

# A Look at Data Management Systems for Climate Research

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# Rise of Machine Learning for Spatial...

The screenshot shows the header of the Nature journal website. The word "nature" is written in a large, lowercase, serif font. Below it, in a smaller, sans-serif font, is the text "International weekly journal of science". A horizontal navigation bar follows, containing links for "Home", "News & Comment", "Research", "Careers & Jobs", "Current Issue", "Archive", and "Audio & Video". Below this bar is another set of links: "Archive", "Volume 548", "Issue 7668", "News", and "Article". The "Article" link is underlined, indicating the current page.

NATURE | NEWS



## How machine learning could help to improve climate forecasts

Mixing artificial intelligence with climate science helps researchers to identify previously unknown atmospheric processes and rank climate models.

**Nicola Jones**

23 August 2017

"If you go to major modeling centers and ask them how they work, the answer won't be machine learning," says Collins. "**But it will get there.**"<sup>1</sup>

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<sup>1</sup> Nicola Jones. "How machine learning could help to improve climate forecasts". In: *Nature* 548.7668 (2017), pp. 379–380.



## APPLIED DATA SCIENCE INVITED TALKS



**Professor Vipin Kumar**

Professor

University of Minnesota

Big Data in Climate: Opportunities and Challenges for Machine Learning

Wednesday 10:00am – 12:00pm, Room 200D

[More Information](#)

*"We discuss the challenges involved in analyzing these **massive data sets** as well as opportunities they present for both advancing machine learning as well as the science of climate change..."<sup>2</sup>*

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<sup>2</sup>[Anuj Karpatne and Vipin Kumar.](#) "Big Data in Climate: Opportunities and Challenges for Machine Learning". In: *KDD* (2017), pp. 21–22.

# ... and Spatial Data Management

## Very Large Databases (**VLDB**)

VLDB 2017 



## 43<sup>rd</sup> International Conference on **Very Large Data Bases**

### Tutorials

- The Era of Big Spatial Data

Ahmed Eldawy (UC Riverside) ([eldawy@cs.ucr.edu](mailto:eldawy@cs.ucr.edu))

Mohamed Mokbel (University of Minnesota) ([mokbel@cs.umn.edu](mailto:mokbel@cs.umn.edu))

[Slides](#)

*"In this tutorial, we present the recent work in the database community for handling **Big Spatial Data**..."<sup>3</sup>*

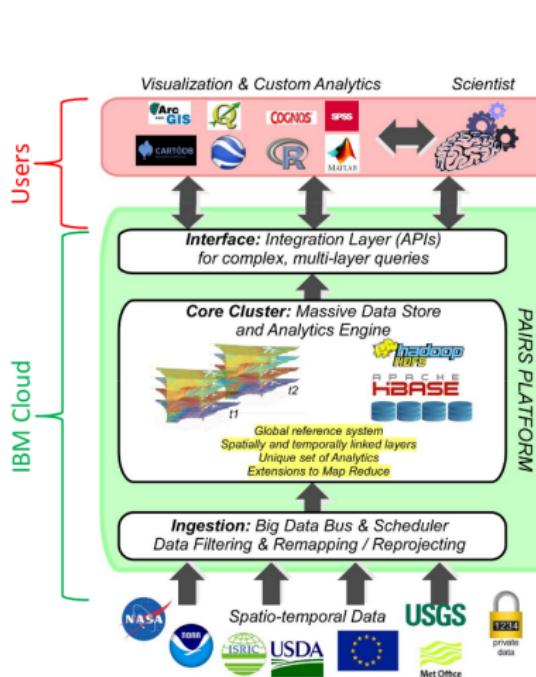
<sup>3</sup> Ahmed Eldawy and Mohamed Mokbel. "The Era of Big Spatial Data". In: VLDB (2017).

# Why is everyone so excited?

Meteorology in the eyes of data scientists:

- ▶ Terabytes of Earth imagery data get generated *per day*
  - ▶ Static analysis “macroscale” (e.g. deforestation, land-use, urban expansion, ...)
  - ▶ Real-time analysis “microscale” (e.g. flood monitoring, wildfires, landslides, tornados, ...)
  - ▶ Somewhere in between “mesoscale” (e.g. soil moisture, crop production, water availability, ...)
- ▶ Discovery and Monitoring: mine data to discover new climate relationships (e.g. teleconnections, tripoles), monitor real-time interactions (e.g. wildfires, floods, effects on rainfall...)
- ▶ Model Improvement: refine climate models to improve predictive and reproductive power
- ▶ **Data Management:** ingest, store, query, ... to efficiently support data applications

# Example: IBM Physical Analytics Integrated Repository and Services (PAIRS)<sup>4</sup>



- ▶ E.g. “Give me two years of temperature in Kyoto, plus 2-week forecast”

<sup>4</sup>Siyuan Lu. “IBM PAIRS - A Big Physical Data Service to Accelerate Analytics and Discovery”. In: (2017).

