

ACCU
2023

CONCURRENCY APPROACHES

PAST, PRESENT, AND FUTURE

LUCIAN RADU TEODORESCU

Concurrency Approaches: past, present, and future

LUCIAN RADU TEODORESCU
GARMIN

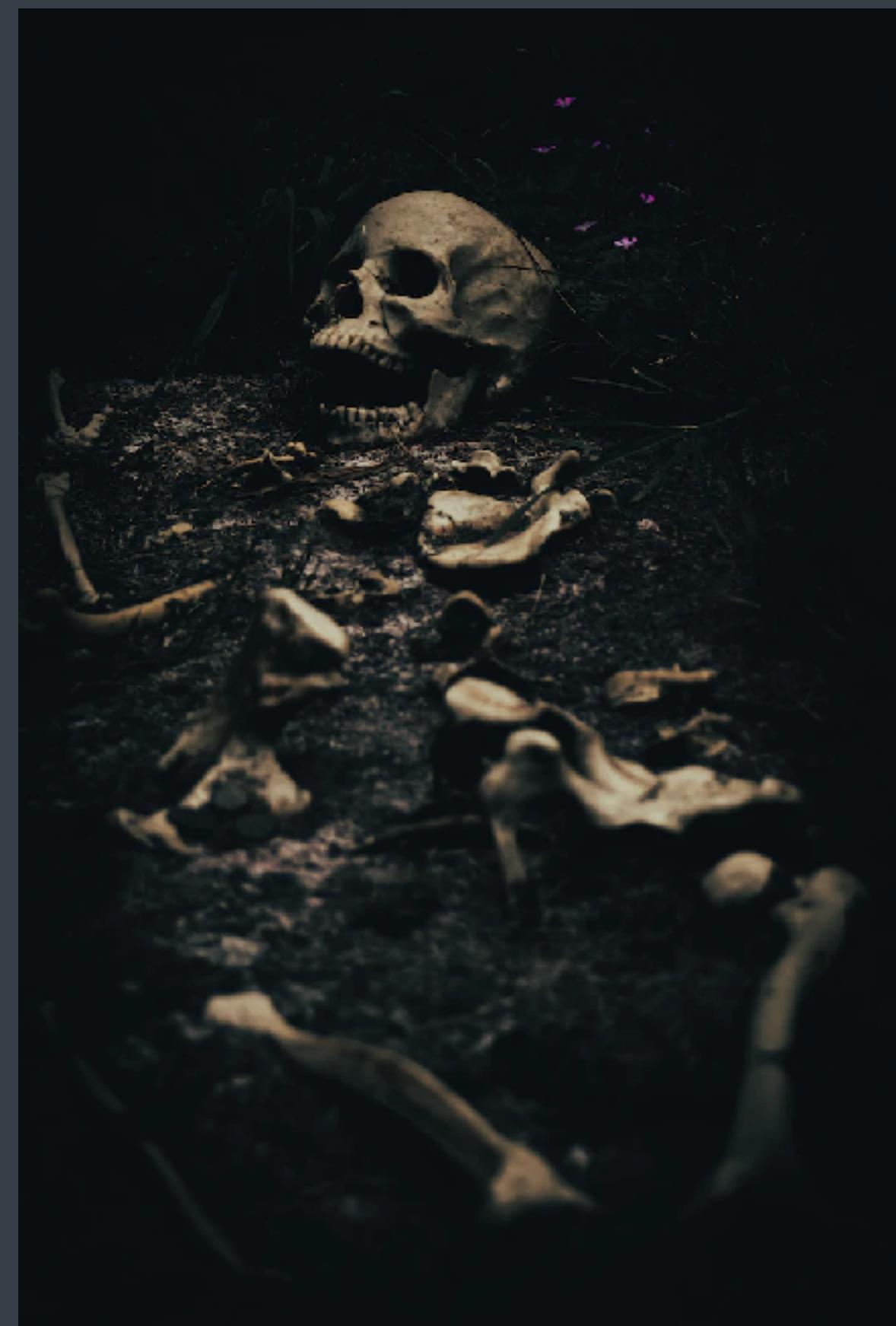
spoiler

concurrency is **HARD**

harder than **public speaking**



public speaking feared more than **death**



spoiler

concurrency is **HARD**
very

overload 174

APRIL 2023 £4.50

In Search of a Better Concurrency Model

Lucian Radu Teodorescu presents current plans for concurrency in the VAL programming language

Drawing a Line Under Aligned Memory

Paul Floyd reminds us about various aligned memory functions

C++20 Concepts: Testing Constrained Functions

Andreas Fertig gives a worked example of testing constraints on functions or classes and other template constructs

