

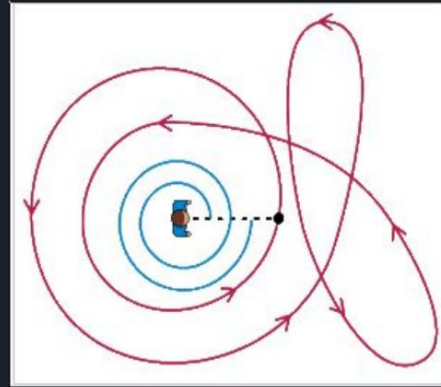
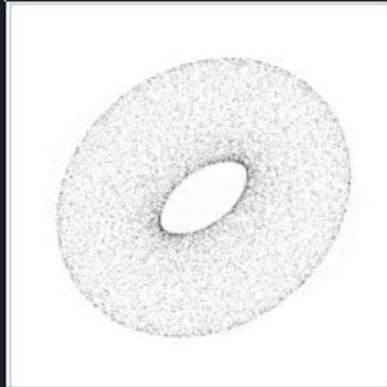
A decorative graphic on the left side of the slide consisting of two overlapping parallelograms. The front one is blue and the back one is a light green. They are positioned diagonally, with the blue one partially covering the green one.

Fast Winding Numbers for Soups and Clouds

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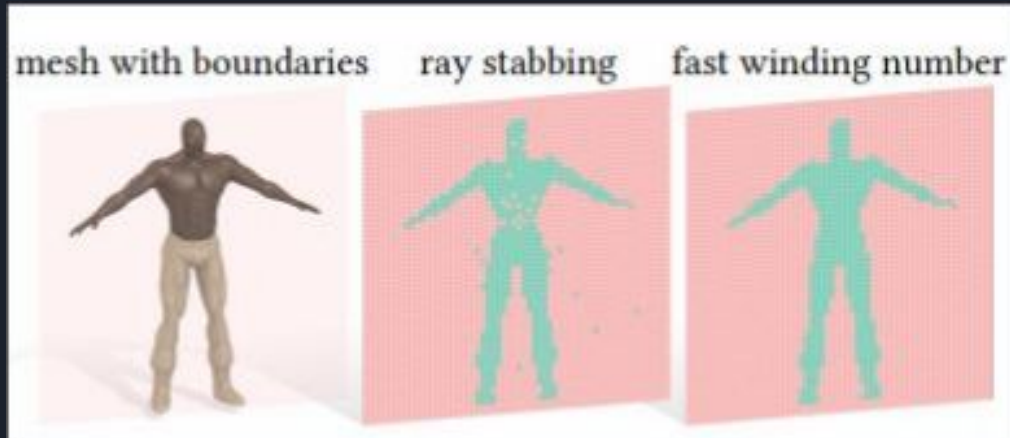
Definitions

- Point Cloud: is a set of data points in space. Point clouds are generally produced by 3D scanners, which measure a large number of points on the external surfaces of objects around them.
- Winding Number: The winding number of the curve is equal to the total number of counterclockwise turns that the object makes around the origin.



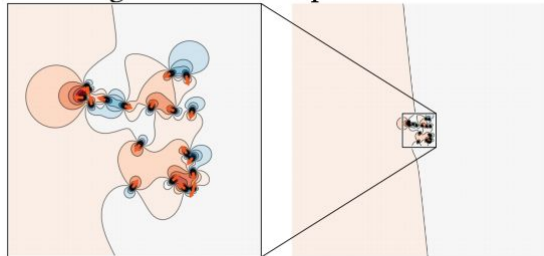
Advantages

- The winding number in a variety of new applications: voxelization, signing distances, generating 3D printer paths, defect-tolerant mesh booleans and point set surfaces.
- Determines how many times a planar curve encircles a query point.
- For overlapping regions, the winding number measures how many times the region is inside the surface.

