High Performance Fortran

Tutorial

at

Europar'99

by Fabien COELHO

Fabien Coelho

- Engineering degree in 1993
- PhD on HPF Compilation in 1996
- researcher at cole des mines de Paris, France
 - algebraic transformations to improve performance
 - focus on scientific applications
 - PhD Student: Julien Zory
- visitor at IMEC, Belgium
 - Francky Catthoor's team
 - program transformations for low power
 - focus on embedded applications

Sources

- HPF language specifications
 http://dacnet.rice.edu/Depts/CRPC/HPFF/
- The HPF Handbook
 by Koelber, Loveman, Schreiber, Steele and Zosel
 MIT Press

Tutorial outline

- What is HPF? MPI? OMP?
- History of HPF
- Presentation of HPF
- Examples of codes
- Conclusion

MPI vs HPF vs OMP

What they have in common:

- 3 letters acronymes
 - Message Passing Interface
 - High Performance Fortran
 - Open Multi Processing
- to express parallelism
- based on standard languages
 - Fortran
 - (

Why Parallelism?

So as to improve performance of applications

- inside processors:
 - super-scalar/VLIW
 different instructions at the same time
 - pipelines, vector processors
 start next computation before previous is finished
 - replication of functionnal units
- network of processors (slow network/fast processors, 2 orders!)

Parallel Hardwares: Multiple Data

- Shared Memory vs Distributed Memory
 - SM: communications/addressing by OS/HW
 - DM: YOU take care (user/compiler, OS...)
 communication and addressing are hard
 management of data distribution among processors...
- Control flow: Single vs Multiple Instruction
 - SIMD: one processor + many computing elements
 - MIMD: multiple processors

Examples of parallel machines

general purpose machines: MIMD survived

- SM-MIMD: shared memory multi-processors
 SUN servers, Intel-based mp
 not scalable wrt #procs
- DM-MIMD: IBM SP2, Cray T3E, SGI Origin 2000 scalable wrt #procs

special purpose machines: often SIMD (DSP, ASICS)

DM-SIMD: CM2, MMX (in processor)

Parallel Programming

- Control Parallelism (based on threads/processes...)
 fork/join model, synchronizations...
 - MPI: processes are spawned
 - OMP: multi threading on loops
 - Tera computer...
- Data Parallelism
 - same computation on different elements
 - Fortran 90, HPF
 - simpler codes (one flow, SIMD-like)

Data Parallelism

- manipulation of arrays as scalars
- operation can be applied in parallel

```
program dataparallel
integer, parameter:: n=1000
real, dimension(n,n):: A, B, C
A = \ldots, B = \ldots
C = 0.0
do while (MAXVAL(ABS(C-B))>0.1)
   C = B ! old B is saved
   B = MATMUL(A, B) + B
end do
end
```

Control parallelism

- fork parallel communicating 'processes'
- beware of data sharing!
- join back...
- can be hidden in the language (OMP)
 the memory is shared among processors
- addressing issues handled by OS/HW

Control parallel example

```
program controlparallelism
      real, dimension(n):: v
      integer:: count
      count = 0
!$omp parallel, do
     do i=1, n
         if (v(i).gt.floor) then
! synchronization needed...
!$omp
        atomic
            count = count + 1
         endif
      enddo
!$omp end parallel
```

Control parallelism: message passing

- fork PROCESSES (a la Unix)
- library to send and receive messages
- addressing managed by programmer

MP example

```
program messagepassing
     real, dimension(n/2):: v ! v(n)
     spawn(2,...)
     if (i am process 1) then
! v holds first half v(1) is v(1)
       senddatato(2)
     else if (i am process 2) then
! v holds second half, v(1) is V((n/2)+1)
       receivedatafrom(1)
     endif
```

Control vs Data parallelism

A dual philosophy:

- control parallel
 - distribute iterations
 - data will follow
- data parallel
 - distribute data
 - iterations will follow

Parallel programming and architectures

- Shared Memory, MIMD
 - OMP is the only standard! things are simple!
- Distributed Memory, MIMD
 - MPI: low level, everything by hand!
 - HPF: put the trouble on the compiler!

What is the trouble?

On a Distributed Memory parallel machine

- find available parallelism
 - program semantics must not be changed!
 - automatically detection of parallelism?
 - hints: language? assertions?
- must distribute data on processors!
 - otherwise no parallelism!
 - how? when? hints?
 - impact on addressing? on performance?
- non local data communication and store?

High Performance Fortran

- Data Parallel language
- based on Fortran 90/95
 the language of scientific computing
- plus directives (i.e. hints as comments)
 YOU are going to help!
 program still in Fortran
- a de facto standard
 not a ISO or ANSI standard
 however only one available!
 contribution to the standard (FORALL)

Fortran standard...

- dusty deck Fortran IV/66/77 code legacy
- engineers are used to Fortran
- software engineering improvements with Fortran 90
 - structured programming
 - data structures, allocate
 - better library of functionse.g. TRANSPOSE MATMUL
- other option: SISAL (functionnal)
- new Fortran 95: small improvements

History of HPF: the Forum

- DEC initiative end of 1991
- composition: vendors, researchers, users
- aim: to propose an industry standard
 - offer portability for users
 - as a commercial argument
- meetings

HPF versions...

- 03/1992-03/1993: HPF 1.0 + subset
 subset: selection of features, Fortran 77 based
- \bullet 04/1994-10/1994: HPF 1.1 + subset
- 01/1995-11/1996: HPF 2.0 + approved extensions
 AE: how should extensions should look like
 no more subset
- kernel HPF: efficient part of the language

Organization

- regular meetings in the USA
- specialized work groups AND mailing lists
 - data distribution, forall
 - Fortran 90, intrinsics, I/O...
- one regularly-participating organisation, one vote
- language defined by the majority...

Existing base

- Standards
 - Fortran 77
 - Fortran 90 (long awaited Fortran 8x)
 - PC Fortran (for share memory)
- Vendors products
 - CM-Fortran from Thinking Machine
 - MPP-Fortran
- Research prototypes
 - Fortran D (Ken Kennedy, Rice University)
 - Vienna Fortran (Hans Peter Zima, Vienna

University)

Ideas taken

For this existing base

- compilation directives seen as comments
- sequential program + data mapping
- data parallel array sections (F90)
- concepts of alignment and distribution
 i.e. two level mapping
- two semantics of parallel loops

Life of the Forum

- great battles: Templates or not Templates
- lost projects: I/O
- forgotten principles
 not only directives
 but extensions included in Fortran 95!
- forgotten dead lines;-)

HPF 1.0, May 3 1993

- a Fortran for parallel computers
- data parallel model
- based on Fortran 90
- data mapping directives
- parallel constructions
- intrinsic functions (reductions)
- official subset...

HPF 2.0, October 19, 1996

- hopefully the same language!
- specification fixes
- features deleted... since in ISO Fortran!
- features moved as approved extensions
- new approved extensions developped
- new feature included: REDUCTION
- no more subset

HPF Programming Model

- 2 factors influence parallel performance
 - available parallelism
 - needed communications
- HPF tunes the program:
 - parallel constructs
 - data mapping directives

Presentation of the language 2.0

- data mapping directives
 - data parallel application are data-driven
 - communications cost is prohibitive
 - emphasize data mapping
- parallelism!
 - implicit parallelism
 - parallel loopS
 - intrinsics
- other issues

Directives

- considered as Fortran comments
- can be declarative or executable
- can expand on several lines
 with F90 continuations &
- spaces are meaningful!

```
program directive
!hpf$ here is a directive
    print *, 'hello'
*hpf$ another directive &
chpf$ spawned on 2 lines
    end
```

Data Mapping Model

- Group lead by Guy L. Steele Jr
 - was Thinking Machine Corp.
 - then at SUN for JAVA
- Aims
 - specify the mapping of arrays
 - onto the processors of a DM machine.
 - understandable intuitive for users
 - optimisable for compilers
 - suitable to applications

2/3 level mapping

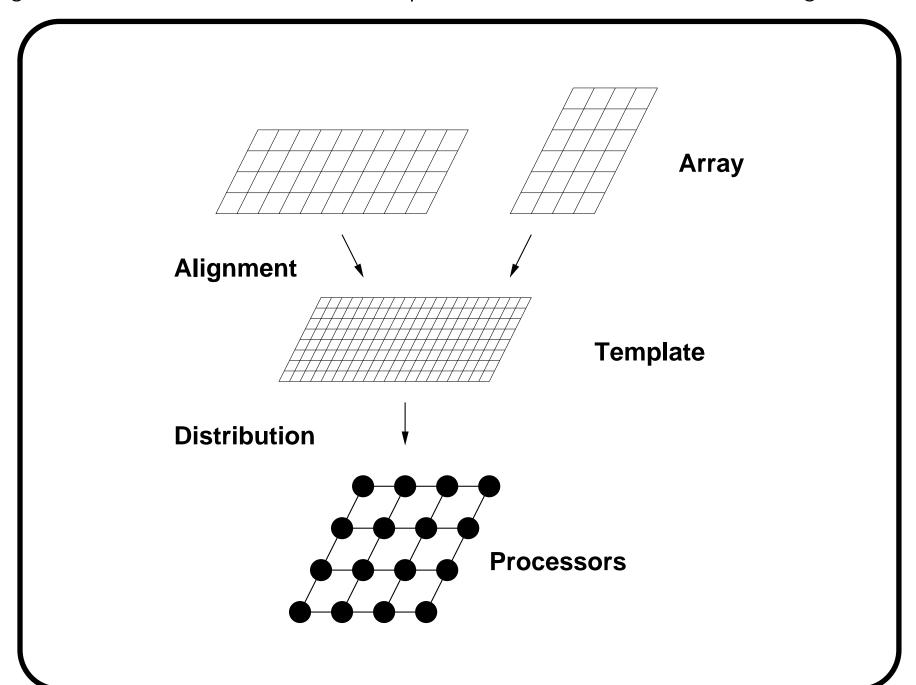
- Arrays are aligned with Templates
- Templates are distributed on Processors
- Processors may be mapped on physical processors

