Data Science for Actuaries (ACT6100)

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Rappels # 4.5 (Linear Programming)

automne 2020

The beer problem: we want to produce beer, either blonde, or brown. We produce q_{blond} and q_{brown} barrels.



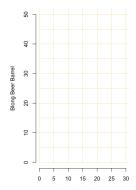






hops: 100kg barley: 280kg

We want to find $q_{brown}(x)$ and $q_{blond}(y)$ that is feaseable and maximize profit...



The beer problem: we want to produce beer, either blonde, or brown. We produce q_{blond} and q_{brown} barrels.

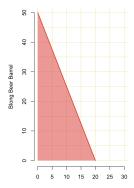
hops: 5kg barley: 10kg price: 40\$ | hops: 2kg barley: 14kg price: 30\$





hops: 100kg barley: 280kg

(1) Our hops stock is 100kg, it takes 5kg per barrel of blonde it takes 2kg per barrel of brown so 5x + 2y < 100



The beer problem: we want to produce beer, either blonde, or brown. We produce q_{blond} and q_{brown} barrels.

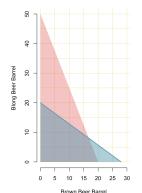
hops: 5kg barley: 10kg barley: 14kg price: 40\$





hops: 100kg barley: 280kg

(1) Our barley stock is 280kg, it takes 10kg per barrel of blonde it takes 14kg per barrel of brown so $10x + 14y \le 280$



The beer problem: we want to produce beer, either blonde, or brown. We produce q_{blond} and q_{brown} barrels.





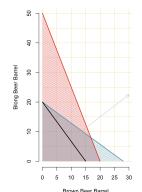




hops: 100kg barley: 280kg

Under the two feasibility conditions,

(2) we want to maximize our profit $\max\{40x + 30y\}$



The beer problem: we want to produce beer, either blonde, or brown. We produce q_{blond} and q_{brown} barrels.



hops: 5kg barley: 10kg price: 40\$ | hops: 2kg barley: 14kg price: 30\$

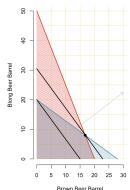




hops: 100kg barley: 280kg

Our problem is

$$\max \{40x + 30y\}$$
s.t. $10x + 14y \le 280$
 $5x + 2y \le 100$
 $x, y \ge 0$



Our problem is here

$$\max \{ \boldsymbol{c}^{\top} \boldsymbol{x} \}$$
 s.t. $\boldsymbol{A} \boldsymbol{x} \leq \boldsymbol{b}$ $\boldsymbol{x} \geq 0$

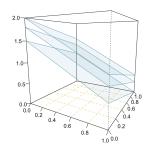
e.g.

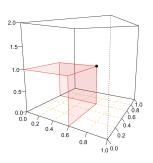
$$\max \{ax + by + cz\}$$

subject to

$$\begin{cases} x \le \alpha \\ y \le \beta \\ z \le \gamma \end{cases}$$

The red volume is the set of feasible points (x, y, z)





Our problem is here

$$\max \{ \boldsymbol{c}^{\top} \boldsymbol{x} \}$$

s.t. $\boldsymbol{A} \boldsymbol{x} \leq \boldsymbol{b}$
 $\boldsymbol{x} \geq 0$

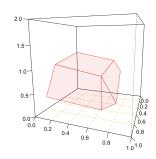
e.g., more generally

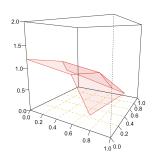
$$\max \left\{ ax + by + cz \right\}$$

subject to

$$\begin{cases} \alpha_1 x + \beta_1 y + \gamma_1 z \leq \delta_1 \\ \alpha_2 x + \beta_2 y + \gamma_2 z \leq \delta_2 \\ \vdots \\ \alpha_k x + \beta_k y + \gamma_k z \leq \delta_k \end{cases}$$

which is a (convex) polyhedron.





Simplex

Our problem is here

$$\max \{ \boldsymbol{c}^{\top} \boldsymbol{x} \}$$
 s.t. $\boldsymbol{A} \boldsymbol{x} \leq \boldsymbol{b}$ $\boldsymbol{x} > 0$

First step: enlarge the parameter space, $10x_1 + 14x_2 \le 280$ becomes $10x_1 + 14x_2 - u_1 = 280$ (so called slack variables)

$$\max \{40x_1 + 30x_2\}$$
s.t.
$$10x_1 + 14x_2 + u_1 = 280$$

$$2x_1 + 5x_2 + u_2 = 100$$

$$x_1, x_2, u_1, u_2 \ge 0$$

which is a problem of the general form

$$\max \{ \boldsymbol{c}^{\top} \boldsymbol{x} \}$$

s.t. $\boldsymbol{A} \boldsymbol{x} = \boldsymbol{b}$
 $\boldsymbol{x} \ge 0$

Simplex Method

A linear program in standard form can be written

$$\max \{ \boldsymbol{c}^{\top} \boldsymbol{x} \}$$

s.t. $\boldsymbol{A} \boldsymbol{x} = \boldsymbol{b}$
 $\boldsymbol{x} \ge 0$

and it can be represented as a tableau (matrix) of the form

$$\begin{bmatrix} \mathbf{1} & -\boldsymbol{c}^\top & \mathbf{0} \\ \mathbf{0} & \boldsymbol{A} & \boldsymbol{b} \end{bmatrix}$$