Meta Verse

Teedy Deigh turns on, jacks in, and checks out the immersive experience

materials



Agenda

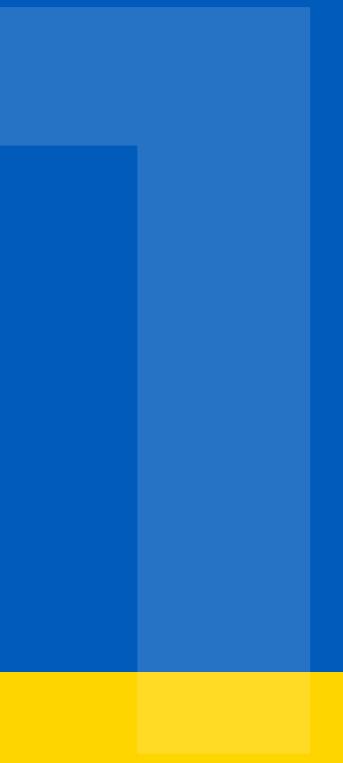
Concurrency models

Concurrency goals

Past and present models

A possible future

Concurrency models



concurrency

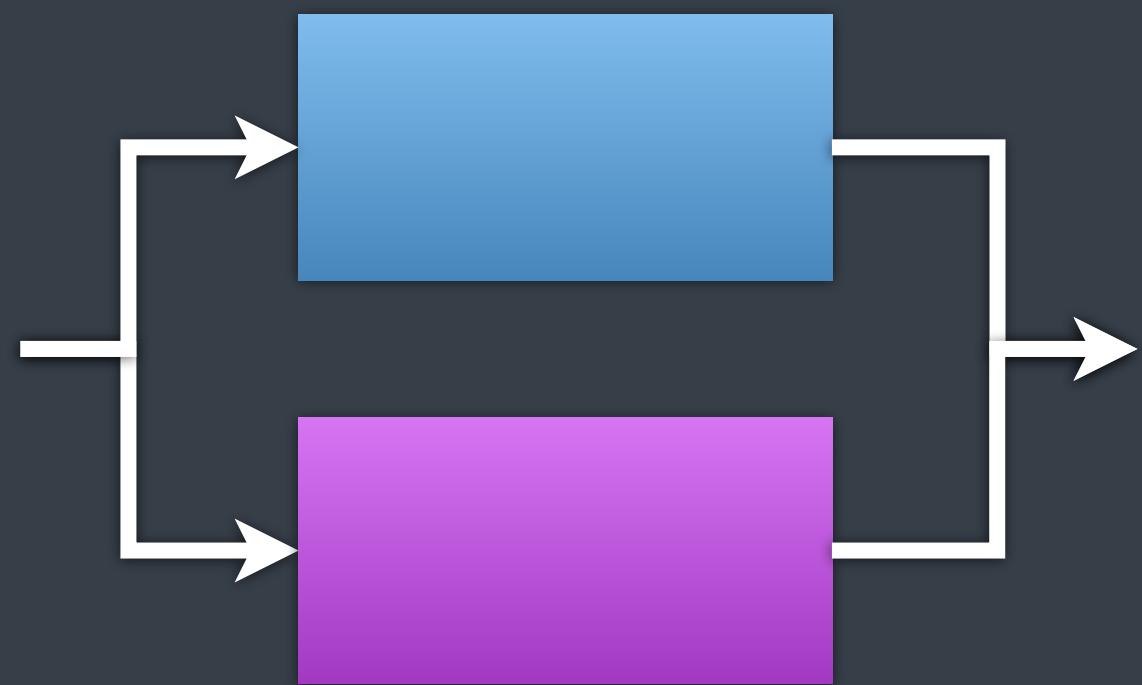
set of rules

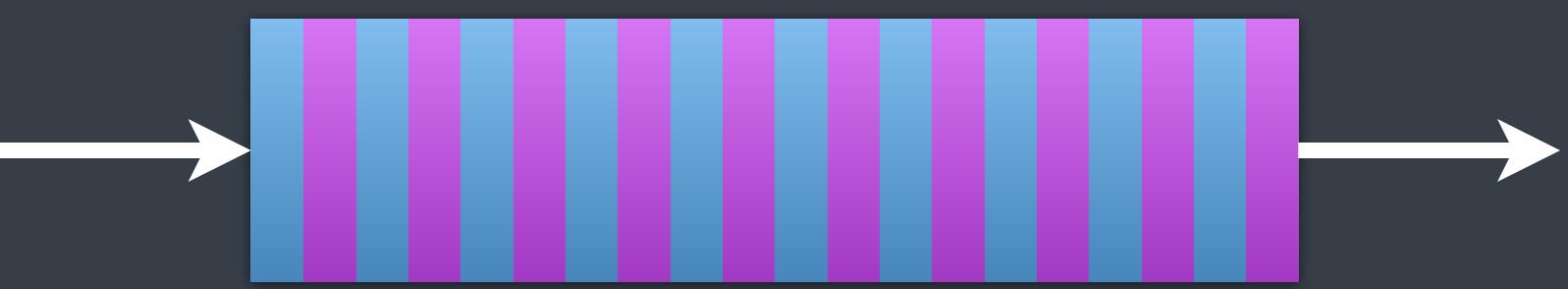
multiple activities in overlapping time periods

partial order over activities

parallelism

simultaneous execution of activities





concurrency

parallelism

design time

run time

independent of hardware

dependent on hardware

hardware threads

enable parallelism

software threads

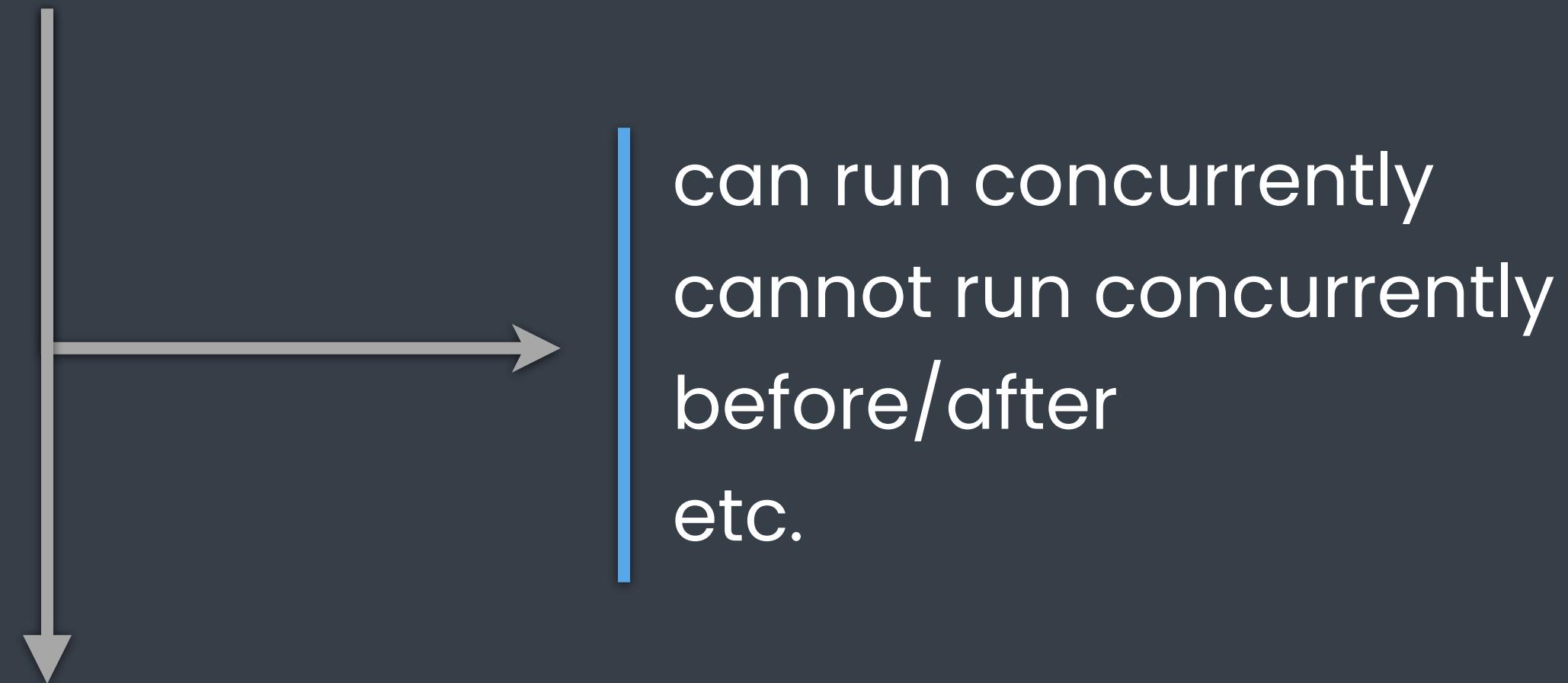
enable expressing concurrency

focus on **concurrency**



focus on **concurrency**

rules overlapping activities



how?

example

pthread_create()
pthread_join()

pthread_mutex_init()
pthread_mutex_destroy()
pthread_mutex_lock()
pthread_mutex_unlock()

example

can **directly** express:

run activity concurrently
join concurrent activities

don't run concurrently

example

can **indirectly** express:
activity before other activity

example

guarantees:

none

example

possible **problems**:

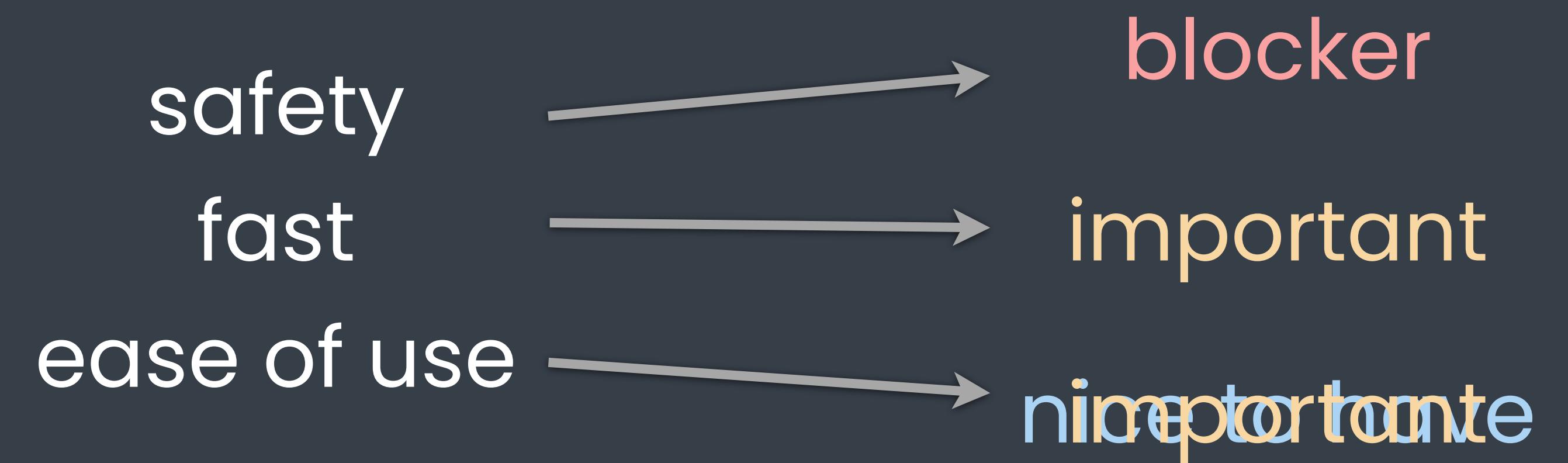
- deadlocks,
- malign race conditions,
- blocking threads more than needed,
- oversubscription,
- high latency,
- etc.