Only the 2 first levels are specified and used

4 directives:

ALIGN TEMPLATE DISTRIBUTE PROCESSORS

HPF Mapping Model



2 Level Mapping

load balancing and reduced communications

- Affine Alignment: ALIGN
 - relative data mapping'this array must be mapped as this one'
 - close to the application/problem decide data locality
- Regular distribution: DISTRIBUTE
 - distribute data on the processors, BLOCK or CYCLIC
 - close to the machine/architecture

Templates

TEMPLATE is a declaration directive

- array of nothing
- not allocated!
- used as an alignment target
- an array can be used as a template

```
!hpf$ template t(100), t2(n,m)
!hpf$ dimension(n,n), template :: &
!hpf$ domain
```

Alignment

ALIGN is a declaration directive

- arrays are aligned with templates
- 2 array elements aligned with the same template element will be on the same processor
- data locality

```
!hpf$ ALIGN A(I,...) WITH T(affine(I),...)
```

Constraints

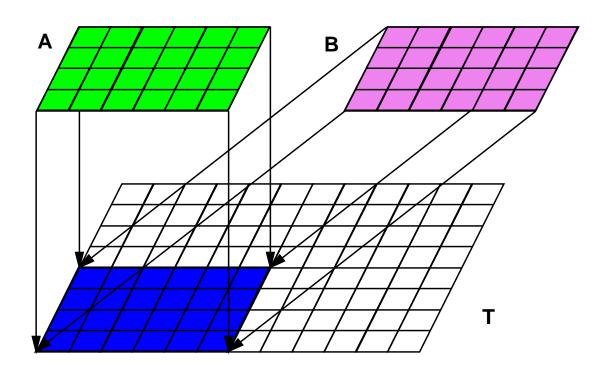
- Affine function
 - one dummy variable a*I+b
 - a and b are integers
- one use of dummy variables

NO: ALIGN A(I) WITH T(I,I)...

Alignment expressivity

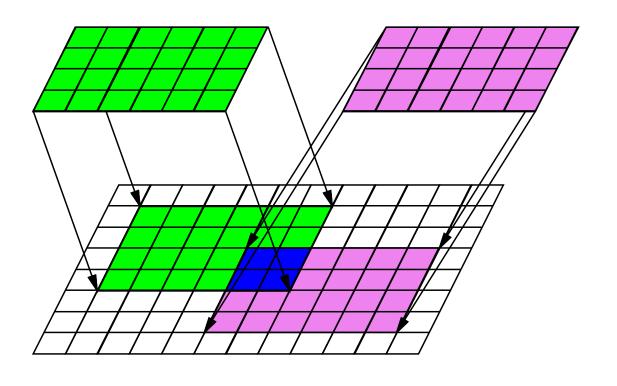
- direct mapping
- transposition
- reversing
- replication
- collapsing
- ...

Direct mapping



align A(i,j) with T(i,j)
align B with T

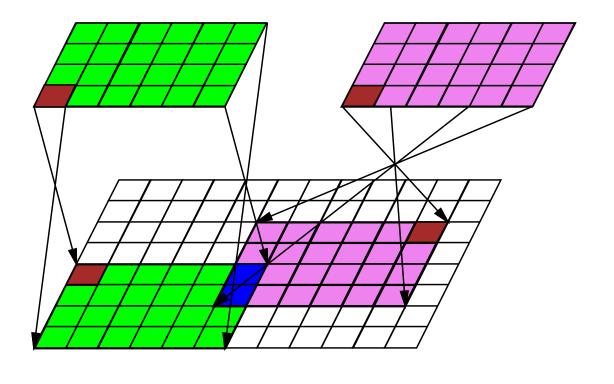




```
align A(i,j) with T(i+1,j+3)
```

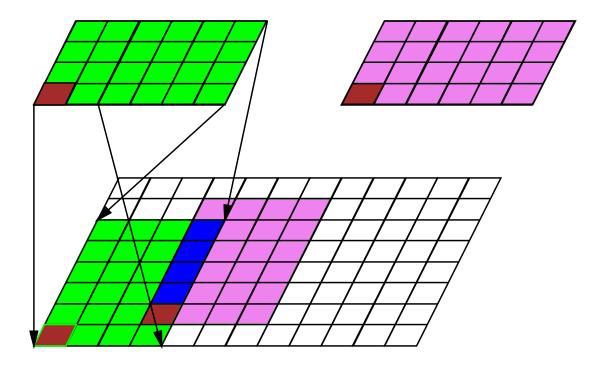
align B(:,:) with T(6:11,2:5)

Reversing



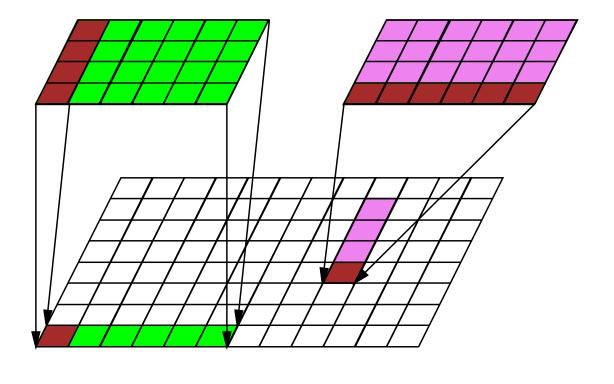
```
align A(i,j) with T(i,5-j) align B(:,:) with T(11:6:-1,6:3:-1)
```





```
align A(i,j) with T(j,i)
align B(i,j) with T(j+3,i+1)
```

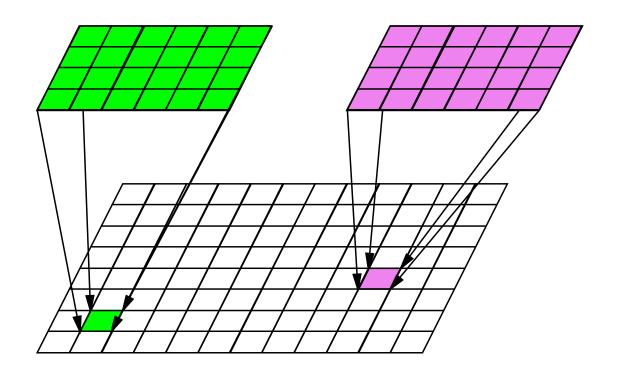




```
align A(i,*) with T(i,1)
```

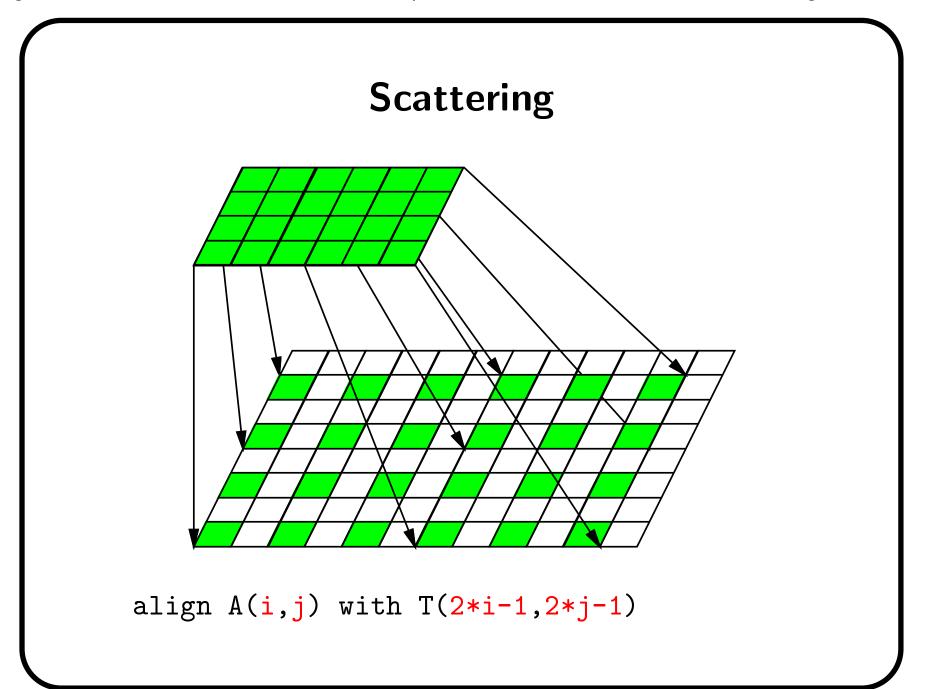
align B(*,:) with T(9,4:7)

More collapsing

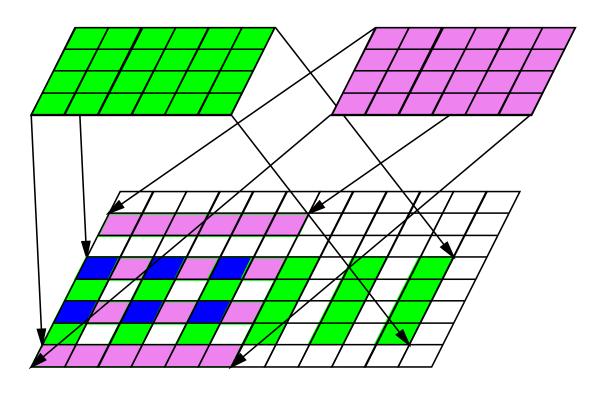


align A(*,*) with T(2,2)