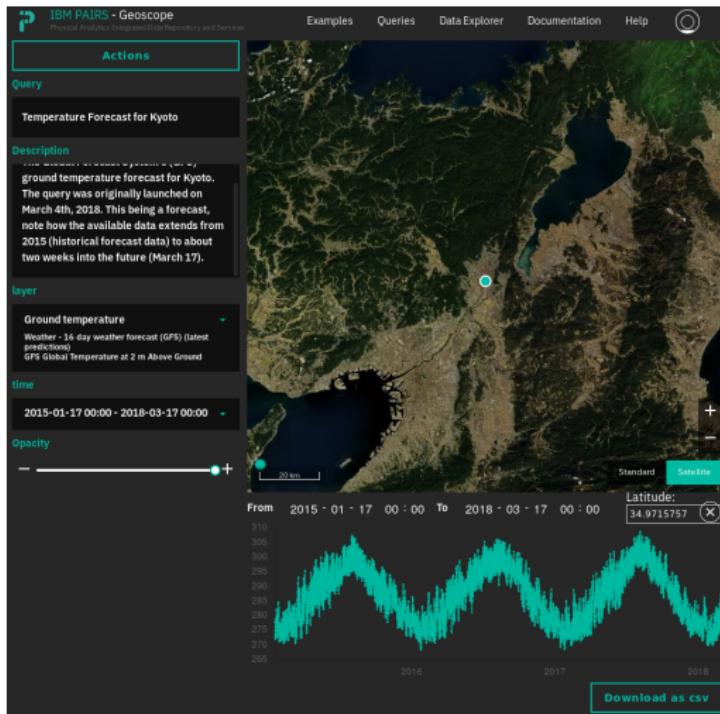


Figure: The IBM Pairs web interface.

- ▶ Good for simple retrievals; limited processing capability
- ▶ No real-time monitoring
- ▶ Limited modeling capability (limited forecasting ability)

# Example: ESRI ArcGIS<sup>5</sup>

**Policymaker:** “How is land use changing in my district?”

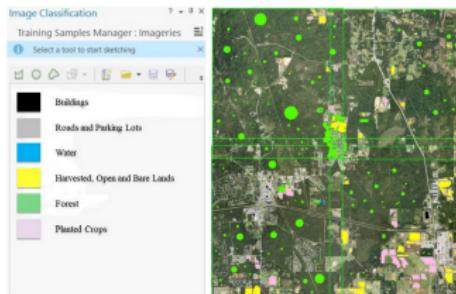


Figure 4. Training Samples Manager and Sampled Labels

**Figure:** Train a convolution neural network based on expert labels, then feed new images per pixel and collect the predicted land use.



Figure 1. Raw NAIP Images and Classified Images

Class	Precision	Recall
<b>Buildings</b>	82.50	81.28
<b>Roads</b>	84.78	85.13
<b>Waters</b>	86.14	85.55
<b>Harvested</b>	90.38	91.88
<b>Planted</b>	89.05	88.19
<b>Forest</b>	91.46	92.65

Table 1. Accuracy Assessment of U-Net Model (Precision and Recall in %)

<sup>5</sup>Amin Tayyebi. “High-Resolution Land Cover Mapping using Deep Learning Models.” (2018).

# Overview

## Introduction

## Background: Filesystems and Data Management

## Management Systems for Spatial Data

Array-Based: SciDB and friends

Spark-Based: GeoTrellis and ClimateSpark

## Scalable Machine Learning

## Summary

# Background: Filesystems and Data Management

Why use data management software (databases) anyway? Isn't the filesystem good enough?

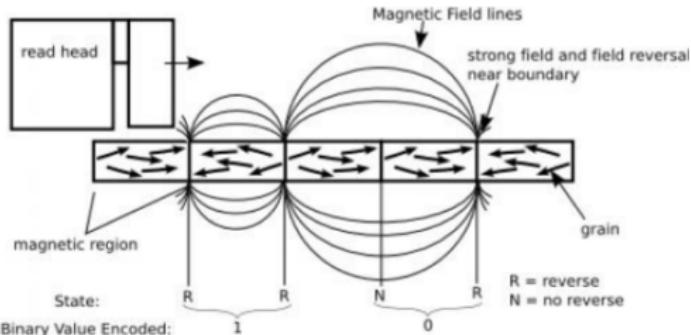


Figure: Magnetic hard drives store bits (0's and 1's) inside magnetic fields. A chunk of data (typically 512 bytes) is called a *sector* and is the minimum size the drive can write.

- ▶ Reading a particular “file” involves spinning the disk to visit sectors containing the file contents.
- ▶ Sometimes sectors belonging to one file are not physically near each other on the disk, known as “fragmentation”.

# Comparison of Read Latencies

- ▶ **Magnetic disk** at 15,000 RPM: up to 4 milliseconds (4,000,000 ns)
- ▶ **Solid-state drive** (NAND): 50,000–100,000 ns
- ▶ **RAM**: about 60 ns
- ▶ CPU L2 cache: about 10 ns
- ▶ CPU access: about 1 ns

Reading from disk is about  $10^6$  times slower than reading from RAM. In other words, relative difference between **1 second** and **11 days**.

