1 Syntax

1.1 Source Syntax

$$
\begin{array}{llll}
\text{Types} & T & ::= & \alpha \mid \top \mid \tau_1 \to \tau_2 \mid \forall \alpha.\, \tau \mid \tau_1 \,\&\, \tau_2 \mid \{l{:}\tau\} \\
\text{Expressions} & E & ::= & x \mid \top \mid \lambda(x{:}\tau).\, e \mid e_1\ e_2 \mid \Lambda\alpha.\, e \mid e\ \tau \mid e_1 ,, e_2 \mid \{l = e\} \mid e.l \mid e \setminus l \\
 & & \mid & \texttt{sig}\ s[\overline{\alpha}]\ \texttt{where}\ \overline{l : \tau}\ \texttt{in}\ e \\
 & & \mid & \texttt{sig}\ s_1[\overline{\alpha_1}]\ \texttt{extends}\ \overline{s_2[\overline{\alpha}]}\ \texttt{where}\ \overline{l : \tau}\ \texttt{in}\ e \\
 & & \mid & \texttt{algebra}\ x\ \texttt{implements}\ s[\overline{\tau}]\ \texttt{where}\ \overline{l@(l_1\ \overline{x_1}) = e_1}\ \texttt{in}\ e \\
 & & \mid & \texttt{algebra}\ x\ \texttt{extends}\ \overline{x_0}\ \texttt{implements}\ s[\overline{\tau}]\ \texttt{where}\ \overline{l@(l_1\ \overline{x_1}) = e_1}\ \texttt{in}\ e \\
 & & \mid & \texttt{data}\ d\ \texttt{from}\ s[\overline{\alpha_0}].\alpha_1\ \texttt{in}\ e \\
 & & \mid & \texttt{let}\ x\ (\overline{x_1} : \overline{\tau_1})\ (\overline{x_2} : d[\overline{\tau}]) : d[\overline{\tau}] = e_1\ \texttt{in}\ e \\
 & & \mid & e[\overline{\tau}]\texttt{<}\overline{x}\texttt{>} \\
\text{Contexts} & \Gamma & ::= & \epsilon \mid \Gamma, \alpha \mid \Gamma, x{:}\tau \\
 & & \mid & \Gamma, s[\overline{\alpha}] \to \overline{l : \tau} \\
 & & \mid & \Gamma, x \multimap \overline{s[\overline{\tau}]} \\
 & & \mid & \Gamma, d \rightsquigarrow s[\overline{\alpha_0}].\alpha_1 : \tau \\
\text{Labels} & l & & \text{(fields)} \\
 & s & & \text{(interfaces)} \\
 & d & & \text{(datatypes)} \\
\text{Syntactic sugars} & \circ & ::= & s[\overline{\alpha_0}] \\
 & \bullet & ::= & [\overline{\alpha_0}/\overline{\alpha}]\Gamma(s) \\
 & \circ & ::= & d(\overline{\tau_0}) \\
 & \bullet & ::= & [\overline{\tau_0}/(\overline{\alpha_0}\setminus\alpha_1)]\Gamma(d)
\end{array}
$$

1.2 Target Syntax

$$
\begin{array}{llll}
\text{Types} & T & ::= & \alpha \mid \top \mid \tau_1 \to \tau_2 \mid \forall \alpha.\, \tau \mid \tau_1 \,\&\, \tau_2 \mid \{l{:}\tau\} \\
\text{Expressions} & E & ::= & x \mid \top \mid \lambda(x{:}\tau).\, e \mid e_1\ e_2 \mid \Lambda\alpha.\, e \mid e\ \tau \mid e_1 ,, e_2 \mid \{l = e\} \mid e.l \mid e \setminus l \\
\text{Contexts} & \Gamma & ::= & \epsilon \mid \Gamma, \alpha \mid \Gamma, x{:}\tau \\
\text{Labels} & l & & \\
\text{Syntactic sugars} & \circ & ::= & \texttt{let}\ x : \tau = e_1\ \texttt{in}\ e_2 \\
 & \bullet & ::= & (\lambda(x{:}\tau).\, e_2)\ e_1
\end{array}
$$

