

Lecture 19 – Closure Properties of Context-Free Languages

COSE215: Theory of Computation

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- A **context-free language (CFL)** is defined in three different ways:
 - A **context free grammar (CFG)**
 - A **pushdown automaton (PDA)** with **final states**
 - A **pushdown automaton (PDA)** with **empty stacks**
- We have learned that the class of **regular languages** is **closed** under various operations. (**Closure Properties**)
- For which operations is the class of **CFLs** closed?

1. Closure Properties of Context-Free Languages

- Union

- Concatenation

- Kleene Star

- Homomorphism

- Reversal

2. Non-Closure Properties of Context-Free Languages

- Intersection

- Complement and Difference

3. Closure Properties of CFLs with Regular Languages

- Intersection with Regular Languages

- Difference with Regular Languages

Definition (Closure Properties)

The class of CFLs is **closed** under an n -ary operator op if and only if $op(L_1, \dots, L_n)$ is context-free for any CFLs L_1, \dots, L_n . We say that such properties are **closure properties** of CFLs.

The class of CFLs is closed under the following operations:

- Union
- Concatenation
- Kleene Star
- Homomorphism
- Reverse

Theorem (Closure under Union)

If L_1 and L_2 are context-free languages, then so is $L_1 \cup L_2$.

Proof) For given two CFLs L_1 and L_2 , we can always construct two CFGs:

$$G_1 = (V_1, \Sigma, S_1, R_1)$$

$$G_2 = (V_2, \Sigma, S_2, R_2)$$

such that $L_1 = L(G_1)$ and $L_2 = L(G_2)$.

Note that the variables of G_1 and G_2 should be disjoint. (i.e., $V_1 \cap V_2 = \emptyset$)

Then, $L_1 \cup L_2$ is accepted by the CFG $G = (V, \Sigma, S, R)$ where:

- $V = V_1 \cup V_2 \cup \{S\}$
- S is a new start variable (i.e., $S \notin V_1 \cup V_2$)
- $R = R_1 \cup R_2 \cup \{S \rightarrow S_1, S \rightarrow S_2\}$



For example, consider the following two CFLs:

$$L_1 = \{ab^n \mid n \geq 0\} \quad L_2 = \{ac^n \mid n \geq 0\}$$

Then, L_1 is accepted by:

$$S_1 \rightarrow aX \quad X \rightarrow bX \mid \epsilon$$

and L_2 is accepted by:

$$S_2 \rightarrow aX \quad X \rightarrow cX \mid \epsilon$$

But, the same variable X is used in both grammars.

So, we need to rename it to different variables, such as B and C .

For example, consider the following two CFLs:

$$L_1 = \{ab^n \mid n \geq 0\} \quad L_2 = \{ac^n \mid n \geq 0\}$$

Then, L_1 is accepted by:

$$S_1 \rightarrow aB \quad B \rightarrow bB \mid \epsilon$$

and L_2 is accepted by:

$$S_2 \rightarrow aC \quad C \rightarrow cC \mid \epsilon$$

Then, $L_1 \cup L_2$ is accepted by the following CFG:

$$\begin{aligned} S &\rightarrow S_1 \mid S_2 \\ S_1 &\rightarrow aB & B &\rightarrow bB \mid \epsilon \\ S_2 &\rightarrow aC & C &\rightarrow cC \mid \epsilon \end{aligned}$$

Theorem (Closure under Concatenation)

If L_1 and L_2 are context-free languages, then so is $L_1 \cdot L_2$.

Proof) For given two CFLs L_1 and L_2 , we can always construct two CFGs:

$$G_1 = (V_1, \Sigma, S_1, R_1)$$

$$G_2 = (V_2, \Sigma, S_2, R_2)$$

such that $L_1 = L(G_1)$ and $L_2 = L(G_2)$.