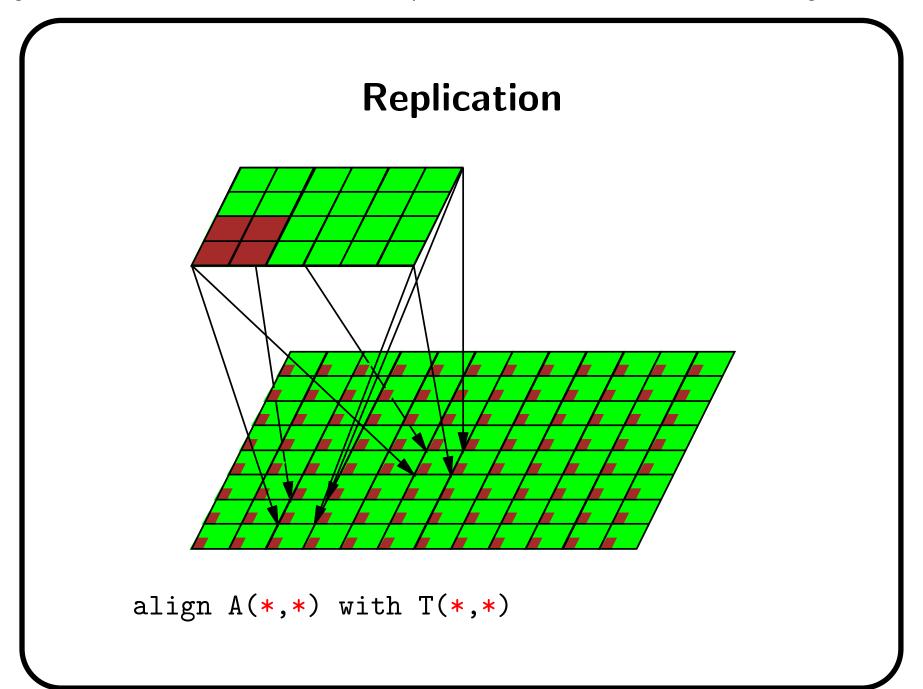
align B(*,*) with T(10,4)

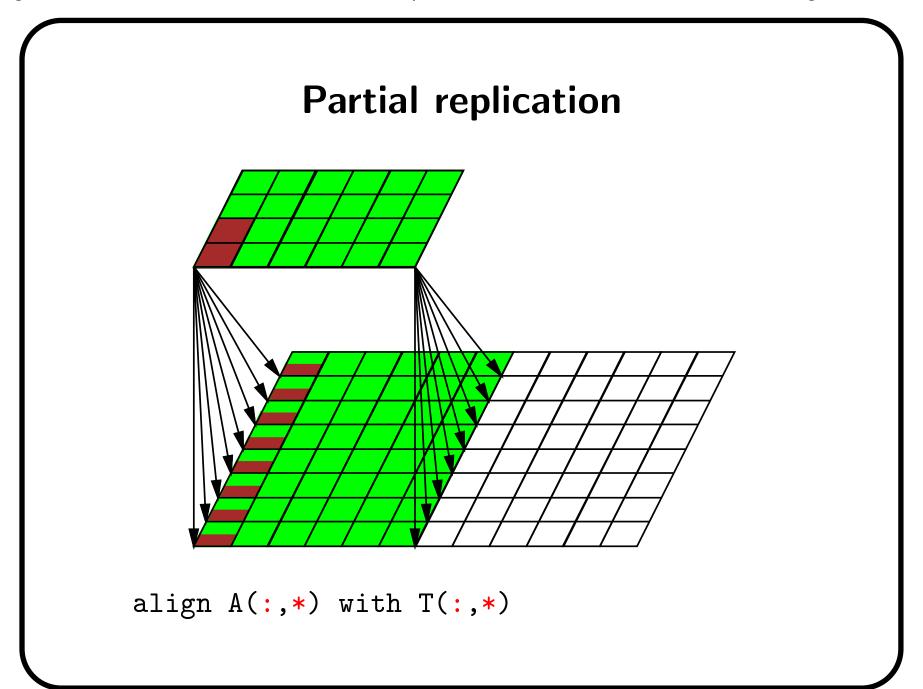


Scattering

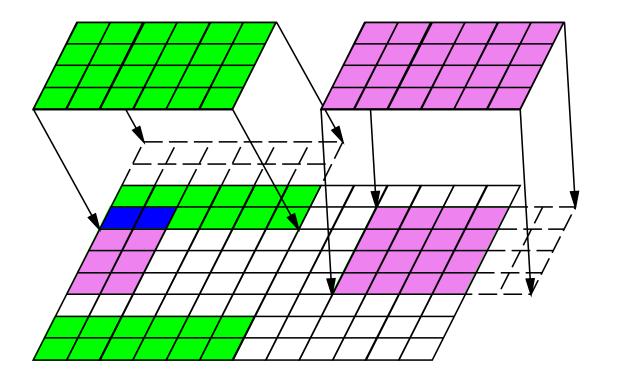


```
align A(i,j) with T(2*i-1,j+1) align B(i,j) with T(i,2*j-1)
```





NO overlaps



```
align A(i,j) with T(i,j+6)! non
```

align B(:,:) with T(9:,4:7) ! non

Processors

PROCESSORS is a declaration directive

- declares abstract processor arrays
- often correspond to actual processors
- can query the number of processors:

```
NUMBER_OF_PROCESSORS()
```

```
!hpf$ PROCESSORS P(10), Q(4,4)
```

```
!hpf$ PROCESSORS Z(2:3,8)
```

!hpf\$ PROCESSORS R(NUMBER_OF_PROCESSORS())

Distribution

DISTRIBUTE is a declaration directive

- templates are distributed on processors
- template elements are grouped together on a processor element

```
!hpf$ DISTRIBUTE T(BLOCK,*) ONTO P
```

!hpf\$ DISTRIBUTE T2(CYCLIC,CYCLIC(10)) ONTO Q

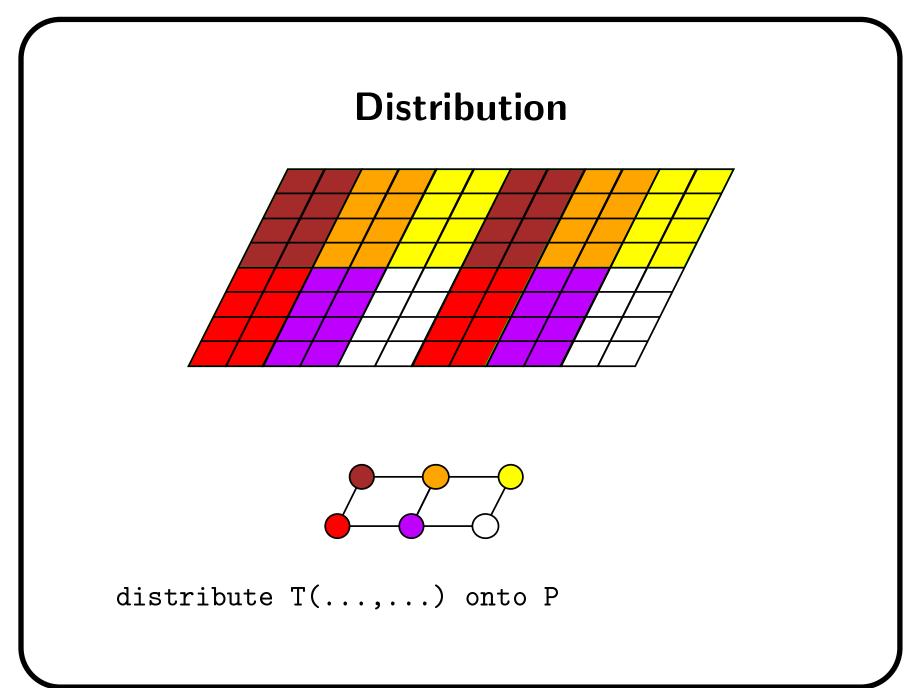
Use of BLOCK and CYCLIC

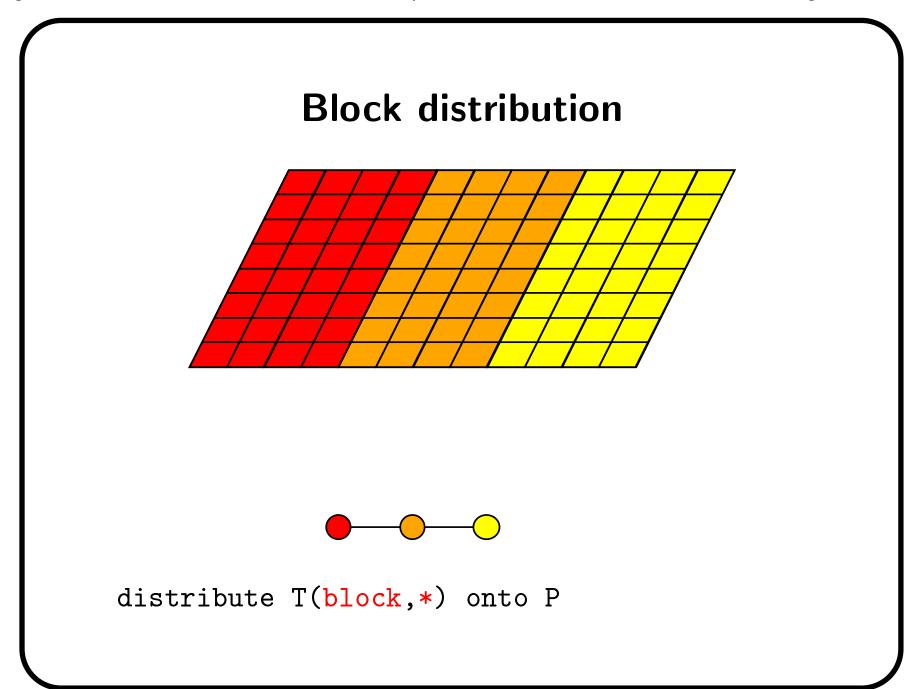
BLOCK regular distribution for neighbor comms typical of stencil computations

