

# Searching

## Sequential / Linear Search

- Search from beginning to the end
- Compare every element in list to the required item
- Match --> found; no match --> go to next element
- Advantages:
  - Elements can be in any order
- Challenges:
  - Require large amount of memory
  - May cause stack overflow and/or memory problems
- Code:

```
array = [ ... ]  
target = input("enter item to be searched")
```

```
for i in range( len(array) ):  
    if array[i] is target:  
        return i
```

## Binary Search

- Search in order e.g. ascending, descending
- Steps:
  1. Arrange list of elements in order
  2. Compare key with mid-index value of list
  3. Determine which side the key is on the list
  4. E.g. key < mid-index element:
    - ◆ Change high-index to mid-1
    - ◆ Compare key with new mid
  5. E.g. key > mid-index element:
    - ◆ Change low-index to mid+1
    - ◆ Compare key with new mid
  6. Repeat until found or not found, low-index > high-index
- Advantages:
  1. Faster search by eliminating half the elements at once
  2. Uses less memory
- Challenges:
  1. Elements must be arranged in specific order: ascending or

descending

- Iterative code:

```
target = input("enter item to be searched")
array = [...] # array of
sorted elements
```

```
def binarySearch(array, target):
    low, high = 0, len(array) - 1
```

```
    while low < high:
        mid = (low + high) // 2
        pivot = array[mid]
```

```
        if pivot == target:
            return mid
        elif target < pivot:
            low = mid + 1
        else:
            high = mid - 1
```

- Recursive code:

```
target = input("enter item to be searched")
array = [...] # array of
sorted elements
n = len(array) - 1
```

```
if len(array) <= 1: # array is
already sorted
    return
else:
    rec_binary_search(target, array, 0, n)
```

```
def rec_binary_search(target, array, low, high):
    if low < high:
        return None # target not
found
```

```
    mid = (low + high) // 2
```

```
pivot = array[mid]
```

```
if pivot == target:
```

```
    return mid
```

```
elif pivot < target:
```

```
    rec_binary_search(target, array, low, mid - 1)
```

```
else:
```

```
    rec_binary_search(target, array, mid + 1, high)
```

## Hash Table Search

- **Hash Function:**

- Location of an item is determined directly as a function of the time itself rather than by a sequence of trial-and-error comparisons
- Only one location is required to be examined
- Time required to locate the item is constant and doesn't depend on the number of items stored

- **Collision Strategies:**

- Collision: Many values may have the same hash value and is tried to be stored at the same location
- **Linear Probing:**
  - ◆ Linear search of the table begins at the location where collision occurs and continues until an empty slot is found in which the item can be stored
  - ◆ Determining if an element is in the hash table:
    - ◆ Apply hash function to compute the position of the value
    - ◆ Three cases to consider:
      1. If location is empty, the value is not in the table
      2. If location contains the specified value, search is immediately successful
      3. If location contains a value other than the specified value, begin a "circular" linear search until the item is found or reach empty location or the starting location, indicating item isn't in the table
- **Chaining:**
  - ◆ All elements that are stored at the same location are chained together
  - ◆ Advantages:

- ◆ Fast searching
- ◆ Challenges:
  - ◆ Collisions occur, causing some elements to occupy locations reserved for other hash values
  - ◆ Hash table may not have enough space to store all elements