Lecture 16 – Equivalence of Pushdown Automata and Context-Free Grammars

COSE215: Theory of Computation

Jihyeok Park

PLRG

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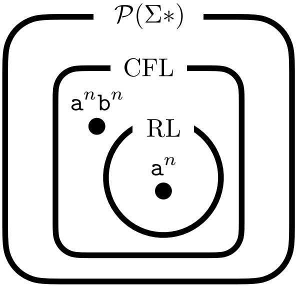
Recall

A **context-free grammar** is a 4-tuple:

$$G = (V, \Sigma, S, P)$$

A **pushdown automaton (PDA)** is a finite automaton with a **stack**.

- Acceptance by **final states**
- Acceptance by **empty stacks**



$\mathrm{PDA}_{\mathrm{FS}}$
(by final states)

$\parallel$ - - - - - ?

$\mathrm{PDA}_{\mathrm{ES}}$
(by empty stacks)

$\parallel$ - - - - - ?

CFG

Contents

1. Equivalence of PDA by Final States and Empty Stacks
    - $\mathsf{PDA}_{\mathsf{FS}}$ to $\mathsf{PDA}_{\mathsf{ES}}$
    - $\mathsf{PDA}_{\mathsf{ES}}$ to $\mathsf{PDA}_{\mathsf{FS}}$
2. Equivalence of PDA and CFGs
    - CFGs to $\mathsf{PDA}_{\mathsf{ES}}$
    - $\mathsf{PDA}_{\mathsf{ES}}$ to CFGs

$$\underset{\text{(by final states)}}{\mathrm{PDA}_{\mathrm{FS}}} \rightleftarrows \underset{\text{(by empty stacks)}}{\mathrm{PDA}_{\mathrm{ES}}} \rightleftarrows \mathrm{CFG}$$

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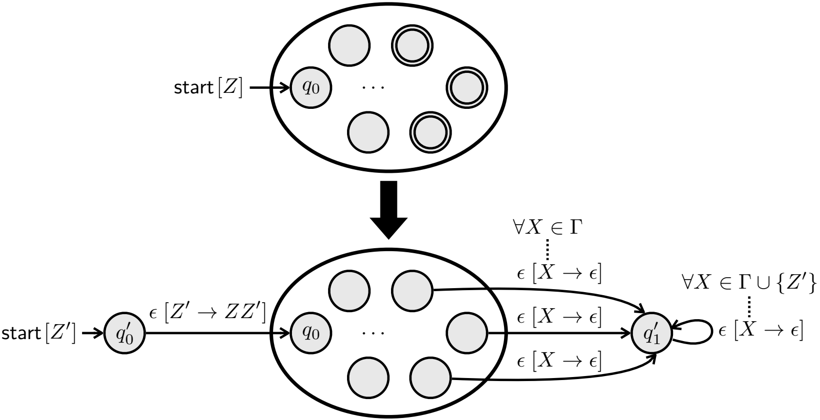
$$\underset{\text{(by final states)}}{\mathrm{PDA_{FS}}} \rightleftarrows \underset{\text{(by empty stacks)}}{\mathrm{PDA_{ES}}} \rightleftarrows \mathrm{CFG}$$

PDA_{FS} to PDA_{ES}

Theorem (PDA_{FS} to PDA_{ES})

For a given PDA $P = (Q, \Sigma, \Gamma, \delta, q_0, Z, F)$, $\exists$ PDA P'. $L_F(P) = L_E(P')$.



PDA_{FS} to PDA_{ES}

Theorem (PDA_{FS} to PDA_{ES})

For a given PDA $P = (Q, \Sigma, \Gamma, \delta, q_0, Z, F)$, $\exists$ PDA P'. $L_F(P) = L_E(P')$.

