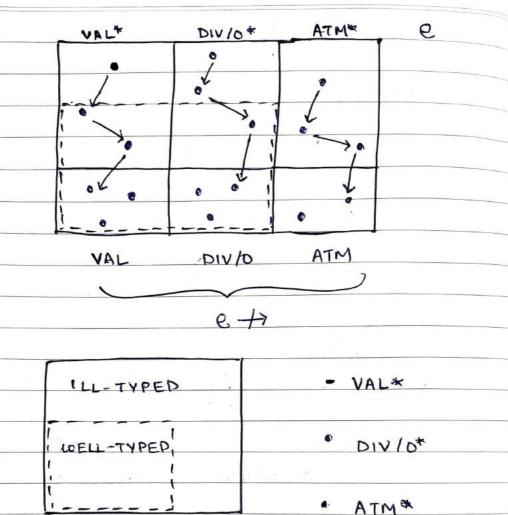


	TYPE SAFEJY
-	.3.
	preservation: if e -> e' and
	wer, then
	ω e' τ
	ω
	progress=: if to e ? then
	Ιω
	a) either & VAL, or
	CYC A *
	(b) e DIV/0, or
	(c) & man ->

Progress: + e ?

either e val\*, or

e Div/o\*

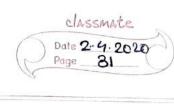


PRESERVATION

et and ete' => e' t

a well-typed expression preserves its well-typedness property (as also its type) upon reduction.

· (ill-typed) (well-typed)



PROOF OF PRESERVATION

given, e → e' , e ?

show e' ?

by induction on ->

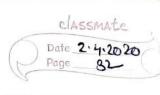
(a) e =  $\bar{n}_1(\bar{+}) \bar{n}_2$ 

N, NOM NOM NOM NOM No NUM ~ C=NUM

(b) n1 + n2 NUM

- NUM NI+N2 NUM

i. e NUM -> e' NUM



provided no \$ 0 similar to (1) (if) true ez ez -> ez e' by inversion, e2 & ez & , done e<sub>1</sub> → e'

e<sub>1</sub> ← e<sub>2</sub>

e<sub>1</sub> ← e<sub>2</sub>

e<sub>1</sub> ← e<sub>2</sub>

e<sub>1</sub> ← e<sub>2</sub> 8:20 by inversion & pattern matching ei + ez num plus 1

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6).  $e_1 \rightarrow e_1'$  DIV-LEFT  $e_1 \bigcirc e_2 \rightarrow e_1' \bigcirc e_2$  Simplan to  $\bigcirc$ .

© e1 VAL e2 → e2 PLUS-RIGHT

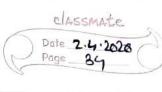
e1 + e2 → e1+ e2

e'

by inversion & pattern matching,

7 = NUM, C1 NUM

by IH, e' NUM



-	.: EI NOM C2' NOM
	e, Fez NUM
	6,
(1)	$\frac{e_1 \vee AL}{e_1 \otimes e_2} \xrightarrow{e_2} \frac{e_2'}{e_1 \otimes e_2'} = \frac{e_1 \vee e_2'}{e_2'}$
	e, (1) e2 -> e, (1) e2'
	Similar to (E)
(8)	$e_1 \rightarrow e_1'$
۳	(if) e1 e2 e3 → (if) e1 e2 e3
	e e'
	ę, ·
	A A
	e c
	by inversion & pattern matching
	e1 BOOL :
	e2 ℃ (b) ei e2 e3 ℃
	e <sub>3</sub> °C
	by IH,
	ei Bool

QED



-	
	thus we have proved PRESERVATION
	e 7 and e > e'
	か e' で
	PRO GRESS
	if e & then,
	e val, or
	& DIV/O, or
	е →
	PROOF BY INDUCTION ON & &
0	NUM clearly C VAL
	n NUM
(2	
	5 BOOL

e

EI NUM E2 NUM PLUS eiflez NUM by IH, e, VAL, or e, DIV/D or  $e_1 \rightarrow$ (9 cases) ez VAL, or ez Divlo or  $e_2 \rightarrow$  $e_1 \rightarrow$ then let, e, -> e,  $e_1 \rightarrow e_1' \quad e_2$ 

·· e ->

e, DIV/O

3.3.1) e2 ->

(3.2)

then we have a DIV/o deduction

· e ->

e2 DIV/O

e1 1 e2 DIV/0

8.3.2

e, VAL 3 cases for e2

e, (+) e2 DIV/O

e, DIV/O : LEFT

then we have a deduction in ->

eIVAL e2 -> e2 P-LEFT

then we have a derivation in DIVO

e, VAL e2 DIV/O DIV/O: RIGHT

e1+e2 → e1+e2

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(3.3.3) e2 VAL

4 cases

 $\begin{cases} e_1 = \overline{n}, \\ e_1 = \overline{b}, \end{cases} \qquad \begin{cases} e_2 = \overline{n}_2 \\ e_2 = \overline{b}_2 \end{cases}$ 

e = b is not possible : e NUM

of these 4 cases

 $e_2 = \overline{n}_2$ 

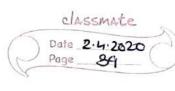
62 = 6 for the same reason

: this leaves only one are

e1 = n, : e1 + e2 = n, + n2

· e ->

done.



\_ OF-TEST

ENDMENUM DIV

ENDRED

EXERCISE

(1) ENDOL ENDRED

(2) ENDRED

(3) ENDOL ENDRED

(4) ENDRED

(4) ENDRED

(5) ENDRED

(6) ENDRED

(7) ENDRED

(7) ENDRED

(8) ENDRED

(8) ENDRED

(9) ENDRED

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(9) ENDRED

(1) ENDRED

(2) ENDRED

(3) ENDRED

(4) ENDRED

(5) ENDRED

(6) ENDRED

by inversion & pattern Matching

e, Bool e2 7 e3 8

by 14, there are 3 cases

e1 VAL

e, DIV/0

----

then we have a derivation

··· e -> done.

(s.i) e, ->



(5.2) e, DIV/O

e, DIV/O : IF ( ) e1 e2 e3 DIV/O

:. e DIV/0

(63) e, VAL

(531) e1 = true, then

sine e BOOL, it follows that

IF-TRUE

(if etrue ez ez -> ez

·. e→

5.3.2) e1 = false

similar

QED