Note that the variables of G_1 and G_2 should be disjoint. (i.e., $V_1 \cap V_2 = \emptyset$)

Then, $L_1 \cdot L_2$ is accepted by the CFG $G = (V, \Sigma, S, R)$ where:

- $V = V_1 \cup V_2 \cup \{S\}$
- S is a new start variable (i.e., $S \notin V_1 \cup V_2$)
- $R = R_1 \cup R_2 \cup \{S \rightarrow S_1 S_2\}$



For example, consider the following two CFLs:

$$L_1 = \{ab^n \mid n \geq 0\} \quad L_2 = \{ac^n \mid n \geq 0\}$$

Then, L_1 is accepted by:

$$S_1 \rightarrow aB \quad B \rightarrow bB \mid \epsilon$$

and L_2 is accepted by:

$$S_2 \rightarrow aC \quad C \rightarrow cC \mid \epsilon$$

Then, $L_1 \cdot L_2$ is accepted by the following CFG:

$$\begin{aligned} S &\rightarrow S_1 S_2 \\ S_1 &\rightarrow aB & B &\rightarrow bB \mid \epsilon \\ S_2 &\rightarrow aC & C &\rightarrow cC \mid \epsilon \end{aligned}$$

Theorem (Closure under Kleene Star)

If L is a context-free language, then so is L^ .*

Proof) For a given CFL L , we can always construct a CFG:

$$G = (V, \Sigma, S, R)$$

such that $L = L(G)$.

Then, L^* is accepted by the CFG $G' = (V', \Sigma, S', R')$ where:

- $V' = V \cup \{S'\}$
- S' is a new start variable (i.e., $S' \notin V$)
- $R' = R \cup \{S' \rightarrow \epsilon, S' \rightarrow SS'\}$



Closure under Kleene Star – Example

For example, consider the following CFL:

$$L = \{a^n b^n \mid n \geq 0\}$$

Then, L is accepted by:

$$S \rightarrow \epsilon \mid aSb$$

Then, L^* is accepted by the following CFG:

$$\begin{aligned} S' &\rightarrow \epsilon \mid SS' \\ S &\rightarrow \epsilon \mid aSb \end{aligned}$$

Theorem (Homomorphism)

Suppose Σ and Γ are two finite sets of symbols. Then, a function $h : \Sigma \rightarrow \Gamma^*$ is called a **homomorphism**. For a given word $w = a_1 a_2 \cdots a_n$,

$$h(w) = h(a_1)h(a_2) \cdots h(a_n)$$

For a language L , $h(L) = \{h(w) \mid w \in L\}$.

For example, $h : \{0, 1\} \rightarrow \{a, b\}^*$ be a homomorphism such that:

$$h(0) = ab \quad h(1) = a$$

Then,

$$h(10) = aab \quad h(010) = abaab \quad h(1100) = aaabab$$

Theorem (Closure under Homomorphism)

If h is a homomorphism and L is a context-free language, then so is $h(L)$.

Proof) For a given CFL L , we can always construct a CFG:

$$G = (V, \Sigma, S, R)$$

such that $L = L(G)$.

Then, for a given homomorphism $h : \Sigma \rightarrow \Gamma^*$, $h(L)$ is accepted by the CFG $G' = (V', \Gamma, S, R')$ where:

- $V' = V \cup \{X_a \mid a \in \Sigma\}$
- $R' = \{Y \rightarrow Y'_1 \cdots Y'_n \mid Y \rightarrow Y_1 \cdots Y_n \in R\} \cup \{X_a \rightarrow h(a) \mid a \in \Sigma\}$

$$\text{where } \forall 1 \leq i \leq n. Y'_i = \begin{cases} Y_i & \text{if } Y_i \in V \\ X_a & \text{if } Y_i = a \in \Sigma \end{cases}$$



For example, consider the following CFL:

$$L = \{ww^R \mid w \in \{0, 1\}^*\}$$

Then, L is accepted by:

$$S \rightarrow \epsilon \mid 0S0 \mid 1S1$$

If a homomorphism $h : \{0, 1\} \rightarrow \{a, b\}^*$ is defined as follows:

$$h(0) = ab \quad h(1) = a$$

Then, $h(L)$ is accepted by the following CFG:

$$S \rightarrow \epsilon \mid X_0SX_0 \mid X_1SX_1$$

$$X_0 \rightarrow ab$$

$$X_1 \rightarrow a$$

Theorem (Closure under Reverse)

If L is a context-free language, then so is L^R .

Proof) For a given CFL L , we can always construct a CFG:

$$G = (V, \Sigma, S, R)$$

such that $L = L(G)$.

