

# Graph Theory

## An Introduction

[https://github.com/mohawk2/  
presentations](https://github.com/mohawk2/presentations)

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# Introduction

- Things
- Connections between things
- Not charts



# Introduction

- Why
- What
- How (in computing)
- Algorithms (why bought ticket)
- Resources
- Exercises



# Why

- Network routing (paths)
- Social media
- Machine learning
- Genetics
- Software build systems such as make (example of hypergraph)



# What

- Directed
- Undirected
- Hypergraphs (both above)



# UNEXPECTED MATHS

## Graph [\[ edit \]](#)

In one restricted but very common sense of the term,<sup>[1][2]</sup> a **graph** is an **ordered pair**  $G = (V, E)$  comprising:

- $V$ , a **set** of **vertices** (also called **nodes** or **points**);
- $E \subseteq \{\{x, y\} \mid x, y \in V \text{ and } x \neq y\}$ , a **set** of **edges** (also called **links** or **lines**), which are **unordered pairs** of vertices (that is, an edge is associated with two distinct vertices).

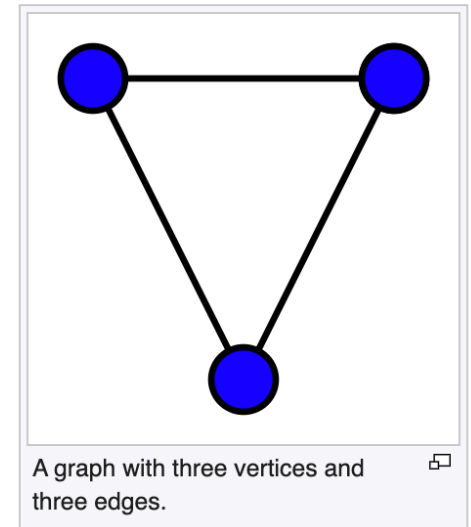
To avoid ambiguity, this type of object may be called precisely an **undirected simple graph**.

In the edge  $\{x, y\}$ , the vertices  $x$  and  $y$  are called the **endpoints** of the edge. The edge is said to **join**  $x$  and  $y$  and to be **incident** on  $x$  and on  $y$ . A vertex may exist in a graph and not belong to an edge. **Multiple edges**, not allowed under the definition above, are two or more edges that join the same two vertices.

In one more general sense of the term allowing multiple edges,<sup>[3][4]</sup> a **graph** is an ordered triple  $G = (V, E, \phi)$  comprising:

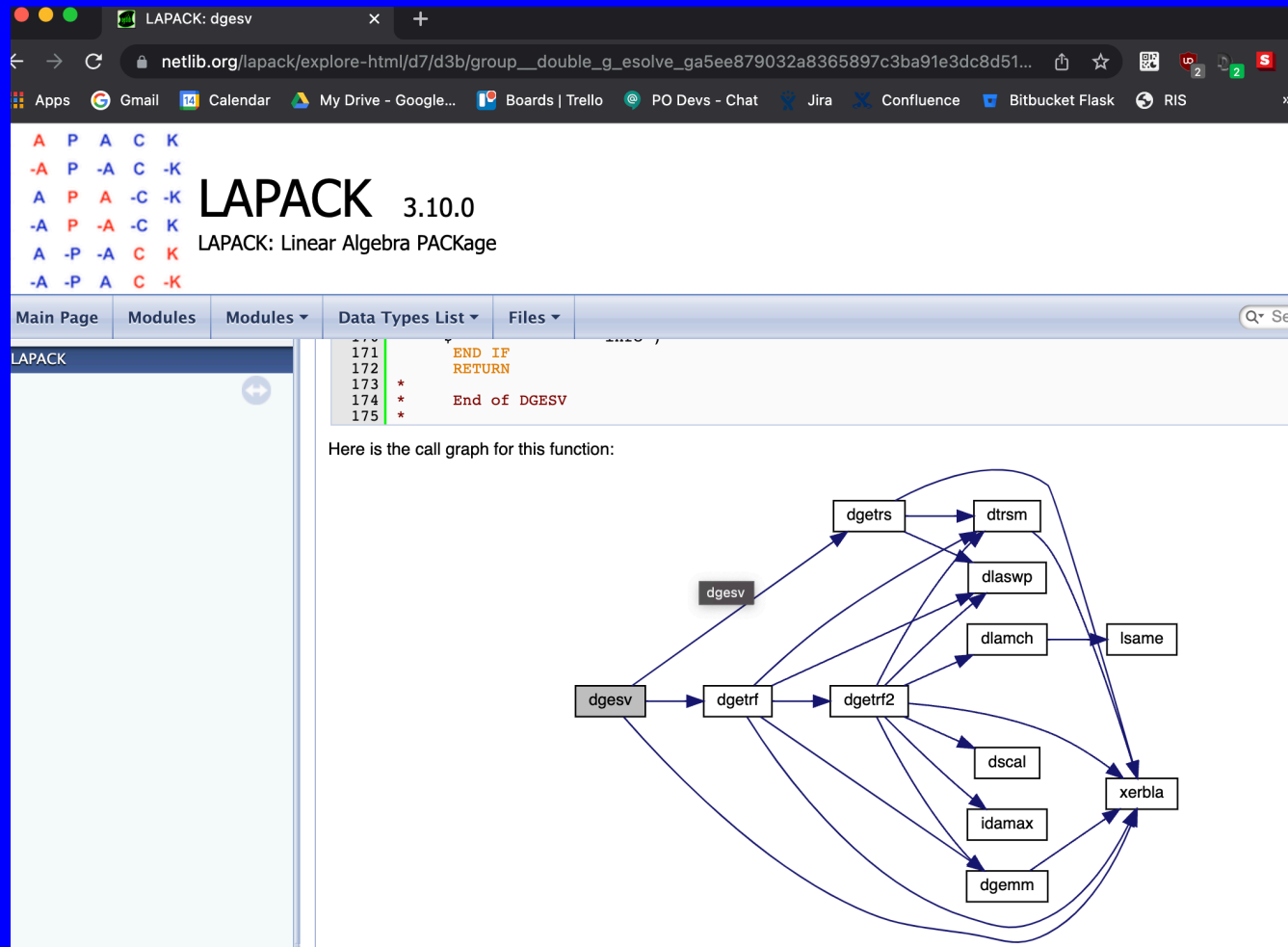
- $V$ , a **set** of **vertices** (also called **nodes** or **points**);
- $E$ , a **set** of **edges** (also called **links** or **lines**);
- $\phi : E \rightarrow \{\{x, y\} \mid x, y \in V \text{ and } x \neq y\}$ , an **incidence function** mapping every edge to an **unordered pair** of vertices (that is, an edge is associated with two distinct vertices).

To avoid ambiguity, this type of object may be called precisely an **undirected multigraph**.



[https://en.wikipedia.org/wiki/Graph\\_theory](https://en.wikipedia.org/wiki/Graph_theory)