2 Translation Rules

$$
\boxed{\Gamma \vdash e : \tau \Rightarrow E}
\qquad
\frac{\Gamma, s[\overline{\alpha}] \to \overline{l : \tau} \ \vdash\ e \ :\ \tau_* \ \Rightarrow\ E}{\Gamma \ \vdash\ \texttt{sig}\ s[\overline{\alpha}]\ \texttt{where}\ \overline{l : \tau}\ \texttt{in}\ e \ :\ \tau_* \ \Rightarrow\ \texttt{let}\ merge_s : \ldots = \ldots\ \texttt{in}\ E}
$$

$$
\frac{\overline{\Gamma \vdash s_2[\overline{\alpha_2}]} \qquad \Gamma, s_1[\overline{\alpha_1}] \to \mathbf{U}_{\varnothing} \overline{[\overline{\alpha}/\overline{\alpha_2}]\Gamma(s_2)}\ \mathbf{U}_{\leftarrow} \overline{l : \tau} \ \vdash\ e \ :\ \tau_* \ \Rightarrow\ E}{\Gamma \ \vdash\ \texttt{sig}\ s_1[\overline{\alpha_1}]\ \texttt{extends}\ \overline{s_2[\overline{\alpha}]}\ \texttt{where}\ \overline{l : \tau}\ \texttt{in}\ e \ :\ \tau_* \ \Rightarrow\ \texttt{let}\ merge_{s_1} : \ldots = \ldots\ \texttt{in}\ E}
$$

$$
\frac{\overline{\Gamma \vdash s[\overline{\alpha}]} \qquad \overline{\Gamma, \overline{x_1} : gen2_A(l_1) \ \vdash\ e_1 \ :\ \tau_1 \ \Rightarrow\ E_1} \qquad \Gamma, x : \&\overline{[\overline{\tau}/\overline{\alpha}]\Gamma(s)}, x \multimap \overline{s[\overline{\tau}]} \ \vdash\ e \ :\ \tau_* \ \Rightarrow\ E}{\begin{array}{c}\Gamma \ \vdash\ \texttt{algebra}\ x\ \texttt{implements}\ \overline{s[\overline{\tau}]}\ \texttt{where}\ \overline{l@(l_1\ \overline{x_1}) = e_1}\ \texttt{in}\ e \ :\ \tau_* \ \Rightarrow \\ \texttt{let}\ x : \&\overline{[\overline{\tau}/\overline{\alpha}]\Gamma(s)} = \overline{\{l_1 = \lambda(\overline{x_1}{:}gen2_A(l_1)).\{l = E_1\}\}}\ \texttt{in}\ E\end{array}}
$$
NEED TCHECK

$$
\frac{\overline{\Gamma \vdash s[\overline{\alpha}]} \qquad \overline{\Gamma \vdash x_0} \qquad \overline{\Gamma, \overline{x_1} : gen2_A(l_1) \ \vdash\ e_1 \ :\ \tau_1 \ \Rightarrow\ E_1} \qquad \Gamma, x : \&\overline{[\overline{\tau}/\overline{\alpha}]\Gamma(s)}, x \multimap \overline{s[\overline{\tau}]} \ \vdash\ e \ :\ \tau_* \ \Rightarrow\ E}{\begin{array}{c}\Gamma \ \vdash\ \texttt{algebra}\ x\ \texttt{extends}\ \overline{x_0}\ \texttt{implements}\ \overline{s[\overline{\tau}]}\ \texttt{where}\ \overline{l@(l_1\ \overline{x_1}) = e_1}\ \texttt{in}\ e \ :\ \tau_* \ \Rightarrow \\ \texttt{let}\ x : \&\overline{[\overline{\tau}/\overline{\alpha}]\Gamma(s)} = \overline{x_0} ,, \overline{\{l_1 = \lambda(\overline{x_1}{:}gen2_A(l_1)).\{l = E_1\}\}}\ \texttt{in}\ E\end{array}}
$$
NEED TCHECK