Then, L^R is accepted by the CFG $G' = (V, \Sigma, S, R')$ where:

- $R' = \{X \rightarrow \alpha^R \mid X \rightarrow \alpha \in R\}$



Closure under Reverse – Example

For example, consider the following CFL:

$$L = \{(ab)^n c^n d^m \mid n, m \geq 0\}$$

Then, L is accepted by:

$$\begin{aligned} S &\rightarrow X \mid Sd \\ X &\rightarrow \epsilon \mid abXc \end{aligned}$$

Then, L^R is accepted by the following CFG:

$$\begin{aligned} S &\rightarrow X \mid dS \\ X &\rightarrow \epsilon \mid cXba \end{aligned}$$

Definition (Closure Properties)

The class of CFLs is **closed** under an n -ary operator op if and only if $op(L_1, \dots, L_n)$ is context-free for any CFLs L_1, \dots, L_n . We say that such properties are **closure properties** of CFLs.

The class of CFLs is **NOT** closed under the following operations:

- Intersection
- Complement
- Difference

We will prove it by using the fact that the following language is not a CFL:

$$L = \{a^n b^n c^n \mid n \geq 0\}$$

We will learn how to prove that L is not a CFL in the next lecture (Pumping Lemma for CFLs).

Theorem (Non-Closure under Intersection)

*The class of CFLs is **NOT** closed under intersection.*

Proof) Consider the following two languages:

$$L_1 = \{a^n b^n c^m \mid n, m \geq 0\} \quad L_2 = \{a^m b^n c^n \mid n, m \geq 0\}$$

Then, L_1 is accepted by:

$$S_1 \rightarrow X \mid S_1 c \quad X \rightarrow \epsilon \mid aXb$$

and L_2 is accepted by:

$$S_2 \rightarrow Y \mid aS_2 \quad Y \rightarrow \epsilon \mid bYc$$

Thus, they are both CFLs. However, their intersection is not a CFL:

$$L_1 \cap L_2 = \{a^n b^n c^n \mid n \geq 0\}$$



Definition (Non-Closure under Complement)

The class of CFLs is **NOT** closed under complement.

Proof) Assume that the class of CFLs is closed under complement. Then, for any two CFLs L_1 and L_2 , $L_1 \cap L_2$ is also a CFL:

$$L_1 \cap L_2 = \overline{\overline{L_1} \cup \overline{L_2}}$$

However, we have already proved that the class of CFLs is not closed under intersection. Thus, the class of CFLs is not closed under complement. \square

Definition (Non-Closure under Difference)

The class of CFLs is **NOT** closed under difference.

Proof) Similarly, we can prove it using the following fact:

$$L_1 \cap L_2 = L_1 \setminus (L_1 \setminus L_2)$$

 \square

Definition (Closure Properties)

The class of CFLs is **closed** under an n -ary operator op if and only if $op(L_1, \dots, L_n)$ is context-free for any CFLs L_1, \dots, L_n . We say that such properties are **closure properties** of CFLs.

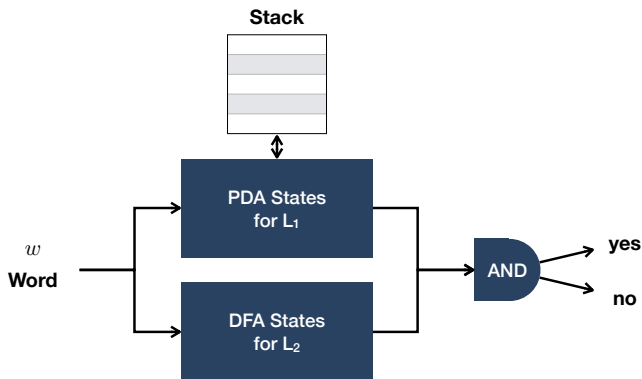
The class of CFLs is closed under the following operations with RLs:

- Intersection
- Difference

Theorem (Closure under Intersection with RLs)

If L_1 is a CFL and L_2 is a RL, then $L_1 \cap L_2$ is a CFL.

There exists a PDA P that accepts L_1 by final states and a DFA D that accepts L_2 . We will construct a PDA P' that accepts $L_1 \cap L_2$ as follows:



Theorem (Closure under Intersection with RLs)

If L_1 is a CFL and L_2 is a RL, then $L_1 \cap L_2$ is a CFL.

