Last time:  $f(x): \mathbb{R}^n \to \mathbb{R}^m$ 

The condition number  $K(x) := \frac{\|Df(x)\|}{\|f(x)\|}$ 

Today: Apply it to Ax = b

• Motivation:  $A \in \mathbb{R}^{n \times n}$ ,  $b \in \mathbb{R}^n$ , A invertible

Solving Ax = b exactly:  $f(b) = A^{-1}b$ 

A backward stable algorithm satisfies  $f(b) = f(b) = A'(b + \Delta b)$ with 11 2611/116(1 = O(Emach)

To analyze forward error, need to ampute KA  $K_A(b) = \frac{\|Df\| \|b\|}{\|A^{-1}b\|} = \frac{\|A^{-1}\| \|b\|}{\|A^{-1}b\|}$ 

We may want a condition number that is independent of input b,  $K(A) := \sup_{b \in \mathbb{R}^{n}} \frac{\|A^{-1}\| \|b\|}{\|A^{-1}b\|} = \|A^{-1}\| \|A\|$   $= \|A^{-1}\| \|A\| \|A^{-1}b\|$   $= \|A^{-1}\| \|A\| \|A^{-1}b\|$ 

"=" Cattainable
for some be R"

K(A) is called the condition number of A

- · Properties of condition #5 (Assume 11.11 is subordinate norm in the following) 1) K(A) >1, K(A-1) = K(A)
  - 2) If A is normal with eigenvalues 12,13 1213... 3 12190 then  $K_2(A) := ||A||_2 ||A^{-1}||_2 = \frac{||A||}{||A||}$

For general A, we have  $K(A) > \frac{|X_1|}{|X_1|}$ 

Pf: When A is wormal, 
$$A = U^*\Lambda U$$
.

 $IIA II_2 = \sqrt{\lambda_{\text{max}}} (A^*A) = \sqrt{\lambda_{\text{max}}} (U^*\Lambda^2 U) = 1\lambda.1$ 

Similarly  $IIA^{-1}II_2 = |\Lambda_1|^{-1}$ 

For general A, we need the following relation between horm and maximum eigenvalue

I emma: The spectral radius of A, denote by

 $Q(A) = \max\{I\lambda I : \lambda \text{ is an eigenvalue of } A\}$ 

Satisfies  $Q(A) \leq IIAII$ 

Proof of lemma: exercise  $IIAII$ 

Using this lemma, we have  $IIAII \geq I\lambda.1$ ,  $IIA^{-1}II \geq IJ.1I^{-1}$ 

wind the lower bound of A follows  $IIIAII^{-1}$ 

omd the lower bound of A follows

3) condition #s of different norms are equivalent.

In particular,  $\frac{1}{n}K_2(A) \leq K_1(A) \leq n K_2(A)$ 

4) 2- condition # measures the inverse of the relative distance to the newest singular matrix.

Prop. Let  $A \in \mathbb{R}^{n \times n}$ , A not singular, then

 $\frac{1}{K_2(A)} = \min_{K_2(A)} \int_{II} \frac{II K_2II_2}{IIAII_2}$ 

Singular

Pf: exercise.  $III$ 

We have the following thum that tells you how the solution is perturbed given an inaccurate right hand side.

Thm 1 Let Ax = b and  $r = b - A\hat{x}$ 1\_residual then we have  $\frac{1}{K(A)} \frac{||r||}{||b||} \leq \frac{||\hat{x}-x||}{||x||} \leq K(A) \frac{||r||}{||b||}$ Pf: Let bx = x-x 110x 11 = 11A-1711 < 11A-111 11711 use 11 b1 = 11 A x 1 \le 11 A 1 | 11 x | we obtain inequ on the right. The left inegn can be proved similarly - exercise 12 K(A) seems to only quantify the error in octput for a perturbation in b, but it turns out that it also quantify the porturbation in A form of equ we get in bakward Thm let Ax = b,  $(A + \triangle A) \hat{x} = b + \triangle b$ error analysis of many algo for where ILDAII & ELIAH, NDBII & ELIBII solving Ax-b and EK(A) < 1. then  $\frac{\|\mathbf{x} - \hat{\mathbf{x}}\|}{\|\mathbf{x}\|} \leq \frac{\mathcal{K}(\mathbf{A})}{1 - \mathcal{K}(\mathbf{A}) \cdot \|\Delta\mathbf{A}\|/\|\mathbf{A}\|} \left(\frac{\|\Delta\mathbf{A}\|}{\|\mathbf{A}\|} + \frac{\|\Delta\mathbf{b}\|}{\|\mathbf{b}\|}\right)$ = O(K(A) E) for  $E << \frac{1}{K(A)}$  $Pf: \hat{\chi} = (A + \Delta A)^{-1}(b + \Delta b)$  $x = (A + \Delta A)^{-1}(A + \Delta A) \times = (A + \Delta A)^{-1}(b + (\Delta A)x)$  $\hat{x} - x = (A + \Delta A)^{-1} (\Delta b - (\Delta A)x)$ 