Concurrency goals

2



multiple levels





SAFETY

S1: race conditions

The concurrency model
shall not allow undefined behaviour
caused by **race conditions.**

example

shared state

S2: deadlocks

The concurrency model
shall not allow deadlocks.

example

acquiring two resources

starvation?

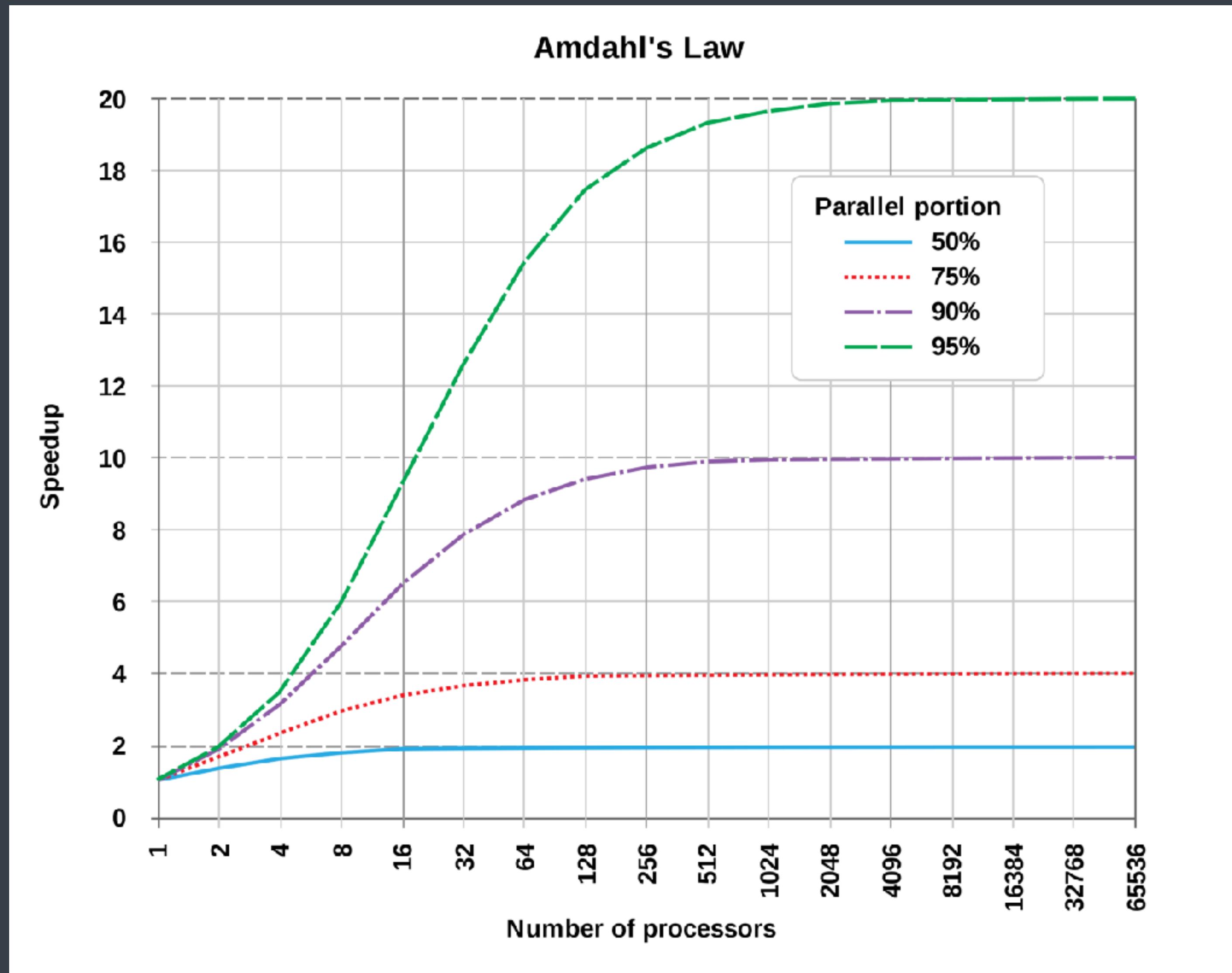
cannot guarantee lack of starvation



FAST

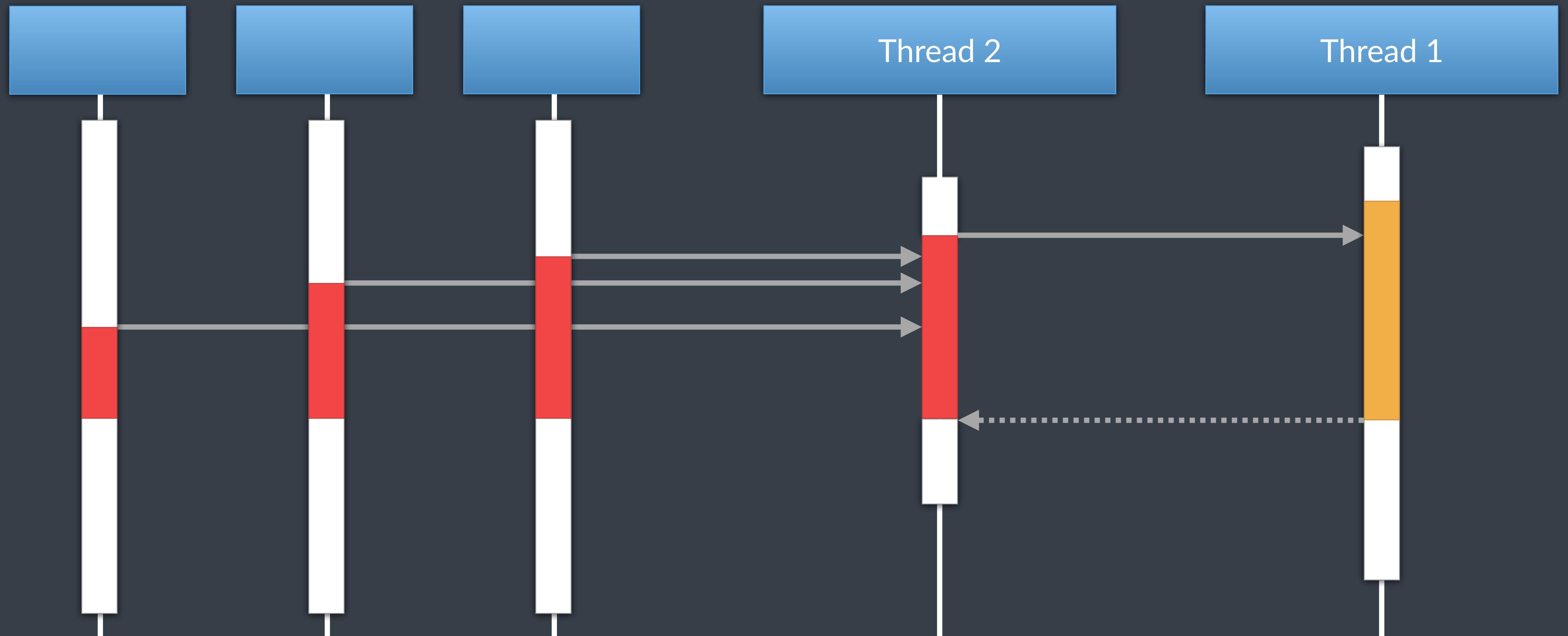
F1: scalability

For applications that express enough concurrent behaviour,
the concurrency model shall guarantee that
the **performance** of the application
scales with the number of **hardware threads**.



