

TYPE INFERENCE ALGORITHM

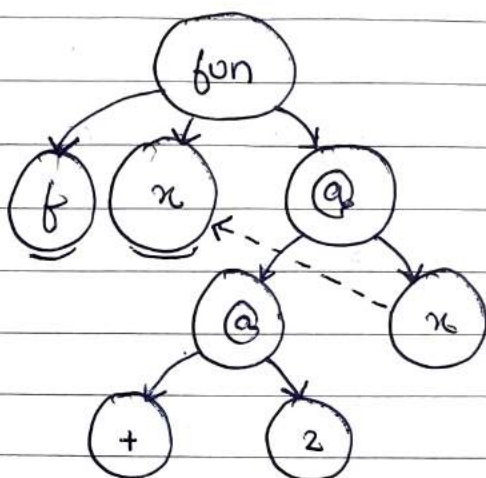
- parse program to build parse tree
- assign type variables to nodes in tree
- generate constraints
 - from environment: literals (2),
built-in operators (+),
known functions (tail).
 - from form of parse tree: application,
abstraction nodes.
- solve constraints using unification
- determine types of top-level declarations.

STEP 1: PARSE PROGRAM

- parse program text to construct parse tree.

let $f\ x = 2 + x$

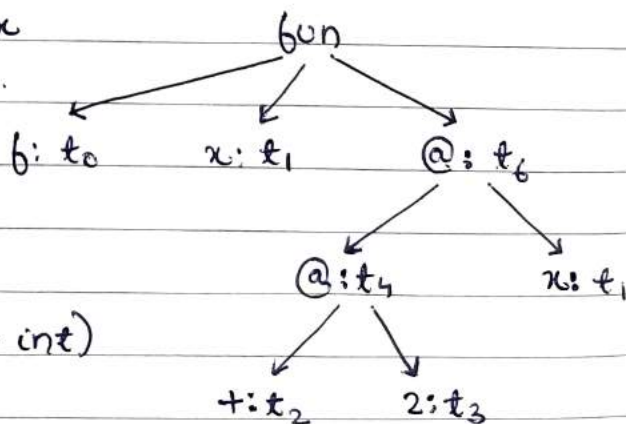
infix operators
are converted to
curried function
application during
parsing:



$2 + x \rightarrow (+) 2\ x$

STEP 3: ADD CONSTRAINTS

let $f\ x = 2 + x$



$t_3 = \text{int}$

$t_2 = \text{int} \rightarrow (\text{int} \rightarrow \text{int})$

$t_2 = t_3 \rightarrow t_4$

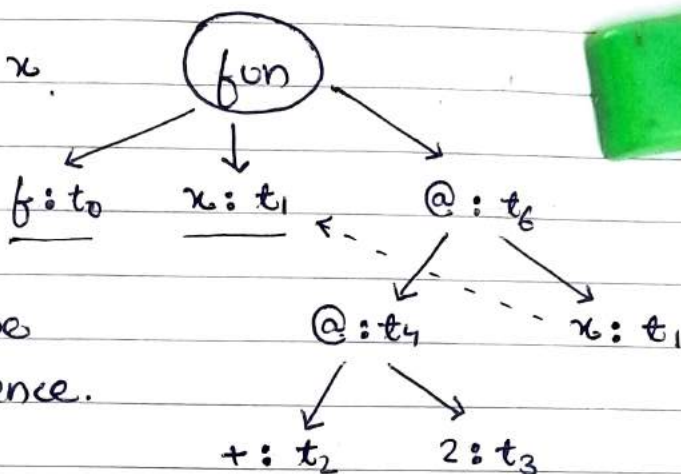
$t_4 = t_1 \rightarrow t_6$

$t_0 = t_1 \rightarrow t_6$

STEP 2: ASSIGN TYPE VARIABLES TO NODES.

let $f\ x = 2 + x$.

variables are
given same type
as binding occurrence.



STEP 4: SOLVE CONSTRAINTS.

① $t_0 = t_1 \rightarrow t_6$

$t_4 = t_1 \rightarrow t_6$

$t_2 = t_3 \rightarrow t_4$

$t_2 = \text{int} \rightarrow (\text{int} \rightarrow \text{int})$

$t_3 = \text{int}$

$t_3 \rightarrow t_4$

$= \text{int} \rightarrow (\text{int} \rightarrow \text{int})$

 \Downarrow

$t_3 = \text{int}$

$t_4 = \text{int} \rightarrow \text{int}$

② $t_0 = t_1 \rightarrow t_6$

$t_4 = t_1 \rightarrow t_6$

$t_4 = \text{int} \rightarrow \text{int}$

$t_1 \rightarrow t_6$

$= \text{int} \rightarrow \text{int}$

$\Rightarrow t_1 = \text{int}, t_6 = \text{int}$

③ $t_0 = \text{int} \rightarrow \text{int}$

STEP 5: DETERMINE TYPE OF DECLARATION

let $f \ x = 2 + x$

$t_0 = \text{int} \rightarrow \text{int}$

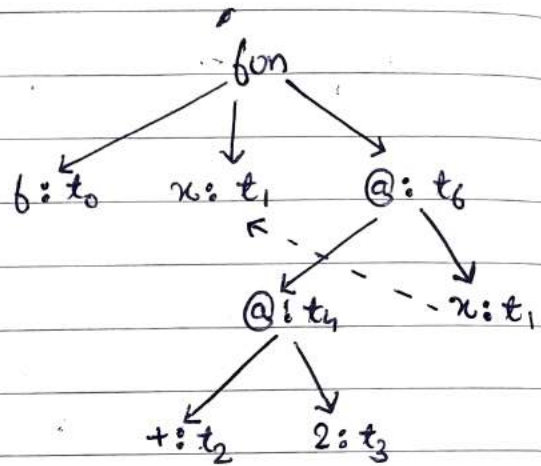
$t_1 = \text{int}$

$t_2 = \text{int} \rightarrow \text{int} \rightarrow \text{int}$

$t_3 = \text{int}$

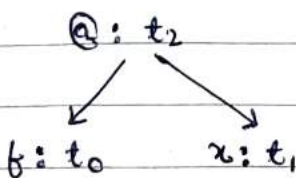
$t_4 = \text{int} \rightarrow \text{int}$

$t_6 = \text{int} \rightarrow \text{int}$



val $f: \text{int} \rightarrow \text{int} = \langle \text{fun} \rangle$

CONSTRAINTS FROM APPLICATION NODES



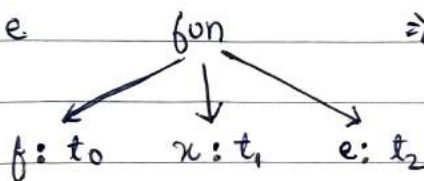
$f \ x$

$\Rightarrow t_0 = t_1 \rightarrow t_2$

domain range

CONSTRAINTS FROM ABSTRACTION NODES.

let $f \cdot x = e$



$\Rightarrow t_0 = t_1 \rightarrow t_2$

domain range