# Filesystems and Data Management

Both filesystems and databases deal with storage and I/O:

- ▶ Filesystems: ext4 (Unix), NTFS (Windows), APFS (Apple)
- ▶ Fairly uniform use case: fast file access, security, data integrity, compression, paging, volume resizing, locking
- ▶ E.g. journaling systems such as NTFS keep a master log file to support data restore in case of crashes
- ▶ Databases: Oracle, SQL Server, Postgres, Teradata, Vertica, MongoDB, MonetDB, kdb+ ...
- ▶ **Data contents are related**, e.g. “find total revenue of the Rome store over all Mondays of last year”
- ▶ **“No one size fits all”**: different approaches for different use cases

# No One Size Fits All: Row-Store vs. Column-Store

Consider the I/O cost of reading/writing this table:

<i>Item</i>	<i>Category</i>	<i>Revenue</i>
<i>Glove</i>	<i>Sport</i>	500
<i>Cap</i>	<i>Sport</i>	200
<i>Chair</i>	<i>Housing</i>	450
<i>Table</i>	<i>Housing</i>	100
<i>Shoe</i>	<i>Sport</i>	600

---

Row-store (e.g. H-Store<sup>6</sup>): {Glove, Sport, 500, Cap, Sport, 200, Chair, Housing, ...}

---

Column-store (e.g. MonetDB<sup>7</sup>): {Glove, Cap, Chair, ..., Sport, Sport, Housing, ..., 500, 200, 450, ...}

<sup>6</sup> Robert Kallman et al. "H-store: a high-performance, distributed main memory transaction processing system". In: VLDB 1.2 (2008), pp. 1496–1499.

<sup>7</sup> Stratos Idreos et al. "MonetDB: Two Decades of Research in Column-oriented Database Architectures". In: IEEE Data Eng. Bull. 35.1 (2012), pp. 40–45.

# Management Systems for Spatial Data: Array Databases

Consider the cost of the blue subarray query:

## Relational Database

<i>i</i>	<i>j</i>	value
0	0	32.5
1	0	90.9
2	0	42.1
3	0	96.7
0	1	46.3
1	1	35.4
2	1	35.7
3	1	41.3
0	2	81.7
1	2	35.9
2	2	35.3
3	2	89.9
0	3	53.6
1	3	86.3
2	3	45.9
3	3	27.6

## Array Database

32.5	46.3	81.7	53.6
90.9	35.4	35.9	86.3
42.1	35.7	35.3	45.9
96.7	41.3	89.9	27.6

## Traditional RDBMS:

- ▶ Basic types (integer, float, string, ...)
- ▶ Relational operators (select, join, group by, ...)
- ▶ Good for queries such as “What are the phone numbers for all employees in Delaware?”

<i>ID</i>	<i>Name</i>	<i>Location</i>	<i>PhoneNumber</i>
1	<i>JohnSmith</i>	<i>Delaware</i>	123 – 456 – 7890
2	<i>JaneDoe</i>	<i>Delaware</i>	123 – 456 – 7891
...	...	...	...

## Array-Based DBMS:

- ▶ **Array** data model
- ▶ Array operators (structural operators, content operators)
- ▶ Good for multidimensional queries such as point time-range query, spatial aggregation query

2	5	4
2	1	8
...	...	...

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<sup>8</sup>The SciDB Development Team. “Overview of SciDB”. In: *SIGMOD* (2010).

# SciDB: Chunk Storage

[ 0 ]	[ 1 ]	[ 2 ]	[ 3 ]	[ 4 ]	
[ 0 ]	( 2, 0.7 )	( 5, 0.5 )	( 4, 0.9 )	( 2, 0.8 )	( 1, 0.2 )
[ 1 ]	( 5, 0.5 )	( 3, 0.5 )	( 5, 0.9 )	( 5, 0.5 )	( 5, 0.5 )
[ 2 ]	( 4, 0.3 )	( 6, 0.1 )	( 6, 0.5 )	( 2, 0.1 )	( 7, 0.4 )
[ 3 ]	( 4, 0.25 )	( 6, 0.45 )	( 6, 0.3 )	( 1, 0.1 )	( 0, 0.3 )
[ 4 ]	( 6, 0.5 )	( 1, 0.6 )	( 5, 0.5 )	( 2, 0.15 )	( 2, 0.4 )

Step 1: Vertically partition *attributes* in the *logical array*.

{ A }					{ B }				
2	5	4	2	1	0.7	0.5	0.9	0.8	0.2
5	3	5	5	5	0.5	0.5	0.9	0.5	0.5
4	6	6	2	7	0.3	0.1	0.5	0.1	0.4
4	6	6	1	0	0.25	0.45	0.3	0.1	0.3
6	1	5	2	2	0.5	0.6	0.5	0.15	0.4

Step 2: Decompose each attribute array into equal sized, and potentially overlapping, *chunks*.

