

Macro Assembler and Utilities

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Preface

This manual describes how to use the A51, AX51, and A251 macro assemblers and the related utilities to translate assembly source code into executable programs for the 8051 and variants like the Philips 80C51MX and Intel/Atmel WM 251 devices. This manual assumes you are familiar with the Windows operating system and know how to program microcontrollers.

- "Chapter 1. Introduction," provides an overview of the different assembler variants and describes the basics of assembly language programming.
- "Chapter 2. Architecture," contains an overview of the 8051, extended 8051, Philips 80C51MX, and Intel/Atmel WM 251 devices.
- "Chapter 3. Writing Assembly Programs," describes assembler statements and the rules for arithmetic and logical expressions.
- "Chapter 4. Assembler Directives," describes how to define segments and symbols and how to use all directives.
- "Chapter 5. Assembler Macros," describes the function of the standard macros and contains information for using standard macros.
- "Chapter 6. Macro Processing Language," defines and describes the use of the Intel Macro Processing Language.
- "Chapter 7. Invocation and Controls," describes how to invoke the assembler and how to control the assembler operation.
- "Chapter 8. Error Messages," contains a list of all assembler error messages and describes their causes and how to avoid them.
- "Chapter 9. Linker/Locator," includes reference section of all linker/locater directives, along with examples and detailed descriptions.
- "Chapter 10. Library Manager," shows you how to create and maintain a library.
- "Chapter 11. Object-Hex Converter," describes how to create Intel HEX files.

The Appendix contains program examples, lists the differences between assembler versions, and contains other items of interest.

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Document Conventions

This document uses the following conventions:

Examples	Description		
README.TXT	Bold capital text is used for the names of executable programs, data files, source files, environment variables, and commands you enter at the Windows command prompt. This text usually represents commands that you must type in literally. For example:		
	CLS	DIR	BL51.EXE
	Note that you are not reqletters.	uired to enter these	e commands using all capital
Courier	Text in this typeface is us screen or prints at the pri		ormation that displays on
	This typeface is also use command line items.	d within the text wh	nen discussing or describing
Variables	Text in italics represents information that you must provide. For example, <i>projectfile</i> in a syntax string means that you must supply the actual project file name.		
	Occasionally, italics are also used to emphasize words in the text.		
Elements that repeat	Ellipses () are used in examples to indicate an item that may be repeated.		
Omitted code	Vertical ellipses are used in source code examples to indicate that a fragment of the program is omitted. For example:		
	void main (void)	{	
•	•		
	•		
	while (1);		
[Optional Items]	Optional arguments in codouble brackets. For exa		ption fields are indicated by
	C51 TEST.C PRINT	$\llbracket ext{(filename)} brace$	
{ opt1 opt2 }	group of items from which	n one must be chooning	a vertical bar represents a sen. The braces enclose all e the choices. One item in
Keys	Text in this sans serif typ For example, "Press Ent	•	ctual keys on the keyboard.

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Chapter 1. Introduction

This manual describes the macro assemblers and utilities for the classic 8051, extended 8051, and 251 microcontroller families and explains the process of developing software in assembly language for these microcontroller families.

A brief overview of the classic 8051, the extended 8051, and the 251 architectures may be found in "Chapter 2. Architecture Overview" on page 27. In this overview, the differences between the classic 8051, the extended 8051 variants and the 251 processors are described. For the most complete information about the microcontroller architecture refer to the hardware reference manual of the microcontroller derivative that your are using.

For optimum support of the different 8051 and 251 variants, Keil provides the following development tools:

Development Tools	Support Microcontrollers, Description
A51 Macro Assembler BL51 Linker/Locater LIB51 Library Manager	Development Tools for classic 8051 . Includes support for 32 x 64KB code banks.
AX51 Macro Assembler LX51 Extended Linker/Locater LIBX51 Library Manager	Development Tools for classic and extended 8051 versions (Philips 80C51MX , Dallas 390 , etc.) Supports up to 16MB code and xdata memory.
A251 Macro Assembler L251 Linker/Locater LIB251 Library Manager	Development Tools for Intel/Atmel WM 251.

The AX51 and A251 assemblers are supersets of the A51 assembler. This user's guide therefore covers all development tools variants. Whenever a feature or an option is available in one specific tool chain only, it is clearly marked.

For general reference to all tool variants and microcontroller architectures the terms listed in the following table are used:

Term	Refers to
Ax51 Macro Assembler	A51, AX51 and A251 Macro Assembler
Cx51 Compiler	C51, CX51 and C251 ANSI C Compiler
Lx51 Linker/Locator	BL51, LX51 and L251 Linker/Locator
LIBx51 Library Manager	LIB51, LIBX51 and LIB251 Library Manager
OHx51 Object-Hex Converter	OH51, OHX51 and OH251 Object-Hex Converter
x51 Architecture or x51 Device	All classic 8051, extended 8051 and 251 device variants.

How to Develop A Program

This section presents an overview of the Ax51 macro assembler, Lx51 linker/locater and how it is used.

What is an Assembler?

An assembler is a software tool designed to simplify the task of writing computer programs. It translates symbolic code into executable object code. This object code may then be programmed into a microcontroller and executed. Assembly language programs translate directly into CPU instructions which instruct the processor what operations to perform. Therefore, to effectively write assembly programs, you should be familiar with both the microcomputer architecture and the assembly language.

Assembly language operation codes (mnemonics) are easily remembered (MOV for move instructions, ADD for addition, and so on). You can also symbolically express addresses and values referenced in the operand field of instructions. Since you assign these names, you can make them as meaningful as the mnemonics for the instructions. For example, if your program must manipulate a date as data, you can assign it the symbolic name DATE. If your program contains a set of instructions used as a timing loop (a set of instructions executed repeatedly until a specific amount of time has passed), you can name the instruction group TIMER LOOP.

An assembly program has three constituent parts:

- Machine instructions
- Assembler directives
- Assembler controls

A machine instruction is a machine code that can be executed by the machine. Detailed discussion of the machine instructions can be found in the hardware manuals of the 8051 or derivative microcontroller. Appendix A provides an overview about machine instructions.

Assembler directives are used to define the program structure and symbols, and generate non-executable code (data, messages, etc.). Refer to "Chapter 4. Assembler Directives" on page 99 for details on all of the assembler directives.

Assembler controls set the assembly modes and direct the assembly flow. "Chapter 7. Invocation and Controls" on page 195 contains a comprehensive guide to all the assembler controls.

Modular Programming

Many programs are too long or complex to write as a single unit. Programming becomes much simpler when the code is divided into small functional units. Modular programs are usually easier to code, debug, and change than monolithic programs.

The modular approach to programming is similar to the design of hardware that contains numerous circuits. The device or program is logically divided into "black boxes" with specific inputs and outputs. Once the interfaces between the units have been defined, the detailed design of each unit can proceed separately.

The benefits of modular programming are:

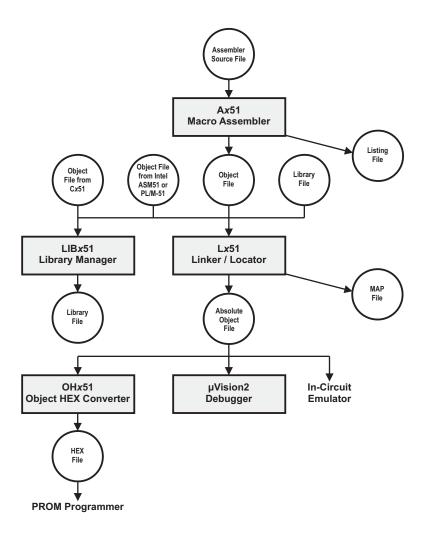
Efficient Program Development: programs can be developed more quickly with the modular approach since small subprograms are easier to understand, design, and test than large programs. With the module inputs and outputs defined, the programmer can supply the needed input and verify the correctness of the module by examining the output. The separate modules are then linked and located by the linker into an absolute executable single program module. Finally, the complete module is tested.

Multiple Use of Subprograms: code written for one program is often useful in others. Modular programming allows these sections to be saved for future use. Because the code is relocatable, saved modules can be linked to any program which fulfills their input and output requirements. With monolithic programming, such sections of code are buried inside the program and are not so available for use by other programs.

Ease of Debugging and Modifying: modular programs are generally easier to debug than monolithic programs. Because of the well defined module interfaces of the program, problems can be isolated to specific modules. Once the faulty

module has been identified, fixing the problem is considerably simpler. When a program must be modified, modular programming simplifies the job. You can link new or debugged modules to an existing program with the confidence that the rest of the program will not change.

The following figure shows an overview of the steps involved in creating a program for the *x*51.



Modular Program Development Process

This section is a brief discussion of the program development process with the relocatable **Ax51** assembler, **Lx51** Linker/Locator, and the **OHx51** code conversion program.

Segments, Modules, and Programs

In the initial design stages, the tasks to be performed by the program are defined, and then partitioned into subprograms. Here are brief introductions to the kinds of subprograms used with the **Ax51** assembler and **Lx51** linker/locator.

A segment is a block of code or data memory. A segment may be relocatable or absolute. A relocatable segment has a name, type, and other attributes. Segments with the same name, from different modules, are considered part of the same segment and are called *partial segments*. Several *partial segments* with the same name are combined into one segment by the **Lx51** linker/locater. An absolute segment cannot be combined with other segments.

A module contains one or more segments or partial segments. A module is a source code unit that can be translated independently. It contains all symbol definitions that are used within the module. A module might be a single ASCII text file that is created by any standard text editor. However, with you may use the *include* assembler directive to merge several text files. The **Ax51** assembler translates a source file into an object file. Each object file is one module.

After assembly of all modules of the program, **Lx51** processes the object module files. The **Lx51** linker/locator assigns absolute memory locations to all the relocatable segments, combining segments with the same name and type. **Lx51** also resolves all references between modules. **Lx51** outputs an absolute object module file with the completed program, and a map file that lists the results of the link/locate process.

Translate and Link Process

Typically you will use the **Ax51** assembler and the tools within the μ Vision2 IDE. For more information on using the μ Vision2 IDE refer to the User's Guide μ Vision2: Getting Started for 8051.

However, you may invoke the **Ax51** assembler also from the command line. Simply type the name of the assembler version that you want to use, for example **A51** at the Windows command prompt. On this command line, you must include the name of the assembler source file to be translated, as well as any other necessary control directives required to translate your source file. Example:

```
A51 DEMO.A51
```

The assembler output for this command line is:

```
A51 MACRO ASSEMBLER V6.00
ASSEMBLY COMPLETE. 0 WARNING(S), 0 ERROR(S)
```

After assembly of all your program modules, the object modules are linked and all variables and addresses are resolved and located into an executable program by the **Lx51** linker. The following example shows a simple command line for the linker:

```
BL51 DEMO.OBJ, PRINT.OBJ
```

The linker generates an absolute object file as well as a map file that contains detailed statistic information and screen messages. The output of the linker is:

```
BL51 LINKER/LOCATER V4.00
LINK/LOCATE RUN COMPLETE. 0 WARNING(S), 0 ERROR(S)
```

Then you might convert the executable program into an Intel HEX file for PROM programming. This is done with the **OH***x***51** hex conversion utility with the following invocation:

```
OH51 DEMO
```

The output of the hex conversion utility is:

```
OBJECT TO HEX FILE CONVERTER OH51 V2.40 GENERATING INTEL HEX FILE: DEMO.HEX OBJECT TO HEX CONVERSION COMPLETED.
```

1

An example listing file generated by the assembler is shown on the following page.

```
A51 MACRO ASSEMBLER ASSEMBLER DEMO PROGRAM
                                                     07/07/2000 18:32:30 PAGE 1
MACRO ASSEMBLER A51 V6.01
OBJECT MODULE PLACED IN demo.OBJ
ASSEMBLER INVOKED BY: C:\KEIL\C51\BIN\A51.EXE DEMO.A51 DEBUG
LOC OBJ
                  LINE
                          SOURCE
                         $title (ASSEMBLER DEMO PROGRAM)
                     1
                     2
                          ; A simple Assembler Module for Demonstration
                         ; Symbol Definition
                               EQU 13
 0000
                     5
                         CR
                                               ; Carriage?Return
  000A
                                 EQU 10
                     6
                          LF
                                                ; Line?Feed
                     7
                     8
                          ; Segment Definition
                     9
                          ?PR?DEMO SEGMENT CODE ; Program Part
                         ?CO?DEMO SEGMENT CODE ; Constant Part
                    10
                    11
                          ; Extern Definition
                    12
                         EXTRN CODE (PRINTS, DEMO)
                    13
                    14
                    15
                         ; The Program Start
                    16
                                 CSEG AT 0
                                                ; Reset Vector
0000 020000 F
                   17
                                 JMP Start
                   18
                    19
                                 RSEG ?PR?DEMO ; Program Part
0000 900000
                        START: MOV DPTR, #Txt; Demo Text
                    20
                    21
                                  CALL PRINTS ; Print String
0003 120000 F
                   22
0006 020000 F
                                 JMP DEMO
                   23
                                               ; Demo Program
                    24
                         ; The Text Constants
                    25
                   26
                                RSEG ?CO?DEMO ; Constant Part
0000 48656C6C
                  27 Txt: DB 'Hello World', CR, LF, 0
0004 6F20576F
0008 726C640D
000C 0A00
                    28
                    29
                                 END
                                               : End of Module
SYMBOL TABLE LISTING
NAME
                  TYPE VALUE ATTRIBUTES
?CO?DEMO . . . . C SEG
                          H3000
                                    REL-UNIT
?PR?DEMO . . . . C SEG
                          0009H
                                    REL=UNIT
CR . . . . . . N NUMB
                          000DH A
DEMO . . . . . . . C ADDR LF . . . . . . N NUMB
                                     EXT
                          000AH A
PRINTS . . . . . C ADDR
START. . . . . . C ADDR
                          0000H R SEG=?PR?DEMO
TXT. . . . . . . C ADDR
                          0000H R SEG=?CO?DEMO
REGISTER BANK(S) USED: 0
ASSEMBLY COMPLETE. 0 WARNING(S), 0 ERROR(S)
```

Filename Extensions

Typically, the filename extension is used to indicate the contents of each file. The following table lists the file name extensions that are used in the 8051 tool chain.

Extension	Content and Description
.A51 .ASM .SRC	Source code file: contains ASCII text that is the input for the Ax51 assembler.
.C .C51	C source code file: contains ASCII text that is the input for the Cx51 ANSI C compiler.
.INC .H	Include file: contains ASCII text that is merged into an source code file with the include directive. Also these files are input files for Ax51 or Cx51 .
.OBJ	Relocatable object file: is the output of the Ax51 or Cx51 that contains the program code and control information. Several relocatable object files are typically input files for the Lx51 Linker/Locater.
.LST	Listing object file: is generated by Ax51 or Cx51 to document the translation process. A listing file typically contains the ASCII program text and diagnostic information about the source module. Appendix F describes the format of the Ax51 listing file.
. (none)	Absolute object file: is the output of the Lx51 . Typically it is a complete program that can be executed on the x51 CPU.
.M51 .MAP	Linker map file: is the listing file generated from Lx51 . A map file contains information about the memory usage and other statistic information.
.HEX .H86	Hex file: is the output file of the OHx51 object hex converter in Intel HEX file format. HEX files are used as input file for PROM programmers or other utility programs.

Program Template File

The following code template contains guidelines and hints on how to write your own assembly modules. This template file TEMPLATE.A51 is provided in the folder \C51\ASM or \C251\ASM.

```
$NOMOD51
                        ; disable predefined 8051 registers
#include <reg52.h>
                      // include CPU definition file (for example, 8052)
; Change names in lowercase to suit your needs.
; This assembly template gives you an idea of how to use the A251/A51
; Assembler. You are not required to build each module this way-this is only
; an example.
; All entries except the END statement at the End Of File are optional.
; If you use this template, make sure you remove any unused segment declarations,
; as well as unused variable space and assembly instructions.
; This file cannot provide for every possible use of the A251/A51 Assembler.
; Module name (optional)
               module name
; Here, you may import symbols form other modules.
EXTRN CODE (code_symbol) ; May be a subroutine entry declared in
                                ; CODE segments or with CODE directive.
EXTRN DATA (data_symbol) ; May be any symbol declared in DATA segments
                                ; segments or with DATA directive.
EXTRN BIT (bit symbol)
                                ; May be any symbol declared in BIT segments
                                 ; or with BIT directive.
EXTRN XDATA (xdata_symbol) ; May be any symbol declared in XDATA segments
                                 ; or with XDATA directive.
EXTRN NUMBER (typeless_symbol); May be any symbol declared with EQU or SET
; You may include more than one symbol in an EXTRN statement.
EXTRN CODE (sub routine1, sub routine2), DATA (variable 1)
; Force a page break in the listing file.
$EJECT
; Here, you may export symbols to other modules.
PUBLIC data variable
PUBLIC code entry
PUBLIC typeless_number
PUBLIC xdata_variable
PUBLIC bit_variable
```

```
; You may include more than one symbol in a PUBLIC statement.
PUBLIC data variable1, code table, typeless num1, xdata variable1
; Put the STACK segment in the main module.
                                ; ?STACK goes into IDATA RAM.
; switch to ?STACK segment.
; reserve your stack space
                SEGMENT IDATA
?STACK
                RSEG ?STACK
                       5
                DS
                                      ; 5 bytes in this example.
$EJECT
; Put segment and variable declarations here.
; DATA SEGMENT--Reserves space in DATA RAM--Delete this segment if not used.
data_seg_name SEGMENT DATA ; segment for DATA RAM.

RSEG data_seg_name ; switch to this data segment
data_variable: DS 1 ; reserve 1 Bytes for data_variable
data variable1: DS
                                      ; reserve 2 Bytes for data variable1
; XDATA SEGMENT--Reserves space in XDATA RAM--Delete this segment if not used.
xdata_seg_name SEGMENT XDATA ; segment for XDATA RAM
RSEG xdata_seg_name ; switch to this xdata segment
xdata_variable: DS 1 ; reserve 1 Bytes for xdata_variable xdata_array: DS 500 ; reserve 500 Bytes for xdata_array
; INPAGE XDATA SEGMENT--Reserves space in XDATA RAM page (page size: 256 Bytes)
; INPAGE segments are useful for @RO addressing methodes.
; Delete this segment if not used.
; reserve 1 Bytes for xdata_variable1
xdata_variable1:DS
; ABSOLUTE XDATA SEGMENT--Reserves space in XDATA RAM at absolute addresses.
; ABSOLUTE segments are useful for memory mapped I/O.
; Delete this segment if not used.
      XSEG AT 8000H ; switch absolute XDATA segment @ 8000H
DS 1 ; reserve 1 Bytes for XIO port
             DS 1
DS 1
XCONFIG: DS
                                      ; reserve 1 Bytes for XCONFIG port
; BIT SEGMENT--Reserves space in BIT RAM--Delete segment if not used.
; segment for BIT RAM.
               RSEG bit_seg_name ; switch to this bit segment
DBIT 1 : reserve 1 Bit for bit veri
bit_variable: DBIT
bit_variable: DBIT 1 ; reserve 1 Bit for bit_variable bit_variable1: DBIT 4 ; reserve 4 Bits for bit_variable1
; Add constant (typeless) numbers here.
SEJECT
```

```
; Provide an LJMP to start at the reset address (address 0) in the main module.
; You may use this style for interrupt service routines.
                   AT 0 ; absolute Segment at Address 0
              CSEG
              LJMP
                     start
                                   ; reset location (jump to start)
; CODE SEGMENT--Reserves space in CODE ROM for assembler instructions.
code_seg_name SEGMENT CODE
              RSEG
                                  ; switch to this code segment
                     code_seg_name
              USTNG
                                   ; state register_bank used
                                   ; for the following program code.
              MOV
                    SP, #?STACK-1 ; assign stack at beginning
start:
; Insert your assembly program here. Note, the code below is non-functional.
                                   ; enable interrupt system (timer 0)
              ORL IE, #82H
                   TR0
              SETB
                                   ; enable timer 0
repeat_label:
             MOV
                    A,data_symbol
              ADD
                    A, #typeless symbol
              CALL
                    code_symbol
              MOV
                     DPTR, #xdata symbol
              MOVX
                     A,@DPTR
              MOV
                    R1,A
                    AR1
              PIISH
              CALL
                    sub routine1
              POP
                     AR1
              ADD
                    A,R1
              JMP
                    repeat label
code_entry:
              CALL
                   code_symbol
              RET
code_table:
             שמ
                     repeat_label
              DW
                     code entry
              DB
                     typeless number
$EJECT
; To include an interrupt service routine, provide an LJMP to the ISR at the
; interrupt vector address.
              CSEG
                    AT OBH ; OBH is address for Timer O interrupt
              LJMP
                     timer0int
; Give each interrupt function its own code segment.
RSEG int0_code_seg ; switch to this code segment
              USING 1
                                  ; register bank for interrupt routine
timer0int:
              PUSH
                    PSW
              MOV
                     PSW, #08H
                                   ; register bank 1
              PUSH
                     ACC
              MOV
                     R1, data variable
              MOV
                     DPTR, #xdata variable
              MOVX
                     A,@DPTR
              ADD
                     A,R1
              MOV
                     data_variable1,A
              CLR
              ADD
                     A,#0
```

1

```
MOV data_variable1+1,A
POP ACC
POP PSW
RETI

; The END directive is ALWAYS required.
; END ; End Of File
```

Chapter 2. Architecture Overview

This chapter gives you an overview of the 8051 architecture and the variants of the 8051. It reviews the memory layout of the classic 8051, extended 8051 variants, the Philips 80C51MX, and the 251 architecture. Also described are the register sets and the CPU instructions of the various CPU variants.

Memory Classes and Memory Layout

This section introduces the different memory classes (also known as memory types) that are used during programming of the 8051 and variants. Memory classes are used to identify the different physical memory regions of the microcontroller architecture that can be represented in a memory layout.

An overview of the different physical memory regions in an *x*51 system is provided below:

Program Memory: in the classic 8051 this is a 64KB space that is called CODE. This region is typically a ROM space that is used for program code and constants. With the BL51 you may expand the physical program code memory to 32 code banks with 64KB each. Constants are fetched with the MOVC instruction. In extended 8051 variants and the 251 you may have program memory of up to 16MB that is called ECODE and HCONST.

Internal Data Memory: in the classic 8051 this is the on-chip RAM space with a maximum of 256 Bytes that contains register banks, BIT space, direct addressable DATA space, and indirect addressable IDATA space. This region should be used for frequently used variables. In the 80C51MX and the 251 this space is expanded to up to 64KB with an EDATA space.

External Data Memory: in classic 8051 devices this area, called XDATA, is off-chip RAM with a space of up to 64KB. However several new 8051 devices have additional on-chip RAM that is mapped into the XDATA space. Usually you need to enable this additional on-chip RAM via dedicated SFR registers. In extended variants and the 251 you may have external data memory of up to 16MB that is called HDATA.

Classic 8051

The following table shows the memory classes used for programming the classic 8051 architecture. These memory classes are available when you are using the A51 macro assembler and the **BL51** linker/locater.

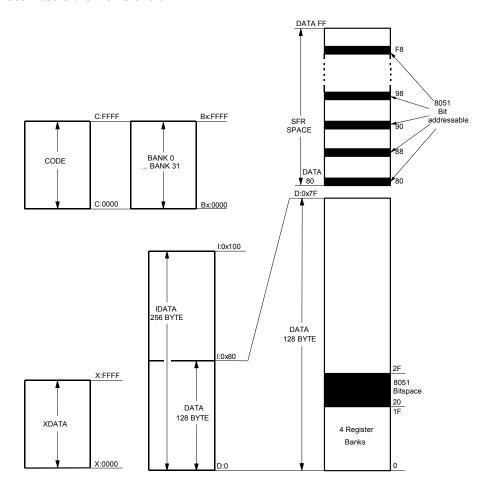
Memory Class	Address Range	Description
DATA	D:00 – D:7F	Direct addressable on-chip RAM.
BIT	D:20 – D:2F	bit-addressable RAM; accessed bit instructions.
IDATA	I:00 – I:FF	Indirect addressable on-chip RAM; can be accessed with @R0 or @R1.
XDATA	X:0000 – X:FFFF	64 KB RAM (read/write access). Accessed with MOVX instruction.
CODE	C:0000 – C:FFFF	64 KB ROM (only read access possible). Used for executable code or constants.
BANK 0 BANK 31	B0:0000 – B0:FFFF B31:0000 – B31:FFFF	Code Banks for expanding the program code space to 32 x 64KB ROM.

NOTE

The memory prefix D: I: X: C: B0: .. B31: cannot be used at Ax51 assembler or BL51 linker/locater level. The memory prefixes are only listed for better understanding. Several Debugging tools, for example the μV ision2 Debugger, are using memory prefixes to identify the memory class of the address.

Classic 8051 Memory Layout

The classic 8051 memory layout, shown in the following figure, is familiar to 8051 users the world over.



The memory code banks overlap the CODE space. The size of the code banks is selected with the **Lx51** directive BANKAREA.

Extended 8051 Variants

Several new variants of the 8051 extend the code and/or xdata space of the classic 8051 with address extension registers. The following table shows the memory classes used for programming the extended 8051 devices. These memory classes are available for classic 8051 devices when you are using memory banking with the LX51 linker/locater. In addition to the code banking known from the BL51 linker/locater, the LX51 linker/locator supports also data banking for xdata and code areas with standard 8051 devices.

Memory Class	Address Range	Description
DATA	D:00 – D:7F	Direct addressable on-chip RAM.
BIT	D:20 – D:2F	bit-addressable RAM; accessed bit instructions.
IDATA	I:00 – I:FF	Indirect addressable on-chip RAM; can be accessed with @R0 or @R1.
XDATA	X:0000 – X:FFFF	64 KB RAM (read/write access). Accessed with MOVX instruction.
HDATA	X:0000 – X:FFFFFF	16 MB RAM (read/write access). Accessed with MOVX instruction and extended DPTR.
CODE	C:0000 – C:FFFF	64 KB ROM (only read access possible). Used for executable code or constants.
ECODE	C:0000 – C:FFFFFF	16 MB ROM (only read access possible). Used for constants. In some modes of the Dallas 390 architecture also program execution is possible.
BANK 0 BANK 31	B0:0000 – B0:FFFF B31:0000 – B31:FFFF	Code Banks for expanding the program code space to 32 x 64KB ROM.

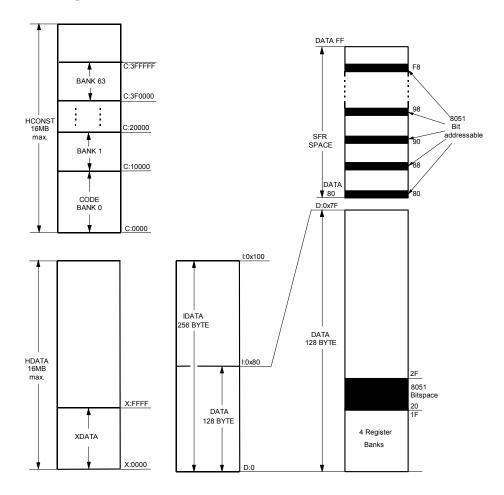
NOTES

The memory prefixes D: I: X: C: B0: .. B31: cannot be used at Ax51 assembler level. The memory prefix is only listed for better understanding. The Lx51 linker/locater and several Debugging tools, for example the μ Vision2 Debugger, are using memory prefixes to identify the memory class of the address.

If you are using the Dallas 390 contiguous mode the address space for CODE can be C:0000 - C:0xFFFFFF.

Extended 8051 Memory Layout

The extended 8051 memory layout, shown in the following figure, expands the address space for variables to a maximum of 16MB.



In several variants the DPTR register is expanded to a 24-bit register with an **DPX** SFR. Fox example, the Dallas 390 and provides new operating modes where this addressing is enabled. You may even use the HCONST and HDATA memory classes with classic 8051 devices by using the memory banking available in LX51.

Philips 80C51MX

The Philips 80C51MX provides a unified 16 MB address space. New instructions can access up to 16MB memory whereby CODE and XDATA space are mapped into one single address space. The stack pointer can be configured as 16-Bit stack pointer that addresses the on-chip RAM area in the EDATA memory class. The following table shows the memory classes used for programming the 80C51MX architecture. These memory classes are available when you are using the AX51 macro assembler and the LX51 linker/locater.

Memory Class	Address Range	Description
DATA	7F:0000 – 7F:007F	Direct addressable on-chip RAM.
BIT	7F:0020 – 7F:002F	Bit-addressable RAM; accessed bit instructions.
IDATA	7F:0000 – 7F:00FF	Indirect addressable on-chip RAM; can be accessed with @R0 or @R1.
EDATA	7F:0000 – 7F:FFFF	Complete on-chip RAM; can be used as stack space or can be accessed with @PR0 or @PR1.
XDATA	00:0000 – 00:FFFF	64 KB RAM (read/write access). Accessed with MOVX instruction.
HDATA	00:0000 – 7F:FFFF	8 MB RAM (read/write access). Accessed with MOVX instruction and extended DPTR.
CODE	80:0000 – 80:FFFF	Classic 8051 compatible 64 KB ROM (only read access possible). Used for executable code or constants.
ECODE	80:0000 - 80:FFFF	8 MB ROM (only read access possible).
HCONST	80:0000 – 80:FFFF	8 MB ROM. Same as ECODE, this class is used by the CX51 Compiler for constants.
BANK 0 BANK 63	80:0000 – 0xBF:FFFF B0:0000 – B63:FFFF	Used by the CX51 Compiler to expand the program memory to more than 64KB. Refer to "Philips 80C51MX" on page 277 for more information.

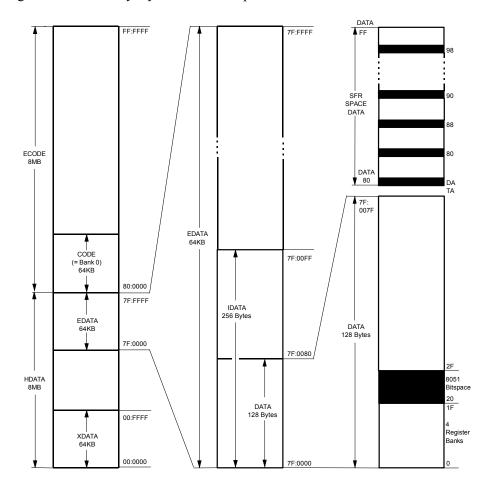
NOTES

Colons are used to improve the readability only. The addresses are entered in the tools as numbers without colon.

The memory prefixes D: I: X: C: B0: .. B31: cannot be used at Ax51 assembler level. The memory prefix is only listed for better understanding. The Lx51 linker/locater and several Debugging tools, for example the μ Vision2 Debugger, are using memory prefixes to identify the memory class of the address.

80C51MX Memory Layout

The Philips 80C51MX memory layout, shown in the following figure, provides a universal memory map that includes all memory types in a single 16MB address region. The memory layout of the Philips 80C51MX is shown below:



The 80C51MX offers new CPU instructions that provide a new addressing mode, called Universal Pointer addressing. Two Universal Pointer registers PR0 and PR1 are available. PR0 is composed of registers R1, R2, and R3. PR1 is composed of registers R5, R6, and R7. These new Universal Pointer registers hold a 24-bit address that is used together with the EMOV instruction to address the complete 16MB memory.

Intel/Atmel WM 251

Also the 251 architecture is a superset of the classic 8051 architecture. The 251 is the most advanced variant and provides the following key features:

- Completely code compatible with the standard 8051 microcontroller.
- Powerful 8/16/32-bit instructions and flexible 8/16/32-bit registers.
- 16MB linear address space and CPU support for 16-bit and 32-bit pointers.
- True stack-oriented instructions with 16-bit stack pointer.

The following table shows the memory classes used for programming a 251 microcontroller. These memory classes are available when you are using the A251 macro assembler and the L251 linker/locater.

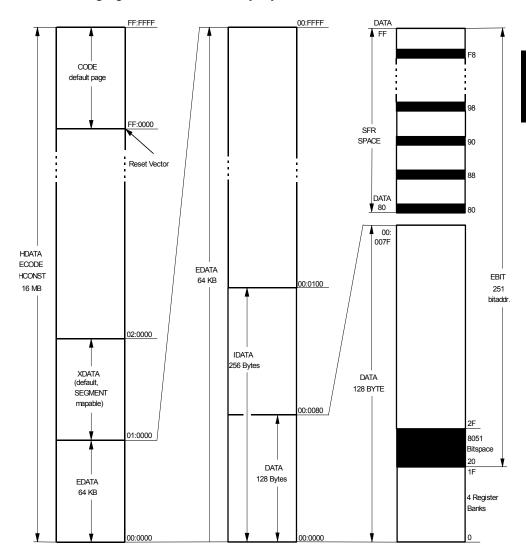
Memory Class	Address Range	Description
DATA	00:0000 - 00:007F	Direct addressable on-chip RAM.
BIT	00:0020 - 00:002F	8051 compatible bit-addressable RAM; can be accessed with short 8-bit addresses.
IDATA	00:0000 - 00:00FF	Indirect addressable on-chip RAM; can be accessed with @R0 or @R1.
EDATA	00:0000 - 00:FFFF	Extended direct addressable memory area; can be accessed with direct 16-bit addresses available on the 251.
ECONST	00:0000 - 00:FFFF	Same as EDATA - but allows the definition of ROM constants.
EBIT	00:0020 - 00:007F	Extended bit-addressable RAM; can be accessed with the extended bit addressing mode available on the 251.
XDATA	01:0000 – 01:FFFF (default space)	8051 compatible DATA space. Can be mapped on the 251 to any 64 KB memory segment. Accessed with MOVX instruction.
HDATA	00:0000 – FF:FFFF	Full 16 MB address space of the 251. Accessed with MOV @DRK instructions. This space is used for RAM areas.
HCONST	00:0000 – FF:FFFF	Same as HDATA - but allows the definition of ROM constants.
ECODE	00:0000 – FF:FFFF	Full 16 MB address space of the 251; executable code accessed with ECALL or EJMP instructions.
CODE	FF:0000 - FF:FFFF (default space)	8051 compatible CODE space; used for executable code or RAM constants. Can be located with L251 to any 64 KB segment
CONST	FF:0000 - FF:FFFF (default space)	Same as CODE - but can be used for ROM constants only.

Colons are used to improve the readability only.

The addresses are entered in the tools as numbers without colon.

251 Memory Layout

The following figure shows the memory layout of the 251 architecture.



The 251 completely supports all aspects of the classic 8051 memory layout and instruction set. Existing 8051 programs can be directly execute on the 251. The four 8051 memory spaces (DATA, IDATA, CODE and XDATA) are mapped into specific regions of the 16 MB address space.

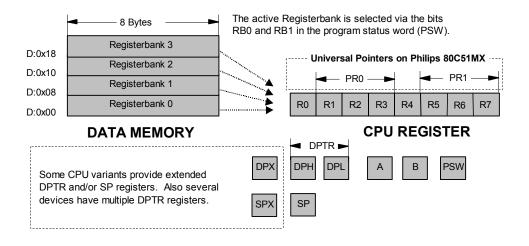
CPU Registers

The following section provides an overview of the CPU registers that are available on the x51 variants.

In addition to the CPU registers R0 - R7, all x51 variants have an SFR space that is used to address on-chip peripherals and I/O ports. In the SFR area also reside the CPU registers SP (stack pointer), PSW (program status word), A (accumulator, accessed via the SFR space as ACC), B, DPL and DPH (16-bit register DPTR).

CPU Registers of the 8051 Variants

The classic 8051 provides 4 register banks of 8 registers each. These register banks are mapped into the DATA memory area at address 0 – 0x1F. In addition the CPU provides a 8-bit A (accumulator) and B register and a 16-bit DPTR (data pointer) for addressing XDATA and CODE memory. These registers are also mapped into the SFR space as special function registers.

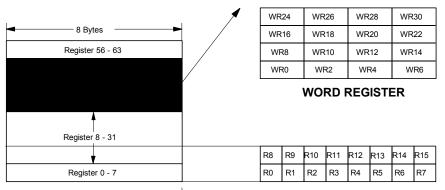


CPU Registers of the Intel/Atmel WM 251

The 251 architecture supports an extra 32 bytes of registers in addition to the 4 banks of 8 registers found in the classic 8051. The lower 8 byte registers are mapped between locations 00:00 - 00:0x1F. The lower 8 byte registers are mapped in this way to support 8051 microcontroller register banking. The register file can be addressed in the following ways:

- Register 0 15 can be used as either byte, word, or double word (Dword) registers.
- Register 16 31 can be addressed as either word or Dword registers.
- Register DR56 and DR60 can be addressed only as Dword registers.
- There are 16 possible byte registers (R0 R15), 16 possible word registers (WR0 WR30) and 10 possible Dword registers (DR0 DR28, DR56 DR60) that can be addressed in any combination.
- All Dword registers are Dword aligned; each is addressed as DRk with "k" being the lowest of the 4 consecutive registers. For example, DR4 consists of registers 4 - 7.
- All word registers are word aligned; each is addressed as WRj with "j" being the lower of the 2 consecutive registers. For example WR4 consists of registers 4 5.
- All byte registers are inherently byte aligned; each is addressed as Rm with "m" being the register number. For example R4 consists of register 4.

The following figure shows the register file format for the 251 microcontroller.



MEMORY

The 8051 CPU Registers A, B, DPL, DPH and SP are mapped into the 251 register file. The

identical with the 8051 CPU registers:

8051 Register 251 Register A (ACC) R11 B R10 DPL low byte of DR56

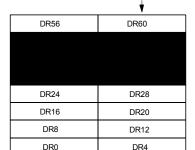
following 251 CPU registers are

DPH 2nd byte of DR56 SP low byte of DR60

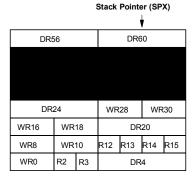
The stack pointer register (DR60) of the 251 CPU is a 16-bit register. DR60 is used for all stack operations (PUSH, POP, CALL, RET, ect.) and can be also used for indirect addressing. Therefore the 251 provides efficient stack addressing modes for reentrant functions.

BYTE REGISTER

Stack Pointer (SPX)



DWORD REGISTER



EXAMPLE OF MIXED USAGE

Program Status Word (PSW)

The Program Status Word (PSW) contains status bits that reflect the current CPU state. The 8051 variants provide one special function register called PSW with this status information. The 251 provides two additional status flags, Z and N, that are available in a second special function register called PSW1.

PSW Register (all 8051 and 251 variants)

Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	Bit 0
CY	AC	F0	RS1	RS0	OV	UD	Р

Additional PSW1 Register (on Intel/Atmel WM 251 only)

Bit	7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	Bit 0
C,	Υ	AC	N	RS1	RS0	OV	Z	-

The following table describes the status bits in the PSW.

	Symbol	Function
	CY	Carry flag
	AC	Auxiliary Carry flag (For BCD Operations)
	F0	Flag 0 (Available to the user for General Purpose)
	RS1, RS0	Register bank select: RS1 RS0 Working Register Bank and Address 0 0 Bank0 (D:0x00 - D:0x07) 0 1 Bank1 (D:0x08 - D:0x0F) 1 0 Bank2 (D:0x10 - D:0x17) 1 1 Bank3 (D:0x18H - D:0x1F)
	ov	Overflow flag
	UD	User definable flag
	P	Parity flag
	-	Reserved for future use
251 ONL)	Z	Zero flag
	N	Negative flag

Instruction Sets

This section lists the instructions of all *x*51 CPU variants in alphabetical order. The following terms are used in the descriptions.

	Identifier	Explanation
	Α	Accumulator
	AB	Register Pair A & B
	В	Multiplication Register
	С	Carry Flag
	DPTR	Data pointer
	PC	Program Counter
	Rn	Register R0 - R7 of the currently selected Register Bank.
	dir8	8-bit data address; Data RAM (D:0x00 - D:0x7F) or SFR (D:0x80 - D:0xFF)
	#data8	8-bit constant included in instruction.
	#data16	16-bit constant included in instruction.
	addr16	16-bit destination address.
	addr11	11-bit destination address. Used by ACALL & AJMP. The branch will be within the same 2KByte block of program memory of the first byte of the following instruction.
	rel	Signed (two's complement) 8-bit offset. Used by SJMP and conditional jumps. Range is -128 +127 bytes relative to the first byte of the following instruction.
	bit8	Direct addressed bit in Data RAM Location.
	Rm	Register R0 - R15 of the currently selected Register File.
	WRj	Register WR0 - WR30 of the currently selected Register File.
	DRk	Register DR0 - DR28, DR56, DR60 of the currently selected Register File.
	dir16	16-bit data address; Data RAM location (00:00 - 00:FFFF).
NLY	@WRj	Data RAM location (0 - 64K) addressed indirectly via WR0 - WR30.
251 ONLY	@DRk	Data RAM location (0 - 16M) addressed indirectly via DR0 - DR28, DR56, DR60.
25	#short	constant 1, 2 or 4 included in instruction.
	bit11	Direct addressed bit in Data RAM or Special Function Register.
	@Wrj+dis	Data RAM location (0 - 64K) addressed indirectly via WR0 - WR30 + displacement.
	@DRk+dis	Data RAM location (0 - 16M) addressed indirectly via DR0 - DR28, DR56, DR60+ 16-bit signed displacement.
	EPTR	23-bit extended data pointer register.
ΓY	PR0, PR1	Universal Pointer Register (PR0 represents R1,R2,R3; PR1 represents R5,R6,R7)
51MX ONLY	@PR0+d2 @PR1+d2	Universal memory location (0 - 16M) addressed indirectly via PR0 or PR1+ 2-bit displacement (+0, +1, +2, +3).
51	#data2	2-bit constant included in instruction (value range: #1, #2, #3, #4).
	addr23	23-bit destination address for HDATA or ECODE

ACALL		Al	bsolute Subroutine CALL CY -			AC N		z
Mnemonic			Description			ytes nary	Byt Sou	
ACALL	addr11		Absolute Subroutine Call			2	2	
DAL		AL	LAS 24-Bit Contiguous Address Mode ONI		Υ			
ACALL	addr19		Absolute Subroutine Call			3		

	ADE)		DD destination, source ddition	CY X	AC X	N X	ov X	Z X		
	Mnemo	nic		Description	-	tes nary	Byte Sou				
	ADD	A,Rn		Add register to accumulator			1	2			
	ADD	A,dir8		Add direct byte to accumulator			2	2			
	ADD	A,@Ri		Add indirect RAM to accumulator			1	2			
	ADD	DD A,#data8		Add immediate data to accumulator			2	2			
	ADD	Rm,Rm		Rm,Rm Add byte register to byte register					3	2	
	ADD	WRj,WRj		Add word register to word register			3				
	ADD	DRk,DRk		Add double word register to dword regis	ter		3	2			
	ADD	Rm,#data8		Add 8 bit data to byte register			4	3			
	ADD	Wrj,#data16		Add 16 bit data to word register			5	4			
ONLY	ADD	Drk,#data16		Add 16 bit unsigned data to dword register			5	4			
251 (ADD	Rm,dir8		Add direct address to byte register			4	3			
2	ADD	WRj,dir8		Add direct address to word register			4	3			
	ADD	Rm,dir16		Add direct address (64K) to byte registe	er		5	4			
	ADD	WRj,dir16		Add direct address (64K) to word regist	er		5	4			
	ADD	Rm,@WRj		Add indirect address (64K) to byte regis	ter		4	3			
	ADD			Add indirect address (16M) to byte regis	Add indirect address (16M) to byte register		4	3			
51	ADD	PR0,#data2	Add immediate data to PR0			2					
MX51	ADD	PR1,#data2		Add immediate data to PR1		2					

		ADDC destination, source Addition with Carry	CY X	AC X	N X	ov X	Z X
Mnemonic		Description	Description			Bytes Source	
ADDC A,Rn		Add register to accumulator with carry fl	Add register to accumulator with carry flag			2	
ADDC	A,dir8	Add direct byte to accumulator with carr	Add direct byte to accumulator with carry flag			2	
ADDC	A,@Ri	Add indirect RAM to accumulator with c flag	Add indirect RAM to accumulator with carry flag			2	
ADDC	A,#data8	Add immediate data to accumulator with flag	n carry		2	2	

AJMP		Al	bsolute JUMP CY			<u>N</u>	ov —	z
Mnemo	Mnemonic		Description			Bytes Binary		es rce
AJMP	AJMP addr11		Absolute Jump			2	2	
DAL		AL	LAS 24-Bit Contiguous Address Mode	ONL	Y			
AJMP	addr19		Absolute Jump			3		

		AND destination, source Logical AND	CY —	AC —	N X	ov —	Z X
Mnemonic		Description			tes nary	Byte Sou	
ANL	A,Rn	AND register to accumulator			1	2	
ANL	A,dir8	AND direct byte to accumulator			2	2	
ANL	A,@Ri	AND indirect RAM to accumulator			1	2	
ANL	A,#data8	AND immediate data to accumulator			2	2	
ANL	dir,A	AND accumulator to direct byte			2	2	
ANL dir,#data8		AND immediate data to direct byte			3	3	

	ANL	-		AND destination, source CY Logical AND —			N X	ov —	Z X
	Mnemonic			Description	Description		Bytes Binary		es rce
	ANL Rm,Rm ANL WRj,WRj			AND byte register to byte register		3	2		
				AND word register to word register		3			
	ANL	Rm,#data8		AND 8 bit data to byte register			4	3	
	ANL	Wrj,#data16		AND 16 bit data to word register			5	4	
ONLY	ANL	Rm,dir8		AND direct address to byte register			4	3	
251 (ANL	Wrj,dir8		AND direct address to word register			4		
2	ANL	Rm,dir16		AND direct address (64K) to byte regist	er		5	4	
	ANL	ANL Wrj,dir16		AND direct address (64K) to word regis	ter		5	4	
	ANL Rm,@WRj			AND indirect address (64K) to byte register			4	3	
	ANL	Rm,@DRk		AND indirect address (16M) to byte regi		4			

ANL	_	ANL destination, source Logical AND for bit variables	CY X	AC —	N —	ov —	z
Mnemo	onic	Description	Description		tes nary	Byt Sou	
ANL	C,bit8	AND direct bit to carry; from BIT space			2	2	
		Intel/Atmel WM 251 ONLY					
ANL	C,bit11	AND direct bit to carry; from EBIT space	9		4	3	

/ / / / / /		ANL/ destination, source CY A Logical AND for bit variables X			AC N		Z X
Mnemonic		Description			Bytes Binary		es rce
ANL	C,/bit8	AND complement of dir bit to carry; BIT	space		2	2	
		Intel/Atmel WM 251 ONLY					
ANL C,/bit11		AND complement of dir bit to carry; EBIT space		4		3	

		COMPARE destination, source and jump if not equal	CY X	AC —	N X	ov —	Z X
Mnemo	onic	Description		_	tes nary	Byte Soul	
CJNE	A,dir8,rel	Compare dir byte to acc. and jump if no	t equal		3	3	
CJNE	/A,#data8,rel	Compare imm. data to acc. and jump if equal	not		3	3	
CJNE	Rn,#data8,rel	Compare imm. data to reg and jump if r equal	not		3	4	
CJNE	@Ri,#data8,re	Compare imm. data to indir and jump if equal	not		3	4	

CLF	₹	CLEAR Operand	CY —	AC —	N X	ov —	Z X
Mnemonic		Description		_	tes nary	Byte Sou	
CLR A Clear accumulator				1	1		

CLF	₹	CLEAR Bit Operand	CY —	AC —	N —	ov —	z _
Mnemo	onic	Description		Bytes Binary		Byte Sou	
CLR	С	Clear carry			1	1	
CLR	bit8	Clear direct bit from BIT space			2	2	
		Intel/Atmel WM 251 ONLY					
CLR	bit11	Clear direct bit from EBIT space			4	3	

	СМІ)	C	OMPARE Operands	CY X	AC X	N X	ov X	Z X			
	Mnemo	nic		Description		,	tes ary	Byte				
	CMP	Rm,Rm		Compare registers	mpare registers							
	CMP	WRj,WRj		Compare word registers		;	3	2				
	CMP	DRk,DRk		Compare double word registers			3	2				
NLY	CMP	Rm,#data8		Compare register with immediate data			4	3				
251 ONLY	CMP	Wrj,#data16		Compare word register with immediate data			5	4				
25	CMP	Drk,#00		Compare dword reg with zero extended data			5	4				
	CMP	Drk,#ff		Compare dword reg with one extended	data		5	4				
	CMP	Rm,dir8		Compare register with direct byte		4		3				
	CMP	WRj,dir8		Compare word register with direct word		4	4	3				
	CMP	Rm,dir16		Compar register with direct byte (64K)	5		5	4				
	CMP	WRj,dir16		Compare word register with direct word	ord register with direct word (64K)			(64K) 5		5	4	
	CMP	Rm,@WRj		Compare register with indirect address	register with indirect address (64K)			3				
	CMP	Rm,@DRk		Compare register with indirect address	(16M)	4	4	3				

CPL	-	COMPLEMENT Operand	CY —	AC —	N X	ov —	Z X
Mnemo	nic	Description			tes nary	Byte Sou	
CPL	Α	Complement accumulator			1	1	

CPL	CPL c		FBit Operand	CY —	AC —	N —	ov —	z —
Mnemo	nic	Description			_	tes nary	Byt Sou	
CPL	С	Complement	carry			1	1	
CPL	bit8	Complement	direct bit from BIT space			2	2	
		Intel/A	tmel WM 251 ONLY					
CPL	bit11	Complement	direct bit from EBIT space	•	T	4	3	

1 1 2 2			ECIMAL ADJUST Accumulator or Addition	CY X	AC —	N X	ov —	Z X
Mnemo	Mnemonic		Description		, ,	tes nary	Byt Sou	
DA	DA A Decimal adjust accumulator		1		1			

		DECREMENT Operand with a constant	CY —	AC —	N X	ov —	Z X
Mnemo	nic	Description		_	tes nary	Byte Sou	
DEC	Α	Decrement accumulator			1	1	
DEC	Rn	Decrement register			1	2	
DEC	dir	Decrement dir byte			2	2	
DEC	@Ri	Decrement indir RAM			1	2	
		Intel/Atmel WM 251 ONLY					
DEC	Rm,#short	Decrement byte register with 1, 2 or 4			3	2	
DEC	WRj,#short	Decrement word register with 1, 2 or 4		3		2	
DEC	DRk,#short	Decrement double word register with 1, 2	or 4		3	2	

DIV		DIVIDE Operands		AC —	N X		Z X		
Mnemo	nic	Description		Bytes Binary				Byt Sou	
DIV	AB	Divide A by B			1	1			
		Intel/Atmel WM 251 ONLY							
DIV	Rm,Rm	Divide byte register by byte register			3	2			
DIV WRj,WRj		Divide word register by word register			3	2			

		DECREMENT Operand and Jump if Not Zero	CY —	AC —	N X	ov —	Z X
Mnemonic		Description	Description		tes nary	Byte Sou	
DJNZ Rn,rel		Decrement register and jump if not zero	Decrement register and jump if not zero		2	3	
DJNZ dir8,rel Decre		Decrement direct byte and jump if not z	crement direct byte and jump if not zero		3	3	

ECA	ALL	Extended Subroutine CALL	CY —	AC —	N —	ov —	Z
Mnemo	nic	Description		Bytes Binary		Byt Sou	
		Intel/Atmel WM 251 ONLY					
ECALL	addr24	Extended subroutine call			5	4	
ECALL	DRk	Extended subroutine call			3	2	
		Philips 80C51MX ONLY					
ECALL	addr23	Extended subroutine Call			5		

EJN	1 P	Extended JUMP	CY —	AC —	N —	ov —	Z
Mnemo	nic	Description			tes nary	Byte	
		Intel/Atmel WM 251 ONLY					
EJMP	addr24	Extended jump			5	4	
EJMP	DRk	Extended jump			3	2	
		Philips 80C51MX ONLY					
EJMP	addr23	Extended jump			5		

			OV destination, source ove data via Unversal Pointer	CY —	AC —	N —	ov —	z _
CONLY			Description		_	tes nary	Byte	
51MX	EMOV	A,@PR0+d2	Move indirect (16M) via Universal Point	er to A		2		
2	EMOV	A,@PR1+d2	Move indirect (16M) via Universal Point	er to A		2		
	EMOV	@PR0+d2,A	Move A to indirect (16M) via Universal F	Pointer		2		
	EMOV	@PR1+d2,A	Move A to indirect (16M) via Universal F	Pointer		2		

(ONLY	8 EKE I			FURN from extended proutine	CY —	AC —	N —	ov —	Z
251& 51MX	Mnemo	nic	D	Description			tes nary	Byte	
25	ERET		R	Return from subroutine			2	1	

				CREMENT Operand with a onstant	CY —	AC —	N X	ov —	Z X			
	Mnemonic			Description	escription		Bytes Binary		es rce			
	INC	Α		Increment accumulator			1	1				
	INC	Rn		Increment register	register					1	2	
	INC	dir		Increment direct byte			2	2				
	INC	@Ri		ncrement indirect RAM			1	2				
	INC	DPTR		Increment Data Pointer	ent Data Pointer		er		1	1		
ΙΓΥ	INC	Rm,#short		Increment byte register with 1, 2 or 4			3	2				
I ONLY	INC	WRj,#short		Increment word register with 1, 2 or 4			3	2				
251	INC	Drk,#short		Increment double word register with 1, 2	ister with 1, 2 or 4		3	2				
				Philips 80C51MX ONLY								
	INC	NC EPTR Increment Enhanced Data Pointer			2							

JB	JB JI		JMP if Bit is set		AC —	<u>N</u>	ov —	Z
Mnemonic			Description			ytes nary	Byt Sou	
JB	bit8,rel		Jump if dir bit (from BIT space) is set			3	3	
			Intel/Atmel WM 251 ONLY					
JB	bit11,rel		Jump if dir bit (from EBIT space) is set			5	4	

JBC		JL	UMP if Bit is set and clear bit CY -		AC —	N —	ov —	z —
Mnemonic			Description		_	tes nary	Byt Sou	
JBC	JBC bit8,rel		Jump if dir bit (BIT space) is set and clear bit			3	3	
	Intel/Atmel WM 251 ONLY							
JBC	JBC bit11,rel		Jump if dir bit (EBIT space) is set and clear bit			5	4	

JC /	JL	JUMP if Carry is set JUMP if less than	CY —	AC —	N —	ov —	Z
Mnemonic		Description			/tes nary	Byt Sou	
JC	rel	Jump if carry is set			2	2	
		Intel/Atmel WM 251 ONLY					
JL	rel	Jump if less than			2	2	

	> J⊏		Jl	JMP if equal	CY —	AC —	N —	ov —	<u>z</u>
251 ON	Mnemonic			Description			tes nary	Byte	
	JE	rel		Jump if equal			3	2	

	JG		Jl	JMP if greater than	CY —	AC —	N —	ov —	z
251 ON	Mnemonic			Description		_	tes nary	Byte	
	JG	rel		Jump if greater than		;	3	2	

ONLY	JLE		JUMP if less than or equal	CY —	AC —	N —	ov —	z
251 ON	Mnemo	nic	Description			tes nary	Byte	
	JLE	rel	Jump if less than or equal			3	2	

JMP		Jl	JMP indir relative to DPTR	CY —	AC —	N —	ov —	z —
Mnemonic			Description			tes nary	Byte Sou	
JMP	@A+DPTR	Jump indir relative to DPTR				1	1	
			Philips 80C51MX ONLY					
JMP	@A+EPTR		JUMP indirect relative to EPTR			2		

JNB		JL	JMP if Bit is Not set	CY —	AC —	<u>N</u>	ov —	z
Mnemonic			Description		_	tes nary	Byt Sou	
JNB	bit8,rel		Jump if dir bit (from BIT space) is not set			3	3	
	Intel/Atmel WM 251 ONLY							
JNB	bit11,rel		Jump if dir bit (from EBIT space) is not set			5	4	

JNC	/JGE	JUMP if Carry is Not set JUMP if greater than or equal	CY —	AC —	N —	ov —	z —
Mnemonic		Description			tes nary	Byte Sou	
JNC	rel	Jump if carry is not set			2	2	
		Intel/Atmel WM 251 ONLY					
JGE	rel	Jump if greater than or equal			2		

ONLY	JNE		Jl	JMP if Not Equal	CY —	AC —	N —	ov —	z
251 ON	Mnemo	Mnemonic		Description			tes nary	Byte	
	JNE	rel		Jump if not equal			3	2	

JNZ		JUMP if Accumulator is Not Zero	CY —	AC —	N —	ov —	z _
Mnemo	onic	Description			tes nary	Byte Sou	
JNZ rel Jump if accumulator is not zero		rel Jump if accumulator is not zero			2	2	

4	JSG		Jl	JMP if greater than (Signed)	CY —	AC —	N —	ov —	Z
251 ON	Mnemonic			Description		,	tes nary	Byte	
	JSG	rel		Jump if greater than (signed)			3	2	

-	131-6		JMP if greater than or Equal igned)	CY —	AC —	N —	ov —	z
251 ON	Mnemo	nic	Description			tes ary	Byte	
	JSGE	rel	Jump if greater than or equal (signed)			3	2	

	JSL		JL	JMP if Less than (Signed)	CY —	AC —	N —	ov —	<u>z</u>
251 ON	Mnemonic			Description			tes ary	Byte	
	JSL	rel		Jump if less than (signed)			3	2	

	JOLE (S		JMP if Less than or Equal igned)	CY —	AC —	N —	ov —	z
251 ON	Mnemonic Mnemonic		Description		,	tes nary	Byte	
	JSLE	rel	Jump if less than or equal (signed)			3	2	

JZ		JUMP if Accumulator is Zero	CY —	AC —	N —	ov —	z —
Mnemo	nic	Description			tes nary	Byt Sou	
JZ rel Jump if accumulator is zero				2	2		

LCA	ALL	Long Subroutine CALL CY A			N —	ov —	z —		
Mnemo	nic	Description			Bytes Binary		-		es rce
LCALL	addr16	Long Subroutine Call			3	3			
		Intel/Atmel WM 251 ONLY							
LJMP	@WRj	Long Jump indirect via word register			3	2			
	D	ALLAS 24-Bit Contiguous Address Mod	e ONL	Υ					
LCALL	addr24	Absolute Subroutine Call			4				

LJN	I P	Long JUMP CY —			<u>N</u>	ov —	z
Mnemo	nic	Description	escription			Byte Sou	
LJMP	addr16	Long Jump			3		
		Intel/Atmel WM 251 ONLY					
LJMP	@WRj	Long Jump indirect via word register			3		
	DALLAS 24-Bit Contiguous Address Mode ONL						
LJMP	addr24	Absolute Subroutine Call			4		

MC	V	MOV destination, source Move data	CY —	AC —	N —	ov —	Z																	
Mnem	nonic	Description		_	tes nary	Byte Soul																		
MOV	A,Rn	Move register to accumulator			1	2																		
MOV	A,dir8	Move direct byte to accumulator			2	2																		
MOV	A,@Ri	Move indirect RAM to accumulator			1	2																		
MOV	A,#data8	Move immediate data to accumulator			2	2																		
MOV	Rn,A	Move accumulator to register			1	2																		
MOV	Rn,dir8	Move direct byte to register			2	3																		
MOV	Rn,#data8	Move immediate data to register			2	3																		
MOV	dir8,A	Move accumulator to direct byte			2	2																		
MOV	dir8,Rn	Move register to direct byte			2	3																		
MOV	dir8,dir8	Move direct byte to direct byte		;	3	3																		
MOV	dir8,@Ri	Move indirect RAM to direct byte		2		2		3																
MOV	dir8,#data8	Move immediate data to direct byte		;	3	3																		
MOV	@Ri,A	Move accumulator to indirect RAM			1	2																		
MOV	@Ri,dir8	Move direct byte to indirect RAM			2	3																		
MOV	@Ri,#data8	Move immediate data to indirect RAM		:	2	3																		
MOV	DPTR,#data10	Load Data Pointer with 16-bit constant		;	3	3																		
MOV	C,bit8	Move dir bit to carry			2	2																		
MOV	bit8,C	Move carry to dir bit			2																			
MOV	Rm,Rm	Move byte register to byte register		:	3	2																		
MOV	WRj,WRj	Move word register to word register			3	2																		
MOV	DRk,DRk	Move dword register to dword register			3	2																		
MOV	Rm,#data8	Move 8 bit data to byte register			4	3																		
MOV	WRj,#data16	Move 16 bit data to word register			5	4																		
MOV	DRk,#0data16	Move 16 bit zero extended data to dword	reg.		5	4																		
MOV	DRk,#1data16	Move 16 bit one extended data to dword	reg.		5	4																		
MOV	Rm,dir8	Move dir address to byte register			4	3																		
MOV	WRj,dir8	Move direct address to word register			register 4	ss to word register 4	4	4	4	4	4	4	4	4		4	4	4	4	4	4	4	3	
MOV	DRk,dir8	Move direct address to dword register			4	3																		
MOV	Rm,dir16	Move direct address (64K) to byte registe	er		5	4																		
MOV	WRj,dir16	Move direct address (64K) to word registe			5	4																		
MOV	DRk,dir16	Move direct address (64K) to dword regis			5	4																		
MOV	Rm,@WRj	Move indirect address (64K) to byte regis			4	3																		
MOV	Rm,@DRk	Move indirect address (16M) to byte regis		- -	4	3																		
	WRj,@WRj	Move indirect address (64K) to word regis		-	4	3																		
MOV		intro man out addition (o m) to mora region																						

МО	V	MOV destination, source Move data	CY —	AC —	<u>N</u>	ov —	<u>z</u>
Mnemo	onic	Description			/tes nary	Byt Sou	
MOV	dir8,Rm	Move byte register to direct address			4	3	
MOV	dir8,WRj	Move word register to direct address			4	3	
MOV	dir8,DRk	Move dword register to direct address			4	3	
MOV	dir16,Rm	Move byte register to direct address (64	łK)		5	4	
MOV	dir16,WRj	Move word register to direct address (6	4K)		5	4	
MOV	dir16,DRk	Move dword register to direct address (64K)		5	4	
MOV	@WRj,Rm	Move byte register to direct address (64	ss (64K)		4	3	
MOV	@DRk,Rm	Move byte register to indirect address (16M)		4		
MOV	@WRj,WRj	Move word register to indirect address	(64K)		4		
MOV	@DRk,WRj	Move word register to indirect address	(16M)		4		
MOV	Rm,@WRj+dis	Move displacement address (64K) to by	te reg.		5	4	
MOV	WRj,@WRj+di	s Move displacement address (64K) to w	ord reg		5	4	
MOV	Rm,@DRk+dis	Move displacement address (16M) to b	yte reg.		5	4	
MOV	WRj,@DRk+di	s Move displacement address (16M) to w reg.	ord		5	4	
MOV	@WRj+dis,Rm	Move byte reg. to displacement addres	s (64K)		5	4	
MOV	@WRj+dis,WF	Rj Move word reg. to displacement addres	s (64K))	5	4	
MOV	@DRk+dis,Rm	Move byte reg. to displacement addres	s (16M)		5	4	
MOV	@DRk+dis,WF	Move word reg. to displacement addres (16M)	S		5	4	
MOV	C,bit11	Move dir bit from 8 bit address location carry	to		2	2	
MOV	bit11,C	Move carry to dir bit from 16 bit address location	S		5	4	
		Philips 80C51MX ONLY					
MOV	EPTR,#adr23	Load extended data pointer with consta	nt		5		

