

# **10.007: Modelling the Systems World**

**SUTD 2020 1D Project Class 19F07 Group 4**

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## PART A

**A0)** Our group members list is as follows:

Member Name	Student ID
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Thus, our average student ID value is 1004382.333 and our value of  $X \approx 219.1$ .

**A1)** Taking the natural logarithm of both sides of the equation  $F = \alpha e^{\beta C}$ , we would obtain:

$$\ln(F) = \beta C + \ln(\alpha)$$

To obtain the values of  $\alpha$  and  $\beta$ , we would use regression. By forming the matrices:

$$A = \begin{bmatrix} 1 & 8 \\ 1 & 10 \\ 1 & 12 \\ 1 & 14 \\ 1 & 16 \\ 1 & 18 \end{bmatrix}, \text{ and } \vec{b} = \begin{bmatrix} \ln(8.2) \\ \ln(19.1) \\ \ln(37.6) \\ \ln(70.4) \\ \ln(145.9) \\ \ln(219.1) \end{bmatrix},$$

$$\text{then } (A^T A)^{-1} \approx \begin{bmatrix} 2.5810 & -0.1857 \\ -0.1857 & 0.0143 \end{bmatrix},$$

$$\text{and } \begin{bmatrix} \ln(\alpha) \\ \beta \end{bmatrix} = \begin{bmatrix} c_0 \\ c_1 \end{bmatrix} = \vec{c} = (A^T A)^{-1} A^T \vec{b} \approx \begin{bmatrix} -0.416 \\ 0.331 \end{bmatrix}.$$

Therefore,  $\ln(\alpha) \approx -0.416$  and  $\beta \approx 0.331$ . Hence,  $\alpha \approx 0.660$  and  $\beta \approx 0.331$ .

**A2)** By substituting the values of  $\alpha$  and  $\beta$  into the equation, we would get the model as:

$$\ln(F) \approx 0.331C - 0.416$$

Substituting in  $C = 15$ , we would obtain  $\ln(F) \approx 4.55$  and thus,  $F \approx 94.3$ .

**A3)** From the graph,  $M$  is the 'plateau position' where  $F_i$  is levelling off, which is 450.0. By substituting  $C_i = d$ , we would obtain  $F_i = \frac{M}{2}$ . By using the graph to find the value of  $C_i$  where  $F_i = \frac{450.0}{2} = 225.0$ , we would obtain  $d = C_i \approx 18.3$ . Hence,  $M = 450.0$  and  $d \approx 18.3$ .

**A4)** We have the original equation as

$$F_i \approx \frac{M}{1 + e^{-k(C_i - d)}}, \text{ for } i \in \{1, 2, 3, \dots, 15\}.$$

We bring over the denominator from the right-hand side to the left-hand side, obtaining

$$F_i \left( 1 + e^{-k(C_i - d)} \right) \approx M$$

We would then rearrange the equation to obtain

$$e^{-k(C_i-d)} \approx \frac{M}{F_i} - 1, \text{ where (expression 1) } = \frac{M}{F_i} - 1 \text{ and (expression 2) } = -k(C_i - d)$$

Taking the natural logarithm of both sides, we would obtain

$$-k(C_i - d) \approx \ln\left(\frac{M}{F_i} - 1\right)$$

Hence, we would get  $\epsilon_i(k) = (c_i - d)k + \ln\left(\frac{M}{F_i} - 1\right)$ , where  $a_i = c_i - d$  and  $b_i = \ln\left(\frac{M}{F_i} - 1\right)$ .

**A5)** We could reformulate the non-linear constraint  $z \geq |f(k)|$  as two linear constraints,  $z \geq f(k)$  and  $z \geq -f(k)$ . Hence, expression 1 is  $f(k)$  and expression 2 is  $-f(k)$ .