{ A <sub>1</sub> }	{ A <sub>2</sub> }	{ A <sub>3</sub> }	{ A <sub>4</sub> }
2 5 4	4 2 1	4 6 6	6 2 7
5 3 5	5 5 5	4 6 6	6 1 0
4 6 6	6 2 7	6 1 5	5 2 2

SciDB  
Figure 5: SciDB Storage Manager

- ▶ Designed to support spatial operators (e.g. Gaussian smoothing)
- ▶ Chunk partitioning motivated by typical access patterns

# Benchmarks<sup>9</sup>

**MySQL** compared with **SciDB**; Cluster size 10 nodes, 2GB Ram + 3.2 GHz CPU per node

DBMS	Dataset	Loading/Cooking [min]				Query Runtimes [min]									
		Load	Obsv	Group	Total	Q1	Q2	Q3	Q4	Q5	Q6	Q7	Q8	Q9	Total
MySQL	<i>small</i>	760	110	2	872	123	21	393	0.4	0.36	0.6	0.6	49	50	638
	<i>normal</i> (scaleup)	770	200	90	1060	54	44	161	50	32	51	52	395	395	1234
SciDB	<i>small</i>	34	1.6	0.6	36	8.2	0.2	3.7	0.007	0.01	0.01	0.01	1.8	1.9	16
	<i>normal</i> (scaleup)	67	1.9	15	84	3.6	0.07	1.7	0.015	0.017	0.02	0.11	2.2	2.3	10
(MySQL /SciDB)	<i>small</i>	(22)	(69)	(3.3)	(24)	(15)	(105)	(106)	(57)	(36)	(60)	(60)	(27)	(26)	(40)
	<i>normal</i>	(12)	(105)	(6)	(13)	(15)	(630)	(95)	(3330)	(1880)	(2550)	(470)	(180)	(170)	(120)

TABLE I

BENCHMARK RESULTS. (*num*) IS A RATIO OF RUNTIMES, EITHER *normal* VS. *small* (SCALEUP) OR MySQL VS. SciDB.

- ▶ “small”: single-machine, 160 3750x3750 images, 99 GB
- ▶ “normal”: distributed, 400 7500x7500 images, 990 GB

**Takeaway:** SciDB achieves 2-orders speedup compared to MySQL across the 9 benchmark queries on the cluster

<sup>9</sup>Philippe Cudre-Mauroux et al. “SS-DB: a standard science DBMS benchmark”. In: *XLDB* (2012).

# Application: EarthDB (MODIS)<sup>10</sup>

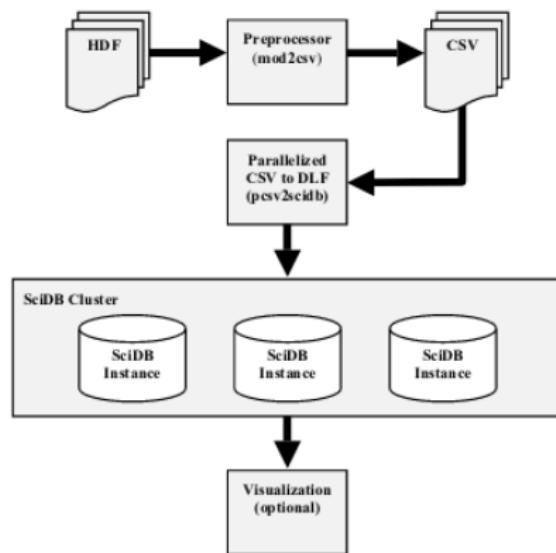


Figure 3: EarthDB High-Level Data Flow

- ▶ Extends SciDB to support MODIS data

<sup>10</sup> Gary Planthaber, Michael Stonebraker, and James Frew. "EarthDB: Scalable Analysis of MODIS Data using SciDB". In: *SIGSPATIAL* (2012), pp. 11–19.

# Application: AscotDB (Astronomy)<sup>11</sup>

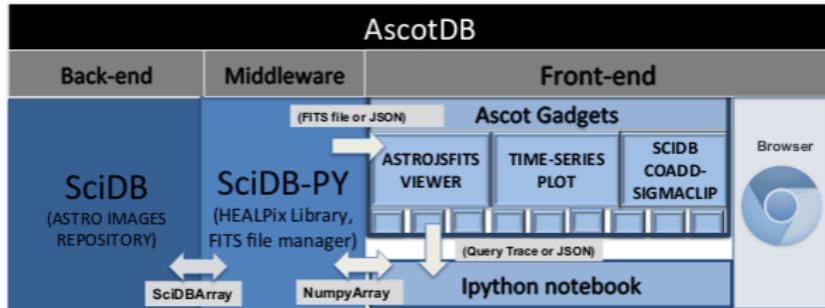


Figure 1: AscotDB architecture: SciDB as back-end, python middleware, Ascot and IPython as front-ends.