MO	VC	- Indive Code Byte —			N —	ov —	z —
Mnemo	nic	Description			tes nary	Byte Sou	
MOVC	A,@A+DPTR	Move code byte relative to DPTR to accumulator			1	1	
MOVC	A,@A+PC	Move code byte relative to PC to accum	nulator		1	1	
		Philips 80C51MX ONLY					
MOVC	A,@A+EPTR	Move code byte relative to EPTR to accumulator			2		

Υ.	MO	VH		OVH destination, source ove data to high word of DR	CY —	AC —	N —	ov —	<u>z</u>
251 ONL	Willemonic			Description			tes nary	Byte	
	MOVH DRk,#data16			Move 16bit imm. data to high word of di reg.	word		3	2	

NLY			MOVS destination, source Move byte to word (signed ext.)		AC —	N —	ov —	z
251 ON	Mnemonic		Description		,	tes nary	Byte	
	MOVS WRj,Rm		Move byte register to word register			3	2	

MO	VX	MOV destination, source External RAM access				
Mnemo	nic	Description		Bytes Binary	Bytes Source	
MOVX	VX A,@Ri Move xdata RAM (8 bit address) to accumulator				2	
MOVX	A,@DPTR	Move xdata RAM (16 bit address) to accumulator		1	1	
MOVX	@Ri,A	Move accumulator to xdata RAM (8 bit address)		1	2	
MOVX	@DPTR,A	Move accumulator to xdata RAM (16 bit address)		1	1	

IV/I		MOV destination, source External RAM access	CY —	AC —	<u>N</u>	ov —	z —
Mnemonic		Description		-	/tes nary	Byt Sou	
		Philips 80C51MX ONLY					_
MOVX	@EPTR,A	Move accu to xdata RAM (23 bit addres	Move accu to xdata RAM (23 bit address)		2		
MOVX	A,@EPTR	Move xdata RAM (23 bit address) to accu			2		_

ONLY			OV destination, source ove byte to word (zero ext.)	CY —	AC —	N —	ov —	<u>z</u>
251 ON	Mnemonic		Description		,	tes nary	Byte	
	MOVZ WRj,Rm		Move byte reg. to word reg. (zero exten	ided)		3	2	

MU	L	MULTIPLY Operands	CY 0	AC —	N X	ov X	Z X
Mnemo	onic	Description			tes nary	Byt Sou	
MUL	AB	Multiply A and B			1	1	
		Intel/Atmel WM 251 ONLY					
MUL	Rm,Rm	Multiply byte register with byte register	Multiply byte register with byte register			2	
MUL	WRj,WRj	Multiply word register with word register			3	2	

NOF	•	No Operation	CY —	AC —	<u>N</u>	ov —	z
Mnemo	nic	Description	·		tes nary	Byt Sou	
NOP		No operation			1	1	

	ORI	-	ORL destination, source Logical OR	CY —	AC —	N X	ov —	Z X
	Mnemo	nic	Description		_	tes nary	Byte	
	ORL	A,Rn	OR register to accumulator			1	2	
	ORL	A,dir8	OR dir byte to accumulator	dir byte to accumulator		2	2	
	ORL	A,@Ri	OR indir RAM to accumulator			1	2	
	ORL	A,#data8 OR immediate data to accumulator			2			
	ORL	dir,A	,A OR accumulator to dir byte			2		
	ORL	dir,#data8	OR immediate data to dir byte			3		
	ORL	Rm,Rm	OR byte register to byte register			3	2	
	ORL	WRj,WRj	OR word register to word register			3	2	
	ORL	Rm,#data8	OR 8 bit data to byte register			4	3	
	ORL	WRj,#data16	OR 16 bit data to word register			5	4	
251 ONLY	ORL	Rm,dir8	OR dir address to byte register			4	3	
51 (ORL	WRj,dir8	OR dir address to word register			4	3	
2	ORL	Rm,dir16	OR dir address (64K) to byte register			5	4	
	ORL	WRj,dir16	OR dir address (64K) to word register			5	4	
	ORL	Rm,@WRj	OR indir address (64K) to byte register			4	3	
	ORL	Rm,@DRk	OR indir address (16M) to byte register			4	3	

ORI	<u>_</u>		tination, source OR for bit varial		CY X	AC —	N —	ov —	z —
Mnemo	onic	Descri	ption			_	tes nary	Byt Sou	
ORL	C,bit8	OR dire	OR direct bit to carry; from BIT space				2	2	
			Intel/Atmel WM 251 ONLY						
ORL	C,bit11	OR dire	OR direct bit to carry; from EBIT space				4	3	

OR	L/	ORL/ destination, source Logical OR with Complement	CY X	AC —	N X	ov —	Z X
Mnemo	onic	Description		_	tes nary	Byt Sou	
ORL	C,/bit8	OR complement of direct bit to carry; B space	IT		2	2	
	Intel/Atmel WM 251 ONLY						
ORL	ORL C,/bit11 OR complement of dir bit to carry; EBIT space				4	3	

	POP		POP Operand from Stack CY A — -		AC —	N —	ov —	z _
	Mnemonic		Description	Description			Bytes Source	
	POP	dir8	Pop direct byte from stack			2	2	
LΥ	POP	Rm	Pop byte register from stack			3	2	
J ONE	POP	WRj	Pop word register from stack			3	2	
7 21	POP	DRk	Pop double word register from stack			3	2	

	PUS	SH	PI	JSH Operand onto Stack	CY —	AC N	ov z – –
	Mnemonic			Description	Bytes Binary	Bytes Source	
	PUSH	dir8		Push direct byte onto stack		2	2
	PUSH	Rm		Push byte register onto stack		3	2
ΙLΥ	PUSH	WRj		Push word register onto stack		3	2
251 ONLY	PUSH	DRk		Push double word register onto stack		3	2
25	PUSH	#data8		Push immediate data onto stack		4	3
	PUSH	#data16		Push immediate data (16 bit) onto stack	(5	4

RET '		RETURN from Subroutine	CY —	AC —	AC N		z _
Mnemonic		Description			tes nary	Byte Sou	
RET		Return from subroutine			1	1	

RETI		RETURN from Interrupt	CY —	AC —	AC N		z
Mnemonic		Description		Bytes Binary		Byt Sou	
RETI		Return from interrupt			Bytes Binary		

RL R		ROTATE Accumulator Left	CY —	AC —	N X	ov —	Z X
Mnemonic		Description				Byte Sou	
RL	RL A Rotate accumulator left			1		1	

RLC	;	ROTATE Accumulator Left through the Carry	CY X	AC —	N X	ov —	Z X
Mnemonic		Description			tes nary	Byte Sou	
RLC A Rotate accumulator left through the carry				1		1	

RR		ROTATE Accumulator Right	CY —	AC —	AC N		Z X
Mnemonic		Description	Bytes Binary			Byte Sou	
RR A		Rotate accumulator right			1	1	

		ROTATE Accumulator Right through the Carry	CY X	AC —	N X	ov —	Z X
Mnemonic		Description			tes nary	Byte Sou	
RRC A		Rotate accumulator right through the c	arry		1	1	

SETB		SET Bit Operand	CY —	AC —	N —	ov —	z —
Mnemonic		Description		By Bir			
SETB	С	Set carry			1	1	
SETB	bit8	Set direct bit from BIT space			2		
		Intel/Atmel WM 251 ONLY					
SETB	bit11	Set direct bit from EBIT space			5	4	

SJMP		Short JUMP		CY —	AC —	N —	ov —	Z
Mnemonic		Description			Bytes Binary		Byt Sou	
SJMP	SJMP rel Short jump (relative address)			2		2		

	SLL		SI	HIFT Register Left	CY X	AC —	N X	ov —	Z X
251 ONL				Description		Bytes Binary		Byte	
	SLL	Rm		Shift byte register left			3	2	
	SLL	WRj		Shift word register left			3	2	

>				HIFT Register Right (arithmet.) gn extended	CY X	AC —	N X	ov —	Z X
251 ONL				Description			tes nary	Byte Soul	
	SRA	Rm		Shift byte register right; sign extended		:	3	2	
	SRA WRj		Shift word register right; sign extended			3	2		

>			HIFT Register Right (logic) zero ktended	CY X	AC —	N X	ov —	Z X
251 ONL			Description		,	tes ary	Byte	
	SRL	Rm	Shift byte register right; zero extended		:	3	2	
	SRL	WRj	Shift word register right; zero extended		;	3	2	

_	SUE	3		JB destination, source ubtraction	CY X	AC X	N X	ov X	Z X
	Mnemo	nic		Description			tes nary	Byte Soul	
	SUB	,		Subtract byte register from byte register	te register from byte register			2	
	SUB	WRj,WRj		Subtract word register from word register	er		3	2	
	SUB	DRk,DRk		Subtract dword register from dregister	word register from dregister			2	
ONLY	SUB	Rm,#data		Subtract 8 bit data from byte register	oyte register			3	
Q	SUB	Wrj,#data16		Subtract 16 bit data from word register			5	4	
251	SUB	Drk,#data16		Subtract 16 bit unsigned data from dword reg.			5	4	
	SUB	Rm,dir		Subtract direct address from byte register			4	3	
	SUB	Wrj,dir		Subtract direct address from word regis	ter		4	3	
	SUB	Rm,dir16		Subtract direct address (64K) from byte register		5		4	
	SUB	Wrj,dir16		Subtract direct address (64K) from word register	d		5	4	
	SUB	Rm,@WRj		Subtract indirect address (64K) from by	te reg.		4	3	
	SUB	Rm,@DRk		Subtract indirect address (16M) from by	te reg.		4	3	

SUE	3B	SUBB destination, source Subtraction with Borrow	CY X	AC X	N X	ov X	Z X
Mnemo	nic	Description		-	tes nary	Byt Sou	
SUBB	A,Rn	Subtract register from accumulator with	borrov	v	1	2	
SUBB	A,dir8	Subtract direct byte from accumulator v borrow	vith		2	2	
SUBB	A,@Ri	Subtract indirect byte from accumulator borrow	with		1	2	
SUBB	A,#data8	Subtract immediate data from accumula with borrow	ator		2	2	

SWAP Nibbles within the Accumulator		CY —	AC —	N X	ov —	Z X	
Mnemo	nic	Description			tes nary	Byte Soul	
SWAP	Α	Swap nibbles within the accumulator			1	1	

- ITA	TRA	\ P	JL	JMP to the Trap Interrupt	CY —	AC —	N —	ov —	z
251 ON	Mnemo	nic		Description		,	tes nary	Byte	
	TRAP			Jumps to the trap interrupt vector			2	1	

XCI	Н	E	XCHANGE Operands	CY —	AC —	N —	ov —	<u>z</u>
Mnemo	onic		Description			tes nary	Byt Sou	
XCH	A,Rn		Exchange register with accumulator			2	2	
XCH	A,dir8		Exchange direct byte with accumulator			2	2	
XCH	A,@Ri		Exchange indirect byte with accumulate	or		1	2	

XCH	1D	EXCHANGE Digit	CY —	AC —	N —	ov —	Z
Mnemo	nic	Description		_	/tes nary	Byt Sou	
XCHD	A,@Ri	Exchange low-order digit in indir. RAM accumulator	with		1	2	

	XRL	•	KCLOR destination, source orgical Exclusive-OR	CY —	AC —	N X	ov —	Z X
	Mnemo	nic	Description			ytes nary	Byte Soul	
	XRL	A,Rn	Exclusive-OR register to accumulator			1	2	
	XRL	A,dir8	Exclusive-OR direct byte to accumulate	r		2	2	
	XRL	A,@Ri	Exclusive-OR indirect byte to accumula	tor		1	2	
	XRL	A,#data8	Exclusive-OR immediate data to accum	ulator		2	2	
	XRL	dir8,A	Exclusive-OR accumulator to direct byte	е		2	2	
	XRL	dir8,#data8	Exclusive-OR immediate data to direct	byte		3	3	
	XRL	Rm,Rm	Exclusive-OR byte register to byte regis	ter	_	3	2	
	XRL	WRj,WRj	Exclusive-OR word register to word reg	ister		3	2	
	XRL	Rm,#data8	Exclusive-OR 8 bit data to byte register			4	3	
	XRL	WRj,#data16	Exclusive-OR 16 bit data to word registe	er		5	4	
	XRL	Rm,dir8	Exclusive-OR direct address to byte reg	gister		4	3	
NL)	XRL	WRj,dir8	Exclusive-OR direct address to word re	gister		4	3	
251 ONLY	XRL	Rm,dir16	Exclusive-OR direct address (64K) to be	yte reg.		5	4	
2	XRL	WRj,dir16	Exclusive-OR direct address (64K) to w reg.	ord		5	4	
	XRL	Rm,@WRj	Exclusive-OR indirect address (64K) to reg.	byte		4	3	
	XRL	Rm,@DRk	Exclusive-OR indirect address (16M) to reg.	byte		4	3	

Opcode Map

The following opcode maps provide an overview of the instruction encoding for the 8051, the 80C51MX, and the 251 architecture. It is arranged as separate maps as described below:

8051 Instructions: these opcode are available on all *x***51** variants. Both the Philips 80C51MX and the Intel/Atmel WM 251 use an **OPCODE PREFIX** byte with the encoding A5 to extend the classic 8051 instruction set. The additional 251 an 80C51MX instructions are described in the following tables.

Additional 251 Instructions: if the 251 is configured in binary mode the **8051 instructions** are the default opcode map and the **OPCODE PREFIX** is the first opcode byte for the additional 251 instructions. If the 251 is configured in source mode the additional 251 instructions are the default opcode map and the **OPCODE PREFIX** is the first op-code byte when the 251 should execute standard 8051 instructions that are encode with the byte values x6-xF.

Additional 80C51MX Instructions via Prefix A5: contains the 80C51MX instructions that require the **OPCODE PREFIX** byte. The Philips 80C51MX provides instructions for addressing the 16MB address space and the extended SFR area that are listed in this table.

2

8051 Instructions

Binary Mode	x0	x1	x2	х3	x4	х5	x6-x7	x8-xF
Source Mode	x0	x1	x2	х3	x4	x5	A5x6-A5x7	A5x8-A5xF
0x	NOP	AJMP adr11	LJMP adr16	RR A	INC A	INC dir	INC @Ri	INC Rn
1x	JBC	ACALL	LCALL	RRC	DEC	DEC	DEC	DEC
	bit,rel	adr11	adr16	A	A	dir	@Ri	Rn
2x	JB bit,rel	AJMP adr11	RET	RL A	ADD A,#data	ADD A,dir	ADD A,@Ri	ADD A,Rn
3x	JNB bit,rel	ACALL adr11	RETI	RLC A	ADDC A,#data	ADDC A,dir	ADDC A,@Ri	ADDC A,Rn
4x	JC	AJMP	ORL	ORL	ORL	ORL	ORL	ORL
	rel	adr11	dir,A	dir,#data	A,#data	A,dir	A,@Ri	A,Rn
5x	JNC	ACALL	ANL	ANL	ANL	ANL	ANL	ANL
	rel	adr11	dir,A	dir,#data	A,#data	A,dir	A,@Ri	A,Rn
6x	JZ	AJMP	XRL	XRL	XRL	XRL	XRL	XRL
	rel	adr11	dir,A	dir,#data	A,#data	A,dir	A,@Ri	A,Rn
7x	JNZ	ACALL	ORL	JMP	MOV	MOV	MOV	MOV
	rel	adr11	c,bit	@A+DPTR	A,#data	dir,#data	@Ri,#data	Rn,#data
8x	SJMP	AJMP	ANL	MOVC	DIV	MOV	MOV	MOV
	rel	adr11	C,bit	A,@A+PC	AB	dir,dir	dir,@Ri	dir,Rn
9x	MOV	ACALL	MOV	MOVC	SUBB	SUBB	SUBB	SUBB
	DPTR,#d16	adr11	bit,c	A,@A+DPTR	A,#data	A,dir	A,@Ri	A,Rn
Ax	ORL	AJMP	MOV	INC	MUL	OPCODE	MOV	MOV
	C,/bit	adr11	C,bit	DPTR	AB	PREFIX	@Ri,dir	Rn,dir
Вх	ANL	ACALL	CPL	CPL	CJNE	CJNE	CJNE	CJNE
	C,/bit	adr11	bit	C	A,#d8,rel	A,dir,rel	@Ri,#d8,rel	Rn,#d8,rel
Cx	PUSH	AJMP	CLR	CLR	SWAP	XCH	XCH	XCH
	dir	adr11	bit	C	A	A,dir	A,@Ri	A,Rn
Dx	POP	ACALL	SETB	SETB	DA	DJNZ	XCHD	DJNZ
	dir	adr11	bit	C	A	dir,rel	A,@Ri	Rn,rel
Ex	MOVX A,@DPTR	AJMP adr11		OVX @Ri	CLR A	MOV A,dir	MOV A,@Ri	MOV A,Rn
Fx	MOV @DPTR,A	ACALL adr11		OVX Ri,A	CPL A	MOV dir,A	MOV @Ri,A	MOV Rn,A

Additional 251 Instructions

Binary Mode	A5x8	A5x9	A5xA	A5xB	A5xC	A5xD	A5xE	A5xF
Source Mode	x8	x9	хА	хВ	хC	хD	хE	хF
0	JSLE rel	MOV Rm @WRj+dis	MOVZ WRj,Rm	INC Rm/WRj/ Drk,#short MOV reg,ind			SRA reg	
1	JSG rel	MOV@WRj +dis,Rm	MOVS WRj,Rm	DEC Rm/WRj/ Drk,#short MOV ind,reg			SRL reg	
2	JLE rel	MOV Rm, @DRk+dis			ADD Rm,Rm	ADD WRj,WRj	ADD reg,op2	ADD DRk,DRk
3	JG rel	MOV@DRk +dis,Rm					SLL reg	
4	JSL rel	MOV Wrj, @WRjj+dis			ORL Rm,Rm	ORL WRj,WRj	ORL reg,op2	
5	JSGE rel	MOV@WRj + dis,WRj			ANL Rm,Rm	ANL WRj,WRj	ANL reg,op2	
6	JE rel	MOV Wrj, @DRk+dis			XRL Rm,Rm	XRL WRj,WRj	XRL reg,op2	
7	JNE rel	MOV @Drk +dis,WRj	MOV op1,reg		MOV Rm,Rm	MOV WRj,WRj	MOV reg,op2	MOV DRk,DRk
8		LJMP@WRj EJMP@DRk	EJMP addr24		DIV Rm,Rm	DIV WRj,WRj		
9		LCALL@WR ECALL@DRk	ECALL addr24		SUB Rm,Rm	SUB WRj,WRj	SUB reg,op2	SUB DRk,DRk
A		BIT instructions	ERET		MUL Rm,Rm	MUL WRj,WRj		
В		TRAP			CMP Rm,Rm	CMP WRj,WRj	CMP reg,op2	CMP DRk,DRk
С			PUSH op1					
D			POP op1					
E								
F								

A,@EPTR

MOV

@EPTR,A

Fx

2

A5x4 A5x1 A5x8-A5xF A5x0 A5x2 A5x3 A5x5 A5x6-A5x7 A5x8-A5xF **EJMP** INC 0x adr23 esfr JBC **ECALL** DEC 1x esbit,rel adr23 esfr JB ADD 2x esbit,rel A,esfr JNB **ADDC** 3x esbit,rel A,esfr **EMOV** ORL ORL ORL **EMOV** A,@PR0+d2 A,@PR1+d2 esfr,A esfr,#data A,esfr ANL ANL ANL **EMOV EMOV** 5x esfr,A esfr,#data A,esfr @PR0+d2,A @PR1+d2,A XRL **XRL** XRL ADD ADD 6x esfr,A esfr,#data A,esfr PR0,#data2 PR1,#data2 **ORL EJMP** MOV 7x c,esbit @A+EPTR dir,#data MOV MOV MOV ANL 8x C.esbit esfr,esfr esfr,@Ri esfr.Rn MOV MOV MOVC **SUBB** 9x EPTR,#d23 esbit,c A,@A+EPTR A,esfr ORL MOV INC MOV MOV Ax C./esbit C.esbit **EPTR** @Ri,esfr Rn.esfr ANL CPL CJNE Вх C,/esbit esbit A,esfr,rel **PUSH** XCH CLR Сх esbit A,esfr esfr POP **SETB** DJNZ Dx esfr esbit esfr,rel MOVX MOV Ex

A,esfr

MOV

esfr,A

Additional 80C51MX Instructions via Prefix A5

Chapter 3. Writing Assembly Programs

The **Ax51** macro assembler is a multi pass assembler that translates **x51** assembly language programs into object files. These object files are then combined or linked using the **Lx51** Linker/Locator to form an executable, ready to run, absolute object module. As a subsequent step, absolute object modules can be converted to Intel HEX files suitable for loading onto to your target hardware, device programmer, or ICE (In-Circuit Emulator) unit.

The following sections describe the components of an assembly program, and some aspects of writing assembly programs. An assembly program consists of one or more statements. These statements contain directives, controls, and instructions.

Assembly Statements

Assembly program source files are made up of statements that may include assembler controls, assembler directives, or *x*51 assembly language instructions (mnemonics). For example:

```
$TITLE(Demo Program #1)

CSEG AT 0000h

JMP $

END
```

This example program consists of four statements. **\$TITLE** is an assembler control, **CSEG** and **END** are assembler directives, and **JMP** is an assembly language instruction.

Each line of an assembly program can contain only one control, directive, or instruction statement. Statements must be contained in exactly one line. Multiline statements are not allowed.

Statements in *x*51 assembly programs are not column sensitive. Controls, directives, and instructions may start in any column. Indentation used in the examples in this manual, is done for program clarity and is neither required nor expected by the assembler. The only exception is that arguments and instruction operands must be separated from controls, directives, and instructions by at least one space.

All x51 assembly programs must include the END directive. This directive signals to the assembler that this is the end of the assembly program. Any instructions, directives, or controls found after the END directive are ignored. The shortest valid assembly program contains only an END directive.

Directives

Assembler directives instruct the assembler how to process subsequent assembly language instructions. Directives also provide a way for you to define program constants and reserve space for variables.

"Chapter 4. Assembler Directives" on page 99 provides complete descriptions and examples of all of the assembler directives that you may include in your program. Refer to this chapter for more information about how to use directives.

Controls

Assembler controls direct the operation of the assembler when generating a listing file or object file. Typically, controls do not impact the code that is generated by the assembler. Controls can be specified on the command line or within an assembler source file.

The conditional assembly controls are the only assembler controls that will impact the code that is assembled by the **Ax51** assembler. The **IF**, **ELSE**, **ENDIF**, and **ELSEIF** controls provide a powerful set of conditional operators that can be used to include or exclude certain parts of your program from the assembly.

"Chapter 7. Invocation and Controls" on page 195 describes the available assembler controls in detail and provides an example of each. Refer to this chapter for more information about control statements.

Instructions

Assembly language instructions specify the program code that is to be assembled by the **Ax51** assembler. The **Ax51** assembler translates the assembly

instructions in your program into machine code and stores the resulting code in an object file.

Assembly instructions have the following general format:

[label:] mnemonic	operand [, operand] [, operand] [; comment]
where	
label	is a symbol name that is assigned the address at which the instruction is located.
mnemonic	is the ASCII text string that symbolically represents a machine language instruction.
operand	is an argument that is required by the specified mnemonic.
comment	is an optional description or explanation of the instruction. A comment may contain any text you wish. Comments are ignored by the assembler.

The "Instruction Sets" of the *x***51** microcontrollers are listed on page 40 by mnemonic and by machine language opcode. Refer to this section for more information about assembler instructions.

Comments

Comments are lines of text that you may include in your program to identify and explain the program. Comments are ignored by the **Ax51** assembler and are not required in order to generate working programs.

You can include comments anywhere in your assembler program. Comments must be preceded with a semicolon character (;). A comment can appear on a line by itself or can appear at the end of an instruction. For example:

```
;This is a comment NOP ;This is also a comment
```

When the assembler recognizes the semicolon character on a line, it ignores subsequent text on that line. Anything that appears on a line to the right of a semicolon will be ignored by the **Ax51** assembler. Comments have no impact on object file generation or the code contained therein.

Shaded directives and options are available only in AX51 and A251.

Symbols

A symbol is a name that you define to represent a value, text block, address, or register name. You can also use symbols to represent numeric constants and expressions.

Symbol Names

Symbols are composed of up to 31 characters from the following list:

A - Z, a - z, 0 - 9, _, and ?

A symbol name can start with any of these characters except the digits 0 - 9.

Symbols can be defined in a number of ways. You can define a symbol to represent an expression using the **EQU** or **SET** directives:

NUMBER_FIVE	VE EQU	5
TRUE FLAG	SET	1
FALSE FLAG	.G SET	0

You can define a symbol to be a label in your assembly program:

LABEL1: DJNZ RO, LABEL1

You can define a symbol to refer to a variable location:

SERIAL_BUFFER DATA 99h

Symbols are used throughout assembly programs. A symbolic name is much easier to understand and remember than an address or numeric constant. The following sections provide more information about how to use and define symbols.

Labels

A label defines a "place" (an address) in your program or data space. All rules that apply to symbol names also apply to labels. When defined, a label must be the first text field in a line. It may be preceded by tabs or spaces. A colon character (:) must immediately follow the symbol name to identify it as a label. Only one label may be defined on a line. For example:

LABEL1:	DS	2	
LABEL2:			;label by itself
NUMBER:	DB	27, 33, 'STRING', 0	; label at a message
COPY:	MOV	R6, #12H	; label in a program

In the above examples, LABEL1, LABEL2, NUMBER, and COPY are all labels.

When a label is defined, it receives the current value of the location counter of the currently selected segment. Refer to "Location Counter" on page 87 for more information about the location counter.

You can use a label just like you would use a program offset within an instruction. Labels can refer to program code, to variable space in internal or external data memory, or can refer to constant data stored in the program or code space.

You can use a label to transfer program execution to a different location. The instruction immediately following a label can be referenced by using the label. Your program can jump to or make a call to the label. The code immediately following the label will be executed.

You can also use labels to provide information to simulators and debuggers. A simulator or debugger can provide the label symbols while debugging. This can help to simplify the debugging process.

Labels may only be defined once. They may not be redefined.

Operands

Operands are arguments, or expressions, that are specified along with assembler directives or instructions. Assembler directives require operands that are constants or symbols. For example:

VVV	EQU	3		
	DS	10h		

Assembler instructions support a wider variety of operands than do directives. Some instructions require no operands and some may require up to 3 operands. Multiple operands are separated by commas. For example:

The number of operands that are required and their types depend on the instruction or directive that is specified. In the following table the first four operands can also be expressions. Instruction operands can be classified as one the following types:

Operand Type	Description
Immediate Data	Symbols or constants the are used as an numeric value.
Direct Bit Address	Symbols or constants that reference a bit address.
Program Addresses	Symbols or constants that reference a code address.
Direct Data Addresses	Symbols or constants that reference a data address.
Indirect Addresses	Indirect reference to a memory location, optionally with offset.
Special Assembler Symbol	Register names.

Special Assembler Symbols

The **Ax51** assembler defines and reserves names of the **x51** register set. These predefined names are used in **x51** programs to access the processor registers. Following, is a list of the each of the 8051, 80C51MX, and 251 registers along with a brief description:

	Register	Description			
	Α	Represents the 8051 Accumulator. It is used with many operations including multiplication and division, moving data to and from external memory, Boolean operations, etc.			
	DPTR	The DPTR register is a 16-bit data pointer used to address data in XDATA or CODE memory.			
	PC	The PC register is the 16-bit program counter. It contains the address of the next instruction to be executed.			
	С	The Carry flag; indicates the status of operations that generate a carry bit. It is also used by operations that require a borrow bit.			
	AB	The A and B register pair used in MUL and DIV instructions.			
	R0 – R7	The eight 8-bit general purpose 8051 registers in the currently active register bank. A Maximum of four register banks are available.			
	AR0 – AR7	Represent the absolute data addresses of R0 through R7 in the current register bank. The absolute address for these registers will change depending on the register bank that is currently selected. These symbols are only available when the USING directive is given. Refer to the USING directive for more information on selecting the register bank. These representations are suppressed by the NOAREGS directive. Refer to the NOAREGS directive for more information.			
51MX ONLY	PR0, PR1	Universal Pointer Registers of the 80C51MX architecture. Universal Pointer can access the complete 16MB address space of the 80C51MX. PR0 is composed of registers R1, R2, and R3. PR1 is composed of registers R5, R6, and R7.			
51N	EPTR	Additional extended data pointer register of the 80C51MX architecture. EPTR may be used to access the complete memory space.			
	R8 – R15	Additional eight 8–bit general purpose registers of the 251.			
251 ONLY	WR0 – WR30	Sixteen 16–bit general purpose registers of the 251. The registers WR0 - WR14 overlap the registers R0 - R15. Note that there is no WR1 available.			
25	DR0 - DR28 DR56, DR60	Ten 32-bit general purpose registers of 251. The registers DR0 - DR28 overlap the registers WR0 - WR30. Note that there is no DR1, DR2 and DR3 available.			

Immediate Data

An immediate data operand is a numeric expression that is encoded as a part of the machine language instruction. Immediate data values are used literally in an instruction to change the contents of a register or memory location. The pound (or number) sign (#) must precede any expression that is to be used as an immediate data operand. The following shows some examples of how the immediate data is typically used:

MY_VAL	EQU	50H	;	an equate symbol
	MOV	A,IO_PORT2	;	direct memory access to DATA
	MOV	A,#0E0h	;	load 0xE0 into the accumulator
	MOV	DPTR,#0x8000	;	load 0x8000 into the data pointer
	ANL	A,#128	;	AND the accumulator with 128
	XRL	A,#0FFh	;	XOR A with Offh
	MOV	R5,#MY VAL	;	load R5 with the value of MY VAL

Memory Access

A memory access reads or writes a value in to the various memory spaces of the x51 system.

Direct memory access encodes the memory address in the instruction that to reads or writes the memory. With direct memory accesses you can access variables in the memory class DATA and BIT. For the 251 also the EDATA memory class is addressable with direct memory accesses.

Indirect memory accesses uses the content of a register in the instruction that reads or writes into the memory. With indirect address operands it is possible to access all memory classes of the *x*51.

The following examples show how to access the different memory classes of an **x51** system.

DATA

Memory locations in the memory class DATA can be addressed with both: direct and indirect memory accesses. Special Function Registers (SFR) of the *x*51 have addresses above 0x80 in the DATA memory class. SFR locations can be addressed only with direct memory accesses. An indirect memory access to SFRs is not supported in the *x*51 microcontrollers.

Example for all 8051 variants

```
?DT?myvar SEGMENT DATA
                              ; define a SEGMENT of class DATA
          RSEG
                ?DT?myvar
VALUE:
                              ; reserve 1 BYTE in DATA space
                 1
IO PORT2
         DATA
                 OAOH
                              ; special function register
VALUE2
          DATA
                 20H
                              ; absolute memory location
?PR?myprog SEGMENT CODE
                              ; define a segment for program code
          RSEG ?PR?myprog
          MOV
                A,IO PORT2
                              ; direct memory access to DATA
          ADD
                A, VALUE
          MOV
                 VALUE2, A
          MOV
                R1,#VALUE
                              ; load address of VALUE to R1
          ADD
               A,@R1
                             ; indirect memory access to VALUE
```

BIT

Memory locations in the memory class BIT are addressed with the bit instructions of the 8051. Also the Special Function Registers (SFR) that are located bit-addressable memory locations can be addressed with bit instructions. Bit-addressable SFR locations are: 80H, 88H, 90H, 98H, 0A0H, 0A8H, 0B0H, 0B8H, 0C0H, 0C8H, 0D0H, 0D8H, 0E0H, 0E8H, 0F0H, and 0F8H.

Example for all 8051 variants

```
?BI?mybits SEGMENT BIT
                               ; define a SEGMENT of class BIT
          RSEG
                ?BI?mybits
FLAG:
          DBIT
                  1
                              ; reserve 1 Bit in BIT space
          DATA
                  90H
                              ; 8051 SFR PORT1
P1
GREEN LED BIT
                  P1.2
                              ; GREEN LED on I/O PORT P1.2
?PR?myprog SEGMENT CODE
                              ; define a segment for program code
          RSEG
                  ?PR?myprog
          SETB
                  GREEN LED
                              ; P1.2 = 1
                              ; direct memory access to DATA
          JB
                  FLAG, is on
          SETB
                  FLAG
                  ACC.5
          CLR
                               ; reset bit 5 in register A
           :
is on:
          CLR
                  FLAG
          CLR
                  GREEN LED
                                ; P1.2 = 0
```

Shaded directives and options are available only in AX51 and A251.

EBIT (only on Intel/Atmel WM 251)

The 251 provides with the EBIT memory class an expanded bit-addressable memory space that is addressed with extended bit instructions. Also all Special Function Registers (SFR) in the 251 can be addressed with extended bit instructions.

Example for Intel/Atmel WM 251

```
?EB?mybits SEGMENT EBIT
                                ; define a SEGMENT of class EBIT
           RSEG
                  ?EB?mybits
           DBIT
FLAG:
                  1
                                ; reserve 1 Bit in BIT space
CMOD
           DATA
                  0D9H
                                ; PCA Counter Modes
CPS0
           BIT
                   CMOD.1
                                ; CPS0 bit
?PR?myprog SEGMENT CODE
                                ; define a segment for program code
           RSEG
                   ?PR?myprog
           JB
                   FLAG, is on
                                ; direct memory access to DATA
           SETB
                  FLAG
           :
is on:
           CLR
                  FLAG
           CLR
                   CPS0
                                 ; CMOD.1 = 0
```

IDATA

Variables in this memory class are accessed via registers R0 or R1.

Example for all 8051 variants

```
?ID?myvars SEGMENT IDATA
                                ; define a SEGMENT of class IDATA
           RSEG
                   ?EB?mybits
BUFFER:
           DS
                   100
                                ; reserve 100 Bytes
?PR?myprog SEGMENT CODE
                                ; define a segment for program code
           RSEG
                   ?PR?myprog
           MOV
                   RO, #BUFFER ; load the address in RO
           MOV
                   A,@R0
                                ; read memory location buffer
           INC
                   R0
                                ; increment memory address in R0
           MOV
                   @RO,A
                                ; write memory location buffer+1
```

EDATA (Intel/Atmel WM 251, Philips 80C51MX only)

The EDATA memory is only available in the Philips 80C51MX and the Intel/Atmel WM 251 architecture.

In the Philips 80C51MX, the EDATA memory can be accessed via EPTR or the Universal Pointers PR0 and PR1. Universal Pointers can access any memory location in the 16MB address space.

Example for Philips 80C51MX

```
?ED?my seg SEGMENT EDATA
                                        ; define a SEGMENT of class EDATA
          RSEG
                  ?ED?my_seg
STRING:
          DS
                  100
                                        ; reserve 100 Bytes
?PR?myprog SEGMENT CODE
                                        ; define a segment for program code
          RSEG ?PR?myprog
          MOV
                  R1, #BYTE0 STRING
                                        ; load address of STRING in PRO
          MOV
                  R2, #BYTE1 STRING
          MOV
                  R3, #BYTE2 STRING
          VOM
                   A,@PRO
                                        ; load first byte of STRING in A
```

In the 251, EDATA memory can be accessed with direct memory addressing or indirect via the registers WR0 .. WR30. Also the memory class IDATA and DATA can be access with this addressing mode.

Example for Intel/Atmel WM 251

```
?ED?my seg SEGMENT EDATA
                                       ; define a SEGMENT of class EDATA
                  ?ED?my_seg
          RSEG
STRING:
          DS
                                       ; reserve 100 Bytes
                  100
?PR?myprog SEGMENT CODE
                                       ; define a segment for program code
          RSEG ?PR?myprog
          MOV
                  R11,STRING+2
                                       ; load character at STRING[2]
          MOV
                   WR4, #STRING
                                       ; load address of STRING
          MOV
                   R6,@WR4
                                       ; indirect access
          MOV
                   @WR4+2,R6
                                       ; access with constant offset
```

XDATA

The XDATA memory class can be accessed with the instruction MOVX via the register DPTR. A single page of the XDATA memory can be also accessed or via the registers R0, R1. At the C Compiler level this memory type is called **pdata** and the segment prefix ?PD? is used. The high address for this pdata page is typically set with the P2 register. But in new 8051 variants there are also dedicated special function registers that define the XDATA page address.

Example for all 8051 variants

```
?XD?my seg SEGMENT XDATA
                                      ; define a SEGMENT of class XDATA
          RSEG ?ED?my_seg
XBUFFER:
        DS
                 100
                                      ; reserve 100 Bytes
?PD?myvars SEGMENT XDATA INPAGE
                                      ; define a paged XDATA segment
         RSEG ?PD?myvars
VAR1:
          DS
                                      ; reserve 1 byte
?PR?myprog SEGMENT CODE
                                      ; define a segment for program code
          RSEG ?PR?myprog
          MOV
                P2, #HIGH ?PD?myvars ; load page address register
           :
                DPTR, #XBUFFER
                                     ; load address
          MOV
          MOVX
                 A,@DPTR
                                     ; access via DPTR
          MOV
                 R1,#VAR1
                                      ; load address
          MOVX
                 @R1,A
                                      ; access via R0 or R1
```

CODE and **CONST**

CODE or CONST memory can be accessed with the instruction MOVC via the DPTR register. The memory class CONST not possible with A51 and BL51.

Example for all 8051 variants

```
?CO?my seg SEGMENT CODE
                                      ; define a SEGMENT of class CODE
          RSEG ?CO?my seg
                                      ; a table with constant values
TABLE:
                  1,2,4,8,0x10
          DB
?PR?myprog SEGMENT CODE
                                      ; define a segment for program code
                  ?PR?myprog
          RSEG
          MOV
                  DPTR, #TABLE
                                      ; load address of table
          MOV
                 A,#3
                                      ; load offset into table
          MOVC A,@A+DPTR
                                      ; access via MOVC instruction
```

HDATA and **HCONST**

The HDATA and HCONST memory can be accessed with CPU instructions only in the Philips 80C51MX and the 251 architecture. HDATA and HCONST memory is simulated with memory banking on classic 8051 devices. The HDATA and HCONST memory class is not possible with A51 and BL51.

In the Philips 80C51MX, the HDATA and HCONST memory can be accessed via EPTR or the Universal Pointers PR0 and PR1. Universal Pointers can access any memory location in the 16MB address space.

Example for Philips 80C51MX

```
?HD?my seg SEGMENT HDATA
                                      ; define a SEGMENT of class HDATA
          RSEG
                  ?HD?my seg
ARRAY:
          DS
                  100
                                      ; reserve 100 Bytes
?PR?myprog SEGMENT CODE
                                      ; define a segment for program code
          RSEG ?PR?myprog
          MOV R1, #BYTEO ARRAY
                                      ; load address of ARRAY in PRO
          MOV R2, #BYTE1 ARRAY
          MOV R3, #BYTE2 ARRAY
          MOV A,@PRO
                                      ; load first byte of ARRAY in A
```

In the 251, HDATA and HCONST memory can be accessed via the registers DR0 .. DR28 and DR56. Any memory location can be accessed with this registers.

Example for Intel/Atmel WM 251

```
?HD?my seg SEGMENT HDATA
                                      ; define a SEGMENT of class HDATA
                 ?HD?my_seg
          RSEG
ARRAY:
          DS
                  100
                                      ; reserve 100 Bytes
?PR?myprog SEGMENT CODE
                                      ; define a segment for program code
          RSEG ?PR?myprog
                                      ; load address of ARRAY
          MOV
                  WR8, #WORD2 ARRAY
                  WR10, #WORD0 ARRAY
          MOV
                                      ; into DR8
          MOV R4,@DR8
                                      ; indirect access
          MOV @DR8+50H,R4
                                      ; access with constant offset
```

Program Addresses

Program addresses are absolute or relocatable expressions with the memory class CODE or ECODE. Typically program addresses are used in jump and call instructions. For indirect jumps or calls it is required to load a program address in a register or a jump table. The following jumps and calls are possible:

SJMP JZ JNZ Relative jumps include conditional jumps (CJNE, DJNZ, JB, JBC, JC, ...) and the unconditional SJMP instruction. The addressable offset is –128 to +127 bytes from the first byte of the instruction that follows the relative jump. When you use a relative jump in your code, you must use an expression that evaluates to the code address of the jump destination. The assembler does all the offset computations. If the address is out of range, the assembler will issue an error message.

ACALL AJMP

In-block jumps and calls permit access only within a 2KByte block of program space. The low order 11 bits of the program counter are replaced when the jump or call is executed. For Dallas 390 contiguous mode the block size is 512KB or 19 bits. If **ACALL** or **AJMP** is the last instruction in a block, the high order bits of the program counter change and the jump will be within the block following the **ACALL** or **AJMP**.

LCALL LJMP

Long jumps and calls allow access to any address within a 64KByte segment of program space. The low order 16 bits of the program counter are replaced when the jump or call is executed. For Dallas 390 contiguous mode: the block size is 16MB or 24 bits. One Philips 80C51MX and Intel/Atmel WM 251: if LCALL or LJMP is the last instruction in a 64KByte segment, the high order bits of the program counter change and the jump will into the segment following the LCALL or LJMP.

ECALL EJMP **Extended jumps and calls** allow access within the extended program space of the Intel/Atmel WM 251 or Philips 80C51MX.

CALL JMP Generic jumps and calls are two instruction mnemonics that do not represent a specific opcode. JMP may assemble to SJMP, AJMP, LJMP or EJMP. CALL may assemble to ACALL, LCALL or ECALL. These generic mnemonics always evaluate to an instruction, not necessarily the shortest, that will reach the specified program

address operand.

Example for all 8051 Variants

```
EXTRN CODE (my function)
          CSEG
                  AT
                        3
          JMP
                  ext int
                                       ; an interrupt vector
?PR?myintr SEGMENT CODE
                                       ; define a segment for program code
          RSEG ?PR?myintr
ext int:
          JB
                FLAG, flag OK
          INC my var
flag_OK:
          CPL
                  FLAG
          RETI
?PR?myprog SEGMENT CODE INBLOCK
                                       ; a segment within a 2K block
          RSEG ?PR?myprog
func1:
          CALL sub func
                                       ; will generate ACALL
                                       ; external function -> LCALL
loop:
          CALL
                 my function
          VOM
                  A,my var
          JNZ
                  loop
          RET
sub func:
          CLR
                  FLAG
          VOM
                  R0,#20
loop1:
          CALL
                  my function
                  R0,loop1
          DJNZ
          RET
```

Example with EJMP, ECALL for Philips 80C51MX and Intel/Atmel WM 251

```
EXTRN ECODE: FAR (my farfunc)
           EQU
                   ECODE OFF0000H
                                        ; Reset location on 251
Reset
?PR?my_seg SEGMENT ECODE
                                        ; define a SEGMENT of class EDATA
           RSEG
                   ?PR?my seg
func1
           PROC
                  FAR
                                        ; far function called with ECALL
           CALL
                   func2
                                        ; generates LCALL
           CALL
                 my farfunc
                                        ; generates ECALL
           JNB
                  Flag, mylab
           EJMP
                   Reset
mylab:
           ERET
           ENDP
func2
                   NEAR
           PROC
           CALL
                   my farfunc
                                        ; generates ECALL
           RET
           ENDP
```

Expressions and Operators

An operand may be a numeric constant, a symbolic name, a character string or an expression.

Operators are used to combine and compare operands within your assembly program. Operators are not assembly language instructions nor do they generate *x*51 assembly code. They represent operations that are evaluated at assembly-time. Therefore, operators can only handle calculations of values that are known when the program is assembled.

An expression is a combination of numbers, character string, symbols, and operators that evaluate to a single 32-bit binary number (for A51: 16-bit binary number). Expressions are evaluated at assembly time and can, therefore, be used to calculate values that would otherwise be difficult to determine beforehand.

The following sections describe operators and expressions and how they are used in x51 assembly programs.

Numbers

Numbers can be specified in hexadecimal (base 16), decimal (base 10), octal (base 8), and binary (base 2). The base of a number is specified by the last character in the number. A number that is specified without an explicit base is interpreted as decimal number.

The following table lists the base types, the base suffix character, and some examples:

Base	Suffix	Legal Characters	Examples
Hexadecimal	H, h	0 – 9, A – F, a – f	0x1234 0x99 1234H 0A0F0h 0FFh
Decimal	D, d	0 – 9	1234 65590d 20d 123
Octal	O, o, Q, q	0 – 7	177o 25q 123o 177777q
Binary	B, b	0 and 1	10011111b 101010101b

The first character of a number must be a digit between 0 and 9. Hexadecimal numbers which do not have a digit as the first character should be prefixed with a 0. The **Ax51** assembler supports also hex numbers written in C notation.

The dollar sign character (\$) can be used in a number to make it more readable, however, the dollar sign character cannot be the first or last character in the number. A dollar sign used within a number is ignored by the assembler and has no impact on the value of the number. For example:

1111\$0000\$1010\$0011b	is equivalent to	1111000010100011B
1\$2\$3\$4	is equivalent to	1234

Colon Notation for Numbers (A251 only)

The **A251** assembler supports the notation *page:number* for specifying absolute addresses. Numbers specified with this notation receive the memory type EDATA when page is 0 or ECODE for all other pages. In this way, you can use such numbers for referencing any memory location. For example:

```
ABSVAL1 EQU
              0:20H
                           ; symbol to address location 20H
ABSVAL2 EQU 0:80H ; symbol to address location 80H in EDATA space PORTO EQU S:80H ; symbol to SFR space 80H
         EQU 10:2000H ; entry point at location 102000H
ENTRY
         MOV
              WR0,ABSVAL1
         MOV
              R1,ABSVAL2
              PORTO,R1
         MOV
         EJMP ENTRY
         MOV WR0,0:20H ; access to ABSVAL1
         MOV R1,0:80H ; access to ABSVAL2
         MOV S:80H,R1
         EJMP 10:2000H
```

The colon notation is accepted in several A251 controls and is converted as described

Number in Colon Notation	Replaced with
VAL1 EQU 0:20H	VAL1 EQU EDATA 20H
VAL2 EQU 0FF:1000H	VAL2 EQU ECODE 0FF1000H
ORG 0FE:2000H	?modulename?number SEGMENT ECODE AT 0FE2000H RSEG ?modulename?number
ORG 0:400H	?modulename?number SEGMENT EDATA AT 400H RSEG ?modulename?number
CSEG AT 0FE:2000H	?modulename?number SEGMENT ECODE AT 0FE2000H RSEG ?modulename?number
BVAR1 BIT 0:20H.1	BVAR1 BIT 20H.1
BVAR1 BIT 0:30H.1	BVAR1 EQUEBIT 30H.1
PUSH.B #13	PUSH BYTE #13
PUSH.W #13	PUSH WORD #13

NOTE

The colon notation is provided for source compatibility with other 251 macro assemblers. If you do not need to port your code to other assemblers, it is recommended to use directly the replacement sequence in your assembler source file.

Characters

The **Ax51** assembler allows you to use ASCII characters in an expression to generate a numeric value. Up to two characters enclosed within single quotes (') may be included in an expression. More than two characters in single quotes in an expression will cause the **Ax51** assembler to generate an error. Following are examples of character expressions:

```
'A' evaluates to 0041h
'AB' evaluates to 4142h
'a' evaluates to 0061h
'ab' evaluates to 6162h
'' null string evaluates to 0000h
'abc' generates an ERROR
```

Characters may be used anywhere in your program as a immediate data operand. For example:

LETTER_A	EQU	'A'
TEST:	MOV SUBB	@RO, #'F' A, #'O'

Character Strings

Character strings can be used in combination with the **DB** directive to define messages that are used in your *x*51 assembly program. Character strings must be enclosed within single quotes ('). For example:

```
KEYMSG: DB 'Press any key to continue.'
```

generates the hexadecimal data (50h, 72h, 65h, 73h, 73h, 20h, ... 6Eh, 75h, 65h, 2Eh) starting at **KEYMSG**. You can mix string and numeric data on the same line. For example:

```
EOLMSG: DB 'End of line', 00h
```

appends the value 00h to the end of the string 'End of line'.

Two successive single quote characters can be used to insert a single quote into a string. For example:

```
MSGTXT: DB 'ISN''T A QUOTE REQUIRED HERE?'.
```

Location Counter

The **Ax51** assembler maintains a location counter for each segment. The location counter contains the offset of the instruction or data being assembled and is incremented after each line by the number of bytes of data or code in that line.

The location counter is initialized to 0 for each segment, but can be changed using the **ORG** directive.

The dollar sign character (\$) returns the current value of the location counter. This operator allows you to use the location counter in an expression. For example, the following code uses \$ to calculate the length of a message string.

```
MSG: DB 'This is a message', 0
MSGLEN EQU $ - MSG
```

You can also use in an instruction. For example, the following line of code will repeat forever.

JMP (\$) ; repeat forever

Operators

The **Ax51** assembler provides several classes of operators that allow you to compare and combine operands and expressions. These operators are described in the sections that follow.

Arithmetic Operators

Arithmetic operators perform arithmetic functions like addition, subtraction, multiplication, and division. These operators require one or two operands depending on the operation. The result is always a 16-bit value. Overflow and underflow conditions are not detected. Division by zero is detected and causes an assembler error.

Operator	Syntax	Description
+	+ expression	Unary plus sign
-	expression	Unary minus sign
+	expression + expression	Addition
-	expression – expression	Subtraction
*	expression * expression	Multiplication
1	expression / expression	Integer division
MOD	expression MOD expression	Remainder
(and)	(expression)	Specify order of execution

Binary Operators

Binary operators are used to complement, shift, and perform bit—wise operations on the binary value of their operands.

Operator	Syntax	Description
NOT	NOT expression	Bit-wise complement
SHR	expression SHR count	Shift right
SHL	expression SHL count	Shift left
AND	expression AND expression	Bit-wise AND
OR	expression OR expression	Bit-wise OR
XOR	expression XOR expression	Bit-wise exclusive OR

Relational Operators

The relational operators compare two operands. The result of the comparison is a TRUE or FALSE. A FALSE result has a value of 0000h. A TRUE result has a non-zero value.

The following table lists the relational operators and provides a brief description of each.

Operator	Syntax	Result
GTE	expression1 GTE expression2	True if expression1 is greater than or equal to expression2
LTE	expression1 LTE expression2	True if expression1 is less than or equal to expression2
NE	expression1 NE expression2	True if expression1 is not equal to expression2
EQ	expression1 EQ expression2	True if expression1 is equal to expression2
LT	expression1 LT expression2	True if expression1 is less than expression2
GT	expression1 GT expression2	True if expression1 is greater than expression2
>=	expression1 >= expression2	True if expression1 is greater than or equal to expression2
<=	expression1 <= expression2	True if expression1 is less than or equal to expression2
<>	expression1 <> expression2	True if expression1 is not equal to expression2
=	expression1 = expression2	True if expression1 is equal to expression2
<	expression1 < expression2	True if expression1 is less than expression2
>	expression1 > expression2	True if expression1 is greater than expression2

Class Operators

The class operator assigns a memory class to an expression. This is how you associate an expression with a class. The **Ax51** assembler generates an error message if you use an expression with a class on an instruction which does not support this class, for example, when you use an **XDATA** expression as a direct address.

The following table lists the class operators and provides a brief description of each.

Operator	Syntax	Description
BIT	BIT expression	Assigns the class BIT to the expression.
CODE	CODE expression	Assigns the class CODE to the expression.
CONST	CONST expression	Assigns the class CONST to the expression.
DATA	DATA expression	Assigns the class DATA to the expression.
EBIT	EBIT expression	Assigns the class EBIT to the expression.
ECODE	ECODE expression	Assigns the class ECODE to the expression.
ECONST	ECONST expression	Assigns the class ECONST to the expression.
EDATA	EDATA expression	Assigns the class EDATA to the expression.
IDATA	IDATA expression	Assigns the class IDATA to the expression.
HCONST	HCONST expression	Assigns the class HCONST to the expression.
HDATA	HDATA expression	Assigns the class HDATA to the expression.
XDATA	XDATA expression	Assigns the class XDATA to the expression.

Type Operators

The type operator assigns a data type to an expression. The A251 assembler generates an error if you attempt to use an instruction with the incorrect data type. For example, this happens when you use a WORD expression as an argument in a byte-wide instruction of the 251.

Operator	Syntax	Description
BYTE	BYTE expression	Assigns the type BYTE to the expression.
WORD	WORD expression	Assigns the class WORD to the expression.
DWORD	DWORD expression	Assigns the class DWORD to the expression.
NEAR	NEAR expression	Assigns the class NEAR to the expression.
FAR	FAR expression	Assigns the class FAR to the expression.

Miscellaneous Operators

Ax51 provides operators that do not fall into the previously listed categories. These operators are listed and described in the following table.

Operator	Syntax	Description
LOW	LOW expression	Low-order byte of expression
HIGH	HIGH expression	High-order byte of expression
BYTE0	BYTE0 expression	Byte 0 of expression. See table below. (identical with LOW).
BYTE1	BYTE1 expression	Byte 1 of expression. See table below. (identical with HIGH).
BYTE2	BYTE2 expression	Byte 2 of expression. See table below.
BYTE3	BYTE3 expression	Byte 3 of expression. See table below.
WORD0	WORD0 expression	Word 0 of expression. See table below.
WORD2	WORD2 expression	Word2 of expression. See table below.
MBYTE	MBYTE expression	AX51 only: memory type information for C51 run-time libraries. Returns the memory type that is used in the C51 run-time library to access variables defined with the far memory type.

The following table shows how the byte and word operators impact a 32-bit value.

MSB			LSB
BYTE3	BYTE2	BYTE1	BYTE0
WORD2		WO	RD0
		HIGH	LOW

The following table shows how the byte and word operators impact a 32-bit value.

Operator Precedence

All operators are evaluated in a certain, well—defined order. This order of evaluation is referred to as operator precedence. Operator precedence is required in order to determine which operators are evaluated first in an expression. The following table lists the operators in the order of evaluation. Operators at level 1 are evaluated first. If there is more than one operator on a given level, the leftmost operator is evaluated first followed by each subsequent operator on that level.

Level	Operators
1	()
2	NOT, HIGH, LOW, BYTE0, BYTE1, BYTE2, BYTE3, WORD0, WORD2
3	BIT, CODE, CONST, DATA, EBIT, EDATA, ECONST, ECODE, HCONST, HDATA, IDATA, XDATA
4	BYTE, WORD, DWORD, NEAR, FAR
5	+ (unary), – (unary)
6	*, /, MOD
7	+, -
8	SHR, SHL
9	AND, OR, XOR
10	>=, <=, =, <>, <, >, GTE, LTE, EQ, NE, LT, GT

Expressions

An expression is a combination of operands and operators that must be calculated by the assembler. An operand with no operators is the simplest form of an expression. An expression can be used in most places where an operand is required.

Expressions have a number of attributes that are described in the following sections.

Expression Classes

Expressions are assigned classes based on the operands that are used. The following classes apply to expressions:

Expression Class	Description	
N NUMB	A classless number.	
C ADDR	A CODE address symbol.	
D ADDR	A DATA address symbol.	
I ADDR	An IDATA address symbol.	
X ADDR	An XDATA address symbol.	
B ADDR	A BIT address symbol.	
CO ADDR	A CONST address symbol.	
EC ADDR	An ECONST address symbol.	
CE ADDR	An ECODE address symbol.	
ED ADDR	An EDATA address symbol.	
EB ADDR	An EBIT address symbol.	
HD ADDR	An HDATA address symbol.	
HC ADDR	An HCONST address symbol.	

Typically, expressions are assigned the class **NUMBER** because they are composed only of numeric operands. You may assign a class to an expression using a class operand. An address symbol value is automatically assigned the class of the segment where it is defined. When a value has a class, a few rules apply to how expressions are formed:

- 1. The result of a unary operation has the same class as its operand.
- 2. The result of all binary operations except + and will be a **NUMBER** type.
- 3. If only one of the operands of an addition or subtraction operation has a class, the result will have that class. If both operands have a class, the result will be a **NUMBER**.

This means that a class value (i.e. an addresses symbol) plus or minus a number (or a number plus a class value) give a value with class.

Examples

```
data_address - 10 gives a data_address value
10 + edata_address gives an edata_address value
(data_address - data_address) gives a classless number
code address + (data address - data address) gives a code address value
```

Expressions that have a type of **NUMBER** can be used virtually anywhere. Expressions that have a class can only be used where a class of that type is valid.

Relocatable Expressions

Relocatable expressions are so named because they contain a reference to a relocatable or external symbol. These types of expressions can only be partially calculated by the assembler since the assembler does not know the final location of relocatable segments. The final calculations are performed by the linker.

A relocatable expression normally contains only a relocatable symbol, however, it may contain other operands and operators as well. A relocatable symbol can be modified by adding or subtracting a constant value.

Examples for valid relocatable expression

- relocatable symbol + absolute expression
- relocatable symbol absolute expression
- absolute_expression + relocatable_symbol

There are two basic types of relocatable expressions: simple relocatable expressions and extended relocatable expressions.

Simple Relocatable Expressions

Simple relocatable expressions contain symbols that are defined in a relocatable segment. Segment and external symbols are not allowed in simple relocatable expressions.

Simple relocatable expression can be used in four contexts:

- 1. As an operand to the ORG directive.
- 2. As an operand to a symbol definition directive (i.e. EQU, SET)
- 3. As an operand to a data initialization directive (DB, DW or DD)
- 4. As an operand to a machine instruction

Examples for simple relocatable expressions

```
REL1 + ABS1 * 10
REL2 - ABS1
REL1 + (REL2 - REL3) assuming REL2 and REL3 refer to the same segment.
```

Invalid form of simple relocatable expressions

```
(REL1 + ABS1) * 10relocatable value may not be multiplied.(EXT1 - ABS1)this is a general relocatable expressionREL1 + REL2you cannot add relocatable symbols.
```

Extended Relocatable Expressions

The extended relocatable expressions have generally the same rules that apply to simple relocatable expressions. Segment and external symbols are allowed in extended relocatable expressions. Extended relocatable expressions may be used only in statements that generate code as operands; these are:

- As an operand to a data initialization directive (DB, DW or DD)
- As an operand to a machine instruction

Examples for extended relocatable expressions

```
REL1 + ABS1 * 10
EXT1 - ABS1
LOW (REL1 + ABS1)
WORD2 (SEG1)
```

Invalid form of simple relocatable expressions

```
(SEG1 + ABS1) * 10 relocatable value may not be multiplied.

(EXT1 - REL1) you can add/subtract only absolute quantities

LOW (REL1) + ABS1 LOW may be applied only to the final relocatable expression
```

Examples with Expressions

```
EXTRN CODE (CLAB)
                                 ; entry in CODE space
EXTRN DATA (DVAR)
                                 ; variable in DATA space
MSK
       EQU
             0F0H
                                 ; define a symbol to replace 0xF0
VALUE
       EQU MSK - 1
                                 ; another constant symbolic value
LVAL
       EQU 12345678H
                                ; LVAL get the value 12345678H
?PR?FOO SEGMENT CODE
       RSEG ?PR?FOO
ENTRY: MOV A, #40H
                                 ; load register with constant
       MOV R5, #VALUE
                               ; load constant symbolic value
       MOV R3, #(0x20 AND MASK) ; examples for calculations
       MOV R7, #LOW (VALUE + 20H)
       MOV R6,#1 OR (MSK SHL 4)
       MOV R0, DVAR+20
                                ; load content from address DVAR+20
       MOV R1, #LOW (CLAB+10) ; load low byte of address CLAB+10
       MOV WR4, #WORD2 (LVAL) ; load high word of LVAL
       MOV DR0, #ENTRY
                                ; load low word of addr. ENTRY to DRO
       MOVH DRO, #WORD2 (ENTRY) ; load high word of addr. ENTRY to DRO
       MOV R4,@DR0
                                ; load content of ENTRY to R4
       MOV R5,80H
                                ; load DATA addr. 80H (= SFR P0) to R5
       MOV R5, EDATA 80H
                                ; load EDATA address 80H to R5
       SETB 30H.2
                                ; set bit at 30H.2 (long address)
       SETB 20H.2
                                 ; set bit at 20H.2 (short address)
       END
```

Chapter 4. Assembler Directives

This chapter describes the assembler directives. It shows how to define symbols and how to control the placement of code and data in program memory.

Introduction

The **Ax51** assembler has several directives that permit you to define symbol values, reserve and initialize storage, and control the placement of your code.

The directives should not be confused with instructions. They do not produce executable code, and with the exception of the DB, DW and DD directives, they have no direct effect on the contents of code memory. These directives change the state of the assembler, define user symbols, and add information to the object file

The following table provides an overview of the assembler directives. Page refers to the page number in this user's guide where you can find detailed information about the directive.

Directive	/ Page	Format	Description
BIT	114	symbol BIT bit_address	Define a bit address in bit data space.
BSEG	111	BSEG [AT absolute address]	Define an absolute segment within the bit address space.
CODE	114	symbol CODE code_address	Assign a symbol name to a specific address in the code space.
CSEG	111	CSEG [AT absolute address]	Define an absolute segment within the code address space.
DATA	114	symbol DATA data_address	Assign a symbol name to a specific on-chip data address.
DB	119	[label:] DB expression [, expr]	Generate a list of byte values.
DBIT	122	[label:] DBIT expression	Reserve a space in bit units.
DD	121	[label:] DD expression [, expr]	Generate a list of double word values.
DS	123	[label:] DS expression	Reserve space in byte units.
DSB	124	[label:] DSB expression	Reserve space in byte units.
DSD	126	[label:] DSD expression	Reserve space in double word units.
DSEG	111	DSEG [AT absolute address]	Define an absolute segment within the indirect internal data space.

