

Program Verification Exam

PROJECT A: IVL1 → DSA
PROJECT B: BINARY SEARCH TREES



PROJECT A:
IVL1 → DSA

Every assignment on a path through C assigns to a fresh variable.

So, we need to:

- Introduce **multiple versions** of each variable
 - **Left-side** of **assignment** is always **fresh**
 - Make sure to always **read latest version**
- In $C1 \parallel C2$, **synchronize choices** on the last version of the variables

```

method foo(x: Int) returns (y: Int)
  requires x > 0
  ensures y > x
{
  var z: Int := x + 1

  if(z > 20){
    y := z + 1 }
  } else {
    z := z + z }
    y := z + 2 }
    y := y + 3 }
  }
  y := y + 4 }
}

```



```

-----IVL1-----
method foo{
  my_x_foo: Int
  my_y_foo: Int
  orig_my_x_foo: Int := my_x_foo
  assume my_x_foo > 0
  my_z_foo: Int := my_x_foo + 1
  {assume my_z_foo > 20
  my_y_foo := my_z_foo + 1}
  []
  {assume !(my_z_foo > 20)
  my_z_foo := my_z_foo + my_z_foo
  my_y_foo := my_z_foo + 2
  my_y_foo := my_y_foo + 3}
  my_y_foo := my_y_foo + 4
  assert my_y_foo > orig_my_x_foo
}

```



```

-----DSA-----
method foo{
  my_x_foo_0: Int
  my_y_foo_0: Int
  orig_my_x_foo_0: Int
  my_z_foo_1: Int
  my_y_foo_1: Int
  my_y_foo_2: Int
  my_z_foo_0: Int
  orig_my_x_foo_0 := my_x_foo_0
  assume my_x_foo_0 > 0
  my_z_foo_0 := my_x_foo_0 + 1
  {assume my_z_foo_0 > 20
  my_y_foo_0 := my_z_foo_0 + 1
  my_z_foo_1 := my_z_foo_0
  my_y_foo_1 := my_y_foo_0}
  []
  {assume !(my_z_foo_0 > 20)
  my_z_foo_1 := my_z_foo_0 + my_z_foo_0
  my_y_foo_0 := my_z_foo_1 + 2
  my_y_foo_1 := my_y_foo_0 + 3}
  my_y_foo_2 := my_y_foo_1 + 4
  assert my_y_foo_2 > orig_my_x_foo_0
}

```

To keep track of the variables versions:

id counter



`varsMap: Map<string, int>`

```
// translates a IVL1 statement into a DSA statement
let rec IVL1_stmt_to_DSA (cmd: v1Statement, varsMap: Map<string, int>) : DSASTatement list * Map<string, int> =
  match cmd with
  | V1Assignment(i, e) ->

    let (e', varsMap') = update_vars_in_expression (Some(i), e, varsMap)

    let count =
      if Map.containsKey i varsMap' then
        Map.find i varsMap' + 1
      else
        0

    let varsMap' = Map.add i count varsMap

    ([ DSAAssignment("${i}_{count}", e') ], varsMap')
```

```
// update the vars ids in an expression, according to the varsMap, to translate them from IVL1 into DSA
let rec update_vars_in_expression
(
  idAssign: ident option,
  e: expr,
  varsMap: Map<string, int>
) : expr * Map<string, int> =
match e with
| Ref id ->
  let count =
    if Map.containsKey id varsMap then
      Map.find id varsMap
    else
      0

  let varsMap' = Map.add id count varsMap

  (Ref("${id}_{count}"), varsMap')
```

```
let rec fix_choice_branches (c1VarsMap, c2VarsMap) =  
  let c1Seq = Map.toSeq c1VarsMap |> List.ofSeq // convert maps into seq to iterate over them  
  let c2Seq = Map.toSeq c2VarsMap |> List.ofSeq  
  
  let c1StmtList = add_assignments_to_choice_branch1 (c1Seq, c2Seq, [])  
  let c2StmtList = add_assignments_to_choice_branch2 (c1Seq, c2Seq, [])  
  
  (c1VarsMap, c2VarsMap, c1StmtList, c2StmtList)
```

```
// add the assignments needed in the first branch of the choice to align the variables counter in the branches
let rec add_assignments_to_choice_branch1 (c1Seq, c2Seq, c1StmtList) =
  match c1Seq, c2Seq with
  | _, (id2, c2Value) :: rest2 ->
    match find_value_by_id id2 c1Seq with
    | Some c1Value when c2Value > c1Value ->
      let c1Stmt = DSAAssignment($"{id2}_{c2Value}", Ref($"{id2}_{c1Value}"))
      add_assignments_to_choice_branch1 (c1Seq, rest2, c1Stmt :: c1StmtList)

    | Some c1Value when c2Value <= c1Value -> add_assignments_to_choice_branch1 (c1Seq, rest2, c1StmtList)

    | _ ->
      let c1Stmt = DSAAssignment($"{id2}_{c2Value}", Ref($"{id2}_{0}"))

      add_assignments_to_choice_branch1 (
        c1Seq,
        rest2,
        if (c2Value <> 0) then c1Stmt :: c1StmtList else c1StmtList
      )
  | _ -> (c1StmtList)
```