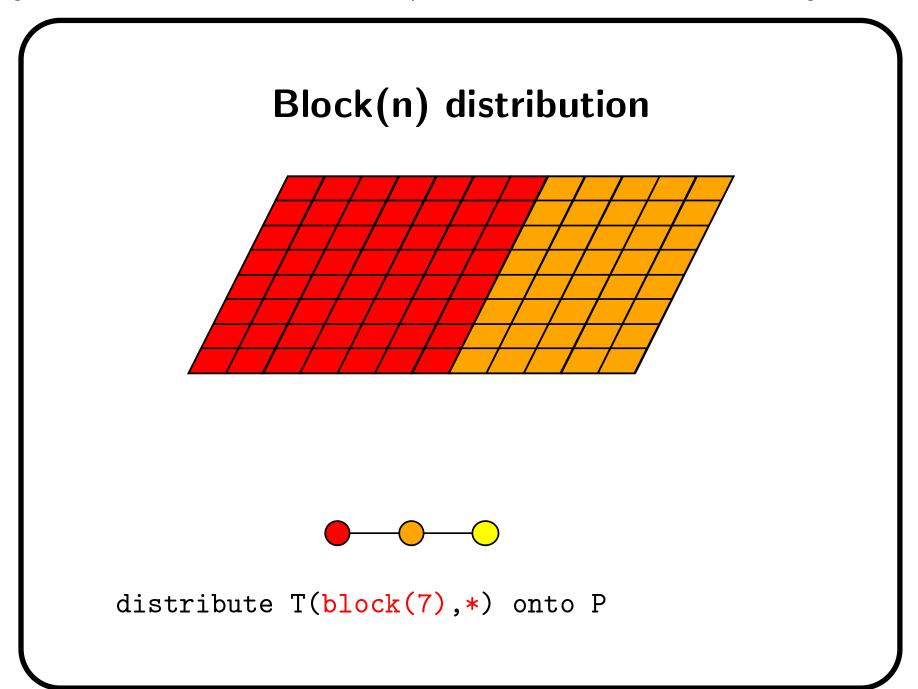
BLOCK(n) adjust block size

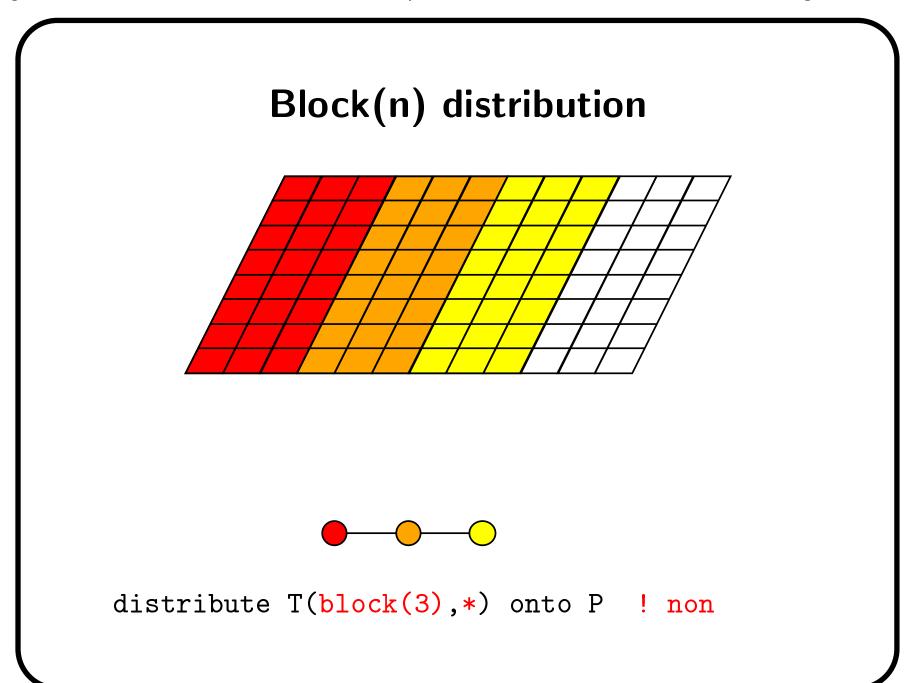
CYCLIC distribute load

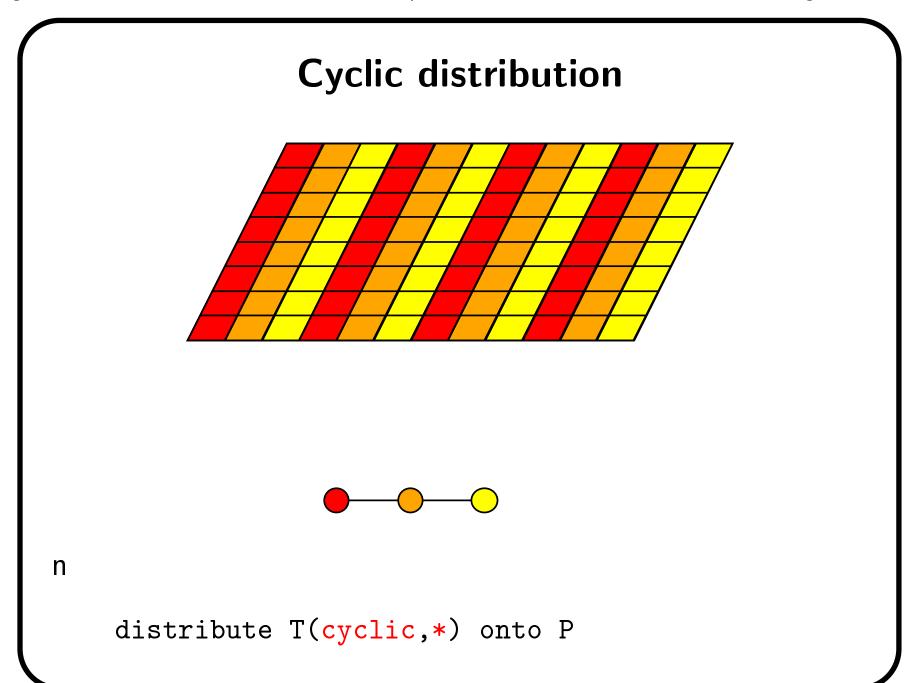
CYCLIC(n) both load balance and neightbor comms.

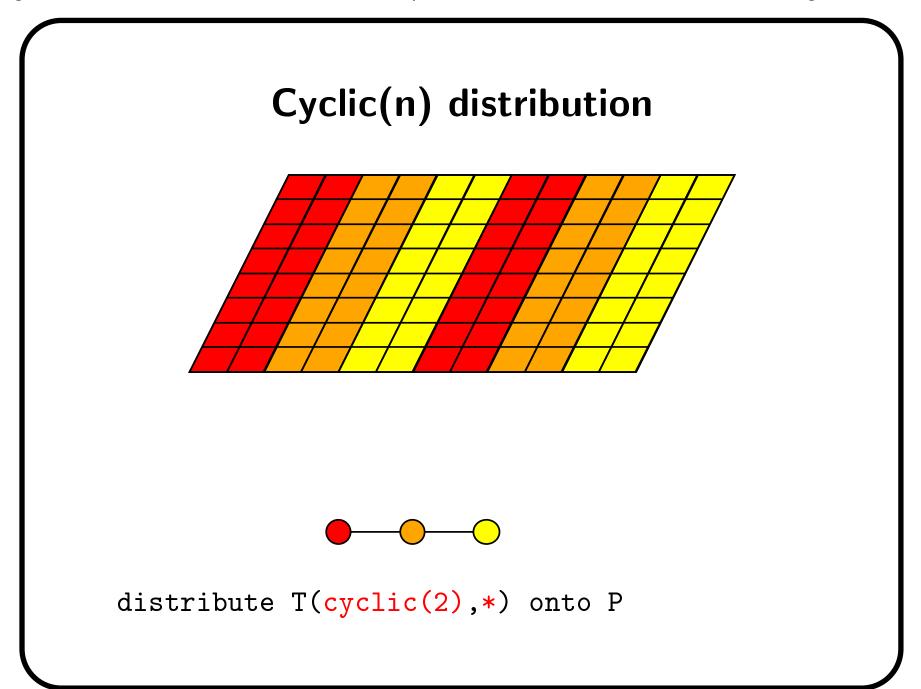




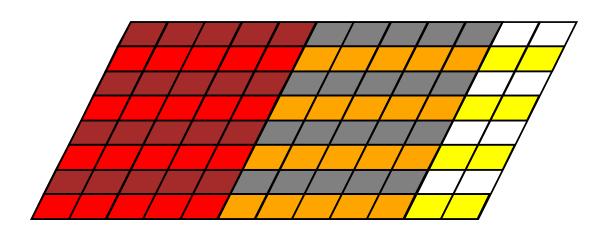


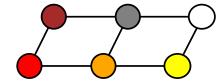






Multi dimensional distributions





distribute T(block(5), cyclic(1)) onto P

The template debate...

Heavy discussions on the mailing list

- Hans Zima, Vienna
 - Vienna Fortran does not include templates
 - templates do not improve expressivity
 - one more non user friendly concept
- Ken Kennedy, Rice
 - Fortran D includes similar decompositions
 - they can be larger than any array
 - easier port, clearer for the user

Other mapping directives

REALIGN executable, dynamic alignment

REDISTRIBUTE executable, dynamic distribution

DYNAMIC attribute to an object which can be RE...

Moved as approved extensions...