Define a PDA

$$P' = (Q \cup \{q'_0, q'_1\}, \Sigma, \Gamma \cup \{Z'\}, \delta', q'_0, Z', \varnothing)$$

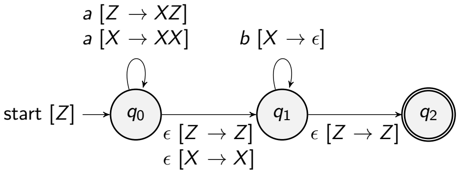
where

$$\begin{array}{lcl}
\delta'(q'_0, \epsilon, Z') & = & \{(q_0, ZZ')\} \\[2ex]
\delta'(q \in Q, a \in \Sigma, X \in \Gamma) & = & \delta(q, a, X) \\[2ex]
\delta'(q \in Q, \epsilon, X \in \Gamma) & = & \begin{cases} \delta(q, \epsilon, X) \cup \{(q'_1, \epsilon)\} & \text{if } q \in F \\ \delta(q, a, X) & \text{otherwise} \end{cases} \\[2ex]
\delta'(q'_1, \epsilon, X \in \Gamma \cup \{Z'\}) & = & \{(q'_1, \epsilon)\}
\end{array}$$

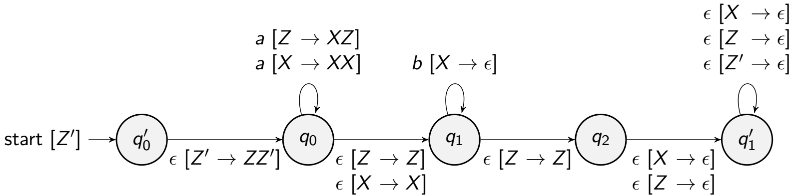
PDA_{FS} to PDA_{ES} – Example

$$L_F(P) = L_E(P') = \{\mathtt{a}^n\mathtt{b}^n \mid n \geq 0\}$$

$P =$

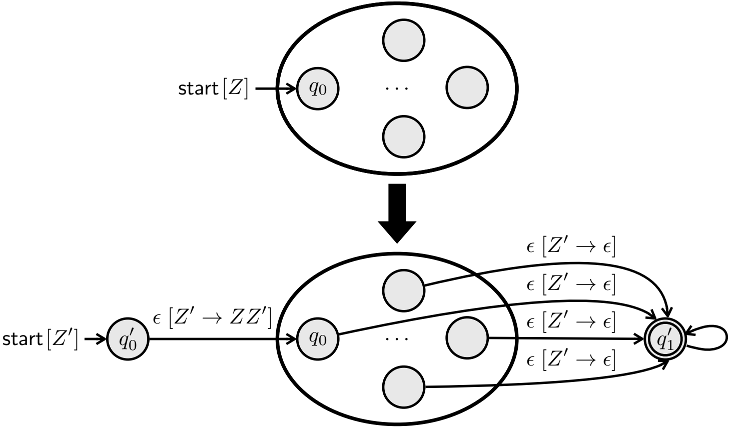


$P' =$



Theorem ($\mathrm{PDA_{ES}}$ to $\mathrm{PDA_{FS}}$)

For a given PDA $P = (Q, \Sigma, \Gamma, \delta, q_0, Z, F)$, $\exists$ *PDA* P'. $L_E(P) = L_F(P')$.



PDA$_{ES}$ to PDA$_{FS}$

Theorem (PDA$_{ES}$ to PDA$_{FS}$)

For a given PDA $P = (Q, \Sigma, \Gamma, \delta, q_0, Z, F)$, $\exists$ *PDA* P'. $L_E(P) = L_F(P')$.

Define a PDA

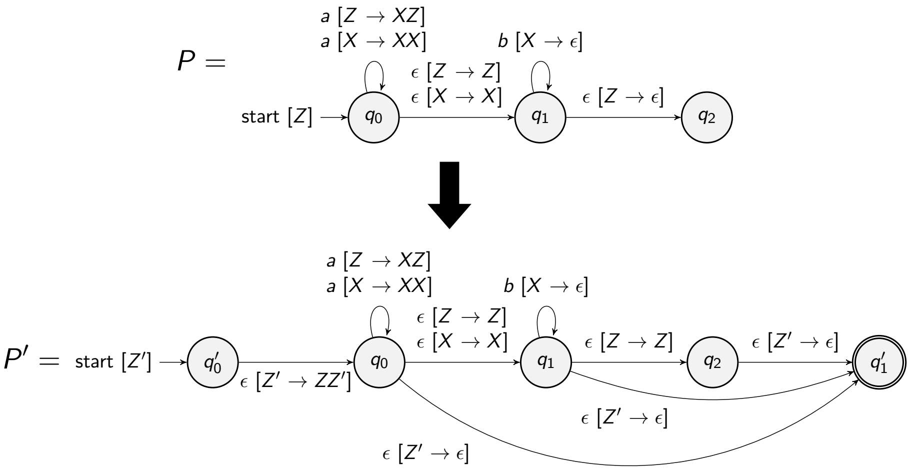
$$P' = (Q \cup \{q'_0, q'_1\}, \Sigma, \Gamma \cup \{Z'\}, \delta', q'_0, Z', \{q'_1\})$$

