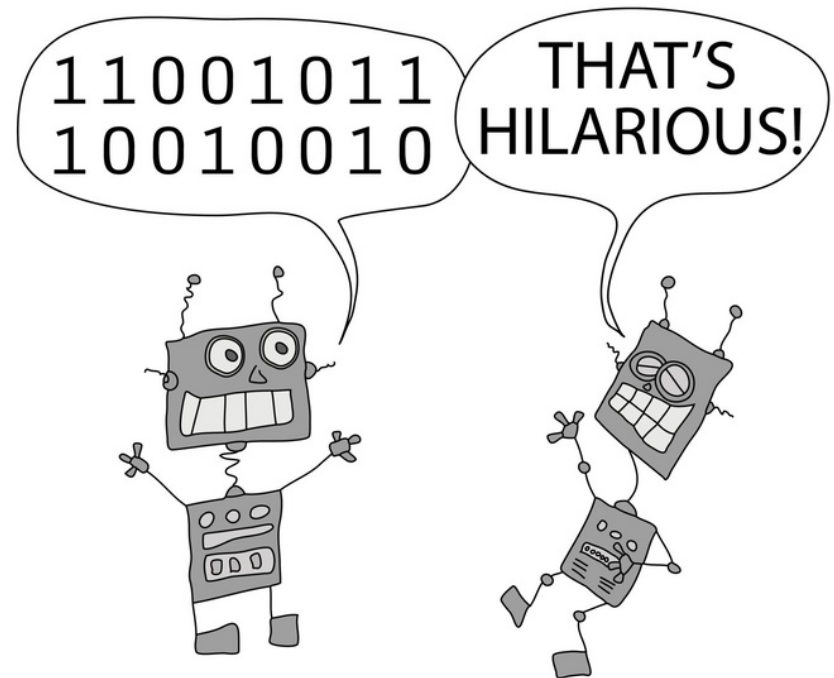


# Binary Search Trees



# Trees can be very efficient

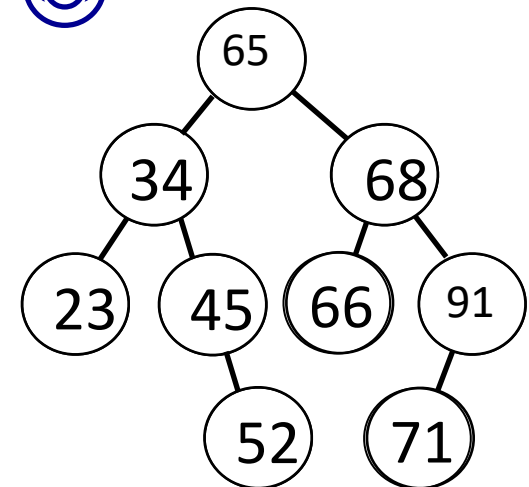
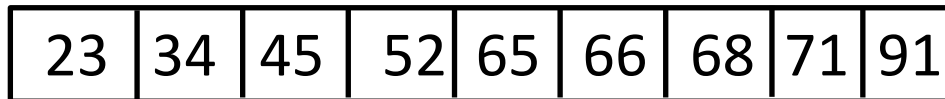
Trees are efficient. There are many algorithms which work on trees in  $O(\log n)$  time.

Usually efficiency depends on the height of the tree.

We want to make use of this efficiency and use binary trees for searching / sorting etc. – how can we do this?

OBSERVATION: For a sorted (ordered) list we could very efficiently find a key using a divide and conquer technique.