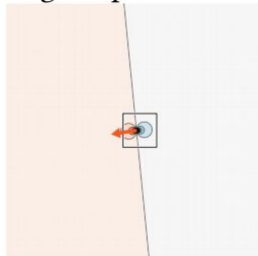


Surface Points

Winding number of 20 points



Single representative



Log of absolute value

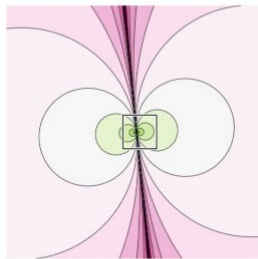
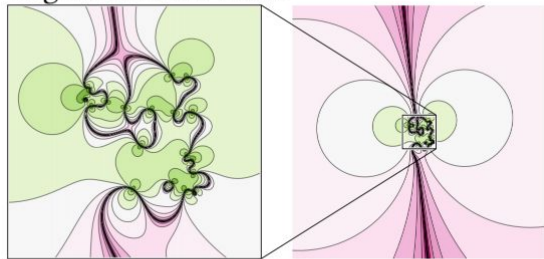
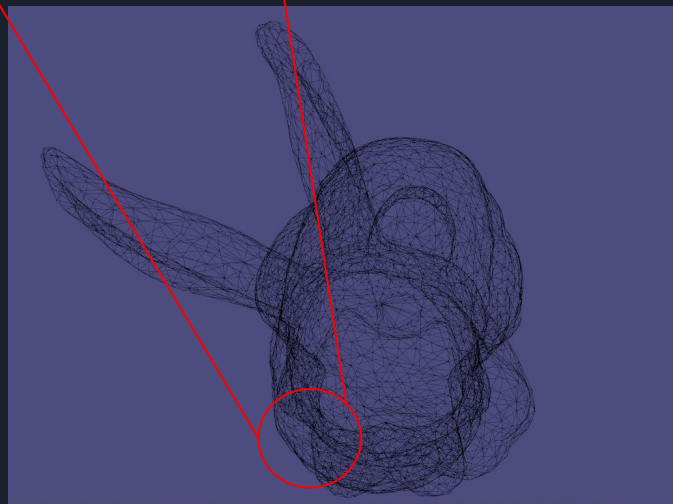
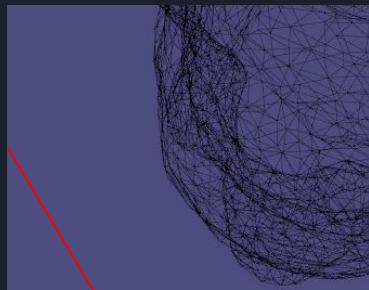


Fig. 6. A cluster of 20 dipoles has an intricate winding number field nearby (left), but far away their function is quite tame (middle) and well approximated by a single, *stronger* dipole (right).



Algorithm to calculate Winding Numbers

$$\frac{(\mathbf{x} - \mathbf{q}) \cdot \hat{\mathbf{n}}}{4\pi \|\mathbf{x} - \mathbf{q}\|^3} = \nabla \left(\frac{-1}{4\pi \|\mathbf{x} - \mathbf{q}\|} \right) \cdot \hat{\mathbf{n}} =: G_{\hat{\mathbf{n}}}(\mathbf{q}, \mathbf{x})$$
$$w(\mathbf{q}) = \sum_{i=1}^m a_i \frac{(\mathbf{p}_i - \mathbf{q}) \cdot \hat{\mathbf{n}}_i}{4\pi \|\mathbf{p}_i - \mathbf{q}\|^3} \approx \frac{(\tilde{\mathbf{p}} - \mathbf{q}) \cdot \tilde{\mathbf{n}}}{4\pi \|\tilde{\mathbf{p}} - \mathbf{q}\|^3} =: \tilde{w}(\mathbf{q})$$
$$\tilde{\mathbf{n}} = \sum_{i=1}^m a_i \hat{\mathbf{n}}_i, \quad \tilde{\mathbf{p}} = \frac{\sum_{i=1}^m a_i \mathbf{p}_i}{\sum_{i=1}^m a_i},$$

```
// Extract interior tets
MatrixXi CT((W.array()>0.5).count(),4);
{
    size_t k = 0;
    for(size_t t = 0;t<T.rows();t++)
    {
        if(W(t)>0.5)
        {
            CT.row(k) = T.row(t);
            k++;
        }
    }
}
// find boundary facets of interior tets
igl::boundary_facets(CT,G);
// boundary_facets seems to be reversed...
G = G.rowwise().reverse().eval();

// normalize
W = (W.array() - W.minCoeff())/(W.maxCoeff()-W.minCoeff());
```