More about HPF mapping directives

- Fortran 90 syntactic sugar
- use of intrinsics in directives
- default mapping
- some strange mappings
- HPF mapping limits

Syntactic sugar

```
!hpf$ processors, dimension(10) :: PROC
!hpf$ template, &
!hpf$ distribute(block) onto PROC, &
!hpf$ dimension(100), &
!hpf$ dynamic :: T1,T2
!hpf$ dynamic, align (:) with T1(:) :: A, B
```

Use of intrinsics

- NUMBER_OF_PROCESSORS()
 number of physical processors available
- PROCESSORS_SHAPE()
 processor number of dimension

```
!hpf$ PROCESSORS P(NUMBER_OF_PROCESSORS()), &
!hpf$ AP(NUMBER_OF_PROCESSORS()/10,10)
```

Default mappings

- direct alignement
- 1D vector of physical processors
- block distribution

```
!hpf$ align A with B
```

!hpf\$ distribute B

HPF mapping limits

No arbitrary irregular mapping

- ! From Vienna Fortran
 ALIGN A(I) WITH B(f(I))
 - f is an arbitrary function
 - compiler optimizations?

Mapping at subprogram interfaces

prescriptive mapping to inforce (with a remapping)

descriptive specify an expected mapping

transcriptive unspecified mapping

Prescriptive mapping

• just like a static mapping

```
subroutine CALCUL(A)
real A(100,100)
!hpf$ distribute A(block,block)
...
```

- remapping on entry and exit (?)
- runtime cost

Descriptive mapping

- similar to prescriptive
- assertion to the compiler
- favor optimisations
- all callers must conform, otherwise errors!

```
!hpf$ align A with * T
!hpf$ distribute * T(block)
```

Transcriptive mapping

- no information for optimisations?
- template inheritance concept
- syntaxes

```
subroutine CALCUL(A)
real A(100,100)
!hpf$ distribute A *
!hpf$ inherit A
```

Conclusion to HPF mapping

- close to Fortran D
- replication
- suits arrays
- should favor optimizations
- directives = advice!

So what?

- array distributed on processors
- the data are shared between them...
- we have now to share the computation
- without changing the semantics
- let's look for parallelism!
 - array expressions
 - parallel loops
 - dataparallel loops
 - and sequential loops...

3 loops in HPF

- Group of Charles Koelbel, Rice University
- sequential do loop: DO
- parallel loop: DO + INDEPENDENT
 iterations to not dependent of the other
- dataparallel loop: FORALL

3 loops in HPF

```
do i=1, 10

a(i) = a(i-1) + a(i) + a(i+1)

enddo
```

```
!hpf$ independent
    do i=1, 10
        a(i) = 1.0
    enddo
```

```
forall (i=1:10)

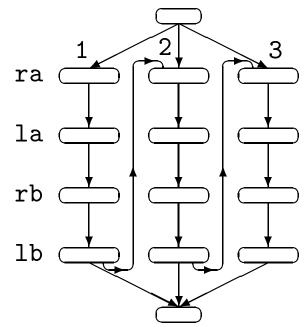
a(i) = a(i-1) + a(i) + a(i+1)

end forall
```

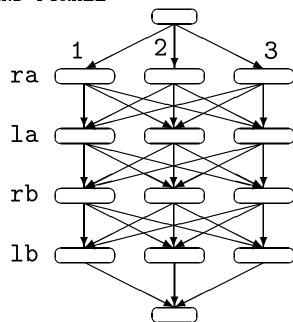
DO et FORALL

D0 i = 1, 3
 lhsa(i) = rhsa(i)
 lhsb(i) = rhsb(i)

END DO



FORALL (i = 1:3)
 lhsa(i) = rhsa(i)
 lhsb(i) = rhsb(i)
END FORALL



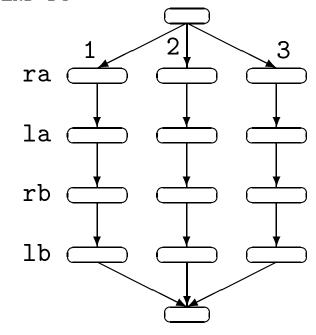
INDEPENDENT

!HPF\$ INDEPENDENT

D0 i = 1,
$$3$$

$$lhsb(i) = rhsb(i)$$

END DO

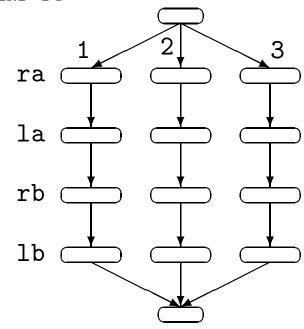


!HPF\$ INDEPENDENT

$$FORALL(i = 1:3)$$

$$lhsb(i) = rhsb(i)$$

END DO



FORALL: 3 dataparallel loops

• implicit with array sections/expressions

```
A = B + C ! matrices

D(1:20) = D(2:21) ! shift
```

generalization: FORALL statement index explicitely named

FORALL (i=10:n, j=10:m)
$$a(i,j) = 0.0$$

• extension: FORALL construct

FORALL statement

- generalize array sections
- dataparallel semantics (SIMD)
- only one assignment
- optional mask
- can call PURE functions

FORALL(i=1:n, ..., mask) &
 assignment

```
! transpose
forall(i=1:n, j=1:n) x(i,j) = y(j,i)
forall(i=1:n, y(i).ne.0.0) &
  x(i) = 1.0 / y(i)
! relaxation
forall(i=2:n-1) &
  x(i) = 0.5 * x(i) + 0.25 * (x(i-1)+x(i+1))
forall(i=1:n, y(i).ge.0.0) &
  x(i) = SQRT(y(i))
```

FORALL construct

- extension of FORALL
- includes
 - assignements
 - WHEREs or FORALLs

```
FORALL(i=1:n,...,mask)
WHERE or FORALL or assigns...
ENDFORALL
```

```
FORALL(i=1:n)
 A(i,i) = SQRT(A(i,i))
 B(i) = LOG(B(i,i)) + i
ENDFORALL
FORALL(i=1:n, j=1:n)
 WHERE (b(i,j)>0.0)
    a(i,j) = 1.0
 ELSEWHERE
    a(i,j) = -1.0
 ENDWHERE
ENDFORALL
```

PURE attribute to functions

- user functions which behaves like intrinsics
 - no I/Os, no STOP, no PAUSE...
 - no access to commons or save data...
- must be declared as such
 - in the function interface
 - in the function itself

```
pure function f(x)
  real :: f
  real, intent(in):: x
  f = (4.3*x+2.1)*x+1.7
end function
```

INDEPENDENT directive

- a parallel loop!
- iterations can be executed in any order
- can specify private variables
 variables private to iterations
- HPF 2.0: can specify reductions
- can apply to DO and FORALL

```
!hpf$ INDEPENDENT, NEW(var-list), REDUCTION(...)
DO i=1,n ... / FORALL(i=1:n)
```

Examples of INDEPENDENT

```
!hpf$ INDEPENDENT
     DO i=1, n
       A(P(i)) = B(i) ! P is a permutation
     ENDDO
!hpf$ INDEPENDENT, NEW(x)
     DO i=1, n
       x = A(i)**3
       B(i) = x*x + x
      ENDDO
```

```
! squere matrix multiplication
!hpf$ INDEPENDENT, NEW(j)
     D0 i=1, n
!hpf$ INDEPENDENT, NEW(k,t)
       DO j=1, n
         t = 0.0
!hpf$ INDEPENDENT, REDUCTION(t)
         DO k=1, n
           t = t + A(i,k)*B(k,j)
         ENDDO
         C(i,j) = t
       ENDDO
     ENDDO
```

Efficient parallel loop?

- assertion about parallel semantics
- compiler may not know how to exploit it!
- additionnal constraints helps:
 - block distributions
 - all data accessed in an iteration are aligned
 - access to neighbor processors
 - remote data stored in overlap areas

Parallel library functions

- reductions applied on arrays
 - Fortran 90: SUM PROD ALL MAXVAL...
 - HPF/F95: ALL PARITY IALL MAXLOC...
 - include options (DIM, MASK...)
- scans...
 - SCATTER
 - _PREFIX
 - _SUFFIX
- others: SORT_UP SORT_DOWN GRADE_DOWN GRADE_UP

Reductions

associative/commutative operations parallelism induce numerical issues...

- Fortran 90
 - arithmetic: SUM (+) PROD (*)
 - logical: ALL (.and.) ANY (.or.) COUNT (.true.)
 - order: MAXVAL (max) MINVAL (min)
 - location: MAXLOC MINLOC
- HPF
 - logical: PARITY (.xor.)
 - bitwise: IALL IANY IPARITY

SCATTER functions

- reordering in Fortran 90
- based on permutations
- extended with _SCATTER
- how to combine indexed elements...

```
! A = [ 10 20 30 ]
! V = [ 1 2 3 4 ]
! IDX = [ 3 2 1 ] ! IDX is a permutation
        V[IDX] = A
! V = [ 30 20 10 4 ]
```

SCATTER functions...

- base array to be modified
- contributing array
- integer index arrays
 - as many as DIM(base)
 - conformant to contributing array
- optionnal mask (logical array)

```
! A = [ 10. 20. 30. 40. ]
! IDX = [ 1 2 1 2 ]
! BASE = [ 1. 2. 3. ]

X = SUM_SCATTER(A, BASE, IDX)
! X = [ 41. 62. 3. ]
```

PREFIX and SUFFIX functions

- accumulation of elements
- forward or backward
- optional MASK, SEGMENT, DIM, EXCLUSIVE

```
! A = [ 1. 2. 3. 4. ]

S = SUM_PREFIX(A)

! S = [ 1. 3. 6. 10. ]

P = PROD_SUFFIX(A)

! P = [ 24. 24. 12. 4. ]
```

PREFIX and DIM option

- DIM: operation performed on it
- else array is linearized...

```
! A = [ 1 2 3 ]
! [ 4 5 6 ]
S = SUM_PREFIX(A, DIM=2) ! S = [ 1 3 6 ]
! [ 4 9 15 ]
P = PROD_SUFFIX(A, DIM=1) ! P = [ 1 10 18 ]
! [ 4 5 6 ]
T = SUM_PREFIX(A) ! T = [ 1 7 15 ]
! [ 5 12 21 ]
```

PREFIX options

- MASK: element contributes if true
- EXCLUSIVE: current element does not contribute
- SEGMENT: grouping on linearized array

options and array expressions

 MASK, SEGMENT, indexes can be the result of array op.

