Comparing functional Embedded Domain-Specific Languages for hardware description

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Domain-Specific Languages

A computer language (turing-complete or *not*) targeting a specific application domain.

Example DSLs:

- SQL (database queries)
- CSS (document formatting)
- MATLAB (Matrix programming)
- VHDL (Hardware description)

Domain-Specific Languages

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Domain-Specific Languages

A computer language (turing-complete or *not*) targeting a specific application domain.

Example DSLs:

- SQL (database queries)
- CSS (document formatting)
- MATLAB (Matrix programming)
- VHDL (Hardware description)

A DSL can also be *embedded* in a general-purpose language.

Example EDSLs:

- ▶ Boost.Proto (C++ / parser combinators)
- Diagrams (Haskell / programmatic drawing)
- Parsec (Haskell / parser combinators)

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Example of an EDSL: Parsec

A simple parser for a "Game of Life"-like input format:

```
dead, alive :: Parser Bool
dead = fmap (const False) (char '.')
alive = fmap (const True) (char '*')
line :: Parser [Bool]
line = many1 (dead <|> alive)
board :: Parser [[Bool]]
board = line 'endBy1' newline
parseBoardFromFile :: FilePath -> IO [[Bool]]
parseBoardFromFile filename = do
    result <- parseFromFile board filename
    return $ either (error . show) id result
```

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Hardware EDSLs

An EDSL used for hardware design-related tasks. Can encompass:

- Modelling / description
- Simulation (validation)
- Formal verification
- Synthesis to other (lower-level) languages

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Example of a hardware EDSL

Some Lava code...

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Chosen EDSLs

The language we chose to evaluate, with the respective host language, were:

- ▶ Lava (Haskell chalmers-lava dialect)
- ForSyDe (Haskell)
- Coquet (Coq interactive theorem prover)

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Evaluation criteria

- Simulation
- Verification
- Genericity
- Depth of embedding
- Tool integration
- Extensibility

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Choice criteria

- Not too simple, not too complex
- Familiar to any hardware designer
 - No signal processing, etc.
- Well-defined, pre-specification
 - · Results to verify the models against

Choice

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Chosen circuits

We cherry-picked circuits from the book "Elements of Computing Systems", as they satisfied all of our demands.

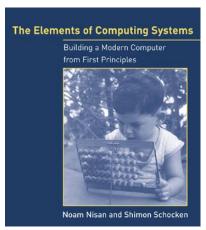


Figure: "Elements of Computing Systems" - Nisan, Schocken, available at http://www.nand2tetris.org.

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Chosen circuits

Circuit 1 A 2-input, 16-bit-wide, simple ALU

Circuit 2 A 64-word long, 16-bit wide memory block

Circuit 3 An extremely reduced instruction set CPU, the Hack CPU.

Let's take a quick look at each of these circuit's specification...

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Circuit 1: ALU

Some of the circuit's key characteristics:

- ▶ 2 operand inputs and 1 operand output, each 16-bit wide
- ▶ 1 output flag
- ▶ Can execute 18 different *functions*, among which:
 - · Addition, subtraction
 - Bitwise AND / OR
 - Constant outputs
 - · Addition of constants to an operand
 - Sign inversion

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Circuit 1: block diagram

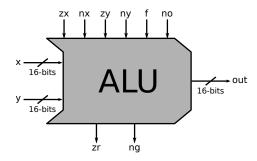


Figure: Input/Output ports of circuit 1, the ALU.

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Circuit 1: specification

The behaviour of the ALU is specified by the values of the *control bits* and *flags*:

Formal definition and test cases in the book.

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Circuit 2: RAM64

Some of the circuit's key characteristics:

- Sequential circuit, with clock input
- ▶ 64 memory words stored, each 16-bit wide
- ▶ Address port has width log₂ 64 = 6 bit

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Circuit 2: block diagram

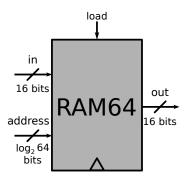


Figure: Input/Output ports of *circuit 2*, the RAM64 block.

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Circuit 2: specification

- The output "out" holds the value at the memory line indicated by "address".
- ► Iff "load" = 1, then the value at input "in" will be loaded into memory line "address".
- ► The loaded value will be emitted on "out" at the *next* clock cycle.

