

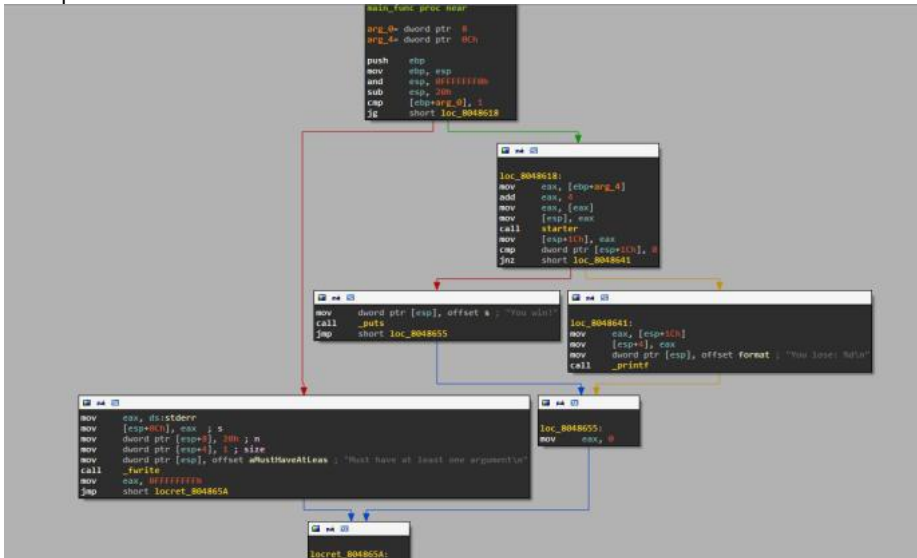
1. Start

We are given a 32-bit executable.

04:48 AM in30:Aerosol_Can-500>file aerosol_can
aerosol_can: ELF 32-bit LSB executable, Intel 80386, version 1 (SYSV), dynamically linked (uses shared libs), for GNU/Linux 2.6

2. Understand

Lets open this in IDA.



In IDA, we see the **main_func** function which checks for the number of arguments passed to the binary and prints "Must have at least one argument\n" string if its less then 2 otherwise it passes the second argument to the function call **starter(0x8053460)** which checks for the correctness of the 2nd argument we passed, which in this case, is our flag string. If we passed the correct flag then function **starter** returns the 0 otherwise any other negative value.

lets go inside starter function...

```

var_1A4= dword ptr -1A4h
var_18C= dword ptr -18Ch
var_16C= dword ptr -16Ch
stack_base= dword ptr -10h
stack_top= dword ptr -0Ch
arg_0= dword ptr 4

push    edi
push    esi
sub     esp, 1C4h
mov     [esp+1CCh+var_1AC.rlim_max], 0
mov     [esp+1CCh+var_1AC.rlim_cur], 0
lea     eax, [esp+20h]
mov     [esp+1CCh+rlimits], eax ; rlimits
mov     [esp+1CCh+resource], 3 ; RLIMIT_STACK resource
call    _getrlimit
mov     edi, [esp+1CCh+var_1AC.rlim_cur]
mov     [esp+1CCh+rlimits], edi ; len
mov     [esp+1CCh+offset], 0 ; offset
mov     [esp+1CCh+fd], 0FFFFFFFh ; fd
mov     [esp+1CCh+flags], 20022h ; flags
mov     [esp+1CCh+prot], 3 ; prot
mov     [esp+1CCh+resource], 0 ; addr
call    _mmap
mov     esi, eax
mov     eax, [esp+1CCh+arg_0]
mov     [esi+edi-34h], eax ; input string
mov     [esp+1CCh+stack_top], esi
lea     eax, [esi+edi]
mov     [esp+1CCh+stack_base], eax
lea     eax, [esi+edi-38h]
mov     [esp+1CCh+var_16C], 0
mov     [esp+1CCh+var_18C], eax
lea     eax, [esp+1CCh+var_1A4]
mov     [esp], eax ; DEVIL argument
call    DEVIL
sub     esp, 4
mov     [esp+1CCh+var_184.rlim_max], 0
mov     [esp+1CCh+var_184.rlim_cur], 0
lea     eax, [esp+1CCh+var_184]
mov     [esp+1CCh+rlimits], eax ; rlimits
mov     [esp+1CCh+resource], 3 ; resource
call    _getrlimit
mov     eax, [esp+1CCh+var_184.rlim_cur]
mov     [esp+1CCh+rlimits], eax ; len
mov     [esp+1CCh+resource], esi ; addr
call    _munmap
mov     eax, [esp+1CCh+var_184] ; check_flag
add     esp, 1C4h
pop     esi
pop     edi
ret
starter endp
  
