Homotopy Type System with Strict Equality

Namdak Tonpa
Groupoid Infinity

Siegmentation Fault
Groupoid Infinity

Abstract

Here is presented a reincarnation of **cubicaltt** called **anders**.

2012 ACM Subject Classification Theory of computation → Lambda calculus

Keywords and phrases Homotopy Type System, Cubical Type Theory

1 Motivation

The **HTS**[1] language proposed by Voevodsky exposes two different presheaf models of type theory: the inner one is homotopy type system presheaf that models HoTT and the outer one is traditional Martin-Löf type system presheaf that models set theory with UIP. The motivation is to have an ability to express semisemplicial types. Theoretical work was continued in **2LTT**[2].

Our aim here is to preserve cubicaltt programs and remain language implementation compatible with original publications. While we are on our road to Lean-like tactic language, currently we are at the stage of regular cubical type checker with CHM-style primitives. You may try it at Github [3] or install through OPAM: **opam install anders**.

2 Syntax

The syntax resembles original syntax of the reference **CCHM** type checker cubicaltt, is slightly compatible with Lean syntax and contains the full set of Cubical Agda[4] primitives. Here is given the mathematical pseudo-code notation of the language expressions that come immediately after parsing. The core syntax definition of **HTS** language E corresponds to exp type defined in expr.ml OCaml module.

$$
\begin{aligned}
\text{cosmos} &:= \mathbf{U}_j \mid \mathbf{V}_k \\
\text{var} &:= \mathbf{var}\ \mathit{name} \mid \mathbf{hole} \\
\text{pi} &:= \Pi\ \mathit{name}\ E\ E \mid \lambda\ \mathit{name}\ E\ E \mid E\ E \\
\text{sigma} &:= \Sigma\ \mathit{name}\ E\ E \mid (E, E) \mid E.1 \mid E.2 \\
\text{id} &:= \mathbf{Id}\ E \mid \mathbf{ref}\ E \mid \mathbf{idJ}\ E \\
\text{path} &:= \mathbf{Path}\ E \mid E^i \mid E\ @\ E \\
\text{I} &:= \mathbf{I} \mid 0 \mid 1 \mid E \bigvee E \mid E \bigwedge E \mid \neg E \\
\text{part} &:= \mathbf{Partial}\ E\ E \mid [\ (E = I) \to E, \ldots\ ] \\
\text{sub} &:= \mathbf{inc}\ E \mid \mathbf{ouc}\ E \mid E\ [\ I\ \mapsto\ E\ ] \\
\text{kan} &:= \mathbf{transp}\ E\ E \mid \mathbf{hcomp}\ E \\
\text{glue} &:= \mathbf{Glue}\ E \mid \mathbf{glue}\ E \mid \mathbf{unglue}\ E\ E
\end{aligned}
$$

[1] https://www.math.ias.edu/vladimir/sites/math.ias.edu.vladimir/files/HTS.pdf
[2] https://arxiv.org/pdf/1705.03307.pdf
[3] https://github.com/groupoid/anders
[4] https://staff.math.su.se/anders.mortberg/papers/cubicalagda.pdf

Further Menhir BNF notation will be used to describe the top-level language E parser.

$$
\begin{aligned}
\text{E} :=\ & \textit{cosmos} \mid \textit{var} \mid \textit{MLTT} \mid \textit{CCHM} \mid \textit{HIT} \\
\text{HIT} :=\ & \textbf{inductive}\ E\ E \mid \textbf{ctor}\ \textit{name}\ E \mid \textbf{match}\ E\ E \\
\text{CCHM} :=\ & \textit{path} \mid I \mid \textit{part} \mid \textit{sub} \mid \textit{kan} \mid \textit{glue} \\
\text{MLTT} :=\ & \textit{pi} \mid \textit{sigma} \mid \textit{id}
\end{aligned}
$$

Keywords. The words of a top-level language (file or repl) consist of keywords or identifiers. The keywords are following: **module**, **where**, **import**, **option**, **def**, **axiom**, **postulate**, **theorem**, **(**, **)**, [,], $<$, $>$, /, **.1**, **.2**, Π, Σ, **,**, λ, $\times$, $\to$, **:**, **:=**, $\mapsto$, **U**, **V**, $\bigvee$, $\bigwedge$, **-**, $+$, **@**, **PathP**, **transp**, **hcomp**, **zero**, **one**, **Partial**, **inc**, **ouc**, **interval**, **inductive**, **Glue**, **glue**, **unglue**.

