

0.0.1 Jacobi eigenvalue algorithm

A basic problem in numerical Linear Algebra is to find the eigenvalues and eigenvectors of a real-symmetric $N \times N$ matrix. An old but effective algorithm is the Jacobi eigenvalue algorithm. This algorithm uses planar rotations to systematically decrease the size of off-diagonal elements while increasing the size of diagonal elements. A bit of linear algebra is the foundation of the algorithm.

Planar rotations

Recall that a 2×2 rotation matrix has the form

$$R_2 = \begin{pmatrix} \cos(\theta) & -\sin(\theta) \\ \sin(\theta) & \cos(\theta) \end{pmatrix}.$$

We can restrict θ in the range $-\pi \leq \theta \leq \pi$. Notice that if we apply this matrix to the standard basis vectors we get

$$\begin{aligned} R_2 E_1 &= \begin{pmatrix} \cos(\theta) & -\sin(\theta) \\ \sin(\theta) & \cos(\theta) \end{pmatrix} \begin{pmatrix} 1 \\ 0 \end{pmatrix} = \begin{pmatrix} \cos(\theta) \\ \sin(\theta) \end{pmatrix}, \\ R_2 E_2 &= \begin{pmatrix} \cos(\theta) & -\sin(\theta) \\ \sin(\theta) & \cos(\theta) \end{pmatrix} \begin{pmatrix} 0 \\ 1 \end{pmatrix} = \begin{pmatrix} -\sin(\theta) \\ \cos(\theta) \end{pmatrix}, \end{aligned}$$

so as given this matrix multiplication gives a counterclockwise rotation in the plane.

The next step is to extend this rotation to two of N coordinates in \mathbb{R}^N . Pick two distinct indices m, n , with $1 \leq m < n \leq N$. Define the $N \times N$ matrix $R_{m,n}(\theta)$ which is the identity, except for replacement of the m -th row and the n -th row. For notational convenience we will sometimes suppress the argument θ , writing $R_{m,n}$ for $R_{m,n}(\theta)$. The new m, n -th rows have zero entries except for

$$R_{m,n}(m, m) = R_{m,n}(n, n) = \cos(\theta), \quad R_{m,n}(n, m) = \sin(\theta) = -R_{m,n}(m, n).$$

As an example with $N = 4$,

$$R_{2,4}(\theta) = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & \cos(\theta) & 0 & -\sin(\theta) \\ 0 & 0 & 1 & 0 \\ 0 & \sin(\theta) & 0 & \cos(\theta) \end{pmatrix}.$$

Such rotation matrices are orthogonal. That is, the inverse is the transpose, and multiplication by $R = R_{m,n}(\theta)$ does not change vector length,

$$\|X\|^2 = X \bullet X = R^{-1}RX \bullet X = R^T RX \bullet X = RX \bullet RX = \|RX\|^2.$$

Suppose A , with entries $A(i, j)$ is an $N \times N$ real symmetric matrix, that is $A = A^T$. Since in general $(AB)^T = B^T A^T$ we get

$$[R_{m,n}(\theta)^T A R_{m,n}(\theta)]^T = R_{m,n}(\theta)^T A R_{m,n}(\theta).$$

Moreover, since $R_{m,n}^T = R_{m,n}^{-1}$, this conjugation replaces A with a new symmetric matrix, but leaves the eigenvalues fixed. That is

$$AX = \lambda X, \quad AR_{m,n}R_{m,n}^{-1}X = \lambda X, \quad (R_{m,n}^{-1}AR_{m,n})R_{m,n}^{-1}X = \lambda R_{m,n}^{-1}X$$

In addition, the mapping from $A \rightarrow R_{m,n}(\theta)^T A$ leaves the lengths of the columns unchanged, while the mapping $R_{m,n}(\theta)^T A \rightarrow (R_{m,n}(\theta)^T A)R_{m,n}(\theta)$ leaves the lengths of the rows unchanged, so the square of the matrix norm

$$\|A\|^2 = \sum_{i,j=1}^N A(i, j)^2$$

is unchanged by the conjugation.

Recall that dot products with the standard basis vectors E_1, \dots, E_N give matrix entries,

$$A(i, j) = E_i \bullet AE_j.$$

This fact helps compute the entries of $R_{m,n}(\theta)^T A R_{m,n}(\theta)$. If $i \neq m, n$ and $j \neq m, n$, then multiplication by $R_{m,n}$ leaves E_i and E_j fixed, so

$$E_i \bullet R_{m,n}^T A R_{m,n} E_j = R_{m,n} E_i \bullet A R_{m,n} E_j = E_i \bullet A E_j = A(i, j).$$

That is, there is no effect on the matrix entries unless the row or column is m or n . It may help to use the 4×4 case as an illustration.