$$= [A(I+A^{-1}AA)]^{-1} (Ab - (AA) \times)$$

$$= (I+A^{-1}AA)^{-1} A^{-1} (Ab - (AA) \times)$$

$$\|(\hat{X}-X)\| \leq \|(I+A^{-1}AA)^{-1}\| \|A^{-1}\| (\|Ab\| + \|AA\| \cdot \|X\|)$$

$$\|(X-X)\| = \|(I+A^{-1}AA)^{-1}\| K(A) (\frac{\|Ab\|}{\|A\| \cdot \|X\|} + \frac{\|AA\|}{\|A\|})$$

$$\leq \|(I+A^{-1}AA)^{-1}\| K(A) (\frac{\|Ab\|}{\|A\| \cdot \|X\|} + \frac{\|AA\|}{\|A\|}) (\#)$$

$$|Amma: If NB || < 1, then I+B is invertible$$

$$|Amma: If NB || < 1, then I+B is invertible$$

$$|Amma: If NB || < 1, then I+B is invertible$$

$$|Amma: If NB || < 1, then I+B is invertible$$

$$|Amma: If NB || < 1, then I+B is invertible$$

$$|Amma: If NB || < 1, then I+B is invertible$$

$$|Amma: If NB || < 1, then I+B is invertible$$

$$|Amma: If NB || < 1, then I+B is invertible$$

$$|Amma: If NB || < 1, then I+B is invertible$$

$$|Amma: If NB || < 1, then I+B is invertible$$

$$|Amma: If NB || < 1, then I+B is invertible$$

$$|Amma: If NB || < 1, then I+B is invertible$$

$$|Amma: If NB || < 1, then I+B is invertible$$

$$|Amma: If NB || < 1, then I+B is invertible$$

$$|Amma: If NB || < 1, then I+B is invertible$$

$$|Amma: If NB || < 1, then I+B is invertible$$

$$|Amma: If NB || < 1, then I+B is invertible$$

$$|Amma: If NB || < 1, then I+B is invertible$$

$$|Amma: If NB || < 1, then I+B is invertible$$

$$|Amma: If NB || < 1, then I+B is invertible$$

$$|Amma: If NB || < 1, then I+B is invertible$$

$$|Amma: If NB || < 1, then I+B is invertible$$

$$|Amma: If NB || < 1, then I+B is invertible$$

$$|Amma: If NB || < 1, then I+B is invertible$$

$$|Amma: If NB || < 1, then I+B is invertible$$

$$|Amma: If NB || < 1, then I+B is invertible$$

$$|Amma: If NB || < 1, then I+B is invertible$$

$$|Amma: If NB || < 1, then I+B is invertible$$

$$|Amma: If NB || < 1, then I+B is invertible$$

$$|Amma: If NB || < 1, then I+B is invertible$$

$$|Amma: If NB || < 1, then I+B is invertible$$

$$|Amma: If NB || < 1, then I+B is invertible$$

$$|Amma: If NB || < 1, then I+B is invertible$$

$$|Amma: If NB || < 1, then I+B is invertible$$

$$|Amma: If NB || < 1, then I+B is invertible$$

$$|Amma: If NB || < 1, then I+B is invertible$$

$$|Amma: If NB || < 1, then I+B is invertible$$

$$|Amma: If NB || < 1, then I+B is invertible$$

$$|Amma: I$$