Algorithm fast approximation

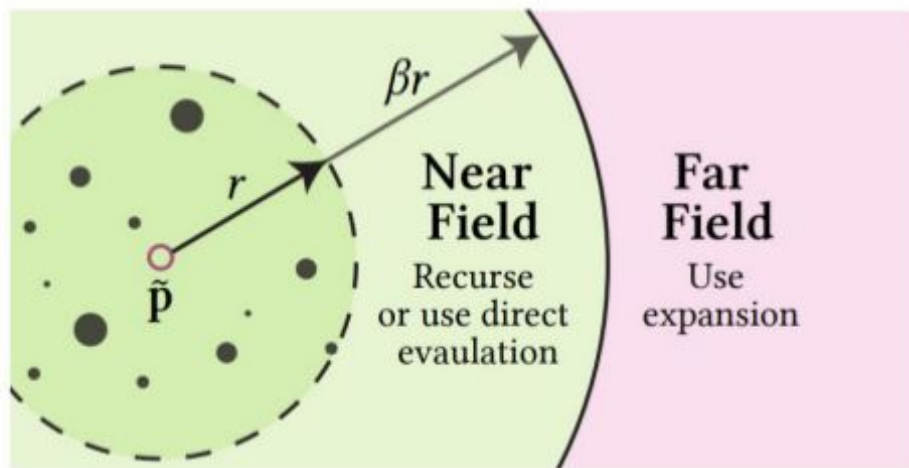


Fig. 7. Our spatial partitioning separates near and far fields, recursively.

Algorithm 1: Fast Approximation of Winding Number $\text{FASTWN}(q, \text{tree})$

Inputs:

q Query point in \mathbb{R}^3
 tree Root of bounding volume hierarchy for points/triangles
 β accuracy parameter

Outputs: scalar winding number of all elements in tree at q
// tree.p : center of tree's winding number approximation, tree.w
// tree.r : maximum distance from tree.p to any of its elements

```
if  $\|q - \text{tree.p}\| > \beta * \text{tree.r}$  then
    //  $q$  is sufficiently far from all elements in tree
    return  $\text{tree.w}(q)$ 
else
    val  $\leftarrow 0$ 
    if tree has no children then
        //  $q$  is nearby; use direct sum for tree's elements
        for each point/triangle  $e$  in tree do
            //  $w_e$ : area-weighted dipole or solid angle
            val  $+= w_e(q)$ 
    else
        for each child of tree do
            // Recursive call
            val  $+= \text{FASTWN}(q, \text{child})$ 
    return val
```

Fast Algorithm to calculate Winding Numbers

$$\begin{aligned} G_{\hat{n}}(\mathbf{q}, \mathbf{x}) &= \hat{n} \cdot \nabla G(\mathbf{q}, \mathbf{x}) \\ &= \hat{n} \cdot \nabla G(\mathbf{q}, \tilde{\mathbf{p}}) \\ &\quad + ((\mathbf{x} - \tilde{\mathbf{p}}) \otimes \hat{n}) \cdot \nabla^2 G(\mathbf{q}, \tilde{\mathbf{p}}) \\ &\quad + \frac{1}{2}((\mathbf{x} - \tilde{\mathbf{p}}) \otimes (\mathbf{x} - \tilde{\mathbf{p}}) \otimes \hat{n}) \cdot \nabla^3 G(\mathbf{q}, \tilde{\mathbf{p}}) \\ &\quad + \text{higher order terms,} \end{aligned}$$

$$\begin{aligned} w(\mathbf{q}) &\approx \left(\sum_{i=1}^m a_i \hat{n}_i \right) \cdot \nabla G(\mathbf{q}, \tilde{\mathbf{p}}) \\ &\quad + \left(\sum_{i=1}^m a_i (\mathbf{p}_i - \tilde{\mathbf{p}}) \otimes \hat{n}_i \right) \cdot \nabla^2 G(\mathbf{q}, \tilde{\mathbf{p}}) \\ &\quad + \frac{1}{2} \left(\sum_{i=1}^m a_i (\mathbf{p}_i - \tilde{\mathbf{p}}) \otimes (\mathbf{p}_i - \tilde{\mathbf{p}}) \otimes \hat{n}_i \right) \cdot \nabla^3 G(\mathbf{q}, \tilde{\mathbf{p}}) \\ &=: \tilde{w}(\mathbf{q}). \end{aligned}$$

```
HDK_Sample::UT_SolidAngle<float,float> solid_angle;

std::vector<HDK_Sample::UT_Vector3T<float> > U(V.rows());
for(int i = 0; i<V.rows(); i++){
    for(int j = 0; j<3; j++){
        U[i][j] = V(i,j);
    }
}
solid_angle.init(F.rows(), F.data(), V.rows(), &U[0], order);

igl::parallel_for(T.rows(), [&](int q)
//for(int q = 0; q<T.rows(); q++)
{
    HDK_Sample::UT_Vector3T<float> Qq;
    Qq[0] = T(q,0);
    Qq[1] = T(q,1);
    Qq[2] = T(q,2);
    Wapprox(q) = solid_angle.computeSolidAngle(Qq, accuracy_scale)/(4.0*M_PI);
}, 1000);
```