# Directed



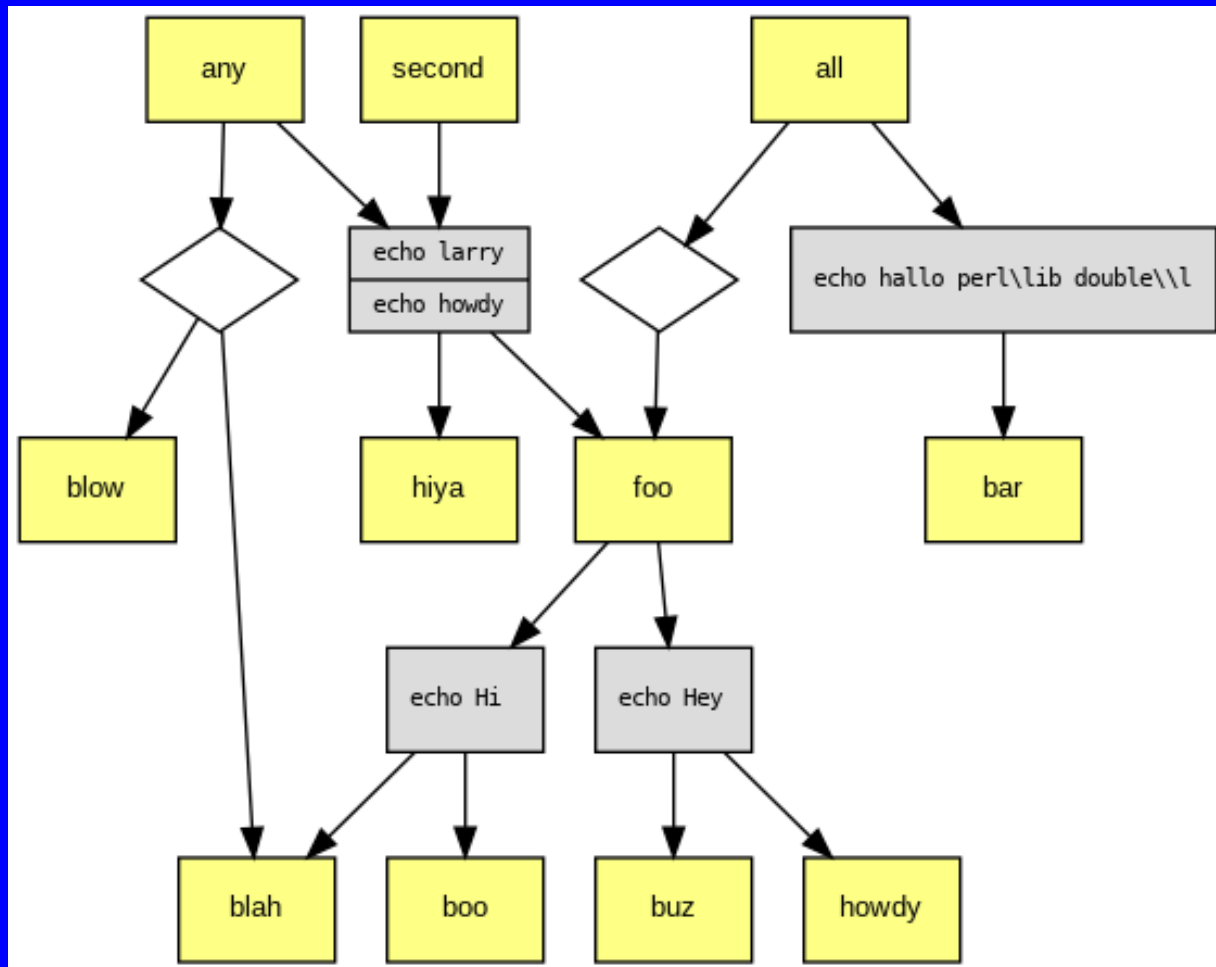
[https://www.netlib.org/lapack/explore-html/d7/d3b/group\\_\\_double\\_g\\_solve\\_ga5ee879032a8365897c3ba91e3dc8d512.html](https://www.netlib.org/lapack/explore-html/d7/d3b/group__double_g_solve_ga5ee879032a8365897c3ba91e3dc8d512.html)

# Undirected





# Directed hypergraph



<https://gitlab.com/graphviz/graphviz/-/issues/1911>

# Graph classification

- Directed? Cyclic? Multi? Hyper?
- Social media
- Transport networks
- Software build systems

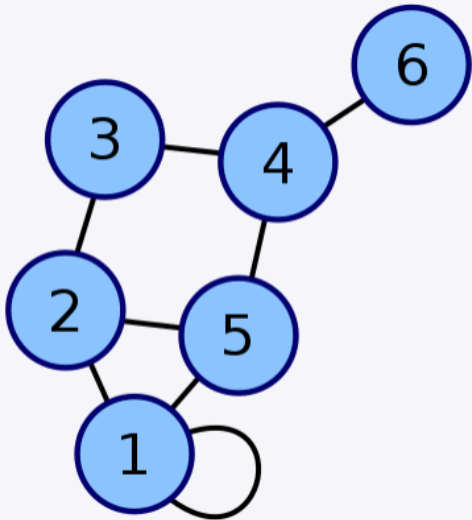


# How

- Matrix: adjacency (undirected = symmetrical)
- Matrix, sparse/hypergraph: incidence
- Linked lists of V, E
- Dictionary of E



# Adjacency matrix

Labeled graph	Adjacency matrix
	$\begin{pmatrix} 2 & 1 & 0 & 0 & 1 & 0 \\ 1 & 0 & 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 & 1 & 1 \\ 1 & 1 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 \end{pmatrix}$ <p>Coordinates are 1–6.</p>

[https://en.wikipedia.org/wiki/Adjacency\\_matrix](https://en.wikipedia.org/wiki/Adjacency_matrix)