PROJECT B: BINARY SEARCH TREES

Binary Search Trees (BST): definition

```

predicate bst(node: Ref) {
  acc(node.elem) && acc(node.left) && acc(node.right) &&

  (node.left != null ==> bst(node.left) &&
   |tree_content(node.left)| > 0 &&
   node.elem > max_seq(tree_content(node.left))) &&

  (node.right != null ==> bst(node.right) &&
   |tree_content(node.right)| > 0 &&
   node.elem < min_seq(tree_content(node.right)))
}

```

```

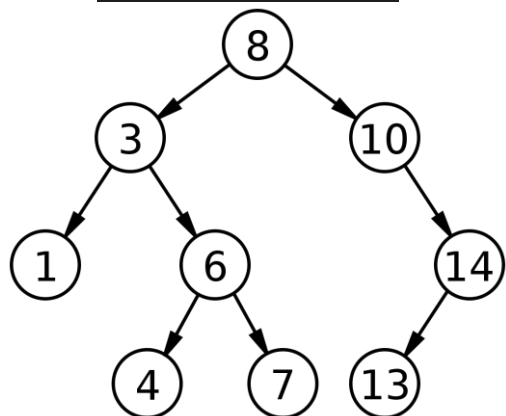
function max_seq(s : Seq[Int]) : Int
  requires |s| > 0
  ensures Seq(s[0]) ++ s[1..] == s
  ensures forall x : Int :: x in s ==> x <= result
  ensures result in s
  {
    |s| == 1 ? s[0] : max(s[0], max_seq(s[1..]))
  }

```

```

field elem : Int
field left : Ref
field right : Ref

```



```


function tree_content(tree: Ref) : Seq[Int]
  requires bst(tree)
  ensures |result| > 0
  {
    unfolding bst(tree) in
    (tree.left == null && tree.right == null) ? Seq(tree.elem) :
    (tree.left == null && tree.right != null) ? Seq(tree.elem) ++ tree_content(tree.right) :
    (tree.left != null && tree.right == null) ? tree_content(tree.left) ++ Seq(tree.elem) :
    tree_content(tree.left) ++ Seq(tree.elem) ++ tree_content(tree.right)
  }

```

```
method bst_insert(tree: Ref, val: Int)
  requires bst(tree)
  requires acc(time_credit(), (height(tree) + 1)/1)
  ensures bst(tree)
  ensures old(seq_to_set(tree_content(tree))) union
    Set(val) == seq_to_set(tree_content(tree))
{
  consume_time_credit()
  unfold bst(tree)

  if(val < tree.elem) {
    if(tree.left == null) {
      var new_tree_left : Ref := new(*)
      new_tree_left.elem := val
      new_tree_left.left := null
      new_tree_left.right := null

      tree.left := new_tree_left
      fold bst(tree.left)
    } else {
      bst_insert(tree.left, val)
    }
  } else {
    if(val > tree.elem) { ... }
  }
  fold bst(tree)
}
```



```
function height(tree: Ref) : Int
  requires bst(tree)
  ensures result >= 0
{
  unfolding bst(tree) in (tree.left == null) ?
    ((tree.right == null) ? 0 : 1 + height(tree.right)) :
    ((tree.right == null) ? 1 + height(tree.left) :
      1 + max(height(tree.left), height(tree.right)))
}
```



THANK YOU