Indentifiers. Identifiers support UTF-8. Indentifiers couldn't start with :, $-$, $\to$. Sample identifiers: ¬-of-⋁, 1→1, is-?, =, $~]!005x, ∞, x→Nat.

Modules. Modules represent files with declarations. More accurate, BNF notation of module consists of imports, options and declarations.

menhir
start <Module.file> file
start <Module.command> repl
repl : COLON IDENT exp1 EOF | COLON IDENT EOF | exp0 EOF | EOF
file : MODULE IDENT WHERE line* EOF
path : IDENT
line : IMPORT path+ | OPTION IDENT IDENT | declarations

Imports. The import construction supports file folder structure (without file extensions) by using reserved symbol / for hierarchy walking.

Options. Each option holds bool value. Language supports following options: 1) girard (enables U : U); 2) pre-eval (normalization cache); 3) impredicative (infinite hierarchy with impredicativity rule); In Anders you can enable or disable language core types, adjust syntaxes or tune inner variables of the type checker. Here is the example how to setup minimal core able to prove internalization of **MLTT-73** variation (Path instead of Id and no inductive types, see base library): In order to turn HIT into ordinary CiC calculus you may say:

option HIT false
option CCHM false
option MLTT true

Declarations. Language supports following top level declarations: 1) axiom (non-computable declaration that breakes normalization); 2) postulate (alternative or inverted axiom that can preserve consistency); 3) definition (almost any explicit term or type in type theory); 4) lemma (helper in big game); 5) theorem (something valuable or complex enough).

axiom isProp (A : U) : U
def isSet (A : U) : U := Π (a b : A) (x y : Path A a b), Path (Path A a b) x y

Sample declarations. For example, signature isProp (A : U) of type U could be defined as normalization-blocking axiom without proof-term or by providing proof-term as definition.

In this example (A : U), (a b : A) and (x y : Path A a b) are called telescopes. Each telescope consists of a series of lenses or empty. Each lense provides a set of variables of the same type. Telescope defines parameters of a declaration. Types in a telescope, type of a declaration and a proof-terms are a language expressions exp1.

menhir

```
ident : IRREF | IDENT
vars : ident+
lense : LPARENS vars COLON exp1 RPARENS
telescope : lense telescope
params : telescope | []
declarations:
  | DEF IDENT params DEFEQ exp1
  | DEF IDENT params COLON exp1 DEFEQ exp1
  | AXIOM IDENT params COLON exp1
```

Expressions. All atomic language expressions are grouped by four categories: exp0 (pair constructions), exp1 (non neutral constructions), exp2 (path and pi applcations), exp3 (neutral constructions).

menhir

```
face : LPARENS IDENT IDENT IDENT RPARENS
partial : face+ ARROW exp1
exp0 : exp1 COMMA exp0 | exp1
exp1: LSQ separated(COMMA, partial) RSQ
  | LAM telescope COMMA exp1       | PI telescope COMMA exp1
  | SIGMA telescope COMMA exp1     | LSQ IRREF ARROW exp1 RSQ
  | LT vars GT exp1                | exp2 ARROW exp1
  | exp2 PROD exp1                 | exp2
```

The LR parsers demand to define exp1 as expressions that cannot be used (without a parens enclosure) as a right part of left-associative application for both Path and Pi lambdas. Universe indicies U_j (inner fibrant), V_k (outer pretypes) and S (outer strict omega) are using unicode subscript letters that are already processed in lexer.

menhir

```
exp2 : exp2 exp3 | exp2 APPFORMULA exp3 | exp3
exp3: LPARENS exp0 RPARENS LSQ exp0 MAP exp0 RSQ
  | HOLE            | PRE            | KAN             | IDJ exp3
  | exp3 FST        | exp3 SND       | NEGATE exp3     | INC exp3
  | exp3 AND exp3   | exp3 OR exp3   | ID exp3         | REF exp3
  | OUC exp3        | PATHP exp3     | PARTIAL exp3    | IDENT
  | IDENT LSQ exp0 MAP exp0 RSQ                      | HCOMP exp3
  | LPARENS exp0 RPARENS                             | TRANSP exp3 exp3
```

