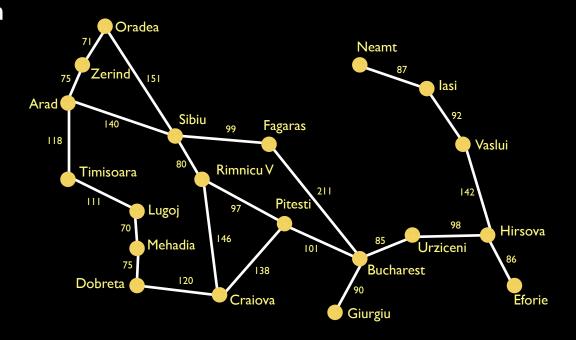
Uninformed Search

Traveling through Romania

- currently in Arad
- flight leaves tomorrow from Bucharest
- goal: be in Bucharest
- problem:
 - states = various cities
 - actions = drive between cities
- solve:
 - find the sequence of cities to drive through: Arad, Sibiu, Fagaras, Bucharest

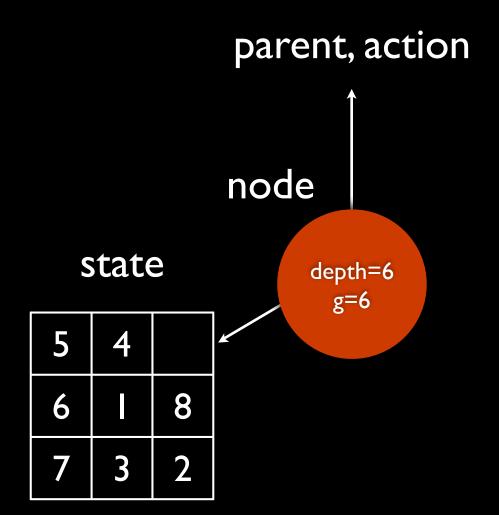


Stating the Problem

- initial state: at Arad
- successor function: succ(state) = { (action, state, cost), ... }
- goal test: goalp(state) = True if state is goal
- path cost: sum of step costs c(x,a,y) = g(n)
- solution: sequence of actions from initial to goal state

States vs Nodes

- state: representation of a problem configuration
- node: data structure in search graph including
 - state
 - parent node
 - action
 - path cost -g(n)
 - depth



```
(defun make-node (&key state (parent nil)
                              (action nil)
                              (path-cost 0) (depth 0))
 (list state parent action path-cost depth))
(defun node-state (node) (car node))
(defun node-parent (node) (cadr node))
(defun node-action (node) (caddr node))
(defun node-path-cost (node) (cadddr node))
(defun node-depth (node) (car (cddddr node)))
OR
(defstruct node state
           (parent nil) (action nil) (path-cost 0) (depth 0))
```

Search Strategies

- defined by picking the order of node expansion
- strategies evaluated by:
 - completeness: does it always find a solution if one exists?
 - time complexity: number of nodes generated
 - space complexity: maximum number of nodes in memory
 - optimality: does it always find a least cost solution?
- time and space complexity are measured in terms of:
 - b: maximum branching factor of the search tree (graph)
 - d: depth of least-cost solution
 - m: maximum depth of the state space (may be ∞)

Uninformed Search

- breadth first
- uniform cost
- depth first
- depth limited
- iterative deepening

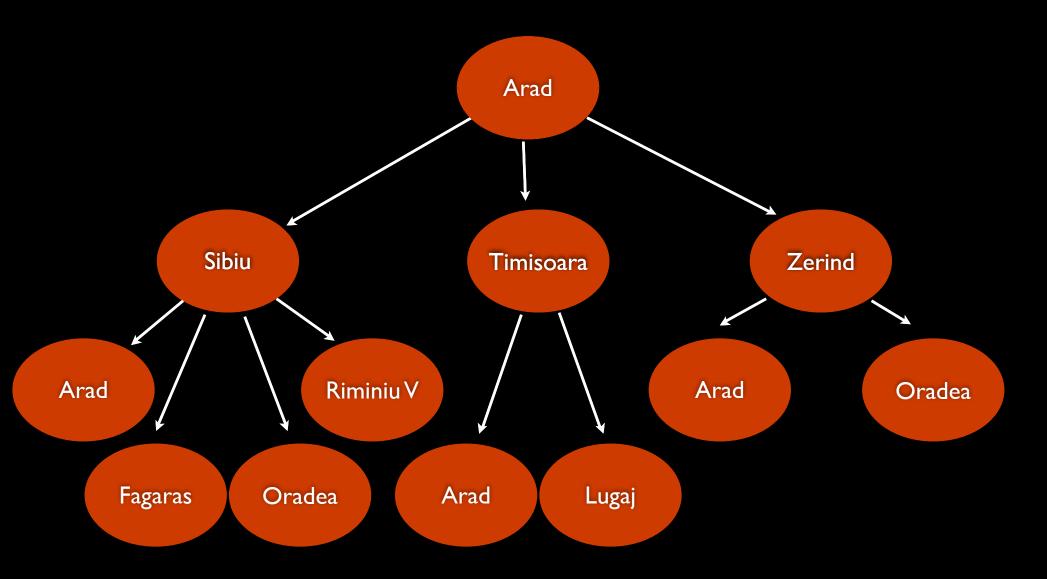
```
(defun expand (successor node)
  (let ((triples (funcall successor (node-state node))))
    (mapcar (lambda (action-state-cost)
              (let ((action (car action-state-cost))
                    (state (cadr action-state-cost))
                    (cost (caddr action-state-cost)))
                (make-node :state state
                           :parent node
                            :action action
                            :path-cost (+ (node-path-cost node)
                                          cost)
                            :depth (1+ (node-depth node)))
            triples)
    ))
```

```
(defstruct node
  (state nil)
  (parent nil)
  (action nil)
  (path-cost 0)
  (depth 0))
; implicitly defines
 (defun make-node (&key (state nil) (parent nil) (action nil)
                         (path-cost 0) (depth 0))...
  (defun node-state (node)...
 (defun node-parent (node)...
 (defun node-action (node)...
 (defun node-path-cost (node)...
; (defun node-depth (node)...
 (defun node-p (x)...
; also, accessors can be used with setf
 (setf (node-path-cost n) 15)
```