Rationale for scans

- SCATTER, PREFIX, SUFFIX
- quite specialized operations
- usefull to some applications
- parallel version can be developped
- but not in HPF!
- thus made available in the library

Sort functions

efficient parallel implementations

- compute permutations
 - GRADE_UP
 - GRADE_DOWN
- performs the sort
 - SORT_UP
 - SORT_DOWN

GRADE example

Miscelaneous scalar functions

- ILEN: number of bits needed
- LEADZ: number leading zero bits
- POPCNT: number of bits to 1
- POPPAR: 1 if POPCNT is even, 0 if odd

HPF inquiry functions

- a set of functions to query about array mappings
- non portable wrt Fortran 90
- examples :
 - HPF_ALIGNMENT
 - HPF_DISTRIBUTION
 - HPF_TEMPLATE

Conclusion about the library

- not portable wrt Fortran 90
- performance depends on implementation
- use them if appropriate!

Extrinsic procedures

- ability to call routine in
 - another language (C, F77, Fortran, HPF-lite)
 - another parallel paradigm (MPI, ...)
- routine called with different conventions:
 - on one pe with all data (serial)
 - on pes... (global, standard HPF!)
 - on every pe with local data (local)
 - **–** . . .

Syntax for extrinsics

- explicit interface in caller
- special declaration in the callee
- with fortran 90 keyword: EXTRINSIC

```
INTERFACE
  EXTRINSIC(HPF_LOCAL) SUBROUTINE FOO(A)
    REAL, DIMENSION(:), INTENT(INOUT):: A
  END SUBROUTINE
END INTERFACE
```

```
!hpf$ distribute X(BLOCK)
CALL FOO(X)
```

function definition

```
EXTRINSIC(HPF_LOCAL) SUBROUTINE FOO(A)
   REAL, DIMENSION(:):: A
   INTEGER I
! query about A LOCAL declaration...
   DO I=LBOUND(A), UBOUND(A)
   ...
   ENDDO
END SUBROUTINE
```

F77 serial model

- call a serial routine in Fortran 77
- for instance, a special I/O routine for X display
- available in ADAPTOR

HPF serial model

- call a serial routine in restricted HPF
- needed data collected before the call
- use Fortran 90 intent declarations!

HPF local model

- routine coded in restricted HPF!
- no data distribution!
- can use message passing such as MPI
- standard inquiry about local data
- inquery about data mapping (HPF 2.0 Ap. Ext.)

```
MY_PROCESSOR()
```

GLOBAL_TO_LOCAL()

HPF 2.0 approved extensions

Many proposal which are not incorporated

- data mapping
- data and task parallelism
- asynchroneous I/Os
- library extensions
- more extrinsics. . .

Approved extensions for data mapping

- REALIGN, REDISTRIBUTE, DYNAMIC...
- mappings on processor subsets
 DISTRIBUTE T(BLOCK) ONTO P(2:4)
- RANGE: list of mappings a subroutine might accept...
- SHADOW: overlap width declaration
 DISTRIBUTE(BLOCK) SHADOW(1:2):: A,
- •

Approved extensions for parallelism

ON directive to locate computations

- TASK_REGION + ON directive
- active processors
- . .

Port of ONDE24 to HPF

- 2D acoustic wave propagation
- from IFP: French Oil Institute
- 1000 lines of Fortran 77
- target: IBM SP2
- stencil-based application!
- direct resolution over time
- order 2 in time, order 4 in space
 Should be a dream!

ONDE24 initial performance

- HPF = High Performance!
- code profile: gprof/tprof...
- can use RS6000 performance counters
- result: 40 Mflop/s
- processor peak performance is 250 Mflop/s
- what is the trouble?

Improving the kernel loop

- cache behavior is ok.
- expression structure:
 - processor needs 4 independent fma: a*b+c
 - more fma can be extracted
 - balanced tree ⇒ more parallelism
- result: 80 Mflop/s!
- with unrolling and scheduling: 110 Mflop/s
 but will not suit HPF...

ONDE24 declaration cleaning

- larger declaration than used...
- impact on distribution and load balance:
 - few processors used if BLOCK
 - or more communications if CYCLIC(n)
- U and V are the big arrays
- dynamic allocation? not F77...

Initial declarations

```
NPMAX, NTMAX, NBTRAC
INTEGER
          (NPMAX
                     = 809)
PARAMETER
PARAMETER (NTMAX
                     = 4000)
PARAMETER (NBTRAC = 50)
                (NPMAX+1, NPMAX, 2)
REAL*8
                (NPMAX+1, NPMAX)
REAL*8
          V
                (NPMAX+1,3,3)
REAL*8
          В
REAL*8
          UB
                (NPMAX+1,3)
REAL*8
          UINT
                (NPMAX)
          SISMO (NBTRAC, NTMAX)
REAL*8
          NP, NT
INTEGER
```

Array declarations

- static allocation
- bounds moved as parameters
- runtime check to match problem
- actual bounds needed in directives
- hence MUST be compile-time constants

Simplified declarations

```
INTEGER NP, NT, NBTRAC
```

PARAMETER (NP = 800)

PARAMETER (NT = 1000)

PARAMETER (NBTRAC = 50)

REAL*8 U(NP,NP,2)

REAL*8 V(NP,NP)

REAL*8 B(NP,3,3)

REAL*8 UB(NP,3)

REAL*8 UINT(NP)

REAL*8 SISMO(NBTRAC,NT)

ONDE24 Data Mapping

- 1D, 2D and 3D arrays are used
- align data used together...
- must examine actual uses!
- then only ONE distribution

!hpf\$ TEMPLATE DOMAIN(NP,NP)

U and V arrays

- scan all uses
- U(I,J,*) and neighbors used with V(I,J)
- suggest direct alignement!

```
!hpf$ ALIGN U(I,J,*) WITH DOMAIN(I,J)
!hpf$ ALIGN V(I,J) WITH DOMAIN(I,J)
!
    or
! or
!hpf$ ALIGN U, V WITH DOMAIN
```

UINT vector first use

```
DO 92 I = 3, NP-2
         UINT(I) = B(I,1,1)*(U(I,NP,KP)+U(I,NP-1,KP))
         [...]
     * + B(I,1,3)*(U(I+1,NP-1,KP)+U(I+1,NP,KP)+
                      U(I-1,NP-1,KP)+U(I-1,NP,KP))
     &
92
      CONTINUE
      DO 93 I = 3, NP-2
        U(I,NP,KP) = UINT(I)
93
      CONTINUE
!hpf$ ALIGN UINT(I) WITH DOMAIN(I,NP)
```

Second use

```
DO 94 I = 3, NP-2
         UINT(I) = B(I,2,1)*(U(NP,I,KP)+U(NP-1,I,KP))
         [...]
           + B(I,2,3)*(U(NP-1,I+1,KP)+U(NP,I+1,KP) +
                       U(NP-1,I-1,KP)+U(NP,I-1,KP))
     &
94
      CONTINUE
      DO 95 I = 3, NP-2
        U(NP,I,KP) = UINT(I)
95
      CONTINUE
!hpf$ ALIGN UINT(I) WITH DOMAIN(NP,I)
```

Third use

```
DO 96 I = 3, NP-2
         UINT(I) = B(I,3,1)*(U(I,1,KP)+U(I,2,KP))
         [...]
         + B(I,3,3)*(U(I+1,2,KP)+U(I+1,1,KP)+
                      U(I-1,2,KP)+U(I-1,1,KP))
     &
96
      CONTINUE
      DO 97 I = 3, NP-2
        U(I,1,KP) = UINT(I)
97
      CONTINUE
!hpf$ ALIGN UINT(I) WITH DOMAIN(I,1)
```

UINT Mapping?

- not used elsewhere
- 3 different alignements?
- values directly reused!
- REALIGN?
 - usually not implemented
 - no need to move the data!
- this suggest 3 different arrays!
 UINTW UINTS UINTE (West, South, East)

UINT* Mapping

```
REAL*8 UINTW(NP), UINTS(NP), UINTE(NP)
```

```
!hpf$ ALIGN UINTE(I) WITH DOMAIN(I,NP)
```

```
!hpf$ ALIGN UINTS(I) WITH DOMAIN(NP,I)
```

!hpf\$ ALIGN UINTW(I) WITH DOMAIN(I,1)

UB use

```
DO 71 I = 2, NP-1
    UB(I,1) = U(I,NP-1,KP)
    UB(I,2) = U(NP-1,I,KP)
71 UB(I,3) = U(I,2,KP)
   DO 92 I = 3, NP-2
    UINTE(I) = U(I,NP,KP) .. U(I,NP-1,KP) .. UB(I,1)
!hpf$ ALIGN UB(I,1) WITH DOMAIN(I,NP-1) ???
!hpf$ ALIGN UB(I,2) WITH DOMAIN(NP-1,I) ???
!hpf$ ALIGN UB(I,3) WITH DOMAIN(I,2)
                                        ???
```

UB Mapping

- dimension (NP,3)
- temporary store for borders in time loop
- cannot map dimensions separatly in HPF
- split UB: UBW, UBS, UBE

UB* Mapping

```
REAL*8 UBE(NP), UBS(NP), UBW(NP)
!hpf$ ALIGN UBE(I) WITH DOMAIN(I,NP-1)
!hpf$ ALIGN UBS(I) WITH DOMAIN(NP-1,I)
!hpf$ ALIGN UBW(I) WITH DOMAIN(I,2)
     DO I = 2, NP-1
       UBE(I) = U(I,NP-1,KP)
       UBS(I) = U(NP-1,I,KP)
       UBW(I) = U(I,2,KP)
     ENDDO
```

