

Denotational Semantics of MicroScala

Abstract Syntax

Syntax Domains:

C: Compilation Unit

M: Main Definition

D: Definition

V: Var Definition

T: Type

P: Formal Parameters

F: Function Body

S: Statement

E: Expression

A: Arguments

L: Literal

Vid: Variable Id

Fid: Function Id

N: Integer

Abstract Syntax Rules:

C ::= **D M** | **M**

M ::= **V S** | **S**

D ::= **D₁ ; D₂** | **Fid (P) : T = F** | **V**

V ::= **V₁ ; V₂** | **Vid : T = L**

T ::= **Int** | **List**

P ::= **P₁ Id : T** | ϵ

F ::= **V S return E** | **S return E** | **V return E** | **return E**

S ::= **S₁ S₂** | **Vid = E** | **if E S** | **if E S₁ else S₂** | **while E S** | **println E**

E ::= **E₁ || E₂** | **E₁ && E₂**

| **E₁ < E₂** | **E₁ <= E₂** | **E₁ > E₂** | **E₁ >= E₂** | **E₁ == E₂** | **E₁ != E₂**

| **+E** | **-E** | **E₁ + E₂** | **E₁ - E₂** | **E₁ * E₂** | **E₁ / E₂**

| **E₁ :: E₂** | **E . head** | **E . tail** | **E . isEmpty**

| **Vid** | **Fid A** | **L**

A ::= **A E** | ϵ

L ::= **N** | **Nil**

Semantic Algebra:

Domain $n \in \mathbb{N} = \{ \dots, -2, -1, 0, 1, 2, \dots \}$ (Integers)

$AO \llbracket + \rrbracket \quad n_1 \ n_2 = (n_1 + n_2)$
 $AO \llbracket - \rrbracket \quad n_1 \ n_2 = (n_1 - n_2)$
 $AO \llbracket * \rrbracket \quad n_1 \ n_2 = (n_1 \times n_2)$
 $AO \llbracket / \rrbracket \quad n_1 \ n_2 = \text{if } (n_2 = 0) \text{ then } \perp \text{ else } (n_1 / n_2)$

$RO \llbracket < \rrbracket \quad n_1 \ n_2 = \text{if } (n_1 < n_2) \text{ then } true \text{ else } false$
 $RO \llbracket \leq \rrbracket \quad n_1 \ n_2 = \text{if } (n_1 \leq n_2) \text{ then } true \text{ else } false$
 $RO \llbracket > \rrbracket \quad n_1 \ n_2 = \text{if } (n_1 > n_2) \text{ then } true \text{ else } false$
 $RO \llbracket \geq \rrbracket \quad n_1 \ n_2 = \text{if } (n_1 \geq n_2) \text{ then } true \text{ else } false$
 $RO \llbracket == \rrbracket \quad n_1 \ n_2 = \text{if } (n_1 = n_2) \text{ then } true \text{ else } false$
 $RO \llbracket != \rrbracket \quad n_1 \ n_2 = \text{if } (n_1 \neq n_2) \text{ then } true \text{ else } false$

Domain $p \in \text{List} = \mathbb{N} * (\text{List of integers})$

nil nil returns an empty list i.e., $\langle | \rangle$
 $cons : n \rightarrow p_1 \rightarrow p_2$ $cons$ takes an integer n and list p_1 ;
 attaches n to the head of p_1 and returns the resultant list p_2
 $hd : p \rightarrow n$ hd takes a list p and returns the head of p
 $tl : p_1 \rightarrow p_2$ tl takes the list p_1 and returns the tail of p_1 as list p_2
 $append : n \rightarrow p_1 \rightarrow p_2$ $append$ takes an integer n and list p_1 ;
 attaches n to the end of p_1 and returns the resultant list p_2
 $eqlist : \text{List} \times \text{List} \rightarrow \text{Integer}$ $eqlist$ compares the two lists for equality
 $neqlist : \text{List} \times \text{List} \rightarrow \text{Integer}$ $neqlist$ compares the two lists for inequality