$$
\frac{\Gamma \vdash s[\overline{\alpha}] \to \overline{l : \tau} \qquad \Gamma, d \rightsquigarrow s[\overline{\alpha_0}].\alpha_1 : \{accept : \forall \alpha_1.\, s[\overline{\alpha_0}] \to \alpha_1\} \ \vdash\ e \ :\ \tau_* \ \Rightarrow\ E}{\Gamma \ \vdash\ \texttt{data}\ d\ \texttt{from}\ s[\overline{\alpha_0}].\alpha_1\ \texttt{in}\ e \ :\ \tau_* \ \Rightarrow\ \texttt{let}\ \overline{gen3(l)} = \ldots\ \texttt{in}\ E}
$$
NOT REVISED YET

$$
\frac{\Gamma \vdash d \qquad \Gamma, \overline{x_1} : \overline{\tau_1}, \overline{x_2} : d[\overline{\tau}] \ \vdash\ e_1 \ :\ d[\overline{\tau}] \ \Rightarrow\ E_1 \qquad \Gamma, x : \overline{\tau_1} \to \overline{d[\overline{\tau}]} \to d[\overline{\tau}] \ \vdash\ e \ :\ \tau_* \ \Rightarrow\ E}{\Gamma \ \vdash\ \texttt{let}\ x\ (\overline{x_1} : \overline{\tau_1})\ (\overline{x_2} : d[\overline{\tau}]) : d[\overline{\tau}] = e_1\ \texttt{in}\ e \ :\ \tau_* \ \Rightarrow\ \texttt{let}\ x = [gen4(d)]E_1\ \texttt{in}\ E}
$$
NOT REVISED YET

$$
\frac{\overline{\Gamma \vdash x \multimap s[\overline{\tau}]} \qquad \Gamma \vdash s[\overline{\alpha}] \qquad \Gamma \ \vdash\ e \ :\ \tau_0 \ \Rightarrow\ E}{\Gamma \ \vdash\ e[\overline{\tau}]\texttt{<}\overline{x}\texttt{>} \ :\ \tau_* \ \Rightarrow\ E.accept[\&\overline{\tau}]\ (merge_s\ \overline{x})}
$$
NEED TCHECK

merge_s: the merge algebra for object algebra interface s.

$\mathit{merge}_s : \forall\overline{\alpha_A}.\,\forall\overline{\alpha_B}.\, s[\overline{\alpha_A}] \to s[\overline{\alpha_B}] \to s[\overline{\alpha_A \,\&\, \alpha_B}] = \Lambda\overline{\alpha_A}.\,\Lambda\overline{\alpha_B}.\,\lambda(alg_1 : s[\overline{\alpha_A}]).\,\lambda(alg_2 : s[\overline{\alpha_B}]).\,\{\overline{l = [\overline{\alpha_A \,\&\, \alpha_B}/\overline{\alpha}]\mathit{gen}(l)}\}$

$\mathit{gen}(l)$: $\lambda(\overline{x} : \ldots).\, alg1.\overline{x}, , alg2.\overline{x}$.

$\mathit{gen2}(l)$: get the type from context $\Gamma(s).l$.

$\mathit{gen3}(l)$: for each case l, generate an auxiliary function for building structures. Only consider those with return type α_1 in $\texttt{data}\ d\ \texttt{from}\ s[\overline{\alpha_0}].\alpha_1\ \texttt{in}\ e$.

$[\mathit{gen4}(d)]$: $[\overline{l[\overline{\tau}]/\overline{l}}]$. Only when $d \rightsquigarrow s[\overline{\alpha_0}].\alpha_1$, l is a constructor in s, and $\mathit{gen3}(l)$ exists.

3 Auxiliary Rules for Expanding Types

4 Notes

The rules should support both special syntax for algebras and common syntax.

- **sig:** (1) in the environment; (2) as a type synonym.
- **alg:** (1) in the environment; (2) as a function.
- **data:** (1) in the environment; (2) as a type synonym.

Each datatype has only one sort. And instantiation only works for datatypes.

Type and consistency check need. Like in the declaration of an algebra, the label l should be consistent. And in the instantiation $e[\overline{\tau}]<\overline{x}>$, it requires $\overline{x \multimap s[\overline{\tau}]}$ with the same s.