- ▶ Adds spherical support to SciDB's Cartesian operators
- ▶ Adds efficient spherical operators
- ▶ Adds Python bindings and graphical interface
- ▶ Interdisciplinary collaboration between astronomers and database experts

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<sup>11</sup> Jacob Vanderplas et al. "Squeezing a Big Orange into Little Boxes: The AscotDB System for Parallel Processing of Data on a Sphere". In: (2015).

# TileDB<sup>12</sup>

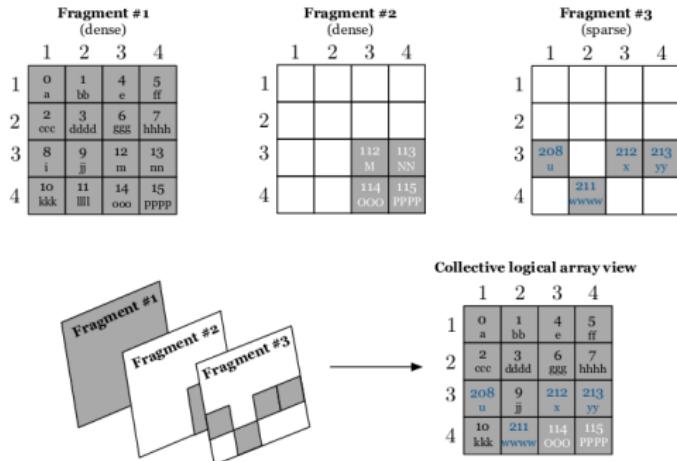


Figure 4: Fragment examples

- ▶ Better support for sparse arrays
- ▶ Avoids fixed-dimension chunking and full chunk reads
- ▶ Key idea: convert random-access writes into sequential appends by storing update **fragments**

<sup>12</sup>Stavros Papadopoulos et al. "The TileDB Array Data Storage Manager". In: VLDB (2016).

space tile extents: 2x2  
tile order: row-major  
cell order: row-major

	1	2	3	4
1	o a	1 bb	4 e	5 ff
2	2 ccc	3 dddd	6 ggg	7 hhhh
3	8 i	9 jj	12 m	13 nn
4	10 kkk	11 lll	14 ooo	15 ppp

**Files**  
(binary format)

a1.tdb [ 0 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 ]

a2.tdb [ 0 1 3 6 10 11 13 16 20 21 23 26 30 31 33 36 ]

a2\_var.tdb [ a bb ccc dddd e ff ggg hhhh i jj kkk llll m ... ]

**Figure 6: Physical organization of dense fragments**

space tile extents: 2x2  
tile order: row-major  
cell order: row-major

	1	2	3	4
1	o a	1 bb		2 ccc
2			3 dddd	
3	4 e		6 ggg	7 hhhh
4		5 ff		

**Files**  
(binary format)

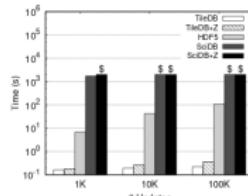
a1.tdb [ 0 1 2 3 4 5 6 7 ]

a2.tdb [ 0 1 3 6 10 11 13 16 ]

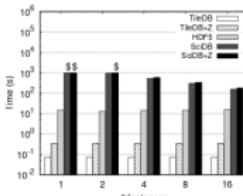
a2\_var.tdb [ a bb ccc dddd e ff ggg hhhh ]

\_\_coords.tdb [ 1,1 1,2 1,4 2,3 3,1 4,2 3,3 3,4 ]

**Figure 7: Physical organization of sparse fragments**

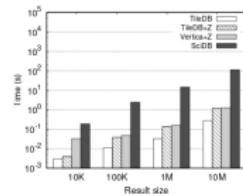


(a) vs. # updates (HDD)

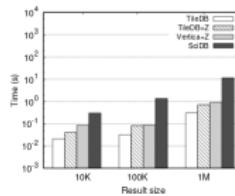


(b) vs. # instances (SSD)

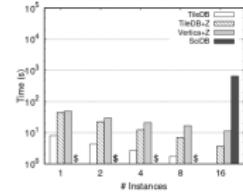
Figure 10: Random update performance of dense arrays



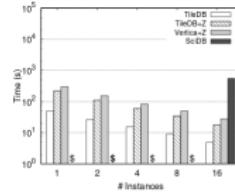
(a) DQ vs. result size (HDD)



(b) SQ vs. # result size (HDD)



(c) DQ vs. # instances (SSD)



(d) SQ vs. # instances (SSD)

Figure 14: Subarray performance for sparse arrays

- ▶ Random-access update on dense arrays is 2-order faster than HDF5, 3-order faster than SciDB
- ▶ Subarray on sparse arrays is 2-order faster than SciDB
- ▶ Dense read comparable to HDF5

# Array DB vs HDF5

## Array DB

- ▶ Is a database
- ▶ Can serve as warehouse of files from various sources
- ▶ Supports native operators, e.g. join, query
- ▶ Different flavors support different use cases, e.g. dense/sparse arrays, writes/reads, ...
- ▶ Easy scale-out on clusters
- ▶ Suitable for frequent, varied access

