, x : \text{Nat} \vdash t_2 \rightsquigarrow t'_2 : A}{\Gamma \vdash \text{case } t \text{ of } t_1 \parallel x.t_2 \rightsquigarrow \text{case } t \text{ of } t_1 \parallel x.t'_2 : A} \text{RD\_CASE3} \\
\\
\frac{\Gamma, x : A_1 \vdash t_2 : A_2 \quad \Gamma \vdash t_1 : A_1}{\Gamma \vdash (\lambda x : A_1.t_2) t_1 \rightsquigarrow [t_1/x]t_2 : A_2} \text{RD\_BETA} \\
\\
\frac{\Gamma \vdash t : A_1 \rightarrow A_2 \quad x \notin \text{FV}(t)}{\Gamma \vdash \lambda x : A_1.t x \rightsquigarrow t : A_1 \rightarrow A_2} \text{RD\_ETA} \\
\\
\frac{\Gamma \vdash t_1 : A_1 \quad \Gamma \vdash t_2 : A_2}{\Gamma \vdash \text{fst}(t_1, t_2) \rightsquigarrow t_1 : A_1} \text{RD\_PROJ1} \\
\\
\frac{\Gamma \vdash t_1 : A_1 \quad \Gamma \vdash t_2 : A_2}{\Gamma \vdash \text{snd}(t_1, t_2) \rightsquigarrow t_2 : A_2} \text{RD\_PROJ2} \\
\\
\frac{\Gamma \vdash t : A_1 \times A_2}{\Gamma \vdash (\text{fst } t, \text{snd } t) \rightsquigarrow t : A_1 \times A_2} \text{RD\_ETAP} \\
\\
\frac{\Gamma, x : A_1 \vdash t \rightsquigarrow t' : A_2}{\Gamma \vdash \lambda x : A_1.t \rightsquigarrow \lambda x : A_1.t' : A_1 \rightarrow A_2} \text{RD\_LAM} \\
\\
\frac{\Gamma \vdash t_1 \rightsquigarrow t'_1 : A_1 \rightarrow A_2 \quad \Gamma \vdash t_2 : A_1}{\Gamma \vdash t_1 t_2 \rightsquigarrow t'_1 t_2 : A_2} \text{RD\_APP1} \\
\\
\frac{\Gamma \vdash t_1 : A_1 \rightarrow A_2 \quad \Gamma \vdash t_2 \rightsquigarrow t'_2 : A_1}{\Gamma \vdash t_1 t_2 \rightsquigarrow t_1 t'_2 : A_2} \text{RD\_APP2} \\
\\
\frac{\Gamma \vdash t \rightsquigarrow t' : A_1 \times A_2}{\Gamma \vdash \text{fst } t \rightsquigarrow \text{fst } t' : A_1} \text{RD\_FST} \\
\\
\frac{\Gamma \vdash t \rightsquigarrow t' : A_1 \times A_2}{\Gamma \vdash \text{snd } t \rightsquigarrow \text{snd } t' : A_2} \text{RD\_SND} \\
\\
\frac{\Gamma \vdash t_1 \rightsquigarrow t'_1 : A_1 \quad \Gamma \vdash t_2 : A_2}{\Gamma \vdash (t_1, t_2) \rightsquigarrow (t'_1, t_2) : A_1 \times A_2} \text{RD\_PAIR1} \\
\\
\frac{\Gamma \vdash t_1 : A_1 \quad \Gamma \vdash t_2 \rightsquigarrow t'_2 : A_2}{\Gamma \vdash (t_1, t_2) \rightsquigarrow (t_1, t'_2) : A_1 \times A_2} \text{RD\_PAIR2} \\
\\
\frac{\Gamma, X : \star \vdash t : B}{\Gamma \vdash [A](\Lambda X.t) \rightsquigarrow [A/X]t : [A/X]B} \text{RD\_TYPEBETA} \\
\\
\frac{\Gamma \vdash t_1 \rightsquigarrow t_2 : \forall X.B}{\Gamma \vdash [A]t_1 \rightsquigarrow [A]t_2 : [A/X]B} \text{RD\_TYPEAPP} \\
\\
\frac{\Gamma, X : \star \vdash t_1 \rightsquigarrow t_2 : A}{\Gamma \vdash \Lambda X.t_1 \rightsquigarrow \Lambda X.t_2 : \forall X.A} \text{RD\_LAM}
\end{array}$$