**A6)** A general non-linear objective would be

$$\min \sum_{i=1}^n |f_i(k)|,$$

which could be linearized as

$$\left. \begin{array}{ll} \min & z = \sum_{i=1}^n f_i(k) \\ \text{subject to:} & \\ \forall j \in \{1, 2, 3, \dots, n\} & z \geq f_j(k) \\ \forall j \in \{1, 2, 3, \dots, n\} & z \geq -f_j(k) \end{array} \right\}$$

Hence, we could rewrite the optimization problem  $\mathcal{P}$  as the linear program:

$$\left. \begin{array}{ll} \min & z = \sum_{i=1}^{15} \epsilon_i(k) \\ \text{subject to:} & \\ \forall j \in \{1, 2, 3, \dots, 15\} & z \geq \epsilon_j(k) \\ \forall j \in \{1, 2, 3, \dots, 15\} & z \geq -\epsilon_j(k) \\ & k \geq 0 \end{array} \right\}$$

By using Excel with  $M = 450.0$  and  $d \approx 18.87$ , we would obtain  $k \approx 0.333$ , with a minimum total absolute error value of  $z \approx 0.457$ .

**A7)** We have the optimization problem

$$\left. \begin{array}{ll} \min & (\max\{|\epsilon_1(k)|, |\epsilon_2(k)|, \dots, |\epsilon_n(k)|\}) \\ \text{subject to:} & k \geq 0 \end{array} \right\},$$

which could be converted to the linear program

$$\left. \begin{array}{ll} \min & \epsilon \\ \text{subject to:} & \\ \forall i \in \{1, 2, 3, \dots, n\} & \epsilon \geq \epsilon_i(k) \\ \forall i \in \{1, 2, 3, \dots, n\} & \epsilon \geq -\epsilon_i(k) \\ & k \geq 0 \end{array} \right\},$$

where  $\epsilon$  is just a random variable that we have defined, instead of it being the sum of all  $\epsilon_i(k) \forall i \in \{1, 2, 3, \dots, n\}$ , which is subtly different from the value of  $z$  in the case of question **A6**.

**A8)** From the graph, we could estimate the value of  $\gamma$  to be the value of  $F_i$  where the value of  $\frac{dF_i}{dC_i}$  is at its maximum, which happens to be our value of  $\textcolor{red}{X}$  in this case. Hence,  $\gamma \approx 219.1$ .

We could rearrange the equation

$$F = (\alpha C + \beta)^{1/3} + \gamma$$

to become

$$F - \gamma = (\alpha C + \beta)^{1/3}$$

Raising both sides of the equation to the third power will yield us

$$(F - \gamma)^3 = \alpha C + \beta$$

To obtain the values of  $\alpha$  and  $\beta$ , we would use regression. By using our estimate of the value of  $\gamma$  and by forming the matrices:

$$X = \begin{bmatrix} 1 & 8 \\ 1 & 10 \\ 1 & 12 \\ 1 & 14 \\ 1 & 16 \\ 1 & 18 \\ 1 & 20 \\ 1 & 22 \\ 1 & 24 \\ 1 & 26 \\ 1 & 28 \\ 1 & 30 \\ 1 & 32 \\ 1 & 34 \\ 1 & 36 \end{bmatrix}, \text{ and } \vec{y} = \begin{bmatrix} (8.2 - \gamma)^3 \\ (19.1 - \gamma)^3 \\ (37.6 - \gamma)^3 \\ (70.4 - \gamma)^3 \\ (145.9 - \gamma)^3 \\ (219.1 - \gamma)^3 \\ (296.7 - \gamma)^3 \\ (367.8 - \gamma)^3 \\ (378.2 - \gamma)^3 \\ (403.3 - \gamma)^3 \\ (424.5 - \gamma)^3 \\ (436.6 - \gamma)^3 \\ (445.1 - \gamma)^3 \\ (446.8 - \gamma)^3 \\ (448.0 - \gamma)^3 \end{bmatrix},$$

$$\text{then } (X^T X)^{-1} \approx \begin{bmatrix} 0.498810 & -0.019643 \\ -0.019643 & 0.000893 \end{bmatrix},$$

$$\text{and } \begin{bmatrix} \beta \\ \alpha \end{bmatrix} = \begin{bmatrix} c_0 \\ c_1 \end{bmatrix} = \vec{c} = (X^T X)^{-1} X^T \vec{y} \approx \begin{bmatrix} -15068952.7 \\ 809891.1 \end{bmatrix}.$$