Domain $\text{varid} \in \text{Variable Id}$

Domain $\text{funcid} \in \text{Function Id}$

Domain $\text{id} \in \text{Id} = \text{Variable Id} + \text{Function Id}$

Domain $t \in \text{Type} = \{ integer, boolean, list \}$

Domain $v \in \text{Value} = \mathbb{N} + \text{List}$

Domain $e \in \text{Expressible-value} = (\text{Type} \times \text{Value}^\circ)$

$Exprval : \text{Type} \rightarrow \text{Value}^\circ \rightarrow \text{Expressible-value}$
 $Exprval \ t \ v = (t, v)$ “constructs the expressible value tuple”
 $type : \text{Expressible-value} \rightarrow \text{Type}$ | $value : \text{Expressible-value} \rightarrow \text{Value}^\circ$
 $type \ e = (e \downarrow 1)$ | $value \ e = (e \downarrow 2)$
 “selects the type component” | “selects the value component”

Domain $\text{fb} \in \text{Function-body} = \text{Conf} \rightarrow (\text{Conf} \times \text{Expressible-value})$

Domain $f \in \text{Function-denot} = \text{Formal Parameters} \times \text{Type} \times \text{Function-body}$

$\text{Function-denot} : \text{Formal Parameters} \rightarrow \text{Type} \rightarrow \text{Function-body} \rightarrow \text{Function-denot}$

$\text{Function-denot } fp \ t \ fb = (fp, t, fb)$ “constructs the function denotation triple”

$\text{formpars} : \text{Function-denot} \rightarrow \text{Formal Parameters}$

$\text{formpars } f = (f \downarrow 1)$ “selects the formal parameter declarations”

$\text{type} : \text{Function-denot} \rightarrow \text{Type}$

$\text{type } f = (f \downarrow 2)$

“selects the type component”

$\text{funcbody} : \text{Function-denot} \rightarrow \text{Function-body}$

$\text{funcbody } f = (f \downarrow 3)$

“selects the function body component”

Domain $d \in \text{Denotable-value} = \text{Expressible-value} + \text{Function-denot}$

Domain $c \in \text{Conf} = \{(\text{Env} \times \text{Env} \times \text{File})\}^\circ$

$\text{Conf} : \text{Env} \rightarrow \text{Env} \rightarrow \text{File} \rightarrow \text{Conf}$

$\text{Conf } env_{\text{global}} \ env_{\text{local}} \ fl_{\text{out}} = (env_{\text{global}}, env_{\text{local}}, fl_{\text{out}})$ “constructs the configuration tuple”

$\text{global-env} : \text{Conf} \rightarrow \text{Env}$ | $\text{local-env} : \text{Conf} \rightarrow \text{Env}$ | $\text{outfile} : \text{Conf} \rightarrow \text{File}$

$\text{env } c = (c \downarrow 1)$ | $\text{env } c = (c \downarrow 2)$ | $\text{outfile } c = (c \downarrow 3)$

“selects the environment components” | “selects the output file component”

Domain $fl \in \text{File} = \text{List}$

Domain $env \in \text{Env} = \{(\text{Id} \rightarrow \text{Denotable-value})\}^\circ$

$\text{new-env} : \text{Env}$

$\text{new-env} = \lambda id . \perp$

$\text{access-env} : \text{Id} \rightarrow \text{Env} \rightarrow \text{Denotable-value}$

$\text{access-env } id \ env = env \ id$

$\text{update-env} : \text{Id} \rightarrow \text{Denotable-value} \rightarrow \text{Env} \rightarrow \text{Env}$

$\text{update-env } id \ d \ env =$

let $d = \text{access-env } id \ env$

in if $d \neq \perp$ then \top else $env \ [d / id]$ “error if already defined”

$\text{update-val} : \text{Variable Id} \rightarrow \text{Expressible-value} \rightarrow \text{Env} \rightarrow \text{Env}$