Fast Algorithm to calculate Winding Numbers

Algorithm 1: Fast Approximation of Winding Number FASTWN(q,tree)

Inputs:

q Query point in \mathbb{R}^3
 tree Root of bounding volume hierarchy for points/triangles
 β accuracy parameter

Outputs: scalar winding number of all elements in tree at q

// tree.p: center of tree's winding number approximation, tree.w

// tree.r: maximum distance from tree.p to any of its elements

if $\|q - \text{tree.p}\| > \beta * \text{tree.r}$ then

 // q is sufficiently far from all elements in tree

 return tree.w(q)

else

 val \leftarrow 0

 if tree has no children then

 // q is nearby; use direct sum for tree's elements

 for each point/triangle e in tree do

 // w_e : area-weighted dipole or solid angle

 val += $w_e(q)$

 else

 for each child of tree do

 // Recursive call

 val += FASTWN(q, child)

 return val

T sum = (descend_bits & 1) ? child_data_array[0] : 0;

for (int i = 1; i < nchildren; ++i)

 sum += ((descend_bits > i) & 1) ? child_data_array[i] : 0;

*data_for_parent += sum;

```
for (int i = 0; i < nchildren; ++i)
{
    const LocalData &child_data = child_data_array[i];
    UT_Vector3T<T> displacement = child_data.myAverageP - UT_Vector3T<T>(data_for_parent->myAverageP);
    UT_Vector3T<T> N = child_data.myN;

    // Adjust Nij for the change in centre P
    data_for_parent->myNijDiag += N*displacement;
    T Nxy = child_data.myNxy + N[0]*displacement[1];
    T Nyx = child_data.myNyx + N[1]*displacement[0];
    T Nyz = child_data.myNyz + N[1]*displacement[2];
    T Nzy = child_data.myNzy + N[2]*displacement[1];
    T Nzx = child_data.myNzx + N[2]*displacement[0];
    T Nxz = child_data.myNxz + N[0]*displacement[2];

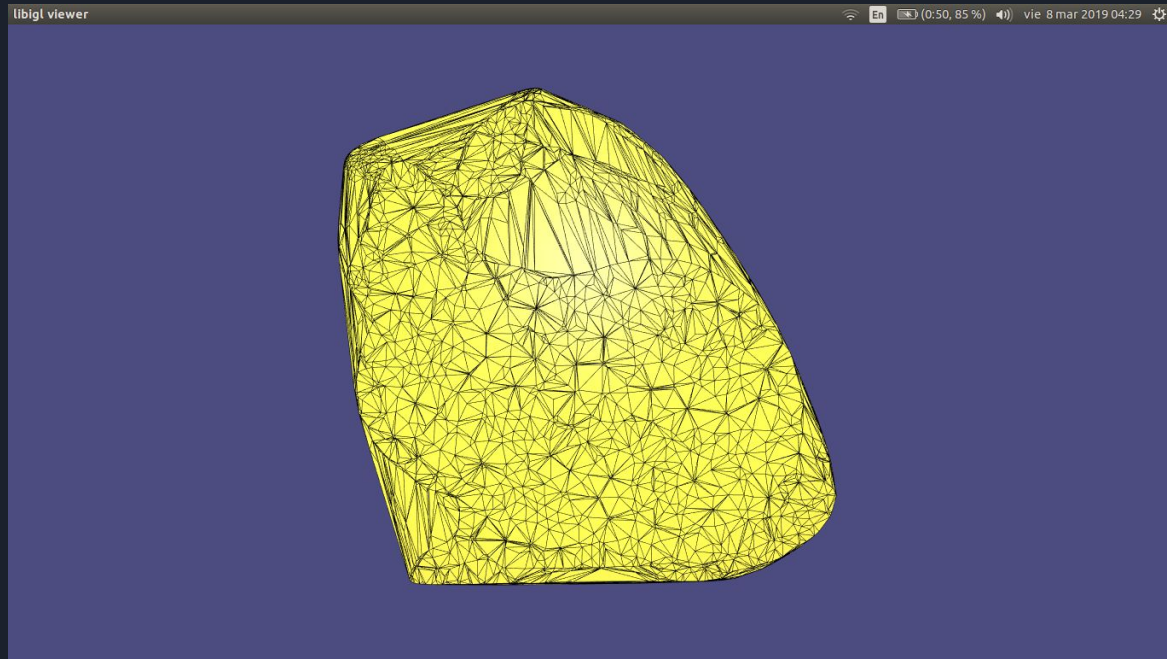
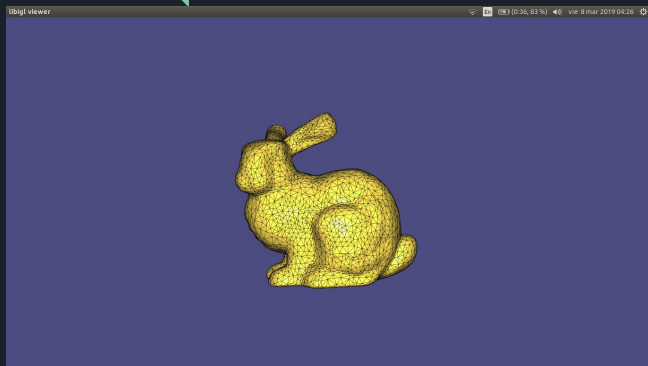
    data_for_parent->myNxy += Nxy;
    data_for_parent->myNyx += Nyx;
    data_for_parent->myNyz += Nyz;