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Circuit 3: Hack CPU

- A very reduced instruction set CPU
 - Only 2 instructions: "C" and "A"
- ▶ Follows the Harvard architecture
 - Separate address spaces for data and instruction memory.
- Instructions are 16-bits wide
 - · As well as the memory input and output
- ► Two internal registers: "D" and "A"

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Circuit 3: block diagram

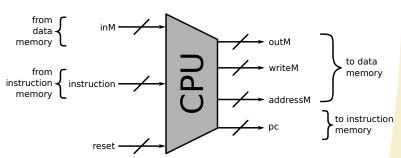


Figure: Input/Output ports of circuit 3, the Hack CPU.

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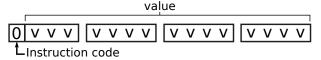
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Circuit 3: specification

Circuit 3 runs "A" and "C" instructions, according to the Hack assembly specification.

▶ The "A" instruction: sets the "A" register.



- The value in "A" can be used:
 - As operand for a subsequent computation
 - As address for jumps

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Circuit 3: specification

Circuit 3 runs "A" and "C" instructions, according to the *Hack* assembly specification.

► The "C" instruction: sets the "C" register, performs computation or jumps.

comp	•	dest	jump
1 x x a c1 c2 c3 c4	c5 c6 d1	d2 d3	j1 j2 j3
LInstruction code	-		

- Some peculiarities:
 - Bits "c1" to "c6" control the ALU
 - · conditional or unconditional jumps
 - destination of the computation result: "A", "D", "M"

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Circuit 3: specification (parts)

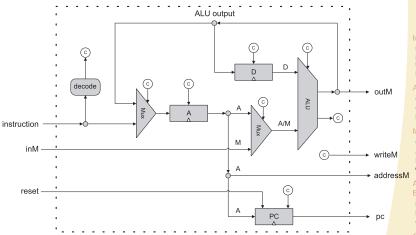


Figure: Parts used to build the *Hack* CPU, and their interconnection.

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Lava: Adders

```
type SB = Signal Bool
halfAdder :: (SB, SB) -> (SB, SB)
halfAdder inputs = (xor2 inputs, and2 inputs)
fullAdder :: (SB, (SB, SB)) -> (SB, SB)
fullAdder (cin, (a, b)) = (s, cout)
   where
      (ab, c1) = halfAdder (a, b)
      (s, c2) = halfAdder (ab, cin)
          = or2 (c1. c2)
      cout
rippleCarryAdder :: [(SB, SB)] -> [SB]
rippleCarryAdder ab = s
    where (s, _) = row fullAdder (low, ab)
```

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Lava: Simulation and verification

A taste of simulation in Lava:

- Cannot be easily automated: equality of Signal is non-trivial
- And verification...

```
prop_FullAdderCommutative :: (SB, (SB, SB)) -> SB
prop_FullAdderCommutative (c, (a, b)) =
   fullAdder (c, (a, b)) <==> fullAdder (c, (b, a))
```

-- satzoo prop_FullAdderCommutative

- Used in conjunction with an external SAT solver (e.g. Satzoo)
- Only verifies instances of specific size

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Lava: ALU

```
type ALUControlBits = (SB, SB, SB, SB, SB, SB)
alu :: ([SB], [SB], ALUControlBits) -> ([SB], SB, SB)
alu (x, y, (zx, nx, zy, ny, f, no)) = (out', zr, ng)
   where x'
                = mux (zx, (x, replicate (length x) low))
                = mux (nx, (x', map inv x'))
          ٧,
                = mux (zy, (y, replicate (length x) low))
               = mux (ny, (y', map inv y'))
          out
                = let xy'' = zip x'' y''
                  in mux (f, (andl xy'', adder xy''))
          out'
                = mux (no, (out, map inv out))
                = foldl (curry and2) low out'
          zr
                = equalBool high (last out')
          ng
          adder = rippleCarryAdder
```

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Remarks

- Cannot introduce new, meaningful datatypes
 - Only Signal Bool is synthesizable
 - Or tuples/lists thereof
- ▶ Input/Output types have to be *uncurried*
- Weak type-safety over the inputs/outputs
 - · Working with tuples is tiresome and has limitations
 - Lists don't enforce size constraints