Directive /	Page	Format	Description
DSW	125	[label:] DSW expression	Reserve space in word units; advances the location counter of the current segment.
DW	120	[label:] DW expression [, expr]	Generate a list of word values.
END	136	END	Indicate end of program.
EQU	113	EQU expression	Set symbol value permanently.
ERROR_	_136	ERROR text	Generate a standard error message.
EVEN	134	EVEN	Ensure word alignment for variables.
EXTRN EXTERN	131	EXTRN class [:type] (symbol [,]) EXTERN class [:type] (symbol [,])	Defines symbols referenced in the current module that are defined in other modules.
IDATA	114	symbol IDATA idata_address	Assign a symbol name to a specific indirect internal address.
ISEG	111	ISEG [AT absolute address]	Define an absolute segment within the internal data space.
LABEL	129	name[:] LABEL [type]	Assign a symbol name to a address location within a segment.
LIT	116	symbol LIT 'literal string'	Assign a symbol name to a string.
NAME	132	NAME modulname	Specify the name of the current module.
ORG	133	ORG expression	Set the location counter of the current segment.
PROC ENDP	127	name PROC [type] name ENDP	Define a function start and end.
PUBLIC	130	PUBLIC symbol [, symbol]	Identify symbols which can be used outside the current module.
RSEG	110	RSEG seg	Select a relocatable segment.
SEGMENT	106	seg SEGMENT class [reloctype] [alloctype]	Define a relocatable segment.
SET	113	SET expression	Set symbol value temporarily.
sfr, sfr16 sbit	116	sfr symbol = address; sfr16 symbol = address; sbit symbol = address;	Define a special function register (SFR) symbol or a SFR bit symbol.
USING	134	USING expression	Set the predefined symbolic register address and reserve space for the specified register bank.
XDATA	114	symbol XDATA xdata_address	Assign a symbol name to a specific off-chip data address.
XSEG	111	XSEG [AT absolute address]	Define an absolute segment within the external data address space.

The directives are divided into the following categories:

Segment Control

Generic Segments: SEGMENT, RSEG

Absolute Segments: CSEG, DSEG, BSEG, ISEG, XSEG

Symbol Definition

Generic Symbols: EQU, SET

Address Symbols: BIT, CODE, DATA, IDATA, XDATA

SFR Symbols: sfr, sfr16, sbit

Text Replacement: LIT

- Memory Initialization DB, DW, DD
- Memory Reservation DBIT, DS, DSB, DSW, DSD
- Procedure Declaration PROC / ENDP, LABEL
- Program Linkage PUBLIC, EXTRN / EXTERN, NAME
- Address Control ORG, EVEN, USING
- OthersEND, __ERROR__

The **Ax51** assembler is a multi-pass assembler. In the first pass, symbol values are determined. In the subsequent passes, forward references are resolved and object code is produced. This structure imposes a restriction on the source program: expressions which define symbol values (refer to "Symbol Definition" on page 113) and expressions which control the location counter (refer to "ORG" on page 133, "DS" on page 123, and "DBIT" on page 122) may not have forward references.

Segment Directives

A segment is a block of code or data memory the assembler creates from code or data in an *x*51 assembly source file. How you use segments in your source modules depends on the complexity of your application. Smaller applications need less memory and are typically less complex than large multi-module applications.

The *x***51** CPU has several specific memory areas. You use segments to locate program code, constant data, and variables in these areas.

Location Counter

Ax51 maintains a location counter for each segment. The location counter is a pointer to the address space of the active segment. It represents an offset for generic segments or the actual address for absolute segments. When a segment is first activated, the location counter is set to 0. The location counter is changed after each instruction by the length of the instruction. The memory initialization and reservation directives (i.e. DS, DB or DBIT) change the value of the location counter as memory is allocated by these directives. The ORG directive sets a new value for the location counter. If you change the active segment and later return to that segment, the location counter is restored to its previous value. Whenever the assembler encounters a label, it assigns the current value of the location counter and the type of the current segment to that label.

The dollar sign (\$) indicates the value of the location counter in the active segment. When you use the \$ symbol, keep in mind that its value changes with each instruction, but only after that instruction has been completely evaluated. If you use \$ in an operand to an instruction or directive, it represents the address of the first byte of that instruction.

The following sections describe the different types of segments.

Generic Segments

Generic segments have a name and a class as well as other attributes. Generic segments with the same name but from different object modules are considered to be parts of the same segment and are called partial segments. These segments are combined at link time by the linker/locator.

Generic segments are created using the **SEGMENT** directive. You must specify the name of the segment, the segment class, and an optional relocation type and alignment type when you create a relocatable segment.

Example

MYPROG SEGMENT CODE

defines a segment named MYPROG with a memory class of CODE. This means that data in the MYPROG segment will be located in the code or program area of the x51. Refer to "SEGMENT" on page 106 for more information on how to declare generic segments.

Once you have defined a relocatable segment name, you must select that segment using the **RSEG** directive. When **RSEG** is used to select a segment, that segment becomes the active segment that **Ax51** uses for subsequent code and data until the segment is changed with **RSEG** or with an absolute segment directive

Example

RSEG MYPROG

will select the MYPROG segment that is defined above.

Typically, assembly routines are placed in generic segments. If you interface your assembly routines to C, all of your assembly routines must reside in separate generic segments and the segment names must follow the standards used by Cx51. Refer to the Compiler User's Guide for more information on interfacing assembler programs to C.

Stack Segment

The x51 architecture uses a hardware stack to store return addresses for CALL instructions and also for temporary storage using the PUSH and POP instructions. An 8051 application that uses these instructions must setup the stack pointer to an area of memory that will not be used by other variables.

For the classic **8051** a stack segment must be defined and space must be reserved as follows. This definition also works for the extended 8051 and the 251, however these controllers typically support stack also in other areas.

STACK	SEGMENT	IDATA		
	RSEG	STACK	;	select the stack segment
	DS	10h	;	reserve 16 bytes of space

Then, you must initialize the stack pointer early in your program.

For the Philips 80C51MX or the Intel/Atmel WM 251 a stack segment may be defined and space must be reserved as follows.

STACK	SEGMENT	EDATA	
	RSEG	STACK	; select the stack segment
	DS	100h	; reserve 256 bytes of space

Then, you must initialize the stack pointer early in your program.

```
CSEG
                  ΑТ
                                             ; RESET Vector
                  TMP
                            STARTUP
                                             ; Jump to startup code
STARTUP:
                                             ; code executed at RESET
; Stack setup for Philips 80C51MX
                           MXCON, #0x02
                                             ; enable extended stack
                  ORL
                  MOV
                            SPE, #HIGH (STACK - 1) ; load Stack high
                            SP, #LOW (STACK - 1) ; load Stack low
                  VOM
; for Intel/Atmel WM 251
                            DR60, #STACK - 1 ; load Stack Pointer
```

If you are interfacing assembly routines to C, you probably do not need to setup the stack. This is already done for you in the C startup code.

Absolute Segments

Absolute segments reside in a fixed memory location. Absolute segments are created using the CSEG, DSEG, XSEG, ISEG, and BSEG directives. These directives allow you to locate code and data or reserve memory space in a fixed location. You use absolute segments when you need to access a fixed memory location or when you want to place program code or constant data at a fixed memory address. Refer to the CSEG, DSEG, ISEG, XSEG, ISEG directives for more information on how to declare absolute segments.

After reset, the 8051 variants begin program executing at CODE address 0. The Intel/Atmel WM 251 starts execution at address FF0000. Some type of program code must reside at this address. You can use an absolute segment to force program code into this address. The following example is used in the **Cx51** startup routines to branch from the reset address to the beginning of the initialization code.

```
.
.
CSEG AT 0
RESET_VEC: LJMP STARTUP
.
.
```

The program code that we place at address 0000h (for 251 at address FF0000h) with the CSEG AT 0 directive performs a jump to the STARTUP label.

AX51 and **A251** supports absolute segment controls for compatibility to **A51**. **AX51** and **A251** translates the CSEG, DSEG, XSEG, ISEG and BSEG directives to a generic segment directive.

Default Segment

By default, **Ax51** assumes that the CODE segment is selected and initializes the location counter to 0000h (FF0000h) when it begins processing an assembly source module. This allows you to create programs without specifying any relocatable or absolute segment directives.

SEGMENT

The **SEGMENT** directive is used to declare a generic segment. A relocation type and an allocation type may be specified in the segment declaration. The **SEGMENT** directive is specified using the following format:

segment	SEGMENT	class	reloctype	alloctype
where				
segment		name is segmen represen	referred by the t symbol name of	assign to the segment. This symbol following RSEG directive. The can be used also in expressions to art address of the combined segment art hker/Locator.
class		is the memory class to use for the specified segment. The class specifies the memory space for the segment. See the table below for more information.		
relocty	pe	what re	location options Locator. Refer	or the segment. This determines is may be performed by the to the table below for more
allocty	pe	relocati	on options may	or the segment. This determines what be performed by the Linker/Locator. w for more information.

Class

The name of each segment within a module must be unique. However, the linker will combine segments having the same segment type. This applies to segments declared in other source modules as well.

The class specifies the memory class space for the segment. The A251 differentiates between basic classes and user-defined classes. The class is used by the linker/locator to access all the segments which belong to that class.

The basic classes are listed below:

Basic Class	Description
BIT	BIT space (address 20H 2FH).
CODE	CODE space
CONST	CONST space; same as CODE but for constant only; access via MOVC.
DATA	DATA space (address 0 to 7FH & SFR registers).
EBIT	Extended 251 bit space (address 20H 7FH)
EDATA	EDATA space
ECONST	ECONST space; same as EDATA but for constants
IDATA	IDATA space (address 0 to 0FFH).
ECODE	Entire Intel/Atmel WM 251 and Philips 80C51MX address space for program code.
HCONST	Entire Intel/Atmel WM 251 and Philips 80C51MX address space for constants.
HDATA	Entire Intel/Atmel WM 251 and Philips 80C51MX address space for data.
XDATA	XDATA space; access via MOVX.

User-defined Class Names (AX51 & A251 only)

User-defined class names are composed of a basic class name and an extension and are enclosed in single quotes ('). They let you access the same address space as basic class names. The advantage is that you may declare several segments with a user-defined class and later use the linker to locate that class (and its segments) at a specific physical address. Refer to the "CLASSES" on page 336 for information on how to locate user defined classes.

Examples

seg1	SEGMENT	'NDATA_FLASH'
seg2	SEGMENT	'HCONST_BITIMAGE'
seg3	SEGMENT	'DATA1'

Relocation Type

The optional relocation type defines the relocation operation that may be performed by the Linker/Locator. The following table lists the valid relocation types:

Relocation Type	Description
AT address	Specifies an absolute segment. The segment will be placed at the specified address.
BITADDRESSABLE	Specifies a segment which will be located within the bit addressable memory area (20H to 2FH in DATA space). BITADDRESSABLE is only allowed for segments with the class DATA that do not exceed 16 bytes in length.
INBLOCK	Specifies a segment which must be contained in a 2048Byte block. This relocation type is only valid for segments with the class CODE .
INPAGE	Specifies a segment which must be contained in a 256Byte page.
OFFS offset	Specifies an absolute segment. The segment is placed at the starting address of the memory class plus the specified <i>offset</i> . The advantage compared to the AT relocation type is that the start address can be modified with the Lx51 linker/locater control CLASSES. Refer to the "CLASSES" on page 336 for more information.
OVERLAYABLE	Specifies that the segment can share memory with other segments. Segments declared with this relocation type can be overlaid with other segments which are also declared with the OVERLAYABLE relocation type. When using this relocation type, the segment name must be declared according to the C251, CX51, C51 or PL/M-51 segment naming rules. Refer to the <i>C Compiler User's Guide</i> for more information.
INSEG	Specifies a segment which must be contained in a 64KByte segment.

Allocation Type

The optional allocation type defines the allocation operation that may be performed by the Linker/Locator. The following table lists the valid allocation types:

Allocation Type	Description
BIT	Specify bit alignment for the segment. Default for all segments with the class BIT.
ВҮТЕ	Specify byte alignment for the segment. Default for all segments except of BIT.
WORD	Specify word alignment for the segment.
DWORD	Specify dword alignment for the segment.
PAGE	Specify a segment whose starting address must be on a 256Byte page boundary.
BLOCK	Specify a segment whose starting address must be on a 2048Byte block boundary.
SEG	Specify a segment whose starting address must be on a 64KByte segment boundary.

Examples for Segment Declarations

Defines a segment with the name IDS and the memory class IDATA.

MYSEG	SEGMENT	CODE	AT OFF20

Defines a segment with the name MYSEG and the memory class CODE to be located at address 0FF2000H.

DSEG SEGMENT HDATA INSEG DWORL

Defines a segment with the name HDSEG and the memory class HDATA. The segment is located within one 64KByte segment and is DWORD aligned.

Defines a segment with the name XDSEG and the memory class XDATA. The segment is PAGE aligned, this means it starts on a 256Byte page.

HCSEG SEGMENT HCONST SEG

Defines a segment with the name HCSEG with the memory class HCONST. The segment is SEGMENT aligned, this means it starts on a 64KByte segment.

RSEG

The **RSEG** directive selects a generic segment that was previously declared using the **SEGMENT** directive. The **RSEG** directive uses the following format:

```
where

segment is the name of a segment that was previously defined using the SEGMENT directive. Once selected, the specified segment remains active until a new segment is specified.
```

```
.
.
.
MYPROG SEGMENT CODE ; declare a segment

RSEG MYPROG ; select the segment

MOV A, #0

MOV PO, A

.
```

BSEG, CSEG, DSEG, ISEG, XSEG

The **BSEG**, **CSEG**, **DSEG**, **ISEG**, **XSEG** directives select an absolute segment. This directives use the following formats:

BSEG	AT	address	defines an absolute BIT segment.	
CSEG	AT	address	defines an absolute CODE segment.	
DSEG	AT	address	defines an absolute DATA segment.	
ISEG	AT	address	defines an absolute IDATA segment.	
XSEG	AT	address	defines an absolute XDATA segment.	

where

address

is an optional absolute base address at which the segment begins. The address may not contain any forward references and must be an expression that can be evaluated to a valid address.

CSEG, DSEG, ISEG, BSEG and XSEG select an absolute segment within the code, internal data, indirect internal data, bit, or external data address spaces. If you choose to specify an absolute address (by including AT address), the assembler terminates the last absolute segment, if any, of the specified segment type, and creates a new absolute segment starting at that address. If you do not specify an address, the last absolute segment of the specified type is continued. If no absolute segment of this type was selected and the absolute address is omitted, a new segment is created starting at location 0. You cannot use any forward references and the start address must be an absolute expression.

The AX51 and A251 Macro Assembler supports the BSEG, CSEG, DSEG, ISEG, and XSEG directives for A51 compatibility.

These directives are converted to standard segments as follows:

A51 Directive	Converted to AX51/A251 Segment Declaration
BSEG AT 20H.1	?BI?modulename?n SEGMENT OFFS 20H.1
CSEG AT 1234H	?CO?modulename?n SEGMENT OFFS 1234H
DSEG AT 40H	?DT?modulename?n SEGMENT OFFS 40H
ISEG AT 80H	?ID?modulename?n SEGMENT OFFS 80H
XSEG AT 5100H	?XD?modulename?n SEGMENT OFFS 5100H

where

modulname is the name of the current assembler module

n is a sequential number incremented for every absolute

segment.

	BSEG AT 30h	; absolute bit segment @ 30h
DEC_FLAG:	DBIT 1	; absolute bit
INC_FLAG:	DBIT 1	
	CSEG AT 100h	; absolute code segment @ 100h
PARITY TAB:	DB 00h	; parity for 00h
	DB 01h	; 01h
	DB 01h	; 02h
	DB 00h	; 03h
•		
	DB 01h	; FEh
	DB 00h	; FFh
	DSEG AT 40h	; absolute data segment @ 40h
TMP A:	DS 2	; absolute data word
TMP B:	DS 4	
_	ISEG AT 40h	; abs indirect data seg @ 40h
TMP IA:	DS 2	
TMP IB:	DS 4	
-		
	XSEG AT 1000h	; abs external data seg @ 1000h
OEMNAME:	DS 25	; abs external data
PRDNAME:	DS 25	
VERSION:	DS 25	

Symbol Definition

The symbol definition directives allow you to create symbols that can be used to represent registers, numbers, and addresses.

Symbols defined by these directives may not have been previously defined and may not be redefined by any means. The **SET** directive is the only exception to this.

EQU, SET

The **EQU** and **SET** directive assigns a numeric value or register symbol to the specified symbol name. Symbols defined with **EQU** may not have been previously defined and may not be redefined by any means. The **SET** directive allows later redefinition of symbols. Statements involving the **EQU** or **SET** directive are formatted as follows:

symbol	EQU	expression
symbol	EQU	register
symbol	SET	expression
symbol	SET	register

where

symbol is the name of the symbol to define. The expression or

register specified in the **EQU** or **SET** directive will be substituted for each occurrence of **symbol** that is used in

your assembly program.

expression is a numeric expression which contains no forward

references, or a simple relocatable expression.

register is one of the following register names: A, R0, R1, R2, R3,

R4, R5, R6, or R7.

Symbols defined with the **EQU** or **SET** directive may be used anywhere in operands, expressions, or addresses. Symbols that are defined as a register name can be used anywhere a register is allowed. A251 replaces each occurrence of the defined symbol in your assembly program with the specified numeric value or register symbol.

Symbols defined with the **EQU** directive may not be changed or redefined. You cannot use the **SET** directive if a symbol was previously defined with **EQU** and you cannot use the **EQU** directive if a symbol which was defined with SET.

Examples

LIMIT	EQU	1200
VALUE	EQU	LIMIT - 200 + 'A'
SERIAL	EQU	SBUF
ACCU	EQU	A
COUNT	EQU	R5
VALUE	SET	100
VALUE	SET	VALUE / 2
COUNTER	SET	R1
TEMP	SET	COUNTER
TEMP	SET	VALUE * VALUE

CODE, DATA, IDATA, XDATA

The **BIT**, **CODE**, **DATA**, **IDATA**, and **XDATA** directives assigns an address value to the specified symbol. Symbols defined with the **BIT**, **CODE**, **DATA**, **IDATA**, and **XDATA** directives may not be changed or redefined. The format of theses directives is:

symbol	BIT	bit_address	defines a BIT symbol
symbol	CODE	code_address	defines a CODE symbol
symbol	DATA	$\mathtt{data}_\mathtt{address}$	defines a DATA symbol
symbol	IDATA	idata_address	defines an IDATA symbol
symbol	XDATA	xdata_address	defines a XDATA symbol

where

symbol	is the name of the symbol to define. The symbol name can be used anywhere an address of this memory class is valid.
bit_address	is the address of a bit in internal data memory in the area 20H 2FH or a bit address of an 8051 bit-addressable SFR.
code_address	is a code address in the range 0000H 0FFFFH.
data_address	is a data memory address in the range 0 to 127 or a special function register (SFR) address in the range 128 255.
idata_address	is an idata memory address in the range 0 to 255.
ydata address	is an xdata memory address in the range 0 to 65535

DATA SEG	SEGMENT I	BITADDRESSABLE		
RSEG	DATA_SEG		;	a bit-addressable rel_seg
CTRL:	DS	1		a 1-byte variable (CTRL)
ALARM	BIT	CTRL.0	;	bit in a relocatable byte
SHUT	BIT	ALARM+1	;	the next bit
ENABLE_FLAG	BIT	60H	;	an absolute bit
DONE_FLAG	BIT	24H.2	;	an absolute bit
P1_BIT2	EQU	90H.2	;	a SFR bit
RESTART	CODE	00H		
INTVEC_0	CODE	RESTART + 3		
INTVEC_1	CODE	RESTART + OBH		
INTVEC_2	CODE	RESTART + 1BH		
SERBUF	DATA	SBUF	;	redfinition of SBUF
RESULT	DATA	40H		
RESULT2	DATA	RESULT + 2		
PORT1	DATA	90H	;	a SFR symbol
BUFFER	IDATA	60H		
BUF_LEN	EQU	20H		
BUF_END	IDATA	BUFFER + BUF_LEN - 1		
XSEG1	SEGMENT :	XDATA		
RSEG	XSEG1			
DTIM:	DS	6	;	reserve 6-bytes for DTIM
TIME	XDATA	DTIM + 0		
DATE	XDATA	DTIM + 3		

<mark>esfr</mark>, sfr, sfr16, sbit

The **sfr**, **sfr16** and **sbit** directives are fully compatible to the **Cx51** compiler and allows you to use a generic SFR register definition file for both: the **Ax51** macro assembler and the **Cx51** compiler. The **esfr** directive defines symbols in the extended SFR space of the Philips 80C51MX architecture. This directive is only available in the AX51 macro assembler. These directives have the following format:

```
sfr sfr_symbol = address;
esfr sfr_symbol = address;
sfr16 sfr_symbol = address; ; ignored by Ax51
sbit sfr symbol = bit-address;
```

where

sfr symbol is the name of the special function register (SFR) symbol to

define.

address is an SFR address in the range 0x80 - 0xFF.

bit-address is address of an SFR bit in the format address hitpos or

sfr_symbol ^ bitpos. address or sfr_symbol refers to an bit-addressable SFR and bitpos specifies the bit position of the

SFR bit in the range 0-7.

Symbols defined with the **esfr**, **sfr**, or **sbit** directive may be used anywhere as address of a SFR or SFR bit.

Example

```
sfr P0 = 0x80;
sfr P1 = 0x90;
sbit P0_0 = P0^0;
sbit P1_1 = 0x90^1;
esfr MXCON = 0xFF;    /* extended Philips 80C51MX SFR */
sfr16 T2 = 0xCC;    /* ignored by Ax51 */
```

NOTE

The Ax51 assembler ignores symbol definitions that start with sfr16. This is implemented for compatibility to the Cx51 compiler.

LIT (AX51 & A251 only)

The **LIT** directive provides a simple text substitution facility. The **LIT** directive has the following format:

symbol LIT 'literal string'
symbol LIT "literal string"

where

symbol is the name of the symbol to define. The literal string

specified in the LIT directive will be substituted for each occurrence of symbol that is used in your assembly

program.

literal string is a numeric expression which contains no forward

references, or a simple relocatable expression.

Every time the symbol is encountered, it is replaced by the literal string assigned to the symbol name. The symbol name follows the same rules as other identifiers, that is, a literal name is not encountered if it does not form a separate token. If a substring is to be replaced, symbol must be enclosed in braces: TEXT{symbol}. The assembler listing shows the expanded lines where literals are substituted.

Example

Source text containing literals before assembly:

```
$INCLUDE (REG51.INC)
REG1 LIT
              'R1'
     LIT
              'A1'
NUM
DBYTE LIT
              "DATA BYTE"
       LIT
FLAG
              'ACC.3'
?PR?MOD SEGMENT CODE
       RSEG ?PR?MOD
              REG1,#5
       MOV
       SETB
              FLAG
             FLAG, LAB {NUM}
       JB
       PUSH DBYTE 0
LAB {NUM}:
       END
```

Assembler listing from previous example:

	1	\$INCLUDE (REG51.INC)
	+1 80	+1 \$RESTORE
	81	
	82	REG1 LIT 'R1'
	83	NUM LIT 'A1'
	84	DBYTE LIT "DATA BYTE"
	85	FLAG LIT 'ACC.3'
	86	12.10 2-1 1.0010
	87	?PR?MOD SEGMENT CODE
		PRIMOD SEGMENT CODE
	88	RSEG ?PR?MOD
	89	
000000 7E1005	90	MOV R1,#5
000003 D2E3	91	SETB ACC.3
000005 20E300	F 92	JB ACC.3,LAB A1
000008 C000	93	PUSH DATA BYTE 0
00000A	94	LAB A1:
	95	_
	96	END

Memory Initialization

The memory initialization directives are used to initialize code or constant space in either word, double-word, or byte units. The memory image starts at the point indicated by the current value of the location counter in the currently active segment.

DB

The **DB** directive initializes code memory with 8-bit byte values. The **DB** directive has the following format:

label:	DB	expression	,	expression		
where						
label		is the sy memory		ol that is give	en the address of the in	itialized
expressi	on	•		lue. Each ex	xpression may be a sy	mbol, a

The **DB** directive can only be specified within a code or const segment. If the **DB** directive is used in a different segment, **Ax51** will generate an error message.

```
REQUEST: DB 'PRESS ANY KEY TO CONTINUE', 0
TABLE: DB 0,1,8,'A','0', LOW(TABLE),';'
ZERO: DB 0,''''
CASE_TAB: DB LOW(REQUEST), LOW(TABLE), LOW(ZERO)
```

DW

The **DW** directive initializes code memory with 16-bit word values. The **DW** directive has the following format:

label: DW	expression , expression
where	
label	is the symbol that is given the address of the initialized memory.
expression	is the initialization data. Each expression may contain a symbol, a character string, or an expression.

The **DW** directive can only be specified within a code or const segment. If the **DW** directive is used in a different segment, **Ax51** will generate an error message.

```
TABLE: DW TABLE, TABLE + 10, ZERO
ZERO: DW 0
CASE_TAB: DW CASE0, CASE1, CASE2, CASE3, CASE4
DW $
```

DD (AX51 & A251 only)

The **DD** directive initializes code memory with 32–bit double word values. The **DD** directive has the following format:

label:	DD	expression	,	expression	2	
where						
label		is the sy memory		U	ve	en the address of the initialized
expressi	on					Each expression may contain a ng, or an expression.

The **DD** directive can only be specified within a code or const segment. If the **DD** directive is used in a different segment, **Ax51** will generate an error message.

TABLE:	DD	TABLE, TABLE + 10, ZERO
	DD	\$
ZERO:	DD	0
LONG VAL:	DD	12345678H, OFFFFFFFH, 1

Reserving Memory

The memory reservation directives are used to reserve space in either word, dword, byte, or bit units. The space reserved starts at the point indicated by the current value of the location counter in the currently active segment.

DBIT

The **DBIT** directive reserves space in a bit or ebit segment. The **DBIT** directive has the following format:

label:	DBIT	expression
where		
label		is the symbol that is given the address of the reserved memory. The label is a symbol of the type BIT and gets the current address value and the memory class of the active segment. The label can only be used where a symbol of this type is allowed.
expressio	on	is the number of bits to reserve. The expression cannot contain forward references, relocatable symbols, or external symbols.

The **DBIT** directive reserves space in the bit segment starting at the current address. The location counter for the bit segment is increased by the value of the *expression*. You should note that the location counter for the bit segment references bits and not bytes.

NOTE

The Ax51 assembler is a two-pass assembler. Symbols are collected and the length of each instruction is determined in the first pass. In the second pass, forward references are resolved and object code is produced. For these reasons, an expression used with the DBIT directive may not contain forward references.

ON_FLAG:	DBIT	1	; reserve 1 bit
OFF_FLAG:	DBIT	1	

DS

The **DS** directive reserves a specified number of bytes in a memory space. The **DS** directive has the following format:

label:	DS	expression
where		
label		is the symbol that is given the address of the reserved memory. The label is a typeless number and gets the current address value and the memory class of the active segment. The label can only be used where a symbol of this type is allowed.
express.	ion	is the number of bytes to reserve. The expression cannot contain forward references, relocatable symbols, or external symbols.

The **DS** directive reserves space in the current segment at the current address. The current address is then increased by the value of the <code>expression</code>. The sum of the location counter and the value of the specified <code>expression</code> should not exceed the limitations of the current address space.

NOTE

The A251 assembler is a two-pass assembler. Symbols are collected and the length of each instruction is determined in the first pass. In the second pass, forward references are resolved and object code is produced. For these reasons, an expression used with the **DS** directive may not contain forward references.

```
GAP: DS (($ + 16) AND 0FFF0H) - $
DS 20
TIME: DS 8
```

DSB (AX51 & A251 only)

The **DSB** directive reserves a specified number of bytes in a memory space. The **DSB** directive has the following format:

label:	DSB	expression
where		is the symbol that is given the address of the reserved memory. The label is a symbol of the type BYTE and gets
		the current address value and the memory class of the active segment. The label can only be used where a symbol of this type is allowed.
expressi	on	is the number of bytes to reserve. The expression cannot contain forward references, relocatable symbols, or external symbols.

The **DSB** directive reserves space in the current segment at the current address. The current address is then increased by the value of the *expression*. The sum of the location counter and the value of the specified *expression* should not exceed the limitations of the current address space.

NOTE

The Ax51 assembler is a two-pass assembler. Symbols are collected and the length of each instruction is determined in the first pass. In the second pass, forward references are resolved and object code is produced. For these reasons, an expression used with the DSB directive may not contain forward references.

DAY:	DSB	1
MONTH:	DSB	1
HOUR:	DSB	1
MIN:	DSB	1

DSW (AX51 & A251 only)

The **DSW** directive reserves a specified number of words in a memory space. The **DSW** directive has the following format:

label:	DSW	expression
where label		is the symbol that is given the address of the reserved memory. The label is a symbol of the type WORD and gets the current address value and the memory class of the active
		segment. The label can only be used where a symbol of this type is allowed.
express	ion	is the number of bytes to reserve. The expression cannot contain forward references, relocatable symbols, or external symbols.

The **DSW** directive reserves space in the current segment at the current address. The current address is then increased by the value of the *expression*. The sum of the location counter and the value of the specified *expression* should not exceed the limitations of the current address space.

NOTE

The Ax51 assembler is a two-pass assembler. Symbols are collected and the length of each instruction is determined in the first pass. In the second pass, forward references are resolved and object code is produced. For these reasons, an expression used with the DSW directive may not contain forward references.

DSD (AX51 & A251 only)

The **DSD** directive reserves a specified number of double words in a memory space. The **DSD** directive has the following format:

label:	DSD	expression
where		
label		is the symbol that is given the address of the reserved memory. The label is a symbol of the type DWORD and gets the current address value and the memory class of the active segment. The label can only be used where a symbol of this type is allowed.
expression	on	is the number of bytes to reserve. The expression cannot contain forward references, relocatable symbols, or external symbols.

The **DSD** directive reserves space in the current segment at the current address. The current address is then increased by the value of the **expression**. The sum of the location counter and the value of the specified **expression** should not exceed the limitations of the current address space.

NOTE

The Ax51 assembler is a two-pass assembler. Symbols are collected and the length of each instruction is determined in the first pass. In the second pass, forward references are resolved and object code is produced. For these reasons, an expression used with the DSD directive may not contain forward references.

SEC_CNT:	DSD	1	
LONG_ARR:	DSD	50	

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Procedure Declaration (AX51 & A251 only)

Ax51 provides procedures to implement the concept of subroutines. Procedures can be executed in-line (control "falls through" to them), jumped to, or invoked by a CALL. Calls are recommended as a better programming practice.

PROC / ENDP (AX51 & A251 only)

The PROC and ENDP directives are used to define a label for a sequence of machine instructions called a procedure. For the Philips 80C51MX and Intel/Atmel WM 251 architecture a procedure may have either the type NEAR or FAR. Depending on the type it is called with LCALL or ACALL (for NEAR) or ECALL (for FAR). Unlike C functions, assembler procedures do not provide local scopes for labels. Identifiers must be unique in A251 because the visibility is module wide. The format of the PROC/ENDP directives is:

```
name PROC [ type ]
    ; procedure text
:
    RET
name ENDP
```

where

name

is the name of the procedure.

type

specifies the type of the procedure, and must be one of the following:

Туре	Description
none	The type defaults to NEAR
NEAR	Defines a near procedure; called with LCALL or ACALL.
FAR	Defines a far procedure; called with ECALL.

You should specify FAR if the procedure is called from a different 64KByte segment. A procedure normally ends with a RET instruction. The software instruction RET will automatically be converted to an appropriate machine return instruction. For example:

RET Return from a near procedure.

ERET Return from a far procedure.

P100	PROC RET ENDP	NEAR	; near return
P200	PROC RET ENDP	FAR	; far return (ERET)
P300	PROC CALL CALL RET ENDP	NEAR P100 P200	; LCALL ; ECALL ; near return

LABEL (AX51 and A251 only)

A label is a symbol name for an address location in a segment. The LABEL directive can be used to define a program label. The label name can be followed by a colon, but it is not required. The label inherits the attributes of the program or code segment currently active. The LABEL directive may therefore never be used outside the scope of a program segment. The syntax of that directive is:

name[:]	LABEL	[type]							
,									
where									
name	is the name	is the name of the label.							
type	specifies th	specifies the type of the label, and must be one of the following:							
	Туре	Description							

Туре	Description
none	The type defaults to NEAR
NEAR	Defines a near label.
FAR	Defines a far label; use ECALL or EJMP.

You should specify FAR if the label will be referenced from a different 64KByte segment. NEAR lets you refer to this label for the current 64KByte segment.

```
RSEG ECODE_SEG1 ; activate an ECODE segment
ENTRY: LABEL FAR ; entry point

RSEG ECODE_SEG2 ; activate another ECODE segment
EJMP ENTRTY ; Jump across 64KB segment
```

Program Linkage

Program linkage directives allow the separately assembled modules to communicate by permitting inter-module references and the naming of modules.

PUBLIC

The **PUBLIC** directive lists symbols that may be used in other object modules. The **PUBLIC** directive makes the specified symbols available in the generated object module. This, in effect, publicizes the names of these symbols. The **PUBLIC** directive has the following format:

PUBLIC symbol , symbol ...

where

symbol

must be a symbol that was defined somewhere within the source file. Forward references to symbol names are permitted. All symbol names, with the exception of register symbols and segment symbols, may be specified with the **PUBLIC** directive. Multiple symbols must be separated with a comma (,).

If you want to use public symbols in other source files, the **EXTRN** or **EXTERN** directive must be used to specify that the symbols are declared in another object module.

Example

PUBLIC PUT_CRLF, PUT_STRING, PUT_EOS
PUBLIC ASCBIN, BINASC
PUBLIC GETTOKEN, GETNUMBER

EXTRN / EXTERN

The **EXTRN** and **EXTERN** directives list symbols (referenced by the source module) that are actually declared in other modules. The format for the **EXTRN** and **EXTERN** directives is as follows:

EXTRN	class	: type	(symbol	,	symbol)
EXTERN	class	: type	(symbol	,	symbol)
where class		may be o	one of the f EBIT, ECC T, IDATA	ollo ON:	owing: B ST, EDA	symbol has been defined and BIT, CODE, CONST, TA, ECODE, HDATA, or NUMBER (to specify a
type		is the syr	nhol type o	of th	ne extern	al symbol and may be one of

is the symbol type of the external symbol and may be one of the following: **BYTE**, **WORD**, **DWORD**, **NEAR**, **FAR**.

symbol is an external symbol name.

The **EXTRN** or **EXTERN** directive may appear anywhere in the source program. Multiple symbols may be separated and included in parentheses following the class and type information.

Symbol names that are specified with the **EXTRN** / **EXTERN** directive must have been specified as public symbols with the **PUBLIC** directive in the source file in which they were declared.

The Linker/Locator resolves all external symbols at link time and verifies that the symbol class and symbol types (specified with the EXTRN / EXTERN and PUBLIC directives) match. Symbols with the class NUMBER match every memory class.

```
EXTRN CODE (PUT_CRLF), DATA (BUFFER)

EXTERN CODE (BINASC, ASCBIN)

EXTRN NUMBER (TABLE_SIZE)

EXTERN CODE:FAR (main)

EXTRN EDATA:BYTE (VALUE, COUNT)

EXTRN NCONST:DWORD (LIMIT)
```

NAME

The **NAME** directive specifies the name to use for the object module generated for the current program. The filename for the object file is not the object module name. The object module name is embedded within the object file. The format for the **NAME** directive is as follows:

NAME modulename

where

modulename is the name to use for the object module and can be up to 40

characters long. The modulename must adhere to the rules

for symbol names.

If a **NAME** directive is not present in the source program, the object module name will be the *basename* of the source file without the extension.

NOTE

Only one NAME directive may be specified in a source file.

Example

NAME PARSERMODULE

Address Control

The following directives allow the control of the address location counter or the control of absolute register symbols.

ORG

The **ORG** directive is used to alter the location counter of the currently active segment and sets a new origin for subsequent statements. The format for the **ORG** statement is as follows:

ORG expression

where

expression

must be an absolute or simple relocatable expression without any forward references. Only absolute addresses or symbol values in the current segment may be used.

When an **ORG** statement is encountered, the assembler calculates the value of the *expression* and changes the location counter for the current segment. If the **ORG** statement occurs in an absolute segment, the location counter is assigned the absolute address value. If the **ORG** statement occurs in a relocatable segment, the location counter is assigned the offset of the specified *expression*.

The **ORG** directive changes the location counter but does not produce a new segment. A possible address gap may be introduced in the segment. With absolute segments, the location counter may not reference an address prior to the base address of the segment.

NOTE

The Ax51 assembler is a multi-pass assembler. Symbols are collected and the length of each instruction is determined in the first pass. In the second pass, forward references are resolved and object code is produced. For these reasons, an expression used with the ORG directive may not contain forward references.

Example

ORG 100H ORG RESTART

Shaded directives and options are available only in AX51 and A251.

```
ORG EXTI1
ORG ($ + 16) AND OFFFOH
```

EVEN (AX51 and A251 only)

The **EVEN** directive ensures that code or data following EVEN is aligned on a word boundary. The assembler creates a gap of one byte if necessary. The content of the byte gap is undefined. The **EVEN** directive has the following syntax:

```
EVEN
```

Example

```
MYDATA SEGMENT DATA WORD ; word alignment
RSEG MYDATA ; activate segment
var1: DSB 1 ; reserve a byte variable
EVEN ; ensure word alignment
var2: DSW 1 ; reserve a word variable
```

USING

The **USING** directive specifies which register bank to use for coding the **AR0** through **AR7** registers. The **USING** directive is specified as follows:

```
where

expression is the register bank number which must be a value between 0 and 3.
```

The **USING** directive does not generate any code to change the register bank. Your program must make sure the correct register bank is selected. For example, the following code can be used to select register bank 2:

```
POP PSW ;restore PSW/register bank
```

The register bank selected by the **USING** directive is marked in the object file and the memory area required by the register bank is reserved by the Linker/Locator.

The value of **AR0** through **AR7** is calculated as the absolute address of **R0** through **R7** in the register bank specified by the **USING** directive. Some 8051 instruction (i.e. PUSH / POP) allow you to use only absolute register addresses. By default register bank 0 is assigned to the symbols **AR0** through **AR7**.

NOTE

When the **EQU** directive is used to define a symbol for an **AR**n register, the address of the register **R**n is calculated when the symbol is defined; not when it is used. If the **USING** directive subsequently changes the register bank, the defined symbol will not have the proper address of the **AR**n register and the generated code is likely to fail.

USING PUSH	3 AR2	;	Push	register	2	in	bank	3
USING PUSH	1 AR2	;	Push	register	2	in	bank	1

Other Directives

END

The **END** directive signals the end of the assembly module. Any text in the assembly file that appears after the **END** directive is ignored.

The **END** directive is required in every assembly source file. If the **END** statement is excluded, **Ax51** will generate a warning message.

Example

END

The __ERROR__ directive generates standard error messages that are report the same style as normal Ax51 assembler errors. The __ERROR__ directive is specified as follows:

```
where

text is the error text that should be displayed in the listing file.

The error text is also displayed on the console if the

"ERRORPRINT" control described on page 205 is used.
```

```
IF CVAR1LEN > 10
__ERROR__ "CVAR1 LEN EXCEEDS 10 BYTES"
ENDIF

$IF TESTVERS AND RELEASE
__ERROR__ "TESTVERS GENERATED IN RELEASE MODE"
$ENDIF
```

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Chapter 5. Assembler Macros

A macro is a name that you assign to one or more assembly statements. For maximum flexibility the **Ax51** macro assembler provides three different macro languages:

- Standard Assembler Macros: are known from many other macro assemblers and allow you to define macros that look like standard assemblers instructions. Refer to "Standard Macro Directives" on page 139 for a detailed description.
- C Macros: are known from ANSI C compilers and allow you to use common header files with constant definitions that can be used on the Ax51 macro assembler as well as on the Cx51 compiler. Refer to "C Macros" on page 156 for more information.
- MPL Macros: are compatible with the Intel ASM-51 and allow you to retranslate existing source files that initially written for this macro assembler. The assembler control MPL enables this macro processor. If you enable MPL macros the C Macros are disabled. Refer to "Chapter 6. Macro Processing Language" on page 163 for a detailed description.

A macro processor enables you to define and to use macros in your **x51** assembly programs. This section describes some of the features and advantages of using macros, lists the directives and operators that are used in macro definitions, and provides a number of example macros.

When you define a macro, you provide text (usually assembly code) that you want to associate with a macro name. Then, when you want to include the macro text in your assembly program, you provide the name of the macro. The **Ax51** assembler will replace the macro name with the text specified in the macro definition

Macros provide a number of advantages when writing assembly programs.

- The frequent use of macros can reduce programmer induced errors. A macro allows you to define instruction sequences that are used repetitively throughout your program. Subsequent use of the macro will faithfully provide the same results each time. A macro can help reduce the likelihood of errors introduced in repetitive programming sequences. Of course, introduction of an error into a macro definition will cause that error to be duplicated where the macro is used.
- The scope of symbols used in a macro is limited to that macro. You do not need to be concerned about utilizing a previously used symbol name.
- Macros are well suited for the creation of simple code tables. Production of these tables by hand is both tedious and error prone.

A macro can be thought of as a subroutine call with the exception that the code that would be contained in the subroutine is included in—line at the point of the macro call. However, macros should not be used to replace subroutines. Each invocation of a subroutine only adds code to call the subroutine. Each invocation of a macro causes the assembly code associated with the macro to be included in—line in the assembly program. This can cause a program to grow rapidly if a large macro is used frequently. In a static environment, a subroutine is the better choice, since program size can be considerably reduced. But in time critical, dynamic programs, macros will speed the execution of algorithms or other frequently called statements without the penalty of the procedure calling overhead.

Use the following guidelines when deciding between macros or subroutines:

- Subroutines are best used when certain procedures are frequently executed or when memory space usage must be kept to a minimum.
- Macros should be used when maximum processor speed is required and when memory space used is of less importance.
- Macros can also be used to make repetitive, short assembly blocks more convenient to enter.

Standard Macro Directives

Ax51 provides a number of directives that are used specifically for defining macros. These directives are listed in the following table:

Directive	Description
ENDM	Ends a macro definition.
EXITM	Causes the macro expansion to immediately terminate.
IRP	Specifies a list of arguments to be substituted, one at a time, for a specified parameter in subsequent lines.
IRPC	Specifies an argument to be substituted, one character at a time, for a specified parameter in subsequent lines.
LOCAL	Specifies up to 16 local symbols used within the macro.
MACRO	Begins a macro definition and specifies the name of the macro and any parameters that may be passed to the macro.
REPT	Specifies a repetition factor for subsequent lines in the macro.

Refer to "Assembler Controls" on page 197 as well as the following sections for more information on these and other directives.

Defining a Macro

Macros must be defined in the program before they can be used. A macro definition begins with the MACRO directive which declares the name of the macro as well as the formal parameters. The macro definition must be terminated with the ENDM directive. The text between the MACRO and ENDM directives is called the macro body.

Example

WAIT	MACRO	X	; macro definition
	REPT	X	; generate X NOP instructions
	NOP		
	ENDM		; end REPT
	ENDM		; end MACRO

In this example, walt is the name of the macro and x is the only formal parameter.

In addition to the **ENDM** directive, the **EXITM** directive can be used to immediately terminate a macro expansion. When an **EXITM** directive is detected, the macro processor stops expanding the current macro and resumes processing after the next **ENDM** directive. The **EXITM** directive is useful in conditional statements.

```
WAIT
                   MACRO
                                           ; macro definition
                      IF NUL X
                                           ; make sure X has a value
                      EXITM
                                           ; if not then exit
                      ENDIF
                      REPT
                                   Х
                                           ; generate X NOP instructions
                      NOP
                                           ; end REPT
                      ENDM
                                           ; end MACRO
                    ENDM
```

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Parameters

Up to 16 parameters can be passed to a macro in the invocation line. Formal parameter names must be defined using the **MACRO** directive.

Example

```
MNAME MACRO P1, P2, P3, P4, P5, P6, P7, P8, P9, P10, P11, P12, P13, P14, P15, P16
```

defines a macro with 16 parameters. Parameters must be separated by commas in both the macro definition and invocation. The invocation line for the above macro would appear as follows:

```
MNAME A,B,C,D,E,F,G,H,I,J,K,L,M,N,O,P
```

where A, B, C, ... O, P are parameters that correspond to the format parameter names P1, P2, P3, ... P15, P16.

Null parameters can be passed to a macro. Null parameters have the value NULL and can be tested for using the NUL operator described later in this chapter. If a parameter is omitted from the parameter list in the macro invocation, that parameter is assigned a value of NULL.

Example

```
MNAME A,,C,,E,,G,,I,,K,,M,,O,
```

P2, P4, P6, P8, P10, P12, P14, and P16 will all be assigned the value NULL when the macro is invoked. You should note that there are no spaces between the comma separators in the above invocation line. A space has an ASCII value of 20h and is not equivalent to a NULL.

Labels

You can use labels within a macro definition. By default, labels used in a macro are global and if the macro is used more than once in a module, **Ax51** will generate an error.

Example

LOC	OBJ	LINE	SOURCE	l				
		1	GLABEL	MACRO				
		2	LOOP:	NOP				
		3		JMP	LOOP			
		4		ENDM				
		5						
		6						
		7		GLABEL				
0000	00	8+1	LOOP:	NOP				
0001	80FD	9+1		JMP	LOOP			
		10		GLABEL				
		11+1	LOOP:	NOP				
***			^					
***	ERROR #	9, LINE #11, AT	TEMPT TO	DEFINE AN	ALREADY	DEFINED	LABEL	
0003	80FB	12+1		JMP	LOOP			
		13						
		14						
		15	END					

Labels used in a macro should be local labels. Local labels are visible only within the macro and will not generate errors if the macro is used multiple times in one source file. You can define a label (or any symbol) used in a macro to be local with the **LOCAL** directive. Up to 16 local symbols may be defined using the **LOCAL** directive.

NOTE

LOCAL must be in the next line after the MACRO definition.

Example

CLRMEM	MACRO	ADDR, LEN
	LOCAL	LOOP
	MOV	R7, #LEN
	MOV	RO, #ADDR
	MOV	A, #0
LOOP:	MOV	@RO, A
	INC	R0
	DJNZ	R7, LOOP
	ENDM	

In this example, the label LOOP is local because it is defined with the LOCAL directive. Any symbol that is not defined using the LOCAL directive will be a global symbol.

Ax51 generates an internal symbol for local symbols defined in a macro. The internal symbol has the form ??0000 and is incremented each time the macro is invoked. Therefore, local labels used in a macro are unique and will not generate errors.

Repeating Blocks

Ax51 provides the ability to repeat a block of text within a macro. The **REPT**, **IRP**, and **IRPC** directives are used to specify text to repeat within a macro. Each of these directives must be terminated with an **ENDM** directive.

REPT

The **REPT** directive repeats a block of text a fixed number of times. The following macro:

DELAY	MACRO		;macro definition
	REPT	5	;insert 5 NOP instructions
	NOP		
	ENDM		;end REPT block
	ENDM		;end macro definition

inserts 5 NOP instructions when it is invoked.

```
NOP
NOP
NOP
NOP
NOP
```

IRP

The **IRP** directive repeats a block once for each argument in a specified list. A specified parameter in the text block is replaced by each argument. The following macro:

```
CLRREGS MACRO ; macro definition
IRP RNUM, <R0,R1,R2,R3,R4,R5,R6,R7>
MOV RNUM, #0
ENDM ; end IRP
ENDM ; end MACRO
```

replaces the argument RNUM with RO, R1, R2, ... R7.

It generates the following code when invoked:

```
MOV
       R0, #0
       R1, #0
MOV
       R2, #0
MOV
MOV
       R3, #0
MOV
       R4, #0
       R5, #0
VOM
       R6, #0
MOV
       R7, #0
MOV
```

IRPC

The **IRPC** directive repeats a block once for each character in the specified argument. A specified parameter in the text block is replaced by each character. The following macro:

```
DEBUGOUT
                MACRO
                                        ; macro definition
                IRPC
                         CHR, <TEST>
                JNB
                         TI, $
                                       ; wait for xmitter
                CLR
                         TI
                        A,#'CHR'
                MOV
                VOM
                         SBUF, A
                                        ; xmit CHR
                 ENDM
                                        ; end IRPC
                                       ; end MACRO
```

replaces the argument CHR with the characters \mathbf{T} , \mathbf{E} , \mathbf{s} , and \mathbf{T} and generates the following code when invoked:

```
TI, $
            ; WAIT FOR XMITTER
JNB
CLR
     TI
MOV
     A,#'T'
MOV
     SBUF, A ; XMIT T
JNB
     TI, $ ; WAIT FOR XMITTER
CLR
     TI
MOV
     A,#'E'
MOV
     SBUF, A ; XMIT E
JNB
     TI, $ ; WAIT FOR XMITTER
CLR
     TI
     A,#'S'
VOM
     SBUF, A ; XMIT S
MOV
JNB
     TI, $ ; WAIT FOR XMITTER
CLR
     TI
     A,#'T'
MOV
MOV
     SBUF, A ; XMIT T
```

Nested Definitions

Macro definitions can be nested up to nine levels deep.

Example

```
L1 MACRO
LOCAL L2
L2 MACRO
INC R0
ENDM
MOV R0, #0
L2
ENDM
```

The macro L2 is defined within the macro definition of L1. Since the LOCAL directive is used to define L2 as a local symbol, it is not visible outside L1. If you want to use L2 outside of L1, exclude L2 from the LOCAL directive symbol list.

Invocation of the L1 macro generates the following:

```
MOV R0, #0
INC R0
```

Nested Repeating Blocks

You can also nest repeating blocks, specified with the **REPT**, **IRP**, and **IRPC** directives

Example

PORTOUT	MACRO IRPC	CHR, <hello></hello>	; macro definition
	REPT NOP	4	; wait for 4 cycles
	ENDM		; end REPT
	MOV MOV	A,#'CHR' P0,A	; write CHR to PO
	ENDM		; end IRPC
	ENDM		; end MACRO

This macro nests a **REPT** block within an **IRPC** block.

Recursive Macros

Macros can call themselves directly or indirectly (via another macro). However, the total number of levels of recursion may not exceed nine. A fatal error will be generated if the total nesting level is greater than nine. The following example shows a recursive macro that is invoked by a non–recursive macro.

RECURSE IF X<>0	MACRO	х	;	recursive macro
	RECURSE	%X-1		
	ADD	A,#X	;	gen add a,#?
ENDIF				
	ENDM			
SUMM	MACRO	X	;	macro to sum numbers
	MOV	A,#0	;	start with zero
IF NUL X			;	exit if null argument
	EXITM			
ENDIF				
IF X=0			;	exit if 0 argument
	EXITM			
ENDIF				
	RECURSE X		;	sum to 0
	ENDM			

Operators

Ax51 provides a number of operators that may be used within a macro definition. The following table lists the operators and gives a description of each.

Operator	Description				
NUL	The NUL operator can be used to determine if a macro argument is NULL. NUL generates a non–zero value if its argument is a NULL. Non–NULL arguments will generate a value of 0. The NUL operator can be used with an IF control to enable condition macro assembly.				
&	The ampersand character is used to concatenate text and parameters.				
<>	Angle brackets are used to literalize delimiters like commas and blanks. Angle brackets are required when passing these characters to a nested macro. One pair of angle brackets is required for every nesting level.				
%	The percent sign is used to prefix a macro argument that should be interpreted as an expression. When this operator is used, the numeric value of the following expression is calculated. That value is passed to the macro instead of the expression text.				
;;	A double semicolon indicates that subsequent text on the line should be ignored. The remaining text is not processed or emitted. This helps to reduce memory usage.				
!	If an exclamation mark is used in front of a character, that character will be literalized. This allows character operators to be passed to a macro as a parameter.				

NUL Operator

When a formal parameter in a macro call is omitted, the parameter is given a value of NULL. You can check for NULL parameters by using the **NUL** operator within an **IF** control statement in the macro. The **NUL** operator requires an argument. If no argument is found, **NUL** returns a value of 0 to the IF control.

For example, the following macro definition:

```
EXAMPLE MACRO X
IF NUL X
EXITM
ENDIF
ENDM
```

when invoked by:

EXAMPLE

will cause the IF NUL x test to pass, process the EXITM statement, and exit the macro expansion.

NOTE

A blank character ('') has an ASCII value of 20h and is not equivalent to a NULL.

& Operator

The ampersand macro operator (&) can be used to concatenate text and macro parameters. The following macro declaration demonstrates the proper use of this operator.

MAK_NOP_LABEL	MACRO	X
LABEL&X:	NOP	
	ENDM	

The MAK_NOP_LABEL macro will insert a new label and a NOP instruction for each invocation. The argument will be appended to the text LABEL to form the label for the line.

Example

LOC	OBJ	LINE	SOURCE		
		1	MAK NOP LABEL	MACRO	x
		2	LABEL&X:	NOP	
		3	ENDM		
		4			
		5			
		6	MAK NOP LABEL	1	
0000	00	7+1	LABEL1:	NOP	
		8	MAK NOP LABEL	2	
0001	00	9+1	LABEL2:	NOP	
		10	MAK NOP LABEL	3	
0002	00	11+1	LABEL3:	NOP	
		12	MAK NOP LABEL	4	
0003	00	13+1	LABEL4:	NOP	
		14			
		15	END		

The MAK_NOP_LABEL macro is invoked in the above example in lines 6, 8, 10, and 12. The generated label and NOP instructions are shown in lines 7, 9, 11, and 13. Note that the labels are concatenated with the argument that is passed in the macro invocation.

< and > Operators

The angle bracket characters (<>) are used to enclose text that should be passed literally to macros. Some characters; for example, the comma; cannot be passed without being enclosed within angle brackets.

The following example shows a macro declaration and invocation passing an argument list within angle brackets.

1	FLAG_CLR	MACRO	FLAGS
2		MOV	A, #0
3		IRP	F, <flags></flags>
4		MOV	FLAG&F, A
5		ENDM	
6		ENDM	
7			
8	DSEG		
0000 9	FLAG1: DS 1		
0001 10	FLAG2: DS 1		
0002 11	FLAG3: DS 1		
0003 12	FLAG4: DS 1		
0004 13	FLAG5: DS 1		
0005 14	FLAG6: DS 1		
0006 15	FLAG7: DS 1		
0007 16	FLAG8: DS 1		
0008 17	FLAG9: DS 1		
18			
19	CSEG		
20			
21	FLAG_CL		<1>
0000 7400 22+1		MOV	A, #0
23+1		IRP	F, <1>
24+1		MOV	FLAG&F, A
25+1		ENDM	
0002 F500 26+2		MOV	FLAG1, A
27	FLAG_CL		<1,2,3>
0004 7400 28+1		MOV	A, #0
29+1		IRP	F, <1,2,3>
30+1		MOV	FLAG&F, A
31+1		ENDM	
0006 F500 32+2		MOV	FLAG1, A
0008 F501 33+2		VOM	FLAG2, A
000A F502 34+2 35	EL VO. CI	MOV	FLAG3, A <1,3,5,7>
000C 7400 36+1	FLAG_CL	MOV	A, #0
37+1		IRP	F, <1,3,5,7>
38+1		MOV	
39+1		ENDM	FLAG&F, A
000E F500 40+2		MOV	FLAG1, A
0010 F502 41+2		MOV	FLAGI, A
0010 F502 41+2 0012 F504 42+2		MOV	FLAG5, A
0012 F504 42+2 0014 F506 43+2		MOV	FLAG7, A
5511 F500 45+2		MOV	I LACI / A
•			
•			

In the previous example, the <code>FLAG_CLR</code> macro is declared to clear any of a number of flag variables. The <code>FLAGS</code> argument specifies a list of arguments that are used by the <code>IRP</code> directive in line 3. The <code>IRP</code> directive repeats the instruction <code>MOV FLAGSF</code>, <code>A</code> for each parameter in the <code>FLAGS</code> argument.

The FLAG_CLR macro is invoked in lines 21, 27, and 35. In line 21, only one parameter is passed. In line 27, three parameters are passed, and in line 35, four parameters are passed. The parameter list is enclosed in angle brackets so that it may be referred to as a single macro parameter, FLAGS. The code generated by the macro is found in lines 26, 32–34, and 40–43.

% Operator

The percent character (%) is used to pass the value of an expression to a macro rather than passing the literal expression itself. For example, the following program example shows a macro declaration that requires a numeric value along with macro invocations that use the percent operator to pass the value of an expression to the macro.

	1 OUTPORT MACRO 2 MOV 3 MOV 4 ENDM 5	N A, #N PO, A
00FF	7 RESET SIG	EQU 0FFh
0000	8 CLEAR SIG	EQU 0
	9 —	
	10	
	11 OUTPORT	%(RESET SIG AND NOT 11110000b)
0000 740F	12+1 MOV	A, #15
0002 F580	13+1 MOV	PO, A
	14	
	15 OUTPORT	%(CLEAR_SIG OR 11110000b)
0004 74F0	16+1 MOV	A, #240
0006 F580	17+1 MOV	PO, A

In this example, the expressions evaluated in lines 11 and 15 could not be passed to the macro because the macro expects a numeric value. Therefore, the expressions must be evaluated before the macro. The percent sign forces **Ax51** to generate a numeric value for the expressions. This value is then passed to the macro.

;; Operator

The double semicolon operator is used to signal that the remaining text on the line should not be emitted when the macro is expanded. This operator is typically used to precede comments that do not need to be expanded when the macro is invoked.

Example

REGCLR	MACRO	CNT		
REGNUM	SET	0		
	MOV	A, #0	;;	load A with 0
	REPT	CNT	;;	rpt for CNT registers
	MOV	R®NUM, A	;;	set R# to 0
	REGNUM	SET %(REGNUM+1)		
	ENDM			
	ENDM			

! Operator

The exclamation mark operator is used to indicate that a special character is to be passed literally to a macro. This operator enables you to pass comma and angle bracket characters, that would normally be interpreted as delimiters, to a macro.

Invoking a Macro

Once a macro has been defined, it can be called many times in the program. A macro call consists of the macro name plus any parameters that are to be passed to the macro.

In the invocation of a macro, the position of the actual parameters corresponds to the position of the parameter names specified in the macro definition. **Ax51** performs parameter substitution in the macro starting with the first parameter. The first parameter passed in the invocation replaces each occurrence of the first formal parameter in the macro definition, the second parameter that is passed replaces the second formal parameter in the macro definition, and so on.

If more parameters are specified in the macro invocation than are actually declared in the macro definition, **Ax51** ignores the additional parameters. If fewer parameters are specified than declared, **Ax51** replaces the missing parameters with a NULL character.

To invoke a macro in your assembly programs, you must first define the macro. For example, the following definition:

```
.
.
DELAY MACRO CNT ; macro definition
REPT CNT ; insert CNT NOP instructions
NOP
ENDM ; end REPT block
ENDM ; end macro definition
.
```

defines a macro called **DELAY** that accepts one argument **CNT**. This macro will generate **CNT** NOP instructions. So, if **CNT** is equal to 3, the emitted code will be:

```
NOP
NOP
```

The following code shows how to invoke the **DELAY** macro from an assembly program.

```
.
.
LOOP: MOV P0, #0 ;clr PORT 0
DELAY 5 ;wait 5 NOPs
MOV P0, #0ffh ;set PORT 0
DELAY 5 ;wait 5 NOPs
JMP LOOP ;repeat
.
```

In this example, a value of 0 is written to port 0. The **DELAY** macro is then invoked with the parameter 5. This will cause 5 NOP instructions to be inserted into the program. A value of 0FFh is written to port 0 and the **DELAY** macro is invoked again. The program then repeats.

C Macros

The **Ax51** macro assembler has a standard C macro preprocessor that is almost identical with the macro preprocessors in the **Cx51** compiler. This allows you to use common header files with constant definitions that can be used in assembler and C source files. The **Ax51** macro assembler accepts also the special function register directives from the **Cx51** compiler. Therefore you may use the same SFR register definition files for both assembler and C source files.

NOTE

C Macros are not available if you have enabled the Intel ASM-51 compatible MPL macro language with the **MPL** assembler control.

C Macro Preprocessor Directives

C macro preprocessor directives must be the first non-whitespace text specified on a line. All directives are prefixed with the pound or number-sign character ('#'). For example:

```
#include <reg51f.h>
#if TEST
    #define DEBUG 1
#endif
```

The following table lists the preprocessor directives and gives a brief description of each.

Directive	Description
define	Defines a preprocessor macro or constant.
elif	Initiates an alternative branch of the if condition, when the previous if, ifdef, ifndef, or elif branch was not taken.
else	Initiates an alternative branch when the previous if, ifdef, or ifndef branch was not taken.
endif	Ends an if, ifdef, ifndef, elif, or else block.
error	Outputs an error message defined by the user. This directive instructs the compiler to emit the specified error message.
ifdef	Evaluates an expression for conditional compilation. The argument to be evaluated is the name of a definition.
ifndef	Same as ifdef but the evaluation succeeds if the definition is not defined.
if	Evaluates an expression for conditional compilation.
include	Reads source text from an external file. The notation sequence determines the search sequence of the included files. Ax51 searches for include files specified with less-than/greater-than symbols ('<' '>') in the include file folder. The include file folder is specified with the INCDIR assembler control and with the environment variable C51INC and is therefore compatible with the Cx51 compiler. Ax51 searches for include files specified with double-quotes (" ") in the current folder, which is typically the folder of the project file.
line	Specifies a line number together with an optional filename. These specifications are used in error messages to identify the error position.
pragma	Allows you to specify assembler controls and are converted into Ax51 control lines. Refer to "Assembler Controls" on page 197 for more information.
undef	Deletes a preprocessor macro or constant definition.

Stringize Operator

The stringize or number-sign operator ('#'), when used within a macro definition, converts a macro parameter into a string constant. This operator may be used only in a macro that has a specified argument or parameter list.

When the stringize operator immediately precedes the name of one of the macro parameters, the parameter passed to the macro is enclosed within quotation marks and is treated as a string literal. For example:

```
#define stringer(x) DB #x, 0x0D, 0x0A
stringer (text)
```

results in the following actual output from the preprocessor.

```
DB "text", 0x0D, 0x0A
```

NOTES

The Ax51 macro assembler does not accept C escape sequences like "\n", "\r" or "\x0d". You need to replace these characters with hex values.