$\text{update-val } varId \ e \ env =$

let $d = \text{access-env } varId \ env$

in if $d = \perp$ then \top else $env \ [e / id]$ “error if not already defined”

Semantic Functions:

$C : \text{Compilation Unit} \rightarrow \text{File}_{\text{output}}$
 $M : \text{Main Definition} \rightarrow \text{Env} \rightarrow \text{Conf} \rightarrow \text{Conf}$
 $D : \text{Definition} \rightarrow \text{Env} \rightarrow \text{Env}$
 $V : \text{Var Definition} \rightarrow \text{Env} \rightarrow \text{Env}$
 $F : \text{Function Body} \rightarrow \text{Conf} \rightarrow (\text{Conf} \times \text{Expressible-value})$
 $S : \text{Statement} \rightarrow \text{Conf} \rightarrow \text{Conf}$
 $E : \text{Expression} \rightarrow \text{Conf} \rightarrow (\text{Conf} \times \text{Expressible-value})$
 $P : \text{Formal Parameters} \rightarrow \text{Env} \rightarrow \text{Arguments} \rightarrow \text{Conf} \rightarrow (\text{Env} \times \text{Conf})$
 $L : \text{Literal} \rightarrow \text{Expressible-value}$

Semantic Equations:

COMPILATION UNIT:

$C \llbracket \mathbf{D} \mathbf{M} \rrbracket =$
 let $\text{env}_{\text{global}} = D \llbracket \mathbf{D} \rrbracket \text{ new-env}$
 in $\text{outfile } M \llbracket \mathbf{M} \rrbracket (\text{Conf } \text{env}_{\text{global}} \text{ new-env nil})$

$C \llbracket \mathbf{M} \rrbracket =$
 let $\text{env}_{\text{global}} = \text{new-env}$
 in $\text{outfile } M \llbracket \mathbf{M} \rrbracket (\text{Conf } \text{env}_{\text{global}} \text{ new-env nil})$

MAIN PROGRAM:

$M \llbracket \mathbf{V} \mathbf{S} \rrbracket \text{ env}_{\text{global}} \mathbf{c} =$
 let $\text{env}_{\text{local}} = V \llbracket \mathbf{V} \rrbracket (\text{new-env } \mathbf{c})$
 in $S \llbracket \mathbf{S} \rrbracket (\text{Conf } \text{env}_{\text{global}} \text{ env}_{\text{local}} (\text{outfile } \mathbf{c}))$

$M \llbracket \mathbf{S} \rrbracket \text{ env}_{\text{global}} \mathbf{c} =$
 let $\text{env}_{\text{local}} = \text{new-env}$
 in $S \llbracket \mathbf{S} \rrbracket (\text{Conf } \text{env}_{\text{global}} \text{ env}_{\text{local}} (\text{outfile } \mathbf{c}))$

DEFINITION:

$D \llbracket \mathbf{D}_1 ; \mathbf{D}_2 \rrbracket \text{ env} = D \llbracket \mathbf{D}_2 \rrbracket (D \llbracket \mathbf{D}_1 \rrbracket \text{ env})$

$D \llbracket \mathbf{Fid} (\mathbf{P}) : \mathbf{T} = \mathbf{F} \rrbracket \text{ env} = \text{let } \mathbf{d} = \text{funcDenot } \llbracket \mathbf{P} \rrbracket (T \llbracket \mathbf{T} \rrbracket) (F \llbracket \mathbf{F} \rrbracket)$
 in $(\text{update-env } \llbracket \mathbf{Fid} \rrbracket \mathbf{d} \text{ env})$

$D \llbracket \mathbf{V} \rrbracket \text{ env} = V \llbracket \mathbf{V} \rrbracket \text{ env}$

