## CS186 Discussion 3

(Join Algorithms)

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# Join Algorithms

#### Cost Notation

- [R] = number of pages in Table R
- p<sub>R</sub> = number of records per page of R
- |R| = number of records in R

$$|R| = p_R^* [R]$$

## Simple Nested Loop Join

#### for record r in R:

for record s in S:

if theta(r, s):

add join(r, s) to result

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## Page Nested Loop Join

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for page p_r in R:
for page p_s in S:
for record r in p_r:
for record s in p_s:
if theta(r, s):
add join(r, s) to result
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$$[R] + [R] * [S]$$

## Chunk Nested Loop Join

```
for chunk c_r in R:
for page p_s in S:
for record r in c_r:
for record s in p_s:
if theta(r, s):
add join(r, s) to result
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[R] = number of pages in Table Rp<sub>R</sub> = number of records per page of R|R| = number of records in R

$$[R] + [R] / (B - 2) * [S]$$

B-2 pages in each chunk

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B-2 pages in each chunk

- 1. External Sort R
- 2. External Sort S
- 3. Merge R and S

[R] = number of pages in Table R

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$$4 * ([R] + [S]) + [R] + [S]$$
  
=  $5 * ([R] + [S])$ 

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- Internal Sort R
   (pass 0 only)
- 2. Internal Sort S (pass 0 only)
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$$2 * ([R] + [S]) + [R] + [S]$$
  
=  $3 * ([R] + [S])$ 

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# Join Algorithms Worksheet #1

1. Describe when you would want to use a chunk nested loop join, a sort-merge join and a hash join:

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#### Chunk Nested Loop Join:

- Cartesian Product
- Non-equality Predicate

#### Sort Merge Join:

- Sorting
- Skewed Data
- Limited Memory

#### Hash Join:

- Hashing
- Uneven Relations
- Hybrid Hashing

## Join Algorithms Worksheet #2

[R] = number of pages in Table R

 $p_R$  = number of records per page of R

|R| = number of records in R

2. We have 15 pages of memory, and we want to join two tables [R] and [S] where [R] is 100 pages and [S] is 50 pages. [R] holds 100 tuples per page and [S] holds 50 tuples per page.

How many disk reads are needed to perform a Simple Nested Loops join?

2. We have 15 pages of memory, and we want to join two tables [R] and [S] where [R] is 100 pages and [S] is 50 pages. [R] holds 100 tuples per page and [S] holds 50 tuples per page.

How many disk reads are needed to perform a Simple Nested Loops join?

Using S as the outer relation yields the lowest I/O count.  $p_S^*[R]^*[S] + [S] = 50*100*50 + 50 = 250050$  I/O's

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How about a Sort Merge Join? (Utilize the optimized version)

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$$3([R] + [S]) = 3*150 = 450 I/O's$$

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How about a Hash Join? (Assume no recursive partitioning and ignore output costs)

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How about a Hash Join? (Assume no recursive partitioning and ignore output costs)

Partitioning Phase: 2([R]+[S])

Matching Phase: [R]+[S]

Total = 3([R] + [S]) = 3\*150 = 450 I/O's