Unlike the **Cx51** compiler, multiple strings are not concatenated to a single string by the **Ax51** macro assembler. Therefore you need to separate multiple items with a comma when using the **Ax51** macro assembler.

Token-pasting Operator

The token-pasting operator (##) within a macro definition combines two arguments. It permits two separate tokens in the macro definition to be joined into a single token.

If the name of a macro parameter used in the macro definition is immediately preceded or followed by the token-pasting operator, the macro parameter and the token-pasting operator are replaced by the value of the passed parameter. Text that is adjacent to the token-pasting operator that is not the name of a macro parameter is not affected. For example:

```
TEST1 EQU 0x10
TEST2 EQU 0x20
#define paster(n) DB TEST##n
paster (2)
```

results in the following actual output from the preprocessor.

DB TEST2

Predefined C Macro Constants

Ax51 provides you with predefined constants to use in preprocessor directives and C code for more portable programs. The following table lists and describes each one.

Constant	Description
A51	Allows you to identify the A51 assembler and returns the version number (for example, 600 for version 6.00). Only defined when using A51 .
AX51	Allows you to identify the AX51 assembler and returns the version number (for example, 100 for version 1.00). Only defined when using AX51 .
A251	Allows you to identify the A251 assembler and returns the version number (for example, 300 for version 3.00). Only defined when using A251 .
DATE	Date when the compilation was started.
FILE	Name of the file being compiled.
KEIL	Defined to 1 to indicate that you are using a development tool from Keil Software.
LINE	Current line number in the file being compiled.
TIME	Time when the compilation was started.
STDC	Defined to 1 to indicate full conformance with the ANSI C Standard.

Examples with C Macros

The following assembler source file shows the usage of C Macros.

The listing file generated by A51 shows the text replacements performed by the C preprocessor:

LOC	OBJ	LINE	SOURCE		
		1			
		4			
		117	\$LIST		
		118			
		119			
		120			
		121			
		122			
0000	48656C6C	123		DB	"Hello World"
0004	6F20576F				
0008	726C64				
		124			
000B	47454E45	125		DB	"GENERATED: ", "Jul 28 2000"
	52415445				
	443A204A				
	756C2032				
	38203230				
	3030				
0011		126			
0021	7864	127		MOV	R0,#10 * 10
0021	7001	128		110 4	Λ0,π10 " 10
		129	END		
		129	FND		

C Preprocessor Side Effects

The integrated C preprocessor in **Ax51** has two side effects. This might cause problems when you translate programs that are written for previous Ax51 versions.

1. If you are using the backslash character at the end of a comment line, the next line will be a comment too.

```
; THIS IS A COMMENT ENDING WITH \
MOV A,#0 DUE TO THE \ IN THE PREVIOUS LINE THE LINES A CONCATINATED
; AND THE MOV INSTRUCTION WILL NOT BE TRANSLATED
```

2. If you are using \$INCLUDE in conditional assembly blocks, the file must exist even when the block will not be assembled.

Chapter 6. Macro Processing Language

The Macro Processing Language (MPL) is a string replacement facility. The macro processing language is enabled with the assembler control MPL and fully compatible to the Intel ASM-51 macro processing language. It permits you to write repeatedly used sections of code once and then insert that code at several places in your program. Perhaps MPL's most valuable capability is conditional assembly-with all microprocessors, compact configuration dependent code is very important to good program design. Conditional assembly of sections of code can help to achieve the most compact code possible.

Overview

The MPL processor views the source file in different terms than the assembler: to the assembler, the source file is a series of lines – control lines, and directive lines. To the MPL processor, the source file is a long string of characters.

All MPL processing of the source file is performed before your code is assembled. Because of this independent processing of the MPL macros and assembly of code, we must differentiate between macro-time and assembly-time. At macro-time, assembly language symbols and labels are unknown. SET and EQU symbols, and the location counter are also not known. Similarly, at assembly-time, no information about the MPL is known.

The MPL processor scans the source file looking for macro calls. A macro call is a request to the processor to replace the macro name of a built-in or user-defined macro by some replacement text.

6

Creating and Calling MPL Macros

The MPL processor is a character string replacement facility. It searches the source file for a macro call, and then replaces the call with the macro's return value. A % character signals a macro call.

The MPL processor function DEFINE creates macros. MPL processor functions are a predefined part of the macro language, and can be called without definition. The syntax for DEFINE is:

%[*]DEFINE (macro name) [parameter-list] (macro-body)

DEFINE is the most important macro processor function. Each of the symbols in the syntax above (macro name, parameter-list, and macro-body) are described in the following.

Creating Parameterless Macros

When you create a parameterless macro, there are two parts to a DEFINE call:

- macro name
 The macro name defines the name used when the macro is called.
- macro body
 The macro-body defines the return value of the call.

The syntax of a parameterless macro definition is shown below:

```
%*DEFINE (macro name) (macro-body)
```

The '%' is the metacharacter that signals a macro call. The '*' is the literal character. The use of the literal character is described later in this part.

Macro names have the following conventions:

- Maximum of 31 characters long
- First character: 'A' 'Z', 'a' 'z', '_', or '?'
- Other characters: 'A' 'Z', 'a' 'z', ', '?', '0' '9'

The macro-body is usually the replacement text of the macro call. However, the macro-body may contain calls to other macros. If so, the replacement text is actually the fully expanded macro-body, including the calls to other macros. When you define a macro using the syntax shown above, macro calls contained in the body of the macro are not expanded, until you call the macro.

The syntax of DEFINE requires that left and right parentheses surround the macro-body. For this reason, you must have balanced parentheses within the macro-body (each left parenthesis must have a succeeding right parenthesis, and each right parenthesis must have a preceding left parenthesis). We call character strings that meet these requirements balanced-text.

To call a macro, use the metacharacter followed by the macro name for the MPL macro. (The literal character is not needed when you call a user-defined macro.) The MPL processor will remove the call and insert the replacement text of the call. If the macro- body contains any call to other macros, they will be replaced with their replacement text.

Once a macro has been created, it may be redefined by a second DEFINE.

6

MPL Macros with Parameters

Parameters in a macro body allow you to fill in values when you call the MPL macro. This permits you to design a generic macro that produces code for many operations.

The term parameter refers to both the formal parameters that are specified when the macro is defined, and the actual parameters or arguments that are replaced when the macro is called.

The syntax for defining MPL macros with parameters is:

```
%*DEFINE (macro-name(parameter-list)) (macro-body)
```

The parameter-list is a list of identifiers separated by macro delimiters. The identifier for each parameter must be unique.

Typically, the macro delimiters are parentheses and commas. When using these delimiters, you would enclose the parameter-list in parentheses and separate each formal parameter with a comma. When you define a macro using parentheses and commas as delimiters, you must use those same delimiters, when you call that macro

The macro-body must be a balanced-text string. To indicate the locations of parameter replacement, place the parameter's name preceded by the metacharacter in the macro-body. The parameters may be used any number of times and in any order within the macro-body. If a macro has the same name as one of the parameters, the macro cannot be called within the macro-body since this would lead to infinite recursion

The example below shows the definition of a macro with three dummy parameters - SOURCE, DESTINATION, and COUNT. The macro will produce code to copy any number of bytes from one part of memory to another.

To call the above macro, you must use the metacharacter followed by the macro's name similar to simple macros without parameters. However, a list of the actual parameters must follow. The actual parameters must be surrounded in the macro definition. The actual parameters must be balanced-text and may optionally contain calls to other macros. A simple program example with the macro defined above might be:

Assembler source text

```
%*DEFINE (BMOVE (src, dst, cnt)) LOCAL lab (
     MOV R0, #%src
     MOV R1,#%dst
     MOV R2, #%cnt
%lab: MOV A,@R0
     MOV @R1,A
     INC RO
     INC R1
     DJNZ R2, %lab
ALEN EQU 10 ; define the array size
DSEC SEGMENT IDATA; define a IDATA segment
PSEC SEGMENT CODE ; define a CODE segment
RSEG DSEC ; activate IDATA segment arr1: DS ALEN ; define arrays
arr2: DS ALEN
     RSEG PSEC ; activate CODE segment
; move memory block
%BMOVE (arr1, arr2, ALEN)
      END
```

The following listing shows the assembler listing of the above source code.

```
LOC
        OBJ
                 LINE
                            SOURCE
                  3 ALEN EQU 10 ; define the array size
4 DSEC SEGMENT IDATA; define a IDATA segment
5 PSEC SEGMENT CODE ; define a CODE segment
A0000A
_____
                    6
                    7 RSEG DSEC ; activate IDATA segment
8 arrl: DS ALEN ; define arrays
9 arr2: DS ALEN
000000
A0000A
                    9
                   10
                                   RSEG PSEC ; activate CODE segment
                   11
                          ; move memory block
                    13
                            ; %BMOVE (arr1,arr2,ALEN)
                    14
                    15
                                    MOV R0, #%src
                           ;
                    16
                          ;
                                   MOV R1,#%dst
                    17
                                  MOV R2, #%cnt
                    18 ; %lab: MOV A,@R0
                    19
                          ; MOV @R1,A
                                  INC RO
                    20
                          ;
                    21
                                  INC R1
                           ;
                    22
                                 DJNZ R2, %lab
                           ;
                    23
                          24
                    25
000000 7E0000 F 26
                    27
                   28 ;
29 MOV R1,#arr2
MOV R2,#%cn
                                              arr2
000003 7E1000 F 29
                   30 ;
                                  MOV R2, #%cnt
31 ; ALI
000006 7E200A 32 MOV R2,#ALEI
33 ; %lab: MOV A,@R0
                          ; ALEI
MOV R2,#ALEN
                                             ALEN
34 ;LAB0
000009 A5E6 35 LAB0: MOV A,@n
00000B A5F7 36 MOV @R1,A
00000D A508 37 INC R0
00000F A509 38 INC D1
39 ; DJNZ R2, %lah
40 ; LAB(
000011 A5DA00 F 41 DJNZ R2, LAB0
                                  DJNZ R2, %lab
                                              LAB0
                    42
                    43
                                   END
```

The example lists an assembled file that contains a macro definition in lines 1 to 9. The macro definition is listed with semicolons at start of each line. These semicolons are added by the assembler to prevent assembly of the definition text which is meaningful to the MPL preprocessor, but not to the remaining assembler phases. The listing only includes macro definitions or macro calls, if the control **GEN** is given.

The macro BMOVE is called in line 12 with three actual parameters. Lines 14 to 20 shows the macro expansion, which is the return value of the macro call. This text will be assembled.

The example will produce assembly errors because no section directives are included in the source file. The purpose here is to show MPL processing, not the assembler semantics.

Local Symbols List

The **DJNZ** instruction in the previous example uses a local label as the target of the branch. If you use a fixed label name (for example xlab, without a leading %), and you use the macro two or more times is the same assembly source file, errors will occur due to multiple definitions of a single name.

Local symbol definitions solve this problem. Local symbols are generated by the MPL processor as *local_symbol_nnn*, whereby *local_symbol* is the name of the local symbol and *nnn* is some number. Each time the macro is called, the number is automatically incremented. The resulting names are unique to each macro invocation.

The MPL processor increments a counter each time your program calls a macro that uses a LOCAL construct. The counter is incremented once for each symbol in the LOCAL list. Symbols in the LOCAL list, when used in the macro-body, receive a one to five digit suffix that is the decimal value of the counter. The first time you call a macro that uses the LOCAL construct, the suffix is 0.

The syntax for the LOCAL construct in the DEFINE functions is shown below:

%*DEFINE (macro-name (parameter-list)) [LOCAL local-list] (macro-body)

The local-list is a list of valid macro identifiers separated by spaces or commas. The LOCAL construct in a macro has no affect on the syntax of a macro call.

Macro Processor Language Functions

The MPL processor has several predefined macro processor functions. These MPL processor functions perform many useful operations that would be difficult or impossible to produce in a user-defined macro. An important difference between a user-defined macro and a MPL processor function is that user-defined macros may be redefined, while MPL processor functions can not be redefined.

We have already seen one of these MPL processor functions, DEFINE. DEFINE creates user defined macros. MPL processor functions are already defined when the MPL processor is started.

Comment Function

The MPL processing language can be very subtle, and the operation of macros written in a straightforward manner may not be immediately obvious. Therefore, it is often necessary to comment macro definitions. The **comment function** has the following syntax:

```
%'text'
%'text end-of-line
```

The comment function always evaluates to the null string. Two terminating characters are recognized, the apostrophe and the end-of-line character. The second form allows you to spread macro definitions over several lines while avoiding unwanted end-of-lines in the return value. In either form, the text or comment is not evaluated for macro calls.

Example

```
%'this is macro comment.'; this is an assembler comment.
%'the complete line including end-of-line is a comment
```

Source text before MPL processing

```
MOV R5, R15 %'the following line will be kept separate'
MOV R1, %'this comment eats the newline character
R12
```

Output text from MPL processor

```
MOV R5, R15
MOV R1, R12
```

Escape Function

Sometimes it is required to prevent the MPL processor from processing macro text. Two MPL processor functions perform this operation:

- escape function
- bracket function

The escape function interrupts scanning of macro text. The syntax of the **escape function** is:

%n text-n-characters-long

The metacharacter followed by a single decimal digit specifies the number of characters (maximum is 9) that are not evaluated. The escape function is useful for inserting a metacharacter (normally the % character), a comma, or a parenthesis.

Example

```
10%1% OF 10 = 1; expands to: 10% OF 10 = 1;
ASM%0251 expands to: ASM251
```

Bracket Function

The other MPL processor function that inhibits the processing of macro text is the bracket function. The syntax of the **bracket function** is:

% (balanced-text)

The bracket function disables all MPL processing of the text contained within the parentheses. However, the escape function, the comment function, and parameter substitution are still recognized.

Since there is no restriction for the length of the text within the bracket function, it is usually easier to use than the escape function.

Example

```
ASM%(251) evaluates to: ASM251
%(1,2,3,4,5) evaluates to: 1,2,3,4,5
```

Macro definition of 'DW'

```
%*DEFINE (DW (LIST, LABEL)) (
%LABEL: DW %LIST
)
```

Macro call to 'DW'

```
%DW (%(120, 121, 122, 123, -1), TABLE)
```

Return value of the macro call to 'DW'

```
TABLE: DW 120, 121, 122, 123, -1
```

The macro above will add word definitions to the source file. It uses two parameters: one for the word expression list and one for the label name. Without the bracket function it would not be possible to pass more than one expression in the list, since the first comma would be interpreted as the delimiter separating the actual parameters to the macro. The bracket function used in the macro call prevents the expression list (120, 121, 122, 123, -1) from being evaluated as separate parameters.

METACHAR Function

The MPL processor function METACHAR allows the programmer to change the character that will be recognized by the MPL processor as the metacharacter. The use of this function requires extreme care.

The syntax of the **METACHAR** function is:

```
%METACHAR (balanced_text)
```

The first character of the balanced text is taken to be the new value of the metacharacter. The characters @, (,), *, blank, tab, and identifier-characters are not allowed to be the metacharacter.

Example

```
%METACHAR (!) ; change metacharacter to '!'
!(1,2,3,4) ; bracket function invoked with !
```

Numbers and Expressions

Balanced text strings appearing in certain places in built-in MPL processor functions are interpreted as numeric expressions:

- The argument to evaluate function 'EVAL'
- The argument to the flow of control functions 'IF', 'WHILE', 'REPEAT' and 'SUBSTR'.

Expressions are processed as follows:

- The text of the numeric expression will be expanded in the ordinary manner of evaluating an argument to a macro function.
- The resulting string is evaluated to both a numeric and character representation of the expressions result. The return value is the character representation.

The following operators are allowed (shown in order of precedence).

- 1. Parenthesized Expressions
- 2. HIGH, LOW
- 3. *, /, MOD, SHL, SHR
- 4. EQ, LT, LE, GT, GE, NE
- 5. NOT
- 6. AND, OR, XOR

The arithmetic is done using signed 16-bit integers. The result of the relational operators is either 0 (FALSE) or 1 (TRUE).

Numbers

Numbers can be specified in hexadecimal (base 16), decimal (base 10), octal (base 8) and binary (base 2). A number without an explicit base is interpreted as decimal, this being the default representation. The first character of a number must always be a digit between 0 and 9. Hexadecimal numbers which do not have a digit as the first character must have a 0 placed in front of them.

Base	Suffix	Valid Characters	Examples	
hexadecimal	H,h	0 - 9, A-F (a - f)	1234H 99H 123H 0A0F0H 0FFH	
		Hexadecimal numbers must be preceded with a 0, if the first digit is in range A to F		
decimal	D,d	0 - 9	1234 65590D 20d 123	
octal	O,o,Q,q	0 - 7	1770 77770 250 1230 1777770	
binary	B,b	0 - 1	1111B 10011111B 101010101B	

Dollar (\$) signs can be placed within the numbers to make them more readable. However a \$ sign is not allowed to be the first or last character of a number and will not be interpreted.

1111\$0000\$1010\$0011B	is equivalent to	1111000010100011B
1\$2\$3\$4	is equivalent to	1234

Hexadecimal numbers may be also entered using the convention from the C language:

0xFE02	0x1234
0X5566	0x0A

Character Strings

The MPL processor allows ASCII characters strings in expressions. An expression is permitted to have a string consisting of one or two characters enclosed in single quote characters (').

'A' evaluates to 0041H
'AB' evaluates to 4142H
'a' evaluates to 0061H
'ab' evaluates to 6162H

" the null string is not valid!

'abc' ERROR due to more than two characters

The MPL processor cannot access the assembler's symbol table. The values of labels, SET and EQU symbols are not known during MPL processing. But, the programmer can define macro-time symbols with the MPL processor function 'SET'.

SET Function

The MPL processor function SET permits you to define macro-time symbols. SET takes two arguments: a valid identifier, and a numeric expression.

The syntax of the **SET** function is:

```
%SET (identifier, expression)
```

SET assigns the value of the numeric expression to the identifier.

The SET function affects the MPL processor symbol table only. Symbols defined by SET can be redefined with a second SET function call, or defined as a macro with DEFINE.

Source text

```
%SET (CNT, 3)
%SET (OFS, 16)
MOV R1,#%CNT+%OFS
%SET (OFS, %OFS + 10)
OFS = %OFS
```

Output text

```
MOV R1,#3+16
OFS = 26
```

The SET symbol may be used in the expression that defines its own value:

Source text

```
%SET (CNT, 10) %' define variable CNT' %SET (OFS, 20) %' define variable OFS'
```

% 'change values for CNT and OFS'

Output text

```
MOV R2,#30 + 40
MOV R5,#30
```

EVAL Function

The MPL processor function EVAL accepts an expression as an argument and returns the decimal character representation of it's result.

The syntax of the **EVAL** function is:

```
%EVAL (expression)
```

The expression arguments must be a legal expression with already defined macro identifiers, if any.

Source text

```
%SET (CNT, 10) %' define variable CNT'
%SET (OFS, 20) %' define variable OFS'

MOV R15,#%EVAL (%CNT+1)

MOV WR14,#%EVAL (14+15*200)

MOV R13,#%EVAL (-(%CNT + %OFS - 1))

MOV R2,#%EVAL (%OFS LE %CNT)

MOV R7,#%EVAL (%OFS GE %CNT)
```

Output text

```
MOV R15,#11

MOV WR14,#3014

MOV R13,#-29

MOV R2,#0

MOV R7,#1
```

Logical Expressions and String Comparison

The following MPL processor functions compare two balanced-text string arguments and return a logical value based on that comparison. If the function evaluates to TRUE, then it returns a value of 1. If the function evaluates to FALSE, then it returns a value of 0. The list of string comparison functions below shows the syntax and describes the type of comparison made for each. Both arguments to these function may contain macro calls. (These MPL calls are expanded before the comparison is made).

%EQS (arg1,arg2)	True if both arguments are identical
%NES (arg1,arg2)	True if arguments are different in any way
%LTS (arg1,arg2)	True if first argument has a lower value than second argument
%LES (arg1,arg2)	True if first argument has a lower value then second argument or if both arguments are identical
%GTS (arg1,arg2)	True if first argument has a higher value than second argument
%GES (arg1,arg2)	True if first argument has a higher value than second argument or if both arguments are identical

Example

ie
r ASCII
nerical
jument
r ne

Conditional MPL Processing

Some MPL functions accept logical expressions as arguments. The MPL uses the value 1 and 0 to determine TRUE or FALSE. If the value is one, then the expression is TRUE. If the value is zero, then the expression is FALSE.

Typically, you will use either the relational operators (EQ, NE, LE, LT, GT, or GE) or the string comparison functions (EQS, NES, LES, LTS, GTS, or GES) to specify a logical value.

IF Function

The IF MPL function evaluates a logical expression, and based on that expression, expands or skips its text arguments. The syntax of the MPL processor function **IF** is:

```
%IF (expression) THEN (balanced-text1) [ ELSE (balanced-text2) ] FI
```

IF first evaluates the expression, if it is TRUE, then balanced-text1 is expanded; if it is FALSE and the optional ELSE clause is included, then balanced-text2 is expanded. If it is FALSE and the ELSE clause is not included, the IF call returns a null string. FI must be included to terminate the call.

IF calls can be nested; when they are, the ELSE clause refers to the most recent IF call that is still open (not terminated by FI). FI terminates the most recent IF call that is still open.

Source text

Output text

ADD I	R15,R3
ב עעא	KID, KO
SUB I	R15,R9

WHILE Function

Often you may wish to perform macro operations until a certain condition is met. The MPL processor function WHILE provides this facility.

The syntax for the MPL processor function WHILE is:

```
%WHILE (expression) (balanced-text)
```

WHILE first evaluates the expression. If it is TRUE, then the balanced-text is expanded; otherwise, it is not. Once the balanced-text has been expanded, the logical argument is retested and if it is still TRUE, then the balanced-text is again expanded. This loop continues until the logical argument proves FALSE.

Since the MPL continues processing until expression evaluates to FALSE, the balanced-text should modify the expression, or the WHILE may never terminate.

A call to the MPL processor function EXIT will always terminate a WHILE function. EXIT is described later.

Source text

Output text

```
ADD R15,R15
ADD R15,R15
ADD R15,R15
ADD R15,R15
ADD R15,R15
ADD R15,R15
```

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REPEAT Function

The MPL processor function REPEAT expands its balanced-text a specified number of times. The syntax for the MPL processor function **REPEAT** is:

```
%REPEAT (expression) (balanced-test)
```

REPEAT uses the expression for a numerical value that specifies the number of times the balanced-text will be expanded. The expression is evaluated once when the macro is first called, then the specified number of iterations is performed.

Source text

```
%REPEAT (5)
( -enter any key to shut down-
)
%REPEAT (5) (+%REPEAT (9) (-))+
```

```
-enter any key to shut down-
```

EXIT Function

The EXIT MPL processor function terminates expansion of the most recently called REPEAT, WHILE or user-defined macro function. It is most commonly used to avoid infinite loops (example: a WHILE that never becomes FALSE, or a recursive user-defined macro that never terminates). It allows several exit points in the same macro.

The syntax for the MPL processor function **EXIT** is:

%EXIT

Source text

```
%SET (count, 0)

%WHILE (1)
(%IF (%count GT 5) THEN (%EXIT)

FI DW %count, -%count
%SET (count, %count + 1))
```

```
DW 0, -0
DW 1, -1
DW 2, -2
DW 3, -3
DW 4, -4
DW 5, -5
```

String Manipulation Functions

The purpose of the MPL is to manipulate character strings. Therefore, there are several MPL functions that perform common character string manipulations.

LEN Function

The MPL processor function LEN returns the length of the character string argument in hexadecimal: The character string is limited to 256 characters.

The syntax for the MPL processor function **LEN** is:

```
%LEN (balanced-text)
```

Source text

```
%LEN (A251) %' len = 4'
%LEN (A251,A251) %' comma counts also'
%LEN ()
%LEN (ABCDEFGHIJKLMNOPQRSTUVWXYZ)
%DEFINE (TEXT) (QUEEN)
%DEFINE (LENGTH) (%LEN (%TEXT))
LENGTH OF '%TEXT' = %LENGTH.
```

```
4
9
0
26
LENGTH OF 'QUEEN' = 5.
```

SUBSTR Function

The MPL processor function SUBSTR returns a substring of the given text argument. The function takes three arguments: a character string to be divided and two numeric arguments.

The syntax for the MPL processor function **SUBSTR** is:

```
%SUBSTR (balanced-text, expression1, expression2)
```

Where balanced-text is any text argument, possibly containing macro calls. Expression1 specifies the starting character of the substring. Expression2 specifies the number of characters to be included in the substring.

If expression1 is zero or greater than the length of the argument string, then SUBSTR returns the null string. The index of the first character of the balanced text is one.

If expression2 is zero, then SUBSTR returns the null string. If expression2 is greater than the remaining length or the string, then all characters from the start character to the end of the string are included.

Source text

```
%DEFINE (STRING) (abcdefgh)
%SUBSTR (%string, 1, 2)
%SUBSTR (%(1,2,3,4,5), 3, 20)
```

Output text

```
ab
2,3,4,5
```

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MATCH Function

The MPL processor function MATCH searches a character string for a delimiter character, and assigns the substrings on either side of the delimiter to the identifiers

The syntax for the MPL processor function **MATCH** is:

```
%MATCH (identifier1 delimiter identifier2) (balanced-text)
```

Identifier1 and identifier2 must be valid macro identifiers. Delimiter is the first character to follow identifier1. Typically, a space or comma is used, but any character that is not a macro identifier character may be used. Balanced-text is the text searched by the MATCH function. It may contain macro calls.

MATCH searches the balanced-text string for the specified delimiter. When the delimiter is found, then all characters to the left are assigned to identifier1 and all characters to the right are assigned to identifier2. If the delimiter is not found, the entire balanced-text string is assigned to identifier1 and the null string is assigned to identifier2.

Source text

```
%DEFINE (text) (-1,-2,-3,-4,-5)
%MATCH (next,list) (%text)
%WHILE (%LEN (%next) NE 0)
( MOV R8,#%next
 MOV @WR2,R8 %MATCH (next,list) (%list)
 INC WR2,#1
)
```

```
MOV
       R8,#-1
MOV
       @WR2,R8
INC
       WR2,#1
       R8.#-2
MOV
VOM
       @WR2,R8
       WR2,#1
INC
MOV
       R8,#-3
MOV
       @WR2,R8
       WR2,#1
INC
MOV
       R8,#-4
MOV
       @WR2,R8
INC
       WR2,#1
       R8,#-5
MOV
       @WR2,R8
MOV
INC
       WR2,#1
```

Console I/O Functions

There are two MPL processor functions that perform console I/O: IN and OUT. Their names describe the function each performs. IN outputs a character '>' as a prompt, and returns the line typed at the console. OUT outputs a string to the console; a call to OUT is replaced by the null string.

The syntax for the MPL processor functions IN and OUT is:

```
%IN
%OUT (balanced-text)
```

Source text

```
%OUT (enter baud rate)
%set (BAUD_RATE, %in)
BAUD_RATE = %BAUD_RATE
```

```
<19200 was entered at the console>
BAUD_RATE = 19200
```

Advanced Macro Processing

The MPL definition function associates an identifier with a functional string. The macro may or may not have an associated pattern consisting of parameters and/or delimiters. Optionally present are local symbols.

The syntax for a macro definition is:

```
%DEFINE (macro_id define_pattern) [LOCAL id_list] (balanced_text)
```

The define_pattern is a balanced string which is further analyzed by the MPL processor as follows:

```
define_pattern = { [parm_id] [delimiter_specifier] }
```

Delimiter specifier is one of the following:

- A string that contains no non-literal id-continuation, logical blank, or at character ('@').
- @delimiter_id

The macro call must have a call pattern which corresponds to the macro define pattern. Regardless of the type of delimiter used to define a macro, once it has been defined, only delimiters used in the definition can be used in the macro call. Macros defined with parentheses and commas require parentheses and commas in the macro call. Macros defined with spaces or any other delimiter require that delimiter when called.

The define pattern may have three kinds of delimiters: implied blank delimiters, identifier delimiters and literal delimiters.

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Literal Delimiters

The delimiters used in user-defined macros (parentheses and commas) are literal delimiters. A literal delimiter can be any character except the metacharacter.

When you define a macro using a literal delimiter, you must use exactly that delimiter when you call the macro. If the specified delimiter is not used as it appears in the definition, a macro error occurs.

When defining a macro, the delimiter string must be literalized, if the delimiter meets any of the following conditions:

- more than one character,
- a macro identifier character (A-Z, 0-9, , or ?),
- a commercial at (@), a space, tab, carriage return, or linefeed.

Use the escape function (%n) or the bracket function (%()) to literalize the delimiter string.

This is the simple form shown earlier:

Before Macro Expansion	After Macro Expansion
%*DEFINE(MAC(A,B))(%A %B)	null string
%MAC(4,5)	4 5

In the following example brackets are used instead of parentheses. The commercial at symbol separates parameters:

```
%*DEFINE (MOV[A%(@)B]) (MOV %A,%B) → null string
%MOV[P0@P1] → MOV P0,P1
```

In the next two examples, delimiters that could be id delimiters have been defined as literal delimiter (the differences are noted):

%*DEFINE(ADD (R10 AND B)) (ADD R10,%B)	→ null string
%ADD (R10 AND #27H)	→ ADD R10,#27H

Spaces around AND are considered as part of the argument string.

Blank Delimiters

Blank delimiters are the easiest to use. Blank delimiter is one or more spaces, tabs or new lines (a carriage-return/linefeed pair) in any order. To define a macro that uses the blank delimiter, simply place one or more spaces, tabs, or new lines surrounding the parameter list.

When the macro defined with the blank delimiter is called, each delimiter will match a series of spaces, tabs, or new lines. Each parameter in the call begins with the first non-blank character, and ends when a blank character is found.

Source text

```
%*DEFINE (X1 X2 X3) (P2=%X2, P3=%X3)
%X1 assembler A251
```

Output text

P2=assembler, P3=A251

Identifier Delimiters

Identifier delimiters are legal macro identifiers designated as delimiters. To define a macro that uses an identifier delimiter, you must prefix the delimiter with the @ symbol. You must separate the identifier delimiter from the macro identifiers (formal parameters or macro name) by a blank character.

When calling a macro defined with identifier delimiters, a blank delimiter is required to precede the identifier delimiter, but none is required to follow the identifier delimiter.

Source text

```
%*DEFINE (ADD X1 @TO X2 @STORE X3)(
   MOV    R1,%X1
   MOV    R2,%X2
   ADD   R1,R2
   MOV   %X3,R1
)
%ADD VAR1 TO VAR2 STORE VAR3
```

```
MOV R1, VAR1
MOV R2, VAR2
ADD R1, R2
MOV VAR3, R1
```

Literal and Normal Mode

In normal mode, the MPL processor scans for the metacharacter. If it is found, parameters are substituted and macros are expanded. This is the usual operation of the MPL processor.

When the literal character (*) is placed in a DEFINE function, the MPL processor shifts to literal mode while expanding the macro. The effect is similar to surrounding the entire call with the bracket function. Parameters to the literalized call are expanded, the escape, comment, and bracket functions are also expanded, but no further processing is performed. If there are any calls to other macros, they are not expanded.

If there are no parameters in the macro being defined, the DEFINE function can be called without literal character. If the macro uses parameters, the MPL processor will attempt to evaluate the formal parameters in the macro-body as parameterless macro calls.

The following example illustrates the difference between defining a macro in literal mode and normal mode:

```
%SET (TOM, 1)
%*DEFINE (AB) (%EVAL (%TOM))
%DEFINE (CD) (%EVAL (%TOM))
```

When AB and CD are defined, TOM is equal to 1. The macro body of AB has not been evaluated due to the literal character, but the macro body of CD has been completely evaluated, since the literal character is not used in the definition. Changing the value of TOM has no effect on CD, but it changes the value of AB:

```
%SET (TOM,2) \rightarrow null string 
%AB \rightarrow 2 
%CD \rightarrow 1 
%*CD \rightarrow 1 
%*AB \rightarrow %EVAL (%TOM)
```

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MACRO Errors

The MPL processor emits error messages if errors occur in the MPL processing phase. Macro errors are displayed like other assembly errors in the listing file. The following table lists the error messages generated by the MPL processor.

Number	Error Message and Description
	· · ·
200	PREMATURE END OF FILE
	The end of the source module was reached while processing some macro call, which requires more input from the source file.
201	' <token>' IDENTIFIER EXPECTED</token>
	The MPL processor expected an identifier while processing some macro. None was found. The unexpected token is displayed with this error message.
202	MPL FUNCTION ' <name>': '<character>' EXPECTED</character></name>
	The context of the MPL processor language requires a specific character from the input given by <character> while processing the built-in function given by <name>.</name></character>
203	<pre><string>: UNBALANCED PARENTHESES</string></pre>
	A balanced string requires the same number of right parentheses and left parentheses.
204	EXPECTED ' <token>'</token>
	The syntax requires a specific token to follow, for example THEN after the balanced text argument to IF.
205	INCOMPLETE MACRO DEFINITION
	The macro definition has not been completely processed due to premature end of input file.
206	FUNCTION 'MATCH': ILLEGAL CALL PATTERN
	The built-in function MATCH was called with an illegal call pattern. The call pattern must consist of some formal name followed by a delimiter specification and another formal name.
207	FUNCTION 'EXIT' IN BAD CONTEXT
	The built-in function EXIT is allowed only in the loop control constructs WHILE and REPEAT.
208	ILLEGAL METACHARACTER ' <character>'</character>
	The first character of the balanced text argument to METACHAR is taken to be the new value of the metacharacter. The characters @, (,), *, blank, tab, and identifier-characters are not allowed to be the metacharacter.
209	CALL PATTERN - DELIMITER ' <delimiter>' NOT FOUND</delimiter>
	The call pattern of some macro does not conform to the define pattern of that macro. The delimiters of the macro call should be checked for conformance.
210	CALL TO UNDEFINED MACRO ' <name>'</name>
	The macro call specifies the name of an undefined macro.

Number

211

212

213

214

Error Message and Description

INVALID MPL COMMAND '%<character>'

A number of an expression contains an invalid digit.

INVALID DIGIT '<character>' IN NUMBER

UNCLOSED STRING OR CHARACTER CONSTANT

INVALID STRING OR CHARACTER CONSTANT

The character following the metacharacter does not form a valid MPL command.

The string representing a number in an expression is invalid. The string must be either one or two characters long. A character constant must not be longer than one character.

Strings or character constants must be enclosed by single or double quotes.

Chapter 7. Invocation and Controls

This chapter explains how to use **Ax51** to assemble **x51** assembly source files and discusses the assembler controls that may be specified on the command line and within the source file.

Using the controls described in this chapter, you can specify which operations are performed by Ax51. For example, you can direct Ax51 to generate a listing file, produce cross reference information, and control the amount of information included in the object file. You can also conditionally assemble sections of code using the conditional assembly controls.

Environment Settings

To run the **Ax51** macro assembler and the utilities from a Windows command prompt, you must create new entries in the environment table. In addition, you must specify a **PATH** for the compiler folder. The following table lists the environment variables, their default paths, and a brief description.

Variable	Path	Description
PATH	KEIL\C51\BIN or KEIL\C251\BIN	This environment variable specifies the path of the Ax51 executable programs.
ТМР		This environment variable specifies which path to use for temporary files generated by the assembler. If the specified path does not exist, the assembler generates an error and aborts compilation.
C51INC	KEIL\C51\INC	This environment variable specifies the location of the standard C51 or CX51 include files.
C251INC	KEIL\C251\INC	This environment variable specifies the location of the standard C251 include files.

Typically, these environment settings are automatically placed in your **AUTOEXEC.BAT** file. However, to put these settings in a separate batch file, use the following example as guideline:

PATH = C:\KEIL\C51\BIN
SET TMP = D:\
SET C51INC = C:\KEIL\C51\INC

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Running Ax51

The Ax51 assembler is invoked by typing the program name at the Windows command prompt. On this command line, you must include the name of the assembler source file to be translated, as well as any other necessary assembler controls required to translate your source file. The format for the Ax51 command line is:

```
A51 sourcefile directives...

AX51 sourcefile directives...

A251 sourcefile directives...
```

or:

```
A51 @commandfile
AX51 @commandfile
A251 @commandfile
```

where

is the name of the source program you want to assemble.

controls are used to direct the operation of the assembler. Refer to

"Assembler Controls" on page 197 for more information.

commandfile is the name of a command input file that may contain

sourcefile and directives. A commandfile is used, when the **Ax51** invocation line gets complex and exceeds the limits of

the Windows command prompt.

The following command line example invokes **A251** macro assembler and specifies the source file **SAMPLE.A51** and uses the controls **DEBUG**, **XREF**, and **PAGEWIDTH**.

```
A251 SAMPLE.A51 DEBUG XREF PAGEWIDTH(132)
```

A251 displays the following information upon successful invocation and assembly.

```
A251 MACRO ASSEMBLER V3.00
ASSEMBLY COMPLETE. 0 ERROR(S) 0 WARNING(S)
```

ERRORLEVEL

After assembly, the number of errors and warnings detected is output to the screen. **Ax51** then sets the **ERRORLEVEL** to indicate the status of the assembly. The ERRORLEVEL values are identical for all the **Ax51** assembler, **Lx51** linker/locater and other **x51** utilities. The values are listed in the following table:

ERROR LEVEL	Meaning
0	No ERRORS or WARNINGS
1	WARNINGS only
2	ERRORS and possibly also WARNINGS
3	FATAL ERRORS

You can access the **ERRORLEVEL** variable in batch files for conditional tests to terminate the batch processing when an error occurs. Refer to the *Windows online help* for more information about **ERRORLEVEL** or batch files.

Output Files

Ax51 generates a number of output files during assembly. By default, these files use the same *basename* as the source file, but with a different file extension. The following table lists the files and gives a brief description of each.

File Extension	Description
basename.LST	Files with this extension are listing files that contain the formatted source text along with any errors detected by the assembler. Listing files may optionally contain symbols used and the generated assembly code. Refer to "PRINT / NOPRINT" on page 221 for more information.
basename.OBJ	Files with this extension are object modules that contain relocatable object code. Object modules can be linked into an absolute object module by the Lx51 Linker/Locator. Refer to "OBJECT / NOOBJECT" on page 219 for more information.

Assembler Controls

Ax51 provides a number of controls that you can use to direct the operation of the assembler. Controls can be specified after the filename on the invocation

line or in a control line within the source file. Control lines are prefixed by the dollar sign character ('\$').

Example

A51 TESTFILE.A51 MPL DEBUG XREF

or

\$MPL \$DEBUG \$XREF

or

SMPL DEBUG XREF

In the above example, MPL, DEBUG, and XREF are all control commands and TESTFILE.A51 is the source file to assemble.

Ax51 has two classes of controls; primary and general. Primary controls are specified in the invocation line on the first few lines of the assembly source file. Primary controls remain in effect throughout the assembly. For this reason, primary controls may be used only in the invocation line or in control lines at the beginning of the program. Only other control lines that do not contain the **INCLUDE** control may precede a line containing a primary control. The **INCLUDE** control marks the end of any primary control specifications.

If a primary control is specified in the invocation line and on the first few lines of the assembly source file, the specification on the invocation line is used. This enables you override primary controls via the invocation line.

The general controls are used to control the immediate action of the assembler. Typically their status is set and modified during the assembly. Control lines containing only general controls may be placed anywhere in your source code.

The table on the next page lists all of the controls, their abbreviations, their default values, and a brief description of each.

NOTE

Some controls like **EJECT** and **SAVE** cannot be specified on the command line. The syntax for each control is the same when specified on the command line or when specified within the source file. **Ax51** will generate a fatal error for controls that are improperly specified.

Directive F	Page	Description
CASE ‡	200	AX51, A251 ONLY: enable case sensitive mode for symbol names.
COND / NOCOND	201	Enable or disable skipped sections to appear in the listing file.
DATE(date) ‡	202	Places a date string in header (9 characters maximum).
DEBUG ‡	203	Outputs debug symbol information to object file.
<u>EJ</u> ECT	204	Continue listing on next page.
ERRORPRINT[(file)] ‡	205	Designates a file to receive error messages in addition to the listing.
FIXDRK ‡	206	A251 ONLY : Replaces INC DRk with ADD DRk for C-Step devices.
GEN ‡	207	Generates a full listing of macro expansions in the listing file.
NOGEN ‡	207	List only the original source text in listing file.
INCDIR(path)	208	Define paths to be searched when a file is included via INCLUDE.
INCLUDE(file)	209	Designates a file to be included as part of the program.
INTR2	210	A251 ONLY: Select 2-Byte interrupt frame size on 251 devices.
<u>LI</u> ST, <u>NOLI</u> ST	211	Print or do not print the assembler source in the listing file.
<u>M</u> OD <u>51</u> ‡	212	AX51 ONLY: Select classic 8051 instruction set (default).
MOD_MX51 ‡	212	AX51 ONLY: Select Philips 80C51MX instruction set).
MOD_CONT ‡	212	AX51 ONLY: Select Dallas 390 contiguous mode instruction set.
MODSRC ‡	213	A251 ONLY: Select Intel/Atmel WM 251 source mode.
MPL ‡	214	Enable Macro Processing Language.
NOLINES	215	Do not generate LINE number information.
NOMACRO ‡	216	Disable Standard Macros
NOMOD51	217	Do not recognize the 8051-specific predefined special register.
NOOBJECT	219	Designates that no object file will be created.
NOREGISTERBANK	222	Indicates that no banks are used.
NOSYMBOLS	218	No symbol table is listed.
NOSYMLIST	225	Do not list the following symbol definitions in the symbol table.
OBJECT[(file)]	219	Designate file to receive object code.
PAGELENGTH(n) ‡	220	Sets maximum number of lines in each page of listing file.
PAGEWIDTH(n) ‡	220	Sets maximum number of characters in each line of listing file.
PRINT[(file)] ‡	221	Designates file to receive source listing.
NOPRINT ‡	221	Designates that no listing file will be created.
REGISTERBANK	222	Indicates one or more banks used in program module.
REGUSE	223	Defines register usage of assembler functions for the C optimizer.
<u>RES</u> TORE	224	Restores control setting from SAVE stack.
SAVE	224	Stores current control setting for GEN, LIST and SYMLIST.
<u>S</u> YM <u>L</u> IST	225	List the following symbol definitions in the symbol table.
<u>TIT</u> LE(string) ‡	226	Places a string in all subsequent page headers.
XREF ‡	227	Creates a cross reference listing of all symbols used in program.

[‡] marks **general controls** that may be specified only once on the command line or at the beginning of a source file in a \$control line. They may not be used more than once in a source file.

CASE (AX51 and A251 only)

Abbreviation: CA

Arguments: None.

Default: No case sensitivity. All characters are converted to

uppercase.

Control Class: Primary

\muVision2 Control: Options – Ax51 – Case sensitive symbols.

Description: The **CASE** control directs the assembler to operate in case

sensitive mode. Without CASE, the assembler operates in case insensitive mode and maps lowercase input characters

to uppercase.

CASE becomes meaningful when modules generated by the assembler are combined with modules generated by the C compiler. Identifiers exported from C modules always

appear in uppercase and lowercase (as written).

Corresponding names used in an assembler module must

match the case of the names from the C module.

Example: \$CASE

AX51 SAMPLE.A51 CASE

COND / NOCOND

Abbreviation: None.

Arguments: None.

Default: COND

Control Class: General

μVision2 Control: Options – Listing – Assembler Listing – Conditional.

Description: The **COND** control directs the **Ax51** assembler to include

unassembled parts of conditional **IF–ELSEIF–ENDIF** blocks in the listing file. Unassembled code is listed without

line numbers.

The **NOCOND** control prevents unassembled portions of **IF-ELSE-ENDIF** blocks from appearing in the listing file.

Examples: AX51 SAMPLE.A51 COND

\$COND

AX51 SAMPLE.A51 NOCOND

\$NOCOND

DATE

Abbreviation: DA

Arguments: A string enclosed within parentheses.

Default: The date obtained from the operating system.

Control Class: Primary

\muVision2 Control: Options – Ax51 – Misc controls: enter the control.

Description: The Ax51 assembler includes the current date in the header

of each page in the listing file. The **DATE** control allows you to specify the date string that is included in the header. The string must immediately follow the **DATE** control and must be enclosed within parentheses. Only the first 8 characters of the date string are used. Additional characters

are ignored.

Example: AX51 SAMPLE.A51 DATE(19JAN00)

\$DATE(10/28/00)

DEBUG

Abbreviation: DB

Arguments: None.

Default: No debugging information is generated.

Control Class: Primary

μVision2 Control: Options – Output – Debug Information

Description: The **DEBUG** control instructs the **Ax51** assembler to

include debugging information in the object file. This information is used when testing the program with an

emulator or simulator.

The **DEBUG** control also includes line number information for source level debugging. Line number information can be

disabled with the **NOLINES** control.

Examples: A51 SAMPLE.A51 DEBUG

\$DEBUG

EJECT

Abbreviation: EJ

Arguments: None

Default: None

Control Class: General

μVision2 Control: This control cannot be specified on the command line.

Description: The **EJECT** control inserts a form feed into the listing file

after the line containing the **EJECT** statement. This control

is ignored if **NOLIST** or **NOPRINT** was previously

specified.

Example: \$EJECT

ERRORPRINT

Abbreviation: EP

Arguments: An optional filename enclosed within parentheses

Default: No error messages are output to the console.

Control Class: Primary

μVision2 Control: This control is used by μVision2 to get the error output. It

should be not specified when you are using the µVision2

IDE.

Description: The **ERRORPRINT** control directs the **Ax51** assembler to

output all error messages either to the console (if no

filename is specified) or to a specified file.

Examples: AX51 SAMPLE.A51 ERRORPRINT (SAMPLE.ERR)

AX51 SAMPLE2.A51 ERRORPRINT

\$EP

FIXDRK (A251 only)

Abbreviation: FD

Arguments: None.

Default: Use the **INC DRk,#const** instruction.

Control Class: Primary

μVision2 Control: Options – A251 – Misc controls: enter the control.

Description: The **FIXDRK** control instructs the assembler to replace the

INC DRk,#const instruction with the ADD DRk,#const

instruction.

You may require this control because the **INC DRk**,#const instruction does not work in the Intel 251SB C-step CPU. Check the stepping level or contact your silicon vendor to find out if you need to use this control. If you are using the Intel 8xC251SB device and if you are in doubt about the stepping code, you should apply this control.

Examples: A251 SAMPLE.A51 FIXDRK

\$FIXDRK

GEN / NOGEN

Abbreviation: GE / NOGE

Arguments: None

Default: NOGEN

Control Class: General

μVision2 Control: Options – Listing – Assembler Listing – Macros.

Description: The **GEN** control directs the **Ax51** assembler to expand or

list, in a listing file, all assembly instructions contained in a

macro.

The **NOGEN** control prevents the **Ax51** assembler from including macro expansion text in the listing file. Only the

macro name is listed.

Examples: A51 SAMPLE.A51 GEN

\$GEN

A51 SAMPLE.A51 NOGEN

\$NOGEN

INCDIR

Abbreviation: ID

Arguments: Path specifications for include files enclosed in parentheses.

Default: None.

Control Class: General

 μ Vision2 Control: Options – Ax51 – Include Paths.

Description: The **INCDIR** control specifies the location of files specified

with the **INCLUDE** control. Multiple path declarations must be separated by semicolon characters (';'). A

maximum of 5 paths may be specified.

When searching for include files, the assembler searches first the current folder, which is typically the folder of the project file. Then, paths specified by **INCDIR** are

searched.

Example: AX51 SAMPLE.A51 INCDIR(C:\AX51\MYINC;C:\CHIP_DIR)

INCLUDE

Abbreviation: IC

Arguments: A filename enclosed within parentheses.

Default: None.

Control Class: General

μVision2 Control: This control cannot be specified on the command line.

Description: The **INCLUDE** control directs the Ax51 assembler to

include the contents of the specified file in the assembly of the program. The include file's contents are inserted immediately following the **INCLUDE** control line. **INCLUDE** files may be nested up to 9 levels deep.

The **INCLUDE** control is usually used to include special function register definition files for *x*51 derivatives. It is also commonly used to include declarations for external routines, variables, and macros. Files containing assembly language code may also be included.

Example: \$INCLUDE (REG51F.INC)

The macro assembler searches the current folder and the folders specified with the **INCDIR** control for include files. If the specified file cannot be found in this folders, the

assembler tries to locate the file in the folder

path_of_the_EXE_file\..\ASM\. In a typical installation of the tool chain this is the correct path for the derivate specific include files. (The \C51\BIN\ or \C251\BIN\ folder contains the macro assembler and the \C51\ASM\ or \C251\ASM\ folder contains the register definition files).

INTR2 (A251 only)

Abbreviation: I2

Arguments: None.

Default: The A251 assembler assumes that an interrupt pushes 4

bytes onto the stack: a 24-bit return address and **PSW1**.

Control Class: General

 μ Vision2 Control: Options – Target – 4 Byte interrupt frame size.

Description: The **INTR2** control informs the A251 assembler and the

L251 linker/locator that the 251 CPU saves the low order 16 bits of the program counter but does not automatically save

PSW1 when entering an interrupt.

The **INTR2** control does not change any assembler code or instruction encoding. It only informs the linker and debugging tools of the interrupt frame size assumed for interrupt functions. The linker uses this information to check the consistency of the interrupt frame sizes between program modules. If the interrupt frame sizes of the object modules differ, the L251 linker/locator generates a warning

message.

Example: A251 SAMPLE.A51 INTR2

\$INTR2

LIST / NOLIST

Abbreviation: LI / NOLI

Arguments: None

Default: LIST

Control Class: General

 μ Vision2 Control: Options – Ax51 – Misc controls: enter the control.

Description: The LIST control directs the Ax51 assembler to include the

program source text in the generated listing file.

The **NOLIST** control prevents subsequent lines of your assembly program from appearing in the generated listing

file.

If a line that would normally not be listed causes an

assembler error, that line will be listed along with the error

message.

Examples: AX51 SAMPLE.A51 LI

\$LIST

AX51 SAMPLE.A51 NOLIST

\$NOLI

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MOD51, MOD CONT, MOD MX51 (AX51 only)

Abbreviation: M51, MC, MX

Arguments: None

Default: MOD51: generate code for classic 8051.

Control Class: Primary

μVision2 Control: Options – Target (mode selection depends on the device).

Description: The **MODxxx** controls selects the instruction set that is used

in the application code.

The **MOD51** control is the default setting of AX51 and instructs the assembler to generate code with that uses only the instructions of the classic 8051.

The MOD_MX51 control enables the instruction set extensions for the Philips 80C51MX architecture. If you are using a device with this architecture, at least one module must be translated with this directive. You can intermix code that has be written for the classic 8051 in a project for the Philips 80C51MX.

The MOD_CONT control enables the 24-bit contiguous address mode that is available on some Dallas devices. If you are using this mode, you need to translate all modules with this directive. It is not possible to use code that has been translated for the classic 8051 when you are using this CPU mode.

Examples: AX51 SAMPLE.A51 MOD_MX51 AX51 SAMPLE.A51 MOD_COND

\$MX ; generate code for Philips 80C51MX architecture \$MC ; generate code for Dallas 24-bit contiguous mode

MODSRC (A251 only)

Abbreviation: MS

Arguments: None

Default: Generate code for binary mode of the Intel/Atmel WM 251

CPU.

Control Class: Primary

μVision2 Control: Options – Target – CPU Mode.

Description: The **MODSRC** control instructs the A251 assembler to

generate code for the Intel/Atmel WM 251 architecture

using the SOURCE mode of this CPU.

Examples: A251 SAMPLE.A51 MODSRC

\$MODSRC

MPL

Abbreviation: None

Arguments: None

Default: The Macro Processing Language is disabled.

Control Class: Primary.

 μ Vision2 Control: Options – Ax51 – Macro processor – MPL.

Description: The MPL control enables the Macro Processing Language.

The **MPL** language is compatible to the Intel ASM51. Refer to "Chapter 6. Macro Processing Language" on page

163 for more information about the MPL processor.

Examples: A251 SAMPLE.A51 MPL

\$MPL

NOLINES

Abbreviation: NOLN

Arguments: None.

Default: Line numbers for source level debugging are generated

when the **DEBUG** control is used.

Control Class: Primary

 μ Vision2 Control: Options – Ax51 – Misc controls: enter the control.

Description: The **NOLINES** control disables the line number information

for source level debugging. This control is useful when the Ax51 assembler is used with very old debugging tools and

very old emulators.

Examples: A251 SAMPLE.A51 NOLINES

\$NOLINES

NOMACRO

Abbreviation: None.

Arguments: None.

Default: Standard Macros are fully expanded.

Control Class: Primary

μVision2 Control: Options – Ax51 – Macro processor – Standard.

Description: The **NOMACRO** control disables the standard macro

facility of the Ax51 assembler so that standard macros are

not expanded.

Examples: A251 SAMPLE.A51 NOMACRO

\$NOMACRO

NOMOD51

Abbreviation: NOMO

Arguments: None.

Default: The A51 assembler pre-defines all special function registers

of the 8051 CPU. The A251 assembler and the AX51 assembler do not pre-define any CPU special function

registers.

Control Class: Primary

 μ Vision2 Control: Options – Ax51 – Special Function Registers – Define 8051

SFR Names.

Description: The **NOMOD51** control prevents the A51 assembler from

implicitly defining symbols for the default 8051 special function registers. This is necessary when you want to include a definition file to declare symbols for the special

function registers of a different 8051 derivative.

The A251 assembler and the AX51 assembler support the **NOMOD51** control only for source compatibility to the A51. However, the 8051 special function registers are not

predefined in A251 or AX51.

Examples: A251 SAMPLE.A51 NOMO

\$NOMOD51

NOSYMBOLS

Abbreviation: SB / NOSB

Arguments: None

Default: The **Ax51** assembler generates a table of all symbols used in

and by the assembly program module. This symbol table is

included in the generated listing file.

Control Class: Primary

μVision2 Control: Options – Listing – Assembler Listing – Symbols

Description: The **NOSYMBOLS** control prevents the **Ax51** assembler

from generating a symbol table in the listing file.

Examples: A251 SAMPLE.A51 SYMBOLS

\$SB

A251 SAMPLE.A51 NOSB

\$NOSYMBOLS

OBJECT / NOOBJECT

Abbreviation: OJ / NOOJ

Arguments: An optional filename enclosed within parentheses.

Default: OBJECT (basename.OBJ)

Control Class: Primary

μVision2 Control: Options – Output – Select Folder for Objects

Description: The **OBJECT** control specifies that the **Ax51** assembler

generate an object file. The default name for the object file is *basename*.OBJ, however, an alternate filename may be specified in parentheses immediately following the

OBJECT control statement.

The **NOOBJECT** control prevents the **Ax51** assembler from

generating an object file.

Examples: A51 SAMPLE.A51 OBJECT (OBJDIR\SAMPLE.OBJ)

OJ(OBJ\SAMPLE.OBJ)

A251 SAMPLE.A51 NOOJ

\$NOOBJECT

PAGELENGTH, PAGEWIDTH

Abbreviation: PL, PW

Arguments: PAGELENGTH accepts a number between 10 and 65535;

PAGEWIDTH accepts a number between 78 and 132

enclosed within parentheses.

Default: PAGELENGTH (60)

PAGEWIDTH (120)

μVision2 Control: Options – Listing – Page Length / Page Width

Description: The **PAGELENGTH** control specifies the number of lines

printed per page in the listing file. The number must be a decimal value between 10 and 65535. The default is 60.

The **PAGEWIDTH** control specifies the maximum number of characters in a line in the listing file. Lines that are longer than the specified width are automatically wrapped around to the next line. The default number of characters

per line is 120.

Example: A251 SAMPLE.A51 PAGELENGTH(132) PAGEWIDTH (79)

\$PL (66) \$PW(132)

PRINT / NOPRINT

Abbreviation: PR / NOPR

Arguments: An optional filename enclosed within parentheses.

Default: PRINT(basename.LST)

Control Class: Primary

μVision2 Control: Options – Listing – Select Folder for List Files

Description: The **PRINT** control directs the **Ax51** assembler to generate

a listing file. The default name for the listing file is **basename.LST**, however, an alternate filename may be specified in parentheses immediately following the **PRINT**

control statement.

The **NOPRINT** control prevents the **Ax51** assembler from

generating a listing file.

Examples: A251 SAMPLE.A51 PRINT

A51TESTPRG.A51 PR(TESTPRG1.LST)

\$PRINT(LPT1)

AX51 SAMPLE.A51 NOPRINT

\$NOPR

REGISTERBANK / NOREGISTERBANK

Abbreviation: RB / NORB

Arguments: Register bank numbers separated by commas and enclosed

within parentheses. For example, **RB** (1,2,3).

Default: REGISTERBANK (0)

Control Class: Primary

 μ Vision2 Control: Options – Ax51 – Misc controls: enter the control.

Description: The **REGISTERBANK** control specifies the register banks

used in a source module. This information is stored in the generated object file for further processing by the Lx51

linker/locator.

The **NOREGISTERBANK** control specifies that the **Ax51** assembler reserve no memory for the register bank. This is useful for assembler modules that should be used in a generic library. Since this library might be called with any active register bank, you may use the NOREGISTERBANK directive. Thus the program that calls the library module must reserve the register bank that is in use.

Examples: A251 RBUSER.A51 REGISTERBANK (0,1,2)

\$RB(0,3)

A51 SAMPLE.A51 NOREGISTERBANK

\$NORB

REGUSE

Abbreviation: RU

Arguments: Name of a PUBLIC symbol and a register list enclosed in

parentheses.

Default: Not applicable.

Control Class: General

μVision2 Control: This control cannot be specified on the command line.

Description: The **REGUSE** control specifies the registers modified

during a function's execution. This control may be used in combination with the C51 Compiler or C251 Compiler to allow global register optimization for functions written in assembly language. For more information about global register optimization refer to the *C Compiler User's Guide*.

The **REGUSE** control may not be specified on the A251

assembler invocation line.

Examples: \$REGUSE MYFUNC (ACC, B, R0 - R7)

\$REGUSE PROCA (DPL, DPH)

\$REGUSE PUTCHAR (R6, R7, CY, ACC)

SAVE / RESTORE

Abbreviation: SA / RS

Arguments: None

Default: None

Control Class: General

μVision2 Control: This control cannot be specified on the command line.

Description: The **SAVE** control stores the current settings of the **LIST**

and GEN controls. Subsequent controls can modify the

LIST and GEN settings.

This control allows these settings to be saved, altered for a number of program lines, and restored using the **RESTORE** control. The **SAVE** control can be nested up to 9 levels

deep.

The **RESTORE** control fetches and restores the values of the **GEN** and **LIST** controls that were stored by the last

SAVE control statement.

Example:

\$SAVE

\$INCLUDE (SAMPLE.INC)

\$RESTORE

•

SYMLIST / NOSYMLIST

Abbreviation: SL/NOSL

Arguments: None.

Default: SYMLIST

Control Class: General

 μ Vision2 Control: Options – Ax51 – Misc controls: enter the control.

Description: The **SYMLIST** control lists symbol definitions in the

symbol table.

The **NOSYMLIST** control prevents the **Ax51** assembler from listing symbol definitions in the symbol table. The **NOSYMLIST** control is useful in special function register definition files or other files where symbols are not desired

in the symbol table.

Examples: A251 SAMPLE.A51 NOSYMLIST

\$NOSYMLIST

\$INCLUDE (REG251S.H)

\$SYMLIST

TITLE

Abbreviation: TT

Arguments: A string enclosed within parentheses.

Default: The basename of the source file excluding the extension.

Control Class: General

 μ Vision2 Control: Options – Ax51 – Misc controls: enter the control.

Description: The **TITLE** control allows you to specify the title to use in

> the header line of the listing file. The text used for the title must immediately follow the TITLE control and must be enclosed in parentheses. A maximum of 60 characters may be specified for the title. If the TITLE control is not used, the module name specified with the "NAME" directive

described on page 132 will be used as title string.

A251 SAMPLE.A51 TITLE(Oven Controller Version 3) **Example:**

\$TT(Race Car Controller)

XREF

Abbreviation:

XR

Arguments:

None.

Default:

No error references are listed.

μVision2 Control: Options – Listing – Assembler Listing – Cross Reference

Description:

The XREF control directs the Ax51 assembler to generate a cross reference table of the symbols used in the source module. The alphabetized cross reference table will be included in the generated listing file. Refer to "Assembler Listing File Format" on page 415 for an example of a cross

reference table.

Example:

AX51 SAMPLE.A51 XREF

\$XREF

Controls for Conditional Assembly

The controls for conditional assembly are General controls—they may be specified any number of times in the body of a source file. Conditional assembly may be used to implement different program versions or different memory models with one source file. You may use conditional assembly to maintain one source module that satisfies several applications.

Conditional text blocks are enclosed by IF, ELSEIF, ELSE and ENDIF.

The **SET** and **RESET** controls may be used in the invocation line of the assembler to set and reset conditions tested by the **IF** and **ELSEIF** controls.

The remaining instructions for conditional assembly are only allowed within the source file and cannot be part of the assembler invocation line.

IF blocks may be nested a maximum of 10 levels deep. If a block is not translated, conditional blocks nested within it are also skipped.

Conditional Assembly Controls

Conditional assembly controls allow you to write *x*51 assembly programs with sections that can be included or excluded from the assembly based on the value of a constant expression. Blocks that are conditionally assembled are enclosed by **IF**, **ELSEIF**, **ELSE**, and **ENDIF** control statements.

The conditional control statements **IF**, **ELSE**, **ELSEIF**, and **ENDIF** may be specified only in the source program. They are not allowed on the invocation line. Additionally, these controls may be specified both with and without the leading dollar sign (\$).

When prefixed with a dollar sign, the conditional control statements may only access symbols defined by the **SET** and **RESET** controls.

When specified without a dollar sign, the conditional control statements may access all symbols except those defined by the **SET** and **RESET** controls. These control statements may access parameters in a macro definition.

The following table lists the conditional assembly control statements.

Directive	Page	Description
IF	233	Translate block if condition is true
ELSE	235	Translate block if the condition of a previous IF is false.
ELSEIF	234	Translate block if condition is true and a previous IF or ELSEIF is false.
ENDIF	236	Marks end of a block.
RESET	232	Set symbols checked by IF or ELSEIF to false.
SET	231	Set symbols checked by IF or ELSEIF to true or to a specified value.