VAR DEFINITION:

$$V[[V_1 ; V_2]] \text{ env} = V[[V_2]] (V[[V_1]] \text{ env})$$
$$\begin{aligned} V[[\mathbf{Vid} : \mathbf{T} = \mathbf{L}]] \text{ env} = \\ \text{let } t = (T[[\mathbf{T}]]) \\ \text{in } \quad \text{let } e = (L[[\mathbf{L}]]) \\ \quad \text{in if } (t = \text{integer and type } e = \text{integer and value } e = 0) \text{ or} \\ \quad \quad (t = \text{list and type } e = \text{list and value } e = \text{nil}) \text{ then} \\ \quad \quad \text{update-env } [[\mathbf{Vid}]] e \text{ env} \\ \quad \text{else } \perp \text{ "type error"} \end{aligned}$$

TYPE:

$$T[[\mathbf{Int}]] = \text{integer}$$
$$T[[\mathbf{List}]] = \text{list}$$

FUNCTION BODY:

$$\begin{aligned} F[[\mathbf{V} \mathbf{S} \text{ return } \mathbf{E}]] c = \\ \text{let env}_{\text{local}} = V[[\mathbf{V}]] \text{ new-env} \\ \text{in } \quad \text{let } c_1 = S[[\mathbf{S}]] (Conf(\text{global-env } c) \text{ env}_{\text{local}} (\text{outfile } c)) \\ \quad \text{in } \quad E[[\mathbf{E}]] c_1 \end{aligned}$$
$$\begin{aligned} F[[\mathbf{S} \text{ return } \mathbf{E}]] c = \\ \text{let env}_{\text{local}} = \text{new-env} \\ \text{in } \quad \text{let } c_1 = S[[\mathbf{S}]] (Conf(\text{global-env } c) \text{ env}_{\text{local}} (\text{outfile } c)) \\ \quad \text{in } \quad E[[\mathbf{E}]] c_1 \end{aligned}$$
$$\begin{aligned} F[[\mathbf{V} \text{ return } \mathbf{E}]] c = \\ \text{let env}_{\text{local}} = V[[\mathbf{V}]] \text{ new-env} \\ \text{in } \quad \text{let } c_1 = Conf(\text{global-env } c) \text{ env}_{\text{local}} (\text{outfile } c) \\ \quad \text{in } \quad E[[\mathbf{E}]] c_1 \end{aligned}$$
$$\begin{aligned} F[[\text{return } \mathbf{E}]] c = \\ \text{let env}_{\text{local}} = \text{new-env} \\ \text{in } \quad \text{let } c_1 = Conf(\text{global-env } c) \text{ env}_{\text{local}} (\text{outfile } c) \\ \quad \text{in } \quad E[[\mathbf{E}]] c_1 \end{aligned}$$

STATEMENT:

$S \llbracket [S_1 \ S_2] \rrbracket c = S \llbracket [S_2] \rrbracket (S \llbracket [S_1] \rrbracket c)$

$S \llbracket [\mathbf{Vid} = \mathbf{E}] \rrbracket c =$

- let $d_{\text{local}} = \text{access-env} \llbracket [\mathbf{Vid}] \rrbracket (\text{local-env } c)$
- in if $d_{\text{local}} \neq \perp$ then
 - let $(c_1, e) = E \llbracket [\mathbf{E}] \rrbracket c$
 - in if $(\text{type } d \neq \text{type } e)$ then \perp “type error”
 - else let $\text{env}_{\text{local}} = (\text{update-val} \llbracket [\mathbf{Vid}] \rrbracket e (\text{local-env } c_1))$
 - in $(\text{Conf } (\text{global-env } c) \text{env}_{\text{local}} (\text{outfile } c_1))$
- else
 - let $d_{\text{global}} = \text{access-env} \llbracket [\mathbf{Vid}] \rrbracket (\text{global-env } c)$
 - in if $d_{\text{global}} \neq \perp$ then
 - let $(c_1, e) = E \llbracket [\mathbf{E}] \rrbracket c$
 - in if $(\text{type } d \neq \text{type } e)$ then \perp “type error”
 - else let $\text{env}_{\text{global}} =$
 $(\text{update-val} \llbracket [\mathbf{Vid}] \rrbracket e (\text{global-env } c_1))$
 - in $(\text{Conf } \text{env}_{\text{global}} (\text{local-env } c_1) (\text{outfile } c_1))$
 - else \perp “undeclared variable error”