B Mapping? first use

```
REAL*8 B(NP,3,3)
     DO 41 I = 2, NP-2
         V(I,NP) = 2/((1/V(I,NP))+(1/V(I,NP-1)))
                 = V(I,NP)*DELTAT
         W
         B(I,1,1) = 1 - (1/(W+H)) * (2*H + W*W/H)
         B(I,1,2) = 2*H/(W+H)
         B(I,1,3) = W*W/(2*H*(W+H))
41
!hpf$ ALIGN B(I,1,*) WITH DOMAIN(I,NP)
```

2nd and 3rd uses

```
DO 42 I = 3, NP-2
        V(NP,I) = 2/((1/V(NP,I))+(1/V(NP-1,I)))
             = V(NP,I)*DELTAT
        W
        B(I,2,1) = 1 - (1/(W+H)) * (2*H + W*W/H)
        B(I,2,2) = 2*H/(W+H)
42
        B(I,2,3) = W*W/(2*H*(W+H))
    DO 43 I = 2, NP-2
        V(I,1) = 2/((1/V(I,1))+(1/V(I,2)))
          = V(I,1)*DELTAT
        W
        B(I,3,1) = 1 - (1/(W+H)) * (2*H + W*W/H)
        B(I,3,2) = 2*H/(W+H)
43
        B(I,3,3) = W*W/(2*H*(W+H))
```

B Mapping

- similar to UB!
- split in UBW, UBS, UBW!

```
REAL*8 BE(NP,3), BS(NP,3), BW(NP,3)
```

```
!hpf$ ALIGN BE(I,*) WITH DOMAIN(I,NP)
```

!hpf\$ ALIGN BS(I,*) WITH DOMAIN(NP,I)

!hpf\$ ALIGN BW(I,*) WITH DOMAIN(I,1)

SISMO Mapping

- stores received signal
- sismograph placed on the field

```
REAL*8 SISMO(NBTRAC,NT)
```

```
DO 81 I = 1,NBTRAC

81 SISMO(I,N) = U(IS, NP - NBTRAC + I, KP)
```

!hpf\$ ALIGN SISMO(I,*) WITH DOMAIN(IS,(NP-NBTRAC)+I)

SISMO Mapping

- array shift is legal!
- but IS is a runtime constant...
- depth of source of in field

!hpf\$ ALIGN SISMO(I,*) WITH DOMAIN(1,(NP-NBTRAC)+I)

1D or 2D distribution on 4 processors?

(BLOCK,*)

- shorter inner loop
- 4*3*n communications per iteration

(*,BLOCK)

- longer inner loop
- 4*3*n communications per iteration

(BLOCK, BLOCK)

- smaller local array
- 4*2*n communications per iteration
- 50% less communications!

ONDE24 HPF Mapping

- alignment and distribution done!
- now let's find some parallelism!
- must scan all computations and loops...
- parallel if independent of execution order
- Berstein's conditions:
 - no data conflicts (R/W, W/R, W/W)
 - between DISTINCTS iterations
- requires dependence analyses

ONDE24 Kernel computation

Potential conflicts

- Write of U(I,J,KP)
- Reads of U array elements:
 - U(I,J,KP): no, same iteration
 - U(I,J,KM): idem
 - U(I[+-][12],J,KM): no if KM and KP differs
 - U(I,J[+-][12],KM): no if KM and KP differs

KM and KP flip-flop

INTEGER KM, KP

KM = 1

KP = 2

DO 70 N = 2,NT

. . .

KM = KP

KP = 3 - KP

70 CONTINUE

Parallel kernel

```
!hpf$ INDEPENDENT, NEW(I)
     DO J = 3, NP-2
!hpf$
      INDEPENDENT
      DO I = 3, NP-2
       U(I,J,KP) =
       (2.*U(I,J,KM)-U(I,J,KP))-V(I,J)*
    $
    $
       (60.*U(I,J,KM)+(((U(I+2,J,KM) + U(I-2,J,KM))
    $
                       + (U(I,J-2,KM) + U(I,J+2,KM)))
                -16. * ((U(I+1,J,KM) + U(I-1,J,KM)))
                       + (U(I,J-1,KM) + U(I,J+1,KM)))
       ENDDO
     ENDDO
```

Parallel loops in ONDE24

- 2 parallel loops found!
- time loop is not parallel (obvious)
- still 33 other loops to consider! really!
- initializations...
- other computations...
- I/O and statistics...

Obviously parallel loops

DO 51 I = 3,NP-1

$$V(I,2) = (V(I,2)*DELTAT/H)**2$$

Other loops (1)

many such loop nests:

```
DO 140 I = 3,NP-1

U(I,2,KP) = 2*U(I,2,KM)-U(I,2,KP)-V(I,2) *

& (4*U(I,2,KM) - (U(I+1,2,KM)+U(I-1,2,KM)

& + U(I,1,KM)+U(I,3,KM)))

140 CONTINUE
```

Parallel version

border stencils similar to kernel stencil

```
!hpf$ INDEPENDENT

DO 140 I = 3,NP-1

U(I,2,KP) = 2*U(I,2,KM)-U(I,2,KP)-V(I,2) *

& (4*U(I,2,KM) - (U(I+1,2,KM)+U(I-1,2,KM)

& + U(I,1,KM)+U(I,3,KM)))

140 CONTINUE
```

Other loops (2)

3 such loops:

41

```
D0 41 I = 2,NP-2
  V(I,NP) = 2/((1/V(I,NP))+(1/V(I,NP-1)))
W = V(I,NP)*DELTAT
BE(I,1) = 1 - (1/(W+H)) * (2*H + W*W/H)
BE(I,2) = 2*H/(W+H)
BE(I,3) = W*W/(2*H*(W+H))
CONTINUE
```

Parallel version

- arrays are aligned
- W is private (defined and not used AFTER)

```
!hpf$ INDEPENDENT, NEW(W)
DO 41 I = 2,NP-2
     V(I,NP) = 2/((1/V(I,NP))+(1/V(I,NP-1)))
W = V(I,NP)*DELTAT
BE(I,1) = 1 - (1/(W+H)) * (2*H + W*W/H)
BE(I,2) = 2*H/(W+H)
BE(I,3) = W*W/(2*H*(W+H))
41 CONTINUE
```

Other loops (3)

- fully parallel... but data not aligned!
- hence it is serialized!
- can be distributed

```
!hpf$ INDEPENDENT
DO 71 I = 2,NP-1
UBE(I) = U(I,NP-1,KP)
UBS(I) = U(NP-1,I,KP)
UBW(I) = U(I,2,KP)
71 CONTINUE
```

Efficient parallel version

```
!hpf$ INDEPENDENT
     DO 71 I = 2, NP-1
        UBE(I) = U(I,NP-1,KP)
71
!hpf$ INDEPENDENT
     DO 72 I = 2, NP-1
        UBS(I) = U(NP-1,I,KP)
72
!hpf$
     INDEPENDENT
     DO 73 I = 2, NP-1
73
        UBW(I) = U(I,2,KP)
```

I/O Loops

not parallel

22

21

```
DO 21 I = 1,NP

DO 22 J = 1,NP

READ (3,*) V(I,J)

CONTINUE

CONTINUE
```

Reduction loops

```
VMAX = V(1,1)
VMIN = V(1,1)
DO 112 I = 1,NP
 DO 122 J = 1,NP
    IF (V(I,J).LT.VMIN) THEN
      VMIN = V(I,J)
    ENDIF
    IF (V(I,J).GT.VMAX) THEN
      VMAX = V(I,J)
    ENDIF
  CONTINUE
CONTINUE
```

122

112

Reduction loops

- REDUCTION directive?
 - operation must appear: MAX/MIN
 - not always available in compilers (HPF 2.0)
- Fortran 90 array intrinsics! MAXVAL/MINVAL

Any other array accesses?

- many individual array accesses
 - signal source
 - initializations and computations in corners
- may lead to many small communications!
- may imply private scalars...
- hints:
 - ON HOME HPF 2.0 directive
 - 1 iteration parallel loop?

Signal source

```
W = F * (N * DELTAT - 1.0/F)
IF (W**2.LT.TMAX) THEN
W = (W*PI)**2
GO = (1.0 - 2.0*W) * EXP (-W)
U(IS,JS,KM) = U(IS,JS,KM) + GO * 12*V(IS,JS) * H**2
ENDIF
```

Efficient version with HPF 2.0

Efficient version with HPF 1.1?

```
!hpf$ INDEPENDENT, NEW(W,GO)
DO I=1,1
W = F * (N * DELTAT - 1.0/F)
IF (W**2.LT.TMAX) THEN
W = (W*PI)**2
GO = (1.0 - 2.0*W) * EXP (-W)
U(IS,JS,KM) = U(IS,JS,KM) + GO * 12*V(IS,JS) * H**2
ENDIF
ENDDO
```

Corner computation

Summary

- all arrays are mapped
- all parallel loops are marked as such
- is that sufficient?
- no, I/O: sequentialized with elementary messages...

```
DO 21 I = 1,NP

DO 21 J = 1,NP

READ (3,*) V(I,J)
```

21 CONTINUE

Efficient I/O Loops

Performance

- with 3 compilers:
 - HPFC (my prototype!)
 - ADAPTOR (Thomas Brandes, GMD)
 - xlhpf (IBM)
- with three different versions...
- 250 Mflop/s on 4 processors