Predefined Constants (A251 only)

The A251 macro assembler provides you with predefined constants to use in conditional \$IF / \$ELSEIF controls for more portable assembler modules. The following table lists and describes each one.

Constant	Description
INTR4	Set to 1 to when A251 assumes 4 byte interrupt frames. If the A251 control INTR2 is used, theINTR4 symbol is not defined.
MODBIN	Set to 1 if the binary mode of 251 CPU is used. If the source mode is specified with the MODSRC control, theMODBIN symbol is not defined.
MODSRC	Set to 1 if the source mode of 251 CPU is specified with the MODSRC control. If the binary mode of the 251 CPU is used theMODSRC symbol is not defined.

SET

Abbreviation: None.

Arguments: A list of symbols with optional value assignments separated

by commas and enclosed within parentheses. For example:

SET (symbol [= number] [, symbol <math>[= number] ...])

Default: None.

Control Class: General

 μ Vision2 Control: Options – Ax51 – Set.

Description: The **SET** control assigns numeric values to the specified

symbols. Symbols that do not include an explicit value assignment are assigned the value 0FFFFh. Symbols that are specified with an equal sign ('=') and a numeric value

are assigned the specified value.

These symbols can be used in **IF** and **ELSEIF** control statements for conditional assembly. They are only used for conditional assembly. These symbols are administered separately and do not interfere with other symbols.

Example: A251 SAMPLE.A51 SET(DEBUG1=1, DEBUG2=0, DEBUG3=1)

\$SET (TESTCODE = 0)

\$SET (DEBUGCODE, PRINTCODE)

RESET

Abbreviation: None.

Arguments: A list of symbols separated by commas and enclosed within

parentheses. For example:

RESET (symbol [, symbol ...])

Default: None

Control Class: General

 μ Vision2 Control: Options – Ax51 – Reset.

Description: The **RESET** control assigns a value of 0000h to the

specified symbols. These symbols may then be used in **IF** and **ELSEIF** control statements for conditional assembly. These symbols are only used for conditional assembly. They are administered separately and do not interfere with

other symbols.

Example: A251 SAMPLE.A51 RESET (DEBUG1, DEBUG2, DEBUG3)

\$RESET (TESTCODE)

\$RESET (DEBUGCODE, PRINTCODE)

IF

Abbreviation: None

Arguments: A numeric expression

Default: None

Control Class: General

μVision2 Control: This control cannot be specified on the command line.

Description: The **IF** control begins an **IF-ELSE-ENDIF** construct that

is used for conditional assembly. The specified numeric expression is evaluated and, if it is non-zero (TRUE), the **IF** block is assembled. If the expression is zero (FALSE), the **IF** block is not assembled and subsequent blocks of the **IF**

construct are evaluated.

IF blocks can be terminated by an ELSE, ELSEIF, or ENDIF control statement.

```
.
.
.
$IF (DEBUG_VAR = 3)
.
.
.
.
Version_3: MOV DPTR, #TABLE_3
.
.
.
$ ENDIF
.
.
```

ELSEIF

Abbreviation: None

Arguments: A numeric expression.

Default: None

μVision2 Control: This control cannot be specified on the command line.

Description: The **ELSEIF** control is used to introduce an alternative

program block after an **IF** or **ELSEIF** control. The **ELSEIF** block is only assembled if the specified numeric expression is non–zero (TRUE) and if previous **IF** and **ELSEIF** conditions in the **IF**–**ELSE**–**ENDIF** construct

were FALSE. **ELSEIF** blocks are terminated by an

ELSEIF, ELSE, or ENDIF control.

ELSE

Abbreviation: None.

Arguments: None.

Default: None.

Control Class: General

 μ Vision2 Control: This control cannot be specified on the command line.

Description: The **ELSE** control is used to introduce an alternative

program block after an **IF** or **ELSEIF** control. The **ELSE** block is only assembled if previous **IF** and **ELSEIF** conditions in the **IF-ELSE-ENDIF** construct were all FALSE. **ELSE** blocks are terminated with an **ENDIF**

control.

ENDIF

Abbreviation: None

Arguments: None

Default: None

Control Class: General

 μ Vision2 Control: This control cannot be specified on the command line.

Description: The ENDIF control terminates an IF-ELSE-ENDIF

> construct. When the Ax51 assembler encounters an ENDIF control statement, it concludes processing the IF block and resumes assembly at the point of the IF block. Since IF blocks may be nested, this may involve continuing in

> another IF block. The ENDIF control must be preceded by

an IF, ELSEIF, or ELSE control block.

```
$IF TEST
SENDIF
```

Chapter 8. Error Messages

This chapter lists the error messages generated by **Ax51**. The following sections include a brief description of the possible error messages along with a description of the error and any corrective actions you can take to avoid or eliminate the error.

Fatal errors terminate the assembly and generate a message that is displayed on the console. Non–fatal errors generate a message in the assembly listing file but do not terminate the assembly.

Fatal Errors

Fatal errors cause immediate termination of the assembly. These errors usually occur as a result of an invalid command line. Fatal errors are also generated when the assembler cannot access a specified source file or when the macros are nested more than 9 deep.

Fatal errors produce a message that conforms to one of the following formats:

```
A251 FATAL ERROR -
                                       <file in which the error occurred>
    FILE:
                                       line in which the error occurred
    LINE:
                                       <corresponding error message>
    ERROR:
A251 TERMINATED.
or
A251 FATAL ERROR -
                                       <error message with description>
    ERROR:
A251 TERMINATED.
where
                   is the name of an input file that could not be opened.
FILE
                   is the line where the error occurred
LINE
                   is the fatal error message text explained below.
ERROR
```

Fatal Error Messages

ATTEMPT TO SHARE FILE

A file is used both for input and output (e.g. list file uses the same name as the source file).

BAD NUMERIC CONSTANT

The numeric argument to the given control is illegal.

CAN'T ATTACH FILE

The given file can't be opened for read access.

CAN'T CREATE FILE

The given file can't be opened for write/update access.

CAN'T HAVE GENERAL CONTROL IN INVOCATION LINE

The given control is allowed in \$control lines within the source file only (for example the **EJECT** control). Some controls are allowed only in the source text and not in the command line. Refer to "Chapter 7. Invocation and Controls" on page 195 for more information about the A251 controls.

CAN'T REMOVE FILE

The given temporary file could not be removed for some reason.

CONFLICTING CONTROL

The given control conflicts with an earlier control (for example **\$NOMOD251 MODSRC**).

CONTROL LINE TOO LONG (500)

A \$-control line has more than 500 characters.

DISK FILE REQUIRED

The given file does not represent a disk file.

ERRORPRINT- AND LIST-FILE CANNOT BE THE SAME

It is illegal to direct the listing file output and the errorprint output to the console at the same time.

EXPECTED DELIMITER '(' AFTER CONTROL

The given control requires a brace enclosed argument

EXPECTED DELIMITER ') ' AFTER ARGUMENT

The given control requires a brace enclosed argument

FILE DOES NOT EXIST

The given file does not exist.

FILE IS READ ONLY

The given file does not permit write/update access.

FILE WRITE ERROR

The given file could not be written to (check free space)

IDENTIFIER EXPECTED

The given control requires an identifier as it's argument, for example **SET** (VAR1=1234H).

ILLEGAL FILE NAME, VOLUME OR DIRECTORY NAME

The name of the file is invalid or designates an invalid file.

INVOCATION LINE TOO LONG

The invocation line is longer than 500 characters.

LIMIT EXCEEDED: BALANCED TEXT LENGTH

The maximum length of a balanced text string is 65000 characters.

LIMIT EXCEEDED: INCLUDE OR MACRO NESTING

The maximum nesting level for MPL-macros is 50. The maximum nesting level of standard macros plus include files is 10.

LIMIT EXCEEDED: MACRO DEFINITION LENGTH

The maximum definition length of a standard macro is 20000 characters. MPL macros are limited to 65000 characters.

THE HACTOS are inniced to 05000 characters

LIMIT EXCEEDED: MORE THAN 16000 SYMBOLS

The number of symbols (labels, equ/set symbols, externals, segment-symbols) must not exceed 16000 per source file.

LIMIT EXCEEDED: SOURCE LINE LENGTH (500)

A single source line must not exceed the 500 characters per line limit.

LIMIT EXCEEDED: TOO MANY EXTERNALS (65535)

The number of external symbols must not exceed 65535 per source module.

LIMIT EXCEEDED: TOO MANY EXTERNALS (65535)

The number of externals must not exceed 65535 per source module.

LIMIT EXCEEDED: TOO MANY SEGMENTS (65535)

The number of segments must not exceed 65535 per source module.

NON-NULL ARGUMENT EXPECTED

The argument to the given control must not be null (for example **\$PRINT()**).

OUT OF MEMORY

The assembler has run out of memory. Remove unnecessary drivers from your system configuration.

OUT OF RANGE NUMERIC VALUE

The numeric argument to the given control is out of range (for example **\$PAGEWIDTH(3000)**).

UNKNOWN CONTROL

The given control is undefined.

Non-Fatal Errors

Non-fatal errors usually occur within the source program and are typically syntax errors. When one of these errors is encountered, the assembler attempts to recover and continue processing the input file. As more errors are encountered, the assembler will produce additional error messages. The error messages that are generated are included in the listing file.

Non–fatal errors produce a message using the following format:

```
*** ERROR number IN line (file, LINE line): error message
or
*** WARNING number IN line (file, LINE line): warning message
where
                   is the error number.
number
                   corresponds to the logical line number in the source file.
line
                   corresponds to the source or include file which contains the
file
                   error.
LINE
                   corresponds to the physical line number in <file>.
                   is descriptive text and depends on the type of error
error message
                   encountered.
```

The logical line number in the source file is counted including the lines of all include files and may therefore differ from the physical line number. For that reason, the physical line number and the associated source or include file is also given in error and warning messages.

Example

```
11 MOV R0,# 25 * | 10
*** ------^
*** ERROR #4 IN 11 (TEST.A51, LINE 11), ILLEGAL CHARACTER
```

The caret character (^) is used to indicate the position of the incorrect character or to identify the point at which the error was detected. It is possible that the position indicated is due to a previous error. If a source line contains more than one error, the additional position indicators are displayed on subsequent lines.

The following table lists the non–fatal error messages that are generated by A251. These messages are listed by error number along with the error message and a brief description of possible causes and corrections.

Number	Non-Fatal Error Message and Description
1	ILLEGAL CHARACTER IN NUMERIC CONSTANT
	This error indicates that an invalid character was found in a numeric constant. Numeric constants must begin with a decimal digit and are delimited by the first non–numeric character (with the exception of the dollar sign). The base of the number decides which characters are valid.
	 Base 2: 0, 1 and the base indicator B Base 8: 0–7 and the base indicator O or Q Base 10: 0–9 and the base indicator D or no indicator Base 16: 0–9, A–F and the base indicator H Base 16: 0xhhhh, 0–9, and A–F
2	MISSING STRING TERMINATOR
	The ending string terminator was missing. The string was terminated with a carriage return.
3	ILLEGAL CHARACTER
	The assembler has detected a character which is not in the set of valid characters for the 51/251 assembler language (for example `).
4	BAD INDIRECT REGISTER IDENTIFIER
'	This error occurs if combined registers are entered incorrectly; e.g.,
	@R7, @R3, @PC+A, @DPTR+A.
5	ILLEGAL USE OF A RESERVED WORD
<u> </u>	This error indicates that a reserved word is used for a label.
6	DEFINITION STATEMENT EXPECTED
	The first symbol in the line must be part of a definition. For example:
	VAR1 EQU 12
_	
7	LABEL NOT PERMITTED
	A label was detected in an invalid context.
8	ATTEMPT TO DEFINE AN ALREADY DEFINED LABEL
	A label was defined more than once. Labels may be defined only once in the source program.
9	SYNTAX ERROR
	Ax51 encountered an error processing the line at the specified token.

Number	Non–Fatal Error Message and Description
10	ATTEMPT TO DEFINE AN ALREADY DEFINED SYMBOL
	An attempt was made to define a symbol more than once. The subsequent definition was ignored.
11	STRING CONTAINS ZERO OR MORE THAN TWO CHARACTERS
	Strings used in an expression can be a maximum of two characters long (16 bits).
12	ILLEGAL OPERAND
	An operand was expected but was not found in an arithmetic expression. The expression is illegal.
13	')' EXPECTED
	A right parenthesis is expected. This usually indicates an error in the definition of external symbols.
14	BAD RELOCATABLE EXPRESSION
	A relocatable expression may contain only one relocatable symbol which may be a segment symbol, external symbol, or a symbol belonging to a relocatable segment. Mathematical operations cannot be carried out on more than one relocatable symbol.
15	MISSING FACTOR
	A constant or a symbolic value is expected after an operator.
16	DIVIDE BY ZERO ERROR
	A division by zero was attempted while calculating an expression. The value calculated is undefined.
17	INVALID BASE IN BIT ADDRESS EXPRESSION
	This error indicates that the byte base in the bit address is invalid. This occurs if the base is outside of the range 20h–2Fh or if it lies between 80h and 0FFh and is not evenly divisible by 8. For the 251 chip, the byte base address must be in range 20H-0FFH with no restrictions. Note that with symbolic operands, the operand specifies an absolute bit segment or an addressable data segment.
18	OUT OF RANGE OR NON-TYPELESS BIT-OFFSET
	The input of the offset (base.offset) in a bit address must be a typeless absolute expression with a value between 0 and 7.
19	INVALID REGISTER FOR EQU/SET
	The registers R0–R7, A and C may be used in SET or EQU directives. No other registers are allowed.

Number	Non-Fatal Error Message and Description
20	INVALID SIMPLE RELOCATABLE EXPRESSION
	A simple relocatable expression is intended to represent an address in a relocatable segment. External symbols as well as segment symbols are not allowed. The expression however may contain more symbols of the same segment. Simple relocatable expressions are allowed in the instructions ORG, EQU, SET, CODE, XDATA, IDATA, BIT, DATA, DB and DW.
21	EXPRESSION WITH FORWARD REFERENCE NOT PERMITTED
'	Expressions in EQU and SET directives may not contain forward references.
22	EXPRESSION TYPE DOES NOT MATCH INSTRUCTION
	The expression does not conform to the x51 conventions. A #, /Bit, register, or numeric expression was expected.
23	NUMERIC EXPRESSION EXPECTED
	A numeric expression is expected. The expression of another type is found.
24	SEGMENT-TYPE EXPECTED
	The segment type of a definition was missing or invalid.
25	RELOCATION-TYPE EXPECTED
	An invalid relocation type for a segment definition was encountered.
26	INVALID RELOCATION-TYPE
	The types PAGE and INPAGE are only allowed for the CODE/ECODE and XDATA segments. INBLOCK/INSEG is only allowed for the CODE/ECODE segments and BITADDRESSABLE is only for the DATA segment (maximum length 16 Bytes). EBITADDRESSABLE is allowed for DATA segments (maximum length 96 Bytes). The type UNIT is the default for all segment types if no input is entered.
27	LOCATION COUNTER MAY NOT POINT BELOW SEGMENT-BASE
	An ORG directive used in a segment defined by the AT address directive may not specify an offset that lies below the segment base. The following example is, therefore, invalid: CSEG AT 1000H
	ORG 800H
28	ABSOLUTE EXPRESSION REQUIRED
	The expression in a DS or DBIT instruction must be an absolute typeless expression. Relocatable expressions are not allowed.

Number	Non-Fatal Error Message and Description
29	SEGMENT-LIMIT EXCEEDED
	The maximum limit of a segment was exceeded. This limit depends on the segment and relocation type. Segments with the attribute DATA should not exceed 128 bytes. BITADDRESSABLE segments should not exceed 16 bytes and INPAGE segments should not exceed 2 KBytes.
30	SEGMENT-SYMBOL EXPECTED
	The operand to an RSEG directive must be a segment symbol.
31	PUBLIC-ATTRIBUTE NOT PERMITTED
	The PUBLIC attribute is not allowed on the specified symbol.
32	ATTEMPT TO RESPECIFY MODULE NAME
	An attempt was made to redefine the name of the module by using a second NAME directive. The NAME directive may only appear once in a program.
33	CONFLICTING ATTRIBUTES
	A symbol may not contain the attributes PUBLIC and EXTRN simultaneously.
34	',' EXPECTED
	A comma is expected in a list of expressions or symbols.
35	'(' EXPECTED
	A left parenthesis is expected at the indicated position.
36	INVALID NUMBER FOR REGISTERBANK
	The expression in a REGISTERBANK control must be an absolute typeless number between 0 and 3.
37	OPERATION INVALID IN THIS SEGMENT
	x51 instructions are allowed only within CODE/ECODE segments.
38	NUMBER OF OPERANDS DOES NOT MATCH INSTRUCTION
	Either too few or too many operands were specified for the indicated instruction. The instruction was ignored.
39	REGISTER-OPERAND EXPECTED
	A register operand was expected but an operand of another type was found.
40	INVALID REGISTER
	The specified register operand does not conform to the <i>x</i> 51 conventions.

Number	Non-Fatal Error Message and Description
41	MISSING 'END' STATEMENT
	The last instruction in a source program must be the END directive. The preceding source is assembled correctly and the object is valid.
42	INTERNAL ERROR (PASS-2), CONTACT TECHNICAL SUPPORT
	Occurs if a symbol in pass 2 contains a value different than in pass 1.
43	RESPECIFIED PRIMARY CONTROL, LINE IGNORED
	A control was repeated or conflicts with a previous control. The control statement was ignored.
44	MISPLACED PRIMARY CONTROL, LINE IGNORED
	A primary control was misplaced. Primary controls may be entered in the invocation line or at the beginning of the source file (as \$ instruction). The processing of primary controls in a source file ends when the first non empty/non comment line containing anything but a primary control is processed.
45	UNDEFINED SYMBOL (PASS-2)
	The symbol is undefined.
46	CODE/ECODE-ADDRESS EXPECTED
	An operand of memory type CODE/ECODE or a typeless expression is expected.
47	XDATA-ADDRESS EXPECTED
	An operand of memory type XDATA or a typeless expression is expected.
48	DATA-ADDRESS EXPECTED
	An operand of memory type DATA or a typeless expression is expected.
49	IDATA-ADDRESS EXPECTED
	An operand of memory type 'IDATA' or a typeless expression is expected.
50	BIT-ADDRESS EXPECTED
	An operand of memory type BIT or a typeless expression is expected.
51	TARGET OUT OF RANGE
	The target of a conditional jump instruction is outside of the +127/–128 range or the target of an AJMP or ACALL instruction is outside the 2 KByte memory block.

Number	Non-Fatal Error Message and Description
52	VALUE HAS BEEN TRUNCATED TO 8 BITS
	The result of the expression exceeds 255 decimal. Only the 8 low–order bits are used for the byte operand.
53	MISSING 'USING' INFORMATION
	The absolute register symbols AR0 through AR7 can be used only if a USING register bank directive was specified. This error indicates that the USING directive is missing and the assembler cannot assign data addresses to the register symbols.
54	MISPLACED CONDITIONAL CONTROL
	An ELSEIF, ELSE, or ENDIF control must be preceded by an IF instruction.
55	BAD CONDITIONAL EXPRESSION
	The expression to the IF or ELSEIF control is invalid. These expressions must be absolute and may not contain relocatable symbols.
	The \$IF and \$ELSEIF can only access symbols defined with the \$SET and \$RESET controls. Both IF and ELSEIF allow access to all symbols of the source program.
56	UNBALANCED IF-ENDIF-CONTROLS
	Each IF block must be terminated with an ENDIF control. This is also true with skipped nested IF blocks.
57	SAVE STACK UNDERFLOW
	A \$RESTORE control instruction is then valid only if a \$SAVE control was previously given.
58	SAVE STACK OVERFLOW
	The context of the GEN, COND , and LIST controls may be stored by the \$SAVE control up to a maximum of 9 levels.
59	MACRO REDEFINITION
	An attempt was made to define an already defined macro.
60	Not generated by Ax51 .
61	MACRO TERMINATED BY END OF FILE, MISSING 'ENDM'
	An attempt was made to define an already defined macro.
62	TOO MANY FORMAL PARAMETERS (16)
	The number of formal parameters to a macro is limited to 16.

Number	Non-Fatal Error Message and Description
63	TOO MANY LOCALS (16)
	The number of local symbols within a macro is limited to 16.
64	DUPLICATE LOCAL/FORMAL
	The number of local or formal identifier must be distinct.
65	IDENTIFIER EXPECTED
	While parsing a macro definition, an identifier was expected but something different was found.
66	'EXITM' INVALID OUTSIDE A MACRO
	The EXITM (exit macro) keyword is illegal outside a macro definition.
67	EXPRESSION TOO COMPLEX
	A too complex expression was encountered. This error occurs, if the number of operands and operators in one expression exceeds 50.
68	UNKNOWN CONTROL OR BAD ARGUMENT(S)
	The control given in a \$-control line or the argument(s) to some control are invalid.
69	MISPLACED ELSEIF/ELSE/ENDIF CONTROL
	These controls require a preceding IF control.
70	LIMIT EXCEEDED: IF-NESTING (10)
	IF controls may be nested up to a level of 10.
71	NUMERIC VALUE OUT OF RANGE
	The value of a numeric expression is out of range (for example \$PAGEWIDTH (2048) where only values in range 80 to 132 are allowed).
72	TOO MANY TOKENS IN SOURCE LINE
	The number of tokens (identifiers, operators, punctuation characters and end of line) exceeds 200. The source line is truncated at 200 tokens.
72	TOO MANY TOKENS IN SOURCE LINE
	The number of tokens (identifiers, operators, punctuation characters and end of line) exceeds 200. The source line is truncated at 200 tokens.
73	TEXT FOUND BEYOND END STATEMENT - IGNORED
	Text following the END directive is not processed by the assembler.

Number	Non-Fatal Error Message and Description
74	REGISTER USAGE: UNDEFINED REGISTER NAME
	A register name argument given to the REGUSE control does not represent the name of a register.
75	'REGISTER USAGE' REQUIRES A PUBLIC CODE SYMBOL
	The register usage value must be assigned to a public symbol, which represents a CODE or ECODE symbol.
76	MULTIPLE REGISTER USES GIVEN TO ONE SYMBOL
	The register usage value may be assigned to a symbol or procedure only once.
77	INSTRUCTION NOT AVAILABLE
	The given instruction is not available in the current mode of operation.
78	Not generated by Ax51 .
79	INVALID ATTRIBUTE
	The OVERLAYABLE attribute given in a segment definition is not valid for code and constant segments.
80	INVALID ABSOLUTE BASE/OFFS VALUE
	The absolute address given in a segment definition does not conform to the memory type of the segment (for example DATA AT 0x1000).
81	EXPRESSION HAS DIFFERENT MEMORY SPACE
	The expression given in a symbol definition statement does not have the memory space required by the directive, for example:
	VAR1 CODE EXPR
	where 'EXPR' has a memory type other than CODE or NUMBER.
82	LABEL STATEMENT MUST BE WITHIN CODE/ECODE SEGMENT
	The LABEL statement is not allowed outside a CODE or ECODE segment.
83	TYPE INCOMPATIBLE WITH GIVEN MEMORY SPACE
	The type given in an external declaration is not compatible to the given memory space. The following examples shows an invalid type since a bit can never reside in code space:
	EXTRN CODE:BIT (bit0, bit1)
84	OPERATOR REQUIRES A CODE/ECODE ADDRESS
	The type override operators NEAR and FAR cannot be applied to addresses with memory type other than CODE and ECODE.

Number	Non–Fatal Error Message and Description
85	INVALID OPERAND TYPE
	An expression contains invalid typed operands to some operator, for example addition/unary minus on bit-type operands.
86	PROCEDURES CAN'T BE NESTED
	A251 does not support nested procedures.
87	UNCLOSED PROCEDURE
	A251 detected an unclosed procedure after scanning the source file.
88	VALUE HAS BEEN TRUNCATED TO 16 BITS
	The displacement value given in a register expression (WRn+disp16, DRk+disp16) has been truncated to 16 bits.
89	Not generated by Ax51 .
90	'FAR' RETURN IN 'NEAR' PROCEDURE
	The return far instruction (ERET) was encountered in a procedure of type NEAR (the code may not work).
91	TYPE MISMATCH
	The operand type of an instruction operand does not match the requested type of the instruction, for example:
	MOV WR10,Byte_Memory_Operand. ; Word/Byte mismatch
	Use a type override to avoid the warning as shown:
	MOV WR10,WORD Byte_Memory_Operand
92, 93	Not generated by Ax51 .
94	VALUE DOES NOT MATCH INSTRUCTION
	The short value given in a INC/DEC Rn,#short is not one of 1,2,4.
95	ILLEGAL MEMORY CLASS SPECIFIER
	The memory class specifier in a segment definition statement does not correspond to one of the predefined memory class names (CODE, ECODE, BIT, EBIT).
96	ACCESS TO MISALIGNED ADDRESS
	A word instruction accesses a misaligned (odd) address. This warning is generated only if the \$WORDALIGN control was given.

Number	Non–Fatal Error Message and Description
97	'FAR' REFERENCE TO 'NEAR' LABEL
	An ECALL/AJMP instruction to some label of type NEAR has been detected.
98	'NEAR' REFERENCE TO 'FAR' LABEL
	An ACALL/AJMP/SJMP or conditional jump instruction to some label of type FAR has been detected.
99	'PROC' NAME REQUIRED
	Ax51 expects the name of the procedure.
100	ILLEGAL CONSTANT VALUE
	The constant value is illegal or has an illegal format.
101	TRAP INSTRUCTION IS RESERVED FOR DEBUGGING TOOLS
	The TRAP instruction should be used for program debugging only.
102	PONTENTIAL ADDRESS OVERLAP
	There is a potential address overlay in your program that is caused by ORG statements.
103	<user error="" text=""></user>
	This error is generated by the C preprocessor #error directive or theERROR directive.
104 - 149	Not generated by Ax51 .
150	PREMATURE END OF FILE ENCOUNTERED
	The MPL macro processor encountered the end of the source file while parsing a macro definition.
151	<name>: IDENTIFIER EXPECTED</name>
	The macro or function given by <name> in the error message expected an identifier but found something else.</name>
152	MPL FUNCTION <name>: <character> EXPECTED</character></name>
	The MPL function <name> expected a specific character in the input stream but found some other character.</name>
153	<name>: UNBALANCED PARENTHESIS</name>
	While scanning balanced text, the macro processor expected a ')' character, but found some other character.

Number	Non–Fatal Error Message and Description
154	EXPECTED <identifier></identifier>
	The macro processor expected some specific identifier (for example ELSE) but found some other text.
155	Not generated by Ax51 .
156	FUNCTION 'MATCH': ILLEGAL CALL PATTERN
	The call pattern to the MPL function match must be a formal parameter followed by a delimiter followed by another formal parameter.
157	FUNCTION 'EXIT' IN BAD CONTEXT
	The EXIT function must not appear outside a macro expansion, %REPEAT or %WHILE.
158	ILLEGAL METACHARACTER <character></character>
	The metacharacter may not be @, (,), *, TAB, EOL, A-Z,a-z, 0-9, _ and ?.
159	CALL PATTERN - DELIMITER <delimiter> NOT FOUND</delimiter>
	The actual parameters in a macro call do not match the call pattern defined in the macro definition of that macro.
160	CALL TO UNDEFINED MACRO <macroname></macroname>
	An attempt to activate an undefined macro has been encountered .
161	ERROR-161
	Not generated by Ax51 .
162	INVALID DIGIT 'character' IN NUMBER
	An ill formed number has been encountered. For numbers, the rules are equal to the numbers in the assembler language with the exception of \$ signs, which are not supported within the MPL.
163	UNCLOSED STRING OR CHARACTER CONSTANT
	A string or character constant is terminated by an end of line character instead of the closing character.
164	INVALID STRING OR CHARACTER CONSTANT
	A string or character constant may contain one or two characters.
165	EVAL: UNKNOWN EXPRESSION IDENTIFIER
	An MPL expression contains an unknown identifier.

Number	Non–Fatal Error Message and Description
166	<token>: INVALID EXPRESSION TOKEN</token>
	An MPL expression contains a token which neither represents an operator nor an operand.
167	<function>: DIV/MOD BY ZERO</function>
	The evaluation of an expression within the MPL function <function> yields a division or modulus by zero.</function>
168	EVAL: SYNTAX ERROR IN EXPRESSION
	An expression is followed by one or more erroneous tokens.
169	CAN'T OPEN FILE <name file="" of=""></name>
	The file given in an \$INCLUDE control cannot be opened.
170	<name file="" of="">: IS NOT A DISK FILE</name>
	An attempt was made to open a file which is not a disk file (for example \$INCLUDE (CON).
171	ERROR IN INCLUDE DIRECTIVE
	The argument to the INCLUDE control must be the brace enclosed name of the file, for example \$INCLUDE (REG251.INC).
172	CAN'T REDEFINE PREDEFINED MACRO `SET'
	The .predefined %SET macro can't be redefined.

Chapter 9. Linker/Locator

The Lx51 linker/locator is used to link or join together object modules that were created using the Ax51 assembler, the Cx51 compiler, and the Intel PL/M-51 compiler. Object modules that are created by these translators are relocatable and cannot be directly executed. They must be converted into absolute object modules. The Lx51 linker/locator does this and much more.

For optimum support of the different 8051 and 251 variants, the following linker/locater variants are available. The **LX51** and **L251** linker/locater provide an improved memory allocation handling and are supersets of the **BL51** linker/locater. The following table gives you an overview of the linker/locater variants along with the translators that are supported.

Linker/Locater	Processes Files from	Description
BL51 Code Banking Linker/Locater	Keil A51 Macro Assembler Keil C51 Compiler Intel ASM51 Assembler Intel PL/M51 Compiler	For classic 8051 . Includes support for 32 x 64KB code banks.
LX51 Extended Linker/Locater	Keil A51 Macro Assembler Keil C51 Compiler Keil AX51 Macro Assembler Keil CX51 Compiler for 80C51MX Intel ASM51 Assembler Intel PL/M51 Compiler	For classic 8051 and extended 8051 versions (Philips 80C51MX, Dallas 390, etc.) Allows code and data banking and supports up to 16MB code and xdata memory.
L251 Linker/Locater	Keil A51 Macro Assembler Keil C51 Compiler Keil A251 Macro Assembler Keil C251 Compiler Intel ASM51 Assembler Intel PL/M51 Compiler	For Intel/Atmel WM 251.

Programs you create using the **Ax51** Assembler and the **Cx51** Compiler must be linked using the **Lx51** linker/locator. You cannot execute or simulate programs that are not linked, even if they consist of only one source module. The **Lx51** linker/locator will link one or more object modules together and will resolve references within these modules. This allows you to create a large program that is spread over a number of source files.

The Lx51 linker/locator provides the following functions:

- Combines several program modules into one module, automatically incorporating modules from the library files
- Combines relocatable partial segments of the same segment name into a single segment
- Allocates and manipulates the necessary memory for the segments with which all relocatable and absolute segments are processed
- Analyzes the program structure and manipulates the data memory using overlay techniques
- Resolves external and public symbols
- Defines absolute addresses and computes the addresses of relocatable segments
- Produces an absolute object file that contains the entire program
- Produces a listing file that contains information about the Link/Locate procedure, the program symbols, and the cross reference of public and external symbol names
- Detects errors found in the invocation line or during the Link/Locate run.
- Supports programs that are larger than 64 Kbytes and applications that are using a Real-Time Multitasking Operating System (RTX51, RTX251, etc.).

All of these operations are described in detail in the remaining sections of this chapter.

"Overview" on page 257 provides you with a summary of the features and capabilities of the **BL51** linker/locator. This chapter introduces the concepts of what a linker is and does.

"Linking Programs" on page 266 describes how to invoke the linker from the command line. The command-line arguments are discussed, and examples are provided.

"Locating Programs to Physical Memory" on page 273 shows how to specify the physical memory available in the target hardware and how to locate segments to specific addresses.

"Data Overlaying" on page 280 explains how the **Lx51** linker/locater creates a call tree for segment overlaying of local variables and discusses how to modify this call tree for applications that use indirect program calls.

"Tips and Tricks for Program Locating" on page 289 shows you several additional features of the Lx51 linker/locater. These features allow you to create in-system programmable applications, to determine the addresses of segments, to use the C251 memory class NCONST without ROM in segment 0, or to locate several segments within a 2KB block.

"Bank Switching" on page 293 describes what bank switching is and how it is implemented by the Lx51 linker/locator. This chapter also shows how to make applications that are larger than 64 KBytes work with code banking.

"Control Summary" on page 305 lists the command-line controls by category and provides you with descriptions of each, along with examples.

"Error Messages" on page 360 lists the errors that you may encounter when you use the **Lx51** linker/locator.

Overview

The **Lx51** linker/locator takes the object files and library files you specify and generates a absolute object file. Absolute object files can be loaded into debugging tools or may be converted into Intel HEX files for PROM programming by OHx51 Object-Hex Converter.

NOTE

Banked object files generated by the **BL51** linker/locater must be converted by the **OC51** Banked Object File Converter into absolute object files (one for each bank) before they can be converted into Intel HEX files by the **OH51** Object-Hex Converter.

While processing object and library files, the Lx51 linker/locator performs the following operations.

Combining Program Modules

The object modules that the Lx51 linker/locator combines are processed in the order in which they are specified on the command line. The Lx51 linker/locator processes the contents of object modules created with the Ax51 assembler or the Cx51 compiler. Library files, however, contain a number of different object modules; and, only the object modules in the library file that specifically resolve external references are processed by the Lx51 linker/locator.

Segment Naming Conventions

Objects generated by the **Cx51** and Intel PL/M-51 compilers are stored in segments, which are units of code or data memory. A segment may be relocatable or may be absolute. Each relocatable segment has a type and a name. This section describes the conventions used for naming these segments.

Segment names include a *module_name*. The *module_name* is the name of the source file in which the object is declared and excludes the drive letter, path specification, and file extension. In order to accommodate a wide variety of existing software and hardware tools, all segment names are converted and stored in uppercase.

Each segment name has a prefix (or in case of PL/M-51 a postfix) that corresponds to the memory type used for the segment. The prefix is enclosed in question marks (?). The following is a list of the standard segment name prefixes.

Segment Prefix	Memory Class	Description
?PR?	CODE	Executable program code
?CO?	CONST	Constant data in program memory
?ED?	EDATA	EDATA memory for near variables
?FD?	HDATA	HDATA memory for far variables
?XD?	XDATA	XDATA memory
?DT?	DATA	DATA memory
?ID?	IDATA	IDATA memory
?BI?	BIT	Bit data in internal data memory
?BA?	DATA	Bit-addressable data in internal data memory
?PD?	XDATA	Paged data in XDATA memory

Combining Segments

A segment is a code or data block that is created by the compiler or assembler from your source code. There are two basic types of segments: absolute and relocatable. Absolute segments reside in a fixed memory location. They cannot be moved by the linker. Absolute segments do not have a segment name and will not be combined with other segments. Relocatable segments have a name and a type (as well as other attributes shown in the table below). Relocatable segments with the same name but from different object modules are considered parts of the same segment and are called partial segments. The linker/locator combines these partial relocatable segments.

The following table lists the segment attributes that are used to determine how to link, combine, and locate code or data in the segment.

Attribute	Description
Name	Each relocatable segment has a name that is used when combining relocatable segments from different program modules. Absolute segments do not have names.
Memory Class	The memory class identifies the address space to which the segment belongs. For BL51 the type can be CODE, XDATA, DATA, IDATA, or BIT. LX51 and L251 support in addition CONST, EBIT, ECONST, EDATA, HDATA, HCODE, HCONST, and user-define memory classes.
Relocation Type	The relocation type specifies the relocation operations that can be performed by the linker/locator. Valid relocation types are AT address, BITADDRESSABLE, INBLOCK, INPAGE, INSEG, OFFS offset, and OVERLAYABLE.
Alignment Type	The alignment type specifies the alignment operations that can be performed by the linker/locator. Valid alignment types are BIT, BYTE, WORD, DWORD, PAGE, BLOCK, and SEG.
Length	The length attribute specifies the length of the segment.
Base Address	The base address specifies the first assigned address of the segment. For absolute segments, the address is assigned by the assembler. For relocatable segments, the address is assigned by the linker/locator.

While processing your program modules, the linker/locator produces a table or map of all segments. The table contains name, type, location method, length, and base address of each segment. This table aids in combining partial relocatable segments. All partial segments having the same name are combined by the linker/locator into one single relocatable segment.

The linker/locator uses the following rules when combining partial segments.

- All partial segments that share a common name must have the same memory class. An error occurs if the types do not correspond.
- The length of the combined segments must not exceed the length of the physical memory area.
- The location method for each of the combined partial segments must correspond.

Absolute segments are not combined with other absolute segments, they are copied directly to the output file.

Locating Segments

After the linker/locator combines partial segments it must determine a physical address for them. The linker/locator processes each memory class separately. Refer to "Memory Classes and Memory Layout" on page 27 for a discussion of the different memory class and the physical address ranges.

After the linker/locator combines partial segments, it must determine a physical address for them. The linker/locator places different segments in each of these memory areas. The memory is allocated in the following order:

- 1. Register Banks and segments with an absolute address.
- 2. Segments specified in Lx51 segment allocation controls.
- 3. Segments with the relocation type BITADDRESSABLE and other BIT segments.
- 4. All other segments with the memory class DATA.
- 5. Segments with the memory class IDATA, EDATA and NCONST.
- 6. Segments with the memory class XDATA.
- 7. Segments with the memory class CODE and the relocation type INBLOCK.
- 8. Other Segments with the memory class CODE and CONST.
- 9. Segments with the memory classes ECODE, HCONST, and HDATA.

Overlaying Data Memory

The stack addressing of the *x***51** CPU is slower compared to accessing fixed, absolute memory locations. For this reason, local variables and function arguments of C and PL/M-51 routines are stored at fixed memory locations rather than on the stack. By using techniques to overlay the parameters and local variables of C and PL/M-51 functions, the linker/locator attempts to maximize the amount of available space.

NOTE

The **Cx51** compiler supports also reentrant functions where the parameters and automatic variables are store on the CPU stack of a simulate stack. For detailed information refer to the **Cx51** User's Guides.

To accomplish overlaying, the linker/locator analyzes all references or calls between the various functions. Using this information, the linker/locator can determine precisely which data and bit segments can be overlaid.

You may use the **OVERLAY** and **NOOVERLAY** control to enable or disable data overlaying. The **OVERLAY** control is the default and allows for very compact data areas. Use the **NOOVERLAY** control to disable the segment overlay function.

Resolving External References

External symbols reference addresses in other modules. A declared external symbol must be resolved with a public symbol of the same name. Therefore, for each external symbol, a public symbol must exist in another module.

The linker/locator builds a table of all public and external symbols that it encounters. External references are resolved with public references as long as the names match and the symbol attributes correspond.

The linker/locator reports an error if the symbol types of an external and public symbol do not correspond. The linker/locator also reports an error if no public symbol is found for an external reference.

The absolute addresses of the public symbols are resolved after the location of the segments is determined.

Absolute Address Calculation

After the segments are assigned fixed memory locations and external and public references are processed, the linker/locator calculates the absolute addresses of the relocatable addresses and external addresses. Symbolic debugging information is also updated to reflect the new addresses.

Generating an Absolute Object File

The linker/locator generates the executable target program in absolute object module format. The generated object module may contain debugging information if the linker/locator is so directed. This information facilitates symbolic debugging and testing. You may use the linker controls to suppress debugging information in the object file.

The output file generated by the linker/locator may be loaded into the μ Vision2 debugger, in-circuit emulators, or may be translated into an Intel HEX file for use with an EPROM programmer. The following table provides an overview of the output format and the processing method for the different linker/locater variants.

Linker/Locater	Output Format	Description
BL51 Linker/Locater	Extended Intel OMF51	Intel OMF51 is the standard format for programming the 8051 and supported by virtually all emulator vendors. Extensions in this format provide symbolic information.
	Banked OMF51	For banked applications the BL51 Linker/Locater generates a banked OMF51 file that can be converted with the OC51 Object File Converter into standard OMF51 files. The OC51 step is also required to convert the file into an Intel HEX file.
LX51 Extended Linker/Locater	Keil OMFx51	The Keil OMFx51 format supports of up to 16MB code and xdata memory. This format is required for extended 8051 versions (Philips 80C51MX , Dallas 390 , etc.). Check with your emulator vendor if this format is supported.
L251 Linker/Locater	Intel OMF251	Intel OMF251 is the standard format for programming the 251 and supported by all 251 emulator vendors.

Generating a Listing File

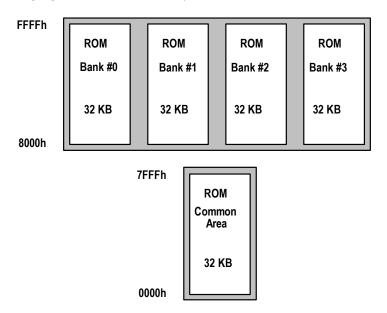
The linker/locator generates a listing file that lists information about each step in the link and locate process. This file also contains information about the symbols and segments involved in the linkage. In addition, the following information may be found in the listing file:

- The filenames and other parameters specified on the command line.
- Filenames and module names of all processed modules.
- A memory allocation table, which contains the location of the segments, the segment type, the location method, and the segment name. This table may be suppressed by specifying the NOMAP control.
- The overlay map which shows the structure of the finished program and lists address information for the local data and bit segments of a function. The overlay map also lists all code segments for which OVERLAYABLE segments exist. You may suppress the overlay map by specifying the NOMAP control.
- LX51 and L251 provide a list of all PUBLIC symbols within a program.
- A list of all errors in segments and symbols. The error causes are listed at the end of the listing file.
- A list of all unresolved external symbols. An external symbol is unresolved if no corresponding public symbol exists in another input file. Each reference to an unresolved external symbol is listed in an error message at the end of the listing file.
- A symbol table, which contains the symbol information from the input files.
 This information consists of the names of the MODULES, SYMBOLS,
 PUBLICS, and LINES. You may selectively suppress the symbolic information with linker controls.
- An alphabetically sorted cross reference report of all PUBLIC and EXTERN symbols in which the memory type and the module names that contain a reference to that symbol are displayed.
- Errors detected during the execution of the linker/locator are displayed on the screen as well as at the end of the listing file. Refer to "Error Messages" on page 360 for a summary of the linker/locator errors and causes.

Bank Switching

The classic 8051 directly supports a maximum of 64 KBytes of code space. The **Lx51** linker/locator allows 8051 programs to be created that are larger than 64 KBytes by using a technique known as code banking or bank switching. Bank switching involves using extra hardware to select one of a number of code banks all of which will reside at a common physical address.

For example, your hardware design may include a ROM mapped from address 0000h to 7FFFh (known as the common area) and four 32K ROM blocks mapped from code address 8000h to 0FFFFh (known as the code bank area). The following figure shows the memory structure.



The code banking facility of **Lx51** is compatible with the **C51** compiler, the **CX51** compiler, and the Intel **PL/M-51** compiler. Modules written using on of these compilers can be easily used in code banking applications. No modifications to the original source files are required.

Refer to "Bank Switching" on page 293 for detailed information about memory banking and instructions for building code banking programs.

Using RTX51, RTX251, and RTX51 Tiny

Programs you create that utilize the RTX51, and RTX51 Tiny Real-Time Operating Systems must be linked using the **BL51** or the **LX51** linker/locator. The **RTX51** and **RTX51TINY** controls enable link-time options that are required to generate RTX51 Full and RTX51 Tiny applications.

Programs that use the RTX251 Full Real-Time Operating Systems must be linked using the **L251** linker/locator. The **RTX251** control enable link-time options that are required to generate RTX251 Full applications.

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Linking Programs

The **Lx51** linker/locater is invoked by typing the program name at the Windows command prompt. On this command line, you must include the name of the assembler source file to be translated, as well as any other necessary assembler controls required to translate your source file. The format for the **Lx51** command line is:

```
BL51 [inputlist] [TO outputfile] [controls]

LX51 [inputlist] [TO outputfile] [controls]

L251 [inputlist] [TO outputfile] [controls]
```

or

```
BL51 @commandfile

LX51 @commandfile

L251 @commandfile
```

where

inputlist

is a list of the object files, separated by commas, for the linker/locator to include in the outputfile. The inputlist can contain files from Ax51, Cx51, PL/M-51 and library files. For library files you may force the inclusion of modules by specifying the module names in parentheses. The format of the inputlist is described below.

outputfile

is the name of the absolute object file that the linker/locator creates. If no <code>outputfile</code> is specified on the command line, the first filename in the input list is used. The basename of the <code>outputfile</code> is also the default name for the map file.

controls

are commands and parameters that control the operation of the Lx51 linker/locator.

commandfile

is the name of a command input file that may contain an inputlist, outputfile, and controls. The text in a commandfile has the same format as the standard command line and is produced by any standard ASCII text editor. Newline characters and comments a commandfile are ignored. Lx51 interprets the first filename preceded by an at sign (@) as a commandfile.

The *inputlist* uses the following general format:

where

filename is the name of an object file created by Ax51, Cx51, or Intel

PL/M-51 or a library file created by the LIBx51 library manager. The filename must be specified with its file

extension. Object files use the extension .OBJ. Library files use

the extension .LIB.

modulename is the name of an object module in the library file. The

modulename may only be used after the name of a library file. The modulenames must be specified in parentheses after the filename. Multiple modulenames may be separated by commas.

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Command Line Examples

The following examples are proper command lines for the Lx51 linker/locator.

```
BL51 C:\MYDIR\PROG.OBJ TO C:\MYDIR\PROG.ABS
```

In this example, only the input file, C:\MYDIR\PROG.OBJ, is processed and the absolute object file generated is stored in the output file C:\MYDIR\PROG.ABS.

```
LX51 SAMPLE1.OBJ, SAMPLE2.OBJ, SAMPLE3.OBJ TO SAMPLE.ABS
```

In this example, the files **SAMPLE1.OBJ**, **SAMPLE2.OBJ**, and **SAMPLE3.OBJ** are linked and absolute object file that is generated is stored in the file **SAMPLE.ABS**.

```
L251 PROG1.OBJ, PROG2.OBJ, UTILITY.LIB
```

In this example, unresolved external symbols are resolved with the public symbols from the library file **UTILITY.LIB**. The modules required from the library are linked automatically. Modules from the library that are not referenced are not included in the generated absolute object file.

```
BL51 PROG1.OBJ, PROG2.OBJ, UTILITY.LIB (FPMUL, FPDIV)
```

In this example, unresolved external symbols are resolved with the public symbols from the library file UTILITY.LIB. The modules required from the library are linked automatically. In addition, the FPMUL and FPDIV modules are included whether they are needed or not. Other modules from the library that are not referenced are not included in the generated absolute object file.

```
LX51 @PROJECT.LIN
```

Content of the file PROJECT.LIN:

```
PROG1.OBJ, /* Program Module 1 */
PROG2.OBJ, // program module 2
UTILITY.LIB (FPMUL, FPDIV) ; include always FPMUL and FPDIV
```

This is example is the same as the example before, but uses a command input file that includes comments.

Control Linker Input with µVision2

The C and assembler source files that are part of a μ Vision2 project are translated when you build your application. The object files generated are then supplied as linker input file by the μ Vision2 build process. However you may also include object and library files as part of a μ Vision2 project in the same way as you include source files. You may set additional linker options for a file or file group using the **Options** – **Properties** dialog. For detailed information refer to the *Getting Started and Creating Applications User's Guide*.

ERRORLEVEL

After linking, the Lx51 linker/locator sets the ERRORLEVEL to indicate the status of the linking process. The Lx51 linker/locater and the other utilities generate the same ERRORLEVEL values as the Ax51 macro assembler. Refer to "ERRORLEVEL" on page 197 for more information.

Output File

The **Lx51** linker/locator creates an output file using the input object files that you specify on the command line. The output file is an absolute object file that may be loaded into debugging tools like the μ Vision2 Debugger or may be converted into a Intel HEX for PROM programming.

Linker/Locater Controls

Controls for the **Lx51** linker/locater may be entered after the output file specification. Multiple controls must be separated by at least one space character (). Each control may be entered only once on the command line. If a control is entered twice, the **Lx51** linker/locator reports an error.

The following table lists all Lx51 linker/locator controls and a brief description. The controls of the BL51 linker/locater are listed in the first table. The controls of the extended LX51 linker/locater and L251 linker/locater are listed in the second table. LX51 and L251 provide the same sets of controls.

The "Control Summary" on page 305 explains the command-line controls in detail. Refer to page number provided in the tables for quick reference to descriptions and examples for each control.

NOTE

Underlined characters denote the abbreviation for the particular control.

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BL51 Controls

Controls F	age	Description
<u>BA</u> NKAREA	332	Specifies the address range where the code banks are located.
<u>B</u> ANK <u>x</u>	333	Specifies the start address and segments for code banks.
<u>BI</u> T	334	Locates and orders BIT segments.
<u>CO</u> DE	338	Locates and orders CODE segments.
<u>DA</u> TA	339	Locates and orders DATA segments.
<u>D</u> ISABLE <u>W</u> ARNING	307	Disables specified warning messages.
<u>IB</u> ANKING	323	Generate bank switch code for Infineon TV TEXT devices.
<u>ID</u> ATA	340	Locates and orders IDATA segments.
<u>IX</u> REF	308	Includes a cross reference report in the listing file.
<u>NA</u> ME	323	Specifies a module name for the object file.
<u>NOAJ</u> MP	325	Generate bank switch code without AJMP instructions.
NODEBUGLINES	325	Excludes line number information from the object file.
NODEBUGPUBLICS	325	Excludes public symbol information from the object file.
NODEBUGSYMBOLS	325	Excludes local symbol information from the object file.
NODEFAULTLIBRARY	352	Excludes modules from the run-time libraries.
NOINDIRECTCALL	327	Do not generate bank switch code for indirectly called functions.
NOJMP <u>T</u> AB	328	Do not generate bank switch code.
NOLINES	310	Excludes line number information from the listing file.
<u>NOMA</u> P	311	Excludes memory map information from the listing file.
NOOVERLAY	353	Prevents overlaying or overlapping local bit and data segments.
<u>NOPR</u> INT	315	Disables generation of a listing file.
NOPUBLICS	312	Excludes public symbol information from the listing file.
NOSORTSIZE	341	Disable size sorting for segments before allocating the memory.
NOSYMBOLS	313	Excludes local symbol information from the listing file.
<u>O</u> VER <u>L</u> AY	354	Modifies call tree for data overlaying of local data & bit segments.
PAGELENGTH	314	Sets maximum number of lines in each page of listing file.
PAGEWIDTH	314	Sets maximum number of characters in each line of listing file.
PDATA	341	Specifies the starting address for PDATA segments.
PRECEDE	343	Locates segments that precede others in the DATA memory.
<u>PR</u> INT	315	Specifies the name of the listing file.
RAMSIZE	344	Specifies the size of the on-chip data memory.
<u>R</u> ECURSIONS	356	Allows analyze of the call tree of complex recursive applications.
REGFILE	356	Specifies the register usage information file generated by Lx51.
RTX51	358	Includes support for the RTX-51 full real-time kernel.
RTX51TINY	358	Includes support for the RTX-51 tiny real-time kernel.
<u>SP</u> EEDOVL	359	Ignore during overlay analysis references from constant segments.
<u>ST</u> ACK	349	Locates and orders STACK segments.
<u>XD</u> ATA	350	Locates and orders XDATA segments.

LX51 and L251 Controls

Controls	Page	Description
<u>AS</u> SIGN	322	Defines public symbols on the command line.
<u>BA</u> NKAREA	332	Specifies the address range where the code banks are located.
<u>CL</u> ASSES	336	Specifies a physical address range for segments in a memory class.
<u>D</u> ISABLE <u>W</u> ARNING	307	Disables specified warning messages.
<u>IX</u> REF	308	Includes a cross reference report in the listing file.
<u>NA</u> ME	323	Specifies a module name for the object file.
<u>NOAJ</u> MP	325	Generate bank switch code without AJMP instructions.
NOCOMMENTS	309	Excludes comment information from listing file and the object file.
NODEFAULT <u>LIB</u> RAR	Y 352	Excludes modules from the run-time libraries.
NOINDIRECTCALL	327	Do not generate bank switch code for indirectly called functions.
<u>NOLI</u> NES	310	Excludes line number information from listing file and object file.
<u>NOMA</u> P	311	Excludes memory map information from the listing file.
NOOVERLAY	353	Prevents overlaying or overlapping local bit and data segments.
<u>NOPR</u> INT	315	Disables generation of a listing file.
NOPUBLICS	312	Excludes public symbol information from the listing and object file.
NOSYMBOLS	313	Excludes local symbol information from the listing file.
NOSORTSIZE	341	Disable size sorting for segments before allocating the memory.
<u>NOTY</u> PE	327	Excludes type information from the listing file and the object file.
OBJECTCONTROLS	330	Excludes specific debugging information from the object file.
<u>O</u> VER <u>L</u> AY	354	Modifies call tree for data overlaying of local data & bit segments.
PAGELENGTH	314	Sets maximum number of lines in each page of listing file.
PAGEWIDTH	314	Sets maximum number of characters in each line of listing file.
<u>PR</u> INT	315	Specifies the name of the listing file.
PRINTCONTROLS	316	Excludes specific debugging information from the listing file.
<u>PU</u> RGE	317	Excludes all debugging information from the listing and object file.
RECURSIONS	356	Allows analyze the call tree of complex recursive applications.
REGFILE	356	Specifies the register usage information file generated by Lx51.
<u>RE</u> SERVE	345	Reserves memory and prevent Lx51 from using memory areas.
RTX251	358	Includes support for the RTX-251 full real-time kernel.
RTX51	358	Includes support for the RTX-51 full real-time kernel.
RTX51TINY	358	Includes support for the RTX-51 tiny real-time kernel.
<u>SE</u> GMENTS	346	Defines physical memory addresses and orders for segments.
<u>S</u> EG <u>S</u> IZE	348	Specifies memory space used by a segment.
<u>W</u> ARNING <u>L</u> EVEL	318	Controls the types and severity of warnings generated.

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Locating Programs to Physical Memory

This section describes with examples how you locate your application into the physical memory space for the different *x*51 variants. Refer to "Segment and Memory Location Controls" on page 331 for a detailed description of the linker/locater controls used in the examples below.

The linker/locator determines the physical memory range for relocatable segments based on the memory class that is assigned to the segment. Refer to "Memory Classes and Memory Layout" on page 27 for more information. However, it is also possible to specify a fixed address for a segment using linker/locater controls.

Classic 8051

The classic 8051 provides three different memory areas: on-chip RAM that covers the DATA, BIT and IDATA memory, XDATA memory, and CODE memory. The "Classic 8051 Memory Layout" is shown on page 29.

Classic 8051 without Code Banking

The following examples illustrate how to setup the linker/locater. For the BL51 linker/locater the physical memory is defined with the **RAMSIZE**, **XDATA** and **CODE** control. For the **LX51** linker/locater the **CLASSES** control is used to specify the available physical memory.

The following example assumes the following memory areas.

Memory Type	Address Range	Used by
ON-CHIP RAM	D:0 - D:0x7F (128 Bytes)	registers, bits, variables, etc.
XDATA RAM	X:0 - X:0x7FFF, X:0xF800 - X:0xFFFF	space for variables.
CODE ROM	C:0 - C:0x7FFF	program code and constant area.

To specify this memory layout BL51 should be invoked with as follows:

```
BL51 PROG.OBJ XDATA (0-0x7FFF, 0xF800-0xFFFF) CODE (0-0x7FFF) RAMSIZE(128)
```

You may also you the LX51 linker/locater. The **CLASSES** directive should have the following settings:

NOTE

You need not to define the address range for the memory classes DATA and BIT since the LX51 default setting already covers the correct physical address range.

Classic 8051 with Code Banking

The following example uses classic 8051 with a code banking hardware. This hardware has the following memory resources:

Memory Type	Address Range	Used by
ON-CHIP RAM	I:0 - I:0xFF (256 Bytes)	registers, bits, variables, etc.
XDATA RAM	X:0 - X:0xEFFF	space for variables.
CODE ROM	C:0 - C:0x7FFF (common area) B0:0x8000 - B3:0xFFFF (four banks)	program code and constant area.

Parts of your program will be located into banks using **BANK***x* in the *inputlist* portion of the **L***x***51** linker/locater command-line. Refer to "Bank Switching" on page 293 for more information. In addition you must specify the size of the common area and the other memory resources of your hardware. For this memory layout, the BL51 linker/locater should be invoked with as follows:

```
BL51 BANKO {A.OBJ}, BANK1 {B.OBJ}, BANK2 {C.OBJ}, BANK3 {D.OBJ}
XDATA (0-0xEFFF) BANKAREA (0x8000 - 0xFFFF) RAMSIZE(256)
```

The LX51 linker/locater needs to be invoked as follows:

```
LX51 BANKO {A.OBJ}, BANK1 {B.OBJ}, BANK2 {C.OBJ}, BANK3 {D.OBJ}

CLASSES (IDATA (I:0-I:0xff), XDATA (X:0-X:0xEFFF),

CODE (C:0-C:0xffff)) BANKAREA (0x8000-0xffff)
```

Extended 8051 Variants

Some extended 8051 variants expand the external data and program memory to up to 16MB. The additional memory space is addressed with the memory classes HDATA and HCONST. The "Extended 8051 Memory Layout" is shown on page 31. Only the LX51 linker/locater supports this expanded memory space. The following example shows assumes the following memory areas.

Memory Type	Address Range	Used by
ON-CHIP RAM	D:0 – D:0xFF (256 Bytes)	registers, bits, variables.
XDATA RAM	X:0 – X:0x1FFFF (128 KB)	space for variables.
CODE ROM	C:0 - C:0xFFFFF (1 MB)	program code and constant area.

To specify this memory layout LX51 should be invoked with the following CLASSES directive

```
LX51 MYPROG.OBJ CLASSES (HDATA (X:0 - X:0x1FFFF),
HCONST (C:0 - C:0xFFFFF))
```

NOTE

You need not to define the address range for the memory classes DATA, IDATA, BIT, CODE, CONST, and XDATA since the LX51 default already covers the correct physical address ranges for these memory classes.

The memory classes HDATA and HCONST are used for constants or variables only. Program code is located into the expanded program memory with the same code banking mechanism as described above under "Classic 8051 with Code Banking". A command line example that locates also program code into the expanded program memory will look as follows:

```
LX51 BANKO {A.OBJ}, BANK1 {B.OBJ}, BANK2 {C.OBJ}, BANK3 {D.OBJ}

CLASSES (IDATA (I:0-I:0xff), XDATA (X:0-X:0xEFFF),

HDATA (X:0-X:0x1FFFF), HCONST (C:0-C:0xFFFFF)

CODE (C:0-C:0xFFFF)) BANKAREA (0x8000-0xFFFF)
```

There are several *Keil Application Notes* available that show how to create programs for extended 8051 devices. Check www.keil.com or the Keil development tools CD-ROM for *Keil Application Notes* that explain how to setup the tools for extended 8051 devices.

Philips 80C51MX

The Philips 80C51MX has a linear 16MB address space that includes the standard 8051 memory areas DATA/IDATA, CODE, and XDATA. In addition both the external data space and the program space can be up to 8 MB. The "80C51MX Memory Layout" is shown on page 33. The LX51 linker/locater is used for the Philips 80C51MX microcontroller family. The following example shows assumes the following memory areas.

Memory Type	Address Range	Used by
ON-CHIP RAM	7F:0000H 7F:03FFH	registers, bits, variables.
RAM	00:0000H 01:FFFFH	EDATA space for variables.
ROM	80:0000H 83:FFFFH	program code and constant area.

To specify this memory layout LX51 should be invoked with the following CLASSES directive.

```
LX51 MYPROG.OBJ CLASSES (HDATA (0 - 0x1FFFF),

EDATA (0x7F0000 - 0x7F03FF),

ECODE (0x800000 - 0x83FFFF),

HCONST (0x800000 - 0x83FFFF))
```

NOTE

You need not to define the address range for the memory classes DATA, IDATA, BIT, CODE, CONST, and XDATA since the LX51 default already covers the correct physical address ranges for these memory classes.

In the AX51 assembler it is possible to use the ECODE class and therefore the complete 8MB code address space for program code. However, the CX51 compiler uses code banks to allocate parts of your program into the extended program memory. Therefore you must use same technique as described above under "Classic 8051 with Code Banking" to locate parts of your program into the ECODE space. A command line example will look as follows:

```
LX51 BANKO {A.OBJ}, BANK1 {B.OBJ}, BANK2 {C.OBJ}, BANK3 {D.OBJ}

CLASSES (HDATA (0 - 0x1FFFF), EDATA (0x7F0000 - 0x7F03FF),

ECODE (0x800000 - 0x83FFFF), HCONST (0x800000 - 0x83FFFF))
```

Intel/Atmel WM 251

The Intel/Atmel WM 251 has like the Philips 80C51MX a linear 16MB address space that includes all the memory classes. The "251 Memory Layout" is shown on page 35. The following examples show you the invocation of the L251 linker/locater that is used for the Intel/Atmel WM 251 microcontroller family.

Example 1: The following example assumes the following memory areas.

Memory Type	Address Range	Used by
ON-CHIP RAM	00:0000H 00:041FH	registers, bits, variables.
RAM	00:8000H 00:FFFFH	EDATA space for variables.
ROM	FF:0000H FF:7FFFH	program code and constant area.

To specify this memory layout L251 should be invoked with the following CLASSES directive.

```
L251 MYPROG.OBJ CLASSES (EDATA (0 - 0x41F, 0x8000 - 0xFFFF),

CODE (0xFF0000 - 0xFF7FFF),

CONST (0xFF0000 - 0xFF7FFF))
```

NOTES

You need not to define the address range for the memory classes DATA, IDATA, BIT and EBIT since the L251 default already covers the correct physical address ranges for these memory classes.

This example assumes that the memory classes XDATA, HDATA, HCONST, HCODE, and NCONST are not used in your application.

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Example 2: In addition to the example above, the next system contains a third RAM for the memory class XDATA. In addition the ROM space is increased.

Memory Type	Address Range	Used by
ON-CHIP RAM	00:0000H 00:041FH	registers, bits, variables.
ROM	00:0420H 00:7FFFH	NCONST space.
RAM	00:8000H 01:7FFFH	EDATA/HDATA space for variables.
ROM	FE:0000H FF:FFFFH	program code and constant area.