3 Semantics

The idea is to have a unified layered type checker, so you can disbale/enable any MLTT-style inference, assign types to universes and enable/disable hierachies. This will be done by providing linking API for pluggable presheaf modules. We selected 5 levels of type checker awareness from universes and pure type systems up to synthetic language of homotopy type theory. Each layer corresponds to its presheaves with separate configuration for universe hierarchies. We want to mention here with homage to its authors all categorical models of

$$
\begin{array}{lll}
\textbf{inductive}\ \mathrm{lang} : \mathbf{U} & := & \text{UNI: cosmos} \rightarrow \text{lang} \\
 & \mid & \text{PI: pure lang} \rightarrow \text{lang} \\
 & \mid & \text{SIGMA: total lang} \rightarrow \text{lang} \\
 & \mid & \text{ID: uip lang} \rightarrow lang \\
 & \mid & \text{PATH: homotopy lang} \rightarrow \text{lang} \\
 & \mid & \text{GLUE: gluening lang} \rightarrow \text{lang} \\
 & \mid & \text{HIT: hit lang} \rightarrow \text{lang}
\end{array}
$$

dependent type theory: Comprehension Categories (Grothendieck, Jacobs), LCCC (Seely), D-Categories and CwA (Cartmell), CwF (Dybjer), C-Systems (Voevodsky), Natural Models (Awodey). While we can build some transports between them, we leave this excercise for our mathematical components library. We will use here the Coquand's notation for Presheaf Type Theories in terms of restriction maps.

3.1 Universe Hierarchies

Language supports Agda-style hierarchy of universes: fibrant (U), interval pretypes (V) and strict omega with explicit level manipulation. All universes are bounded with preorder

$$Fibrant_j \prec Pretypes_k \prec Strict_l, \tag{1}$$

in which j, k, l are bounded with equation:

$$j < k < l. \tag{2}$$

Large elimination to upper universes is prohibited. This is extendable to Agda model:

$$
\begin{array}{lll}
\textbf{inductive}\ \mathrm{cosmos} : \mathbf{U} & := & \text{prop: nat} \rightarrow \text{cosmos} \\
 & \mid & \text{fibrant: nat} \rightarrow \text{cosmos} \\
 & \mid & \text{pretypes: nat} \rightarrow \text{cosmos} \\
 & \mid & \text{strict: nat} \rightarrow \text{cosmos} \\
 & \mid & \text{omega: cosmos} \\
 & \mid & \text{lock: cosmos}
\end{array}
$$

3.2 Dependent Types

▶ **Definition 1** (Type). *A type is interpreted as a presheaf A, a family of sets A_I with restriction maps $u \mapsto u\ f, A_I \to A_J$ for $f : J \to I$. A dependent type B on A is interpreted by a presheaf on category of elements of A: the objects are pairs (I, u) with $u : A_I$ and morphisms $f : (J, v) \to (I, u)$ are maps $f : J \to$ such that $v = u\ f$. A dependent type B is thus given by a family of sets $B(I, u)$ and restriction maps $B(I, u) \to B(J, u\ f)$.*

We think of A as a type and B as a family of presheves $B(x)$ varying $x : A$. The operation $\Pi(x : A)B(x)$ generalizes the semantics of implication in a Kripke model.