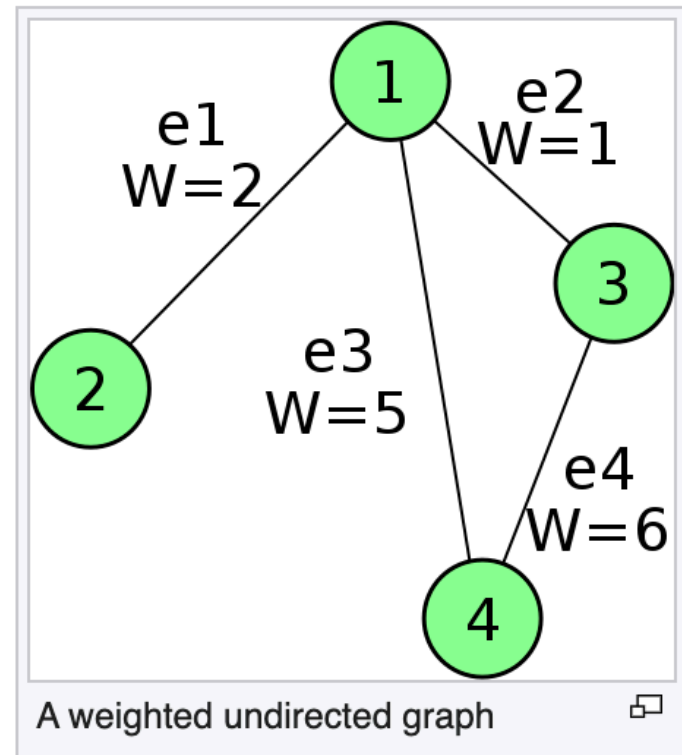
# Incidence matrix

## Weighted graphs [edit]

A weighted graph can be represented using the weight of the edge in place of a 1. For example, the incidence matrix of the graph to the right is:

	$e_1$	$e_2$	$e_3$	$e_4$
1	2	1	5	0
2	2	0	0	0
3	0	1	0	6
4	0	0	5	6

$$= \begin{bmatrix} 2 & 1 & 5 & 0 \\ 2 & 0 & 0 & 0 \\ 0 & 1 & 0 & 6 \\ 0 & 0 & 5 & 6 \end{bmatrix}.$$



[https://en.wikipedia.org/wiki/Incidence\\_matrix](https://en.wikipedia.org/wiki/Incidence_matrix)

## Graph as dictionary (pseudocode)

```
class Graph:
    nodes = []
    edges = {}
    def add_node(v):
        nodes.push(v)
    def add_edge(v1, v2):
        if not v1 in node:
            throw("invalid node {v1}")
        if v2 in nodes:
            throw("invalid node {v2}")
        edges[v1][v2] = 1
```

# Algorithms

- Floyd-Warshall
- Topological sorting



# Floyd-Warshall

```
let dist be a  $|V| \times |V|$  array of minimum distances
    initialized to  $\infty$  (infinity)
for each edge  $(u, v)$  do
     $\text{dist}[u][v] \leftarrow w(u, v)$  // The weight of the edge  $(u, v)$ 
for each vertex  $v$  do
     $\text{dist}[v][v] \leftarrow 0$ 
for  $k$  from 1 to  $|V|$ 
    for  $i$  from 1 to  $|V|$ 
        for  $j$  from 1 to  $|V|$ 
            if  $\text{dist}[i][j] > \text{dist}[i][k] + \text{dist}[k][j]$ 
                 $\text{dist}[i][j] \leftarrow \text{dist}[i][k] + \text{dist}[k][j]$ 
            end if
        end if
    end for
end for
```

[https://en.wikipedia.org/wiki/Floyd%E2%80%93Warshall\\_algorithm](https://en.wikipedia.org/wiki/Floyd%E2%80%93Warshall_algorithm)

[https://metacpan.org/pod/PDL::MATLAB#COMPARISON:-  
FLOYD-WARSHALL-ALGORITHM](https://metacpan.org/pod/PDL::MATLAB#COMPARISON:-FLOYD-WARSHALL-ALGORITHM)



# Topological sorting

```
 $L \leftarrow$  Empty list that will contain the sorted elements  
 $S \leftarrow$  Set of all nodes with no incoming edge
```

```
while  $S$  is not empty do  
    remove a node  $n$  from  $S$   
    add  $n$  to  $L$   
    for each node  $m$  with an edge  $e$  from  $n$  to  $m$  do  
        remove edge  $e$  from the graph  
        if  $m$  has no other incoming edges then  
            insert  $m$  into  $S$ 
```

```
if graph has edges then  
    return error    (graph has at least one cycle)  
else  
    return  $L$       (a topologically sorted order)
```

