## Gradual Typing from a Categorical Perspective

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## References

## A The Complete Spec of Grady

```
termvar, x, z
index, k
t
      ::=
                                  _{\rm term}
                                     variable
               triv
                                     unit
                                     injection of the retract
               squash
               split
                                     surjection of the retract
                                     generalize to the untyped universe
               box
               \mathsf{unbox}_T
                                     specialize the untyped universe to a specific type
               \lambda x : A.t
                                     \lambda-abstraction
                                     function application
               t_1 t_2
               (t_1, t_2)
                                     pair constructor
               \mathsf{fst}\ t
                                     first projection
               \mathsf{snd}\; t
                                     second projection
                                     successor function
               \mathsf{succ}\ t
               0
                                     zero
                            S
               (t)
T
                                  terminating types
               1
                                     unit type
               \mathbb{N}
                                     natural number type
              T_1 \rightarrow T_2
                                     function type
               T_1 \times T_2
                                     cartesian product type
                                  type
                                     terminating type
                                     untyped universe
                                     function type
```

$$\begin{array}{c} \Gamma \vdash t : ? \rightarrow ? \\ \hline \Gamma \vdash \text{split} \left( \text{squash } t \right) \leadsto t : ? \rightarrow ? \\ \hline \\ \Gamma \vdash \text{split} \left( \text{squash } t \right) \leadsto t : ? \rightarrow ? \\ \hline \\ \frac{\Gamma \vdash t : A_1 \rightarrow A_2 \quad x \not\in \text{FV}(t)}{\Gamma \vdash \lambda x : A_1 \cdot t \, x \leadsto t : A_1 \rightarrow A_2} \quad \text{RD\_ETA} \\ \hline \\ \frac{\Gamma, x : A_1 \vdash t_2 : A_2 \quad \Gamma \vdash t_1 : A_1}{\Gamma \vdash (\lambda x : A_1 \cdot t_2) \, t_1 \leadsto [t_1/x] t_2 : A_2} \quad \text{RD\_BETA} \\ \hline \\ \frac{\Gamma \vdash t_1 : A_1 \quad \Gamma \vdash t_2 : A_2}{\Gamma \vdash \text{fst} \left( t_1, t_2 \right) \leadsto t_1 : A_1} \quad \text{RD\_PROJ1} \\ \hline \\ \frac{\Gamma \vdash t_1 : A_1 \quad \Gamma \vdash t_2 : A_2}{\Gamma \vdash \text{snd} \left( t_1, t_2 \right) \leadsto t_2 : A_2} \quad \text{RD\_PROJ2} \\ \hline \\ \frac{\Gamma \vdash t : A_1 \times A_2}{\Gamma \vdash \text{snd} \left( t_1, t_2 \right) \leadsto t_2 : A_2} \quad \text{RD\_ETAP} \\ \hline \\ \frac{\Gamma, x : A_1 \vdash t \leadsto t' : A_1 \times A_2}{\Gamma \vdash (\text{fst} \, t, \text{snd} \, t) \leadsto t : A_1 \times A_2} \quad \text{RD\_ETAP} \\ \hline \\ \frac{\Gamma, x : A_1 \vdash t \leadsto t' : A_1}{\Gamma \vdash t_1 t_2 \leadsto t'_1 t_2 : A_2} \quad \text{RD\_APP1} \\ \hline \\ \frac{\Gamma \vdash t_1 \leadsto t'_1 : A_1 \rightarrow A_2 \quad \Gamma \vdash t_2 \leadsto t'_2 : A_1}{\Gamma \vdash t_1 t_2 \leadsto t_1 t'_2 : A_2} \quad \text{RD\_APP2} \\ \hline \\ \frac{\Gamma \vdash t \leadsto t' : A_1 \times A_2}{\Gamma \vdash \text{fst} \, t \leadsto \text{fst} \, t' : A_1} \quad \text{RD\_APP2} \\ \hline \\ \frac{\Gamma \vdash t \leadsto t' : A_1 \times A_2}{\Gamma \vdash \text{snd} \, t \leadsto \text{snd} \, t' : A_2} \quad \text{RD\_SND} \\ \hline \\ \frac{\Gamma \vdash t \leadsto t'_1 : A_1 \quad \Gamma \vdash t_2 : A_2}{\Gamma \vdash \text{snd} \, t \leadsto \text{snd} \, t' : A_2} \quad \text{RD\_PAIR1} \\ \hline \\ \frac{\Gamma \vdash t_1 \leadsto t'_1 : A_1 \quad \Gamma \vdash t_2 \leadsto t'_2 : A_2}{\Gamma \vdash (t_1, t_2) \leadsto (t'_1, t_2) : A_1 \times A_2} \quad \text{RD\_PAIR2} \\ \hline \\ \frac{\Gamma \vdash t_1 : A_1 \quad \Gamma \vdash t_2 \leadsto t'_2 : A_2}{\Gamma \vdash (t_1, t_2) \leadsto (t'_1, t'_2) : A_1 \times A_2} \quad \text{RD\_PAIR2} \\ \hline \end{array}$$