    data_for_parent->myNzy += Nzy;
    data_for_parent->myNzx += Nzx;
    data_for_parent->myNxz += Nxz;

    // Adjust Nijk for the change in centre P
    data_for_parent->myNijkDiag += T(2)*displacement*child_data.myNijDiag + displacement*displacement*child_data.myN;
    data_for_parent->mySumPermuteNxyz += (displacement[0]*(Nyz+Nzy) + displacement[1]*(Nzx+Nxz) + displacement[2]*(Nxy+Nyx));
    data_for_parent->my2Nxxx_Nyxx +=
        2*(displacement[1]*child_data.myNijDiag[0] + displacement[0]*child_data.myNxy + N[0]*displacement[0]*displacement[1])
        + 2*child_data.myNyx*displacement[0] + N[1]*displacement[0]*displacement[0];
    data_for_parent->my2Nxxz_Nzxx +=
        2*(displacement[2]*child_data.myNijDiag[0] + displacement[0]*child_data.myNxz + N[0]*displacement[0]*displacement[2])
        + 2*child_data.myNzx*displacement[0] + N[2]*displacement[0]*displacement[0];
    data_for_parent->my2Nyyz_Nzyy +=
        2*(displacement[2]*child_data.myNijDiag[1] + displacement[1]*child_data.myNyz + N[1]*displacement[1]*displacement[2])
        + 2*child_data.myNzy*displacement[1] + N[2]*displacement[1]*displacement[1];
    data_for_parent->my2Nyxx_Nxyx +=
        2*(displacement[0]*child_data.myNijDiag[1] + displacement[1]*child_data.myNyx + N[1]*displacement[1]*displacement[0])
        + 2*child_data.myNxy*displacement[1] + N[0]*displacement[1]*displacement[1];
    data_for_parent->my2Nxxz_Nzxx +=
        2*(displacement[0]*child_data.myNijDiag[2] + displacement[2]*child_data.myNzx + N[2]*displacement[2]*displacement[0])
        + 2*child_data.myNxz*displacement[2] + N[0]*displacement[2]*displacement[2];
    data_for_parent->my2Nzyz_Nyzz +=
        2*(displacement[1]*child_data.myNijDiag[2] + displacement[2]*child_data.myNzy + N[2]*displacement[2]*displacement[1])
        + 2*child_data.myNyz*displacement[2] + N[1]*displacement[2]*displacement[2];
}
```