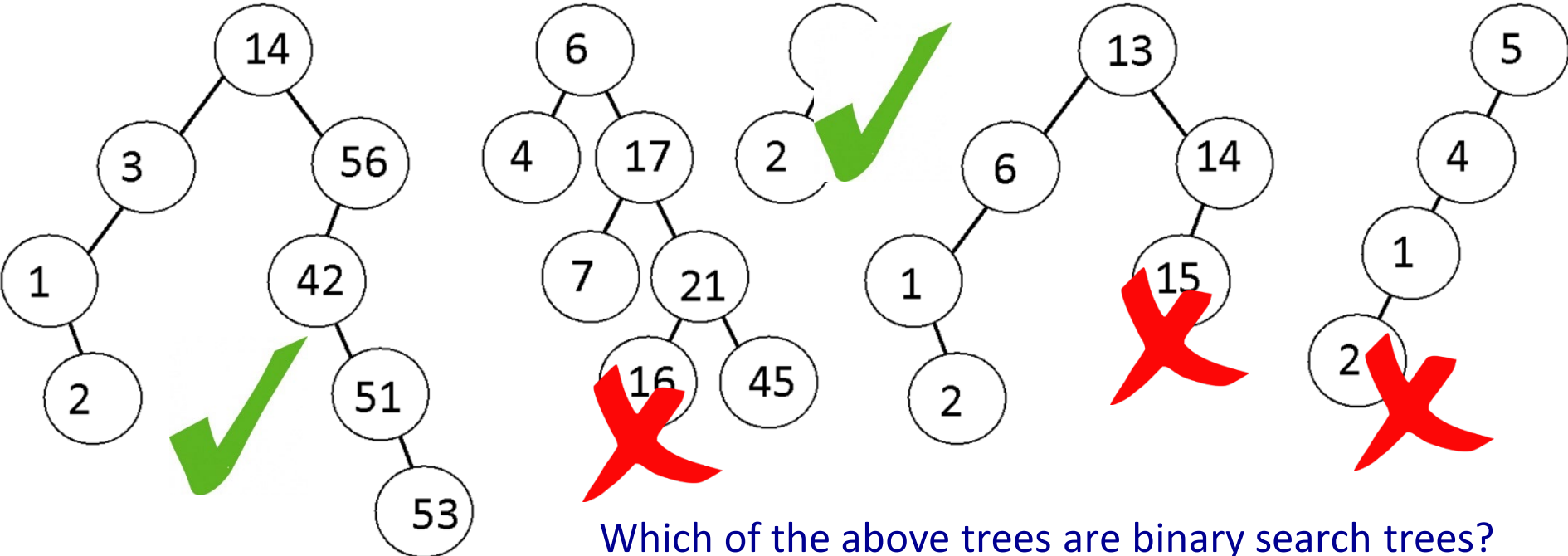
IDEA: Design trees which define an order ☺



# Binary search trees

Binary search trees are trees which have the following properties:

- For all nodes the values in the left subtree of that node are smaller than the value of the node
- For all nodes the values in the right subtree of that node are greater than the value of the node

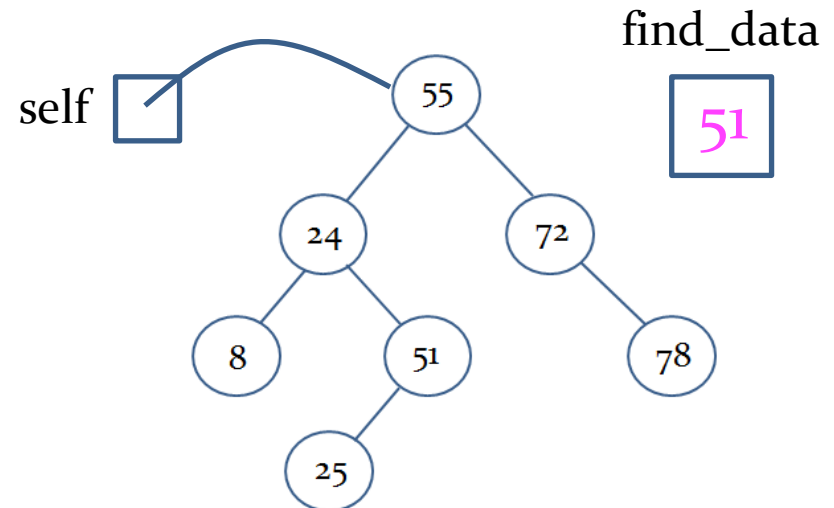


# Binary search trees – Searching

- Implement the search function for a binary search tree

**Compare the find\_data with the current root node data**

- If the two are equal, we have found the data.  
Return True**
- If the find\_data is less than the current node data, we know the find\_data either isn't in the tree or is in the left sub-tree.  
Search the left sub-tree.**
- If the find\_data is greater than the current node data.  
Search the right sub-tree**
- If we reach a node with no children and we haven't found a match for the data, it isn't in the tree.  
Return False**



