

Sorting

Bubble Sort

- Maximum number of passes: number of elements - 1
- For each pass, compare adjacent elements, exchange their places if out of place
- Rear of array is sorted first
 - Largest element put to last place
 - Second largest element put to second last place
 - ...
 - Smallest element put to first place
- Process:
 - Pass 0:
 - ◆ Compares adjacent elements (A[0], A[1]), (A[1], A[2]) ... (A[n-2], A[n-1])
 - ◆ For each pair of (A[j], A[j+1]):
 - ◆ Exchange their values if $A[j] > A[j+1]$
 - ◆ Set lastExchangeIndex = j
 - ◆ Largest element is A[n-1]
 - ◆ Front of array (A[0] - A[lastExchangeIndex]) is unordered
 - ◆ Rear of array (A[lastExchangeIndex] - A[n-1]) is ordered
 - Subsequent passes:
 - ◆ Compare adjacent terms in sublist of (A[0] - A[lastExchangeIndex])
 - Terminates when lastExchangeIndex = 0
- Code:

```
def bubbleSort(array):
```

```
    n = len(array)
    isSorted = False
```

```
    while not isSorted and n > 0:
```

```
        # assume sorted
        isSorted = True
```

```
        for i in range(n - 1):
```

```
            # wrong order
```

```
            if array[i] > array[i + 1]:
                array = swap(array, i, i+1)
```

isSorted = False

n -= 1
return arr

- Benefits:
 - Stable
 - Memory efficient
- Drawback: Inefficient

Insertion Sort

- Process:
 - Pass 0:
 - ◆ First element stays at position 0
 - ◆ Compare second element A[1] with first A[0]
 - ◆ Swap A[0] and A[1] if A[0] > A[1]
 - Subsequent passes:
 - ◆ For every target element A[i], compare element down the list of elements (A[i-1], A[i-2] ... A[0])
 - ◆ Stop comparison at first element A[j] if (A[j] <= A[i]), or beginning of array (A[0]) is reached
 - ◆ Shift every element to the right after comparing (A[j] = A[j-1])
 - ◆ Insert target element at correct position (j) after sliding other elements
 - ◆ Sublist (A[0] - A[j]) is ordered
- Code:

```
def insertionSort(array):  
    for i in range(1, len(array)):  
        target = array[i]  
        j = i  
  
        # locate insertion point  
        while j > 0 and array[j] < array[j-1]:  
  
            # free up space to insert  
            array = swap(array, j, j-1)  
            j -= 1
```

```
array[j] = target
```

```
return array
```

- Benefits:
 - Efficient for small sets of data
 - Easily implemented
- Drawbacks: Inefficient on large lists and arrays
- Example:

| Array Index | A[0] | A[1] | A[2] | A[3] | A[4] |
|-------------|------|------|------|------|------|
| Original | 50 | 20 | 40 | 75 | 35 |
| PASS 0 | 20 | 50 | 40 | 75 | 35 |
| PASS 1 | 20 | 40 | 50 | 75 | 35 |
| PASS 2 | 20 | 40 | 50 | 75 | 35 |
| PASS 3 | 20 | 35 | 40 | 50 | 75 |

Quick Sort

- Fastest sorting algorithm
- Uses partition approach to sort array
- Process:
 - Array is sorted
- Simple quick sort code:

```
def quickSort(array):  
  
    # list of 0 or 1 element is already sorted  
    if len(array) <= 1:  
        return array  
  
    else:  
        # select & remove pivot value (any index  
value)  
        pivot_value = array.pop(0)  
        less = []
```

```

        greater = []

        # append each item into appropriate
array
        for item in array:
            if item < pivot_value:
                less.append(item)
            else:
                greater.append(item)

        return quickSort(less) + [pivot_value] +
quickSort(greater)
        # note: [pivot_value] is a list

```

- Hoare's partition code (*complicated*):

```

def quickSort(A, low, high):

    # array has more than one element
    if low < high:
        pos = partition(A, low, high)
        # splits array into two sublists, pos is
final position of pivot

        # quick sort lower sublist
        quickSort(A, low, pos-1)

        # quick sort higher sublist
        quickSort(A, pos+1, high)

def partition(A, low, high):
    pivot = A[low]
    left, right = low, high -1

    # infinite loop
    while True:

        # increment left pointer
        while A[left] < pivot:

```

```

    left = left + 1

    # decrement right pointer
    while A[right] > pivot:
        right = right - 1

    # swap pointers
    if left < right:
        swap(A, left, right)

    # left and right pointer meets
    else:
        return right

```

- Benefit: Efficient for any array
- Drawback: Unstable

Selection Sort (Not In Syllabus)

- Total number of passes: number of elements - 1
- For each pass, find smallest element to exchange with first element in selected group
- Ignore first element of selected group after every pass to get next selected group
- Code:

```

def selectionSort (A, n):
    for i in range(0, n):
        startIndex = i
        minIndex = i
        for j in range(i+1, n):
            if A[j] < A[minIndex]:
                minIndex = j
        swap(A, minIndex, startIndex)

```

- Benefit: Simple
- Drawback: Inefficient

Merge Sort (Not in Syllabus)

- **Merging:**

- Combines 2 arrays that are already sorted
- Outputs 3rd array that is sorted
- Process:
 - ◆ Assume sorted arrays array A and array B
 - ◆ Initialise empty array C
 - ◆ Read first element x from A
 - ◆ Read first element y from B
 - ◆ If $x < y$:
 - ◆ Write x into C
 - ◆ Read new x from A
 - ◆ Else ($y < x$):
 - ◆ Write y into C
 - ◆ Read new y from B
 - ◆ If end of A is reached:
 - ◆ Copy all remaining elements from B into C
 - ◆ If end of B is reached:
 - ◆ Copy all remaining elements from A into C
- Code:

```
# Assume sorted arrays A, B
def mergeSort (A, B):
    res = [ ]
    i = 0
    j = 0

    # loop when end of array A and B are not
reached
    while (i < len(A)) and (j < len(B)):
        if A[i] < B[j]:
            res.append (A[i])
            i = i + 1

        else:
            res.append (B[j])
            j = j +1

    # end of array B reached
    while i < len(A):
```

```

resultant          # append remaining A into
                   res.append (A[i])
                   i = i + 1

                   # end of array A reached
while j < len(B):
    resultant      # append remaining B into
                   res.append (B[j])
                   j = j + 1

```

● **Straight Merge Sort:**

- Process:
 - ◆ Start with array A = [6, 5, 3, 1, 8, 7, 2, 4]
 - ◆ Separate all elements into single-element subarrays A[0] - A[n]:
 - ◆ First pass: [6], [5], [3], [1], [8], [7], [2], [4]
 - ◆ Compare and sort adjacent subarrays, write into resultant array:
 - ◆ Second pass: [5, 6], [1, 3], [7, 8], [2, 4]
 - ◆ Third pass: [1, 3, 5, 6], [2, 4, 7, 8]
 - ◆ Fourth pass: [1, 2, 3, 4, 5, 6, 7, 8]
- Efficiency increases as length of subarrays increases