Therefore,  $\alpha \approx 809891.1$ ,  $\beta \approx -15068952.7$  and  $\gamma \approx 219.1$ .

## PART B

- B1)** Let  $A_i$  and  $C_i$  be the binary variables which is 1 if the  $i$ -th base of the 'toy' primer is  $A$  and  $C$  respectively, and 0 otherwise. Let  $S$  be the set  $\{1, 2, 3, \dots, 10\}$ . The integer linear program would be

$$\left. \begin{array}{ll} \max & 0 \\ \text{subject to:} & \\ \text{(Each base is either an } A \text{ or a } C\text{)} & A_i + C_i = 1 \quad \forall i \in S \\ \text{(The number of } C\text{'s is between 4 and 6)} & 4 \leq \sum_{i=1}^{10} C_i \leq 6 \\ \text{(There are 2 } C\text{'s among the last 5 bases)} & \sum_{i=6}^{10} C_i = 2 \\ \text{(Primer has 10 bases)} & A_i, C_i \in \{0, 1\} \quad \forall i \in S \end{array} \right\}$$

An example solution of the aforementioned integer linear program is *CCAAAAACAC*, which is attached in the Excel file.

- B2)** Let  $A_i$ ,  $C_i$ ,  $T_i$  and  $G_i$  be the binary variables which is 1 if the  $i$ -th base of the forward primer is  $A$ ,  $C$ ,  $G$  and  $T$  respectively, and 0 otherwise. Let  $S$  be the set  $\{1, 2, 3, \dots, 18\}$ . The integer linear program would be

$$\left. \begin{array}{ll} \max & 0 \\ \text{subject to:} & \\ \text{(Each base is either an } A, \text{ a } C, \text{ a } G \text{ or a } T\text{)} & A_i + C_i + G_i + T_i = 1 \quad \forall i \in S \\ \text{(Total } C\text{'s and } G\text{'s is between 7 and 11)} & 7 \leq \left( \sum_{i=1}^{18} C_i + \sum_{i=1}^{18} G_i \right) \leq 11 \\ \text{(At most 3 } C\text{'s and } G\text{'s in the last 5 bases)} & \left( \sum_{i=14}^{18} C_i + \sum_{i=14}^{18} G_i \right) \leq 3 \\ \text{(Similar annealing temperature)} & 4 \times \left( \sum_{i=1}^{18} C_i + \sum_{i=1}^{18} G_i \right) + 2 \times \left( \sum_{i=1}^{18} A_i + \sum_{i=1}^{18} T_i \right) = 58 \\ \text{(Primer has 18 bases)} & A_i, C_i \in \{0, 1\} \quad \forall i \in S \end{array} \right\}$$

A solution forward primer would be *CCCCCCCCCCCCAAAAAA*, which is attached in the Excel file.

- B3)** We decided to implement Constraints 1 and 3 in Excel, with a 'lucky bonus' of Constraint 2 being automatically satisfied without explicitly stating it in Excel's Solver.

Constraint 1 dictates that the forward primer has no 'runs' of length more than 4. We can impose this constraint in Excel by setting every set of 5 consecutive bases to be less than or equal to 4. Let  $K$  be the set  $\{1, 2, 3, \dots, 14\}$ . Thus,

$$\sum_{i=k}^{k+4} B_i \leq 4 \quad \forall k \in K, \forall B \in \{A, C, G, T\}$$

This would yield us an additional 56 constraints.