F2: blocking

The concurrency model
shall not require **blocking threads**
(keeping the threads idle for longer periods of time).



can dramatically
reduce performance

mutex == bottleneck

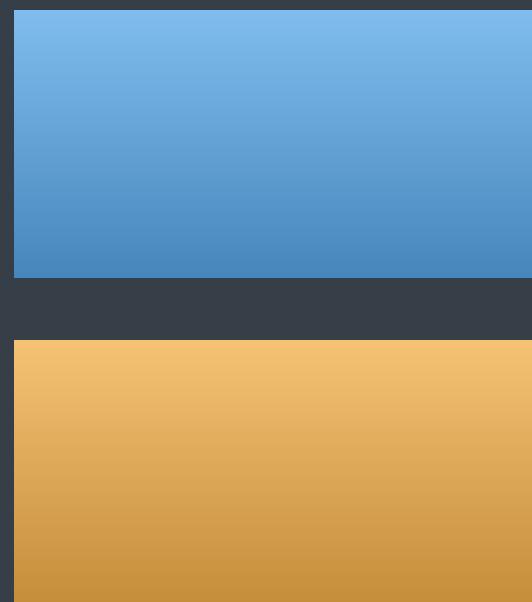
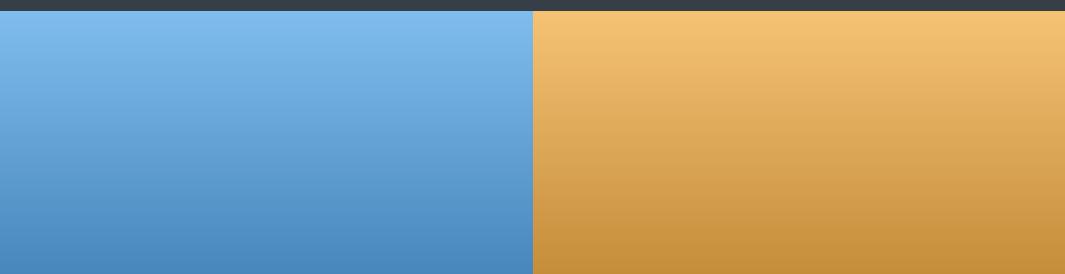
F3: oversubscription

The concurrency model shall allow
limiting the oversubscription on hardware threads.

