Cheatsheet - (First-order) Predicate & Propositional Logic

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1. Intro

Predicates describe properties of objects.

For example:

$$\text{odd}(3)$$

$\text{odd}(3)$ means 3 is an odd number. odd is a predicate, 3 is an object. Predicates take arguments and become **propositions**. A proposition is a statement that can be either *true* or *false*. It must be one or the other, and it cannot be both.

Connectives can be applied:

$$\text{odd}(3) \wedge \text{prime}(3)$$

This means that 3 is odd but also prime.

2. Syntax

Propositions are denoted by capital letters, such as $P, Q, \ldots$. General statements are denoted by lowercase letters, such as $p, q, \ldots$

3. Connectives

Logical NOT: $\neg p$ is true if and only if p is false (also called *negation*).

Logical OR: $p \vee q$ is true if and only if at least one of p or q is true or if both p and q are true (also called *disjunction*).

Logical AND: $p \wedge q$ is true if and only if both p and q are true (also called *conjunction*).

Logical IF…THEN: $p \rightarrow q$ is true if and only if either p is false or q is true (also called *conditional* or *implication*). p is the premise, q is the conclusion.

Logical IF and only IF: $p \leftrightarrow q$ is true if and only if both p and q are true (also called *bi-conditional*).

Exclusive OR: XOR: $p \oplus q$ is true if p or q is true but not both.

3.1. Translation to Connectives

As an example, lets consider the propositions:

- $P = \text{I study 20 hours a week}$
- $R = \text{I will pass the exam}$
- $S = \text{I will be happy}$
- $Q = \text{I attend all the lectures}$

And the following connectives:

$$(P \vee Q) \rightarrow (R \wedge S)$$

which is a translation of: "**If** I study 20 hours a week **or** attend all the lectures, **then** I will pass the exam **and** I will be happy."

4. Truth Tables

4.1. Negation: $\neg$

$$true = \neg false$$
$$false = \neg true$$

4.2. Conjunction: $\wedge$

$$true \wedge true = true$$
$$true \wedge false = false$$
$$false \wedge true = false$$
$$false \wedge false = false$$

4.3. Disjunction: $\vee$

$$true \vee true = true$$

$$true \lor false = true$$

$$false \lor true = true$$

$$false \lor false = false$$

4.4. Implication: $\rightarrow$

$$true \rightarrow true = true$$

$$true \rightarrow false = false$$

$$false \rightarrow true = true$$

$$false \rightarrow false = true$$

> **NOTE** This can seem weird at first, this answer helps: https://math.stackexchange.com/a/100288
>
> *If you start out with a false premise, then, as far as implication is concerned, you are free to conclude anything. (This corresponds to the fact that, when p is false, the implication $p \rightarrow q$ is true no matter what q is.)*
>
> *If you start out with a true premise, then the implication should be true only when the conclusion is also true. (This corresponds to the fact that, when p is true, the truth of the implication is the same as the truth of q.)*

Additionally, let p and q be propositions and A the conditional statement:

$$p \rightarrow q$$

then:

- p is called the **hypothesis** (or antecedent or premise) and q is called the **conclusion** (or consequence).
- The proposition $q \rightarrow p$ is the **converse** of A.
- The proposition $\neg q \rightarrow \neg p$ is the **contrapositive** of A.

4.5. Bi-conditional: $\leftrightarrow$

$$1 \leftrightarrow 1 = 1$$

$$1 \leftrightarrow 0 = 0$$

$$0 \leftrightarrow 1 = 0$$

$$0 \leftrightarrow 0 = 1$$

4.6. Exclusive or: XOR, $\oplus$

$$1 \oplus 1 = 0$$

$$1 \oplus 0 = 1$$

$$0 \oplus 1 = 1$$

$$0 \oplus 0 = 0$$

5. Operator Precedence

Operators are applied in the following order (ascending):

1. $\neg$
2. $\land$
3. $\lor$
4. $\rightarrow$
5. $\leftrightarrow$

For example:

$$p \rightarrow p \land \neg q \lor s \equiv (p \rightarrow ((p \land (\neg q)) \lor s))$$

6. Descriptors

A formula that's always true is called a **tautology**. A formula that is true for at least one scenario is **consistent**. A formula that's never true is **inconsistent**. A formula can also result in a **contradiction**.

7. Equivalances

Formulas are equivalanent if they result in the same logical outcomes.

For example (*De Morgan's Laws*):

$$\neg(p \wedge q) \equiv \neg p \vee \neg q$$
$$\neg(p \vee q) \equiv \neg p \wedge \neg q$$

For example:

$$\neg(\mathit{true} \wedge \mathit{true}) \equiv \mathit{false} \vee \mathit{false} \equiv \mathit{false}$$
$$\neg \mathit{true} \vee \neg \mathit{true} \equiv \neg(\mathit{true} \wedge \mathit{true}) = \neg \mathit{true} = \mathit{false}$$

8. Quantifiers

We use the symbol $\exists$ to indicate the existence of something (**existential quantifier**).

$$\exists x \quad \text{odd}(x)$$

This means that there exists some x that is odd.

We denote the **universal quantifier** as $\forall$.

$$\forall x(\text{odd}(x) \vee \text{even}(x))$$

This meant that **for all** x the number is either even or odd.

Other examples, "All Ps are Qs":

$$\forall x(P(x) \rightarrow Q(x))$$

And "No Ps are Qs":

$$\forall x(P(x) \rightarrow \neg Q(x))$$

8.1. Quantifiers to Connectives

$\exists x, P(x)$ where $x \in \{x_1, x_2, ..., x_n\}$ means that there exists some x for which $P(x)$ is true.

Denoted alternatively:

$$\exists x, P(x) \equiv P(x_1) \vee P(x_2) \vee ... \vee P(x_3)$$

We can also conclude:

$$\neg \exists x, P(x) \equiv \neg(P(x_1) \vee P(x_2) \vee ... \vee P(x_3))$$
$$\neg \exists x, P(x) \equiv \neg P(x_1) \wedge \neg P(x_2) \wedge ... \wedge \neg P(x_3)$$
$$\neg \exists x, P(x) \equiv \forall x, \neg P(x)$$

8.2. De Morgan's Law for Negation

$$\neg \forall x P(x) \equiv \exists x \neg P(x)$$
$$\neg \exists P(x) \equiv \forall x \neg P(x)$$

9. Laws of Propositional Logic

9.1. Logic 1

	Disjunction	**Conjunction**
idempotent laws	$p \vee p \equiv p$	$p \wedge p \equiv p$
commutative laws	$p \vee q \equiv q \vee p$	$p \wedge q \equiv q \wedge p$
associative laws	$(p \vee q) \vee r \equiv p \vee (q \vee r)$	$(p \wedge q) \wedge r \equiv p \wedge (q \wedge r)$
distributive laws	$p \vee (q \wedge r) \equiv (p \vee q) \wedge (p \vee r)$	$p \wedge (q \vee r) \equiv (p \wedge q) \vee (p \wedge r)$
identity laws	$p \vee F \equiv p$	$p \wedge T \equiv p$
domination laws	$p \vee T \equiv T$	$p \wedge F \equiv F$

9.2. Logic 2

	Disjunction	Conjunction
De Morgan’s laws	$\neg(p \vee q) \equiv \neg p \wedge \neg q$	$\neg(p \wedge q) \equiv \neg p \vee \neg q$
absorption laws	$p \vee (p \wedge q) \equiv p$	$p \wedge (p \vee q) \equiv p$
negation laws	$p \vee \neg p \equiv T$	$p \wedge \neg p \equiv F$
double negation law	$\neg\neg p \equiv p$	

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