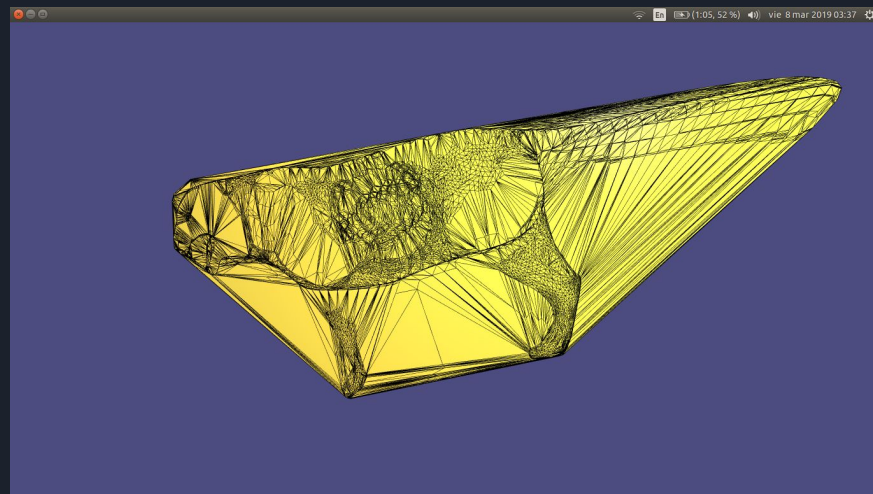
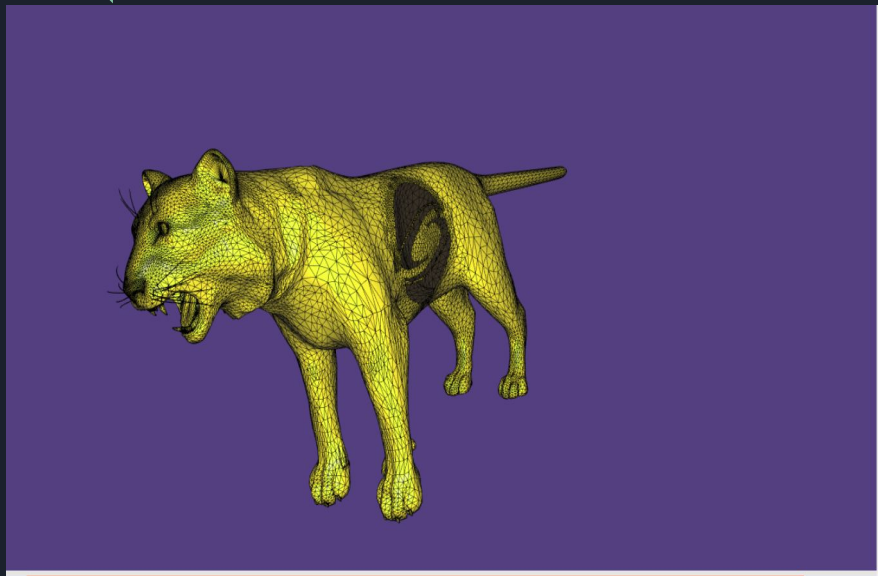