multiple threads on the same core

2 x 1 second

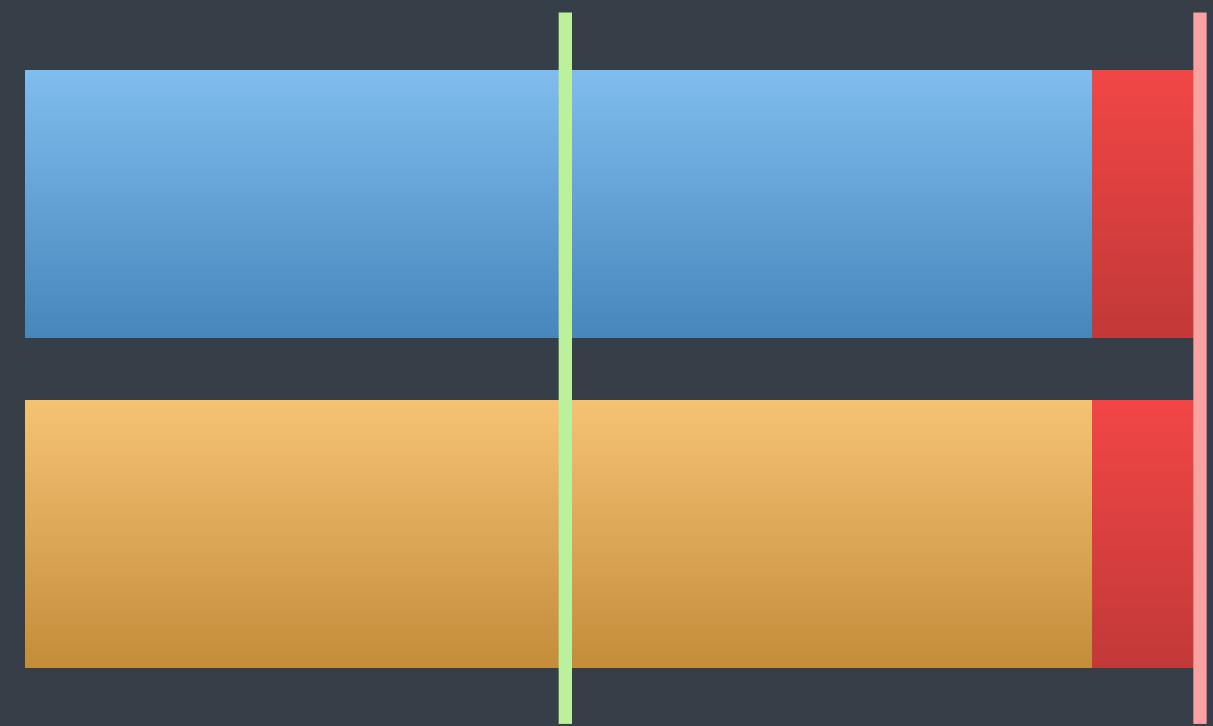
1 core



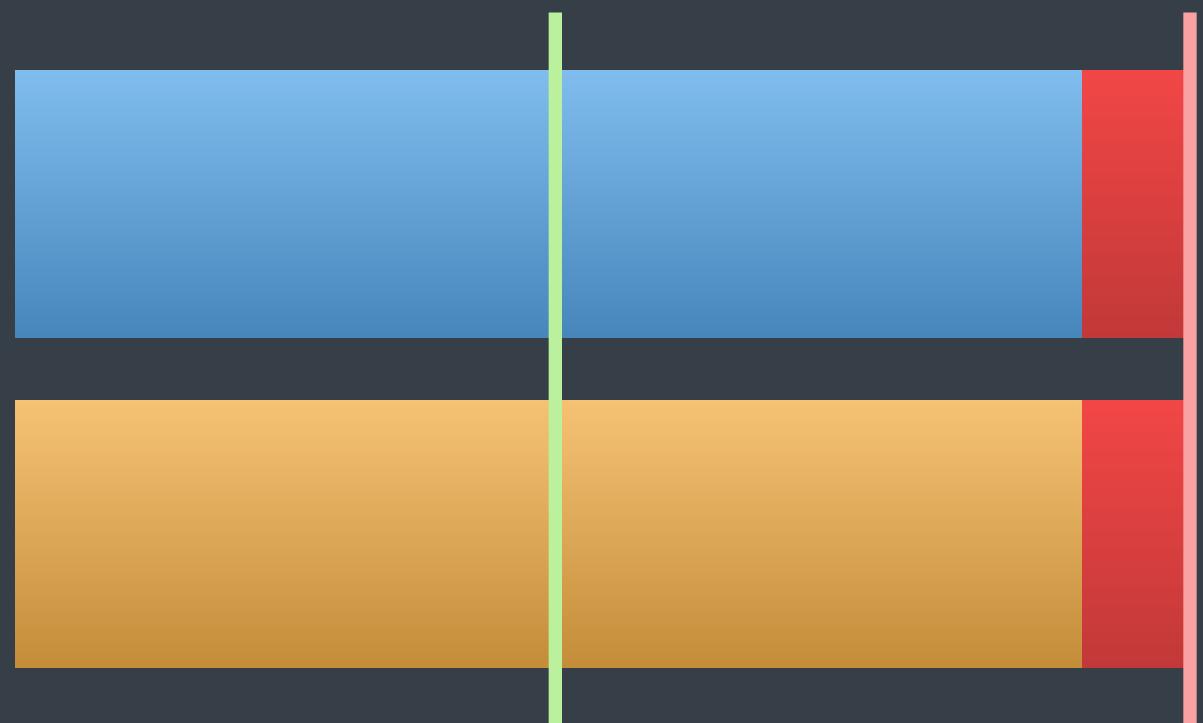
2 threads, 1 core



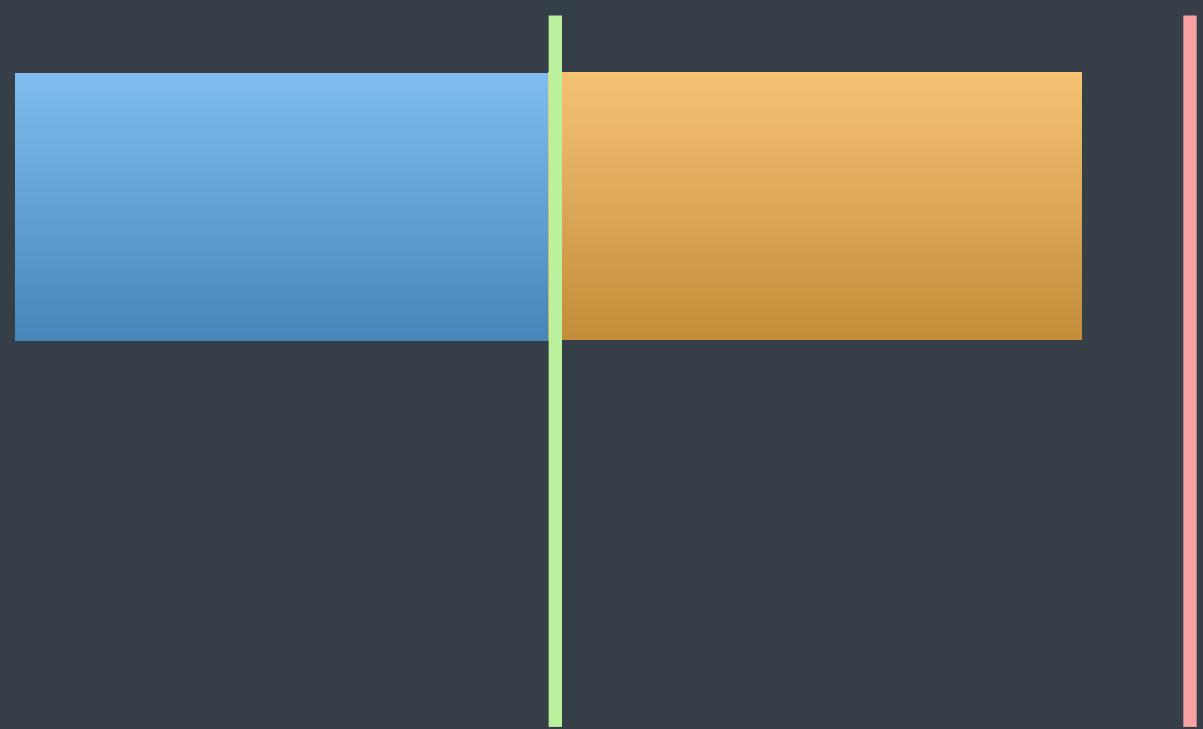
2 threads, 1 core



good ?



alternative



live demo

for
a
variety
of
reasons,
this
is
a
book
eagerly
awaited

to
say
that
their
patience
has
been
rewarded
would
be
an
understatement

for
a
variety
of
reasons,
this
is
a
book
eagerly
awaited
to
say
that
their
patience
has
been
rewarded
would
be
an
understatement

ready?

for
a
variety
of
reasons,
this
is
a
book
eagerly
awaited
to
say
that
their
patience
has
been
rewarded
would
be
an
understatement

important !

no blocking
no oversubscription

F4: synchronization

The concurrency model
shall not require any synchronisation code
during the execution of the tasks
(except in the concurrency-control code).

not just blocking

all synchronisation costs

F5: memory allocation

The concurrency model
shall not require **dynamic memory allocation**
(unless type-erasure is requested by the user).

A painting of a Venetian canal scene. In the foreground, several gondolas filled with people are moving through the water. One gondola has a person rowing. In the background, there is a large, ornate building with a dome and a tall, thin tower. The sky is blue with some clouds.

EASE OF USE

E1: structured

The concurrency model shall match the description of
structured concurrency.

accu
2022

STRUCTURED CONCURRENCY

LUCIAN RADU TEODORESCU



structured concurrency

1. abstractions as building blocks
2. recursive decomposition
3. local reasoning
4. single entry, single exit point
5. soundness and completeness

negative example

threads and locks

E2: same syntax/semantics

Concurrent code shall be expressed using the
same syntax and semantics
as non-concurrent code.

negative example

```
auto work1() -> int;
auto work2() -> double;
auto work3() -> std::string;
auto combine_res(int i, double d, const std::string& s) -> int;

auto compute_in_parallel() -> int {
    static_thread_pool pool{8};
    ex::scheduler auto sched = pool.get_scheduler();

    ex::sender auto work =
        ex::when_all(
            ex::schedule(sched) | ex::then(work1),
            ex::schedule(sched) | ex::then(work2),
            ex::schedule(sched) | ex::then(work3)
        );
    auto [i, d, s] = std::this_thread::sync_wait(std::move(work)).value();
    return combine_res(i, d, s);
}
```

E3: function colouring

Function colouring shall not be required
for expressing concurrent code.

negative example

```
public int RegularFunction() { ... }
```

```
public async Task<int> AsyncFunction() { ... }
```

negative example

```
int regular_function(int input) { ... }
```

```
task<int> coroutine_is_colored(int input) { ... }
```

main problem

it's pervasive

E4: extra code

Except for concurrency-control code,
the user shall **not** be required to
add any extra code in concurrent code,
versus non-concurrent code.

synchronisation

only present in concurrent-control code

E5: minimal rules

The concurrency model should have a **minimum set of rules** for the user to follow to stay within the model.

traditional textbook

creating threads
problems with threading
creating locks
more problems with threading
avoiding these problems
tips for efficiency

...

Past and present models

3



concurrency models evaluation

	S1	S2	F1	F2	F3	F4	F5	E1	E2	E3	E4	E5
threads and locks	✗	✗	✗	✗	✗	✗	!	✗	✗	✓	✗	✗
tasks	!	✓	✓	✓	✓	✓	✗	✗	✗	✓	✓	✗
C# asynchronous model	!	✓	✓	✗	✗	✓	✗	✓	!	✗	✓	✗
C++ coroutines	!	✓	✓	!	✓	✓	✓	✗	!	✓	✓	✗
senders/receivers	!	✓	✓	✓	✓	✓	✓	!	✗	✓	✓	✗
Rust's fearless concurrency	✓	!	!	✗	✗	✗	✗	✗	✗	✓	✗	✗

locks and threads

	S1	S2	F1	F2	F3	F4	F5	E1	E2	E3	E4	E5
threads and locks	✗	✗	✗	✗	✗	✗	!	✗	✗	✓	✗	✗
tasks	!	✓	✓	✓	✓	✓	✗	✗	✗	✓	✓	✗
C# asynchronous model	!	✓	✓	✗	✗	✓	✗	✓	!	✗	✓	✗
C++ coroutines	!	✓	✓	!	✓	✓	✓	✓	!	✓	✓	✗
senders/receivers	!	✓	✓	✓	✓	✓	✓	✓	!	✓	✓	✗
Rust's fearless concurrency	✓	!	!	✓	✓	✓	✓	✓	!	✓	✓	✓

E3 = function colouring

tasks

	S1	S2	F1	F2	F3	F4	F5	E1	E2	E3	E4	E5
threads and locks	✗	✗	✗	✗	✗	✗	!	✗	✗	✓	✗	✗
tasks	!	✓	✓	✓	✓	✓	✗	✗	✗	✓	✗	
C# asynchronous model	!	✓	✓	✗	✗	✓	✗	✓	!	✗	✓	✗
C++ coroutines	!	✓	✓	!	✓	✓	✓	✓	!	✓	✓	✗
senders/receivers	!	✓	✓	✓	✓	✓						
Rust's fearless concurrency	✓	!	!									