$\boxed{\Gamma \vdash t_1 \rightsquigarrow t_2 : A}$ 

Reduction for annotated Siek16

$$\begin{array}{c}
 \frac{\Gamma \vdash v : A}{\Gamma \vdash v \rightsquigarrow v : A} \quad \text{RDA\_VALUES} \\
 \frac{\Gamma \vdash \text{drop-cast } v : C}{\Gamma \vdash \langle C \rangle v \rightsquigarrow \text{drop-cast } v : C} \quad \text{RDA\_CASTA} \\
 \frac{\Gamma \vdash t : ?}{\Gamma \vdash \langle \text{Nat} \rangle (\text{succ } t) \rightsquigarrow \text{succ } \langle \text{Nat} \rangle t : \text{Nat}} \quad \text{RDA\_CASTNAT} \\
 \frac{\Gamma \vdash t : A_1 \rightarrow B_1 \quad (A_1 \rightarrow B_1) \sim (A_2 \rightarrow B_2)}{\Gamma \vdash \langle A_2 \rightarrow B_2 \rangle t \rightsquigarrow \lambda y : A_2. \langle B_2 \rangle (t \langle A_1 \rangle y) : A_2 \rightarrow B_2} \quad \text{RDA\_CASTARROW} \\
 \frac{\Gamma \vdash t : A_1 \times B_1 \quad (A_1 \times B_1) \sim (A_2 \times B_2)}{\Gamma \vdash \langle A_2 \times B_2 \rangle t \rightsquigarrow (\langle A_2 \rangle (\text{fst } t), \langle B_2 \rangle (\text{snd } t)) : A_2 \times B_2} \quad \text{RDA\_CASTPAIR} \\
 \frac{\Gamma \vdash t_1 \rightsquigarrow t_2 : A \quad A \sim B}{\Gamma \vdash \langle B \rangle t_1 \rightsquigarrow \langle B \rangle t_2 : B} \quad \text{RDA\_CAST} \\
 \frac{\Gamma, x : A_1 \vdash t_2 : A_2 \quad \Gamma \vdash t_1 : A_1}{\Gamma \vdash (\lambda x : A_1. t_2) t_1 \rightsquigarrow [t_1/x] t_2 : A_2} \quad \text{RDA\_BETA} \\
 \frac{\Gamma \vdash t : A_1 \rightarrow A_2 \quad x \notin \text{FV}(t)}{\Gamma \vdash \lambda x : A_1. t x \rightsquigarrow t : A_1 \rightarrow A_2} \quad \text{RDA\_ETA} \\
 \frac{\Gamma, x : A_1 \vdash t \rightsquigarrow t' : A_2}{\Gamma \vdash \lambda x : A_1. t \rightsquigarrow \lambda x : A_1. t' : A_1 \rightarrow A_2} \quad \text{RDA\_LAM} \\
 \frac{\Gamma \vdash t_1 \rightsquigarrow t'_1 : A_1 \rightarrow A_2 \quad \Gamma \vdash t_2 : A_1}{\Gamma \vdash t_1 t_2 \rightsquigarrow t'_1 t_2 : A_2} \quad \text{RDA\_APP1} \\
 \frac{\Gamma \vdash t_1 : A_1 \rightarrow A_2 \quad \Gamma \vdash t_2 \rightsquigarrow t'_2 : A_1}{\Gamma \vdash t_1 t_2 \rightsquigarrow t_1 t'_2 : A_2} \quad \text{RDA\_APP2} \\
 \frac{\Gamma \vdash t \rightsquigarrow t' : A_1 \times A_2}{\Gamma \vdash \text{fst } t \rightsquigarrow \text{fst } t' : A_1} \quad \text{RDA\_FST} \\
 \frac{\Gamma \vdash t \rightsquigarrow t' : A_1 \times A_2}{\Gamma \vdash \text{snd } t \rightsquigarrow \text{snd } t' : A_2} \quad \text{RDA\_SND} \\
 \frac{\Gamma \vdash t : A_1 \times A_2}{\Gamma \vdash (\text{fst } t, \text{snd } t) \rightsquigarrow t : A_1 \times A_2} \quad \text{RDA\_ETAP} \\
 \frac{\Gamma \vdash t_1 \rightsquigarrow t'_1 : A_1 \quad \Gamma \vdash t_2 : A_2}{\Gamma \vdash (t_1, t_2) \rightsquigarrow (t'_1, t_2) : A_1 \times A_2} \quad \text{RDA\_PAIR1} \\
 \frac{\Gamma \vdash t_1 : A_1 \quad \Gamma \vdash t_2 \rightsquigarrow t'_2 : A_2}{\Gamma \vdash (t_1, t_2) \rightsquigarrow (t_1, t'_2) : A_1 \times A_2} \quad \text{RDA\_PAIR2}
 \end{array}$$