where

$$\begin{array}{rcl} \delta'(q'_0, \epsilon, Z') & = & \{(q_0, ZZ')\} \\ \delta'(q \in Q, a \in \Sigma, X \in \Gamma) & = & \delta(q, a, X) \\ \delta'(q \in Q, \epsilon, X \in \Gamma) & = & \delta(q, a, X) \\ \delta'(q \in Q, \epsilon, Z') & = & \{(q'_1, \epsilon)\} \end{array}$$

PDA_{ES} to PDA_{FS} – Example

$$L_E(P) = L_F(P') = \{\mathtt{a}^n\mathtt{b}^n \mid n \geq 0\}$$



Contents

$$\underset{\text{(by final states)}}{\mathrm{PDA}_{\mathrm{FS}}} \rightleftarrows \underset{\text{(by empty stacks)}}{\mathrm{PDA}_{\mathrm{ES}}} \rightleftarrows \mathrm{CFG}$$

CFGs to PDA_{ES}

Theorem (CFGs to PDA_{ES})

For a given CFG $G = (V, \Sigma, S, R)$, $\exists$ PDA P. $L(G) = L_E(P)$.

Define a PDA

$$P = (\{q\}, \Sigma, V \cup \Sigma, \delta, q, S, \varnothing)$$

where

$$\begin{array}{rcl} \delta(q, \epsilon, A \in V) & = & \{(q, \alpha) \mid A \to \alpha \in R\} \\ \delta(q, a \in \Sigma, a \in \Sigma) & = & \{(q, \epsilon)\} \end{array}$$

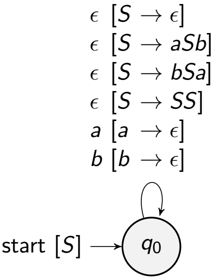
CFGs to PDA_{ES} – Example

$$
\begin{array}{lcl}
\delta(q, \epsilon, A \in V) & = & \{(q, \alpha) \mid A \to \alpha \in R\} \\
\delta(q, a \in \Sigma, a \in \Sigma) & = & \{(q, \epsilon)\}
\end{array}
$$

Consider the following CFG:

$$S \to \epsilon \mid \mathtt{a}S\mathtt{b} \mid \mathtt{b}S\mathtt{a} \mid SS$$

Then, the equivalent PDA (by empty stacks) is:



$$
\begin{array}{lcl}
(q, \mathtt{abab}, S) & \vdash & (q, \mathtt{abab}, \mathtt{a}S\mathtt{b}) \\
& \vdash & (q, \mathtt{bab}, S\mathtt{b}) \\
& \vdash & (q, \mathtt{bab}, \mathtt{b}S\mathtt{ab}) \\
& \vdash & (q, \mathtt{ab}, S\mathtt{ab}) \\
& \vdash & (q, \mathtt{ab}, \mathtt{ab}) \\
& \vdash & (q, \mathtt{b}, \mathtt{b}) \\
& \vdash & (q, \epsilon, \epsilon)
\end{array}
$$

Theorem (PDA$_{\text{ES}}$ to CFGs)

For a given PDA $P = (Q = \{q_0, \cdots, q_{n-1}\}, \Sigma, \Gamma, \delta, q_0, Z, F)$,
$\exists$ *CFG* G. $L_E(P) = L(G)$.