```

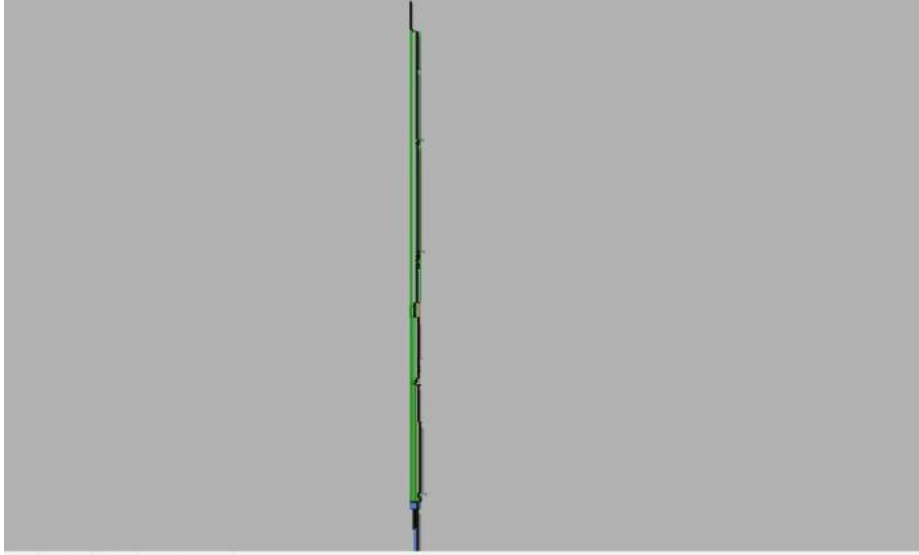
What this function does is :

- Calls **getrlimit** function to get the stack limit(size) for the program.
- Calls **mmap** function to map virtual stack of the same size as above.(this will later be used extensively)
- Call function **DEVIL(0x8048660)** to do some processing on our input string.

- Again Call *getrlimit* function to get the stack limit(size) of the program.
- Call *munmap* to unmap the previously mapped stack.

3. Obfuscation

DEVIL function is passed some address(*esp+0x11c-0x1A4*) of the original stack, but before that it puts an address(base address of new stack - *0x38*) at [*esp+0x11c-0x18C*] which is at an offset *0x18* of the address which is passed in **DEVIL** function. It also puts the address of the input string at (base address of new stack - *0x34*) which is just below the address we put at [*esp+0x11c-0x18C*] above. So lets see the **DEVIL** function...



Its a *thousands of lines of code* which is quite an *Obfuscation*. And It took me lot a of time to figure out what to do with this binary.

First, let me tell you what the *obfuscation* is and then we will rip this binary apart :P

Actually, the obfuscation is a play between *Original Stack*, *New Stack* and *local variables* of **DEVIL** function on stack. **DEVIL** function allocates *0x1DC* size memory for its local variables on the stack in it prologue, then it takes the argument we passed, takes random values from different offset of it and puts it in random address on new stack and local variables to obfuscate the code. I could have done dead code analysis on the binary but it does use those dead code(not so dead actually :P) to pass the arguments to the functions and for the control flow checks.

What I did to figure out a way out of this mind f**king Obfuscation is:

1. follow the [*ebx+18h*] value(which contains some address of the new stack and is used to get reference to our input string) in the code
2. 2. follow the function calls and what their arguments and what they are returning (i tried to understand the functions but they are obfuscated too :())
3. 3. follow the control flow checks, what check leads to which path.

See their are many different solution possible for this , and each solution leads to return 0, and , you can follow your own path to make it return 0. Following is just my way of doing that.

4. Functions

Lets start getting the flag. I will reduce my explanation to just checks only ,otherwise it will take soo much time and space...

I will explain the different function call used during the checks in **DEVIL** function.