Object representation



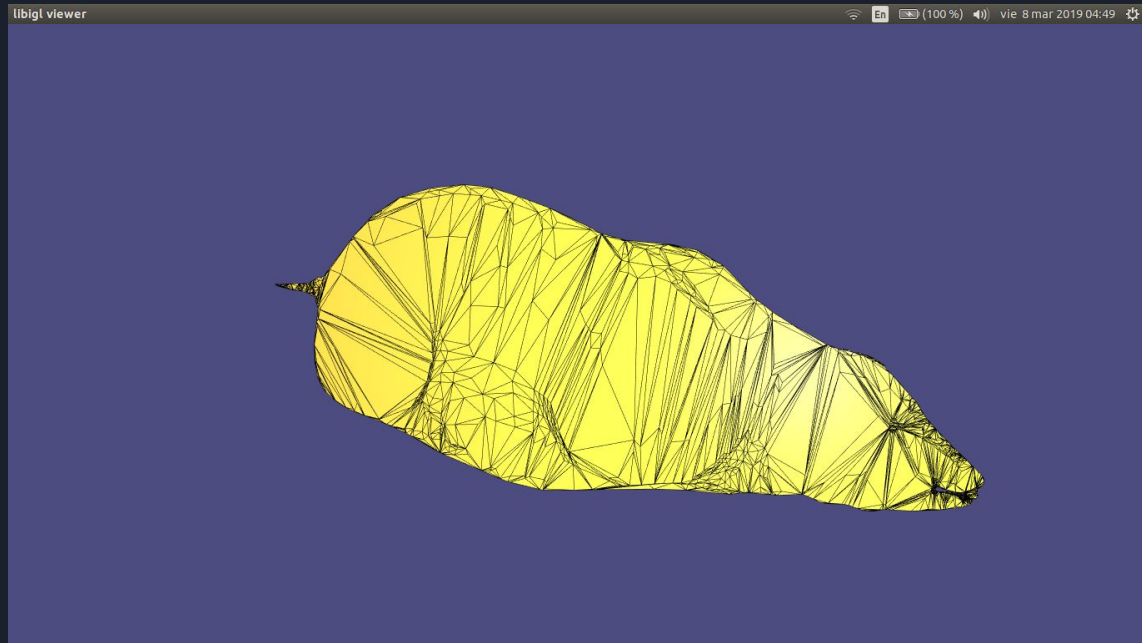
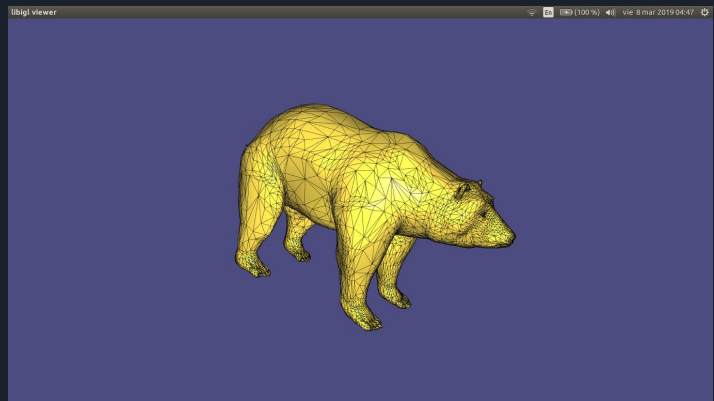
Results - Rabbit



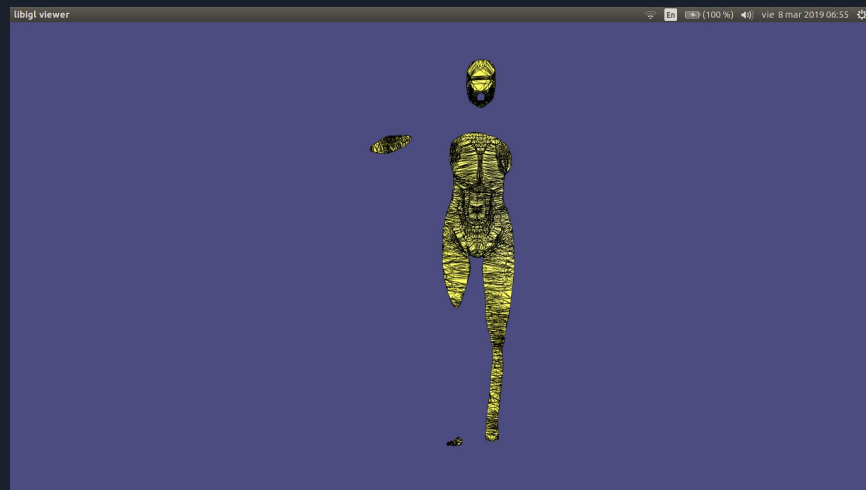
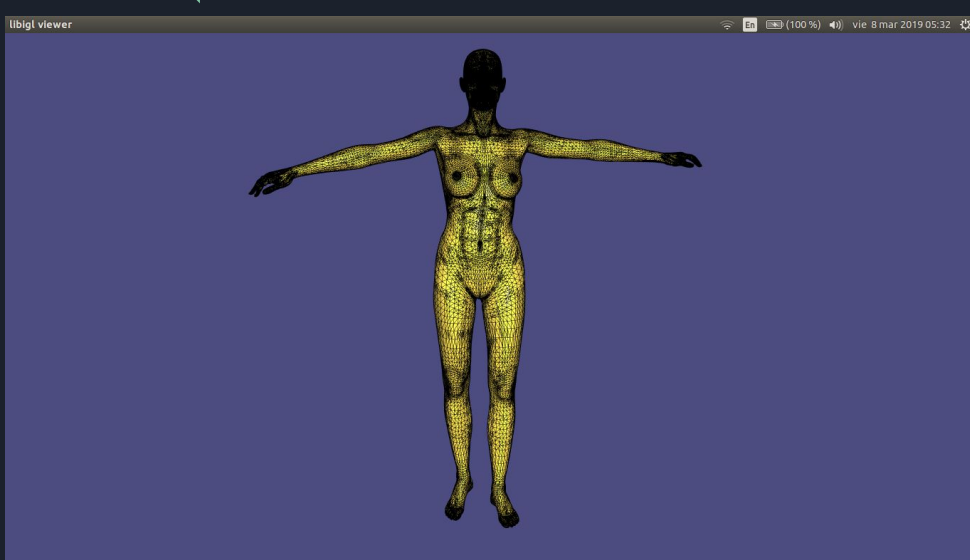
Results - Cat