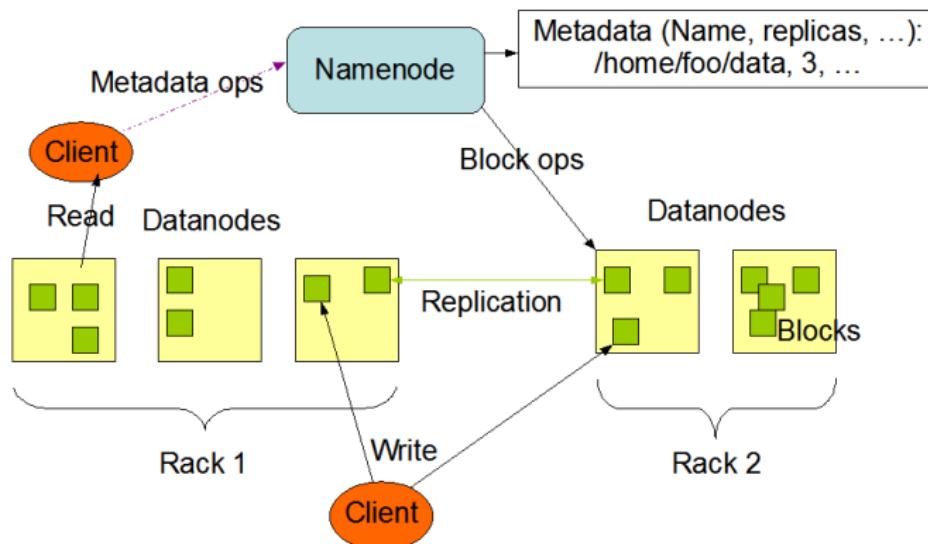
## HDF5

- ▶ Is a file format and library
- ▶ Typically stores one “data product” per file
- ▶ Operators are typically at the application level, e.g. a Python program to join two datasets
- ▶ Not meant to be distributed in a cluster
- ▶ Suitable for one-shot processing

# GeoTrellis (<http://geotrellis.io>)

- ▶ In-memory big raster analysis on top of Apache Spark
- ▶ Data is managed by Hadoop Filesystem<sup>13</sup> (HDFS) for fault-tolerance and to support map-reduce style computation
- ▶ “Moving computation is cheaper than moving data”

HDFS Architecture



<sup>13</sup> “HDFS Architecture Guide”. In: () .

# GeoTrellis: NDVI Example

$NDVI = (NIR - Red) / (NIR + Red)$  In GeoTrellis:

```
rdd.mapValues { (red, nir) => (nir - red) /  
(nir + red) } .reduceByKey(_.localMax(_))
```

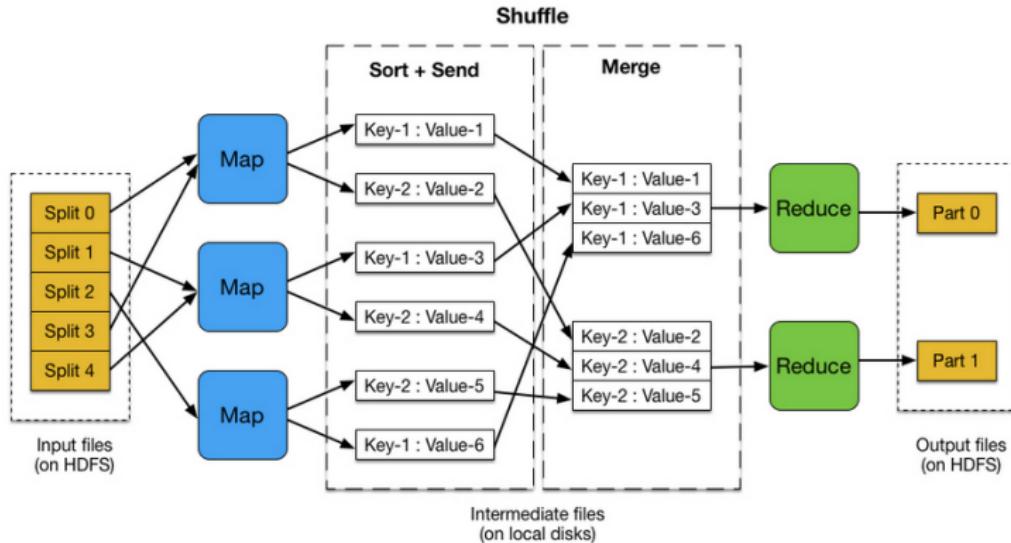
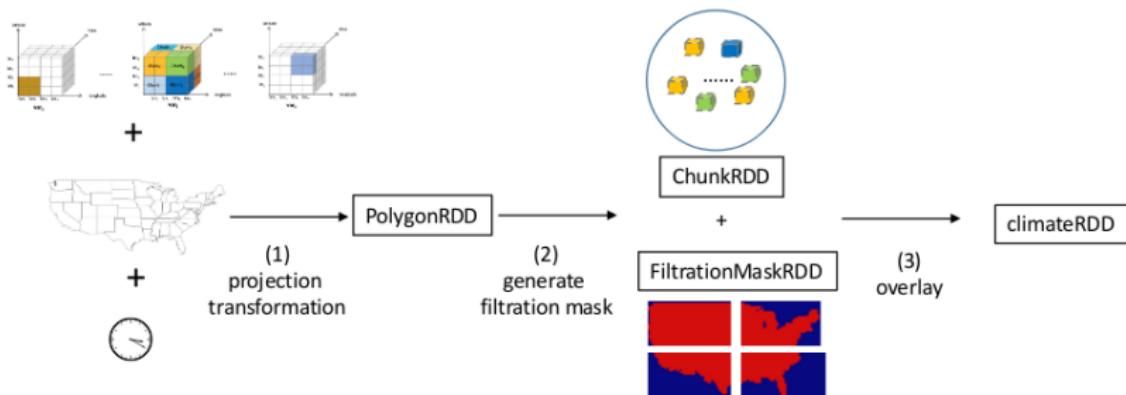