- | | | |
|--|---|--|
| 1. str_to_hex (<i>0x8051bd0</i>) | takes string address as argument | convert/decode string like "abcd" into 0xabcd , "efgh" into 0xef00 |
| (using the following table for conversion) | | |
| 2. hex_to_1char (<i>0x80512e0</i>) | takes input a character address | convert/decode that character according to above table |
| 3. sum_natural_num (<i>0x8050030</i>) | takes an integer(n) as argument | returns the sum of first n natural numbers (1 + 2 + 3.... n) |
| 4. divide_ (<i>0x08053540</i>) | takes 4 integer as argument(a,b,c,d) | returns a/(b*c*d) |
| 5. sum_4hexchar (<i>0x0804DAB9</i>) | takes (<i>esp+0x11c-0x1A4</i>)(a) as argument | returns sum of 8bit values from [a] to [a+8] |

TABLE :

• 0x20 – 0x2f	->	0x0 – 0xf
• 0x3A – 0x46	->	0x3 – 0xf
• 0x30 – 0x39	->	0x0 – 0x9
• 0x47 – 0x57	->	0x0 – 0xf
• 0x58 – 0x66	->	0x0 – 0xf
• 0x66 onward	->	0x0

5. Flag Checks

1. Code first checks for the length of the string we passed

```

1
2.text:0804866A      mov     ebx, [esp+1ECh+arg_0]    ; new stack address
3.text:080486A4      mov     ebp, [ebx+18h]
4.text:080489D5      mov     eax, [ebp+4]            ; input string
4.text:080489DE      mov     [esp+1ECh+s], eax ; s

```

```

5.text:080489E1      call    _strlen
6.text:080489E6      mov     [ebp-8Ch], eax    ; string length
7.text:08048A96      cmp     eax, 26h

```

It must be equal to 0x26.

2. Then it checks for the first 5 characters of the string flag[0:5]

```

1
2 .text:08048A9F      mov     eax, [ebp+4]
3 .text:08048AA2      mov     ecx, 0FFFFFFFh
4 .text:08048AA7      cmp     byte ptr [eax], 'f'
5 .text:08048AAA      jnz     loc_80492FF
6 .text:08048AB0      cmp     byte ptr [eax+1], 'l'
7 .text:08048AB4      jnz     loc_80492FF
8 .text:08048ABA      cmp     byte ptr [eax+2], 'a'
9 .text:08048ABE      jnz     loc_80492FF
10 .text:08048AC4      cmp     byte ptr [eax+3], 'g'
11 .text:08048AC8      jnz     loc_80492FF
12 .text:08048AE0      cmp     byte ptr [eax+4], '{'
13 .text:08048AE4      jnz     loc_80492FF

```

first 5 chars must be **flag[0:5] = "flag{"**

3. Check for next 4 char

```

1 .text:08048B02      add     ecx, 5            ; string pointer added 5 for flag{
2 .text:08048B1A      mov     [ebx], ecx
3 .text:08048ECD      mov     [esp+1ECh+s], ebx    ; input_string[5:9]
4 .text:08048ED0      call    str_to_hex
5 .text:0804921F      mov     ecx, [esp+1ECh+var_174] ; str_to_hex returned value
6 .text:0804922A      movzx   eax, word ptr [edx-9Ch] ; constant 0x0E607
7 .text:08049235      xor     ecx, eax
8 .text:080492DC      cmp     cx, 0FEEDh

```

I choosed **flag[5:9] = "1_ea"**

4.Check for 6th character

```

1 .text:08049AF1      mov     [esp+1ECh+s], ebx    ; input_string[6]
2 .text:08049AF4      call    hex_to_1char
3 .text:0804A270      mov     [esp+1ECh+s], ebx    ; return value from above hex_to_1char
4 .text:0804A273      call    sum_natural_num
5 .text:0804A278      sub     esp, 4
6 .text:0804A40C      mov     al, [eax+ecx+1]      ; ecx is address if input string , eax is the return value of
7 .text:0804A66B      cmp     byte ptr [esp+1ECh+var_E8], '{'

```

flag[0x25] = "}"

5. Check for 9th char

```

1 .text:0804A69D      movsx   eax, byte ptr [eax+9] ; flag[9]
2 .text:0804A6A1      mov     cl, al
3 .text:0804A6A3      and     cl, 0Fh
4 .text:0804A6A6      cmp     cl, 6              ; check for last 4 bits of flag[9]
5 .text:0804A6AD      jnz     loc_804A756
6 .text:0804A6B3      mov     ecx, eax
7 .text:0804A6B5      and     ecx, 0FCh
8 .text:0804A6BB      cmp     ecx, 34h          ; check for first 4 bits of flag[9]