$S \llbracket [\mathbf{if } \mathbf{E} \ \mathbf{S}] \rrbracket c =$

- let $(c_1, e) = E \llbracket [\mathbf{E}] \rrbracket c$
- in if $(\text{type } e \neq \text{boolean})$ then \perp “type error”
- else if $(\text{value } e)$ then $S \llbracket [\mathbf{S}] \rrbracket c_1$
- else c_1

$S \llbracket [\mathbf{if } \mathbf{E} \ \mathbf{S}_1 \ \mathbf{else} \ \mathbf{S}_2] \rrbracket c =$

- let $(c_1, e) = E \llbracket [\mathbf{E}] \rrbracket c$
- in if $(\text{type } e \neq \text{boolean})$ then \perp “type error”
- else if $(\text{value } e)$ then $S \llbracket [\mathbf{S}_1] \rrbracket c_1$
- else $S \llbracket [\mathbf{S}_2] \rrbracket c_1$

$S \llbracket [\mathbf{while } \mathbf{E} \ \mathbf{S}] \rrbracket c =$

- let $(c_1, e) = E \llbracket [\mathbf{E}] \rrbracket c$
- in if $(\text{type } e \neq \text{boolean})$ then \perp “type error”
- else if $(\text{value } e)$ then $S \llbracket [\mathbf{while } \mathbf{E} \ \mathbf{S}] \rrbracket (S \llbracket [\mathbf{S}] \rrbracket c_1)$
- else c_1

$S \llbracket [\mathbf{println } \mathbf{E}] \rrbracket c =$

- let $(c_1, e) = (E \llbracket [\mathbf{E}] \rrbracket c)$
- in let $\text{fl} = (\text{outfile } c_1)$
- in let $\text{fl}_1 = (\text{append } (\text{value } e) \text{fl})$
- in $\text{Conf } (\text{global-env } c_1) (\text{local-env } c_1) \text{fl}_1$

EXPRESSIONS:

$E [[E_1 \ \&\& \ E_2]] \ c =$
 let $(c_1, e_1) = E [[E_1]] \ c$
 in if $(type \ e_1 \neq \text{boolean})$ then \perp “type error”
 else if $(value \ e_1)$ then
 let $(c_2, e_2) = E [[E_2]] \ c_1$
 in if $(type \ e_2 \neq \text{boolean})$ then \perp “type error”
 else $(c_2, value \ e_2)$
 else $(c_1, value \ e_1)$

$E [[E_1 \ || \ E_2]] \ c =$
 let $(c_1, e_1) = E [[E_1]] \ c$
 in if $(type \ e_1 \neq \text{boolean})$ then \perp “type error”
 else if $(value \ e_1)$ then $(c_1, value \ e_1)$
 else let $(c_2, e_2) = E [[E_2]] \ c_1$
 in if $(type \ e_2 \neq \text{boolean})$ then \perp “type error”
 else $(c_2, value \ e_2)$

$E [[E_1 \ \text{rop} \ E_2]] \ c =$
 let $(c_1, e_1) = E [[E_1]] \ c$
 in let $(c_2, e_2) = E [[E_2]] \ c_1$
 in if $(type \ e_1 \neq \text{integer})$ or $(type \ e_2 \neq \text{integer})$ then \perp “type error”
 else let $e_{\text{result}} = Exprval \ \text{boolean} \ (RO[[\text{rop}]] \ (value \ e_1) \ (value \ e_2))$
 in (c_2, e_{result})
 where $\text{rop} \in \{<, >, <=, >=\}$

$E [[E_1 \ == \ E_2]] \ c =$
 let $(c_1, e_1) = E [[E_1]] \ c$
 in let $(c_2, e_2) = E [[E_2]] \ c_1$
 in if $(type \ e_1 = \text{integer})$ and $(type \ e_2 = \text{integer})$ then
 $Exprval \ \text{boolean} \ (RO [[==]] \ (value \ e_1) \ (value \ e_2))$
 else if $(type \ e_1 = \text{list})$ and $(type \ e_2 = \text{list})$ then
 $Exprval \ \text{boolean} \ (eqlist \ (value \ e_1) \ (value \ e_2))$
 else \perp “type error”

