# Searching

## **Sequential / Linear Search**

- Search from beginning to the end
- Compare every element in list to the required item
- Match --> found; no match --> go to next element
- Advantages:
  - Elements can be in any order
- Challenges:
  - Require large amount of memory
  - May cause stack overflow and/or memory problems
- Code:

```
array = [ ... ]
target = input("enter item to be searched")
for i in range( len(array) ):
    if array[i] is target:
        return i
```

## **Binary Search**

- Search in order e.g. ascending, descending
- Steps:
  - 1. Arrange list of elements in order
  - 2. Compare key with mid-index value of list
  - 3. Determine which side the key is on the list
  - 4. E.g. key < mid-index element:
    - ◆ Change high-index to mid-1
    - Compare key with new mid
  - 5. E.g. key > mid-index element:
    - ◆ Change low-index to mid+1
    - Compare key with new mid
  - 6. Repeat until found or not found, low-index > high-index
- Advantages:
  - 1. Faster search by eliminating half the elements at once
  - 2. Uses less memory
- Challenges:
  - 1. Elements must be arranged in specific order: ascending or

descending

#### • Iterative code:

target = input("enter item to be searched")

array = [...] # array of

sorted elements

def binary\_search(array, target):
 low, high = 0, len(array) - 1

while low <= high:

mid = (low + high) // 2 # get pivot

index

pivot = array[mid] # assign

pivot value

if pivot == target: # target

found

return mid

elif target < mid\_value: # target in

lower subarray

low = mid + 1

else: # target in

higher subarray

high = mid - 1

if low > high: # pointers

crossed -- target is not found

return -1 # return

dummy value

#### Recursive code:

target = input("enter item to be searched")

array = [...] # array of

sorted elements

n = len(array) - 1

if len(array) <= 1: # array is

```
already sorted
         return array
    else:
         rec binary search(target, array, 0, n)
    def rec_binary_search(target, array, low, high):
         if low < high:
             return None
                                                      # target not
found
         mid = (low + high) // 2
                                                      # get pivot
index
         pivot = array[mid]
                                                      # assign
pivot value
         if pivot == target:
                                                      # target
found
             return mid
         elif pivot < target:
                                                      # target in
lower subarray
             rec_binary_search(target, array, low, mid - 1)
         else:
                                                      # target in
higher subarray
             rec_binary_search(target, array, mid + 1, high)
```

#### **Hash Table Search**

#### Hash Function:

- Location of an item is determined directly as a function of the time itself rather than by a sequence of trial-and-error comparisons
- Only one location is required to be examined
- Time required to locate the item is constant and doesn't depend on the number of items stored

## • Collision Strategies:

- Collision: Many values may have the same hash value and is tried to be stored at the same location
- Linear Probing:
  - Linear search of the table begins at the location where

collision occurs and continues until an empty slot is found in which the item can be stored

- Determining is an element is in the hash table:
  - Apply hash function to compute the position of the value
  - Three cases to consider:
    - 1. If location is empty, the value is not in the table
    - 2. If location contains the specified value, search is immediately successful
    - 3. If location contains a value other than the specified value, begin a "circular" linear search until the item is found or reach empty location or the starting location, indicating item isn't in the table

### Chaining:

- All elements that are stored at the same location are chained together
- Advantages:
  - Fast searching
- Challenges:
  - Collisions occur, causing some elements to occupy locations reserved for other hash values
  - Hash table may not have enough space to store all elements