Figure: From <http://www.sunlab.org/teaching/cse6250/spring2017/lab/image/post/mapreduce-flow.jpg>

# ClimateSpark<sup>14</sup>

- ▶ Adds spatiotemporal index to improve data access
- ▶ HDF, netCDF files stored inside HDFS, then indexed, chunked, and finally converted to in-memory RDD's
- ▶ RDDs are processed in Spark via map reduce, same as GeoTrellis



<sup>14</sup>Fei Hu et al. "ClimateSpark: An in-memory distributed computing framework for big climate data and analytics". In: *Computers and Geosciences* (2018).

# ClimateSpark

- ▶ Experiment: 20 nodes, 24 CPU@2.35 GHz + 24GB RAM per node;  $1/2^\circ \times 5/8^\circ$  image resolution, roughly 9 TB MERRA2 data spanning 16 years

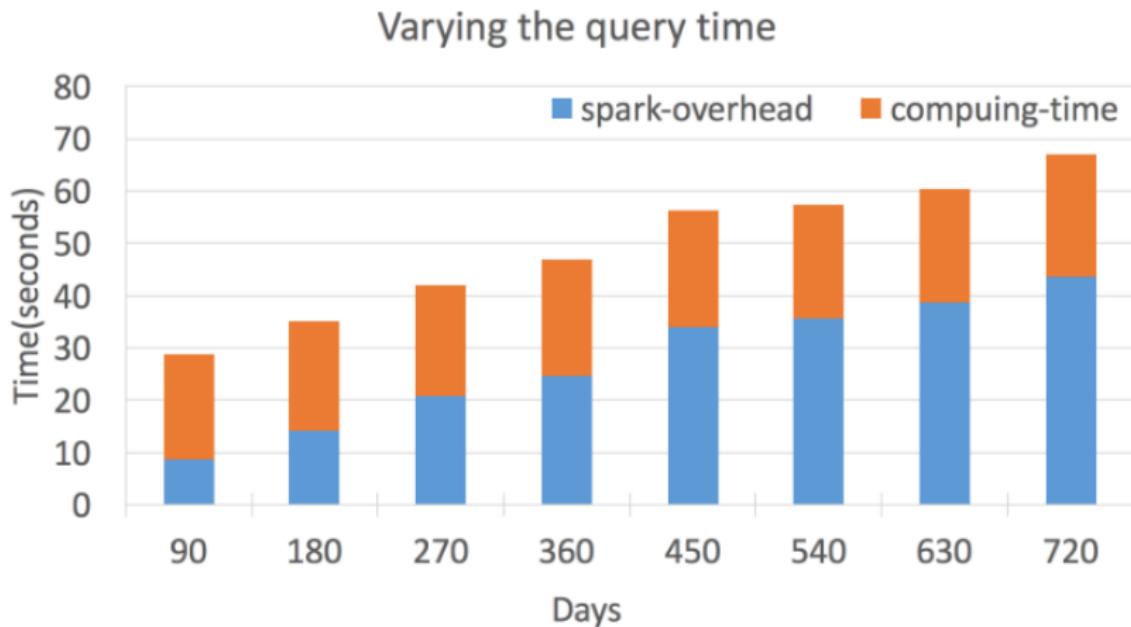


Figure 1. Run time for querying the data in Virginia and computing the monthly mean when varying the query time

# Scalable Machine Learning: MLlib

Many ML operations are a good fit for map reduce. MLLib<sup>15</sup> adds machine learning operators onto Apache Spark. Example, gradient descent:

$$w \leftarrow w - \alpha \cdot \sum_i g(w; x_i, y_i)$$

```
for (i <- 1 to n) {  
    val gradient =  
        points.map { p =>  
            (1/(1 + exp(-p.y*w.dot(p.x)) - 1)*p.y*p.x  
        ).reduce(_ + _)  
    w -= alpha*gradient  
}
```

---

<sup>15</sup>Xiangrui Meng. "MLlib: Scalable Machine Learning on Spark". In: () .