Heview: unitary matrix Q\*Q=I, Q&C"x" Q is unitary  $\iff$   $q_i^*q_j = \delta_{ij}$ ,  $i,j = 1, \dots, n$ i.e. 1 q...., qn3 forms an orthonormal basis Some properties: Let XERM, AECMAN 11 Q X 112 = 11x112 11 QAII2 = 11AQ ||2 = 11AII2 | norm 11Q All= 11AQ11= 11A11F  $Pf: \|Qx\|_2^2 = x^*Q^*Qx = x^*x = \|x\|_2^2$  $||QA||_{2}^{2} = \sqrt{\lambda_{max}(A^{*}Q^{*}QA)} = \sqrt{\lambda_{max}(A^{*}A)} = ||A||_{2}^{2}$  $||QA||_F^2 = \sqrt{tr(A^*Q^*QA)} = \sqrt{tr(A^*A)} = ||A||_F^2$ Other equation can be proved similarly - exercise & In Linear algebra course, we know when  $A \in \mathbb{C}^{n \times n}$  is normal, then A is diagonalizable, i.e., can be factored as  $A = U \Lambda U^*$ where  $U \in \mathbb{C}^{n \times n}$  is unitary,  $\Lambda = \text{cliag}(\lambda_1, \dots, \lambda_n)$ · Q: A is not normal? A ∈ Cmxn? · Singular value Decomposition A E C Mxn SVD: A = U I V\*

where 
$$U = \begin{bmatrix} u_1 & \cdots & u_m \end{bmatrix} \in \mathbb{C}^{m \times m}$$
, us loft singular vector  $\mathbb{Z}^{m} = [v_1 & \cdots & v_n] \in \mathbb{C}^{m \times m}$ , us loft singular vector  $\mathbb{Z}^{m} = [v_1 & \cdots & v_n] \in \mathbb{C}^{m \times m}$ ,  $v_2 = [v_1 & \cdots & v_n] \in \mathbb{C}^{m \times m}$ ,  $v_3 = [v_1 & \cdots & v_n] \in \mathbb{C}^{m \times m}$ ,  $v_4 = [v_1 & \cdots & v_n] \in \mathbb{C}^{m \times m}$ ,  $v_4 = [v_1 & \cdots & v_n] \in \mathbb{C}^{m \times m}$ ,  $v_4 = [v_1 & \cdots & v_n] \in \mathbb{C}^{m \times m}$ .

Existence of  $V = [v_1 & \cdots & v_n] \in \mathbb{C}^{m \times m}$ ,  $v_1 = [v_1 & \cdots & v_n] \in \mathbb{C}^{m \times m}$ .

Existence of  $V = [v_1 & \cdots & v_n] \in \mathbb{C}^{m \times m}$ ,  $v_1 = [v_1 & \cdots & v_n] \in \mathbb{C}^{m \times m}$ .

I.e.  $A^*A = V = [v_1 & \cdots & v_n] \in \mathbb{C}^{m \times m}$ .

I.e.  $A^*A = V = [v_1 & \cdots & v_n] \in \mathbb{C}^{m \times m}$ .

I.e.  $A^*A = [v_1 & \cdots & v_n] \in \mathbb{C}^{m \times m}$ .

I.e.  $A^*A = [v_1 & \cdots & v_n] \in \mathbb{C}^{m \times m}$ .

I.e.  $A^*A = [v_1 & \cdots & v_n] \in \mathbb{C}^{m \times m}$ .

I.e.  $A^*A = [v_1 & \cdots & v_n] \in \mathbb{C}^{m \times m}$ .

I.e.  $A^*A = [v_1 & \cdots & v_n] \in \mathbb{C}^{m \times m}$ .

I.e.  $A^*A = [v_1 & \cdots & v_n] \in \mathbb{C}^{m \times m}$ .

I.e.  $A^*A = [v_1 & \cdots & v_n] \in \mathbb{C}^{m \times m}$ .

I.e.  $A^*A = [v_1 & \cdots & v_n] \in \mathbb{C}^{m \times m}$ .

I.e.  $A^*A = [v_1 & \cdots & v_n] \in \mathbb{C}^{m \times m}$ .

I.e.  $A^*A = [v_1 & \cdots & v_n] \in \mathbb{C}^{m \times m}$ .

I.e.  $A^*A = [v_1 & \cdots & v_n] \in \mathbb{C}^{m \times m}$ .

I.e.  $A^*A = [v_1 & \cdots & v_n] \in \mathbb{C}^{m \times m}$ .

I.e.  $A^*A = [v_1 & \cdots & v_n] \in \mathbb{C}^{m \times m}$ .

I.e.  $A^*A = [v_1 & \cdots & v_n] \in \mathbb{C}^{m \times m}$ .

I.e.  $A^*A = [v_1 & \cdots & v_n] \in \mathbb{C}^{m \times m}$ .

I.e.  $A^*A = [v_1 & \cdots & v_n] \in \mathbb{C}^{m \times m}$ .