[https://en.wikipedia.org/wiki/Topological\\_sorting](https://en.wikipedia.org/wiki/Topological_sorting)

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# Summary

- Graphs: things and connections between things
- What, why, how
- Algorithms

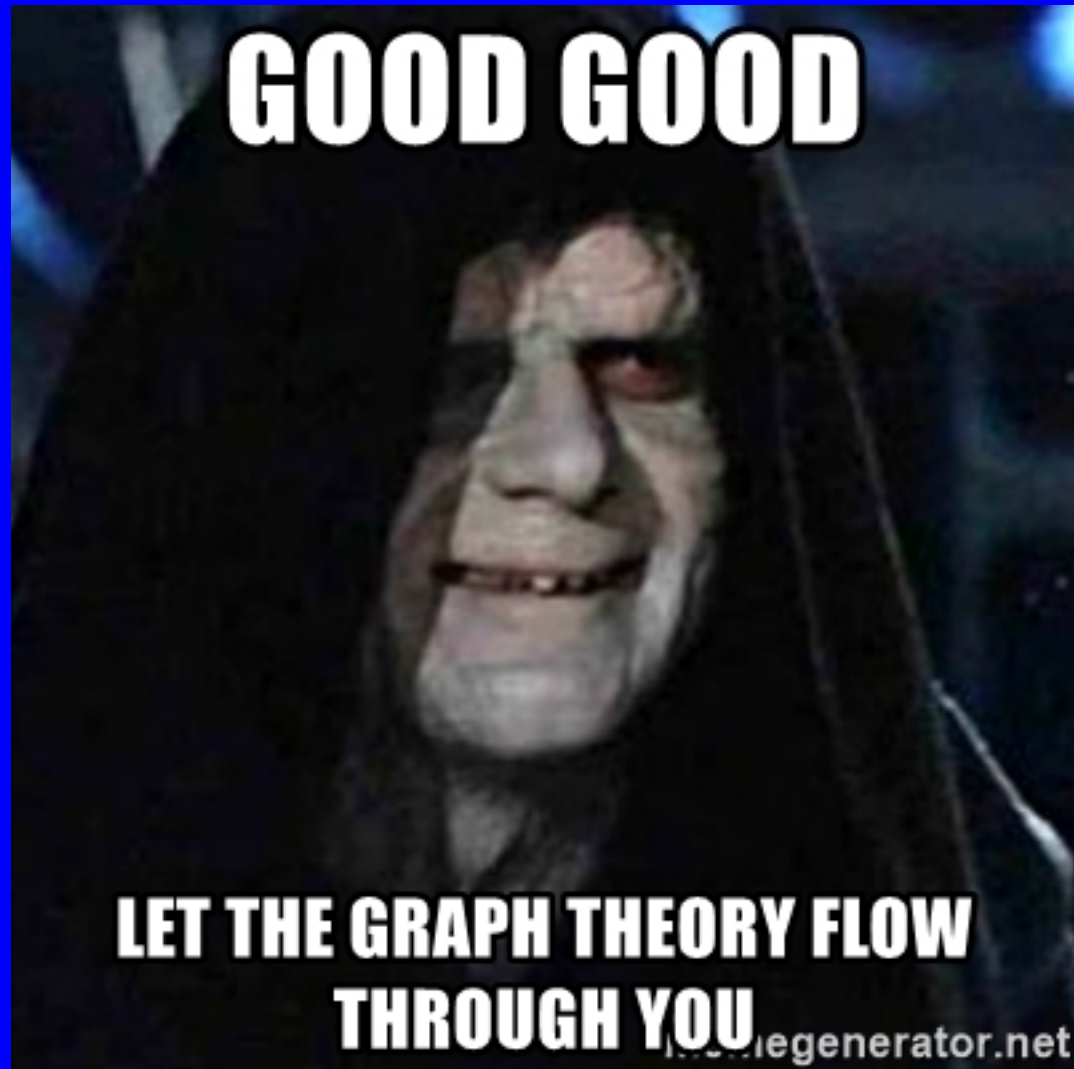


# Resources

- [https://en.wikipedia.org/wiki/Graph\\_theory](https://en.wikipedia.org/wiki/Graph_theory)
- <https://en.wikipedia.org/wiki/Hypergraph>
- <https://graphviz-perl.github.io/>
- <https://github.com/mohawk2/presentations>



# Questions?



# Exercises

- Floyd-Warshall
- Topological sorting