F5 = memory allocation
E4 = extra code

C# asynchronous model

	S1	S2	F1	F2	F3	F4	F5	E1	E2	E3	E4	E5
threads and locks	---	---	---	---	---	---	!	---	---	✓	---	---
tasks	!	✓	✓	✓	✓	✓	---	---	---	✓	✓	---
C# asynchronous model	!	✓	✓	---	---	✓	---	✓	!	---	✓	---
C++ coroutines	!	✓	✓	!	✓	✓	---	✓	!	---	✓	---
senders/receivers	!	✓	✓	✓	✓	✓	---	✓	!	---	✓	---
Rust's fearless concurrency	✓	!	!	---								

F2 = blocking

F3 = oversubscription

F5 = memory allocation

E3 = function colouring

C++ coroutines

F5 = memory allocation
E3 = function colouring

tasks

C# asynchronous model

C++ coroutines

senders/receivers

Rust's fearless concurrency

	S1	S2	F1	F2	F3	F4	F5	E1	E2	E3	E4	E5
tasks	---	---	---	---	---	---	!	---	---	✓	---	---
C# asynchronous model	!	✓	✓	✓	✓	✓	---	---	---	✓	✓	---
C++ coroutines	!	✓	✓	!	✓	✓	---	✓	!	✓	✓	---
senders/receivers	!	✓	✓	✓	✓	✓	!	✓	---	✓	✓	---
Rust's fearless concurrency	✓	!	!	---	---	---	---	---	---	✓	---	---

senders/receivers

		S1	S2	F1	F2	F3	F4	F5	E1	E2	E3	E4	E5
tasks													
C# asynchronous model	!	✓	✓	✓	✓	✓	✓	!	✓	✓	✓	✓	✓
C++ coroutines	!	✓	✓	✓	!	✓	✓	✓	✓	✓	✓	✓	✓
senders/receivers	!	✓	✓	✓	✓	✓	✓	!	✓	✗	✗	✓	✗
Rust's fearless concurrency	✓	!	!	✓	✓	✓	✓	!	✓	✓	✓	✓	✓

Rust's fearless concurrency

S2 = deadlocks

F2 = blocking

F4 = synchronisation

F5 = memory allocation

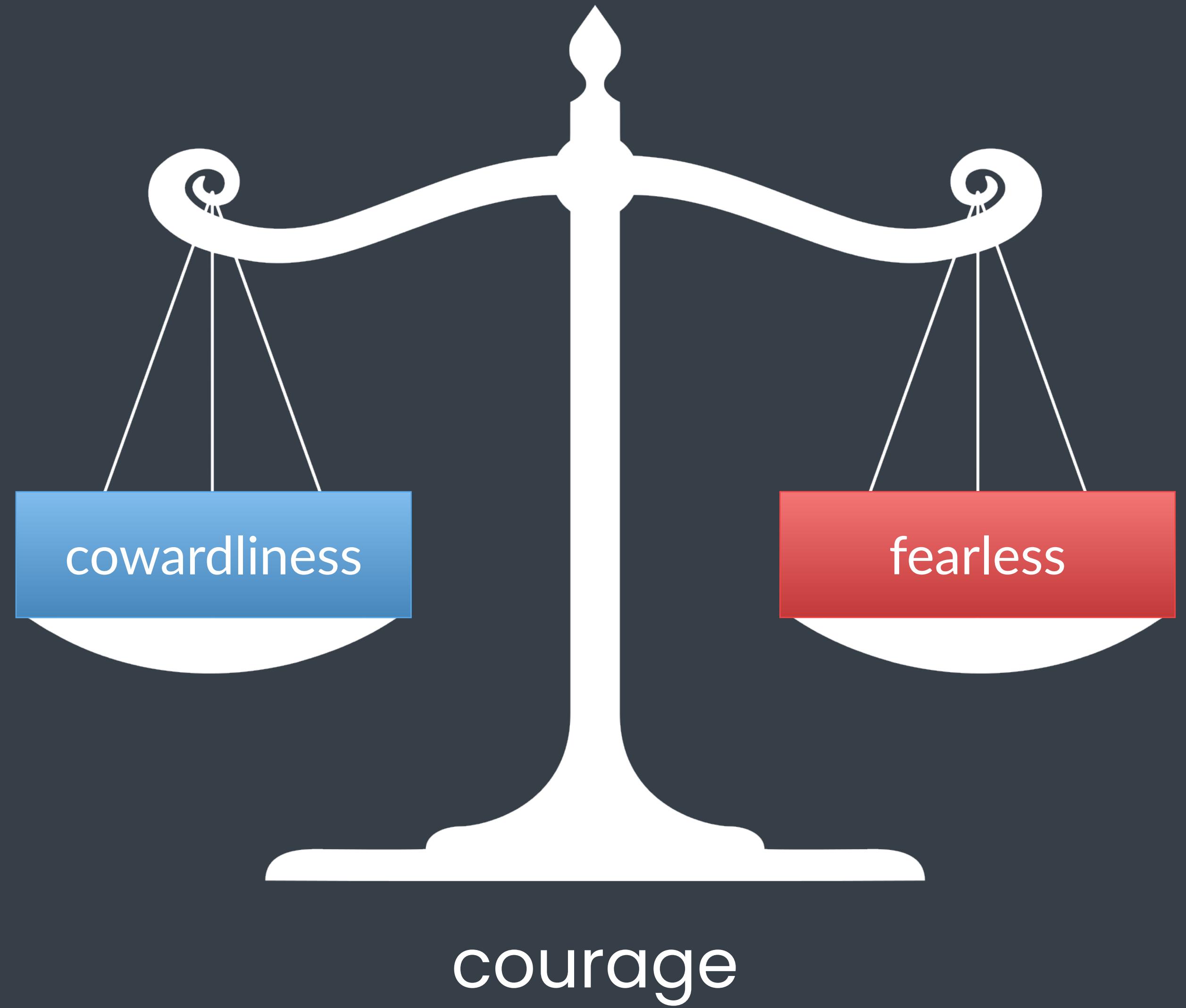
E1 = structured

E2 = same syntax/semantics

E4 = extra code

Rust's fearless concurrency

S1	S2	F1	F2	F3	F4	F5	E1	E2	E3	E4	E5
✗	✗	✗	✗	✗	✗	⚠	✗	✗	✓	✗	✗
⚠	✓	✓	✓	✓	✓	✗	✗	✗	✓	✗	✗
⚠	✓	✓	✗	✗	✓	✗	✓	✗	✓	✗	✗
⚠	✓	✓	⚠	✓	✓	✗	✓	✗	✓	✗	✗
⚠	✓	✓	✓	✓	✓	⚠	✓	✗	✓	✓	✗



A possible future



an experiment
within in an experiment

Val programming language



fast by definition
safe by default
simple

www.val-lang.dev/



concurrency experiment

good concurrency model?