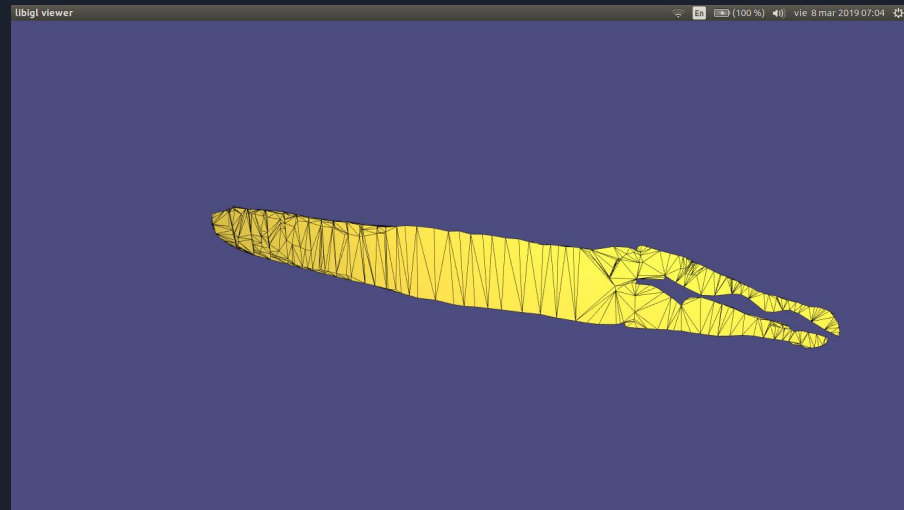
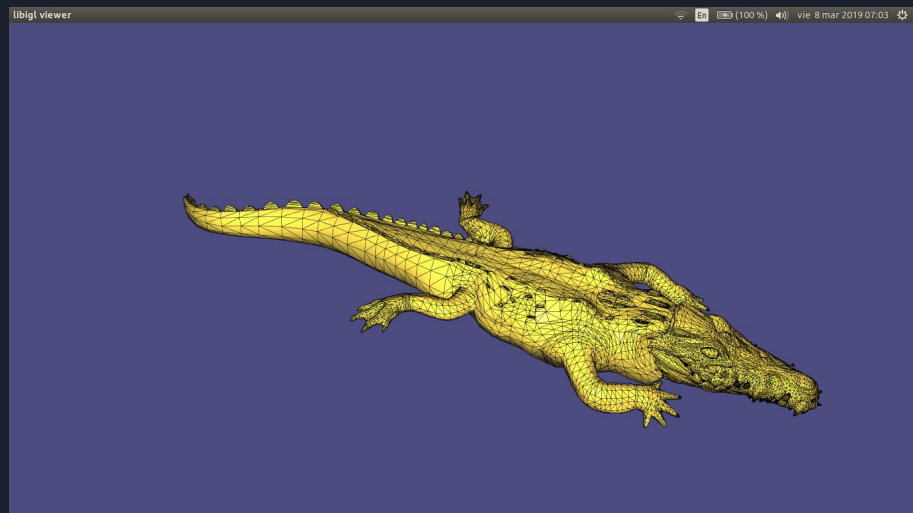
Results - Bear



Results - Girl



Results - Crocodile





Results

Mesh	Winding Numbers	Calculate Winding Numbers (ms)	Calculate Fast Winding Numbers (ms)
Big-Sigcat	164916	107014	3242
Bunny	34055	3794	310
Bear	56605	65600	14769
Crocodile	98719	29242	149
Girl	615313	603822	4954
Beast	192613	93014	5836



Thanks