▶ **Definition 2** (Pi). *An element $w : [\Pi(x : A)B(x)](I)$ is a family of functions $w_f : \Pi(u : A(J))B(J, u)$ for $f : J \to I$ such that $(w_f u)g = w_{f\ g}(u\ g)$ when $u : A(J)$ and $g : K \to J$.*

$$
\begin{array}{lcl}
\textbf{inductive}\ \text{pure (lang : U)} : \mathbf{U} & := & \text{pi: name} \to \text{nat} \to \text{lang} \to \text{lang} \to \text{pure lang} \\
 & | & \text{lambda: name} \to \text{nat} \to \text{lang} \to \text{lang} \to \text{pure lang} \\
 & | & \text{app: lang} \to \text{lang} \to \text{pure lang}
\end{array}
$$

▶ **Definition 3** (Sigma). *The set $\Sigma(x : A)B(x)$ is the set of pairs (u, v) when $u : A(I), v : B(I, u)$ and restriction map $(u, v)\ f = (u\ f, v\ f)$.*

$$
\begin{array}{lcl}
\textbf{inductive}\ \text{total (lang : U)} : \mathbf{U} & := & \text{sigma: name} \to \text{lang} \to \text{total lang} \\
 & | & \text{pair: lang} \to \text{lang} \to \text{total lang} \\
 & | & \text{fst: lang} \to \text{total lang} \\
 & | & \text{snd: lang} \to \text{total lang}
\end{array}
$$

The presheaf with only Pi and Sigma is called **MLTT-72**. Its internalization in **anders** is as follows:

$$
\begin{array}{l}
\textbf{def}\ \text{MLTT-72 (A : U) : U} := \Sigma \\
\quad (\Pi\text{-form} : \Pi\ (\text{B} : \text{A} \to \text{U}),\ \text{U}) \\
\quad (\Pi\text{-ctor}_1 : \Pi\ (\text{B} : \text{A} \to \text{U}),\ \text{Pi A B} \to \text{Pi A B}) \\
\quad (\Pi\text{-elim}_1 : \Pi\ (\text{B} : \text{A} \to \text{U}),\ \text{Pi A B} \to \text{Pi A B}) \\
\quad (\Pi\text{-comp}_1 : \Pi\ (\text{B} : \text{A} \to \text{U})\ (\text{a} : \text{A})\ (\text{f} : \text{Pi A B}),\ \text{Equ (B a)}\ (\Pi\text{-elim}_1\ \text{B}\ (\Pi\text{-ctor}_1\ \text{B f) a) (f a)}) \\
\quad (\Pi\text{-comp}_2 : \Pi\ (\text{B} : \text{A} \to \text{U})\ (\text{a} : \text{A})\ (\text{f} : \text{Pi A B}),\ \text{Equ (Pi A B) f}\ (\lambda\ (\text{x} : \text{A}),\ \text{f x})) \\
\quad (\Sigma\text{-form} : \Pi\ (\text{B} : \text{A} \to \text{U}),\ \text{U}) \\
\quad (\Sigma\text{-ctor}_1 : \Pi\ (\text{B} : \text{A} \to \text{U})\ (\text{a} : \text{A})\ (\text{b} : \text{B a})\ ,\ \text{Sigma A B}) \\
\quad (\Sigma\text{-elim}_1 : \Pi\ (\text{B} : \text{A} \to \text{U})\ (\text{p} : \text{Sigma A B}),\ \text{A}) \\
\quad (\Sigma\text{-elim}_2 : \Pi\ (\text{B} : \text{A} \to \text{U})\ (\text{p} : \text{Sigma A B}),\ \text{B}\ (\text{pr}_1\ \text{A B p})) \\
\quad (\Sigma\text{-comp}_1 : \Pi\ (\text{B} : \text{A} \to \text{U})\ (\text{a} : \text{A})\ (\text{b: B a}),\ \text{Equ A a}\ (\Sigma\text{-elim}_1\ \text{B}\ (\Sigma\text{-ctor}_1\ \text{B a b}))) \\
\quad (\Sigma\text{-comp}_2 : \Pi\ (\text{B} : \text{A} \to \text{U})\ (\text{a} : \text{A})\ (\text{b: B a}),\ \text{Equ (B a) b}\ (\Sigma\text{-elim}_2\ \text{B (a, b)})) \\
\quad (\Sigma\text{-comp}_3 : \Pi\ (\text{B} : \text{A} \to \text{U})\ (\text{p} : \text{Sigma A B}),\ \text{Equ (Sigma A B) p}\ (\text{pr}_1\ \text{A B p},\ \text{pr}_2\ \text{A B p})),\ \text{U}
\end{array}
$$

3.3 Path Equality

3.4 Strict Equality

3.5 Glue Types

3.6 Higher Inductive Types