

# Implementation of the A-star Algorithm

Observations & Analysis

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# Assignment Specifications

- Code A\*; keep it general enough to be able to adapt to any search problem.
- Write modules for open and closed list management. Similarly for parent pointer redirection.
- Verify experimentally the intuition, "better heuristic performs better".
- Verify that "if  $h(n) > h^*(n)$ , for all  $n$ , A\* may find the goal faster, but may discover a suboptimal path".
- Verify that monotone restriction is satisfied, parent pointer redirection for nodes on closed list is not needed.
- Come up with new heuristics for 8-puzzle and missionaries and cannibals; establish their admissibility and monotonicity or otherwise and measure performance.
- Carry out bidirectional A\* search  $S \rightarrow G$  and  $G \rightarrow S$ .

# A\* Pseudo-Code

```
create the open list of nodes, initially containing only our starting node
create the closed list of nodes, initially empty

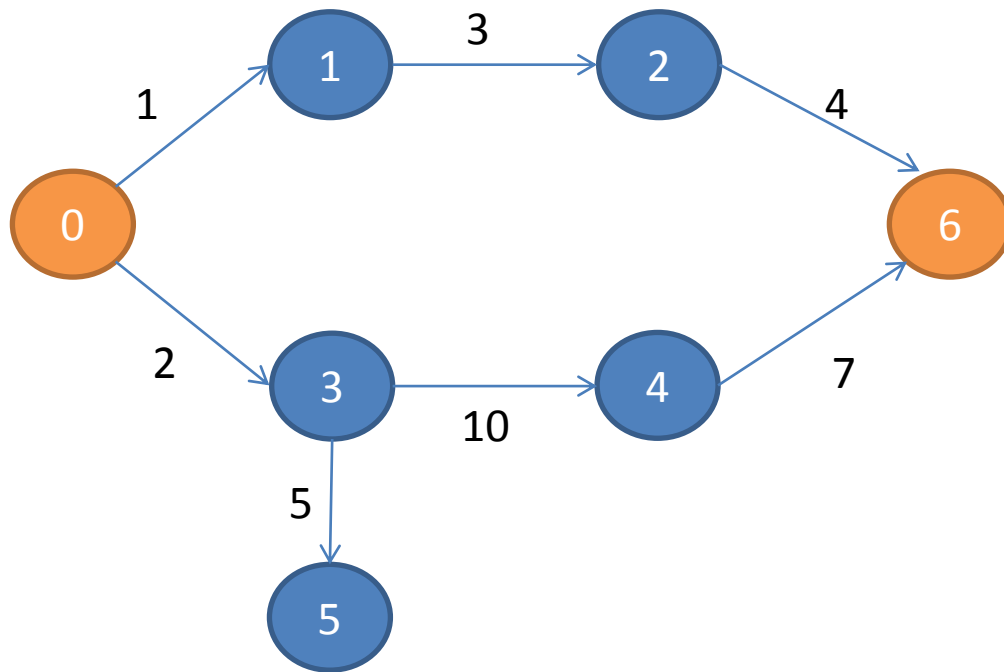
while (we have not reached our goal) {
    consider the best node in the open list (the node with the lowest f value)

    if (this node is the goal) {
        then we're done
    }
    else {
        move the current node to the closed list and consider all of its neighbors

        for (each neighbor) {
            if (this neighbor is in the closed list and our current g value is lower) {
                update the neighbor with the new, lower, g value
                change the neighbor's parent to our current node
            }
            else if (this neighbor is in the open list and our current g value is lower) {
                update the neighbor with the new, lower, g value
                change the neighbor's parent to our current node
            }
            else this neighbor is not in either the open or closed list {
                add the neighbor to the open list and set its g value
            }
        }
    }
}
```