$E [[E_1 \ != \ E_2]] \ c =$
 let $(c_1, e_1) = E [[E_1]] \ c$
 in let $(c_2, e_2) = E [[E_2]] \ c_1$
 in if $(type \ e_1 = \text{integer})$ and $(type \ e_2 = \text{integer})$ then
 $Exprval \ \text{boolean} \ (RO [[!=]] \ (value \ e_1) \ (value \ e_2))$
 else if $(type \ e_1 = \text{list})$ and $(type \ e_2 = \text{list})$ then
 $Exprval \ \text{boolean} \ (neqlist \ (value \ e_1) \ (value \ e_2))$
 else \perp “type error”

$$\begin{array}{l}
E \llbracket \mathbf{aop} \ E \rrbracket \ c = \\
\quad \text{let } (c_1, e_1) = E \llbracket E \rrbracket \ c \\
\quad \text{in} \quad \text{if } (\text{type } e_1 \neq \text{integer}) \text{ then } \top \text{ "type error"} \\
\quad \quad \text{else} \quad \text{let } e_{\text{result}} = \text{Exprval integer } (AO[\llbracket \mathbf{aop} \rrbracket] \ 0 \ (\text{value } e_1)) \\
\quad \quad \quad \text{in } (c_1, e_{\text{result}}) \\
\text{where } \mathbf{aop} \in \{ +, - \}
\end{array}$$
$$\begin{array}{l}
E \llbracket [E_1, \mathbf{aop} E_2] \rrbracket c = \\
\quad \text{let } (c_1, e_1) = E \llbracket [E_1] \rrbracket c \\
\quad \text{in} \quad \text{let } (c_2, e_2) = E \llbracket [E_2] \rrbracket c_1 \\
\quad \quad \text{in} \quad \text{if } (\text{type } e_1 \neq \text{integer}) \text{ or } (\text{type } e_2 \neq \text{integer}) \text{ then } \perp \text{ “type error”} \\
\quad \quad \quad \text{else} \quad \text{let } e_{\text{result}} = \text{Exprval integer } (AO \llbracket [\mathbf{aop}] \rrbracket (\text{value } e_1) (\text{value } e_2)) \\
\quad \quad \quad \quad \text{in } (c_2, e_{\text{result}}) \\
\text{where } \mathbf{aop} \in \{ +, -, *, / \}
\end{array}$$
$$\begin{aligned}
E \llbracket [\mathbf{E}_1 :: \mathbf{E}_2] \rrbracket c = & \\
\text{let } (c_1, e_1) = (E \llbracket [\mathbf{E}_1] \rrbracket c) & \\
\text{in let } (c_2, e_2) = (E \llbracket [\mathbf{E}_2] \rrbracket c_1) & \\
\text{in if } (type\ e_1 \neq integer) \text{ or } (type\ e_2 \neq list) \text{ then } \perp \text{ “type error”} & \\
\text{else } (c_2, Exprval\ list\ (cons\ (value\ e_1)\ (value\ e_2))) &
\end{aligned}$$
$$E \llbracket [E . \text{head}] \rrbracket c = \begin{array}{l} \text{let } (c_1, e_1) = (E \llbracket [E] \rrbracket c) \\ \text{in} \quad \text{if } (\text{type } e_1 \neq \text{list}) \text{ or } (\text{value } e_1 = \text{nil}) \text{ then } \top \text{ “error”} \\ \quad \text{else } (c_1, \text{Exprval integer } (hd (\text{value } e_1))) \end{array}$$
$$E \llbracket \mathbf{E} . \mathbf{tail} \rrbracket c = \text{let } (c_1, e_1) = (E \llbracket \mathbf{E} \rrbracket c) \\ \text{in if } (type\ e_1 \neq list) \text{ or } (value\ e_1 = nil) \text{ then } \top \text{ “error”} \\ \text{else } (c_1, Exprval\ list\ (tl\ (value\ e_1)))$$
$$\begin{array}{l}
E \llbracket \mathbf{E} . \mathbf{isEmpty} \rrbracket \text{ env } c = \\
\quad \text{let } (c_i, e_i) = (E \llbracket \mathbf{E} \rrbracket \text{ env } c) \\
\quad \text{in} \quad \text{if } (type\ e_i \neq list) \text{ then } \top \text{ “type error”} \\
\quad \quad \text{else } (c_i, Exprval\ boolean\ (value\ e_i = nil))
\end{array}$$
$$\begin{array}{l}
E \llbracket \mathbf{Vid} \rrbracket c = \\
\quad \text{let } d_{\text{local}} = \text{access-env} \llbracket \mathbf{Vid} \rrbracket (\text{local-env } c) \\
\quad \text{in} \quad \text{if } d_{\text{local}} \neq \perp \text{ then } (c, d_{\text{local}}) \\
\quad \quad \text{else let } d_{\text{global}} = \text{access-env} \llbracket \mathbf{Vid} \rrbracket (\text{global-env } c) \\
\quad \quad \quad \text{in} \quad \text{if } d_{\text{global}} \neq \perp \text{ then } (c, d_{\text{global}}) \\
\quad \quad \quad \text{else } \top \text{ “undeclared variable error”}
\end{array}$$
$$E \llbracket \mathbf{L} \rrbracket \mathbf{c} = (\mathbf{c}, L \llbracket \mathbf{L} \rrbracket)$$