Constraint 2 dictates that the forward primer has no alternating patterns of the form  $XYXY$ . Although we did not use this constraint in Excel's Solver (cells highlighted in yellow), we could implement it by choosing two types of bases at a time for a specific base pattern and impose that the sum of every 4 consecutive bases alternating between the two types of bases has to be less than or equal to 3. For example, to eliminate the pattern *ACAC*, we choose  $A$  and  $C$  and we set the constraints

$$A_i + C_{i+1} + A_{i+2} + C_{i+3} \leq 3 \quad \forall i \in \{1, 2, 3, \dots, 15\}$$

Similar constraints are set up for the patterns *CACA*, *GTGT*, *AGAG*, *ATAT*, *CGCG*, *CTCT*, *GAGA*, *GCGC*, *TATA*, *TCTC* and *TGTG*.

This would yield us an additional 180 constraints, which is impossible to be considered by Excel. Thus, we do not explicitly add these constraints to Excel's Solver.

Constraint 3 dictates that when the forward and reverse primers are aligned side-by-side, the number of alignments between *A*'s from one primer and *T*'s from the other is less than or equal to 4, and the same goes for *C*'s and *G*'s. Since the reverse primer is a constant, we can use it to build our constraints for the forward primer. First, we note down the positions of *A* and *T* bases in the reverse primer. Then, we find the sum of the binary variables of the *A* bases in the forward primer in the same positions as the *T* bases in the reverse primer, as well as the sum of the binary variables of the *T* bases in the forward primer in the same positions as the *A* bases in the reverse primer. We add these sums together which would give us the total number of *A-T* alignments between the forward and reverse primers. Then, we impose that this final sum should be less than or equal to 4. Repeat for *G-C* alignments.

Hence, we would obtain

$$A_5 + A_8 + A_{12} + A_{14} + T_3 + T_{11} + T_{17} \leq 4$$

and

$$C_1 + C_2 + C_6 + C_9 + C_{15} + G_4 + G_7 + G_{10} + G_{16} + G_{18} \leq 4$$

A possible solution for the forward primer under these 3 constraints would be *CCCCACCCCTCCCTAAAA*.

- B4)** Instead of having the constraint  $A_{18} + C_{18} + G_{18} + T_{18} = 1$ , replace it with  $A_{18} + C_{18} + G_{18} + T_{18} \leq 1$ .
- B5)** Let  $Ar_i$ ,  $Cr_i$ ,  $Tr_i$  and  $Gr_i$  be the binary variables which is 1 if the  $i$ -th base of the reverse primer is *A*, *C*, *T* and *G* respectively, and 0 otherwise.

To form a constraint for this model, firstly we multiply the binary decision variables for *A* in the forward primer with the binary decision variables for *T* in the reverse primer at the respective positions. Then, we sum up the products at each position. For this, we would get

$$\sum_{i=1}^{18} (A_i Tr_i)$$

Similarly, we multiply the binary decision variables for *T* in the forward primer with the binary decision variables for *A* in the reverse primer at the respective positions. Then, we sum up the products at each position. For this, we would get

$$\sum_{i=1}^{18} (T_i Ar_i)$$

We would then require the sum of the two sums produced to be less than or equal to 4, i.e.:

$$\sum_{i=1}^{18} (A_i Tr_i) + \sum_{i=1}^{18} (T_i Ar_i) \leq 4$$

These steps are similarly repeated for the alignment between the *C* and *G* bases.

Thus, to model the alignment constraint, we would have to implement the following two constraints:

$$\begin{aligned} \sum_{i=1}^{18} (A_i Tr_i) + \sum_{i=1}^{18} (T_i Ar_i) &\leq 4 \\ \sum_{i=1}^{18} (C_i Gr_i) + \sum_{i=1}^{18} (G_i Cr_i) &\leq 4 \end{aligned}$$