To specify this memory layout L251 should be invoked with the following CLASSES directive.

```
L251 MYPROG.OBJ CLASSES (EDATA (0 - 0x41f, 0x8000 - 0xffff),

NCONST (0x420 - 0x7fff),

HDATA (0x8000-0x1ffff),

HCONST (0xfe0000 - 0xffffff),

ECODE (0xfe0000 - 0xffffff))
```

NOTE

You need not to define the address range for the memory classes DATA, IDATA, BIT, EBIT, CODE, CONST, and XDATA since the L251 default already covers the correct physical address ranges for these memory classes.

Data Overlaying

Because of the limited amount of stack space available on the *x*51, local variables and function arguments of C and PL/M-51 routines are stored at fixed memory locations rather than on the stack. Normally, the L*x*51 linker/locator analyses the program structure of your application, creates a call tree and overlays the data segments that contain local variables and function arguments.

This technique usually works very well and provides the same efficient use of memory than a conventional stack frame would. Therefore this technique is also known as *compiled-time* stack since the stack layout is fixed during the Compiler and Linker run. However, in certain situations, this can be undesirable. You may use the **NOOVERLAY** control to disable overlay analysis and data overlaying. Refer to "NOOVERLAY" on page 353 for more information about this control.

By default, the Lx51 linker/locator overlays the memory areas for local variables and function arguments with the same memory areas of other functions, provided that the functions do no call each other. Therefore the Lx51 linker/locator analyses the program structure and creates a function call tree.

The local data and bit segments that belong to a function are determined by their segment names. The local data and bit segments of a function are overlaid with other function's data and bit segments under the following conditions:

- No call references may exist between the functions. The Lx51 linker/locator considers direct calls between functions, as well as calls via other functions.
- The functions may be invoked by only one program event: main or interrupt. It is not possible to overlay data areas if a function is called by an interrupt and during the main program. The same is true when the function is called by several interrupts that might be nested.
- The segment definitions must conform to the rules described below.

NOTE

Typically, the **Lx51** linker/locator generates overlay information that is accurate. However, in some instances the default analysis of the call tree is ineffective or incorrect. This occurs with indirectly called functions through function pointers and functions that are called by both the main program and an interrupt function.

Program and Data Segments of Functions

For correct data overlaying the **Lx51** linker/locater must know the function code and the local variable space that belongs to this function. The program and data segments that belong together are determined by standard segment naming convention used by the **Cx51** compiler and PL/M-51 compiler. Therefore, segments used in assembler programs should be constructed according to the following rules.

Segment Content	Cx51 Segment Name	PL/M-51 Segment Name
Program CODE	?PR?functionname?modulename	?modulename?PR
Local BIT space	?BI?functionname?modulename	?modulename?BI
Local DATA space	?DT?functionname?modulename	?modulename?DT
Local IDATA space	?ID?functionname?modulename	_
Local XDATA space	?XD?functionname?modulename	_
Local PDATA space	?PD?functionname?modulename	_
Local EDATA space	?ED?functionname?modulename	_
Local EBIT space	?EB?functionname?modulename	_
Local HDATA space	?HD?functionname?modulename	_

?PR?, ?BI?, ?DT?, ?XD?, ?ID?, ?PD?, ?EB?, and ?HD? is derived from the memory class. In addition each bit and data segment must have the relocation type OVERLAYABLE.

The **Cx51** compiler and PL/M-51 compiler define automatically local data segments according to these rules. However, if you use overlayable segments in your assembly modules, you must follow these naming conventions. Refer to "SEGMENT" on page 106 for information on how to declare segments.

Example for Segment declaration in assembly language:

More information can be found in the *Keil Application Note 149: Data Overlaying and Code Banking with Assembler Modules* that is available on www.keil.com or the Keil development tools CD-ROM.

Using the Overlay Control

In most cases, the **Lx51** data overlay algorithm works correct and does not require any adjustments. However, in some instances the overlay algorithm cannot determine the *real* structure of your program and you must adjust the function call tree with the **OVERLAY** control. This is the case when your program uses function pointers or contains virtual program jumps as it is the case in the scheduler of a real-time operating system.

NOTE

The **Lx51** linker/locater recognizes correctly the program structure and the call tree of applications that are using the RTXx51 real-time operating system. You need not to use the OVERLAY control to specify the task functions of the RTXx51 application, since this is automatically performed by the **Lx51** linker/locater.

The **OVERLAY** control allows you to change the call references used by the **Lx51** linker/locator during the overlay analysis. Using the **OVERLAY** control is easy when you know the structure of your program. The program structure or call tree is reflected in the segments listed in the OVERLAY MAP of the listing file created by the **Lx51** linker/locater.

The following application examples show situations where the **OVERLAY** control is required to correct the call tree. In general, a modification of the references (calls) is required in the following cases:

- When a pointer to a function is passed or returned as function argument.
- When a pointer to a function is contained in initialized variables.
- When your program includes a real-time operating system.

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Disable Data Overlaying

If you are in doubt about whether certain segments should be overlaid or not, you may disable overlaying of those segments. Segment overlaying can be disabled at Cx51 compiler level or with the OVERLAY control at the Lx51 linker/locator command line as follows:

- You can invoke the Lx51 linker/locator with the NOOVERLAY option to disable data overlaying for the entire application.
- The Lx51 linker/locator control **OVERLAY** (*sfname*!*) disables data overlaying for the function specified by *sfname*.
- C code that is translated with the Cx51 compiler directive OPTIMIZE (1) does not use the relocation type OVERLAYABLE. Therefore the local data segments of this code portions cannot be overlaid.

Pointer to a Function as Function Argument

In the following example indirectfunc1 and indirectfunc2 are indirectly called through a function pointer in execute. The value of the function pointer is passed in the function main. Since main contains the reference, the Lx51 linker/locator thinks that main calls indirectfunc1 and indirectfunc2. But this is incorrect, since the real function call is in the function execute.

Following is a program listing for this example.

```
:
bit indirectfunc1 (void) {
                             /* indirect function 1 */
  unsigned char n1, n2;
  return (n1 < n2);
bit indirectfunc2 (void) { /* indirect function 2 */
 unsigned char a1, a2;
 return ((a1 - 0x41) < (a2 - 0x41));
void execute (bit (*fct) ()) {     /* sort routine */
  unsigned char i;
  for (i = 0; i < 10; i++) {
   if (fct ()) i = 10;
void main (void) {
                               /* switch: defines function */
  if (SWITCH)
    execute (indirectfunc1);
    execute (indirectfunc2);
}
```

The following listing file shows the overlay map for the program before making adjustments with the **OVERLAY** control.

OVERLAY MAP OF MODULE: OVL1	(OVL1)			
SEGMENT +> CALLING SEGMENT	BIT-0	GROUP LENGTH		-GROUP LENGTH
?C_C51STARTUP +> ?PR?MAIN?OVL1				
<pre>?PR?MAIN?OVL1 +> ?PR?INDIRECTFUNC1?OVL1 +> ?PR?EXECUTE?OVL1 +> ?PR?INDIRECTFUNC2?OVL1</pre>				
?PR?INDIRECTFUNC1?OVL1			0008н	0002Н
?PR?EXECUTE?OVL1			0008н	0004H
?PR?INDIRECTFUNC2?OVL1			18000	0002Н

The entry for ?PR?MAIN?OVL1 shows a call to ?PR?INDIRECTFUNC1?OVL1, ?PR?EXECUTE?OVL1, and ?PR?INDIRECTFUNC2?OVL1. However, only the function execute is called from main. The other references are results from using the function pointer fct, which is passed to execute. The function call to indirectfunc1 and indirectfunc2 takes place in execute, not in main where the functions are referenced

In this situation, the linker/locator cannot locate the actual function calls. Therefore, the Lx51 linker/locator incorrectly overlays the local segments of the functions execute, indirectfunc1, and indirectfunc2. This might result in a data overwrites of the variables i and fct.

You can use the **OVERLAY** control to correct the function call tree as it seen by the linker. For this example, you must remove the references from main to indirectfunc1 and indirectfunc2. Do this with main ~ (indirectfunc1, indirectfunc2). Then, add the actual function call from execute to indirectfunc1 and indirectfunc2 with executed! (indirectfunc1, indirectfunc2). The following shows the complete linker invocation line for this example.

```
Lx51 OVL1.OBJ OVERLAY (main ~ (indirectfunc1, indirectfunc2), execute ! (indirectfunc1, indirectfunc2))
```

With this Lx51 invocation the overlay map shows the correct references.

OVERLAY MAP OF MODULE: OV	7L1 (OVL1)		
SEGMENT +> CALLING SEGMENT		GROUP LENGTH	DATA-GROUP START LENGTH	
?C_C51STARTUP +> ?PR?MAIN?OVL1				
?PR?MAIN?OVL1 +> ?PR?EXECUTE?OVL1				
<pre>?PR?EXECUTE?OVL1 +> ?PR?INDIRECTFUNC1?OV +> ?PR?INDIRECTFUNC2?OV</pre>			0008н 0004н	
?PR?INDIRECTFUNC1?OVL1			000CH 0002H	
?PR?INDIRECTFUNC2?OVL1			000CH 0002H	

Pointer to a Function in Arrays or Tables

Another typical scenario is an array that contains a pointer to a function. This is typical for applications with function tables. In the following example, func1 and func2 are called indirectly by main but the entry points are stored as constant values in the table functab. This table is located in the segment ?CO?modulname. Therefore, the ?CO?OVL2 segment contains references to func1 and func2.

In reality, however, the calls are executed from the main function. But, the Lx51 linker/locator assumes that func1 and func2 are recursive called, because in func1 and func2 constant strings are used. These constants strings are also stored in the segment ?CO?OVL2. The result is that the Lx51 linker/locator reports warnings which indicate recursive calls from the segment ?CO?OVL2 to func1 and func2.

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The following listing shows part of the OVL2 program.

The Lx51 linker/locater generates typically warnings when you generate programs that contain a table with pointer to functions. Although the program can be executed correct in this example above, the references should be adjusted to the real calls. In the real application the functions func1 and func2 are called by the main function.

If you are using the BL51 linker/locater in the default configuration you need to delete the references from the code segment that contains the tables with ?CO?OVL2 ~ (func1, func2). Then you need to add the calls from main to func1 and func2 with main! (func1, func2):

```
BL51 OVL2.OBJ OVERLAY (?CO?OVL2~(func1, func2), main!(func1, func2))
```

The SPEEDOVL control of **BL51** ignores all references from constant segments to program code. This is also the operation mode of LX51 and L251. Therefore you need only to add the calls from main to func1 and func2 with main! (func1, func2):

```
BL51 OVL2.OBJ OVERLAY (main ! (func1, func2)) (SPEEDOVL)
```

The LX51 and L251 linker/locater always ignores the references from constant segments to program code and requires only to add the function calls:

```
LX51 OVL2.OBJ OVERLAY (main ! (func1, func2))
```

After this correction the memory usage of your application is typically more efficient and the overlay map shows a call tree that matches your application. Also the linker/locater does not generate any warning messages.

OVERLAY MAP OF MODULE: OVL2	(OVL2)		
SEGMENT +> CALLING SEGMENT	BIT-GI START I	DATA-0	
C_C51STARTUP +> ?PR?MAIN?OVL2		 	
<pre>?PR?MAIN?OVL2 +> ?PR?FUNC1?OVL2 +> ?PR?FUNC2?OVL2</pre>		 	
?PR?FUNC1?OVL2 +> ?PR?PRINTF?PRINTF		 н8000	0001H
?PR?PRINTF?PRINTF +> ?PR?PUTCHAR?PUTCHAR		 0009Н	0014H
?PR?FUNC2?OVL2 +> ?PR?PRINTF?PRINTF		 18000	0001H

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Tips and Tricks for Program Locating

The **Lx51** linker/locator supports several techniques that are required in for special tasks, for example in-system Flash programming or systems that use a **RAM** section for constants. The following section provides examples that show the usage of the **Lx51** linker/locater in such situations.

Locate Segments with Wildcards

The **Lx51** linker/locater allows in the segment controls wildcards for specifying the segment name. For example you may use such segment name specifications to locate all segments within one module into one 2KB block. In this way you can use the ACALL and AJMP instructions for function calls within this module.

```
BL51 myfile.obj CODE (?PR?*?myfile (0x1000))

LX51 myfile.obj SEGMENTS (?PR?*?myfile (C:0x1000))
```

Special ROM Handling (LX51 & L251 only)

The LX51 and L251 linker/locater provide the memory class SROM that is used to handle segments or memory classes that are to be stored in ROM, but copied for execution into RAM areas. This is useful for:

- In-system Flash programming when the Flash ROM contains also the flash programming code. With standard Flash devices it is impossible to fetch program code from the while other parts of the device are erased or programmed. The *Keil Application Note 139: "In-system Flash Programming with 8051 and 251"* that is available on www.keil.com or the Keil development tools CD-ROM contains a program example.
- For using the C251 TINY or XTINY memory model it is required to provide a NCONST memory class in the lowest 64KB memory region. However, if only RAM is mapped into this memory region, you can specify a different storage address for the NCONST memory class and copy the content at the program start into RAM. This allows you to use the efficient TINY or XTINY memory model while the system hardware just provides RAM in the lowest 64KB memory segment. Refer to "Use RAM for the 251 Memory Class NCONST" on page 292 for a program example.

Refer to **Lx51** linker/locater controls "SEGMENTS" on page 346 and "CLASSES" on page 336 for syntax on defining segments and memory classes that have a different storage and execution address.

Segment and Class Information (LX51 & L251 only)

The Lx51 linker/locater creates special symbols that can be used to obtain address and length information for segments or classes used in an application. The information is passed via external variable declarations. The Lx51 linker/locater uses symbols with the notation segmentname_p_ or _classname_p_. Question mark (?) characters in the segment name generated by the Cx51 compiler are replaced with underscore (_) characters. The postfix _p_ specifies the information that should be obtained and is explained in the following table:

Postfix	Description
l	is the length in bytes of the segment or memory class. For a memory class this number includes also any gaps that are required to allocate all segments that belong to this memory class.
s	is the start address of the segment or memory class. For a memory class this number refers to the first segment that belongs to this memory class.
e	is the end address of the segment or memory class. For a memory class this number refers to the last segment that belongs to this memory class.
t	is the target or execution address of the segment or memory class. For a memory class this number refers to the first segment that belongs to this memory class. This information is only available for segments or memory classes which have assigned a different storage and execution address.

Examples:

If ?PR?FUNC1 is the segment name:

```
_PR_FUNC1_L is the length in bytes of the segment ?PR?FUNC1.

_PR_FUNC1_S is the start address of the segment ?PR?FUNC1.

_PR_FUNC1_E is the end address of the segment ?PR?FUNC1.

_PR_FUNC1_T is the target or execution address of the segment ?PR?FUNC1.
```

If **nconst** is the memory class name:

_NCONST_L_	is the length in bytes of the memory class NCONST .
_NCONST_S_	is the start address of the memory class nconst .
_NCONST_E_	is the end address of the memory class NCONST .
_NCONST_T_	is the target or execution address of the memory class NCONST .

You may access this information in **Cx51** applications as shown in the following program example:

```
extern char _PR_FUNC1_L_;
unsigned int get_length (void) {
  return ((unsigned int) &_PR_FUNC1_L_); // length of segment ?PR?FUNC1
}
```

The file **SROM.H** contains macro definitions for accessing segment and class information. Refer to the *Keil Application Note 139: "In-system Flash Programming with 8051 and 251"* for more information.

Use RAM for the 251 Memory Class NCONST

The C251 compiler memory model **TINY** or **XTINY** requires a NCONST memory class in the lowest 64KB memory region. If your hardware provides only RAM in this memory area, you may use the SROM memory class to store the constants somewhere in the 16MB memory space and you may copy the content of the NCONST memory class into a RAM in the lowest 64KB memory. This is shown in the following program example. Refer to the "CLASSES" control on page 336 for more information.

```
#include <string.h>
extern char huge _NCONST_S_;
extern char huge _NCONST_T_;
extern char near _NCONST_L;

const char text [] = "This text is accessed in the NCONST memory class";

void main (void) {
   fmemcpy (&_NCONST_T_, &_NCONST_S_, (unsigned int)&_NCONST_L_);
   :
}
```

The C251 compiler and L251 linker/locater is invoke as follows:

```
C251 SAMPLE.C XTINY DEBUG
L251 SAMPLE.C CLASSES (NCONST (0x2000-0x4000)[]), SROM (0xFE0000-0xFEFFFF)
```

Bank Switching

The **Lx51** linker/locator manages and allows you to locate program code in up to 32 code banks and one common code area. The common code area is always available for all the code banks. The common code area and other aspects of the code banking are described below.

Common Code Area

The common code area can be accessed by all banks. This area usually includes routines and constant data that must always be accessible; for example, interrupt and reset vectors, interrupt routines, string constants, bank switching routines, etc. The following code sections must always be located in the common area:

- Reset and Interrupt Vectors: reset and interrupt jump entries must remain in the common area, since the code bank selected by the x51 program is not known at the time of the CPU reset or interrupt. The Lx51 linker/locator, therefore, locates absolute code segments in the common area in each case.
- Code Constants: constant values (strings, tables, etc.) which are defined in the code area must be stored in the common area unless you guarantee that the code bank containing the constant data is selected at the time they are accessed by program code. You can relocate these segments in code banks by means of control statements.
- Interrupt Functions: generated using the Cx51 compiler must always be located in the common area. However, interrupt functions can call functions in other code banks. The Lx51 linker/locator produces a warning when an attempt is made to locate a Cx51 interrupt function in a code bank.
- Bank Switch Code: is required for switching the code banks as well as the associated jump table are located in the common area since these program sections are required by all banks. By default, the Lx51 linker/locator automatically locates these segments in the common area. Do not attempt to locate these program sections into bank areas.
- **Library Functions:** intrinsic run-time library functions used by the **Cx51** compiler or the PL/M-51 compiler must be located in the common area. It is possible that the bank switch code will use registers that are used to transfer values to such library functions. Therefore, the **Lx51** linker/locator always locates program sections of the runtime library in the common area. Do not locate these program sections in other bank areas.

It is difficult to estimate the size of the common area. The size will always depend on the particular software application and hardware constraints. If the ROM area that is dedicated as common area is not large enough to contain the entire common code, the **Lx51** linker/locator will duplicate the remaining part of the common code area into each code bank. This is also the case, if your hardware does not provide a common code area section in the ROM space.

Code Bank Areas

The classic 8051 only provides 16 address lines for accessing code memory. With 16 address lines, only 64 KBytes of code space can be accessed. Code banks are addressed using up to five additional address lines that must originate from 8051 I/O ports or from external hardware devices (latch or port I/O device) that are mapped into the XDATA space. A particular code bank is selected by controlling the state of the additional address lines. Up to 32 banks can be used.

Code banking applications must include the assembly file L51_BANK.A51 that is located in the folder LIB. This source module contains the code that is invoked to switch code banks. You must configure this source file to match the bank switching technique used by your target hardware. Refer to "Bank Switching Configuration" on page 297 for a description of this source file.

Optimum Program Structure with Bank Switching

The **Lx51** linker/locator automatically generates a jump table for all functions, which are stored in the bank area and are called from the common area or from other banks. The **Lx51** linker/locator only uses bank switching when the program section called actually lies in another memory bank or when it can be called from the common area. This improves performance and prevents bank switching from significantly impacting the performance of your application program. Additionally, the memory and stack requirements for this bank switching technique are considerably smaller than other alternative solutions.

Each bank switch takes on a classic 8051 approximately 50 CPU cycles and requires two additional bytes in the stack area. Bank switches are relatively fast, however, programs should be structured so that bank switches are seldom required to achieve maximum performance. This means that functions that are frequently invoked and functions that are called from multiple code banks should be located in the common code area.

Program Code in Bank and Common Areas

The Lx51 linker/locator provides the BANKAREA, BANKx, and COMMON controls to specify the location and size of the bank switching area and to locate segments in code banks or the common area.

When you generate a code banking application, you must specify the modules you want located in a code bank or common area. This is accomplished using **BANK***x* or **COMMON** in the *inputlist* portion of the **Lx51** linker/locator command line.

BANK*x* in the *inputlist* specifies the code bank for object and library files. The *x* in the **BANK***x* keyword specifies a bank number from 0 to 31. For example, **BANK**0 for code bank number 0, **BANK**1 for code bank number 1, and so on. All program code segments contained in these modules will be located in the specified code bank. A program code segment is determined by it prefix or postfix ?**PR**?.

COMMON locates program code into the common area. This is also the default for modules that are not explicitly located with **BANK***x*.

The general format for **BANK***x* and **COMMON** in the *inputlist* are:

```
BANKx { filename [(modulename)] [, filename ...]} [, ...]

COMMON { filename [(modulename)] [, filename ...]} [, ...]

where

x is the bank number to use and can be a number from 0 to 15.

{ and } are used to enclose object files or library files.

filename is the name of an object module in a library file.

modulename is the name of an object module in a library file.
```

The start and end address of the area where the code banks are located is specified with the **BANKAREA** control. These address range should reflect the space where the code bank ROMs are physically mapped. All program code segments that are assigned to a bank will with the **BANK** keyword in the <code>inputlist</code> will be located within this address range. Refer to "BANKAREA" on page 332 for more information.

Command-Line Example:

A typical Lx51 linker/locator command line appears as follows. More "Error! Reference source not found." can be found on page Error! Bookmark not defined.

```
LX51 COMMON{C_ROOT.OBJ},

BANK0{C_BANK0.OBJ, MODUL1.OBJ},

BANK1{C_BANK1.OBJ},

BANK2{C_BANK2.OBJ}

TO MYPROG.ABS &

BANKAREA(8000H,0FFFFH)
```

Segments in Bank Areas

In the *controls* portion of the **BL51** linker/locator command line you can use the **BANK***x* control to locate or order segments within a code bank. Refer to "BANK*x*" on page 333 for more information.

For the **LX51** linker/locator the **SEGMENTS** control allows you also to locate or order segments within a code banks. Refer to "SEGMENTS" on page 346 for more information.

With these controls you may place constants in code banks. You can use this technique to locate arrays or large tables in code banks other than the one in which your program resides. However, in your **C***x***51** programs, you must manually ensure that the proper code bank is used when accessing that constant data. You can do this with the **switchbank** function which is defined in the **L51_BANK.A51** module. "BANK_EX2 – Banking with Constants" on page 407 shows a complete example program.

Bank Switching Configuration

When you create a code banking application, you must specify the number of code banks your hardware provides as well as how the code banks are switched. This is done with definitions in an assembler source file. For the classic 8051 devices the bank switch configuration is defined in the file L51_BANK.A51 found in the \C51\LIB\ subdirectory. For the Philips 80C51MX the file MX51BANK.A51 is used.

NOTE

For extended 8051 devices there are several **Keil Application Notes** available on www.keil.com or the Keil development tools CD-ROM that explain the banking and memory configuration for these devices.

The following table explains the definitions in the Bank Configuration File:

Name	Description			
?B_NBANKS	number of banks to be supported. The following values are allowed: 2, 4, 8, 16, and 32. Two banks require one additional address (or I/O Port) line; four banks require two lines; eight banks require three lines; sixteen banks require four lines, and thirty-two banks require five address lines.			
?B_MODE	indicates the way how the address extension is done. 0 for using an standard 8051 I/O Port, 1 for using an XDATA port; 2 for using 80C51MX address line; and 4 for user provide bank switch code.			
?B_RTX	specifies if the application uses RTX-51 Full. Only ?B_MODE 0 and 1 are supported by RTX51 Full.			
?B_VARBANKING	Enables variable banking in XDATA and CODE memory. Variable banking requires the LX51 linker/locater. It is not supported by BL51 . Refer to "Banking With Common Area" on page 303 for an example on how to setup the LX51 linker/locater.			
?B_RST_BANK	Specifies the default bank that is selected after CPU reset. This setting is used by the LX51 linker/locater to reduce the entries in the INTERBANK CALL TABLE. The value 0xFF disables this optimization. This value is not used by BL51 .			
For ?B_MODE = 0 (b	ank switching via 8051 I/O Port) define the following:			
?B_PORT	used to specify the address of the internal data port. The SFR address of an internal data port must be specified. (For example: P1 as for port 1).			
?B_FIRSTBIT	indicates which bit of the 8051 I/O port is to be assigned first. The value 3 indicates that port bit 3 is used as first port line for the address extension. If, for example, two address lines are used, P1.3 and P1.4 are allocated in this case. The remaining lines of the 8051 I/O port can be used for other purposes.			

Name	Description			
For ?B_MODE = 1 (b	For ?B_MODE = 1 (bank switching via xdata mapped port) define the following:			
?B_XDATAPORT	specifies the XDATA memory address used to select the bank address and defines the address of an external data port. Any XDATA address can be specified (address range 0H to 0FFFFH) under which a port can be addressed in the XDATA area. 0FFFFH is defined as the default value. In this mode ?B_CURRENTBANK and ?B_XDATAPORT are initialized with the value 0 at the start of the program by the C51 startup code.			
?B_FIRSTBIT	indicates which bit of the defined port is to be assigned first. Other than with ?B_MODE=0 the remaining bits of the XDATA port cannot be used for other purposes.			
For ?B_MODE = 4 (b	ank switching via user provided code) define the following:			
SWITCHX	For each memory bank a own macro defines the bank switch code. The number of macros must conform with the ?B_NBANKS definition. For example 4 banks require a SWITCH0 , SWITCH1 , SWITCH2 and SWITCH3 macro. Each macro must generate exactly the same number of code bytes. If the size is different you use should NOP instructions to make it identical. You must also ensure that the CPU has selected a defined state at startup. The RTX51 real-time operating system does not support this banking mode.			

The **Ax51** assembler is required to assemble **L51_BANK.A51** or **MX51BANK.A51**. The source file should be copied as part of your project file. Public Symbols in L51_BANK.A51

Additional PUBLIC Symbols are provided in L51_BANK.A51 for your convenience. They are described below.

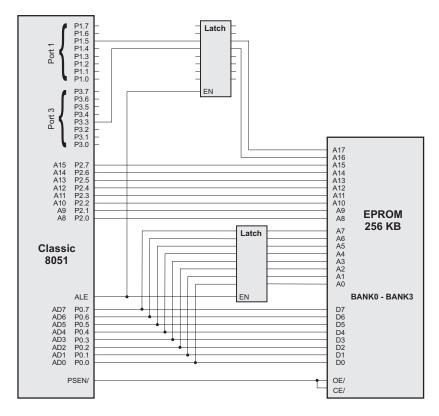
Name	Description
?B_CURRENTBANK	is a memory location in the DATA or SFR memory, which contains the currently selected memory bank. This memory location can be read for debugging. A modification of the memory location, however, does not cause a bank switching in most cases. Only required bits based on setting of ?B_NBANKS and ?B_FIRSTBIT are valid in this memory location. The bits, which are not required, must be masked out with a corresponding mask.
_SWITCHBANK	is a Cx51 compatible function, which allows the bank address to be selected by the user program. This function can be used for bank switching if the constant memory is too small. Note that this C function can be called only from code in the common area. The function is accessed as follows:
	extern void switchbank (unsigned char bank_number); : : switchbank (0);

Configuration Examples

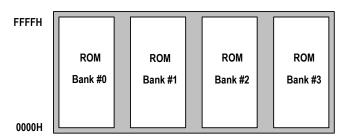
The following examples demonstrate how to configure code banking.

Banking With Four 64 KByte Banks

This example demonstrates the configuration required to bank switch using one 256KB EPROM. The following figure illustrates the hardware schematic.



The following figure illustrates the memory map for this example.



One 256KB EPROM is used in this hardware configuration. The bank switching can be implemented by using two bank select address lines (in this example Port 1.5 and Port 3.3). L51_BANK.A51 can be configured as follows for this hardware configuration.

```
?N_BANKS EQU 4 ; Four banks are required.
?B_MODE EQU 4 ; user-provided bank switch code is used.
```

The section that starts with IF ?B_MODE = 4 defines the code that switches between the code banks. This section needs to be configured as follows:

```
P1
              DATA
                      90H
                              ; I/O Port Addresses
Р3
              DATA
                      0B0H
                              ; Switch to Memory Bank #0
SWITCH0
              MACRO
                      P3.3 ; Clear Port 3 Bit 3
              CLR
              CLR
                      P1.5
                              ; Clear Port 1 Bit 5
              ENDM
SWITCH1
                              ; Switch to Memory Bank #1
              MACRO
                      P3.3
                              ; Set Port 3 Bit 3
              SETB
                      P1.5
                              ; Clear Port 1 Bit 5
              CLR
              ENDM
SWITCH2
                              ; Switch to Memory Bank #2
              MACRO
              CLR P3.3
                             ; Clear Port 3 Bit 3
              SETB P1.5
                             ; Set Port 1 Bit 5
              ENDM
SWITCH3
              MACRO
                              ; Switch to Memory Bank #3
                             ; Set Port 3 Bit 3
                     P3.3
              SETB
                      P1.5
              SETB
                              ; Set Port 1 Bit 5
              ENDM
```

You need to ensure that the CPU starts in a defined state at reset. Therefore the following code needs to be added to the STARTUP.A51 file of your application:

```
MOV
               SP,#?STACK-1
; added for bank switching
P1
       DATA
               90H
                                    ; I/O Port Addresses
ÞЗ
       DATA
               0B0H
       EXTRN DATA (?B CURRENTBANK)
            ?B CURRENTBANK, #0 ; select code bank 0
       CLR
              P3.3
                                   ; Clear Port 3 Bit 3
                                   ; Clear Port 1 Bit 5
       CLR
              P1.5
: end
               ?C START
```

The **Lx51** linker/locator automatically places copies of the code and data in the common area into each bank so that the contents of all EPROM banks are identical in the address range of the common area. The **BANKAREA** control is not required since the default setting already defines address range 0 to 0xFFFF as banked area.

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Banking With On-Chip Code ROM

Several device variants offer SFR registers that configure the on-chip code ROM space. You may use this feature for existing hardware designs to introduce code banking. For example, if your hardware uses currently a Dallas 80C320 (ROM-less device) and an external 64KB ROM, you may increase the code space of this existing hardware design with Dallas 80C520 design that offers 16KB on-chip ROM. You may use the **ROMSIZE** SFR register for code bank switching of the 16KB on-chip and off-chip ROM block.

The figure on the right shows this memory layout. For this configuration the following settings in **L51_BANK.A51** are required.

```
Common ROM Area (off-chip)

4000h

ROM Bank #0 (on-chip)

ROM Bank #1 (off-chip)
```

The macros need to be configured as follows:

```
0C2H
ROMSIZE
                DATA
                                       ; SFR Address
SWITCH0
                MACRO
                                       ; Switch to Memory Bank #0
                MOV
                        ROMSIZE, #05H ; Enable on-chip 16KB ROM
                ENDM
SWITCH1
                MACRO
                                       ; Switch to Memory Bank #1
                MOV
                        ROMSIZE, #00H ; Disable on-chip 16KB ROM
                ENDM
```

You need to ensure that the CPU starts in a defined state at reset. Therefore the following code needs to be added to the STARTUP.A51 file of your application:

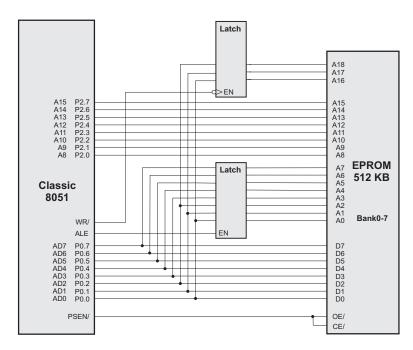
```
SP, #?STACK-1
        MOV
; added for bank switching
ROMSIZE
               DATA
                        0C2H
                                      ; SFR Address
        EXTRN DATA (?B CURRENTBANK)
             ?B CURRENTBANK,#0
        MOV
                                      ; select code bank 0
                ROMSIZE, #05H
        MOV
                                      ; start with on-chip ROM enabled
; end
        JMP
               ?C START
```

The Lx51 linker/locater BANKAREA control should be set as follows:

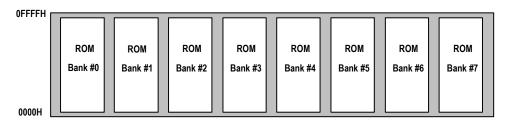
```
BL51 ... BANKAREA (0,0x3FFF)
```

Banking With XDATA Port

You may also use a latch or I/O device that is mapped into the XDATA space to extend the address lines of the 8051 device. The following application illustrates a hardware that uses a latch mapped into the XDATA space to address a 512KB EPROM.



The following figure illustrates the memory map for this example.



For this hardware the L51_BANK.A51 file can be configured as follows:

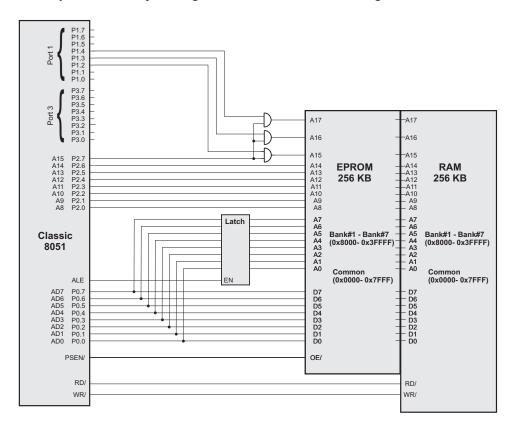
You need no additional configuration in the STARTUP.A51 file. The Lx51 linker/locator automatically places copies of the code and data in the common

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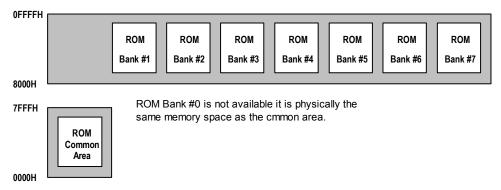
area into each bank so that the contents of all EPROM banks are identical in the address range of the common area. The **BANKAREA** control is not required since the default setting already defines address range 0 to 0xFFFF as banked area.

Banking With Common Area

The following schematic shows a hardware that offers a 32KB common area and seven 32KB code banks. A single EPROM is used to map the complete memory. Due to the address decoding logic, the code bank 0 is identical with the common area and should be therefore not used by your application. The design also provides 256KB xdata memory that is mapped the same way as the code memory. The xdata space might be used for variable banking.



The following figure illustrates the memory map for this example.



For this hardware the L51_BANK.A51 file can be configured as follows:

```
?N BANKS
                EQU
                      8
                               ; Eight banks are required.
?B MODE
                      0
                EQU
                               ; banking via on-chip I/O Port.
?B VAR BANKING EQU
                      1
                               ; you may use also variable banking.
?B PORT
                EQU
                      090H
                               ; Port address of P1.
?B FIRSTBIT
                EQU
                               ; Bit 2 is used as the first address line.
```

You should not use the code bank 0 in your application, since this memory is effectively identical with the common area. Therefore no module of you application should be assigned to code bank 0. The **Lx51** linker/locater **BANKAREA** control should be set as follows:

```
BL51 BANK1 {A.OBJ}, ..., BANK7{G.OBJ} ... BANKAREA (0x8000,0xffff)
```

If you are using variable banking, you need to use LX51 linker/locator. To define the additional memory the **HDATA** and **HCONST** memory classes are used. In this case the memory classes need to be set as follows:

```
LX51 BANK1 {A.OBJ}, ..., BANK7{G.OBJ} ... BANKAREA (0x8000,0xFFFF)

CLASSES (XDATA (X:0-X:0x7FFF),

HDATA (X:0x18000-X:0x1FFFF,X:0x28000-X:0x2FFFF,

X:0x38000-X:0x3FFFF,X:0x48000-X:0x4FFFF,

X:0x58000-X:0x5FFFF,X:0x68000-X:0x6FFFF,

X:0x78000-X:0x7FFFF),

HCONST (C:0x18000-C:0x1FFFF,C:0x28000-C:0x2FFFF,

C:0x38000-C:0x3FFFF,C:0x48000-C:0x4FFFF,

C:0x58000-C:0x7FFFF),

C:0x78000-C:0x7FFFF)
```

Control Summary

This section describes all **Lx51** linker/locator command-line controls. The controls are grouped into the following categories:

- Listing File Controls
- Output File Controls
- Segment and Location Controls
- High-Level Language Controls

Many of the Lx51 linker/locator controls allow you to specify optional arguments and parameters in parentheses immediately following the control. The following table lists the types of arguments that are allowed with certain controls.

Argument	Description		
address	A value representing a memory location. For BL51, L251 and LX51 in Philips 80C51MX mode plain numbers are used to represent an address. LX51 uses for classic 8051 devices a memory prefix in the address specification. For example: D:0x55 refers to DATA memory address 0x55 C:0x8000 refers to CODE memory address 0x8000 B4:0x4000 refers to CODE memory address 0x4000 in code bank 4.		
classname	A name of a memory class. The x51 tools allows basic classes and user defined classes. Refer to "Memory Classes and Memory Layout" the page 27 for more information about memory classes.		
filename	A file name that corresponds to the Windows file name conventions.		
modname	A module name. Can be up to 40 characters long and must start with: A – Z, ?, _, or @; following characters can be: 0 – 9, A – Z, ?, _, or @.		
range	An address range in the format: startaddress — endaddress The startaddress is the first address specified by the range. The endaddress is optional and specifies the last address which is included in the address range.		
segname	A segment name. Can be up to 40 characters long and must start with: A – Z, ?, _, or @; following characters can be: 0 – 9, A – Z, ?, _, or @.		
sfname	A segment or function name.		
value	A number, for example, 1011B, 2048D, 0x1000, or 0D5FFh.		

Listing File Controls

The Lx51 linker/locator generates a listing file that contains information about the link/locate process. This file is sometimes referred to as a map file. The following controls specify the filename, format, and information that is included in the listing file. For a detailed description of each control refer to the page listed in the table.

BL51	LX51, L251	Page	Description
<u>D</u> ISABLE <u>W</u> ARNING	<u>D</u> ISABLE <u>W</u> ARNING	307	Disables specified warning messages.
<u>IX</u> REF	<u>IX</u> REF	308	Includes a cross reference report.
-	NOCOMMENTS	309	Excludes comment information.
<u>NOLI</u> NES	<u>NOLI</u> NES	310	Excludes line number information.
<u>NOMA</u> P	<u>NOMA</u> P	311	Excludes memory map information.
<u>NOPR</u> INT	<u>NOPR</u> INT	315	Disables generation of a listing file.
NOPUBLICS	NOPUBLICS	312	Excludes public symbol information.
NOSYMBOLS	NOSYMBOLS	313	Excludes local symbol information.
PAGELENGTH(n)	PAGELENGTH(n)	314	Sets number of lines in each page.
PAGEWIDTH(n)	PAGEWIDTH(n)	314	Sets number of characters in each line.
<u>PR</u> INT	<u>PR</u> INT	315	Specifies the name of the listing file.
-	PRINTCONTROLS	316	Excludes specific debugging information.
_	<u>PU</u> RGE	317	Excludes all debugging information.
-	<u>W</u> ARNING <u>L</u> EVEL(<i>n</i>)	318	Controls the types and severity of warnings generated.

DISABLEWARNING

Abbreviation: DW

Arguments: DISABLEWARNING (number, | ... |)

Default: All warning messages are displayed.

μVision2 Control: Options – Lx51 Misc – Warnings – Disable Warning

Numbers.

Description: The **DISABLEWARNING** control lets you to selectively

disable Linker warnings. The warning numbers that should

be suppressed are specified in parenthesis.

The following examples disables the report of Warning

Number 1 and 5.

Example: LX51 myfile.obj DISABLEWARNING (1, 5)

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IXREF

Abbreviation: IX

Arguments: IXREF [(NOGENERATED, NOLIBRARIES)]

Default: No cross reference is generated in the listing file.

μVision2 Control: Options – Listing – Linker Listing – Cross Reference

Description: The **IXREF** control instructs the **Lx51** linker/locator to

include a cross reference report in the listing file. The cross reference is an alphabetically sorted list of all PUBLIC and EXTERN symbols in your program along with memory type and module names. The first module name is the module in which the PUBLIC symbol is defined. Further module names show the modules in which the EXTERN symbol is

defined. If no PUBLIC symbol is present,

** UNRESOLVED ** is shown as first module name.

The option **NOGENERATED** suppresses symbols starting with '?'. These question mark symbols are normally produced by the compiler for calling specific C functions or passing parameters.

The option **NOLIBRARIES** suppresses those symbols, which are defined in a library file.

Example: BL51 myfile.obj IXREF

LX51 myfile.obj IXREF (NOGENERATED)

L251 myfile.obj IXREF(NOLIBRARIES, NOGENERATED)

NOCOMMENTS

Restriction: This control is available in **LX51** and **L251 only.**

Abbreviation: NOCO

Arguments: None

Default: include comment information.

Description: The **NOCOMMENTS** control removes the comment

records contained in the input files from the listing file **and** the object output file. Comment records are added to the object module to identify the compiler or assembler that produced the object file. If you want to exclude comment information just from the listing file you have to use the

PRINTCONTROL control.

See Also: OBJECTCONTROLS, PRINTCONTROLS

Example: L251 MYPROG.OBJ NOCOMMENTS

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NOLINES

Abbreviation: NOLI

Arguments: None

Default: Include line number information

μVision2 Control: Options – Listing – Linker Listing – Line Numbers

Description: For the **BL51** linker/locator, the **NOLINES** control excludes

line number information in the listing file.

For the LX51 linker/locater and the L251 linker/locator, the NOLINES control excludes line number information in the listing file and object output file. If you want to exclude line number information just from the listing file you have to use the PRINTCONTROL control.

NOTE

Line numbers are address information about the source code lines and are used for debugging purposes. The Lx51 linker/locator generates line numbers for source modules in your program only, if the Ax51 assembler and Cx51 compiler include that information in the input object files. Refer to the assembler control NOLINES on page 215 and to the Cx51 User's Guide for information on including line number information in the object files.

See Also: NODEBUGLINES, PRINTCONTROLS,

OBJECTCONTROLS

Example: BL51 MYPROG.OBJ NOLINES

NOMAP

Abbreviation: NOMA

Arguments: None

Default: Include a memory map in the listing file.

μVision2 Control: Options – Listing – Linker Listing – Memory Map

Description: The **NOMAP** control prevents the **Lx51** linker/locator from

including the memory map in the listing file.

Example: BL51 MYPROG.OBJ NOMAP

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NOPUBLICS

Abbreviation: NOPU

Arguments: None

Default: Include information about public symbols

μVision2 Control: Options – Listing – Linker Listing – Public Symbols

Description: For the **BL51** linker/locator, the **NOPUBLICS** control

excludes public symbols from the listing file.

For the LX51 linker/locater and the L251 linker/locator, the NOLINES control excludes public symbols from the listing file and object output file. If you want to exclude public symbols just from the listing file you have to use the

PRINTCONTROLS control.

See Also: NODEBUGPUBLICS, PRINTCONTROLS,

OBJECTCONTROLS

Example: BL51 MYPROG.OBJ NOPUBLICS

NOSYMBOLS

Abbreviation: NOSY

Arguments: None

Default: Include information about local program symbols

μVision2 Control: Options – Listing – Linker Listing – Local Symbols

Description: For the **BL51** linker/locator, the **NOSYMBOLS** control

excludes local symbols from the listing file.

For the LX51 linker/locater and the L251 linker/locator, the NOSYMBOLS control excludes local symbols from the listing file and object output file. If you want to exclude local symbols just from the listing file you have to use the

PRINTCONTROLS control.

NOTE

Symbols information is typically used for debugging purposes. The Lx51 linker/locator generates symbol information for source modules in your program only, if the Ax51 assembler and Cx51 compiler include that information in the input object files. Refer to the assembler control DEBUG on page 203 and to the Cx51 User's Guide for information on including symbol information in the object files.

See Also: NODEBUGSYMBOLS, PRINTCONTROLS,

OBJECTCONTROLS

Example: BL51 MYPROG.OBJ NOSYMBOLS

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PAGELENGTH / PAGEWIDTH

Abbreviation: PL

Arguments: PAGELENGTH (value)

PAGEWIDTH (value)

Default: PAGELENGTH (60)

PAGEWIDTH (132)

μVision2 Control: Options – Listing – Page Width, Page Length

Description: The **PAGELENGTH** control sets the maximum number of

lines per page for the listing file.

The **PAGEWIDTH** control defines the maximum width of

lines in the listing file. The page width may be set to a

number in the 72 to 132 range.

Examples: BL51 PROG.OBJ TO PROG.ABS PAGELENGTH (50) PAGEWIDTH (100)

L251 MYPROG.OBJ PAGELENGTH(30000) PAGEWIDTH(120)

0

PRINT / NOPRINT

Abbreviation: PR / NOPR

Arguments: PRINT (filename)

Default: The listing file is generated using the basename of the output

file. BL51 use the extension .M51; LX51 and L251 use the

extension .MAP as default for the listing file.

μVision2 Control: Options – Listing – Select Folder for Listing Files

Description: The **PRINT** control allows you to specify the name of the

listing file that is generated by the Lx51 linker/locator. The

name must be enclosed in parentheses immediately following the **PRINT** control on the command line.

The **NOPRINT** control prevents the linker/locater from

generating a listing file.

Example: LX51 MYPROG.OBJ TO MYPROG.ABS PRINT (OUTPUT.MAP)

PRINTCONTROLS

Restriction: This control is available in **LX51** and **L251 only**.

Abbreviation: PC

Arguments: PRINTCONTROLS (subcontrol[, ...])

Default: all debug information is printed in the listing file.

μVision2 Control: Options – Listing – Linker Listing

Description: The **PRINTCONTROLS** control allows you to remove

specific debug information from the listing file. The *subcontrol* option can be one or more of the following

parameters:

subcontrol	Removes from the listing file		
NOCO MMENTS	comment records.		
<u>NOLI</u> NES	line number information.		
NOPUBLICS	public symbol information.		
NOSYMBOLS	local symbol information.		
<u>PU</u> RGE	complete debug information.		

See Also: NOCOMMENTS, NOLINES, NOPUBLICS,

NOSYMBOLS, OBJECTCONTROLS, PURGE

Example: LX51 MYPROG.OBJ PRINTCONTROLS (NOLINES, NOSYMBOLS)

PURGE

Restriction: This control is available in **LX51** and **L251 only.**

Abbreviation: PU

Default: all debug information is processed.

 μ Vision2 Control: Options – Lx51 Misc – Misc Controls – enter the control.

Description: The **PURGE** control allows you to remove the complete

debug information contained in the input files from the listing file **and** the object output file. **PURGE** has the same

effect as specifying NOCOMMENTS, NOLINES,

NOPUBLICS and **NOSYMBOLS**. The debug information is only required for program debugging and has no influence on the executable code. If you want to exclude line number information just from the listing file you have to use the

PRINTCONTROLS control.

See Also: NOCOMMENTS, NOLINES, NOPUBLICS,

NOSYMBOLS, OBJECTCONTROLS,

PRINTCONTROLS

Example: L251 MYPROG.OBJ PURGE

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WARNINGLEVEL

Restriction: This control is available in **LX51** and **L251 only**.

Abbreviation: WL

Arguments: A number between 0 .. 2.

Default: WARNINGLEVEL (2)

μVision2 Control: Options – Lx51 Misc – Warnings – Warning Level.

Description: The **WARNINGLEVEL** control allows you to suppress

linker warnings. Refer to "Warnings" on page 360 for a full

list of the linker warnings.

Warning Level	Description	
0	Disables almost all linker warnings.	
1	Lists only those warnings that may generate incorrect code, including information about data type mismatches of total different types.	
2 (Default)	Lists all WARNING messages including warnings all data type mismatches.	

Example: LX51 MYFILE.OBJ WL (1)

Example Listing File

The following example includes all optional sections of the listing file.

```
L251 LINKER/LOCATER V3.00
                                                               09/06/2000 12:09:21 PAGE 1
L251 LINKER/LOCATER V3.00, INVOKED BY:
                                                                  The listing file shows the command
E:\L251.EXE MEASURE.OBJ, MCOMMAND.OBJ, GETLINE.OBJ IXREF
                                                                  line that invoked the linker.
CPU MODE:
              251 SOURCE MODE
                                                                  CPU mode, interrupt frame size,
INTR FRAME: 4 BYTES SAVED ON INTERRUPT
                                                                  memory model and floating point
MEMORY MODEL: SMALL WITH FLOATING POINT ARITHMETIC
                                                                  arithmetic are listed.
INPUT MODULES INCLUDED:
                                                                  Object modules that were included
  MEASURE.OBJ (MEASURE)
                                                                  along with translator information
          COMMENT TYPE 0: C251 V3.00
                                                                  are listed.
  MCOMMAND.OBJ (MCOMMAND)
          COMMENT TYPE 0: C251 V3.00
  GETLINE.OBJ (GETLINE)
         COMMENT TYPE 0: C251 V3.00
  C:\KEIL\C251\LIB\C2SFPS.LIB (?C_FPADD)
         COMMENT TYPE 0: A251 V3.00
                                                                  LX51 and L251 list an overview of all
ACTIVE MEMORY CLASSES OF MODULE: MEASURE (MEASURE)
                                                                  memory classes in used.
         START END
                               MEMORY CLASS
______
FF0000H FF0000H FFFFFFH CODE 000000H 000000H 00007FH DATA
000000H 000000H 0000FFH IDATA
010000H 010000H 01FFFFH XDATA
000020H.0 000020H.0 00002FH.7 BIT
000000Н
          000000H 00FFFFH
                                                                  The memory map is included
                                                                  You can disable the memory map
MEMORY MAP OF MODULE: MEASURE (MEASURE)
                                                                  using the NOMAP control.
START STOP
                   LENGTH ALIGN RELOC MEMORY CLASS SEGMENT NAME
______
000000H 000007H 000008H --- AT.. DATA "REG BANK 0" 000008H 000008H 000008H --- AT.. DATA "REG BANK 1"
000010H 000010H 000001H BYTE UNIT DATA
000011H 00001FH 00000FH BYTE UNIT IDATA
000020H.0 000020H.2 000000H.3 BIT UNIT BIT
000020H.3 000020H.3 000000H.1 BIT UNIT BIT
                                                             ?DT?GETCHAR
_IDATA_GROUI
?BI?MEASURE
                                                                 ?DT?GETCHAR
                                                                   IDATA GROUP
                                                                 ?BI?GETCHAR
000020H.4 000021H.6 000001H.3 BIT UNIT BIT
                                                                   BIT GROUP
000022H 000039H 000018H BYTE UNIT DATA
00003AH 000065H 00002CH BYTE UNIT DATA
000066H 000165H 000100H BYTE UNIT EDATA
010000H 011FF7H 001FF8H BYTE UNIT XDATA
                                                                 ?DT?MEASURE
                                                                 STACK
                                                                   DATA GROUP
                                                                 ?XD?MEASURE
FF0000H FF0002H 000003H --- OFFS.. CODE FF0003H FF0008H 000006H BYTE UNIT CODE
                                                                  ?CO?START251?3
                                                                  ?PR?GETCHAR?UNGETCHAR
                                                                  An overlay map is listed after the
OVERLAY MAP OF MODULE: MEASURE (MEASURE)
                                                                  memory map. The overlay map shows
                                                                  the call tree of your application.
FUNCTION/MODULE
                                           BIT GROUP DATA GROUP IDATA GROUP
--> CALLED FUNCTION/MODULE START STOP START STOP START STOP
______
TIMERO/MEASURE
--> SAVE CURRENT MEASUREMENTS/MEASURE
*** NEW ROOT ***************
?C C251STARTUP
--> MAIN/MEASURE
```

```
MAIN/MEASURE
                                   ---- 003AH 003CH 0011H 001FH
--> CLEAR RECORDS/MEASURE
--> PRINTF/PRINTF
--> GETLINE/GETLINE
--> TOUPPER/TOUPPER
--> READ INDEX/MEASURE
--> GETKEY/ GETKEY
--> MEASURE DISPLAY/MCOMMAND
--> SET TIME/MCOMMAND
--> SET INTERVAL/MCOMMAND
PRINTF/PRINTF
                                   20H.4 21H.4 0049H 0064H -----
--> PUTCHAR/PUTCHAR
                                   ---- 003DH 0040H ----
GETLINE/GETLINE
                                                       A list of all public symbols is printed.
PUBLIC SYMBOLS OF MODULE: MEASURE (MEASURE)
                                                       This list can be disabled using the
                                                       NOPUBLICS or PRINTCONTROLS
     VALUE CLASS TYPE
                                PUBLIC SYMBOL NAME
                                                      control.
     -----
     00000021H.2 BIT BIT ?C?ATOFFIRSTCALL
     00FF0F5CH CODE ---
00FF12D3H CODE ---
00000020H.3 BIT BIT
                                 ?C?CASTF
                              ?C?CHARLOADED
                                 ?C?CCASE
     00FF186DH CODE
                       ---
                                 SSCANF
SYMBOL TABLE OF MODULE: MEASURE (MEASURE)
                                                       The symbol table lists the complete
                                                       debug information of your project.
              REP CLASS TYPE SYMBOL NAME
     VALUE
     ______
               MODULE ---
                                --- MEASURE
     00000020H.2 PUBLIC
                                        MEASUREMENT_INTERVAL
                        BIT
DATA ---
BIT
     00000036H PUBLIC
                                          INTERVAL
     00000020H.1 PUBLIC BIT
                                         MDISPLAY
                                                      You can use the PRINTCONTROL
                                BYTE INTCYCLE
WORD SAVEFIRST
     00000035H PUBLIC DATA
00000033H PUBLIC DATA
                                                       control to exclude part of the
                                                      symbol information from the listing.
                                         LVL=0
     00FF000EH BLOCK CODE ---
     00FF000EH LINE CODE ---
00FF000EH LINE CODE ---
                                         #87
                                          #88
               BLOCKEND ---
                                 ---
                                         LVL=0
                                                       The IXREF control instructs Lx51
INTER-MODULE CROSS-REFERENCE LISTING
                                                       to include a cross reference table.
NAME . . . . . . . . . . . . CLASS MODULE NAMES
_____
?C?ATOFFIRSTCALL . . . . BIT
                              ?C ATOF SCANF
?C?CASTF . . . . . . . . CODE
                              ?C CASTF MCOMMAND
?C?CCASE . . . . . . CODE
                               ?C CCASE PRINTF SCANF
                               GETCHAR UNGETC
?C?CHARLOADED. . . . . . BIT
?C?COPY2 . . . . . . . CODE ?C_COPY2 MCOMMAND MEASURE
?C?FCASTC. . . . . . . CODE ?C FCAST ?C ATOF MCOMMAND
                               ?C_FCAST
?C?FCASTI. . . . . . . . . CODE
?C FCAST
                                ?C_FPADD ?C_ATOF ?C_FPCONVERT
?C?FPADD . . . . . . . CODE
                                ?C ATOF SCANF
?C?FPATOF. . . . . . . . CODE
?C?FPCMP . . . . . . . CODE
                                ?C FPCMP
?C?FPCMP3. . . . . . . CODE
                                ?C FPCMP MCOMMAND
?C?FPCONVERT . . . . . . CODE
                                ?C FPCONVERT PRINTF
:
```

Output File Controls

The linker/locator either generates absolute object files or banked object files. Absolute object files contain no relocatable information or external references. Absolute object files can be loaded into debugging tools or may be converted into Intel HEX files for PROM programming by **OHx51** Object-Hex Converter.

Banked object files generated by the **BL51** linker/locater must be converted by the **OC51** Banked Object File Converter into absolute object files (one for each bank) to convert them into Intel HEX files by the **OH51** Object-Hex Converter.

The generated object module may contain debugging information if the linker/locator is so directed. This information facilitates symbolic debugging and testing. You may use the linker controls to suppress debugging information in the object file. The following table provides an overview of the controls that control information in the output file. For a detailed description of each control refer to the page specified in the table.

BL51	LX51, L251	Page	Description
-	<u>AS</u> SIGN	322	Defines public symbols on the command line.
<u>IB</u> ANKING	-	323	Generate bank switch code for Infineon TV TEXT devices SDA555x and SDA30C16x.
<u>NA</u> ME	<u>NA</u> ME	323	Specifies a module name for the object file.
<u>NOAJ</u> MP	NOAJMP	325	Generate bank switch code without AJMP instructions.
NOINDIRECTCALL	NOINDIRECTCALL	327	Do not generate by default bank switch code for indirectly called functions.
NOJMP <u>T</u> AB	-	328	Do not generate bank switch code.
-	<u>NOTY</u> PE	327	Specifies a module name for the object file.
NODEBUGLINES NODEBUGPUBLICS NODEBUGSYMBOLS	<u>OBJECTCONTROLS</u>	325 330	Excludes debug information from the object file.

ASSIGN

Restriction: This control is available in **LX51** and **L251 only**.

Abbreviation: AS

Arguments: ASSIGN (symname (value) [, ...])

Default: None

μVision2 Control: Options – Lx51 Misc – Assign.

Description: ASSIGN defines a PUBLIC symbol whit a numeric value at

Lx51 linker/locater level. The PUBLIC symbol is handled as a number without a specific memory class and matches with an unresolved external symbol with the same name.

Example: L251 MYFILE.OBJ ASSIGN (FUNC (0x2000), BITVAR (20H.2))

In this example the public symbols FUNC and BITVAR are defined. The value 0x2000 is given as value for FUNC. The value 20H.2 is used as bit-address for BITVAR.

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IBANKING

Restriction: This control is available in **BL51 only.**

Abbreviation: IB

Arguments: IBANKING (bank_sfr_address)

Default: The default *bank sfr address* is 0x94. This is also the SFR

address for the support Infineon devices.

μVision2 Control: Options – Lx51 Misc – Misc Controls – enter the control.

Description: With the **BL51** linker/locater control **IBANKING** the linker

uses the on-chip code banking hardware of the Infineon SDA30C16x/26x and SDA555x TV TEXT devices. The **BL51** linker/locater places automatically all code segments in the bank area, which do not have the ?CO? prefix or ?CO postfix. Segments with a ?CO prefix or postfix are placed

into the common area.

The module L51_BANK.A51 is not used when the control IBANKING is used. The BL51 linker/locater generates in this operation mode a jump table with the following format:

MOV bank_sfr, #BANK_NUMBER
LJMP target

NOTE

When you are using this directive, you need also special C51 run-time libraries. Please contact **Keil Software** to obtain these C51 run-time libraries.

See Also: NOAJMP, NOINDIRECTCALL, NOJMPTAB,

BANKAREA

Example: BL51 BANK0 {MODULA.OBJ}, BANK1 {MODULB.OBJ} IBANKING

BL51 BANKO {MODULA.OBJ}, BANK1 {MODULB.OBJ} IB (80H)

NAME

Abbreviation: NA

Arguments: NAME (*modulename*)

Default: Module name of the first object file in the input list is used.

μVision2 Control: Options – Lx51 Misc – Misc Controls – enter the control.

Description: Use the **NAME** control to specify a module name for the

absolute object module that the **BL51** linker/locator generates. The **NAME** control may be accompanied by the module name (in parentheses) that you want to assign. If no module name is specified with the **NAME** control, the name of the first input module is used for the module name.

NOTE

The module name specified with the NAME control is not the filename of the absolute object file. The module name is stored in the object module file and may be accessed only by a program that reads the contents of that file.

Example: BL51 MYPROG.OBJ TO MYPROG.ABS NAME (BIGPROG)

In this example **BIGPROG** is the module name stored in the object file.

NOAJMP

Abbreviation: NOAJ

Default: The Lx51 linker/locater generates for code banking

applications an inter-bank jump table. This bank switch table is used for jumps into a code bank from a different code bank or the common area. Depending on the table size, the linker uses AJMP or LJMP instructions within this bank

switch table.

μVision2 Control: Options – Lx51 Misc – Misc Controls – enter the control.

Description: With the **NOAJMP** control you can disable the AJMP

instruction in the inter-bank jump table. This option is required for 8051 derivatives that are not supporting the

AJMP instruction.

See Also: IBANKING, NOINDIRECTCALL, NOJMPTAB,

BANKAREA

Example: BL51 MYPROG.OBJ NOAJMP

NODEBUGLINES, NODEBUGPUBLICS, NODEBUGSYMBOLS

Restriction: This control is available in **BL51 only**. For **LX51** and **L251**

use the **OBJECTCONTROLS** control.

Abbreviation: NODL, NODP, NODS

Default: Include complete debug information in the output file.

μVision2 Control: Options – Lx51 Misc – Misc Controls – enter the control.

Description: The **NODEBUGLINES** control directs the **BL51**

linker/locator to exclude line number information from the

output object file.

The **NODEBUGPUBLICS** control excludes public symbol

information from the output object file.

The **NODEBUGSYMBOLS** control excludes local symbol

information from the output object file.

NOTE

Line number and symbol information in the absolute object file is used for symbolic debugging in the μV ision2 debugger or in-circuit emulator. If you exclude debug information, source level debugging of your program is no longer possible.

See Also: NOLINES, NOPUBLICS, NOSYMBOLS,

OBJECTCONTROLS, PRINTCONTROLS

Example: BL51 MYPROG.OBJ NODEBUGLINES NODEBUGSYMBOLS

NOINDIRECTCALL

Abbreviation: NOIC

Default: In code banking applications, the Lx51 linker/locater inserts

an interbank CALL for each function that is indirectly called via a function pointer. This is done (by the linker) to ensure that functions which are invoked through a function pointer

are available to all code banks.

μVision2 Control: Options – Lx51 Misc – Misc Controls – enter the control.

Description: The **NOINDIRECTCALL** control allows you to disable the

generation of interbank CALLs (for indirect function calls). This directive is useful if your application uses tables that contain pointers to functions and if you ensure that these

indirect function calls never cross a code bank.

See Also: IBANKING, NOAJMP, NOJMPTAB, BANKAREA

Example: BL51 MYPROG.OBJ NOINDIRECTCALL

NOJMPTAB

Restriction: This control is available in **BL51 only.**

Abbreviation: NOJT

Default: The Lx51 linker/locater generates for code banking

automatically an inter-bank jump table or bank switch table. For each function that is located in a code bank and is called from a different code bank or the common area the linker inserts a bank switch code into the inter-bank jump table

redirects the function call to this table.

μVision2 Control: Options – Lx51 Misc – Misc Controls – enter the control.

Description: When the **NOJMPTAB** control is stated, BL51 no longer

inserts inter-bank calls for program calls. This feature is implemented to use the user-defined bank switch mechanism for code banking. The **NOJMPTAB** directive modifies the

following features of **BL51**:

 The linker no longer needs the bank switch configuration file: L51 BANK.OBJ.

The linker does not modify any jump call instructions.

