

Sorting

Bubble Sort

- Maximum number of passes: number of elements - 1
- For each pass, compare adjacent elements, exchange their places if out of place
- Rear of array is sorted first
 - Largest element put to last place
 - Second largest element put to second last place
 - ...
 - Smallest element put to first place
- Process:
 - Pass 0:
 - ◆ Compares adjacent elements ($A[0], A[1]$), ($A[1], A[2]$) ... ($A[n-2], A[n-1]$)
 - ◆ For each pair of ($A[j], A[j+1]$):
 - ◆ Exchange their values if $A[j] > A[j+1]$
 - ◆ Set `lastExchangeIndex = j`
 - ◆ Largest element is $A[n-1]$
 - ◆ Front of array ($A[0] - A[\text{lastExchangeIndex}]$) is unordered
 - ◆ Rear of array ($A[\text{lastExchangeIndex}] - A[n-1]$) is ordered
 - Subsequent passes:
 - ◆ Compare adjacent terms in sublist of ($A[0] - A[\text{lastExchangeIndex}]$)
 - Terminates when `lastExchangeIndex = 0`
- Code:

```
def bubbleSort(array):  
    n = len( array )  
    isSorted = False                                # create flag  
  
    while not isSorted and n > 0:  
        isSorted = True                            # assume sorted  
  
        for i in range( n - 1 ):  
            if array[i] > array[i + 1]:            # wrong order  
                array = swap( array, i, i+1 )  
                isSorted = False
```

```

        n -= 1
        place after each pass
    return arr

```

- Benefits:
 - Stable
 - Memory efficient
- Drawback: Inefficient

Insertion Sort

- Process:
 - Pass 0:
 - ◆ First element stays at position 0
 - ◆ Compare second element A[1] with first A[0]
 - ◆ Swap A[0] and A[1] if A[0] > A[1]
 - Subsequent passes:
 - ◆ For every target element A[i], compare element down the list of elements (A[i-1], A[i-2] ... A[0])
 - ◆ Stop comparison at first element A[j] if (A[j] <= A[i]), or beginning of array (A[0]) is reached
 - ◆ Shift every element to the right after comparing (A[j] = A[j-1])
 - ◆ Insert target element at correct position (j) after sliding other elements
 - ◆ Sublist (A[0] - A[j]) is ordered
- Code:

```

def insertionSort(array):
    for i in range( 1, len(array) ):
        target = array[i]
        j = i

        while j > 0 and target < array[j-1]:
            array = swap( array, j, j-1 )
            j -= 1

        array[j] = target

```

scans down list

locate insertion point

free up space to insert

insertion

return array

- Benefits:
 - Efficient for small sets of data
 - Easily implemented
- Drawbacks: Inefficient on large lists and arrays
- Sample:

Array Index	A[0]	A[1]	A[2]	A[3]	A[4]
Original	50	20	40	75	35
PASS 0	20	50	40	75	35
PASS 1	20	40	50	75	35
PASS 2	20	40	50	75	35
PASS 3	20	35	40	50	75

Quick Sort

- Fastest sorting algorithm
- Uses partition approach to sort array
- Process:
 - Array is sorted
- Simple quick sort code:

```
def quickSort(array):  
    if len( array ) <= 1:                # list of 0 or 1 element is  
already sorted                          # already sorted  
        return array  
  
    else:  
        pivot_value = array.pop(0)      # select & remove pivot  
value (any index value)  
        less = []  
        greater = []  
  
        for item in array:              # append each item into  
appropriate array
```

```

        if item < pivot_value:
            less.append( item )
        else:
            greater.append( item )

    return quickSort( less ) + [pivot_value] +
quickSort( greater )

# note: [pivot_value] is a
list

```

- Hoare's partition code:

```

def quickSort(A, low, high):
    if low < high:                                # array has more than
one element
        pos = partition( A, low, high ) # splits array into two
sublists, pos is final position of pivot

        quickSort( A, low, pos-1 )    # quick sort lower sublist
        quickSort( A, pos+1, high )    # quick sort higher sublist

def partition(A, low, high):
    pivot = A[low]                                # set pivot value
    left, right = low, high - 1

    while True:                                    # infinite loop
        while A[left] < pivot:                # increment left pointer
            left = left + 1

        while A[right] > pivot:                # decrement right pointer
            right = right - 1

        if left < right:                        # swap pointers
            swap( A, left, right )
        else:                                    # left and right pointer
meets
            return right