 $\boxed{\Gamma \vdash t_1 \Rightarrow t_2 : A}$ 

Cast insertion from Siek16

$$\begin{array}{c}
 \frac{x : A \in \Gamma}{\Gamma \vdash x \Rightarrow x : A} \quad \text{CL\_VAR} \\
 \frac{}{\Gamma \vdash 0 \Rightarrow 0 : A} \quad \text{CL\_ZERO} \\
 \frac{}{\Gamma \vdash \text{triv} \Rightarrow \text{triv} : \text{Unit}} \quad \text{CL\_TRIV}
 \end{array}$$

$$\begin{array}{c}
\frac{\Gamma \vdash t_1 \Rightarrow t_2 : \text{Nat}}{\Gamma \vdash \text{succ } t_1 \Rightarrow \text{succ } t_2 : \text{Nat}} \quad \text{CI\_SUCC} \\
\\
\frac{\Gamma \vdash t_1 \Rightarrow t_3 : A_1 \quad \Gamma \vdash t_2 \Rightarrow t_4 : A_2}{\Gamma \vdash (t_1, t_2) \Rightarrow (t_3, t_4) : A_1 \times A_2} \quad \text{CI\_PAIR} \\
\\
\frac{\Gamma \vdash t_1 \Rightarrow t_2 : A_1 \times B \quad A_1 \sim A_2 \quad A_1 \neq A_2}{\Gamma \vdash \text{fst } t_1 \Rightarrow \text{fst } \langle A_2 \times B \rangle t_2 : A_2} \quad \text{CI\_FST1} \\
\\
\frac{\Gamma \vdash t_1 \Rightarrow t_2 : A \times B}{\Gamma \vdash \text{fst } t_1 \Rightarrow \text{fst } t_2 : A} \quad \text{CI\_FST2} \\
\\
\frac{\Gamma \vdash t_1 \Rightarrow t_2 : A \times B_1 \quad B_1 \sim B_2 \quad B_1 \neq B_2}{\Gamma \vdash \text{snd } t_1 \Rightarrow \text{snd } \langle A \times B_2 \rangle t_2 : B_2} \quad \text{CI\_SND1} \\
\\
\frac{\Gamma \vdash t_1 \Rightarrow t_2 : A \times B}{\Gamma \vdash \text{snd } t_1 \Rightarrow \text{snd } t_2 : B} \quad \text{CI\_SND2} \\
\\
\frac{\Gamma, x : A_1 \vdash t_1 \Rightarrow t_2 : A_2}{\Gamma \vdash \lambda x : A_1. t_1 \Rightarrow \lambda x : A_1. t_2 : A_1 \rightarrow A_2} \quad \text{CI\_LAM} \\
\\
\frac{\Gamma \vdash t_1 \Rightarrow t_3 : ? \quad \Gamma \vdash t_2 \Rightarrow t_4 : A}{\Gamma \vdash t_1 t_2 \Rightarrow (\langle A \rightarrow ? \rangle t_3) t_4 : ?} \quad \text{CI\_APP1} \\
\\
\frac{\Gamma \vdash t_1 \Rightarrow t_3 : A_1 \rightarrow B \quad \Gamma \vdash t_2 \Rightarrow t_4 : A_2 \quad A_1 \sim A_2 \quad A_1 \neq A_2}{\Gamma \vdash t_1 t_2 \Rightarrow t_1 \langle A_1 \rangle t : B} \quad \text{CI\_APP2} \\
\\
\frac{\Gamma \vdash t_1 \Rightarrow t_3 : A_1 \rightarrow A_2 \quad \Gamma \vdash t_2 \Rightarrow t_4 : A_1}{\Gamma \vdash t_1 t_2 \Rightarrow t_3 t_4 : A_2} \quad \text{CI\_APP3} \\
\\
\frac{\Gamma \vdash t_1 \Rightarrow t_2 : ?}{\Gamma \vdash \text{succ } t_1 \Rightarrow \langle ? \rangle \text{succ } \langle \text{Nat} \rangle t_2 : ?} \quad \text{CI\_SUCCU} \\
\\
\frac{\Gamma \vdash t_1 \Rightarrow t_2 : ?}{\Gamma \vdash \text{fst } t_1 \Rightarrow \text{fst } \langle ? \times ? \rangle t_2 : ?} \quad \text{CI\_FSTU} \\
\\
\frac{\Gamma \vdash t_1 \Rightarrow t_2 : ?}{\Gamma \vdash \text{snd } t_1 \Rightarrow \text{snd } \langle ? \times ? \rangle t_2 : ?} \quad \text{CI\_SNDU}
\end{array}$$

Definition rules: 105 good 0 bad

Definition rule clauses: 200 good 0 bad