The linker does not generate any warnings if a jump/call is made to another bank. The user must ensure that the proper bank is selected before a call is made since the BL51 linker/locator no longer selects the bank automatically.

See Also: IBANKING, NOAJMP, NOINDIRECTCALL,

BANKAREA

Example: BL51 MYPROG.OBJ NOJMPTAB

NOTYPE

Restriction: This control is available in **LX51** and **L251 only**.

Syntax: NOTYPE

Abbreviation: NOTY

Default: Include complete type information in the output file.

μVision2 Control: Options – Lx51 Misc – Misc Controls – enter the control.

Description: The **NOTYPE** control removes symbol type information for

debug symbols from the output file. The symbol type information is only required for program debugging and has

no influence on the executable code.

NOTE

Symbol type information in the absolute object file is used for symbolic debugging in the μV ision2 debugger or incircuit emulator. The Cx51 compiler generates complete symbol information up to structure members and parameter passing values. If you exclude symbol type information, you might no be able to display variables during debugging.

See also: OBJECTCONTROL, PURGE

Example: LX51 file1.obj NOTYPE

OBJECTCONTROLS

Restriction: This control is available in **LX51** and **L251 only.**

Abbreviation: OC

Arguments: OBJECTCONTROLS (subcontrol , ...)

Default: Include complete debug information in the output file.

μVision2 Control: Options – Lx51 Misc – Misc Controls – enter the control.

Description: The **OBJECTCONTROLS** control allows you to remove

specific debug information from the object output file. The *subcontrol* option can be one or more of the following

parameters:

subcontrol	Removes from the object output file
NOCOMMENTS	comment records.
<u>NOLI</u> NES	line number information.
NOPUBLICS	public symbol information.
NOSYMBOLS	local symbol information.
<u>PU</u> RGE	complete debug information.

See Also: NOCOMMENTS, NOLINES, NOPUBLICS,

NOSYMBOLS, OBJECTCONTROLS, PURGE

Example: LX51 MYPROG.OBJ OBJECTCONTROLS (NOCOMMENTS)

Segment and Memory Location Controls

The **Lx51** linker/locator allows you to specify the size of the different memory areas or memory classes, the order of the segments within the different memory areas, and the location or absolute memory address of different segments. Also the size of segments can be manipulated and specific memory areas can be excluded from being used. These segment manipulations are performed using the following controls.

BL51	LX51, L251	Page	Description
<u>BA</u> NKAREA	<u>BA</u> NKAREA	332	Specifies the address range where the code banks are located.
<u>B</u> ANK <u>x</u>	-	333	Locates and orders segments in code bank <i>x</i> (where <i>x</i> is a code bank from 0 to 31).
<u>BI</u> T	-	334	Locates and orders BIT segments.
-	<u>CL</u> ASSES	336	Specifies a physical address range for segments in a memory class.
CODE	_	338	Locates and orders CODE segments.
<u>DA</u> TA	_	339	Locates and orders DATA segments.
<u>ID</u> ATA	_	340	Locates and orders IDATA segments.
<u>NOSO</u> RTSIZE	NOSORTSIZE	341	Disable size sorting for segments before allocating the memory.
<u>PD</u> ATA	-	341	Specifies start address for PDATA segments.
PRECEDE	-	343	Locates and orders segments that should precede others in DATA memory.
RAMSIZE	_	344	Specifies size of DATA and IDATA memory.
-	<u>RE</u> SERVE	345	Reserves memory ranges and prevents the linker from using these memory areas.
-	<u>SE</u> GMENTS	346	Defines physical memory addresses and orders for specified segments.
-	<u>S</u> EG <u>S</u> IZE	348	Modifies the size for a specific segment.
<u>ST</u> ACK	-	349	Locates and orders STACK segments.
<u>XD</u> ATA	-	350	Locates and orders XDATA segments.

The **Lx51** linker/locator locates segments in according their memory class and follows a predefined order of precedence. The standard allocation algorithm usually produces the best workable solution and does not requiring you to enter any segment names on the command line. The controls described in this section allow you to define the physical memory layout of your target system and more closely control the location of segments within the different memory classes. Refer to "Locating Programs to Physical Memory" on page 273 for examples on how to define the available memory in your *x51* system.

BANKAREA

Abbreviation: BA

Arguments: BANKAREA (start address, end address)

None **Default:**

μVision2 Control: Options – Target – Code Banking – Bank Area.

Description: Use the **BANKAREA** control to specify the starting and

> ending address of the area where the code banks will be located. The addresses specified should reflect the actual address where the code bank ROMs are physically mapped. All segments that are assigned to a bank will be located within this address range unless they are defined differently using the **BANK***x* control. Refer to "Bank Switching" on page 293 for more information about the code banking

controls.

See Also: **BANK***x*

LX51 COMMON{C ROOT.OBJ}, BANKO{C BANKO.OBJ}, **Example:**

BANKO {C_BANKO.OBJ}, BANK1 {C_BANK1.OBJ},

BANK2 (C_BANK2.OBJ) TO MYPROG.ABS &

BANKAREA (8000H, 0FFFFH)

BANKx

Restriction: This control is available in **BL51 only**. For **LX51** and **L251**

use the **SEGMENTS** control.

Abbreviation: B0, B1, B2, ... B30, B31

Arguments: BANKx ([start_address] [segname [(address)] [, ...]])

μVision2 Control: Options – BL51 Misc – Misc Controls – enter the control.

Description: Use the **BANK***x* control to specify a code bank for segments

(x in **BANK**x is replaced by a bank number). Refer to "Bank Switching" on page 293 for more information about

the code banking controls.

Segments are located in the specified code bank starting at *start_address* or address 0000h if *start_address* is not specified. If an *address* is specified the segment referred by

segname will be located at this address.

If you allocate a constant segment of a **Cx51** program into a code bank, you must manually ensure that the proper code bank is used when accessing that constant data. You can do this with the **switchbank** function that is defined in the **L51_BANK.A51** module. "BANK_EX2 – Banking with Constants" on page 407 shows a complete example program.

See Also: BANKAREA, CODE

Example: This example will locate the segment ?PR?FUNC1?A that

belongs to the C function func1 in module "A.C" into code

bank 1 starting at address 0x8000. The segment ?PR?FUNC1?A will be located at address 0x8200.

BL51 COMMON{A.OBJ}, BANKO{B.OBJ} BANK1(0x8000, ?PR?FUNC1?A, ?PR?FUNC2?B(0x8200))

BIT

Restriction: This control is available in **BL51 only**. For **LX51** and **L251**

use the CLASSES and SEGMENTS control.

Abbreviation: BI

Arguments: BIT ($[start_address]$ [segname [(address)]], ...])

μVision2 Control: Options – BL51 Locate – Bit.

Description: The **BIT** control allows you to specify:

 The starting address for segments placed in the bit-addressable internal data space

 The order of segments within the bit-addressable internal data space

• The absolute memory location of segments in the bit-addressable internal data space.

Addresses that you specify with the **BIT** control are bit addresses. Bit addresses 00h through 7Fh reference bits in DATA memory bytes from byte address 20h to 2Fh (16 bytes of 8 bits each, $16 \times 8 = 128 = 80h$). Bit addresses that are evenly divisible by 8 are aligned on a byte boundary. A DATA segment that is bit-addressable can be located with the BIT control; however, the bit address specified must be byte aligned, that means evenly divisible by 8.

See Also: CODE, DATA, IDATA, XDATA

Examples: The following example specifies that relocatable BIT

segments be located at or after bit address 48 decimal (30 hex) which is equivalent to byte address 26H.0 in the data

memory:

BL51 MYPROG.OBJ BIT(48)

or

BL51 MYPROG.OBJ BIT(0x30)

To specify the order for segments, you must include the segment names, separated by commas. The following example places the PDT?A, PDT?B, and PDT?C segments at the beginning of the bit-addressable DATA memory.

BL51 MYPROG.OBJ, A.OBJ, B.OBJ, C.OBJ BIT (?DT?A, ?DT?B, ?DT?C)

You may also specify the bit address for the segments. The following example places the ?DT?A and ?DT?B segments at 28h and 30h:

BL51 MYPROG.OBJ, A.OBJ, B.OBJ BIT(?DT?A(28h),?DT?B(30h))

CLASSES

Restriction: This control is available in **LX51** and **L251** only.

Abbreviation: CL

Arguments: CLASSES (classname (range [, ...]) [, ...])

Default: None

μVision2 Control: Options – **Lx51** Locate – User Classes.

Description: The CLASSES control specifies the physical address range

control provides an efficient way to define the physical memory layout. If the address limits for a memory class are not specified with the **CLASSES** control, the **Lx51** linker/locater uses the physical address limits of the memory class. The address ranges that are used are listed in the linker MAP file in the section **ACTIVE MEMORY CLASSES**. It is recommended to check this section of the MAP file, since it lists where the Lx51 linker/locater assumes memory in your target hardware. More information about "Locating Programs to Physical Memory" can be found on page 273.

for segments within a memory class. The CLASSES

With the **CLASSES** control the absolute address for segments with the relocation type **OFFS** can be modified. For more information on how to declare such segments refer to "Relocation Type" on page 108. The offset is specified as first address in the range field with a '\$' prefix, for example: **CLASSES** (**CODE** (\$0xFF8000, 0xFF8000 - 0xFFFFFF). In this case all segments that are defined with the OFFS relocation type, are redirected to the address 0xFF8000. Typically the interrupt and reset vectors of a program are defined this way. In this way, you can quickly redirect these vectors, for example, when you are debugging programs with the Monitor-251 installed at address 0xFF0000.

A memory class can be copied into RAM for execution whereas the content is stored in the SROM memory class. In this case you must copy the memory class from ROM to RAM before execution. Empty brackets after the address range are used to store the content of a memory class within the address range of the SROM memory class, for example: CLASSES (NCONST (0xE000 - 0xFFFF)[]). Refer to "Use RAM for the 251 Memory Class NCONST" on page 292 for a program example that uses this feature.

See Also: SEGMENTS

Examples: The following example specifies the address range of the EDATA and CODE memory class:

```
L251 MYFILE.OBJ &
CLASSES (EDATA (0 - 0x41F, 0x2000H - 0x3FFF),
CODE (0xFF0000 - 0xFF7FFF))
```

This example defines the memory classes for a classic 8051 device:

```
LX51 MYFILE.OBJ

CLASSES (IDATA (I:0-I:0xFF), XDATA (X:0-X:0xEFFF),

CODE (C:0-C:0x7FFF, C:0xC000-C:0xFFFF))
```

In this example the user-defined memory class XDATA_FLASH is defined. Refer to "User-defined Class Names" on page 107 for more information.

```
LX51 MYFILE.OBJ
CLASSES (XDATA_FLASH (X:0x8000-X:0xEFFF))
```

CODE

Restriction: This control is available in **BL51 only**. For **LX51** and **L251**

use the **CLASSES** and **SEGMENTS** control.

Abbreviation: CO

Arguments: CODE ([address_range] [segname [(address)] [, ...]])

μVision2 Control: Options – BL51 Locate – Code Range, Code.

Description: The **CODE** control allows you to specify:

 The address range for segments placed in the CODE memory class or CODE memory space.

• The order of segments within the CODE space.

 The absolute memory location of segments in the CODE memory space.

See Also: BIT, DATA, IDATA, XDATA

Examples: The example below specifies that relocatable CODE segments be located in the address space 0 - 0x3FFF and

0x8000 - 0xFFFF:

BL51 MYPROG.OBJ CODE(0 - 0x3FFF, 0x8000 - 0xFFFF)

To specify the order for segments, you must include the names of the segments separated by commas. The following example will place the <code>?PR?FUNC1?A</code> and <code>?PR?FUNC2?A</code> segments at the beginning of the CODE memory:

BL51 A.OBJ CODE(?PR?FUNC1?A, ?PR?FUNC2?A)

You can also specify the memory location for a segment. The example below will place the <code>?PR?FUNC1?A</code> segment at 800h and the <code>?PR?FUNC2?A</code> segment after at this segment:

BL51 A.OBJ CODE(?PR?FUNC1?A (0x800), ?PR?FUNC2?A)

DATA

Restriction: This control is available in **BL51 only**. For **LX51** and **L251**

use the **CLASSES** and **SEGMENTS** control.

Abbreviation: DA

Arguments: DATA ([start_address] [segname [(address)] [, ...]])

μVision2 Control: Options – BL51 Locate – Data.

Description: The **DATA** control allows you to specify:

 The address range for segments placed in the directly-addressable DATA space.

The order of segments within the DATA space.

 The absolute memory location of segments in the directly-addressable internal DATA space.

See Also: BIT, CODE, IDATA, XDATA

Examples: The example below specifies that relocatable DATA

segments be located at or after address 48 decimal (30 hex)

in the on-chip DATA memory:

BL51 MYPROG.OBJ DATA(48)

or

BL51 MYPROG.OBJ DATA(0x30)

To specify the order for segments, you must include the names of the segments separated by commas. The following example will place the ?DT?A, ?DT?B, and ?DT?C segments at the beginning of the DATA memory:

BL51 A.OBJ, B.OBJ, C.OBJ DATA (?DT?A, ?DT?B, ?DT?C)

You can also specify the memory location. The example below will place the PDT?A and PDT?B segments at 28h and 30h in the DATA memory:

BL51 MYPROG.OBJ, A.OBJ, B.OBJ DATA (?DT?A(28h), ?DT?B(30h))

IDATA

Restriction: This control is available in **BL51 only**. For **LX51** and **L251**

use the **CLASSES** and **SEGMENTS** control.

Abbreviation: ID

Arguments: IDATA ([start_address] [segname [(address)] [, ...]))

μVision2 Control: Options – BL51 Locate – Idata.

Description: The **IDATA** control allows you to specify:

 The starting address for segments placed in the indirectly-addressable on-chip IDATA space.

- The order of segments within the IDATA space.
- The absolute memory location of segments in the IDATA memory space.

See Also: BIT, CODE, DATA, XDATA

Examples: The example below specifies that relocatable IDATA segments be located at or after address 64 decimal (40 hex)

in the IDATA memory.

BL51 MYPROG.OBJ IDATA(64)

or

BL51 MYPROG.OBJ IDATA(0x40)

To specify the order for segments, you must include the names of the segments separated by commas. The following example places the ?ID?A, ?ID?B, and ?ID?C segments at the beginning of the IDATA memory:

BL51 A.OBJ, B.OBJ, C.OBJ IDATA (?ID?A, ?ID?B, ?ID?C)

You may also specify the memory location. This example places the ?ID?A and ?ID?B segments at 30h and 40h in the on-chip IDATA memory:

BL51 MYPROG.OBJ, A.OBJ, B.OBJ IDATA (?ID?A(30h),?ID?B(40h))

NOSORTSIZE

Abbreviation: NOSO

Arguments: None

Default: The segments are sorted according their size before the Lx51

linker/locater allocates the memory space. This reduces typically the memory gaps that are required to fulfill the

allocation requirements.

 μ Vision2 Control: Options – Lx51 Misc – Misc Controls: enter the control.

Description: The **NOSORTSIZE** control allows you disable the sorting

algorithm. In this case the linker allocates the memory in

the order the segments appear in the input files.

Example: BL51 MYPROG.OBJ NOSORTSIZE

PDATA

Restriction: This control is available in **BL51 only**. For **LX51** and **L251**

use the **CLASSES** and **SEGMENTS** control.

Abbreviation: None

Arguments: PDATA (address)

μVision2 Control: Options – BL51 Locate – Pdata.

Description: The **PDATA** control allows you to specify the starting

address in external data space for **PDATA** segments. You must enter the starting address immediately following the **PDATA** control on the command line. The address must be

enclosed in parentheses.

In addition to specifying the starting address for PDATA segments on the linker command line, you must also modify the startup code stored in **STARTUP.A51** to indicate that PDATA segments are located at 8000h. Refer to the *C51 User's Guide* for more information about PDATA and

COMPACT model programming.

See Also: XDATA

Example: This example specifies that PDATA segments are to be

located starting at address 8000 hex in the external data

memory.

BL51 MYPROG.OBJ PDATA(0x8000)

PRECEDE

Restriction: This control is available in **BL51 only**. For **LX51** and **L251**

use the CLASSES and SEGMENTS control.

Abbreviation: PC

Arguments: PRECEDE (segname [(address)][, ...])

μVision2 Control: Options – BL51 Locate – Precede.

Description: The **PRECEDE** control allows you to specify segments that

lie in the on-chip DATA memory that should precede other segments in that memory space. Segments that you specify with this control are located after the **BL51** linker/locator has located register banks and any absolute **BIT**, **DATA**, and **IDATA** segments, but before any other segments in the

internal DATA memory.

See Also: DATA, STACK

Examples: You specify segment names with the **PRECEDE** control.

Segment names must be separated by commas and must be

enclosed in parentheses immediately following the **PRECEDE** control. For example:

BL51 MYPROG.OBJ, A.OBJ, B.OBJ PRECEDE (?DT?A,?DT?B)

The segments that you specify are located at the lowest available memory location in the DATA memory in the order that you specify. You may also specify the memory location of the segments you specify with the **PRECEDE** control. The example below places the PDT?A and PDT?B segments at O9h and 13h in the DATA memory:

BL51 A.OBJ, B.OBJ PRECEDE (?DT?A(09h),?DT?B(13h))

RAMSIZE

Restriction: This control is available in **BL51 only**. For **LX51** and **L251**

use the CLASSES and SEGMENTS control.

Abbreviation: RS

Arguments: RAMSIZE (value)

Default: RAMSIZE (128)

μVision2 Control: Generated from the Device Database Information.

Description: The **RAMSIZE** control allows you to specify the number of

bytes of **DATA** and **IDATA** memory that are available in your target 8051 derivative. The number of bytes must be a number between 64 and 256. This number must be enclosed

in parentheses.

NOTE

In the device data sheets the size of the DATA and IDATA memory is usually referred as on-chip RAM size. However, several new 8051 devices have additional on-chip RAM that is mapped into the XDATA space. If the on-chip RAM size of your 8051 derivative is more than 256 bytes, then your device has most likely additional RAM that is accessed as XDATA memory. In this case use **RAMSIZE** (256) to enable the complete DATA and IDATA address space and define the additional on-chip RAM with the **XDATA** control.

Example: BL51 MYPROG.OBJ RAMSIZE(256)

This example specifies there are 256 bytes of on-chip memory for **DATA** and **IDATA** that may be allocated by the linker

RESERVE

Restriction: This control is available in **LX51** and **L251** only.

Abbreviation: RE

Arguments: RESERVE (range[, ...])

Default: no memory areas are reserved.

μVision2 Control: Options – Lx51 Misc – Reserve.

Description: The **RESERVE** directive allows you to prevent **Lx51** from

locating segments in the specified address ranges of the physical memory. The Lx51 linker/locater will not use all

memory address within the specified address range.

If an absolute segment uses a reserved memory area, a warning message is generated. Refer to "Error Messages" on page 360 for more information about this directive.

See Also: CLASSES, SEGMENTS

Example: L251 MYPROG.OBJ RESERVE (0x200 - 0x3FFF,

0xFF8000H - 0xFFBFFFH)

SEGMENTS

Restriction: This control is available in **LX51** and **L251** only.

Abbreviation: SE

Arguments: SEGEMENTS (segname (address) , ...)

μVision2 Control: Options – Lx51 Locate – User Segments.

Description: The **SEGMENTS** control allows you to specify:

- The absolute memory location of a segment. The absolute address can be either a start or an end address.
- The order of segments within the memory. Segments may be located as first or last segment. Segments defined in the SEGMENTS control are allocated sequentially. By default, the first segment is located at the lowest possible address range (specified with the CLASSES control). Subsequent segments are located at ascending addresses. When you are using the keyword LAST in the address field, then the segment is located as last segment for this memory class.
- A segment can be executed in RAM whereas the content is stored in the ROM memory class. Such segments need to be copied from ROM to RAM before execution. You can specify both the ROM address that stores the content and the RAM address that is used to address the segment during program execution. This syntax is:

SEGMENTS Syntax for Store in ROM and Executed in RAM

SEGMENTS (segment_name(exec_address)[store_address], ...)

exec_address specifies the execution address for the segment.
store address is the address where the segment is stored in ROM.

SEGMENTS (segment_name(exec_address)[], ...)

If you specify empty brackets [] for the **store_address** the segment will be stored within the address range of the **SROM** memory class.

SEGMENTS (segment name(exec address)[!store address], ...)

Lx51 does not reserve the space for execution, if the exclamation mark is given before the **store_address**. This is useful, if the segment content is copied over RAM that is temporarily used, for example the stack area.

SEGMENTS (segment_name(exec_address)[!], ...)

If not **store_address** is given the segment is stored within the range of the **SROM** memory class. Also here not space is reserved for execution.

See Also: CLASSES, RESERVE, SEGSIZE

Examples:

The example below will place the ?DT?A and ?DT?B segments at 28h and 30h in the DATA memory:

LX51 A.OBJ, B.OBJ SEGMENTS (?DT?A(D:0x28),?DT?B(D:0x30))

To specify the order for segments, you must include the names of the segment separated by commas. The following example places the ?DT?A, ?DT?B, and ?DT?C segments at the beginning of the memory class. If these segments belong to the DATA memory class they will be places as first segments in the DATA memory class.

L251 A.OBJ, B.OBJ, C.OBJ SEGMENTS (?DT?A,?DT?B,?DT?C)

A segment can be located to a code bank. The next example locates the segment <code>?pr?func2?B</code> into code bank 0 and the segment <code>?pr?func1?A</code> to address <code>0x8000</code> in code bank 1.

```
L251 BANKO {A.OBJ}, BANK1 {B.OBJ}
SEGMENTS(?PR?FUNC2?B (B0:), ?PR?FUNC1?A (B1:0x8000))
```

You can also specify that a segment should be placed as last segment in a memory class by using the LAST keyword as address specification. The following example places the segment ?DT?A as last segment in the DATA memory class:

```
LX51 A.OBJ, B.OBJ, C.OBJ SEGMENTS (?DT?B(LAST))
```

The prefix '^' before the address specifies the end address for a segment. The following command places the segment PPR?SUM?B in memory so that it ends at address 0xFF8000.

```
L251 A.OBJ, B.OBJ SEGMENTS (?PR?SUM?B(^0xFF8000))
```

Next, the segment ?PR?FUNC1?A is assigned an execution address of 0x4000 and a storage address of 0xFF8000.

```
L251 A.OBJ SEGMENTS (?PR?FUNC1?A(0x4000)[0xFF8000])
```

The last example uses only an exclamation point as *store_address*. This means that no memory is reserved at address 0x2000 and the section will be stored within the address range of the SROM memory class.

L251 A.OBJ SEGMENTS (?PR?FUNC1?A(0x2000)[!])

SEGSIZE

Restriction: This control is available in **LX51** and **L251** only.

Abbreviation: SEGSIZE

Arguments: SEGSIZE (segname (size) [, ...]})

μVision2 Control: Options – Lx51 Misc – Misc Controls: enter the control.

Description: The **SEGSIZE** directive allows you to specify the memory

space used by a segment. For BIT segments the *size* may be specified in bits with the "." operator. The *segname* is any

segment contained in the input modules.

The *size* specifies the change of the segment size or segment length. There are three ways of specifying this value:

• '+' indicates that the value should be added to the current segment length.

• '-' indicates that the value should be subtracted from the current segment length.

 No sign indicates that the value should become the new segment length.

See Also: SEGMENTS

Example: L251 MYPROG.OBJ SEGSIZE (?STACK (+200H))

STACK

Restriction: This control is available in **BL51 only**. For **LX51** and **L251**

use the CLASSES and SEGMENTS control.

Abbreviation: ST

Arguments: STACK (segname [(address)]], ...])

μVision2 Control: Options – BL51 Locate – Stack.

Description: The **STACK** control locates segments in the uppermost

IDATA memory space. The segments specified will follow

all other segments in the internal data memory space.

Both, the Cx51 compiler and the PL/M-51 compiler generate a stack segment named ?STACK which is

automatically located at the top of the IDATA memory. The stack pointer is initialized by the startup code to point to the start of this segment. All return addresses and data that are pushed are stored in this memory area. It is not necessary to specifically locate this ?STACK segment. The STACK control is usually used with assembly programs which might

have several stack segments.

NOTE

Use extreme caution when relocating the ?STACK segment. It might result in a target program that does not run since data or idata variables are corrupted.

See Also: DATA, IDATA, PRECEDE

Examples: The segments that you specify are located at the highest

available memory location in the internal data memory in

the order that you specify, for example:

BL51 MYPROG.OBJ, A.OBJ, B.OBJ STACK (?DT?A,?DT?B)

You can also specify the memory location. This example places the PDT?A and PDT?B segments at 69h and 73h:

BL51 MYPROG.OBJ, A.OBJ, B.OBJ STACK (?DT?A(69h), ?DT?B(73h))

XDATA

Restriction: This control is available in **BL51 only**. For **LX51** and **L251**

use the CLASSES and SEGMENTS control.

Abbreviation: XD

Arguments: XDATA ($[address_range][segname[(address)][, ...]]$)

μVision2 Control: Options – BL51 Locate – Xdata Range, Xdata.

Description: The **XDATA** control allows you to specify:

 The starting address for segments placed in the external data space

The order of segments within the external data space

• The absolute memory location of segments in the external data space.

See Also: BIT, CODE, DATA, IDATA, PDATA

Examples: The example below specifies that relocatable XDATA segments be located in the address space 0 - 0x3FF and

0xF800 - 0xFFFF:

BL51 MYPROG.OBJ CODE(0 - 0x3FF, 0xF800 - 0xFFFF)

To specify the order for segments, you must include the names of the segments separated by commas. The following example will place the <code>?XD?MOD1</code> and <code>?XD?MOD2</code> segments at the beginning of the XDATA memory:

BL51 MOD1.OBJ, MOD2.OBJ CODE(?XD?MOD1, ?XD?MOD2)

You can also specify the memory location for a segment. The example below will place the ?xd?mod1 segment at 800h:

BL51 MOD1.OBJ, MOD2.OBJ CODE(?XD?MOD1 (0x800))

High-Level Language Controls

The Lx51 linker/locator provides controls that have to do with the high-level languages Cx51 and PL/M-51 and the real-time operating systems RTXx51. For example, you can control whether or not the Lx51 linker/locator includes automatically the run-time library and whether or overlays the local variable areas of C and PL/M-51 functions.

The following table provides an overview of these controls. For a detailed description of each control refer to the page specified in the table.

BL51	LX51, L251	Page	Description
NODEFAULT LIBRARY	NODEFAULT LIBRARY	352	Excludes modules from the run-time libraries.
NOOVERLAY	NOOVERLAY	353	Prevents overlaying or overlapping local bit and data segments.
<u>O</u> VER <u>L</u> AY	<u>O</u> VER <u>L</u> AY	354	Lets you change call references between functions and segments for data overlaying program flow analysis.
<u>R</u> E <u>C</u> URSIONS	<u>R</u> E <u>C</u> URSIONS	356	Allows you to analyze the call tree of complex recursive applications.
<u>R</u> EG <u>F</u> ILE	<u>R</u> EG <u>F</u> ILE	356	Specifies the name of the generated file to contain register usage information.
-	RTX251	358	Includes support for the RTX-251 full real-time kernel.
RTX51	RTX51	358	Includes support for the RTX-51 full real-time kernel.
RTX51TINY	RTX51TINY	358	Includes support for the RTX-51 tiny real-time kernel.
<u>SP</u> EEDOVL	-	359	Ignore during the overlay analysis references from constant segments to program code.

NODEFAULTLIBRARY

Abbreviation: NLIB

Arguments: None

Default: The run-time libraries of Cx51, RTXx51, and PL/M-51 are

searched to resolve external references in your C or PL/M

programs.

The path of the run-time libraries can be set for BL51 and LX51 with the C51LIB environment variable and for L251 with the C251LIB environment variable. This variable is the defined with the SET command that is typically entered in the AUTOEXEC.BAT batch file as shown below:

SET C51LIB=C:\KEIL\C51\LIB

In μ Vision2 the path for run-time libraries can be specified in the dialog **Project** – **File Extensions**, **Books and Environment** under environment setup – LIB folder.

If no environment variable is set, the linker tries to locate the libraries in: *path_of_the_EXE_file\..\LIB*. In a typical installation of the tool chain this sets the correct path for the run-time libraries to \C51\LIB\ or \C251\LIB. These folders contain the run-time libraries for the Cx51 compiler and the RTXx51 real-time operating system.

The libraries that are automatically searched depend on the memory model and floating-point requirements of your application. For more information refer to the *Cx51 Compiler User's Guide*.

μVision2 Control: Options – Lx51 Misc – Misc Controls: enter the control.

Description: Use the **NODEFAULTLIBRARY** control to prevent the

Lx51 linker/locator from automatically including run-time

libraries.

Example: BL51 MYPROG.OBJ NODEFAULTLIBRARY

NOOVERLAY

Abbreviation: NOOL

Arguments: None

Default: The Lx51 linker/locater is analyzing your program and

overlays the segments of local variables and function

arguments.

μVision2 Control: Options – Lx51 Misc – Misc Controls: enter the control.

Description: The **NOOVERLAY** control disables the overlay analysis

and the data overlaying. When this control is specified, the Lx51 linker/locator does not overlay the data space of local

variables and function arguments.

Examples: LX51 MYPROG.OBJ NOOVERLAY

OVERLAY

Abbreviation: OL

Arguments: OVERLAY (sfname $\{! \mid \sim\}$ sfname [, ...])

OVERLAY (sfname $\{! \mid \sim\}$ (sfname, sfname [, ...]), ...

OVERLAY (sfname!*)
OVERLAY (*!sfname)

Default: The Lx51 linker/locater is analyses the call tree of your

program and assumes normal program flow without indirect

calls via function pointers.

μVision2 Control: Options – **Lx51** Misc – Overlay.

Description: The **OVERLAY** control allows you to modify the call tree

as it is recognized by the **Lx51** linker/locater in the overlay analysis. Adjustments to the program call tree are typically required when your application uses function pointers or contains virtual program jumps as it is the case in the scheduler of a real-time operating system. The different

forms of the overlay control are shown below:

Control Specification	Description
OVERLAY (* ! sfname)	Add new root for sfname.
OVERLAY (sfname!*)	Exclude <i>sfname</i> from the overlay analysis and locate data & bit segments in non-overlaid memory. This does not influence data overlaying of other functions.
OVERLAY (sfname! sfname1) OVERLAY (sfname! (sfname1, sfname2))	Add <i>virtual</i> call references to segments or functions.
OVERLAY (sfname ~ sfname1) OVERLAY (sfname ~ (sfname1, sfname2))	Ignore call references between segments or functions.

sfname can be the name of a function or a segment.

Refer to "Using the Overlay Control" on page 282 for program examples that require the OVERLAY control.

Examples:

If your application uses a real-time operating system, each task function might be an own program path or root and the call tree of that task must be independently analyzed. This is required since the task can be terminated (i.e. by a timeout) and a previously terminated task becomes running again. In the following example **Lx51** handles the functions **TASKO** and **TASK1** as independent programs or call trees:

Identify tasks of a real-time OS

LX51 SAMPLE.OBJ OVERLAY (* ! TASKO, TASK1)

NOTE

Task functions of the RTXx51 real-time operating system are automatically handled as in depended program roots. The OVERLAY control is not required for RTXx51 tasks.

Exclude a function from overlaying

In the next example, the local data and bit segments of **FUNC1** are excluded from data overlaying. This does not influence data overlaying of other functions.

BL51 SAMPLE.OBJ OVERLAY (FUNC1 ! *)

You may add virtual references or functions calls for between segments or functions. In the following example, Lx51 *thinks* during the overlay analysis that function FUNC1 calls FUNC2 and FUNC3 even when no real calls exist.

Add virtual function calls

BL51 CMODUL1.OBJ OVERLAY (FUNC1 ! (FUNC2, FUNC3))

You may delete or remove references between segments or functions. The next example Lx51 ignores during overlay analysis the references to the <code>?PR?MAINMOD</code> segment from <code>FUNC1</code> and <code>FUNC2</code>:

Ignore references or function calls

BL51 MAINMOD.OBJ, TEXTOUT.OBJ &
OVERLAY (FUNC1 ~ ?PR?MAINMOD, FUNC2 ~ ?PR?MAINMOD)

RECURSIONS

Abbreviation: RC

Arguments: RECURSIONS (number of recursions)

Default: RECURSIONS (10)

μVision2 Control: Options – Lx51 Misc – Misc Controls: enter the control.

Description: The **RECURSIONS** control allows you to specify the

number of recursions that are allowed before the Lx51

linker/locator responds with:

FATAL ERROR 232: APPLICATION CONTAINS TOO MANY RECURSIONS.

Each time the linker encounters a recursive call during the overlay analysis of the application, this recursive call is automatically removed from the call tree and the overlay analysis is restarted. You might increase the number of accepted recursions on very complex recursive applications. However this might increase significantly the execution time of the Lx51 linker/locater.

If your application contains many pointer to function tables, you might receive the FATAL ERROR 232 before you have corrected the call tree with the OVERLAY control. The **RECURSIONS** control allows you in such situations to analysis the OVERLAY MAP of your application. Refer to "Pointer to a Function in Arrays or Tables" on page 286 for more information on correcting the call tree.

See Also: OVERLAY, SPEEDOVL

Example: LX51 MYPROG.OBJ, A.OBJ, B.OBJ RECURSIONS (100)

REGFILE

Abbreviation: RF

Arguments: REGFILE (*filename*)

Default: No register usage file is generated.

μVision2 Control: Options – Cx51 Compiler – Global Register Coloring.

Description: The **REGFILE** control allows you to specify the name of

the register usage file generated by the **Lx51** linker/locator. The information in this file is used for global register optimization by the **Cx51** compiler. The register usage information allows the **Cx51** compiler to optimize the use of

registers when calling external functions.

Example: In this example, the LX51 linker/locator generates the file

MYPROG.REG that contains register usage information.

LX51 MYPROG.OBJ, A.OBJ, B.OBJ REGFILE (MYPROG.REG)

RTX251, RTX51, RTX51TINY

Abbreviation: None

Arguments: None

Default: None

μVision2 Control: Options – Target – Operating System.

Description: These controls specify to the Lx51 linker/locator that the

application should be linked for use with the RTXx51 realtime multitasking operating system. This involves resolving references within your program to RTXx51 functions found

in the library of the real-time operating system.

The control that you should use, depends on the real-time operating system that you are using in your application:

Control	Real-Time Operating System used
RTX251	RTX251 Full Multitasking RTOS.
RTX51	RTX51 Full Multitasking RTOS.
RTX51TINY	RTX51 Tiny Multitasking RTOS.

Examples: Linker/Locater invocation for RTX51 Full:

BL51 RTX EX1.OBJ RTX51

Linker/Locater invocation for RTX251 Full:

L251 RTX_EX1.OBJ RTX251

SPEEDOVL

Restriction: This control is available in **BL51 only**. **LX51** and **L251**

always ignore references from constant segments to program

code during the overlay analysis.

Abbreviation: SP

Arguments: None

Default: BL51 does not ignore the references from constant segments

to program code during the overlay analysis.

μVision2 Control: Options – **BL51** Misc – Misc Controls: enter the control.

Description: The **RECURSIONS** control instructs **BL51** to ignore

references from constant segments to program code during the overlay analysis. This improves the execution speed of the **BL51** linker/locater during overlay analysis and is useful for applications that are using "Pointer to a Function in Arrays or Tables" as described on page 286. However, the

usage of the SPEEDOVL control makes the linker incompatible to existing applications that are using the **OVERLAY** control to correct the call tree of the

nnlication

application.

See Also: OVERLAY, SPEEDOVL

Example: BL51 MYPROG.OBJ SPEEDOVL

Error Messages

The **Lx51** linker/locator generates error messages that describe warnings, non-fatal errors, fatal errors, and exceptions.

Fatal errors immediately abort the Lx51 linker/locator operation.

Errors and warnings do not abort the **Lx51** linker/locator operation; however, they may result in an output module that cannot be used. Errors and warnings generate messages that may or may not have been intended by the user. The listing file can be very useful in such an instance. Error and warning messages are displayed in the listing file as well as on the screen.

This section displays all the **Lx51** linker/locator error messages, causes, and any recovery actions.

Warnings

Warning	Warning Message and Description
1	UNRESOLVED EXTERNAL SYMBOL SYMBOL: external-name MODULE: filename (modulename) The specified external symbol, requested in the specified module, has no corresponding PUBLIC symbol in any of the input files.
2	REFERENCE MADE TO UNRESOLVED EXTERNAL SYMBOL: external-name MODULE: filename (modulename) ADDRESS: code-address The specified unresolved external symbol is referenced at the specified code address.
3	ASSIGNED ADDRESS NOT COMPATIBLE WITH ALIGNMENT SEGMENT: segment-name The address specified for the segment is not compatible with the alignment of the segment declaration.
4	DATA SPACE MEMORY OVERLAP FROM: byte.bit address TO: byte.bit address The specified area of the on-chip data RAM is occupied by more than one segment.

Warning	Warning Message and Description
5	CODE SPACE MEMORY OVERLAP FROM: byte address TO: byte address The specified area of the code memory is occupied by more than one segment.
6	XDATA SPACE MEMORY OVERLAP FROM: byte address TO: byte address The specified area of the external data memory is occupied by more than one segment.
7	MODULE: filename (modulename) The specified module name is used for more than one module. The specified module name is not processed.
8	MODULE: filename (modulename) The specified module name is requested in the invocation line of another file that has not yet been processed. The specified module name is not processed.
9	EMPTY ABSOLUTE SEGMENT MODULE: filename (modulename) The specified module contains an empty absolute segment. This segment is not located and may be overlapped with another segment without any additional message.
10	CANNOT DETERMINE ROOT SEGMENT The Linker/Locator has recognized the C51 compiler or PL/M-51 input files and tries to process a flow analysis. However, it is impossible to determine the root segment. This error occurs if the main program is called by an assembly module. In this case, the available references (calls) must be modified with the OVERLAY control.
11	CANNOT FIND SEGMENT OR FUNCTION NAME NAME: overlay-control-name A segment or function name defined in the OVERLAY control cannot be found in the object modules.
12	NO REFERENCE BETWEEN SEGMENTS SEGMENT1: segment-name SEGMENT2: segment-name An attempt was made to delete a reference or call between two non-existent functions or segments, with the OVERLAY control.

Warning	Warning Message and Description
13	RECURSIVE CALL TO SEGMENT SEGMENT: segment-name CALLER: segment-name The specified segment is called recursively from CALLER specified segments. Recursive calls are not allowed in C51 and PL/M-51 programs.
14	INCOMPATIBLE MEMORY MODEL MODULE: filename (modulename) MODEL: memory model The specified module is not compiled in the same memory model as the former compiled modules. The memory model of the improper module is showed by MODEL.
15	MULTIPLE CALL TO SEGMENT SEGMENT: segment-name CALLER1: segment-name CALLER2: segment-name The specified segment is called from two levels, CALLER1, and CALLER2; e.g., main and interrupt program. This has the same effect as a recursive call and may thus lead to the overwriting of parameters or data.
16	UNCALLED SEGMENT, IGNORED FOR OVERLAY PROCESS SEGMENT: segment-name This warning occurs when functions, which were not previously called, are contained in a program (e.g., for test purposes). The function specified is excluded from the overlay process in this case. It is possible that the program then occupies more memory as during a call of the specified segment.
17	INTERRUPT FUNCTION IN BANKS NOT ALLOWED SYMBOL: function-name SPACE: code-bank The specified C function is an interrupt function (a C51 function) that was specified to be located in a code bank. Interrupt functions cannot be located in a code bank.
18	no generated by Lx51
19	COMMON CODE SEGMENTS LOCATED TO BANKED AREA Some segments that are usually located to the common area located into the banked area. This warning just informs you, that you might free up some code space by locating program code into banks. The warning is not generated for the default setting of the BANKAREA (0 - 0xFFFF).
20	L51_BANK.A51: NBANKS < NUMBER OF CODE BANKS The setting for NBANKS in the L51_BANK.A51 module is smaller than the number of banks used in your application.

Warning	Warning Message and Description
21	SEGMENT LOCATED OUTSIDE BANKED AREA
	A segment that should be located in the code banking area is located outside the address range of the BANKAREA directive.
22	SEGMENT SIZE UNDERFLOW: OLD SIZE + CHANGE < 0
	SEGMENT: segment-name The size change specified in the SEGSIZE control causes the segment size to be less than zero.
23	UNRESOLVED EXTERNAL SYMBOL DURING LINK PROCESS During the link run one or more external symbols have no corresponding PUBLIC symbol in any of the input files.
24	INCOMPATIBLE CPU MODE MODULE: module-name
	MODE: cpu-mode
	The specified module is not translated with the same CPU mode as the former Lx51 input modules. The CPU mode of the invalid module is displayed by MODE. The CPU mode of other input modules is displayed in the Lx51 listing file.
25	DATA TYPES DIFFERENT
	SYMBOL: symbol-name MODULE: module-name
	The definition of the specified symbol in the specified module is not identical with the public definition of that symbol. The module which contains the public symbol can be determined with the IXREF listing. This warning is disabled with WARNINGLEVEL (0) control.
26	DATA TYPES SLIGHTLY DIFFERENT
	SYMBOL: symbol-name MODULE: module-name
	The definition of the specified symbol in the specified module is not 100% identical with the public definition of that symbol. This warning is the result when unsigned signed mismatches occur, i.e. unsigned char does not match char. The module which contains the public symbol can be determined with the IXREF listing. This warning is disabled with WARNINGLEVEL (1) control.
27	INCOMPATIBLE INTERRUPT FRAME SIZE MODULE: module-name
	FRAME: frame-size
	The specified module is not translated with the same interrupt frame size assumptions as the former input modules. The frame size of the invalid module is displayed by FRAME. The frame size of other input modules is displayed in the Lx51 listing file.

Warning	Warning Message and Description
28	DECRESING SIZE OF SEGMENT SEGMENT: segment-name The size specified in the SEGSIZE control has caused Lx51 to decrease the size of the specified segment.
29	SEGMENT LOCATED OUTSIDE CLASS AREA SEGMENT: segment-name The specified segment is located outside the memory class limits specified by the CLASSES control.
30	MEMORY SPACE OVERLAP FROM: address TO: address The specified area of the physical memory is occupied by more than one segment.

Non-Fatal Errors

Error	Error Message and Description
101	SEGMENT: segment-name MODULE: filename (modulename) The attributes of the specified partial segment in the specified module cannot be combined with the attributes of the previous defined partial segments of the same name. The partial segment is ignored.
102	EXTERNAL ATTRIBUTE MISMATCH SYMBOL: external-name MODULE: filename (modulename) The attributes of the specified external symbol in the specified module do not match the attributes of the previously defined external symbols. The specified symbol is ignored.
103	EXTERNAL ATTRIBUTE DO NOT MATCH PUBLIC SYMBOL: public-name MODULE: filename (modulename) The attributes of the specified public symbols in the specified module do not match the attributes of the previous defined external symbols. The specified symbol is ignored.
104	MULTIPLE PUBLIC DEFINITIONS SYMBOL: public-name MODULE: filename (modulename) The specified public symbol in the specified module has already been defined in a previously processed file.
105	PUBLIC REFERS TO IGNORED SEGMENT SYMBOL: public-name SEGMENT: segment-name The specified public symbol is defined in the specified segment. It cannot be processed on account of an error. The public symbol is therefore ignored.
106	SEGMENT OVERFLOW SEGMENT: segment-name The specified segment is longer than the limits implied by the memory class to which the segment belongs to.
107	ADDRESS SPACE OVERFLOW SPACE: space-name SEGMENT: segment-name The specified segment cannot be located at the specified address space. The segment is ignored.

Error	Error Message and Description
108	SEGMENT IN LOCATING CONTROL CANNOT BE ALLOCATED SEGMENT: segment-name The specified segment in the invocation line cannot be processed on account of its attributes.
109	EMPTY RELOCATABLE SEGMENT SEGMENT: segment-name The specified segment after combination has a zero size. The specified segment is ignored.
110	CANNOT FIND SEGMENT SEGMENT: segment-name The specified segment is contained in the invocation line but cannot be found in an input module. The specified segment is ignored.
111	SPECIFIED BIT ADDRESS NOT ON BYTE BOUNDARY SEGMENT: segment-name The specified segment contained in the BIT control is a DATA segment. The specified BIT address however is not on a byte boundary. The segment is ignored.
112	SEGMENT TYPE NOT LEGAL FOR COMMAND SEGMENT: segment-name The specified segment cannot be processed because it does not have a legal type.
113	SEGMENT IN LOCATING CONTROL IS ALREADY ABSOLUTE SEGMENT: segment-name The specified segment is already an absolute segment and cannot be located with the SEGMENTS control.
114	SEGMENT DOES NOT FIT SPACE: space-name SEGMENT: segment-name BASE: base-address LENGTH: segment-length The specified segment cannot be located at the base address in the specified address space because of its length. The segment is ignored.
115	INPAGE SEGMENT IS GREATER THAN 256 BYTES SEGMENT: segment-name The specified segment with the attributes PAGE or INPAGE is greater than 256 bytes. The segment is ignored.

Error	Error Message and Description
116	INBLOCK SEGMENT IS GREATER THAN 2048 BYTES SEGMENT: segment-name The specified segment with the attribute INBLOCK is greater than 2048 bytes. The segment is ignored.
117	BIT ADDRESSABLE SEGMENT IS GREATER THAN 16 BYTES SEGMENT: segment-name The specified bit or data segment that was declared with the BITADDRESSABLE attribute is larger than 16 bytes. The segment is not ignored.
118	REFERENCE MADE TO ERRONEOUS EXTERNAL SYMBOL: symbol-name MODULE: filename (modulename) ADDRESS: code-address The specified external symbol that was erroneously processed, is referenced in the specified code address.
119	REFERENCE MADE TO ERRONEOUS SEGMENT SEGMENT: symbol-name MODULE: filename (modulename) ADDRESS: code-address The specified segment processed with an error, is referenced in the specified code address.
120	CONTENT BELONGS TO ERRONEOUS SEGMENT SEGMENT: segment-name MODULE: filename (modulename) A specified segment that was erroneously processed, is referenced at a specific code address. The segment contents are not available.
121	IMPROPER FIXUP MODULE: filename (modulename) SEGMENT: segment-name OFFSET: segment-address After evaluation of absolute fix-ups, an address is not accessible. The improper address along with the specific module name, partial segment, and segment address are displayed. The fix-up command is not processed. This error occurs when an instruction cannot reach the address, i.e. ACALL
	instruction calls a location outside the 2KB block. If you are working with the Cx51 compiler, you have typically selected the ROM(SMALL) option for a program that exceeds the 2KB ROM size. You can locate the instruction, when you open the LST file of the translator and search for the instruction that is located in the offset of the specified segment.
122	CANNOT FIND MODULE MODULE: filename (modulename) The module specified in the invocation line cannot be found in the input file.

Error	Error Message and Description
123	ABSOLUTE DATA/IDATA SEGMENT DOES NOT FIT MODULE: filename (modulename) FROM: byte address TO: byte address An absolute DATA or IDATA segment contained in the specified module is not permissible due to a conflict with the value specified with the RAMSIZE control. The absolute segment cannot be located in the area, which was output.
124	BANK SWITCH MODULE INCORRECT This error message is issued when the bank switch module file (L51_BANK.OBJ) contains invalid information or is not specified.
125	DUPLICATE TASK NUMBER TASK1: function name TASK2: function name TASKID: task-id A task number has been assigned to more than one RTXx51 task function.
126	TASK WITH PRIORITY 3 CANNOT WORK WITH REGISTERBANK 0 TASK: function name TASKID: task-id A task that has priority 3 must have a using attribute that refers to register bank 1, 2, or 3.
127	UNRESOLVED EXTERNAL SYMBOL SYMBOL: external-name MODULE: filename (modulename) The specified external symbol, requested in the specified module, has no corresponding PUBLIC symbol in any of the input files.
128	REFERENCE MADE TO UNRESOLVED EXTERNAL SYMBOL: external-name MODULE: filename (modulename) ADDRESS: code-address The specified unresolved external symbol is referenced at the specified code address.
129	TASK REQUIRES REGISTERBANK TASK: function name TASKID: task-id The task function requires that you assign a register bank with an using attribute.
130	NO MATCHING SEGMENT FOR WILDCARD SEGMENT NAME SEGMENT: segment-name The linker could not find a segment name that matches the wildcard segment name stated in the command line.

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Error	Error Message and Description		
131	ILLEGAL PRIORITY FOR TASK TASK: function name TASKID: task-id You have assigned a priority for an RTX51 Tiny task. RTX51 Tiny does not support priorities.		
132	ILLEGAL TASKID: RTX-51 TINY SUPPORTS ONLY 16 TASKS TASK: function name TASKID: task-id You have assigned a task-id that is higher than 15. RTX51 Tiny tasks supports only 16 tasks.		
133	SFR SYMBOL HAS DIFFERENT VALUES SYMBOL: public-name MODULE: filename (modulename) The specified SFR symbol is defined with different values in several input modules.		
134	ADDRESS SPACE OVERFLOW IN BANKAREA SPACE: space-name SEGMENT: segment-name The specified segment cannot be located in the banked area, since the banked area is already full.		

Fatal Errors

Error	Error Message and Description
201	INVALID COMMAND LINE SYNTAX A syntax error is detected in the command line. The command line is displayed up to and including the point of error.
202	INVALID COMMAND LINE, TOKEN TOO LONG The command line contains a token that is too long. The command line is displayed up to and including the point of error.
203	EXPECTED ITEM MISSING An expected item is missing in the command line. The command line is displayed up to and including the point of error.
204	INVALID KEYWORD The invocation line contains an invalid keyword. The command line is displayed up to and including the point of error.
205	CONSTANT TOO LARGE A constant in the invocation line is larger than 0FFFFH. The command line is displayed up to and including the point of error.
206	INVALID CONSTANT A constant in the invocation line is invalid; e.g., a hexadecimal number with a leading letter. The command line is displayed up to and including the point of error.
207	INVALID NAME A module or segment name is invalid. The command line is displayed up to and including the point of error.
208	INVALID FILENAME A filename is invalid. The command line is displayed up to and including the point of error.
209	FILE USED IN CONFLICTING CONTEXTS FILE: filename A specified filename is used for multiple files or used as an input as well as an output file.
210	I/O ERROR ON INPUT FILE: system error message FILE: filename An I/O error is detected by accessing an input file. A detailed error description of the EXCEPTION messages is described afterwards.

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Error	Error Message and Description
211	I/O ERROR ON OUTPUT FILE: system error message FILE: filename An I/O error is detected by accessing an output file. A detailed error description of the EXCEPTION messages is described afterwards.
212	I/O ERROR ON LISTING FILE: system error message FILE: filename An I/O error is detected by accessing a listing file. A detailed error description of the EXCEPTION messages is described afterwards.
213	I/O ERROR ON WORK FILE: system error message An I/O error is detected by accessing a temporary work file of BL51. A detailed error description of the EXCEPTION messages is described afterwards.
214	INPUT PHASE ERROR MODULE: filename (modulename) This error occurs when BL51 encounters different data during pass two. This error could be the result of an assembly error.
215	CHECK SUM ERROR MODULE: filename (modulename) The checksum does not correspond to the contents of the file.
216	INSUFFICIENT MEMORY The memory available for the execution of BL51 is used up.
217	NO MODULE TO BE PROCESSED No module to be processed is found in the invocation line.
218	NOT AN OBJECT FILE FILE: filename The specified file is not an object file.
219	NOT AN 8051/X51 OBJECT FILE FILE: filename The specified file is not a valid x51 object file.
220	INVALID INPUT MODULE FILE: filename The specified input module is invalid. This error could be the result of an assembler error.

Error	Error Message and Description
221	MODULE SPECIFIED MORE THAN ONCE The invocation line contains the specified module more than once. The command line is displayed up to and including the point of error.
222	SEGMENT SPECIFIED MORE THAN ONCE The invocation line contains the specified segment more than once. The command line is displayed up to and including the point of error.
224	DUPLICATE KEYWORD OR CONFLICTING CONTROL The same keyword is contained in the invocation line more than once or contradicts with other keywords. The command line is displayed up to and including the point of error.
225	SEGMENT ADDRESS ARE NOT IN ASCENDING ORDER The base addresses for the segments are not displayed in ascending order during the location control. The command line is displayed up to and including the point of error.
226	SEGMENT ADDRESS INVALID FOR CONTROL The base addresses for the segments are invalid for the location control. The command line is displayed up to and including the point of error.
227	PARAMETER OUT OF RANGE The specified value for the PAGEWIDTH or PAGELENGTH control is out of the acceptable range. The command line is displayed up to and including the point of error.
228	RAMSIZE PARAMETER OUT OF RANGE The specified value for the RAMSIZE control is out of the acceptable range. The command line is displayed up to and including the point of error.
229	INTERNAL PROCESS ERROR Lx51 detects an internal processing error. Please contact your dealer.
230	START ADDRESS SPECIFIED MORE THAN ONCE The invocation line contains more than one start address for unnamed segment group. The command is displayed up to and including the point of error.
231	ADDRESS RANGE FOR BANKAREA INCORRECT The address space specified with the BANKAREA control is invalid.
232	APPLICATION CONTAINS TOO MANY RECURSIONS The application contains to many recursive calls. Refer to "RECURSIONS" on page 356 for more information.

Error	Error Message and Description
233	ILLEGAL USE OF * IN OVERLAY CONTROL The use of "*! *" or "* ~ *" with the OVERLAY control is illegal.
234	USE RTX251 OR RTX51 CONTROL The application uses RTXx51 tasks.
233	ILLEGAL USE OF * IN OVERLAY CONTROL command line The use of "*! *" or "* ~ *" with the OVERLAY control is illegal.
234	USE RTX-251 SWITCH The application uses a real-time operating system RTX251 Full or RTX251 Tiny. The L251 linker/locater must be invoked with the RTX251 or RTX251TINY control.
235	TOO MANY ADDRESS RANGES You are using to many address ranges.
236	ADDRESSES ARE NOT IN ASCENDING ORDER The address range does not contain addresses in ascending order.
237	INVALID CLASS NAME The class name given in the CLASSES control is not valid.
238	BIT ADDRESS INVALID FOR THIS CLASS TYPE The CLASSES control contains a bit address for a memory class which cannot be used for bit objects.
239	BASE ADDRESS ALREADY GIVEN FOR THIS CLASS The CLASSES control contains a base address, but the class has already a base address specified with a previous CLASSES control.
240	BASE ADDRESS MUST BE THE FIRST ARGUMENT The base address must be the first argument in the CLASSES control.
241	BASE ADDRESS CANNOT BE GIVEN FOR THIS CLASS A base address cannot be given for this memory class in the CLASSES control.
242	WRONG SYNTAX FOR THE EXECUTION ADDRESS The execution address field contains a wrong syntax.
243	EXECUTION ADDRESS REQUIRED IF SPACE IS NOT RESERVED You need to specify an execution address, if the execution space should not be reserved.

Error	Error Message and Description
244	OVERLAPPING CLASS RANGE The address ranges in the classes control are overlapping.
245	ADDRESS RANGE INVALID FOR THIS CLASS TYPE The address range given in the CLASSES control is not valid for this memory class type.
246	SYMBOL SPECIFIED MORE THAN ONCE The symbol name is already used.
249	MODULE USES AN UNKNOWN OMF VERSION MODULE: filename (modulename) The module uses an un-known or unsupported OMF version.
250	CODE SIZE LIMIT IN RESTRICTED VERSION EXCEEDED You are using modules that are created with an evaluation version or a code size limited version and the size limit is exceeded.
251	RESTRICTED MODULE IN LIBRARY NOT SUPPORTED A library contains a module that is created with an evaluation version or a code size limited version. This is not supported.

Exceptions

Exception messages are displayed with some error messages. The **BL51** linker/locator exception messages that are possible are listed below:

Exception	Exception Message and Description
0021H	PATH OR FILE NOT FOUND The specified path or filename is missing.
0026Н	ILLEGAL FILE ACCESS An attempt was made to write to or delete a write-protected file.
0029Н	ACCESS TO FILE DENIED The file indicated is a directory.
002AH	I/O-ERROR The drive being written to is either full or the drive was not ready.
0101H	ILLEGAL CONTEXT An attempt was made to access a file in an illegal context; e.g., the printer was opened for reading.

Chapter 10. Library Manager

The **LIBx51** library manager allows you to create and maintain library files. A library file is a formatted collection of one or more object files. Library files provide a convenient method of referencing a large number of object files and can be used by the **Lx51** linker/locator. The **LIBx51** library manager can be controlled interactively or from the command line.

The following table gives you an overview of the LIBx51 library manager variants along with the translators that are supported.

Library Manager	Processes Files from	Description
LIB51	Keil A51 Macro Assembler Keil C51 Compiler Intel ASM51 Assembler Intel PL/M51 Compiler	For classic 8051 . Includes support for 32 x 64KB code banks.
LIBX51	Keil A51 Macro Assembler Keil C51 Compiler Keil AX51 Macro Assembler Keil CX51 Compiler for 80C51MX Intel ASM51 Assembler Intel PL/M51 Compiler	For classic 8051 and extended 8051 versions (Philips 80C51MX, Dallas 390, etc.) Allows code and data banking and supports up to 16MB code and xdata memory.
LIB251	Keil A51 Macro Assembler Keil C51 Compiler Keil A251 Macro Assembler Keil C251 Compiler Intel ASM51 Assembler Intel PL/M51 Compiler	For Intel/Atmel WM 251.

Using LIBx51

To invoke the **LIBx51** library manager from the command prompt, type the program name along with an optional command. The format for the **LIBx51** command line is:

LIB51	[command]
LIBX51	[command]
LIB251	command

or

LIB51	@commandfile
LIBX51	@commandfile
LIB251	@commandfile

where

command may be a single library manager command.

commandfile is the name of a command input file that may contain a very long

library manager command.

Interactive Mode

If no *command* is entered on the command line, or if the ampersand character is included at the end of the line, the LIB51 library manager enters interactive mode. The **LIBx51** library manager displays an asterisk character (*) to signal that it is in interactive mode and is waiting for input.

Any of the **LIBx51** library manager commands may be entered on the command line or after the * prompt when in interactive mode.

Type **EXIT** to leave the **LIBx51** library manager interactive mode.

Create Library within µVision2

You can directly create a library file from your μVision2 project. Select **Create Library** in the dialog **Options for Target – Output**. μVision2 will call the correct **LIBx51** Library Manager instead of the **Lx51** Linker/Locater. Since the

code in the Library will be not linked and located, the entries in the L51 Locate and L51 Misc options page are ignored.

Command Summary

The following table lists the commands that are available for the LIBx51 library manager. The usage and the syntax of these commands are described in the sections that follow.

NOTEUnderlined characters denote the abbreviation for the particular command.

LIBx51 Command	Description
ADD	Adds an object module to the library file. For example, LIB51 ADD GOODCODE.OBJ TO MYLIB.LIB adds the GOODCODE.OBJ object module to MYLIB.LIB.
<u>C</u> REATE	Creates a new library file. For example, LIB251 CREATE MYLIB.LIB creates a new library file named MYLIB.LIB.
<u>D</u> ELETE	Removes an object module from the library file. For example, LIBX51 DELETE MYLIB.LIB (GOODCODE) removes the GOODCODE module from MYLIB.LIB.
E <u>X</u> TRACT	Extracts an object module from the library file. For example, LIB251 EXTRACT MYLIB.LIB (GOODCODE) TO GOOD.OBJ copies the GOODCODE module to the object file GOOD.OBJ.
<u>E</u> XIT	Exits the library manager interactive mode.
<u>H</u> ELP	Displays help information for the library manager.
<u>L</u> IST	Lists the module and public symbol information stored in the library file. For example, LIB251 LIST MYLIB.LIB TO MYLIB.LST PUBLICS generates a listing file (named MYLIB.LST) that contains the module names stored in the MYLIB.LIB library file. The PUBLICS directive specifies that public symbols are also included in the listing.
REPLACE	Replaces an existing object module to the library file. For example, LIB51 REPLACE GOODCODE.OBJ IN MYLIB.LIB replaces the GOODCODE.OBJ object module in MYLIB.LIB. Note that Replace will add GOODCODE.OBJ to the library if it does not exist.
TRANSFER	Generates a complete new library and adds object modules. For example, LIB251 TRANSFER FILE1.0BJ, FILE2.0BJ TO MYLIB.LIB deletes the existing library MYLIB.LIB, re-creates it and adds the object modules FILE1.0BJ and FILE2.0BJ to that library.

Creating a Library

The **CREATE** command creates a new, empty library file and has the following format:

CREATE libfile

libfile is the name of the library file to create and should include a file extension. Usually, **LIB** is the extension that is used for library files.

Example:

```
LIBX51 CREATE MYFILE.LIB

* CREATE FASTMATH.LIB
```

The **TRANSFER** command creates a new library file and adds object modules. The **TRANSFER** command must be entered in the following format:

TRANSFER filename	[(modulename,)] [,] TO libfile
where	
filename	is the name of an object file or library file. You may specify several files separated by a comma.
modulename	is the name of a module in a library file. If you do not want to add the entire contents of a library, you may select the modules that you want to add. Module names are specified immediately following the filename, must be enclosed in parentheses, and must be separated by commas.
libfile	is the name of the library file that should be created. The LIBx51 library manager will remove a previous version of the library, if this file already exists. The specified object modules are added to the new created library.

Example:

```
LIB251 TRANSFER FILE1.OBJ, FILE2.OBJ TO MYLIB.LIB

LIBX51 @mycmd.lin

--- content of mycmd.lin: ---

TRANSFER FILE1.OBJ, FILE2.OBJ, FILE3.OBJ TO MYLIB.LIB
```

Adding or Replacing Object Modules

The **ADD** command is used to add one or more object modules to an existing library file. The **ADD** command must be entered in the following format:

ADD filename (modulename, ...) | [, ...] TO libfile

where

filename is the name of an object file or library file. You may specify

several files separated by a comma.

modulename is the name of a module in a library file. If you do not want

to add the entire contents of a library, you may select the modules that you want to add. Module names are specified immediately following the filename, must be enclosed in

parentheses, and must be separated by commas.

1ibfile is the name of an existing library file. The specified object

modules are added to this library.