goals

safe concurrency
fast as senders/receivers (or close)
no function colouring

proposed Val concurrency model

	S1	S2	F1	F2	F3	F4	F5	E1	E2	E3	E4	E5
threads and locks	🚫	🚫	🚫	🚫	🚫	🚫	⚠️	🚫	🚫	✓	🚫	🚫
tasks	⚠️	✓	✓	✓	✓	✓	✓	🚫	🚫	✓	✓	🚫
C# asynchronous model	⚠️	✓	✓	🚫	🚫	✓	✓	✓	⚠️	✓	✓	🚫
from C++ coroutines	⚠️	✓	✓	⚠️	⚠️	✓	✓	✓	✓	✓	✓	✓
from senders/receivers	✓	✓	✓	✓	✓	✓	⚠️	✓	🚫	✓	✓	🚫
Rust's fearless concurrency	✓	⚠️	⚠️	🚫	✓	✓	✓	🚫	🚫	✓	✓	🚫

starting example

```
fun long_task(input: Int) -> Int {  
    var result = input  
    for let i in 0 ..< 42 {  
        sleep(1)  
        &result += 1  
    }  
    return result  
}  
  
fun greeting_task() -> Int {  
    print("Hello world! Have an int.")  
    return 13  
}  
  
fun example() -> Int {  
    var handle = long_task(input: 0)  
    let x = greeting_task()  
    let y = handle  
    return x+y  
}
```

starting example

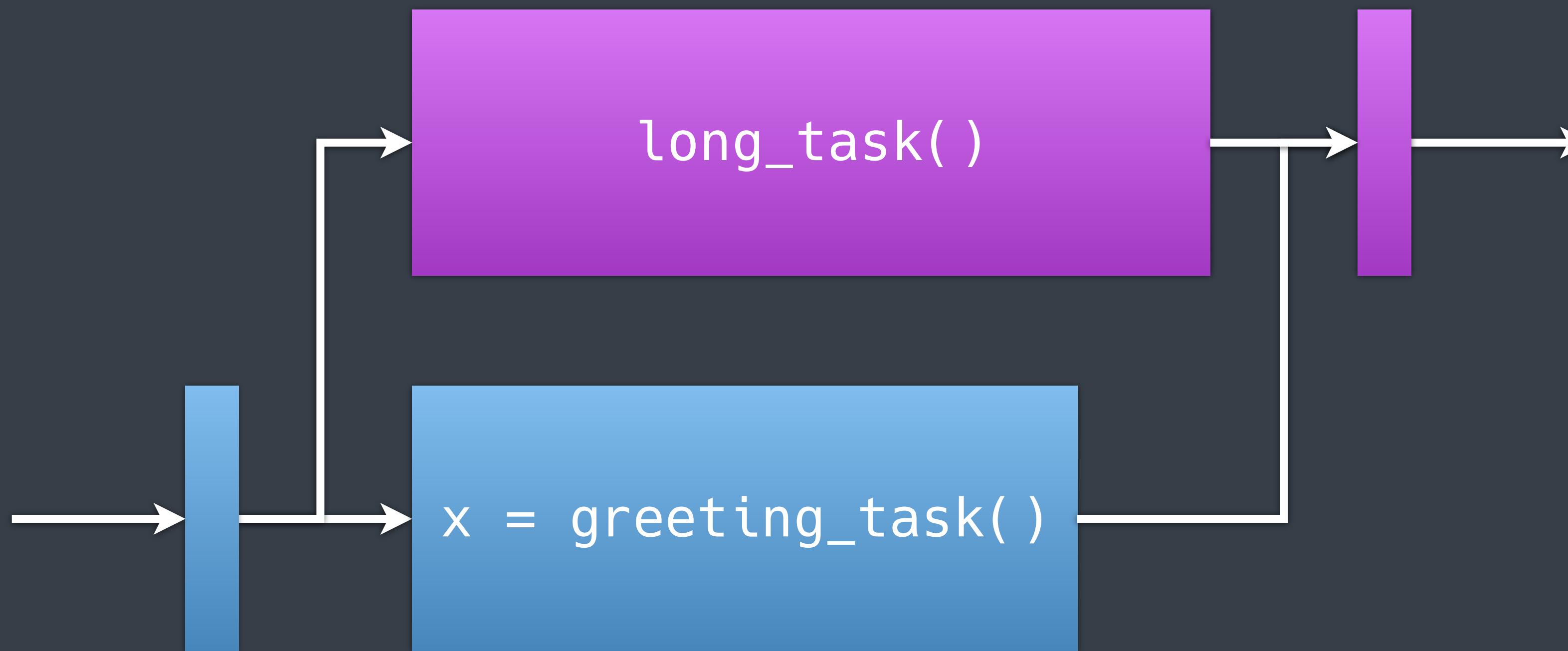
```
fun long_task(input: Int) -> Int {  
    var result = input  
    for let i in 0 ..< 42 {  
        sleep(1)  
        &result += 1  
    }  
    return result  
}  
  
fun greeting_task() -> Int {  
    print("Hello world! Have an int.")  
    return 13  
}  
  
fun concurrency_example() -> Int {  
    var handle = spawn long_task(input: 0) // create concurrent work  
    let x = greeting_task()  
    let y = handle.await() // joining; switching threads  
    return x+y  
}
```

in C++, with coroutines

```
int long_task(int input) {
    int result = input;
    for (int i=0; i<42; i++) {
        sleep(1);
        result += 1;
    }
    return result;
}
task<int> long_task_wrapper(int input) {
    co_await global_task_pool.enter_thread();
    co_return long_task(input);
}

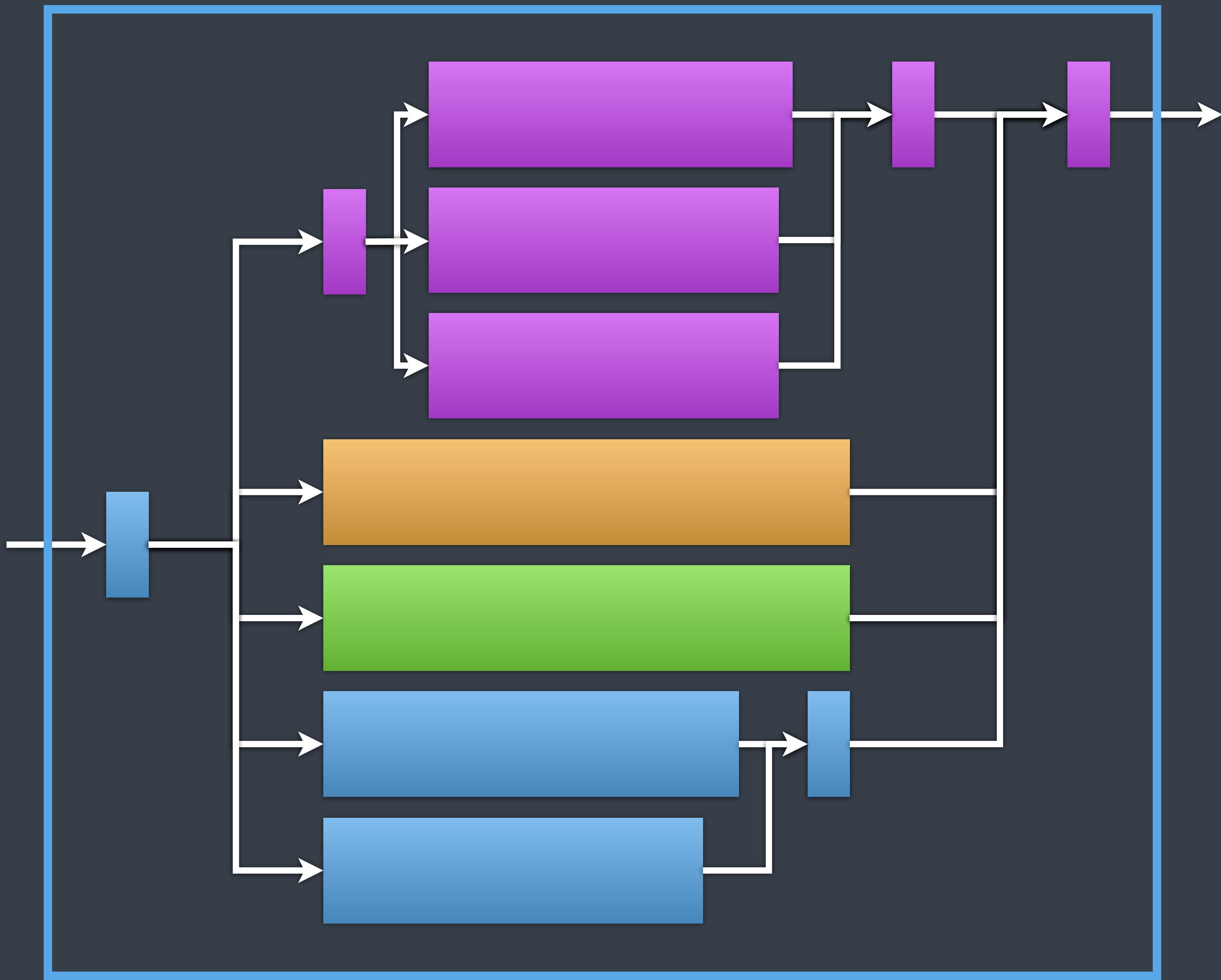
int greeting_task() {
    std::cout << "Hello world! Have an int.";
    return 13;
}

task<int> concurrency_example() {
    auto handle = long_task_wrapper(0);
    auto x = greeting_task()
    auto y = co_await handle;
    co_return x+y;
}
```