Proof) Consider a PDA $P = (Q_P, \Sigma, \Gamma, \delta_P, q_P, Z, F_P)$ and a DFA $D = (Q_D, \Sigma, \delta_D, q_D, F_D)$ such that:

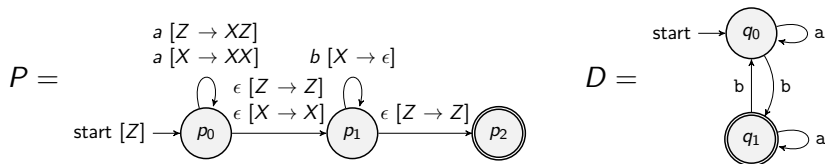
$$L_F(P) = L_1 \quad L(D) = L_2$$

Then, $L_1 \cap L_2$ is accepted by the PDA $P' = (Q, \Sigma, \Gamma, \delta, q_0, Z, F)$ by final states, where:

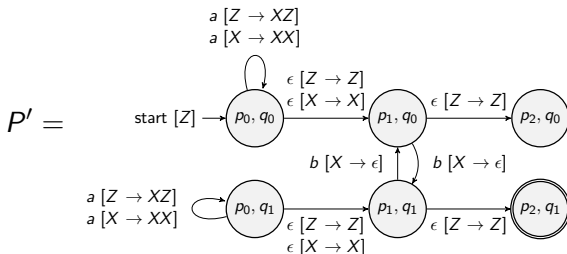
- $Q = Q_P \times Q_D$
- $\delta((p, q), \epsilon, X) = \{((p', q), \alpha) \mid (p', \alpha) \in \delta_P(p, \epsilon, X)\}$
- $\delta((p, q), a, X) = \{((p', q'), \alpha) \mid (p', \alpha) \in \delta_P(p, a, X) \wedge q' = \delta_D(q, a)\}$
- $q_0 = (q_P, q_D)$
- $F = F_P \times F_D$



For example, consider the following PDA P and DFA D :



Then, a PDA P' that accepts $L_F(P) \cap L(D)$ by the final states can be constructed as follows:



Theorem (Closure under Difference with RLs)

If L_1 is a CFL and L_2 is a RL, then $L_1 \setminus L_2$ is a CFL.

Proof) We know the following fact:

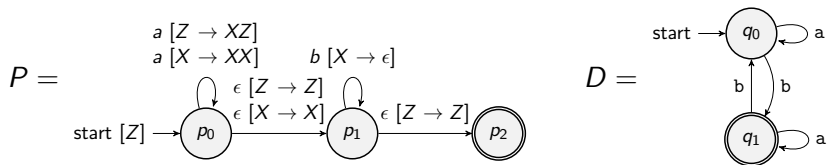
$$L_1 \setminus L_2 = L_1 \cap \overline{L_2}$$

Since the class of RLs is closed under complement, $\overline{L_2}$ is a RL. In addition, we know that the class of CFLs is closed under intersection with RLs.

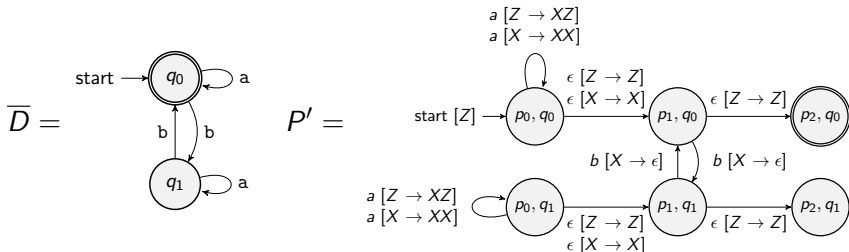
Thus, $L_1 \setminus L_2$ is a CFL. □

Closure under Difference with RLs – Example

For example, consider the following PDA P and DFA D :



Then, a DFA \overline{D} that accepts $\overline{L(D)}$ and a PDA P' that accepts $L_F(P) \setminus L(D) = L_F(P) \cap \overline{L(D)}$ can be constructed as follows:



1. Closure Properties of Context-Free Languages

Union

Concatenation

Kleene Star

Homomorphism

Reversal

2. Non-Closure Properties of Context-Free Languages

Intersection

Complement and Difference

3. Closure Properties of CFLs with Regular Languages

Intersection with Regular Languages

Difference with Regular Languages

- The Pumping Lemma for Context-Free Languages

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