Use some linear algebra, one can solve the optimization problem

Application to median computation

Given $\mathbf{y} = \{y_1, \dots, y_n\}$, the median is a solution to

$$\min_{\mu} \left\{ \sum_{i=1}^{n} |y_i - \mu| \right\}.$$

Equivalently, we want to solve

$$\min_{\mu,\mathbf{a},\mathbf{b}} \left\{ \sum_{i=1}^n a_i + b_i \right\}$$

with $a_i, b_i > 0$ and $y_i - \mu = a_i - b_i, \forall i = 1, \dots, n$.

Heuristically, the idea is to write $y_i = \mu + \varepsilon_i$, and then define a_i 's and b_i 's so that $\varepsilon_i = a_i - b_i$ and $|\varepsilon_i| = a_i + b_i$, i.e.

$$a_i = (\varepsilon_i)_+ = \max\{0, \varepsilon_i\} = |\varepsilon| \cdot \mathbf{1}_{\varepsilon_i > 0}$$

and

$$b_i = (-\varepsilon_i)_+ = \max\{0, -\varepsilon_i\} = |\varepsilon| \cdot \mathbf{1}_{\varepsilon_i < 0}$$

Application to median computation

Thus, set $\mathbf{z} = (\mu^+; \mu^-; \mathbf{a}, \mathbf{b})^\top \in \mathbb{R}^{2n+2}_+$, and then write the constraint as Az = b with b = y and $A = \begin{bmatrix} \mathbf{1}_n, -\mathbf{1}_n, \mathbb{I}_n, -\mathbb{I}_n \end{bmatrix}$ For the objective function $\boldsymbol{c} = (\boldsymbol{0}, \boldsymbol{1}_n, -\boldsymbol{1}_n)^{\top} \in \mathbb{R}^{2n+2}_+$ and our program is min $\{c^{\top}z\}$ s.t. $Az = b, z \ge 0$.

```
1 > n = 101
                     _{1} > X = rep(1,n)
2 > set.seed(1)
3 > y = rlnorm(n)
2 > A = cbind(X, -X, diag(n), -diag(n))
3 > b = y
4 > median(y) 4 > c = c(rep(0,2), rep(1,n), rep(1,n))
5 [1] 1.077415 5 > r = lp("min", c, A, rep("=", n), b)
                      6 > head(r$solution,1)
7 > library(lpSolve) 7 [1] 1.077415
```

since the median is a solution to

$$\min_{\mu} \left\{ \sum_{i=1}^{n} |y_i - \mu| \right\} = \min_{\mu} \left\{ \sum_{i=1}^{n} \max\{(y_i - \mu), -(y_i - \mu) \right\}$$

Application to quantile computation

More generally, if the quantile of order τ is a solution of $\tau \in (0,1)$,

$$\min_{q} \left\{ \sum_{i=1}^{n} \max\{\tau(y_i - \mu), (1 - \tau)(y_i - \mu) \right\}$$

The linear program is now

$$\min_{q^+,q^-,\mathbf{a},\mathbf{b}} \left\{ \sum_{i=1}^n \tau a_i + (1-\tau)b_i \right\}$$

with $a_i, b_i, q^+, q^- \ge 0$ and $y_i = q^+ - q^- + a_i - b_i$, $\forall i = 1, \dots, n$.

Application to quantile regression

In a regression, we use $\mathbf{x}_i^{\top} \boldsymbol{\beta}$ instead of $\boldsymbol{\mu}$. The linear program is

$$\min_{\boldsymbol{\beta}^+, \boldsymbol{\beta}^-, \mathbf{a}, \mathbf{b}} \left\{ \sum_{i=1}^n \tau a_i + (1-\tau)b_i \right\}$$

with $a_i, b_i \geq 0$ and $y_i = \mathbf{x}^{\top} [\boldsymbol{\beta}^+ - \boldsymbol{\beta}^-] + a_i - b_i$, $\forall i = 1, \dots, n$ and $\beta_j^+, \beta_j^- \geq 0 \ \forall j = 0, \dots, k$.

```
1 > n=nrow(Davis)
2 > X = cbind( 1, Davis$height)
3 > y =Davis$weight
4 > K = ncol(X)
5 > N = nrow(X)
6 > A = cbind(X,-X,diag(N),-diag(N))
7 > c = c(rep(0,2*ncol(X)),tau*rep(1,N),(1-tau)*rep(1,N))
8 > b = y
9 > r = lp("min",c,A,rep("=",N),b)
10 > beta = r$sol[1:K] - r$sol[(1:K+K)]
11 > beta
12 [1] -110 1
```

Application to quantile regression

See

to compare with