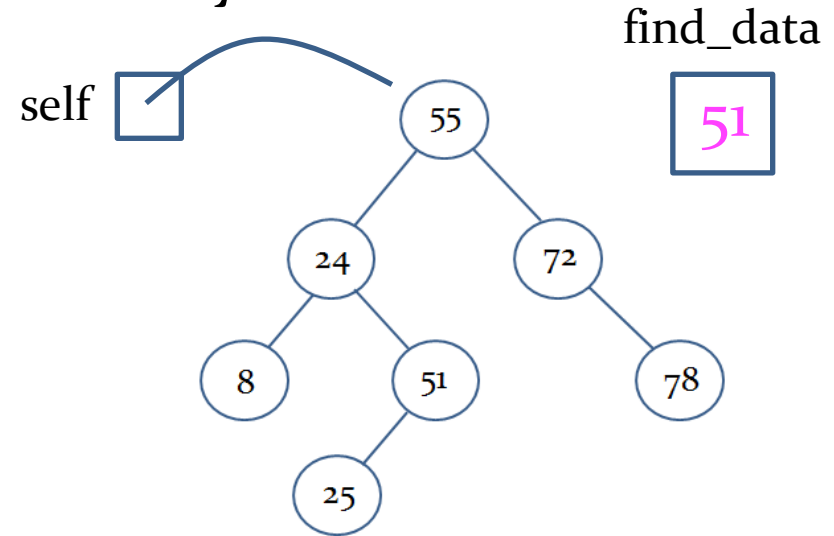
# Binary search trees – Searching

- Implement the search function for a binary search tree

```
class BinarySearchTree:
```

```
    def __init__(self, data):  
        self.__data = data  
        self.__left = None  
        self.__right = None
```

```
    def search(self, find_data):  
        if self.get_data() == find_data:  
            return True  
        elif find_data < self.get_data() and self.get_left() != None:  
            return self.get_left().search(find_data)  
        elif find_data > self.get_data() and self.get_right() != None:  
            return self.get_right().search(find_data)  
        else:  
            return False
```

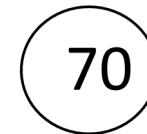


# Binary search trees - insert

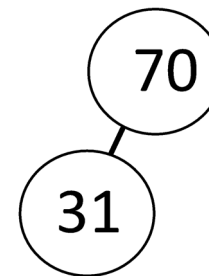
To demonstrate, we add a list of elements in the order they occur and ALWAYS MAINTAIN THE BINARY SEARCH TREE PROPERTY. For example, the following list:

70, 31, 93, 94, 14, 23, 73

**70**, 31, 93, 94, 14, 23, 73



70, **31**, 93, 94, 14, 23, 73



# Binary search trees - insert

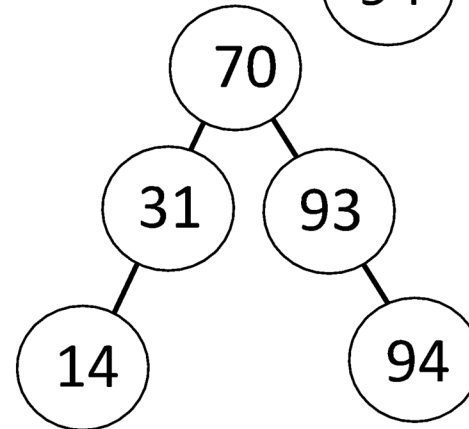
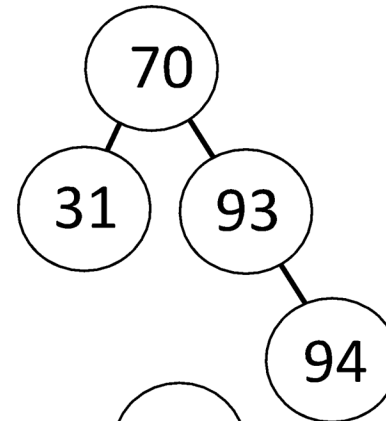
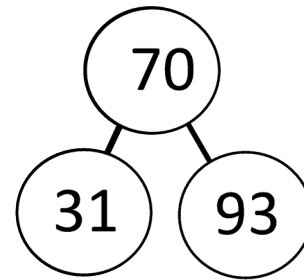
70, 31, **93**, 94, 14, 23, 73