Breadth-First Search

- fringe is FIFO queue
- complete (if b is finite)
- time complexity = $O(b^{d+1})$
- space complexity = $O(b^{d+1})$ keeps all nodes in memory
- optimal if unit cost

Searching the Tree (Graph)



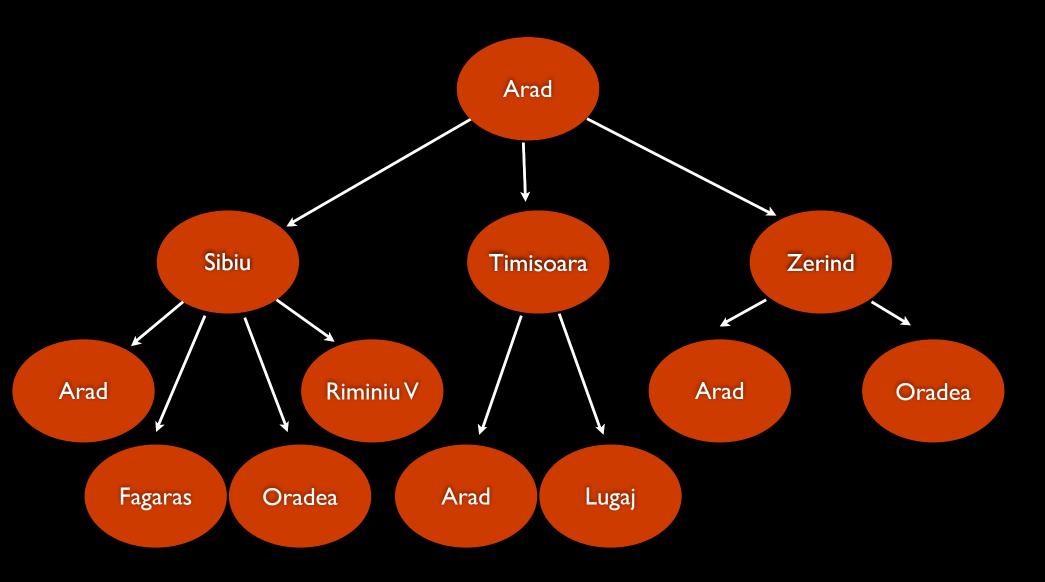
Uniform-Cost Search

- fringe = $\overline{\text{queue ordered by path cost, g(n)}}$
- equivalent to breadth-first if step costs are all equal
- complete if step costs are non-negative
- time & space complexity: similar to breadth-first
- optimal because nodes are expanded in increasing order of path cost, g(n)

Depth First Search

- fringe is LIFO queue (stack)
- incomplete fails in infinite-depth spaces or spaces with loops
 - modify code to avoid repeated states
 - complete in finite spaces
- time complexity is O(b^m) terrible if m >> b, but often much faster than breadth-first
- space complexity is O(bm) linear!!!
- optimal no. why?

Searching the Tree (Graph)



Other Uniformed Searches

- depth-limited is depth-first with a cutoff (nodes at maximum depth have no successors)
- iterative deepening search is iterative calls to depth-limited search, increasing depth cutoff each time
- what's the overhead?
- is it complete? what are the time and space complexities? is it optimal?

Eliminating Repeated States

- failure to detect repeated states turns linear problem into exponential one
- fixing this turns tree search into graph search

```
(defun graph-search (fringe closed successor goalp samep)
  (unless (q-emptyp fringe)
    (let ((node (q-remove fringe)))
      (cond ((funcall goalp (node-state node))
             (action-sequence node))
            ((member (node-state node) closed
                     :test samep :key #'node-state)
             (graph-search fringe closed
                           successor goalp samep))
            (t (let ((successors (expand successor node)))
                 (graph-search (q-insert fringe successors)
                                (cons node closed)
                               successor goalp samep)))
            ))
    ))
```