```

- Benefit: Efficient for any array
- Drawback: Unstable

Selection Sort (Not In Syllabus)

- Total number of passes: number of elements - 1
- For each pass, find smallest element to exchange with first element in selected group
- Ignore first element of selected group after every pass to get next selected group
- Process:
 - Assume n elements in array A , index starts from 0
 - Pass 0:
 - ◆ Select smallest element, exchange with $A[0]$, placing it at beginning of array A
 - ◆ Front of array ($A[0]$) is ordered
 - ◆ Rear of array ($A[1] - A[n-1]$) is unordered
 - Pass 1:
 - ◆ Select smallest element from $A[1] - A[n-1]$, exchange with $A[1]$, placing it at front of sublist ($A[1] - A[n-1]$)
 - ◆ Front of array ($A[0] - A[1]$) is ordered
 - ◆ Rear of array ($A[2] - A[n-1]$) is unordered
 - Pass 2:
 - ◆ Select smallest element from $A[2] - A[n-1]$, exchange with $A[2]$, placing it at front of sublist ($A[2] - A[n-1]$)
 - ◆ Front of array ($A[0] - A[2]$) is ordered
 - ◆ Rear of array ($A[3] - A[n-1]$) is unordered
 - ...

- Code:

```
def selectionSort ( A, n ):      # array A[0...n-1] of n integers
    for i in range( 0, n ):
        startIndex = i          # start of sub-array
        minIndex = i            # check for smallest value in sub-
array
        for j in range( i+1, n ):
            if A[j] < A[minIndex]:
                minIndex = j
        swap(A, minIndex, startIndex)
```

- Benefit: Simple
- Drawback: Inefficient

Merge Sort (Not in Syllabus)

- **Merging:**

- Combines 2 arrays that are already sorted
- Outputs 3rd arrays that is sorted
- Process:ku
 - ◆ Assume sorted arrays array A and array B
 - ◆ Initialise empty array C
 - ◆ Read first element x from A
 - ◆ Read first element y from B
 - ◆ If $x < y$:
 - ◆ Write x into C
 - ◆ Read new x from A
 - ◆ Else ($y < x$):
 - ◆ Write y into C
 - ◆ Read new y from B
 - ◆ If end of A is reached:
 - ◆ Copy all remaining elements from B into C
 - ◆ If end of B is reached:
 - ◆ Copy all remaining elements from A into C
- Code:

```

def mergeSort (A, B):                                     # Assume sorted arrays
A, B
    res = [ ]                                             # Resultant array
    i = 0
    j = 0

    while ( i < len(A) ) and ( j < len(B) ):              # loop when end
of array A and B are not reached
        if A[i] < B[j]:
            res.append ( A[i] )
            i = i + 1
        else:
            res.append ( B[j] )
            j = j + 1
  
```

```

                                while i < len(A):                # end of array B
reached
                                res.append ( A[i] )              # append
remaining A into resultant
                                i = i + 1

```

```

                                while j < len(B):                # end of array A
reached
                                res.append ( B[j] )              # append
remaining B into resultant
                                j = j + 1

```

● **Straight Merge Sort:**

- Process:
 - ◆ Start with array A = [6, 5, 3, 1, 8, 7, 2, 4]
 - ◆ Separate all elements into single-element subarrays A[0] - A[n]:
 - ◆ First pass: [6], [5], [3], [1], [8], [7], [2], [4]
 - ◆ Compare and sort adjacent subarrays, write into resultant array:
 - ◆ Second pass: [5, 6], [1, 3], [7, 8], [2, 4]
 - ◆ Third pass: [1, 3, 5, 6], [2, 4, 7, 8]
 - ◆ Fourth pass: [1, 2, 3, 4, 5, 6, 7, 8]
- Efficiency increases as length of subarrays increases