I.e.  $A^*A = [v_1 & \cdots & v_n] \in \mathbb{C}^{m \times m}$ .

I.e.  $A^*A = [v_1 & \cdots & v_n] \in \mathbb{C}^{m \times m}$ .

I.e.  $A^*A = [v_1 & \cdots & v_n] \in \mathbb{C}^{m \times m}$ .

I.e.  $A^*A = [v_1 & \cdots & v_n] \in \mathbb{C}^{m \times m}$ .

I.e.  $A^*A = [v_1 & \cdots & v_n] \in \mathbb{C}^{m \times m}$ .

I.e.  $A^*A = [v_1 & \cdots & v_n] \in \mathbb{C}^{m \times m}$ .

I.e.  $A^*A = [v_1 & \cdots & v_n] \in \mathbb{C}^{m \times m}$ .

I.e.  $A^*A = [v_1 & \cdots & v_n] \in \mathbb{C}^{m \times m}$ .

I.e.  $A^*A = [v_1 & \cdots & v_n] \in \mathbb{C}^{m \times m}$ .

I.e.  $A^*A = [v_1 & \cdots & v_n] \in \mathbb{C}^{m \times m}$ .

I.e.  $A^*A = [v_1 & \cdots & v_n] \in \mathbb{C}^{m \times m}$ .

I.e.  $A^*A = [v_1 & \cdots & v_n] \in \mathbb{C}^{m \times m}$ .

I.e.  $A^*A = [v_1 & \cdots & v_n] \in \mathbb{C}^{m \times m}$ .

I.e.  $A^*A$ 

Facts about SVD;

- 1) r := rank(A) = # non-zero sv's 6:
- 2) range  $(A) = span \{u, \dots, u_r\}$   $null (A) = span \{vrn, \dots, v_n\}$
- 3) I(Allz = 61,

When A & C "x" invertible, then

 $K_2(A) = 61/6n$ 

Pf:  $||A||_2 = ||U\Sigma V^*||_2 = ||Z||_2 = \sqrt{||\Delta_{max}(Z^*\Sigma)||} = 6$ Similarly  $||A^{-1}||_2 = 6$  and hence  $K_2(A) = 6$  on

$$A = \sum_{i=1}^{r} 6_i u_i v_i^*,$$

Application:

Best rank K approximation:  $A_{k:=} \sum_{i=1}^{k} s_i u_i v_i^*$  (k << m.n)

min  $||A - B||_F = ||A - A_{k}||_F = \sqrt{\sum_{i > k} S_i^2}$  (1)

rank(B)  $\leq r$ 

min  $||A - B||_2 = ||A - A \times ||_2 = 6 \times 1$  (2)  $||A - B||_2 = ||A - A \times ||_2 = 6 \times 1$ 

(Eckart-Young, 1936)

Application: Regularizing ill-conditioned problems

Consider solving the following problem

$$Ax = b$$
 with  $A \in \mathbb{R}^{n \times n}$ ,  $\lambda_i(A) > 0$ 

$$K_2(A) = \frac{\lambda_{max}(A)}{\lambda_{min}(A)} > 1$$

Since the problem is ill-conditioned, any algorithm that solves the problem is difficult.

How do we decrease the condition number?

(regularize the problem)

We may solve instead

$$(A + \alpha I) \hat{x} = b \qquad (a > 0)$$

$$K(A) = \frac{\lambda_{\max}(A) + \alpha}{\lambda_{\min}(A) + \alpha} < K_2(A)$$

How do we choose a?

| 11x1 |
| 11x1 |
| 10x1 |

Since  

$$x - \hat{x} = (A + \alpha I)^{-1} (A + \alpha I - A) A^{-1} b$$

$$= \alpha (A + \alpha I)^{-1} x$$

we have  $||x-\hat{x}||_2 \le \alpha ||(A+\alpha I)^{-1}||_2 ||x||_2 = \frac{\alpha}{\lambda_{min}(A)+\alpha} ||x||_2$ we hope to balance  $\frac{||x-\hat{x}||_2}{||x||_2} \sim ||x|(A+\alpha I)||x||_2$ 

that is 
$$\frac{\alpha}{\lambda \min(A) + \alpha} \sim \frac{\lambda \max(A) + \alpha}{\lambda \min(A) + \alpha} \sim \frac{\lambda \max(A) + \alpha}{\lambda \min(A) + \alpha}$$

that is a ~ lmax (A) Emach