70, 31, 93, **94**, 14, 23, 73

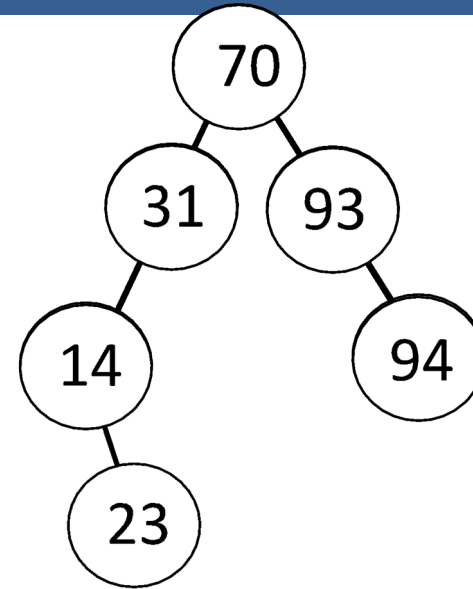


70, 31, 93, 94, **14**, 23, 73

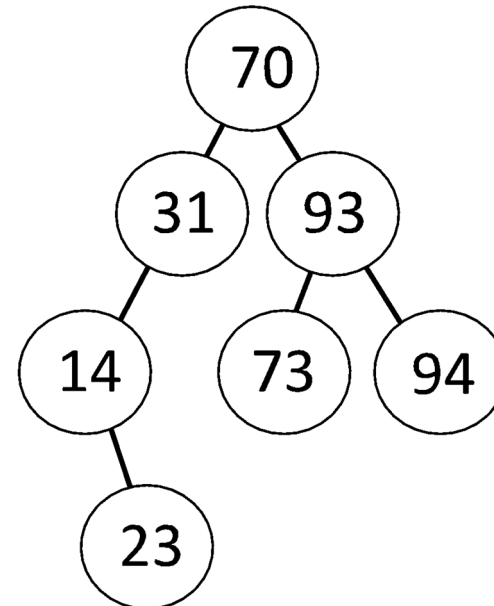


# Binary search trees - insert

70, 31, 93, 94, 14, **23**, 73



70, 31, 93, 94, 14, 23, **73**





# Binary search trees

- Implement the insert function for a binary search tree

```
class BinarySearchTree:
```

```
    def __init__(self, data):  
        self.__data = data  
        self.__left = None  
        self.__right = None
```

```
    def insert(self, new_data):  
        if new_data == self.get_data():  
            return  
        elif new_data < self.get_data():  
            if self.get_left() == None:  
                self.set_left(BinarySearchTree(new_data))  
            else:  
                self.get_left().insert(new_data)  
        else:  
            if self.get_right() == None:  
                self.set_right(BinarySearchTree(new_data))  
            else:  
                self.get_right().insert(new_data)
```

# Performance of BST

NOTE: A tree is balanced if for every node its left and right subtree vary in height by at most one

If **BST is balanced** then the height is  $O(\log n)$  and hence insert, locate, delete are all  $O(\log n)$ !

Can show that the average running times for insert, locate, delete are all  $O(\log n)$ !

Worst case is  $O(n)$  😞

BUT 😊 : Can create tree which is always balanced and hence always  $O(\log n)$  [AVL tree - not part of this lecture]

Another famous tree is the Splay tree, which has an amortised cost of  $O(\log n)$



# Advantages of BST

Compared to unsorted list:

- Insert is slightly slower ( $O(\log n)$  vs.  $O(1)$ ), but delete and find are much faster ( $O(\log n)$  vs.  $O(n)$ )

Compared to sorted list:

- Both have  $O(\log n)$  find operation, but BST can also insert and delete in  $O(\log n)$

NOTE: Can use BST for sorting (Tree Sort):  
Insert  $n$  elements and output in inorder