$E \llbracket \mathbf{Fid} \ A \rrbracket c =$
 let $d = \text{access-env} \llbracket \mathbf{Fid} \rrbracket (\text{global-env } c)$
 in if $d = \perp$ then “undeclared function error”
 else let $(\text{env}_{\text{local}}, c_1) = PM(\text{formpar } d) \text{ new-env } \llbracket A \rrbracket c$ “evaluate actuals”
 in let $(c_2, e_{\text{result}}) =$
 $(\text{funbody } d) (\text{Conf}(\text{global-env } c_1) \text{ env}_{\text{local}} (\text{outfile } c_1))$ “call”
 in $(\text{Conf}(\text{global-env } c_2) (\text{local-env } c_1) (\text{outfile } c_2), e_{\text{result}})$ “restore caller env”

PARAMETER MATCHING:

$P \llbracket \mathbf{T} \ \mathbf{Vid} \ P \rrbracket \text{env}_{\text{local}} \llbracket E \ A \rrbracket c =$
 let $(c_1, e) = E \llbracket E \rrbracket c$
 in if $(\text{type } e \neq (T \llbracket T \rrbracket))$ then “type error”
 else let $\text{env}_1 = V \llbracket \mathbf{T} \ \mathbf{Vid} \rrbracket \text{env}_{\text{local}}$ “add formals into local env”
 in let $\text{env}_2 = (\text{update-val } \llbracket \mathbf{Vid} \rrbracket e \text{ env}_1)$ “set formals to actuals”
 in $P \llbracket P \rrbracket \text{env}_2 \llbracket A \rrbracket c_1$

$P \llbracket \] \rrbracket \text{env}_{\text{local}} \llbracket \] \rrbracket c = (\text{env}_{\text{local}}, c)$

$P \llbracket \] \rrbracket \text{env}_{\text{local}} \llbracket E \ A \rrbracket c = \top$ “number of parameters mismatch”

$P \llbracket \mathbf{T} \ \mathbf{Vid} \ P \rrbracket \text{env}_{\text{local}} \llbracket \] \rrbracket c = \top$ “number of parameters mismatch”

LITERAL:

$L \llbracket N \rrbracket = \text{Exprval integer } (N \llbracket N \rrbracket)$

$L \llbracket \mathbf{Nil} \rrbracket = \text{Exprval list nil}$