Consider the set of variables $V = \{S\} \cup \{ A^X_{i,j} \mid 0 \leq i, j < n \wedge X \in \Gamma\}$.
Then, define a CFG:

- For all $0 \leq j < n$,

$$S \rightarrow A^Z_{0,j}$$

- For all $q_i \in Q$, $a \in \Sigma \cup \{\epsilon\}$, and $X \in \Gamma$, consider any $(q_j, X_1 \cdots X_m) \in \delta(q_i, a, X)$ and $0 \leq k_1, \cdots, k_m < n$. Then,

$$A^X_{i,k_m} \rightarrow a\ A^{X_1}_{j,k_1}\ A^{X_2}_{k_1,k_2} \cdots A^{X_m}_{k_{m-1},k_m}$$

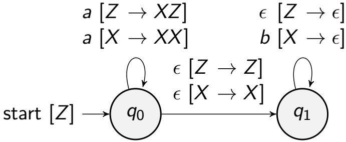
Note that each variable $A^X_{i,j}$ generates all words that cause the PDA to go from state q_i to state q_j by popping X:

$$A^X_{i,j} \Rightarrow^* w \qquad \text{if and only if} \qquad (q_j, w, X) \vdash^* (q_i, \epsilon, \epsilon)$$

PDA_{ES} to CFGs – Example

$$S \to A^{Z}_{0,j} \qquad\qquad A^{X}_{i,k_m} \to a\ A^{X_1}_{j,k_1}\ A^{X_2}_{k_1,k_2} \cdots A^{X_m}_{k_{m-1},k_m}$$

Consider the following PDA (by empty stacks):



Then, the equivalent CFG is:

$$\begin{array}{lcl}
S & \to & A^{Z}_{0,0} \mid A^{Z}_{0,1} \\
A^{Z}_{0,0} & \to & \mathrm{a}\ A^{X}_{0,0}\ A^{Z}_{0,0} \mid \mathrm{a}\ A^{X}_{0,1}\ A^{Z}_{1,0} \mid A^{Z}_{1,0} \\
A^{Z}_{0,1} & \to & \mathrm{a}\ A^{X}_{0,0}\ A^{Z}_{0,1} \mid \mathrm{a}\ A^{X}_{0,1}\ A^{Z}_{1,1} \mid A^{Z}_{1,1} \\
A^{X}_{0,0} & \to & \mathrm{a}\ A^{X}_{0,0}\ A^{X}_{0,0} \mid \mathrm{a}\ A^{X}_{0,1}\ A^{X}_{1,0} \mid A^{X}_{1,0} \\
A^{X}_{0,1} & \to & \mathrm{a}\ A^{X}_{0,0}\ A^{X}_{0,1} \mid \mathrm{a}\ A^{X}_{0,1}\ A^{X}_{1,1} \mid A^{X}_{1,1} \\
A^{Z}_{1,1} & \to & \epsilon \\
A^{X}_{1,1} & \to & b
\end{array}
\qquad
\begin{array}{lcl}
S & \Rightarrow & A^{Z}_{0,1} \\
 & \Rightarrow & \mathrm{a}\ A^{X}_{0,1}\ A^{Z}_{1,1} \\
 & \Rightarrow & \mathrm{aa}\ A^{X}_{0,1}\ A^{X}_{1,1}\ A^{Z}_{1,1} \\
 & \Rightarrow & \mathrm{aa}\ A^{X}_{1,1}\ A^{X}_{1,1}\ A^{Z}_{1,1} \\
 & \Rightarrow & \mathrm{aab}\ A^{X}_{1,1}\ A^{Z}_{1,1} \\
 & \Rightarrow & \mathrm{aabb}\ A^{Z}_{1,1} \\
 & \Rightarrow & \mathrm{aabb}
\end{array}$$

Summary

1. Equivalence of PDA by Final States and Empty Stacks
   - $\mathsf{PDA}_{\mathsf{FS}}$ to $\mathsf{PDA}_{\mathsf{ES}}$
   - $\mathsf{PDA}_{\mathsf{ES}}$ to $\mathsf{PDA}_{\mathsf{FS}}$

2. Equivalence of PDA and CFGs
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Next Lecture

- Deterministic Pushdown Automata (DPDA)

Jihyeok Park
jihyeok_park@korea.ac.kr
https://plrg.korea.ac.kr