Example:

LIB51 ADD MOD1.OBJ, UTIL.LIB(FPMUL, FPDIV) TO NEW.LIB

With the **REPLACE** command you can update an existing object module in a library file. The **REPLACE** command will the object module to the library if it does not exist. The format is:

REPLACE filename IN libfile

* ADD FPMOD.OBJ TO NEW.LIB

where

filename is the name of an object file you want to update.

libfile is the name of an existing library file. The object module is

replaced in this library.

Example:

LIBX51 REPLACE MOD1.OBJ IN MYLIB.LIB

* REPLACE FPMOD.OBJ TO FLOAT.LIB

Removing Object Modules

The **DELETE** command removes object modules from a library file. This command must be entered in the following format:

```
DELETE libfile (modulename [, modulename ...])

where
```

libfile is the name of an existing library file. The specified object

modules are removed from this library.

modulename is the name of a module in the library file that you want to

remove. Module names are entered in parentheses and are

separated by commas.

Example:

```
LIB51 DELETE NEW.LIB (MODUL1)

* DELETE NEW.LIB (FPMULT, FPDIV)
```

Extracting Object Modules

The **EXTRACT** command creates a standard object module for a specified module in a library file. This command must be entered in the following format:

```
where

libfile is the name of an existing library file. For the specified object module a standard object module will be created.

modulename is the name of a module in the library file. Only one module name can be entered in parentheses.

filename is the name of the object file that should be created from the library module.
```

Example:

LIBX51 EXTRACT FLOAT.LIB(FPMUL) TO FLOATMUL.OBJ

Listing Library Contents

Use the **LIST** command to direct the **LIBx51** library manager to generate a listing of the object modules that are stored in a library file. **LIST** may be specified on the command line or after the * prompt in interactive mode. This command has the following format:

```
where

libfile is the library file from which a module list is generated.

listfile is the file where listing information is written. If no

listfile is specified, the listing information is displayed on the screen.

PUBLICS specifies that public symbols are included in the listing. Normally, only module names are listed.
```

Example:

```
LIB251 LIST NEW.LIB

* LIST NEW.LIB TO NEW.LST PUBLICS
```

The LIBx51 library manager produces a module listing that appears as follows:

```
LIBRARY: NEW.LIB

PUTCHAR

_PUTCHAR

PRINTF

?_PRINTF517?BYTE

?_SPRINTF?BYTE

?_SPRINTF?BYTE

_PRINTF

_PRINTF

_PRINTF

_PRINTF

_PRINTF

_PRINTF

_PRINTF

_PRINTF517

_SPRINTF517

PUTS

_PUTS
```

In this example, PUTCHAR, PRINTF, and PUTS are module names. The names listed below each of these module names are public symbols found in each of the modules.

Error Messages

This chapter lists the fatal and non-fatal errors that may be generated by the LIB51 library manager during execution. Each section includes a brief description of the message, as well as corrective actions you can take to eliminate the error or warning condition.

Fatal Errors

Fatal errors cause immediate termination of the LIB51 library manager. These errors normally occur as the result of a corrupt library or object file, or as a result of a specification problem involving library or object files.

Error	Error Message and Description
215	CHECK SUM ERROR FILE: filename The checksum for filename is incorrect. This usually indicates a corrupt file.
216	INSUFFICIENT MEMORY There is not enough memory for the LIB51 library manager to successfully complete the requested operation.
217	NOT A LIBRARY FILE: filename The filename that was specified is not a library file.
219	NOT AN 8051 OBJECT FILE FILE: filename The filename that was specified is not a valid 8051 object file.
222	MODULE SPECIFIED MORE THAN ONCE MODULE: filename (modulename) The specified modulename is included on the command line more than once.

Errors

The following errors cause immediate termination of the LIB51 library manager. These errors usually involve invalid command line syntax or I/O errors.

Error	Error Message and Description
201	INVALID COMMAND LINE SYNTAX A syntax error was detected in the command. The command line is displayed up to and including the point of error.
202	INVALID COMMAND LINE, TOKEN TOO LONG The command line contains a token that is too long for the LIB51 library manager to process.
203	EXPECTED ITEM MISSING The command line is incomplete. An expected item is missing.
205	FILE ALREADY EXISTS FILE: filename The filename that was specified already exists. This error is usually generated when attempting to create a library file that already exists. Erase the file or use a different filename.
208	MISSING OR INVALID FILENAME A filename is missing or invalid.
209	UNRECOGNIZED COMMAND A command is unrecognized by the LIB51 library manager. Make sure you correctly specified the command name.
210	I/O ERROR ON INPUT FILE: system error message FILE: filename An I/O error was detected when accessing one of the input files.
211	I/O ERROR ON LIBRARY FILE: system error message FILE: filename An I/O error was detected when accessing a library file.
212	I/O ERROR ON LISTING FILE: system error message FILE: filename An I/O error was detected when accessing a listing file.

Error	Error Message and Description
213	I/O ERROR ON TEMPORARY FILE: system error message FILE: filename An I/O error was detected when a temporary file was being accessed.
220	INVALID INPUT MODULE FILE: filename The specified input module is invalid. This error could be the result of an assembler error or could indicate that the input object file is corrupt.
221	FILE SPECIFIED MORE THAN ONCE FILE: filename The filename specified was included on the command line more than once.
223	CANNOT FIND MODULE MODULE: filename (modulename) The modulename specified on the command line was not located in the object or library file.
224	ATTEMPT TO ADD DUPLICATE MODULE MODULE: filename (modulename) The specified modulename already exists in the library file and cannot be added.
225	ATTEMPT TO ADD DUPLICATE PUBLIC SYMBOL MODULE: filename (modulename) PUBLIC: symbolname The specified public symbolname in modulename already exists in the library file and cannot be added.

Chapter 11. Object-Hex Converter

Program code stored in an absolute object file can be converted an Intel HEX file. Intel HEX files may be used as input files for EPROM programmers or emulators.

For the each Lx51 linker/locater variant a different OHx51 object hex converter is required to create an Intel HEX file. For code banking applications generated with the BL51 linker/locater, you need in addition the OC51 Banked Object File Converter to create HEX files. The following table gives you an overview of the conversion tools required to create an Intel HEX file.

Output from	Object Hex Converter	Description
BL51	OH51	For classic 8051 applications without banking.
BL51 Banked Application	OC51 combined with OH51	For classic 8051 applications with banking.
LX51	OHX51	For classic 8051 and extended 8051 versions.
L251	OH251	For Intel/Atmel WM 251.

In μ Vision2 you enable the generation of an Intel Object file in the dialog **Options** – **Output** with **Create HEX File**. The μ Vision2 project manager selects automatically the correct utility.

The following sections describe how to use the **OHx51** and **OC51** utilities, the command-line options that are available, and any errors that may be encountered during execution.

Using OHx51

To invoke **OHx51** from the command prompt, type the program name along with the name of the absolute object file. The command line format for the OHx51 utilities is:

ОН51	abs_file	HEXFILE	(file)		
OHX51	abs_file	HEXFILE	(file) H386 RANGE	(start-end) OFFSET	(offset)
OH251	abs_file	HEXFILE	(file)][H386][RANGE	$(start-end) \ \ \ \ \ \ \ \ \ \ \ \ \ $	(offset)

where

is the name of the absolute object file that is generated by the abs file

Lx51 linker/locator.

file is the name of the Intel HEX file to generate. By default, the

HEX file name is the name of the abs file with the

extension .HEX.

specifies Intel HEX-386 format for the Intel HEX file. This H386

format is automatically used, if the specified address range is

more than 64KBytes.

specifies the address range of the abs_file that should be start-end

converted to the Intel HEX file. The default range depends on

the device you are using and is listed in the following table:

OHx51 Converter (Architecture)	Address Range
OHX51 (Classic, and Extended 8051) Note: For code banking applications OHX51 the default setting converts the complete content into an Intel HEX-386 format.	C:0x0000 - C:0xFFFF
OHX51 (Philips 80C51MX)	0x800000 - 0x80FFFF
OH251 (Intel/Atmel WM 251)	0xFF0000 - 0xFFFFFF

offset specifies an offset which is added to the address stored in the abs file.

OHx51 Command Line Examples

The following command generates an Intel HEX-386 file for a 251 device. The address range 0xFE0000 - 0xFFFFFF should be converted. The offset 0xFE0000 is subtracted to get an Intel HEX file that can be directly programmed into an EPROM that is mapped to the address space 0xFE0000 - 0xFFFFFF in the 251 address space.

```
OH251 MYPROG RANGE (0xFE0000-0xFFFFFF) OFFSET (-0xFE0000)
```

The next example generates an Intel HEX file for a banked application with a classic 8051 device. Only the code bank 0 should be converted. The file format used will be the standard Intel HEX format.

```
OHX51 PROG RANGE (B:0-B:0xFFFF)
```

The command below generates an Intel HEX-386 file for a Philips 80C51MX device. The **OFFSET** control is used to create an output file that can be directly programmed into an EPROM.

```
OHX51 MYAPP RANGE (0x800000-0x81FFFF) OFFSET (-0x800000)
```

With the next command line, the constants stored in the XDATA space are converted into an Intel HEX file

```
OHX51 MYPROG RANGE (X:0-X:FFFF)
```

Creating HEX Files for Banked Applications

For the **BL51** linker/locater the **OC51** Banked Object File Converter described on page 392 is used to split banked object files into standard object files that contain a 64KB code bank. These files can be converted with OH51 into HEX files that store the content of a 64KB bank. These files are programmed separately into the corresponding physical address space of the EPROM.

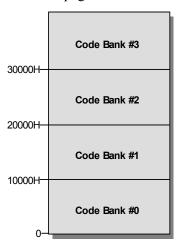
For the extended **LX51** linker/locater the **OHX51** object hex converter generates in the default setting and Intel HEX-386 file that contains the common area and all the code banks. If code bank 0 does not exist in your application, **OHX51** will skip this memory area.

Examples:

The figure below shows the HEX

The following figure shows the

file content for the example "Banking With Four 64 KByte Banks" on page 299.



HEX file content for the example "Banking With Common Area" on page 303.

38000H-	Code Bank #7
	Code Bank #6
30000H	Code Bank #5
28000H—	Code Bank #4
20000H	Code Bank #3
18000H—	Code Bank #2
10000H	Code Bank #1
8000H—	Common Area
0	

OH*x***51 Error Messages**

The following tables list error and warning messages of **OH***x***51**. Each message includes a brief description of the reason for the error or warning condition.

```
Error Message
*** ERROR, INVALID RECORD-TYPE ENCOUNTERED
The absolute object file contains an invalid record type.
*** FATAL, INCONSISTENT OBJECT FILE
The input file has an invalid format.
*** ERROR, ARGUMENT TOO LONG
An argument in the command line is too long.
*** ERROR, DELIMITER '(' AFTER PARAMETER EXPECTED
The command-line parameter must be followed by an argument enclosed in parentheses ().
*** ERROR, DELIMITER ')' AFTER PARAMETER EXPECTED
The command-line parameter must be followed by an argument enclosed in parentheses ().
*** ERROR, UNKNOWN CONTROL:
The specified command-line parameter is unrecognized.
*** ERROR, RESPECIFIED CONTROL, IGNORED
The indicated command-line control was specified twice.
*** ERROR, CAN'T OPEN FILE filename
The specified file cannot be open for read.
*** ERROR, CAN'T CREATE FILE filename
The specified file cannot be open for write.
```

Error Message

*** I/O-ERROR ON FILE filename

A read/write error occurred during access of the specified file.

*** ERROR: PREMATURE END OF FILE ON filename

The input file does not end correctly. This is usually a result of a previous fatal error of an translator or linker/locater.

*** ERROR: MORE THAN 512 CLASSES ON filename

The input file contains more than 512 memory classes. This is the limit of OHx51.

*** ERROR, NON-NULL ARGUMENT EXPECTED

An argument is missing.

Warning Message

WARNING: <PUBDEF> HEX-FILE WILL BE INVALID

The absolute object file still contains public definitions. This warning usually indicates that the object file has not been processed by the **Lx51** linker/locator. The hex file that is produced may be invalid.

WARNING: <EXTDEF> UNDEFINED EXTERNAL

The absolute object file still contains external definitions. This warning usually indicates that the object file has not been processed by the **Lx51** linker/locator. The hex file that is produced may be invalid.

WARNING: <FIXUPP> HEX-FILE WILL BE INVALID

The absolute object file still contains fix-ups. This warning usually indicates that the object file has not been processed by the L251 linker/locator. The hex file that is produced may be invalid.

Using OC51

The **BL51** linker/locator generates a banked object file for programs that use bank switching. Banked object files contain several banks of code that reside at the same physical location. These object files can be converted into standard object files using the **OC51** Banked Object File Converter.

The **OC51** Banked Object File Converter will create an absolute object file for each code bank represented in the banked object file. Symbolic debugging information that was included in the banked object file will be copied to the absolute object modules that are generated. Once you have created absolute object files with OC51, you may use the OH51 Object-Hex Converter to create Intel HEX files for each absolute object file.

The OC51 Banked Object File Converter is invoked from the command prompt by typing OC51 along with the name of the banked object file. The command line format is:

OC51 banked obj file

where

banked_obj_file is the name of the banked object file that is generated by the **BL51** linker/locator.

OC51 will create separate absolute object modules for each code bank represented in the banked object file. The absolute object modules will be created with a filename consisting of the *basename* of the banked object file combined with the file extension Bnn where nn corresponds to the bank number 00 - 31. For example:

OC51 MYPROG

creates the absolute object files MYPROG.B00 for code bank 0, MYPROG.B01 for code bank 1, MYPROG.B02 for code bank 2, etc.

NOTE

Use the **OC51** Banked Object File Converter only if you used the **BANKx** control on the **BL51** linker/locator command line. If your program does not use code banking, do not use **OC51**.

OC51 Error Messages

The following table lists the errors that you may encounter when you use the OC51 Banked Object File Converter. Each message includes a brief description of the message as well as corrective actions you can take to eliminate the error condition.

Error	Error Message and Description
201	FILE ACCESS ERROR ON INPUT FILE FILE: filename An error occurred while reading the specified file.
202	FILE ACCESS ERROR ON OUTPUT FILE FILE: filename An error occurred while writing the specified file.
203	NOT A BANKED 8051 OBJECT FILE The input file is not a banked object file.
204	INVALID INPUT FILE The input file has an invalid format.
205	CHECKSUM ERROR The input file has an invalid checksum. This error is usually caused by an error from BL51. Make sure that your program was linked successfully.
206	INTERNAL ERROR OC51 has detected an internal error. Contact technical support.
207	SCOPE LEVEL ERROR MODULE: modulename The symbolic information in the specified file contains errors. This error message is usually the result of an error at link time. Make sure that your program was linked successfully.
208	PATH OR FILE NOT FOUND FILE: filename The OC51 Banked Object File Converter cannot find the specified file. Make sure the file actually exists.

Intel HEX File Format

The Intel HEX file is an ASCII text file with lines of text that follow the Intel HEX file format. Each line in an Intel HEX file contains one HEX record. These records are made up of hexadecimal numbers that represent machine language code and/or constant data. Intel HEX files are often used to transfer the program and data that would be stored in a ROM or EPROM. Most EPROM programmers or emulators can use Intel HEX files.

Record Format

An Intel HEX file is composed of any number of HEX records. Each record is made up of five fields that are arranged in the following format:

```
:llaaaatt dd... cc
```

Each group of letters corresponds to a different field, and each letter represents a single hexadecimal digit. Each field is composed of at least two hexadecimal digits—which make up a byte—as described below:

:	is the colon that starts every Intel HEX record.	
11	is the record-length field that represents the number of data bytes (ad) in the record.	
aaaa	is the address field that represents the starting address for subsequent data in the record.	
tt	is the field that represents the HEX record type, which may be one of the following:	
	 data record end-of-file record extended 8086 segment address record. extended linear address record. 	
dd	is a data field that represents one byte of data. A record may have multiple data bytes. The number of data bytes in the record must match the number specified by the 11 field.	
cc	is the checksum field that represents the checksum of the	

record. The checksum is calculated by summing the values

of all hexadecimal digit pairs in the record modulo 256 and taking the two's complement.

Data Record

The Intel HEX file is made up of any number of data records that are terminated with a carriage return and a linefeed. Data records appear as follows:

:10246200464C5549442050524F46494C4500464C33

where:

is the number of data bytes in the record.

is the address where the data are to be located in memory.

is the record type 00 (a data record).

464C...464C is the data.

is the checksum of the record.

End-of-File (EOF) Record

An Intel HEX file must end with an end-of-file (EOF) record. This record must have the value 01 in the record type field. An EOF record always appears as follows:

:0000001FF

where:

is the number of data bytes in the record.

is the address where the data are to be located in memory. The

address in end-of-file records is meaningless and is ignored. An

address of 0000h is typical.

os is the record type 01 (an end-of-file record).

is the checksum of the record and is calculated as

01h + NOT(00h + 00h + 00h + 01h).

Extended 8086 Segment Record

The Intel HEX contains extended 8086 segment records when the H86 directive is used. This record is used to specify an address offset (in 8086 paragraph form) for the following data records. Extended 8086 segment records appear as follows:

:02000002F0000C

where:

is the number of data bytes in the record.

is always 0 in a extended 8086 segment record.

is the record type 02 (a extended 8086 segment record).

F000 is the offset in 8086 paragraph notation (= 0x0F0000).

oc is the checksum of the record.

Extended Linear Address Record

The Intel HEX contains extended linear address records when the H386 directive is used. This record is used to specify the two most significant bytes (bits 16 - 31) of the absolute address. This address offset is used for all following data records. Extended linear address records appear as follows:

:0200000400FFFB

where:

is the number of data bytes in the record.

is always 0 in a extended 8086 segment record.

is the record type 04 (a extended linear address record).

is the high word of the address offset (= 0xFF0000).

is the checksum of the record.

Example Intel HEX File

Following is an example of a complete Intel HEX file:

- :0200000400FFFB
- :03000000020003F8
- :10000300758107758920758DFDD28E759852C20052
- :0B00130090001E12003612002B80F53A
- :0D001E00544553542050524F4752414D005D
- :10002B00740D120047740A12004722200004E49357
- :0C003B008001E06006120047A380F02264
- :080047003099FDC299F59922E0
- :0000001FF

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Appendix A. Application Examples

This chapter illustrates project development for the x51 microcontroller family. The sample programs are found in the C:\KEIL\C51\EXAMPLES\ or C:\KEIL\C251\EXAMPLES\ folder. Each sample program is stored in a separate folder along with μ Vision2 project files that help you quickly build each sample program.

The following table lists the sample programs that are discussed in the following section and their folder names.

Example	Description		
ASM	This section describes a short x51 program, developed in assembler language. The program displays the text "PROGRAM TEST" on the serial interface.		
CSAMPLE	Simple addition and subtraction calculator that shows how to build a C application.		
BANK_EX1	C application for a classic 8051 device that shows code banking.		
BANK_EX2	C program for a classic 8051 device that stores constants in different code banks.		
BANK_EX3	Intel PL/M-51 application that uses code banking.		
Philips 80C51MX \ASM	Assembler example that demonstrates the new instructions of the Philips 80C51MX.		
Philips 80C51MX \Banking	C Compiler example that shows how to use the extended address space of the Philips $80\text{C}51\text{MX}.$		

The folder **EXAMPLES** contains several other example programs that are described in the $\mu Vision 2$ for the x51 Family User's Guide.

ASM – Assembler Example

This section shows you how to create a *x*51 program, developed in assembler language. The program outputs the text "PROGRAM TEST" on the serial interface. The program consists of three modules that can be translated using the various tool versions.





Using A51 and BL51

The following commands are required to translate and link the ASM example with the A51 macro assembler and the BL51 linker/locater. The output file can be converted into an Intel HEX file with the OH51 hex file converter.

```
A51 ASAMPLE1.A51 DEBUG XREF
A51 ASAMPLE2.A51 DEBUG XREF
A51 ASAMPLE3.A51 DEBUG XREF
```

The **XREF** control causes the A51 assembler to include in the listing (LST) files a cross reference report of the symbols used in the module. The **DEBUG** control includes complete symbol information in the object file.

After assembly, the files are linked by the **BL51** linker/locator with:

```
BL51 ASAMPLE1.OBJ, ASAMPLE2.OBJ, ASAMPLE3.OBJ PRECEDE (VAR1) IXREF
```

In the above linker command line, the **PRECEDE** control locates the VAR1 segment before other internal data memory segments. The **IXREF** control includes a cross reference report of all public and external symbols in the linker listing (M51) file. The linker creates an absolute object module that is stored in the file **ASAMPLE1**. This file can be used as input for debuggers or may be used to create an Intel HEX file using the **OH51** object hex converter with the following command:

OH51 ASAMPLE1

Using AX51 and LX51

The commands for translating the application with the AX51 macro assembler and the LX51 linker/locater are:

```
AX51 ASAMPLE1.A51 DEBUG XREF
AX51 ASAMPLE2.A51 DEBUG XREF
AX51 ASAMPLE3.A51 DEBUG XREF
```

After assembly, the files are linked by the LX51 linker/locator with:

```
LX51 ASAMPLE1.OBJ, ASAMPLE2.OBJ, ASAMPLE3.OBJ SEGMENTS (VAR1) IXREF
```

The **SEGMENTS** control replaces the PRECEDE control used in the **BL51** command line to locate the VAR1 segment before other internal data memory segments. The **IXREF** control includes a cross reference in the linker listing (MAP) file. The file **ASAMPLE1** is the absolute object module created by the linker. This file can be used as input for debuggers or may be converted into an Intel HEX file using **OHX51** with the following command:

OHX51 ASAMPLE1

Using A251 and L251

The Intel/Atmel WM 251 application is build with the following commands:

```
A251 ASAMPLE1.A51 DEBUG XREF
A251 ASAMPLE2.A51 DEBUG XREF
A251 ASAMPLE3.A51 DEBUG XREF
L251 ASAMPLE3.OBJ, ASAMPLE2.OBJ, ASAMPLE3.OBJ SEGMENTS (VAR1) IXREF
```

The **SEGMENTS** control locates the VAR1 segment before other internal data memory segments. The **IXREF** control includes a cross reference in the linker listing (MAP) file. The file **ASAMPLE1** is the absolute object module created by the linker. This file can be used as input for debuggers or may be converted into an Intel HEX file using **OH251** with the following command:

OH251 ASAMPLE1



CSAMPLE – C Compiler Example

This section describes shows a *x*51 program, developed with the C*x*51 compiler. This program demonstrates the concept of modular programming development and can be translated using the various tool versions.

The program calculates the sum of two input numbers and displays the result. Numbers are input with the **getchar** library function and results are output with the **printf** library function. The program consists of three source modules, which are translated using the following command lines.

Using C51 and BL51

The following commands are required to translate and link the C example with the C51 compiler and the BL51 linker/locater. The output file can be converted into an Intel HEX file with the OH51 hex file converter.

```
C51 CSAMPLE1.C DEBUG OBJECTEXTEND
C51 CSAMPLE2.C DEBUG OBJECTEXTEND
C51 CSAMPLE3.C DEBUG OBJECTEXTEND
```

The **DEBUG** and **OBJECTEXTEND** control directs the compiler to include complete symbol information in the object file.

After compilation, the files are linked using the **BL51** linker/locator:

```
BL51 CSAMPLE1.OBJ, CSAMPLE2.OBJ, CSAMPLE3.OBJ PRECEDE (?DT?CSAMPLE3) IXREF
```

In the above linker command line, the **PRECEDE** control locates the **?DT?CSAMPLE3** segment before other internal data memory segments. The **IXREF** control includes a cross reference report in the linker listing (M51) file. The linker creates an absolute object module that is stored in the file **CSAMPLE1**. This file can be used as input for debuggers or may be used to create an Intel HEX file using the **OH51** object hex converter with the following command:

OH51 CSAMPLE1

Using C51 and LX51

The commands for translating the application with the C51 compiler and the LX51 linker/locater are:

```
C51 CSAMPLE1.C DEBUG OMF251
C51 CSAMPLE2.C DEBUG OMF251
C51 CSAMPLE3.C DEBUG OMF251
```

The **DEBUG** control directs the compiler to include symbol information in the object file. The **OMF251** control generates extended object files that support the extensions of the **LX51** linker/locater. The files are linked with:

```
LX51 CSAMPLE1.OBJ, CSAMPLE2.OBJ, CSAMPLE3.OBJ SEGMENTS (?DT?CSAMPLE3) IXREF
```

The **SEGMENTS** control replaces the PRECEDE control used in the **BL51** command line to locate the **?DT?CSAMPLE3** segment before other internal data memory segments. The **IXREF** control includes a cross reference in the linker listing (MAP) file. The file **CSAMPLE1** is the absolute object module created by the linker. This file can be used as input for debuggers or may be converted into an Intel HEX file using **OHX51** with the following command:

OHX51 CSAMPLE1

Using C251 and L251

The Intel/Atmel WM 251 application is build with the following commands:

```
C251 CSAMPLE1.C DEBUG
C251 CSAMPLE2.C DEBUG
C251 CSAMPLE3.C DEBUG
```

After assembly, the files are linked by the **L251** linker/locator with:

```
L251 CSAMPLE1.OBJ, CSAMPLE2.OBJ, CSAMPLE3.OBJ SEGMENTS (?DT?CSAMPLE3) IXREF
```

The **SEGMENTS** control locates the **?DT?CSAMPLE3** segment before other internal data memory segments. The **IXREF** control includes a cross reference in the linker listing (MAP) file. The file **CSAMPLE1** is the absolute object module created by the linker. This file can be used as input for debuggers or may be converted into an Intel HEX file using **OH251** with the following command:

OH251 CSAMPLE1



BANK_EX1 - Code Banking with C51

The following C51 example shows how to compile and link a program using multiple code banks.

The program begins with function **main** in C_ROOT.C. The **main** function calls functions in other code banks. These functions, in turn, call functions in yet different code banks. The **printf** function outputs the number of the code bank in each function.

Using C51 and BL51

The program can be translated using the following commands:

```
C51 C_ROOT.C DEBUG OBJECTEXTEND
C51 C_BANKO.C DEBUG OBJECTEXTEND
C51 C_BANKI.C DEBUG OBJECTEXTEND
C51 C_BANK2.C DEBUG OBJECTEXTEND
```

C_ROOT.C contains the main function and is located in the common area. C_BANK0.C, C_BANK1.C, and C_BANK2.C contain the bank functions and are located in the bank area. The BL51 linker/locator is invoked as follows:

```
BL51 COMMON{C_ROOT.OBJ}, BANKO{C_BANKO.OBJ},
BANK1{C_BANK1.OBJ}, BANK2{C_BANK2.OBJ}
BANKAREA(8000H,0FFFFH)
```

The **BANKAREA** (8000H, 0FFFFH) control defines the address space 80000H to 0FFFFH as the area for code banks. The **COMMON** control places the **C_ROOT.OBJ** module in the common area. The **BANK0**, **BANK1**, and **BANK2** controls place modules in bank 0, 1, and 2 respectively.

The **BL51** linker/locator creates a listing file, C_ROOT.M51, which contains information about memory allocation and about the intra-bank jump table that is generated. BL51 also creates the output file C_ROOT that in banked object file format. You must use the OC51 banked object file converter to convert this file into standard object files:

```
OC51 C_ROOT
```

For this example program, the OC51 banked object file converter produces three standard object files from C ROOT. They are listed in the following table.

Filename	Contents
C_ROOT.B00	All information (including symbols) for code bank 0 and the common area.
C_ROOT.B01	Information for code bank 1 and the common area.
C_ROOT.B02	Information for code bank 2 and the common area.

You can create Intel HEX files for each of these object files by using the OH51 object to hex converter. The Intel HEX files you create with OH51 contain complete information for each code bank including the common area:

```
OH51 C_ROOT.B00 HEXFILE (C_ROOT.H00)
OH51 C_ROOT.B01 HEXFILE (C_ROOT.H01)
OH51 C_ROOT.B02 HEXFILE (C_ROOT.H02)
```





Using C51 and LX51

When you are using the extended LX51 linker/locater the program is generated as shown below:

```
C51 C_ROOT.C DEBUG OMF251
C51 C_BANKO.C DEBUG OMF251
C51 C_BANK1.C DEBUG OMF251
C51 C_BANK2.C DEBUG OMF251
LX51 COMMON{C_ROOT.OBJ}, BANKO{C_BANKO.OBJ},
BANK1{C_BANK1.OBJ}, BANK2{C_BANK2.OBJ}
BANKAREA(8000H,0FFFFH)
```

The **LX51** linker/locator creates a listing file, **C_ROOT.MAP**, which contains information about memory allocation and about the intra-bank jump table that is generated. The linker output file **C_ROOT** can be directly converted into an Intel HEX file with OHX51:

```
OHX51 C ROOT
```

BANK_EX2 – Banking with Constants

This example shows how to place constants in code banks. You can use this technique to place messages or large tables in code banks other than the one in which your program resides. This example uses three source files: C_PROG.C, C_MESSO.C, and C_MESS1.C.

You use the **LX51** linker/locator to locate constant segments in particular code banks. Segment names for constant data have the general format ?CO?*modulename* where *modulename* is the name of the source file the constant data is declared.

In your C51 programs, when you access constant data that is in a different segment, you must manually ensure that the proper code bank is used when accessing that constant data. You so this with the **switchbank** function. This function is defined in the **L51_BANK.A51** source module.

Using C51 and BL51

These source files are compiled and linked using the following commands.

```
C51 C_PROG.C DEBUG OBJECTEXTEND
C51 C_MESSO.C DEBUG OBJECTEXTEND
C51 C_MESS1.C DEBUG OBJECTEXTEND

BL51 C_PROG.OBJ, C_MESSO.OBJ, C_MESS1.OBJ
BANKAREA(8000H, OFFFFH) &
BANKO(?CO?C_MESSO (8000H)) BANK1(?CO?C_MESS1 (8000H))

OC51 C_PROG
OH51 C_PROG.B00 HEXFILE (C_PROG.H00)
OH51 C_PROG.B01 HEXFILE (C_PROG.H01)
```





Using C51 and LX51

When you are using the extended LX51 linker/locater the program is generated as shown below:

```
C51 C_PROG.C DEBUG OMF251
C51 C_MESSO.C DEBUG OMF251
C51 C_MESS1.C DEBUG OMF251
LX51 C_PROG.OBJ, C_MESSO.OBJ, C_MESS1.OBJ
BANKAREA(8000H,0FFFFH) &
SEGMENTS (?CO?C_MESSO (B0:8000H)) BANK1(?CO?C_MESS1 (B1:8000H))
OHX51 C_PROG
```

BANK_EX3 - Code Banking with PL/M-51

The following PL/M-51 example shows how to compile and link a PL/M-51 program using multiple code banks. The function of this example is similar to that shown in "BANK_EX1 – Code Banking with C51" on page 404.

The program begins with the procedure in P_ROOT.P51. This routine calls routines in other code banks, which, in turn, call routines in yet different code banks.

The PL/M-51 programs are compiled using the following commands.

```
PLM51 P_ROOT.P51 DEBUG
PLM51 P_BANKO.P51 DEBUG
PLM51 P_BANK1.P51 DEBUG
PLM51 P_BANK2.P51 DEBUG
```

In this example, P_ROOT.OBJ is located in the common area and P_BANK0.OBJ, P_BANK1.OBJ, and P_BANK2.OBJ are located in the bank area.

NOTE

The PL/M-51 runtime library, **PLM51.LIB**, must be included in the linkage. You must either specify a path to the directory in which this library is stored, or you must include it directly in the linker command line.

Using BL51

The **BL51** linker/locator is invoked as follows:

```
BL51 COMMON{P_ROOT.OBJ}, BANKO{P_BANKO.OBJ}, &
BANK1{P_BANK1.OBJ}, BANK2{P_BANK2.OBJ} &
BANKAREA(8000H,0FFFFH)
```

The **BANKAREA** (8000H, 0FFFFH) control defines the address space 8000H to 0FFFFH as the area for code banks. The **COMMON** control places the **P_ROOT.OBJ** module in the common area. The **BANK0**, **BANK1**, and **BANK2** controls place modules in bank 0, 1, and 2 respectively.

The **BL51** linker/locator creates a listing file, **P_ROOT.M51**, which contains information about memory allocation and about the intra-bank jump table that is generated. BL51 also creates the output module, **P_ROOT**, which is stored in banked OMF format. You must use the OC51 banked object file converter to convert the banked OMF file into standard OMF files. OMF files may be loaded with the μ Vision2 Debugger/Simulator or an in-circuit emulator. Invoke the OC51 banked object file converter as follows:

```
OC51 P_ROOT
```

For this example program, the OC51 banked object file converter produces three standard OMF-51 files from **P ROOT**. They are listed in the following table.

Filename	Contents
P_ROOT.B00	All information (including symbols) for code bank 0 and the common area.
P_ROOT.B01	Information for code bank 1 and the common area.
P_ROOT.B02	Information for code bank 2 and the common area.

You can create Intel HEX files for each of these OMF-51 files by using the OH51 object to hex converter. The Intel HEX files you create with OH51 contain complete information for each code bank including the common area. Intel HEX files can be generated using the following OH51 object to hex converter command line.

```
OH51 P_ROOT.B00 HEXFILE (P_ROOT.H00)
OH51 P_ROOT.B01 HEXFILE (P_ROOT.H01)
OH51 P_ROOT.B02 HEXFILE (P_ROOT.H02)
```

A



Using C51 and LX51

When you are using the extended LX51 linker/locater the program is generated as shown below:

```
LX51 COMMON {P_ROOT.OBJ}, BANKO {P_BANKO.OBJ},
BANK1 {P_BANK1.OBJ}, BANK2 {P_BANK2.OBJ}
BANKAREA (8000H, 0ffffh)

OHX51 P_ROOT
```

Philips 80C51MX – Assembler Example

The example **Philips 80C51MX\ASM** shows how to use the new instructions of the Philips 80C51MX architecture in assembly language. Segments with the memory class **ECODE** are used to show the ECALL and ERET instructions. Segments with **HCONST** and **HDATA** are used to show how to access memory in the 16MB address space of this architecture.

The example program is build with the **AX51** macro assembler and the **LX51** linker/locater as shown below:

```
AX51 MX_INST.A51 DEBUG MOD_MX51

LX51 MX_INST.OBJ

OHX51 MX_INST
```

Philips 80C51MX – C Compiler Example

The example **Philips 80C51MX\Banking** shows how to create large C programs for the Philips 80C51MX architecture. The program uses the code banking facilities of the **LX51** linker/locater to place program code into the code pages 0x80:0000 (bank 0) and 0x81:0000 (bank 1). The function of this example is similar to that shown in "BANK_EX1 – Code Banking with C51" on page 404. In addition some variables are declared with the *far* memory type to show the usage of the HCONST and HDATA memory class.

The example program is build with the **AX51** macro assembler and the **LX51** linker/locater as shown below:

```
CX51 C_ROOT.C DEBUG
CX51 C_BANKO.C DEBUG
CX51 C_BANKI.C DEBUG
CX51 C_BANKI.C DEBUG
AX51 START_MX.A51 MOD_MX51

LX51 COMMON {C_ROOT.OBJ, START_MX.OBJ},
BANKO {C_BANKO.OBJ}, BANKI {C_BANK1.OBJ}
CLASSES (HCONST (0x810000 - 0x81FFFF), HDATA (0x010000 - 0x01FFFF))

OHX51 C_ROOT
```





Appendix B. Reserved Symbols

The **Ax51** assembler uses predefined or reserved symbols that may not be redefined in your program. Reserved symbol names include instruction mnemonics, directives, operators, and register names. The following lists the reserved symbol names that are found in all **Ax51** variants:

A	DA	INPAGE	MOD	REPT
AB	DATA	INSEG	MOV	RET
ACALL	DB	IRP	MOVC	RETI
ADD	DBIT	IRPC	MOVX	RL
ADDC	DEC	ISEG	MUL	RLC
AJMP				
AND	DIV	JB JBC	NAME NE	RR RRC
	DJNZ			
ANL	DPTR	JC	NOP	RSEG
AR0	DS	JE	NOT	SEG
AR1	DSEG	JG	NUL	SEGMENT
AR2	DW	JLE	NUMBER	SET
AR3	ELSE	JMP	OR	SETB
AR4	ELSEIF	JNB	ORG	SHL
AR5	END	JNC	ORL	SHR
AR6	ENDIF	JNE	OVERLAYABLE	SJMP
AR7	ENDM	JNZ	PAGE	SUB
BIT	ENDP	JSG	PC	SUBB
BITADDRESSABLE	EQ	JSGE	POP	SWAP
BLOCK	EQU	JSL	PUBLIC	UNIT
BSEG	EXITM	JSLE	PUSH	USING
c	EXTRN	JZ	R0	XCH
CALL	GE	LCALL	R1	XCHD
CJNE	GT	LE	R2	XDATA
CLR	HIGH	LJMP	R3	XOR
CMP	IDATA	LOCAL	R4	XRL
CODE	IF	LOW	R5	XSEG
CPL	INBLOCK	LT	R6	
CSEG	INC	MACRO	R7	

The **A51** assembler defines the following additional reserved symbols which are special function registers (SFR) of the classic 8051 CPU. These SFR definitions may be disabled using the **NOMOD51** control. The predefined SFR symbols are reserved symbols and may not be redefined in your program.

AC	IE	PS	SBUF	TI
ACC	IEO	PSW	SCON	TL0
В	IE1	PT0	SM0	TL1
CY	INTO	PT1	SM1	TMOD
DPH	INT1	PX0	SM2	TO
DPL	IT0	PX1	SP	TR0
EA	IT1	RB8	T1	TR1
ES	ov	RD	TB8	TXD
ET0	P	REN	TCON	WR
ET1}	P0	RI	TF0	
EX0	P1	RS0	TF1	
EX1	P2	RS1	TH0	
F0	P3	RXD	TH1	

B

The **AX51** assembler defines the following additional reserved symbols which comprise the additional instructions and registers of the Philips 80C51MX architecture.

AT	DD	EDATA	FAR	OFFS
BYTE	DSB	EJMP	HCONST	PR0
BYTE0	DSD	EMOV	HDATA	PR1
BYTE1	DSW	EPTR	LABEL	PROC
BYTE2	DWORD	ERET	LIT	WORD
BYTE3	ECALL	EVEN	MBYTE	WORD0
CONST	ECODE	EXTERN	NEAR	WORD2

The **A251** assembler defines the following additional reserved symbols which comprise the additional instructions and registers of the Intel and Atmel WM 251 architectures.

AT	DR56	EXTERN	R12	WR12
BYTE	DR60	FAR	R13	WR14
BYTE0	DR8	HCONST	R14	WR16
BYTE1	DSB	HDATA	R15	WR18
BYTE2	DSD	LABEL	R8	WR2
BYTE3	DSW	LIT	R9	WR20
CONST	DWORD	MOVH	SLL	WR22
DD	EBIT	MOVS	SRA	WR24
DR0	EBITADDRESSABLE	MOVZ	SRL	WR26
DR12	ECALL	NCONST	TRAP	WR28
DR16	ECODE	NEAR	WORD	WR30
DR20	EDATA	OFFS	WORD0	WR4
DR24	EJMP	PROC	WORD2	WR6
DR28	ERET	R10	WR0	WR8
DR4	EVEN	R11	WR10	

Appendix C. Listing File Format

This appendix describes the format of the listing file generated by the assembler.

Assembler Listing File Format

The **Ax51** assembler, unless overridden by controls, outputs two files: an object file and a listing file. The object file contains the machine code. The listing file contains a formatted copy of your source code with page headers and, if requested through controls (**SYMBOL** or **XREF**), a symbol table.

Sample Ax51 Listing

```
A251 MACRO ASSEMBLER ASAMPLE1
                                                         25/01/2001 15:02:23 PAGE
A251 MACRO ASSEMBLER V1.40
OBJECT MODULE PLACED IN ASAMPLE1.OBJ
ASSEMBLER INVOKED BY: F:\RK\ZX\ASM\A251.EXE ASAMPLE1.A51 XREF
                        T.TNE
                                  SOURCE
                           1
                                  $NOMOD51
                           2
                                 $INCLUDE (REG52.INC)
                         3 +1 $SAVE
                     +1 106 +1 $RESTORE
                         107
                         108
                                NAME SAMPLE
                          109
                                EXTRN CODE
                                                   (PUT CRLF, PUTSTRING)
                         110
                         111
                                PUBLIC TXTBIT
                         112
                         113
                                 PROG SEGMENT CODE
                        113 PROG SEGMENT CODE
114 PCONST SEGMENT CODE
                        115 VAR1 SEGMENT DATA
                        116 BITVAR SEGMENT BIT
                         117
                                STACK SEGMENT IDATA
                         118
                         119
                                         RSEG STACK
DS 10H ; 16 Bytes Stack

121

000000 122 CSEG AT 0

123 USING 0 ; Register-Bank 0

124 ; Execution starts at address 0 on power-up.

125 JMP START
                                  DS 10H ; 16 Bytes Stack
                         127
                                         RSEG PROG
                         128
                                 ; first set Stack Pointer
000000 758100 F 129
                               START: MOV SP, #STACK-1
                         130
                         131
                                 ; Initialize serial interface
                        131 ; Initialize serial interface
132 ; Using TIMER 1 to Generate Baud Rates
133 ; Oscillator frequency = 11.059 MHz
                       134
000003 758920
                                        MOV TMOD, \#00100000B; C/T = 0, Mode = 2
                                         MOV
000006 758DFD
                         135
                                                TH1.#0FDH
000009 D28E
                         136
                                          SETB TR1
                                         MOV SCON, #01010010B
00000B 759852
                         137
```

```
139
                                     ; clear TXTBIT to read form CODE-Memory
00000E C200
                            140
                                              CLR
                                                    TXTBIT
                            141
                            142
                                    ; This is the main program. It is a loop,
                            143
                                    ; which displays the a text on the console.
000010
                            144
                                     REPEAT:
                            145
                                     ; type message
000010 900000
                   F
                           146
                                             MOV DPTR, #TXT
000013 120000
                 E
                           147
                                              CALL PUTSTRING
000016 120000
                   E
                            148
                                              CALL PUT_CRLF
                            149
                                     ; repeat
000019 8000
                           150
                                             SJMP REPEAT
                    F
                            151
                            152
                                             RSEG PCONST
000000 54455354
                                     TXT:
                                            DB 'TEST PROGRAM',00H
                            153
000004 2050524F
000008 4752414D
00000C 00
                            154
                            155
                                     ; only for demonstration
                                       RSEG VAR1
                            156
                                     DUMMY: DS
000000
                            157
                            158
                            159
                                     ; TXTBIT = 0 read text from CODE Memory
                            160
                                     ; TXTBIT = 1 read text from XDATA Memory
                            161
                                            RSEG BITVAR
0000.0
                            162
                                     TXTBIT: DBIT 1
                            163
                            164
                                              END
XREF SYMBOL TABLE LISTING
                     TYPE VALUE ATTRIBUTES / REFERENCES
NAME
BITVAR . . . . . B SEG 000001H
                                               REL=UNIT, ALN=BIT 116# 161
DUMMY. . . . . . D ADDR 000000H R
                                               SEG=VAR1 157#
PCONST . . . . . C SEG 00000DH
                                               REL=UNIT, ALN=BYTE 114# 152
                                               REL=UNIT, ALN=BYTE 113# 127
PROG . . . . . . C SEG 00001BH

      PUTSTRING.
      C ADDR
      ------
      EXT 110# 147

      PUT_CRLF.
      C ADDR
      ------
      EXT 110# 148

      REPEAT.
      C ADDR
      000010H R
      SEG=PROG
      144# 150

REL=UNIT, ALN=BYTE 117# 119 129

      START.
      .
      .
      .
      C
      ADDR
      000000H R

      TXT.
      .
      .
      .
      C
      ADDR
      000000H R

      TXTBIT
      .
      .
      .
      B
      ADDR
      00000H.0 R

                                               SEG=PROG 125 129#
                                               SEG=PCONST 146 153#
                                               SEG=BITVAR 111 140 162#
VAR1 . . . . . . D SEG 000021H
                                               REL=UNIT, ALN=BYTE 115# 156
REGISTER BANK(S) USED: 0
ASSEMBLY COMPLETE. 0 WARNING(S), 0 ERROR(S)
```

Listing File Heading

Every page has a header on the first line. It contains the words Ax51 MACRO ASSEMBLER followed by the title, if specified. If the title is not specified, then the module name is used. It is derived from the **NAME** directive (if specified), or from the root of the source filename. On the extreme right side of the header, the date (if specified) and the page number are printed.

In addition to the normal header, the first page of listing includes the **Ax51** listing file header. This header shows the assembler version number, the file name of the object file, if any, and the entire invocation line.

C

Source Listing

The main body of the listing file is the formatted source listing. A section of formatted source is shown in the following.

Sample Source Listing

FOC	OBJ	LINE	SOURCE		
000006	758DFD	135	MOV	TH1,#0FDH	

The format for each line in the listing file depends on the source line that appears on it. Instruction lines contain 4 fields. The name of each field and its meanings is shown in the list below:

- Loc shows the location relative or absolute (code address) of the first byte of the instruction. The value is displayed in hexadecimal.
- OBJ shows the actual machine code produced by the instruction, displayed in hexadecimal. If the object that corresponds to the printed line is to be fixed up (it contains external references or is relocatable), an F or E is printed after the OBJ field. The object fields to be fixed up contain zeros.
- LINE shows the **INCLUDE** nesting level, if any, the number of source lines from the top of the program, and the macro nesting level, if any. All values in this field are displayed in decimal numbers.
- source shows the source line as it appears in the file. This line may be extended onto the subsequent lines in the listing file.

DB, **DW**, and **DD** directives are formatted similarly to instruction lines, except the OBJ field shows the data values placed in memory. All data values are shown. If the expression list is long, then it may take several lines in the listing file to display all of the values placed in memory. The extra lines will only contain the LOC and OBJ fields.

The directives that affect the location counter without initializing memory (e.g. **ORG**, **DBIT**, or **DS**) do not use the OBJ field, but the new value of the location counter is shown in the LOC field.

The **SET** and **EQU** directives do not have a LOC or OBJ field. In their place the assembler lists the value that the symbol is set to. If the symbol is defined to equal one of the registers, then REG is placed in this field. The remainder of the directive line is formatted in the same way as the other directives.

C

Macro / Include File / Save Stack Format

In the listing file, the assembler displays the macro nesting level, the include file level, and the level of the **SAVE/RESTORE** stack. These nesting levels are shown before and after the LINE number as shown in the following listing.

```
LOC
       OBJ
                    LINE
                             SOURCE
                             $GEN ; Enable Macro Listing
                                                    ; A sample macro
                             MYMACRO MACRO
                                     INC A
                       4
                                                    ; Macro Level 1
                                      MACRO2
                                      ENDM
                                                    ; Macro 2
                       8
                           MACRO2 MACRO
                       9
                                     NOP
                                                    ; Macro Level 2
                      10
                      11
                      12
                      13 MYPROG SEGMENT CODE
                     14
                                     RSEG MYPROG
                     15
000000 7400
                     16
                                     MOV
                                             A,#0
                      17
                                     MYMACRO
000002 04
                     18+1
                                     INC A
                                                    ; Macro Level 1
                                    MACRO2
                     19+1
000003 00
               20+2 NOP ; Macro Level 2
21 $INCLUDE (MYFILE.INC) ; A include file
+1 22 ; This is a comment ; Include Level 1
                +1 23
                                    MACRO2
              +1 24+1 NOP ; Macro Level 1
25 MOV A,#1
26 +1 $SAVE ; Save Control
000004 00
000005 7401
                                                  ; Save Control
                     26 +1 $SAVE ; Save Control
27 +1 MYMACRO ; SAVE Level 1
28+1+1 INC A ; Macro Level 1
000007 04
                     28+1+1
                     29+1+1 MACRO2
30+2+1 NOP
000008 00
                      30+2+1
                                                    ; Macro Level 2
                      31 +1 $RESTORE
000009 00
                                     NOP
                     32
```

Symbol Table

The symbol table is a list of all symbols defined in the program along with the status information about the symbol. Any predefined symbols used will also be listed in the symbol table. If the XREF control is used, the symbol table will contain information about where the symbol was used in the program.

The status information includes a NAME field, a TYPE field, a VALUE field, and an ATTRIBUTES field.

The TYPE field specifies the type of the symbol: ADDR if it is a memory address, NUMB if it is a pure number (e.g., as defined by EQU), SEG if it is a relocatable segment, and REG if a register. For ADDR and SEG symbols, the segment type is added.

The VALUE field shows the value of the symbol when the assembly was completed. For REG symbols, the name of the register is given. For NUMB and ADDR symbols, their absolute value (or if relocatable, their offset) is given, followed by A (absolute) or R (relocatable). For SEG symbols, the segment size is given here. Bit address and size are given by the byte part, a period (.), followed by the bit part. The scope attribute, if any, is PUB (public) or EXT (external). These are given after the VALUE field.

The **ATTRIBUTES** field contains an additional piece of information for some symbols: relocation type for segments, segment name for relocatable symbols.

Example Symbol Table Listing

SYMBOL TABLE	T.TSTTN	G		
		_		
		_		
NAME		TYPE	VALUE	ATTRIBUTES
BITVAR		B SEG	000001H	REL=UNIT, ALN=BIT
DUMMY		D ADDR	000000H R	SEG=VAR1
PCONST		C SEG	00000DH	REL=UNIT, ALN=BYTE
PROG		C SEG	00001BH	REL=UNIT, ALN=BYTE
PUTSTRING		C ADDR		EXT
PUT_CRLF		C ADDR		EXT
REPEAT		C ADDR	000010H R	SEG=PROG
SAMPLE				
STACK		I SEG	000010H	REL=UNIT, ALN=BYTE
START		C ADDR	000000H R	SEG=PROG
TXT		C ADDR	000000H R	SEG=PCONST
TXTBIT		B ADDR	0000H.0 R	SEG=BITVAR
VAR1		D SEG	000021H	REL=UNIT, ALN=BYTE

If the **XREF** control is used, then the symbol table listing will also contain all of the line numbers of each line of code that the symbol was used. If the value of the symbol was changed or defined on a line, then that line will have a hash mark (#) following it. The line numbers are displayed in decimal.

Listing File Trailer

At the end of the listing, the assembler prints a message in the following format:

```
REGISTER BANK(S) USED: [r r r]

ASSEMBLY COMPLETE. (n) WARNING(S), (m) ERROR(S)

where

r are the numbers of the register banks used.

n is the number of warnings found in the program.

m is the number of errors found in the program.
```

C

C

Appendix D. Assembler Differences

This appendix lists the differences between the Intel ASM-51 assembler, the Keil A51 assembler, and the Keil A251/AX51 assembler.

Differences Between A51 and A251/AX51

Assembly modules written for the A51 assembler may be assembled using the A251/AX51 macro assembler. However, since the A251 macro assembler supports the Intel/Atmel WM 251 architecture and the AX51 macro assembler supports extended 8051 variants like the Philips 80C51MX, the following incompatibilities may arise when A51 assembly modules are assembled with the A251/AX51 assembler.

32-Bit Values in Numeric Evaluations

The A51 assembler uses 16-bit values for all numerical expressions. The A251/AX51 macro assembler uses 32-bit values. This may cause problems when overflows occur in numerical expressions. For example:

Value EQU (8000H + 9000H) / 2

generates the result 800h in A51 since the result of the addition is only a 16-bit value (1000h). However, the A251/AX51 assembler calculates a value of 8800h.

8051 Pre-defined Special Function Register Symbol Set

The default setting of the A51 assembler pre-defines the Special Function Register (SFR) set of 8051 CPU. This default SFR set can be disabled with the A51 control **NOMOD51**. Both A251 and AX51 do not pre-define the 8051 SFR set. The control **NOMOD51** is accepted by A251/AX51 but does not influence any SFR definitions.

More Reserved Symbols

The A251/AX51 macro assembler has more reserved symbols as A51. Therefore it might be necessary to change user-defined symbol names. For example the symbol ECALL cannot be used as label name in A251/AX51, since it is a mnemonic for a new instruction.

D

Object File Differences

Ax51 uses the OMF-251/51MX file format for object files. A51 uses an extended version of the Intel OMF-51 file format. The OMF-51 file format limits the numbers of external symbols and segments to 256 per module. The OMF-251 file format does not have such a limit on the segment and external declarations.

Differences between A51 and ASM51

Assembly modules written for the Intel ASM51 macro assembler can be re-translated with the A51 macro assembler. However you have to take care about the following differences:

Enable the MPL Macro Language

If your assembly module contains Intel ASM51 macros, the A51 MPL macros need to be enable with the **MPL** control.

8051 Pre-defined Interrupt Vectors

The Intel ASM51 pre-defines the following symbol names if **MOD51** is active: RESET, EXTI0, EXTI1, SINT, TIMER0, TIMER1. A51 does not pre-define this symbol names.

More Reserved Symbols

Since the A51 macro assembler supports also conditional assembly and standard macros, A51 has more reserved symbols then Intel ASM51. Therefore it might be necessary to change user-defined symbol names. For example the symbol IF cannot be used as label name in A51, since it is a control for conditional assembly.

Object File Differences

The A51 assembler generates line number information for source level debugging and file dependencies. For compatibility to previous A51 versions and to ASM51, the line number information can be disabled with the A51 control **NOLINES**.

■ C Preprocessor Side Effects

The integrated C preprocessor in **Ax51** has two side effects that are incompatible to the Intel ASM51 macro assembler. If you are using the backslash character at the end of a comment line, the next line will be comment out too. If you are using \$INCLUDE in conditional assembly blocks, the file must exist even when the block will not be assembled.

Differences between A251/AX51 & ASM51

Assembly modules written for Intel ASM51 can be re-translated with the A251 macro assembler. However, since the A251 macro assembler supports additional 251 features, the following incompatibilities can arise when ASM51 modules are re-translated with A251.

32-Bit Values in Numeric Evaluations

The ASM51 assembler uses 16-bit numbers for all numerical expressions. The A251 macro assembler uses 32-bit values. This can cause problems when overflows occur in numerical expressions. For example:

Value EQU (8000H + 9000H) / 2

has the result 800H in ASM51 since the result of the addition is only a 16-bit value (1000H), whereas the A251 calculates Value as 8800H.

8051 Pre-defined Symbols

The default setting of Intel ASM51 pre-defines the Special Function Register (SFR) set and symbol names for reset and interrupt vectors of 8051 CPU. This default symbol set can be disabled with the ASM51 control **NOMOD51**. A251 does not pre-define any of the 8051 SFR or interrupt vector symbols. The control **NOMOD51** is accepted by A251 but does not influence any symbol definitions.

More Reserved Symbols

The A251 macro assembler has more reserved symbols as ASM51. Therefore it might be necessary to change user-defined symbol names. For example the symbol ECALL cannot be used as label name in A251, since the Intel/Atmel WM 251 has a new instruction with that mnemonic.

Enable the MPL Macro Language

If your assembly module contains Intel ASM51 macros, the A251 MPL macros need to be enabled with the **MPL** control

Object File Differences

The A251 assembler uses the Intel OMF-251 file format for object files. The ASM51 assembler uses the Intel OMF-51 file format. The OMF-51 file format limits the numbers of external symbols and segments to 256 per module. The OMF-251 file format does not have such a limit on the segment and external declarations. The ASM51 assembler generates line number information for source level debugging. For compatibility with ASM51, line number information can be disabled with the A251 control **NOLINES**.

D

■ C Preprocessor Side Effects

The integrated C preprocessor in **Ax51** has two side effects that are incompatible to the Intel ASM51 macro assembler. If you are using the backslash character at the end of a comment line, the next line will be comment out too. If you are using \$INCLUDE in conditional assembly blocks, the file must exist even when the block will not be assembled.

D

Glossary

A51

The standard 8051 Macro Assembler.

AX51

The extended 8051 Macro Assembler.

A251

The 251 Macro Assembler.

ANSI

American National Standards Institute. The organization responsible for defining the C language standard.

argument

The value that is passed to a macro or function.

arithmetic types

Data types that are integral, floating-point, or enumerations.

array

A set of elements, all of the same data type.

ASCII

American Standard Code for Information Interchange. This is a set of 256 codes used by computers to represent digits, characters, punctuation, and other special symbols. The first 128 characters are standardized. The remaining 128 are defined by the implementation.

batch file

An ASCII text file containing commands and programs that can be invoked from the command line.

Binary-Coded Decimal (BCD)

A BCD (Binary-Coded Decimal) is a system used to encode decimal numbers in binary form. Each decimal digit of a number is encoded as a binary value 4 bits long. A byte can hold 2 BCD digits – one in the upper 4 bits (or nibble) and one in the lower 4 bits (or nibble).

BL51

The standard 8051 linker/locator

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block

A sequence of C statements, including definitions and declarations, enclosed within braces ({ }).

C51

The Optimizing C Compiler for classic 8051 and extended 8051 devices.

CX51

The Optimizing C Compiler for Philips 80C51MX architecture.

C251

The Optimizing C Compiler for Intel/Atmel WM 251.

constant expression

Any expression that evaluates to a constant non-variable value. Constants may include character and integer constant values.

control

Command line control switch to the compiler, assembler or linker.

declaration

A C construct that associates the attributes of a variable, type, or function with a name.

definition

A C construct that specifies the name, formal parameters, body, and return type of a function or that initializes and allocates storage for a variable.

directive

Instruction or control switch to the compiler, assembler or linker.

escape sequence

A backslash ('\') character followed by a single letter or a combination of digits that specifies a particular character value in strings and character constants.

expression

A combination of any number of operators and operands that produces a constant value.

formal parameters

The variables that receive the value of arguments passed to a function.

function

A combination of declarations and statements that can be called by name to perform an operation and/or return a value.

function body

A block containing the declarations and statements that make up a function.

function call

An expression that invokes and possibly passes arguments to a function.

function declaration

A declaration providing the name and return type of a function that is explicitly defined elsewhere in the program.

function definition

A definition providing the name, formal parameters, return type, declarations, and statements describing what a function does.

function prototype

A function declaration that includes a list of formal parameters in parentheses following the function name.

in-circuit emulator (ICE)

A hardware device that aids in debugging embedded software by providing hardware-level single-stepping, tracing, and break-pointing. Some ICEs provide a trace buffer that stores the most recent CPU events.

include file

A text file that is incorporated into a source file.

keyword

A reserved word with a predefined meaning for the compiler or assembler.

L51

The **old** version of the 8051 linker/locator. L51 is replaced with the **BL51** linker/locater.

LX51

The extended 8051 linker/locator.

L251

The 251 linker/locator.

LIB51, LIBX51, LIB251

The commands to manipulate library files using the Library Manager.

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library

A file that stores a number of possibly related object modules. The linker can extract modules from the library to use in building a target object file.

LSB

Least significant bit or byte.

macro

An identifier that represents a series of keystrokes.

manifest constant

A macro that is defined to have a constant value.

MCS[®] 51

The general name applied to the Intel family of 8051 compatible microprocessors.

MCS[®] 251

The general name applied to the Intel family of 251 compatible microprocessors.

memory model

Any of the models that specifies which memory areas are used for function arguments and local variables.

mnemonic

An ASCII string that represents a machine language opcode in an assembly language instruction.

MON51

An 8051 program that can be loaded into your target CPU to aid in debugging and rapid product development through rapid software downloading.

MON251

A 251 program that can be loaded into your target CPU to aid in debugging and rapid product development through rapid software downloading.

MSB

Most significant bit or byte.

newline character

A character used to mark the end of a line in a text file or the escape sequence ('\n') to represent the newline character.

null character

ASCII character with the value 0 represented as the escape sequence ('\0').

null pointer

A pointer that references nothing. A null pointer has the integer value 0.

object

An area of memory that can be examined. Usually used when referring to the memory area associated with a variable or function.

object file

A file, created by the compiler, that contains the program segment information and relocatable machine code.

OH51, OHX51, OH251

The commands to convert absolute object files into Intel HEX file format.

opcode

Also referred to as operation code. An opcode is the first byte of a machine code instruction and is usually represented as a 2–digit hexadecimal number. The opcode indicates the type of machine language instruction and the type of operation to perform.

operand

A variable or constant that is used in an expression.

operator

A symbol (e.g., +, -, *, /) that specifies how to manipulate the operands of an expression.

parameter

The value that is passed to a macro or function.

PL/M-51

A high-level programming language introduced by Intel at the beginning of the 1980's.

pointer

A variable containing the address of another variable, function, or memory area.

pragma

A statement that passes an instruction to the compiler at compile time.

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preprocessor

The compiler's first pass text processor that manipulates the contents of a C file. The preprocessor defines and expands macros, reads include files, and passes control directives to the compiler.

relocatable

Object code that can be relocated and is not at a fixed address.

RTX51 Full

An 8051 Real-time Executive that provides a multitasking operating system kernel and library of routines for its use.

RTX51 Tiny

A limited version of RTX51.

RTX251 Full

An 251 Real-Time Executive that provides a multitasking operating system kernel and library of routines for its use.

scalar types

In C, integer, enumerated, floating-point, and pointer types.

scope

Sections of a program where an item (function or variable) can be referenced by name. The scope of an item may be limited to file, function, or block.

Special Function Register (SFR)

An SFR or Special Function Register is a register in the 8051 internal data memory space that is used to read an write to the hardware components of the 8051. This includes the serial port, timers, counters, I/O ports, and other hardware control registers.

source file

A text file containing C program or assembly program code.

stack

An area of memory, indirectly accessed by a stack pointer, that shrinks and expands dynamically as items are pushed onto and popped off of the stack. Items in the stack are removed on a LIFO (last-in first-out) basis.

static

A storage class that, when used with a variable declaration in a function, causes variables to retain their value after exiting the block or function in which they are declared.

stream functions

Routines in the library that read and write characters using the input and output streams.

string

An array of characters that is terminated with a null character ($\0^{\circ}$).

string literal

A string of characters enclosed within double quotes ("").

structure

A set of elements of possibly different types grouped together under one name.

structure member

One element of a structure.

token

A fundamental symbol that represents a name or entity in a programming language.

two's complement

A binary notation that is used to represent both positive and negative numbers. Negative values are created by complementing all bits of a positive value and adding 1.

type

A description of the range of values associated with a variable. For example, an **int** type can have any value within its specified range (-32768 to 32767).

type cast

An operation in which an operand of one type is converted to another type by specifying the desired type, enclosed within parentheses, immediately preceding the operand.

μVision2

An integrated software development platform that supports the Keil Software development tools. µVision2 combines Project Management, Source Code Editing, and Program Debugging in one environment.

whitespace character

Characters used as delimiters in C programs such as space, tab, and newline.

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wild card

One of the characters (? or *) that can be used in place of characters in a filename.

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