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Lava: RAM64

```
reg :: (SB, SB) -> SB
reg (input, load) = out
    where dff = mux (load, (out, input))
          out = delay low dff
regN :: Int -> ([SB], SB) -> [SB]
regN n (input, load) = map reg $ zip input (replicate n load)
ram64Rows :: Int -> ([SB], (SB,SB,SB,SB,SB,SB), SB) -> [SB]
ram64Rows n (input, addr, load) = mux64WordN n (addr, regs)
   where memLine sel = regN n (input, sel <&> load)
                      = map memLine (decode6To64 addr)
          regs
```

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Remarks

Positive:

- Uses host language for binding (let/where) and recursion
- Uses host language for structural combinators

Negative:

- Again, weak type-safety of lists, can cause problems at simulation
- ▶ No modularity in the generated VHDL code.

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Remarks

Could benefit from:

- Fixed-length vectors
 - ForSyDe-style or with type-level naturals in recent GHC.

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ForSyDe: ALU (non-synth)

```
type S = Signal
type Word = Int16
data ALUOp = ALUSum | ALUAnd
    deriving (Typeable, Data, Show)
$(deriveLift1 ''ALUOp)
type ALUCtrl = (Bit, Bit, Bit, Bit, ALUOp, Bit)
type ALUFlag = (Bit, Bit)
bo, bb :: Bit -> Bool
bo = bitToBool
bb = boolToBit
```

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ForSyDe: ALU (non-synth)

```
aluFunc :: ProcFunc (ALUCtrl -> Word -> Word -> (Word, ALUFlag))
aluFunc = $(newProcFun [d]
  aluFunc' (zx,nx,zy,ny,f,no) x y =
      (out, (bb (out == 0), bb (out < 0)))
   where
     zfzw = if bo z then 0 else w
     nf n w = if bo n then complement welse w
      (xn, yn) = (nf nx \$ zf zx \$ x, nf ny \$ zf zy \$ y)
     out
              = nf no $ case f of
                         ALUSum -> xn + yn
                         ALUAnd -> xn .&. yn |] )
aluProc :: S ALUCtrl -> S Word -> S (Word,ALUFlag)
aluProc = zipWith3SY "aluProc" aluFunc
```

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ForSyDe: Muxes

```
mux2 :: S Bit -> S Word -> S Word -> S Word
mux2 = zipWith3SY "zipWith3SY" $(newProcFun [d]
  f s x y = if s == L then x else y | ])
mux2SysDef :: SysDef (S Bit -> S Word -> S Word -> S Word)
mux2SysDef = newSysDef mux2 "mux2" ["s","i1","i2"] ["o"]
mux4 :: S (FSVec D2 Bit) -> S (FSVec D4 Word) -> S Word
mux4 ss is = (mux2' "m1") (sv ! d1) m00 m01 where
  mux2' 1 = instantiate 1 mux2SysDef
          = unzipxSY "unzipSel" ss
  sv
  iv
          = unzipxSY "unzipInp" is
  m()()
          = (mux2' "m00") (sv ! d0) (iv ! d0) (iv ! d1)
  m01
          = (mux2' "m01") (sv ! d0) (iv ! d2) (iv ! d3)
mux4SysDef :: SysDef ( S (FSVec D2 Bit) -> S (FSVec D4 Word)
                    -> S Word)
mux4SysDef = newSysDef mux4 "mux4" ["s","is"] ["o"]
```

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ForSyDe: RAM64

```
reg :: S Word -> S Bit -> S Word
reg input load = out where
    out = delaySY "delay" (0 :: WordType) dff
    dff = (instantiate "mux2" mux2SysDef) load out input
ram64 :: S Word -> S (FSVec D6 Bit) -> S Bit -> S Word
ram64 input addr load = mux' addr (zipxSY "zipRows" rs) where
          = instantiate "mux" mux64SysDef
 mux'
  decoder' = instantiate "decoder" decode6To64SysDef
  reg' l = instantiate l regSvsDef
  and' l = instantiate l andSysDef
  r(s,l) = (reg' l) input ((and' (l ++ ":and")) load s)
  rs'
         = unzipxSY "unzipAddr" $ decoder' addr
                                                               ForSyDe
          = V.map r $ V.zip rs' (V.map (\n -> "r" ++ show n)
  rs
                                 (V.unsafeVector d64 [0..63]))
ram64SysDef = newSysDef ram64 "ram64" ["i", "a", "l"] ["o"]
```