# Scalable Machine Learning: MLlib

MLlib is a standard Spark component, thus is tightly integrated with Spark data operations such as Spark SQL:

```
val trainingTable = sql("""SELECT... FROM... JOIN...""")
val training = traingTable.map { row => ... }
val model = SVMWithSGD.train(training)
```

# Scalable Machine Learning

Integrating machine learning with existing data management tools is an exciting area of research<sup>16</sup>.

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<sup>16</sup> Arun Kumar, Matthias Boehm, and Jun Yang. "Data Management in Machine Learning: Challenges, Techniques, and Systems". In: *SIGMOD* (2017).

# Ridesharing: A different kind of spatial problem

17

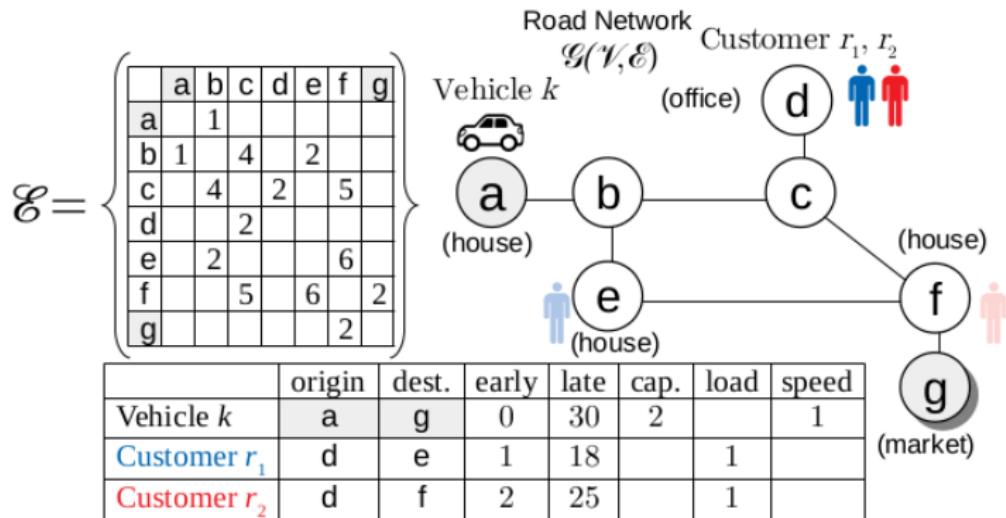


Figure 1: *Ridesharing example.*

Figure: How to optimally match vehicles to customers and design the service route?

<sup>17</sup> James Pan, Guoliang Li, and Juntao Hu. "Ridesharing: Benchmark, Simulator, and Evaluation". In: VLDB (2019).

# Ridesharing

18

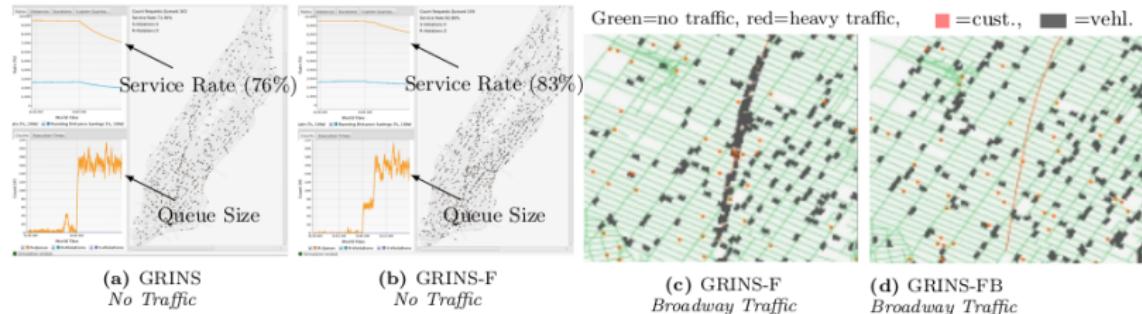


Figure 4: Jargo usage scenario.

**Figure:** How to evaluate ridesharing algorithms? Jargo can achieve real-time simulation of 1,000's vehicles on commodity machine by manipulating data of the system instead of physical simulation.

<sup>18</sup>James Pan, Guoliang Li, and Yong Wang. "Evaluating Ridesharing Algorithms using the Jargo Stochastic Simulator". In: VLDB (2020).

# Summary

Key takeaways:

- ▶ New knowledge hiding in vast amounts of data
- ▶ Data management helps work around physical limitations of storage
- ▶ Array DB (SciDB, EarthDB, TileDB) for new science applications
- ▶ Integration of ML with data management, e.g. Spark MLLib
- ▶ **Very active space**

Thank you!

Q&A