Now consider the new rows m, n and columns m, n . If the row is m and the column is j then the (m, j) entry is

$$R_{m,n} E_m \bullet A R_{m,n} E_j = (\cos(\theta) E_m + \sin(\theta) E_n) \bullet A R_{m,n} E_j.$$

Now if the row is n and the column is j then the (n, j) entry is

$$R_{m,n} E_n \bullet A R_{m,n} E_j = (-\sin(\theta) E_m + \cos(\theta) E_n) \bullet A R_{m,n} E_j.$$

If $j \neq m, n$ then $R_{m,n}E_j = E_j$.

Consider the diagonal entries of $R_{m,n}^T A R_{m,n}$. Using $A(m, n) = A(n, m)$, the (m, m) entry $B_1(\theta)$ is

$$\begin{aligned} R_{m,n}E_m \bullet A R_{m,n}E_m &= (\cos(\theta)E_m + \sin(\theta)E_n) \bullet A(\cos(\theta)E_m + \sin(\theta)E_n) \\ &= \cos^2(\theta)A(m, m) + \sin^2(\theta)A(n, n) + 2\sin(\theta)\cos(\theta)A(m, n). \\ &= A(m, m) + \sin^2(\theta)[A(n, n) - A(m, m)] + 2\sin(\theta)\cos(\theta)A(m, n) \\ &= A(m, m) + \frac{1 - \cos(2\theta)}{2}[A(n, n) - A(m, m)] + \sin(2\theta)A(m, n). \end{aligned}$$

Similarly, the (n, n) entry $B_2(\theta)$ is

$$\begin{aligned} R_{m,n}E_n \bullet A R_{m,n}E_n &= (-\sin(\theta)E_m + \cos(\theta)E_n) \bullet A(-\sin(\theta)E_m + \cos(\theta)E_n) \\ &= A(n, n) + \sin^2(\theta)[A(m, m) - A(n, n)] - 2\sin(\theta)\cos(\theta)A(m, n) \\ &= A(n, n) + \frac{1 - \cos(2\theta)}{2}[A(m, m) - A(n, n)] - \sin(2\theta)A(m, n). \end{aligned}$$

Write the sum of the squares of the new diagonal entries as $B_1(\theta)^2 + B_2(\theta)^2$. Since this function is continuous and periodic with period π , it must have maxima and minima. $B_1(\theta)^2 + B_2(\theta)^2$ is extremized when

$$2B_1 \frac{dB_1}{d\theta} + 2B_2 \frac{dB_2}{d\theta} = 0.$$

The derivative terms are

$$\frac{dB_1}{d\theta} = \sin(2\theta)[A(n, n) - A(m, m)] + 2\cos(2\theta)A(m, n) = -\frac{dB_2}{d\theta}.$$

We want either $dB_1/d\theta = 0$, or $B_1 = B_2$.

If the derivatives are zero we get

$$\tan(2\theta) = 2 \frac{A(m, n)}{A(m, m) - A(n, n)}. \quad (0.0.1)$$

If $B_1 = B_2$ then simplifying leads to

$$\cos(2\theta)[A(n, n) - A(m, m)] = 2\sin(2\theta)A(m, n),$$

and this value of θ then gives

$$B_1 = B_2 = \frac{1}{2}[A(n, n) + A(m, m)],$$

so

$$B_1^2 + B_2^2 = \frac{1}{2}[A(n, n)^2 + 2A(m, m)A(n, n) + A(m, m)^2].$$

From

$$0 \leq (\alpha - \beta)^2 = \alpha^2 - 2\alpha\beta + \beta^2, \quad 2\alpha\beta \leq \alpha^2 + \beta^2,$$

we see that

$$\frac{(A(m, m) + A(n, n))^2}{2} < A(m, m)^2 + A(n, n)^2$$

unless the values are equal. It is (0.0.1) that we want.

Jacobi algorithm

The algorithm is much easier to describe than to justify. Pick an $\epsilon > 0$, for example $\epsilon = 10^{-6}$.

Sweep: First do a *sweep*: for $n = 1, \dots, N$ and $m = 1, \dots, n$, determine θ as in (0.0.1). Replace A by $R_{m,n}(\theta)^T A R_{m,n}(\theta)$.

Iterate: At the end of the sweep, check to see if the sum of the squares of the off-diagonal elements is less than ϵ . If so, stop. Otherwise perform another sweep.

The original version of the algorithm was slightly different. You were supposed to find the maximum $|A(m, n)|$ for off-diagonal entries and then replace A by $R_{m,n}(\theta)^T A R_{m,n}(\theta)$. Then repeat until the sum of the squares of the off-diagonal elements (or some other measure of goodness) is less than ϵ . Since the original algorithm was described in the middle of the 1800's, the matrices were probably small and the maximum $|A(m, n)|$ could be seen by eye. With modern computers, large matrices, and small ϵ , the repeated searches for maxima seems inefficient.