computation

one entry, one exit point
can switch threads in the middle



computations, examples

senders

C++ coroutines

C# coroutines

Val coroutines

accu
2022

STRUCTURED CONCURRENCY

LUCIAN RADU TEODORESCU



EASE OF USE



EI: structured concurrency

use **computations** as the main abstraction

structured concurrency

1. abstractions as building blocks
2. recursive decomposition
3. local reasoning
4. single entry, single exit point
5. soundness and completeness

acscu
2022

STRUCTURED CONCURRENCY

LUCIAN RADU TEODORESCU



main abstraction

function = computation

coroutines

concurrency is embedded
within functions

local reasoning

guaranteed by Val

local reasoning

```
fun long_task(input: Int) -> Int { ... }

fun greeting_task() -> Int { ... }

fun concurrency_example() -> Int {
    var handle = spawn long_task(input: 0)
    let x = greeting_task()
    let y = handle.await()
    return x+y
}
```

E2: same syntax/semantics

just use functions

in C++, with coroutines

```
int long_task(int input) {
    int result = input;
    for (int i=0; i<42; i++) {
        sleep(1);
        result += 1;
    }
    return result;
}
task<int> long_task_wrapper(int input) {
    co_await global_task_pool.enter_thread();
    co_return long_task(input);
}

int greeting_task() {
    std::cout << "Hello world! Have an int.";
    return 13;
}

task<int> concurrency_example() {
    auto handle = long_task_wrapper(0);
    auto x = greeting_task()
    auto y = co_await handle;
    co_return x+y;
}
```

in C++, with coroutines

```
int long_task(int input) {
    int result = input;
    for (int i=0; i<42; i++) {
        sleep(1);
        result += 1;
    }
    return result;
}
task<int> long_task_wrapper(int input) {
    co_await global_task_pool.enter_thread();
    co_return long_task(input);
}

int greeting_task() {
    std::cout << "Hello world! Have an int.";
    return 13;
}

task<int> concurrency_example() {
    auto handle = long_task_wrapper(0);
    auto x = greeting_task()
    auto y = co_await handle;
    co_return x+y;
}
```

in C++, with coroutines

```
int long_task(int input) {
    int result = input;
    for (int i=0; i<42; i++) {
        sleep(1);
        result += 1;
    }
    return result;
}
task<int> long_task_wrapper(int input) {
    co_await global_task_pool.enter_thread();
    co_return long_task(input);
}

int greeting_task() {
    std::cout << "Hello world! Have an int.";
    return 13;
}

task<int> concurrency_example() {
    auto handle = long_task_wrapper(0);
    auto x = greeting_task()
    auto y = co_await handle;
    co_return x+y;
}
```

same syntax & semantics

```
fun long_task(input: Int) -> Int {  
    var result = input  
    for let i in 0 ..< 42 {  
        sleep(1)  
        &result += 1  
    }  
    return result  
}  
  
fun greeting_task() -> Int {  
    print("Hello world! Have an int.")  
    return 13  
}  
  
fun concurrency_example() -> Int {  
    var handle = spawn long_task(input: 0)  
    let x = greeting_task()  
    let y = handle.await() // switching threads?  
    return x+y  
}
```

E3: function colouring

coroutines are pervasive
caller of a coroutine is usually a coroutine

C++ colouring example

```
int f1(int input) { ... }

int f2(int input) { return f1(input); }

int f3(int input) { return f2(input); }

int f4(int input) { return f3(input); }

int main() {
    f4(19);
}
```

C++ colouring example

```
task<int> f1(int input) { ... }

int f2(int input) { return co_await f1(input); }

int f3(int input) { return f2(input); }

int f4(int input) { return f3(input); }

int main() {
    f4(19);
}
```

C++ colouring example

```
task<int> f1(int input) { ... }

task<int> f2(int input) { co_return co_await f1(input); }

int f3(int input) { return co_await f2(input); }

int f4(int input) { return f3(input); }

int main() {
    f4(19);
}
```

C++ colouring example

```
task<int> f1(int input) { ... }

task<int> f2(int input) { co_return co_await f1(input); }

task<int> f3(int input) { co_return co_await f2(input); }

int f4(int input) { return co_await f3(input); }

int main() {
    f4(19);
}
```

C++ colouring example

```
task<int> f1(int input) { ... }

task<int> f2(int input) { co_return co_await f1(input); }

task<int> f3(int input) { co_return co_await f2(input); }

task<int> f4(int input) { co_return co_await f3(input); }

int main() {
    f4(19);
}
```

C++ colouring example

```
task<int> f1(int input) { ... }

task<int> f2(int input) { co_return co_await f1(input); }

task<int> f3(int input) { co_return co_await f2(input); }

task<int> f4(int input) { co_return co_await f3(input); }

int main() {
    sync_wait( f4(19) );
}
```

concurrent intensive app

most of the functions are coroutines

no colouring in Val

```
fun f1(input: Int) -> Int { ... }

fun f2(input: Int) -> Int { f1(input) }

fun f3(input: Int) -> Int { f2(input) }

fun f4(input: Int) -> Int { f3(input) }

fun main() -> Int { f4(19) }
```

E4: extra code

no need for synchronisation code

concurrency-control code

start of activity
end of activity

spawn f

schedule(global_scheduler) | then(f)

or

schedule(global_scheduler) | let_value(f)

h.await()

when_all(*cur_work*, *h_work*)

plus generalisations

different schedulers
spawning/joining multiple activities

E5: minimal rules

The concurrency model should have a **minimum set of rules** for the user to follow to stay within the model.

minimal rules

how work is spanned

how work is joined

minimal rules

no synchronisation
no sync/async distinction



FAST

from coroutines

F1: scalability

F2: blocking

F3: oversubscription

F4: synchronization

memory allocation

C++ coroutines need dynamic memory for stack frames

every call

coroutines all the way down

doesn't work

coroutines

stackless

stores one frame
suspend in called function

C++ coroutines

stackfull

stores the entire stack
suspend anywhere

Boost coroutine2

costs

stackless
calling

stackfull
spawning
exceeding stack space

to be confirmed!

success of the model

cost of using stackfull coroutines



SAFETY

S1: race conditions

by construction of Val
law of exclusivity

S2: deadlocks

no synchronisation primitives
activities don't have cycles

Takeaways

5



a **framework** for evaluating concurrency models

safety / fast / ease of use

proposed concurrency model

big advancements in usability
safety by default
fast?

more research

improving stackfull coroutines

also, on Val

ease of use
safety by default
fast?

on C++

new perspectives

Thank You



@LucTeo@techhub.social
lucteo.ro