```

I choosed **flag[9] = "a"(0x36)**

6. Check for 10,11,12th chars

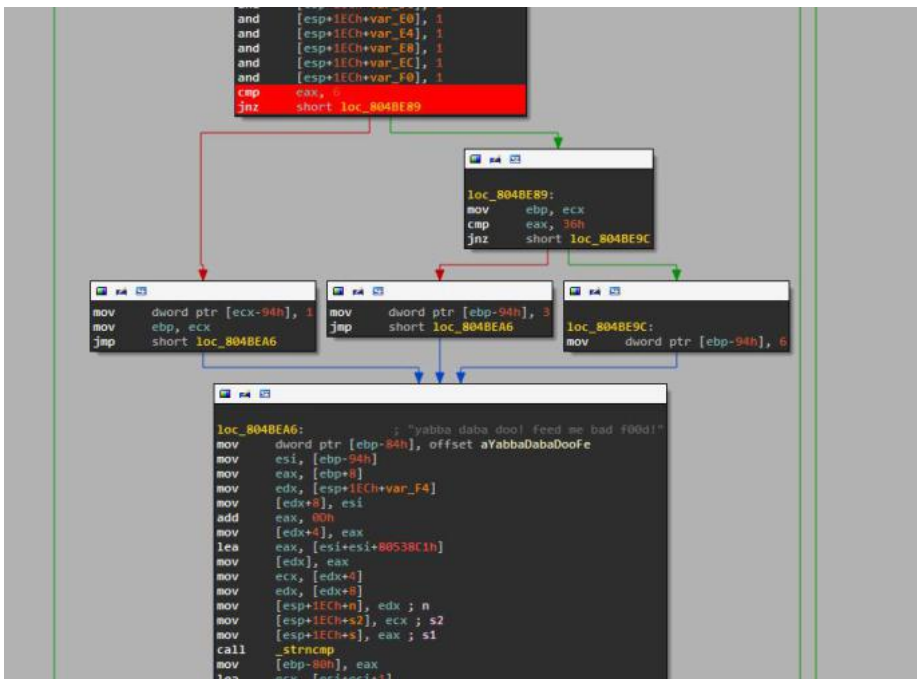
```

1 .text:0804AB6B      call    hex_to_1char        ; input_string[10] (b)
2 .text:0804B2DD      call    hex_to_1char        ; input_string[11] (c)
3 .text:0804BA56      call    hex_to_1char        ; input_string[12] (d)
4 .text:0804BDC5      call    divide_             ; 144 , b,c,d
5 .text:0804BE76      cmp     eax, 6

```

I choosed **flag[10:13] = "lm6"**

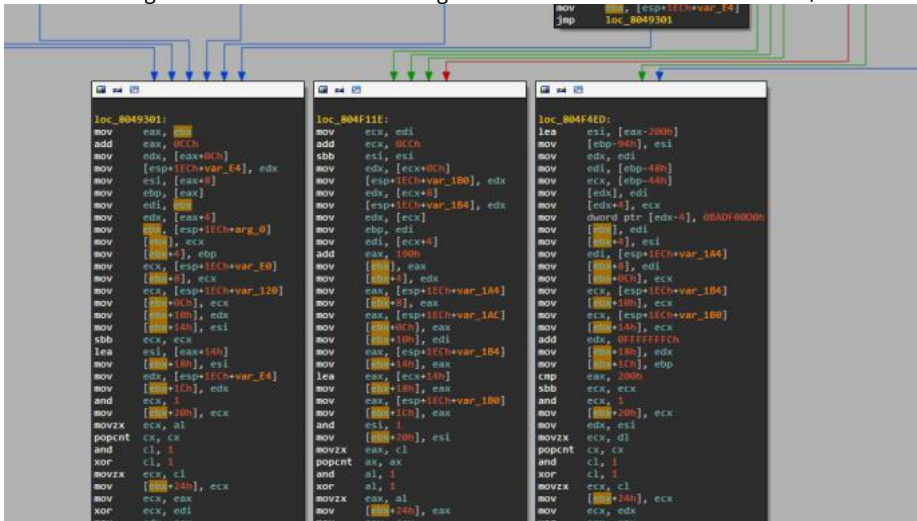
7. Here I tried to get as many character possible from the given string, but only 3 char was possible... :/



I crossed **flag[13:16] = "aba"**

and so on... their are many checks further but i think you got the idea :)

One more thing i found is that in the following 3 blocks at the end of **DEVIL** function, we can only *return 0* if we can get to the block **loc_804F4ED**.



NOTE : If anybody found any better solution or any method to de-obfuscate it using code, please tell me :) I did all of it manually.

NOTE: I haven't been able to confirm the key as nobody is replying on #csaw irc channel. but I am getting "You win!" as output :)

UPDATE 1 : I confirmed with admin. Flag i got is correct :)

Thanks to **Artem** for this amazing problem. I don't think i solved it by the way he was expecting :(but still I got the flag :P

5. Flag

Flag I got is:

flag{1_ea6lm6aba0faad77703151f6666667}