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The Circuit type

```
Context {tech : Techno}
Inductive Circuit: Type -> Type -> Type :=
| Atom : forall {n m : Type} {Hfn : Fin n} {Hfm : Fin m},
             techno n m -> Circuit n m
 Plug : forall {n m : Type} {Hfn : Fin n} {Hfm : Fin m}
            (f : m -> n), Circuit n m
 Ser : forall {n m p : Type},
             Circuit n m -> Circuit m p -> Circuit n p
 Par : forall {n m p q : Type},
             Circuit n p -> Circuit m q
             \rightarrow Circuit (n + m) (p + q)
 Loop: forall {n m p : Type},
             Circuit (n + p) (n + p) -> Circuit n m
```

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```
Definition HADD a b s c: circuit ([:a] + [:b]) ([:s] + [:c])
      Fork2 ([:a] + [:b])
    |> (XOR a b s & AND a b c).
```

: = Chosen EDSLs

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```
Program Definition FADD a b cin sum cout :
    circuit ([:cin] + ([:a] + [:b])) ([:sum] + [:cout]) :=
   (ONE [: cin] & HADD a b "s" "co1")
|> Rewire (* (a, (b,c)) => ((a,b), c) *)
> (HADD cin "s" sum "co2" & ONE [: "co1"])
|> Rewire (* ((a,b), c) => (a, (b,c)) *)
|> (ONE [:sum] & OR "co2" "co1" cout).
Next Obligation. revert H; plug_def. Defined.
Next Obligation. plug_auto. Defined.
Next Obligation. revert H; plug_def. Defined.
Next Obligation. plug_auto.Defined.
```

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Coquet: Meaning relation

```
Inductive Sem : forall {n} {m},
   C \cap m \rightarrow (n \rightarrow Data) \rightarrow (m \rightarrow Data) \rightarrow Prop :=
| KAtom: forall n m {Hfn: Fin n} {Hfm: Fin m}
           (t: techno n m) i o, spec t i o -> Sem (Atom t) i o
| KSer: forall n m p (x: C n m) (y: C m p) i mid o,
         Sem x i mid -> Sem y mid o -> Sem (Ser x y) i o
 KPar: forall n m p q (x: C n p) (y: C m q) i o
             Sem x (select_left i) (select_left o)
         -> Sem y (select_right i) (select_right o)
         -> Sem (Par x v) i o
| KPlug: forall n m {Hfn: Fin n} {Hfm: Fin m} (f: m -> n) i,
          Sem (Plug f) i (Data.lift f i)
 KLoop: forall n m l (x: C (n + l) (m + l)) i o ret,
              Sem x (Data.app i ret) (Data.app o ret)
          -> Sem (Loop x) i o
```

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Coquet: Specification

```
Context {n m N M : Type}
        (Rn : Iso (n \rightarrow T) N) (Rm : Iso (m \rightarrow T) M).
Class Realise (c : Circuit n m) (R : N -> M -> Prop) :=
  realise: forall i o, Semantics c i o -> R (iso i) (iso o)
Class Implement (c : Circuit n m) (f : N -> M) :=
  implement: forall i o, Semantics c i o -> iso o = f (iso i)
```

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Coquet: Correctness proofs

```
Instance HADD_Implement {a b s c}:
    Implement (HADD a b s c) _ _
        (fun (x : bool * bool) =>
            match x with (a,b) => (xorb a b, andb a b) end).
Proof.
    unfold HADD; intros ins outs H; tac.
Qed.
```

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Coquet: How to prove correctness

```
Ltac tac :=
    rinvert;
                             (* destruct the circuit *)
    realise_all;
                             (* use the hint data-base *)
    unreify_all bool;
                             (* unreify *)
    destruct_all;
                             (* destruct the booleans *)
    intros_all;
    clear;
```

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boolean_eq.

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Introduction

Hardware design Domain-Specific Languages Hardware EDSL

Analyzed El

Choice criteria Chosen EDSLs Evaluation criteria

Modeled Circuit

Choice ALU Memory bank

Analysis of the EDSLs

Lava ForSyDe Coquet

Conclusio

Results

Future work



Thank you!

Questions?