ONDE24 Conclusion

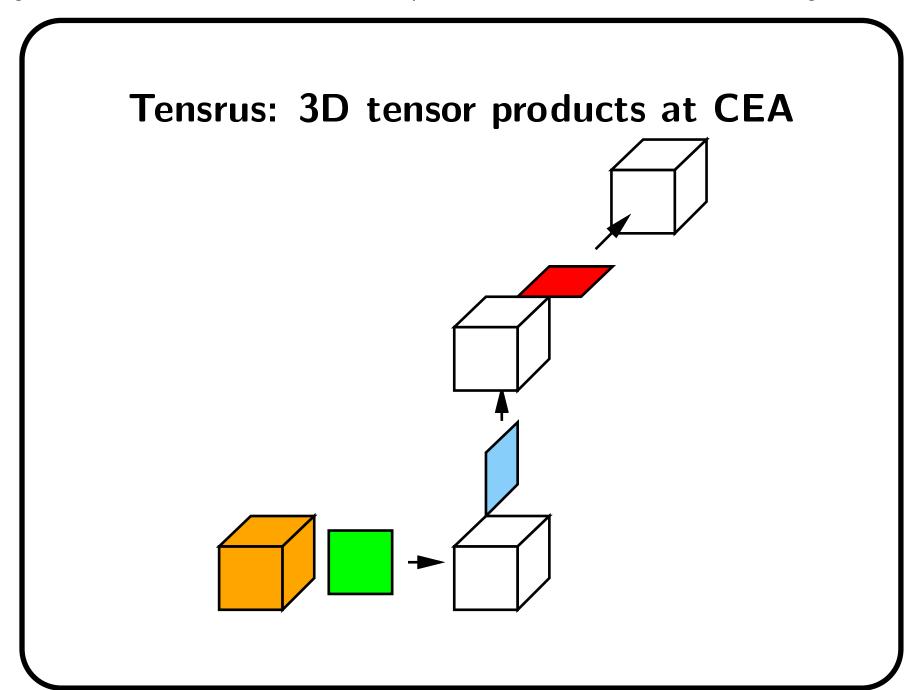
- alignment suits the parallelisation process
 - focus on relative accesses
 - some data need spliting
- HPF parallelism ok for stencils
 - cannot apply on unrolled version
 - limited to 80 Mflop/s/node
- details:
 - array scalar accesses...
 - I/O...

Tensrus

- solver for elliptic differential equations
- for Poisson, Helmoltz equations
- code features:
 - 3D domain, non uniform cart. grid
 - B-spline collocation method of odd order
 - Multipole expansion for boundary conditions
 - cost: N^4 (FFT is $N^3 \ln N$)
- Authors: Jean-Yves Berthou and Laurent Plagne
- CEA: French Nuclear Research Institute

Computation kernel

- basically matrix products!
- large 3D arrays
 - double float type
 - -N*N*N, N=64, 128... 512
 - up to 1GB per matrix!



Initial code

```
SUBROUTINE tensrus(N, MX, MY, MZ, V, T)
REAL, DIMENSION(N,N) :: MX, MY, MZ
REAL, DIMENSION(N,N,N):: T, V
REAL, DIMENSION(N,N,N):: G, F
DO i = 1, N
  DO j = 1, N
    DO k = 1, N
      F(i,j,k) = V(i,k,j)
    ENDDO
  ENDDO
ENDDO
```

```
DO i = 1, N
  DO j = 1, N
    DO c = 1, N
      G(i, j, c) = 0.0
      DO k = 1, N
        G(i,j,c) = G(i,j,c) + MZ(k,c) * F(i,j,k)
      ENDDO
    ENDDO
  ENDDO
ENDDO
! and 2 other products with MY and MZ
! final result stored in T
```

Improved sequential code

- reorder all matrix dimensions before product
- loop order take care of cache
- otherwise basically the same

Better sequential code

```
SUBROUTINE tensrus(N, MX, MY, MZ, V, T)
REAL, DIMENSION(N,N) :: MX, MY, MZ
REAL, DIMENSION(N,N,N):: T, V
REAL, DIMENSION(N,N,N):: G, F
DO i = 1, N
 DO j = 1, N
   DO k = 1, N
      F(k,j,i) = V(i,j,k)
    ENDDO
  ENDDO
ENDDO
```

```
DO i = 1, N
  DO j = 1, N
    DO k = 1, N
      G(k,j,i) = 0.0
      DO c = 1, N
        G(k,j,i) = G(k,j,i) + MZ(c,k)*F(c,j,i)
      ENDDO
    ENDDO
  ENDDO
ENDDO
! and 2 other products with MY and MZ
! final result stored in T
```

HPF version with remappings (hpfc)

- descriptive mappings + interface in caller
- (BLOCK,BLOCK,BLOCK) distribution (better comms)
- redistributions AND reorderings...
 no commercial HPF compiler implements RE*...
- partial replication!
- no reductions: in local memory only

Parallel code

```
SUBROUTINE tensrus(MX, MY, MZ, V, T)
     REAL, DIMENSION(N,N) :: MX, MY, MZ
     REAL, DIMENSION(N,N,N):: T, V
! prescriptive/descriptive mappings:
!hpf$ TEMPLATE D(N,N,N)
!hpf$ PROCESSORS P(4,4,4)
!hpf$ DISTRIBUTE D(BLOCK, BLOCK, BLOCK) ONTO P
!hpf$ ALIGN MX(*,I) WITH D(I,*,*)
!hpf$ ALIGN MY(*,J) WITH D(*,J,*)
!hpf$ ALIGN MZ(*,K) WITH D(*,*,K)
!hpf$ ALIGN WITH D:: V, T
```

```
REAL, DIMENSION(N,N,N):: G, F
!hpf$ ALIGN F(k,j,i), G(k,j,i) WITH D(i,j,k)
!hpf$ DYNAMIC F, G
! reordering
!hpf$ INDEPENDENT
     DO i = 1, N
!hpf$ INDEPENDENT
       DO j = 1, N
!hpf$ INDEPENDENT
         DO k = 1, N
           F(k,j,i) = V(i,j,k)
         ENDDO
       ENDDO
     ENDDO
```

```
! redistribute with partial replication
!hpf$ REALIGN F(*,j,i) WITH D(i,j,*)
!hpf$ INDEPENDENT
     DO i = 1, N
!hpf$ INDEPENDENT
       DO j = 1, N
!hpf$
         INDEPENDENT
         DO k = 1, N
           G(k,j,i) = 0.0
           DO c = 1, N
             G(k,j,i) = G(k,j,i) + MZ(c,k)*F(c,j,i)
            ENDDO
         ENDDO
       ENDDO
     ENDDO
```

CGS

- Conjugate Gradient with Neuman Preconditionning
- generated by some applications
- large sparse matrices involved...
- does it suits HPF?
- Catherine Gaudart and Jean-Yves Berthou, CEA

Computation kernel

```
! \quad Y = D * X
     SUBROUTINE matvect(n, ndiag, irf, d, x, y)
     INTEGER n, ndiag, irf(n, ndiag)
     REAL*8 d(n, ndiag), x(n), y(n)
     DO 10 i = 1, n
     y(i) = 0.0
10
     D0 20 i = 1, n
20
       DO 20 j = 1, ndiag
          y(i) = y(i)+d(i,j)*x(irf(i,j))
20
     END
```

Comments about matvect

- n = 216,000, ndiag = 7
- optimization:
 - n and ndiag compile time constants!
 - loop unrolling?
 - array transposition?
- parallel on the first large dimension

Improved matvect

```
SUBROUTINE matvect(irf, d, x, y)
parameter (n=216000)
INTEGER irf(7,n)
REAL*8 d(7,n), x(n), y(n)
D0 i = 1, n
   y(i) = d(1,i)*x(irf(1,i))
        + d(2,i)*x(irf(2,i))
        + d(3,i)*x(irf(3,i))
        + d(4,i)*x(irf(4,i))
        + d(5,i)*x(irf(5,i))
        + d(6,i)*x(irf(6,i))
        + d(7,i)*x(irf(7,i))
ENDDO
```

HPF version

- from call site and loops
 - Y and X BLOCK distributed
 - pointer passed, switch to F90 array sections?
- in matvect: prescriptive mappings
- indirect access to X: full replication

HPF version

```
SUBROUTINE matvect(irf, d, x, y)
     parameter (n=216000)
     INTEGER irf(7,n)
     REAL*8 d(7,n), x(n), y(n)
!hpf$ TEMPLATE T(n)
!hpf$ DISTRIBUTE T(BLOCK)
!hpf$ ALIGN WITH T(i):: IRF(*,i), D(*,i), Y(i)
!hpf$ ALIGN WITH T(*):: X(*)
!hpf$ INDEPENDENT
     DO 10 i = 1, n
        y(i) = d(1,i)*x(irf(1,i)) + ...
10
```

Communication analyses

Let us assume p processors

- 216000/p elements per processor from X and Y
- 13*216000/p ops per processor in matvect
- thus at most 13*216000/p elements of X used
- remapping: ((p-1)/p)*216000 data transfered!
- p grows: constant communications, reduced ops!!
- useless elements transfered!!!
- p=4: 702 Kops/p, 162 Ktransfers/p
- p=32: 87 Kops/p, 209 Ktransfers/p

How to improve the situation?

- IRF does not change over matvect calls...
- inspect once to know which data to transfer
- execute the communications many times!
- a lot of reuse: really reduced communications!
- smaller X can be allocated
- implemented by hand in MPI, called from HPF

Conclusion about HPF

- efficiency?
 - for some applications
 - fine tunning useful
- portability?
 - if same compiler used, yes!
 - otherwise, different versions often needed...
- simplicity?
 - mappings require geometric training;-)
 - YES if compared to MPI!

HPF vs MPI

- Jean-Yves Berthou:
 - high degree of technicity in both
 - few hours vs few days
 - few weeks vs few months
- HPF always make sense as a first try
 - very good performance on some codes
 - helps understanding the application

HPF Future?

- processor performance increases a lot
- shared memory architecture bottleneck: 8 processors
- some applications/institution requires large machines
- vendor support? only PGI?

Pointers on the web

- HPF: http://dacnet.rice.edu/Depts/CRPC/HPFF/
- ADAPTOR:

http://www.gmd.de/SCAI/lab/adaptor/

- PGI: http://www.pgroup.com/
- HPFC: http://www.cri.ensmp.fr/~coelho/hpfc
- a tutorial: http://qqq.npac.syr.edu/hpfa/