# A-star Algorithm On a Simple Graph

- **Without parent pointer redirection**



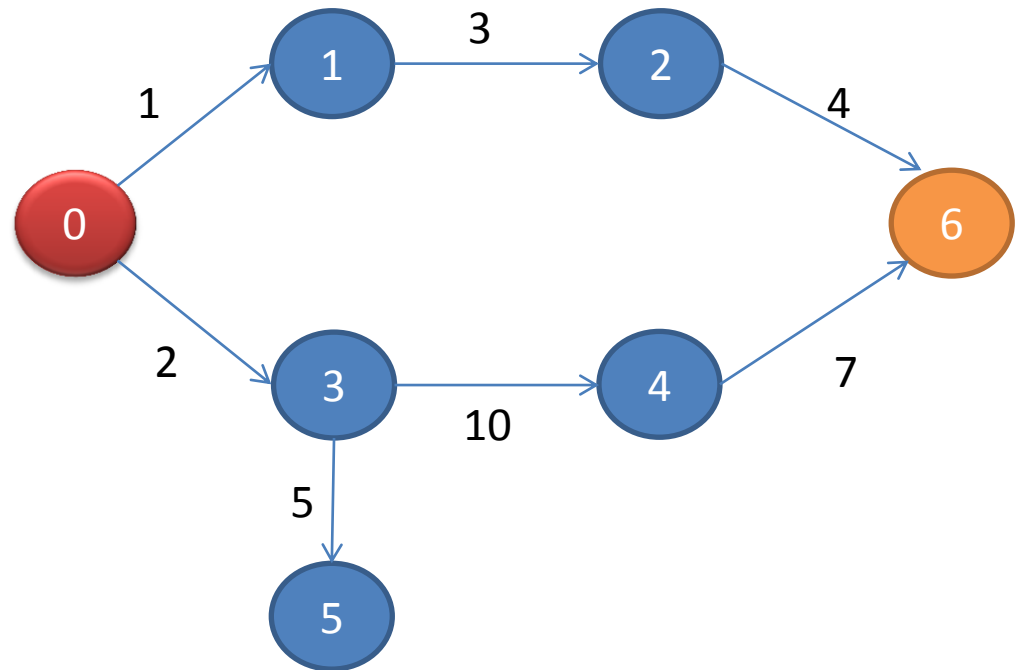
Start Node: 0

Goal Node: 6

$h(i) = 0$  for  $i=1,2,\dots,7$

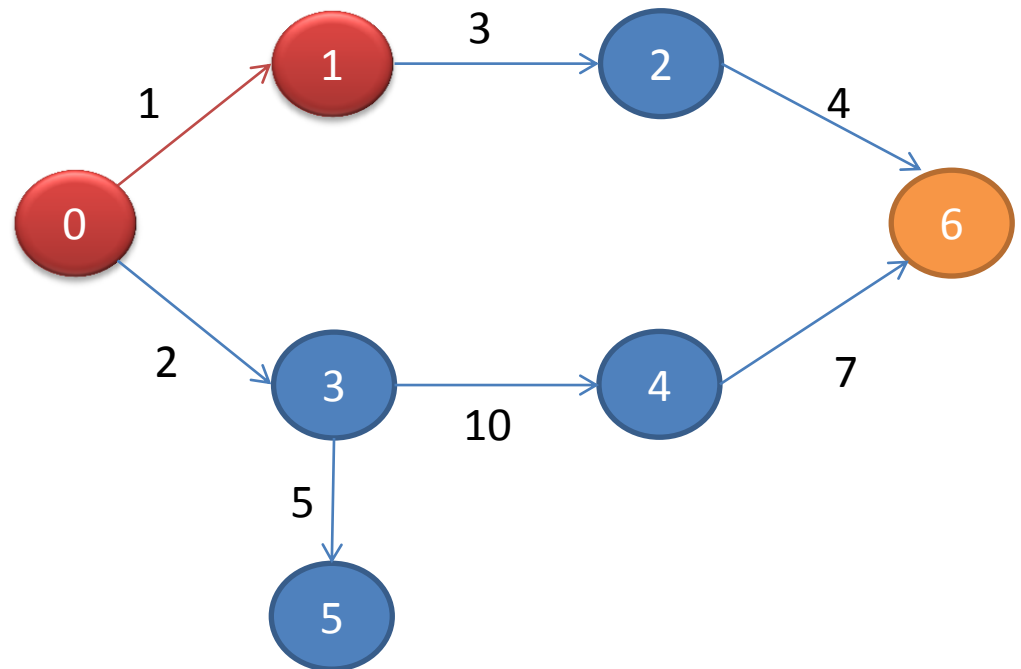
# A-star Algorithm On a Simple Graph

- **Node picked by algorithm from open list:**
  - Iteration 1: Node 0



# A-star Algorithm On a Simple Graph

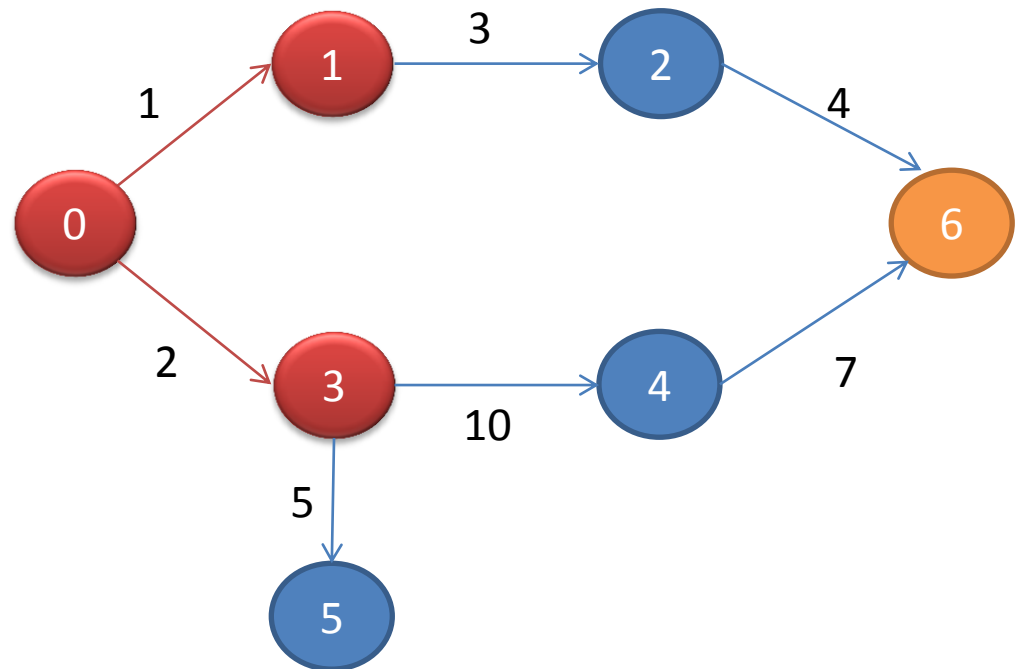
- **Node picked by algorithm from open list:**
  - Iteration 1: Node 0
  - Iteration 2 : Node 1



# A-star Algorithm On a Simple Graph

- **Node picked by algorithm from open list:**

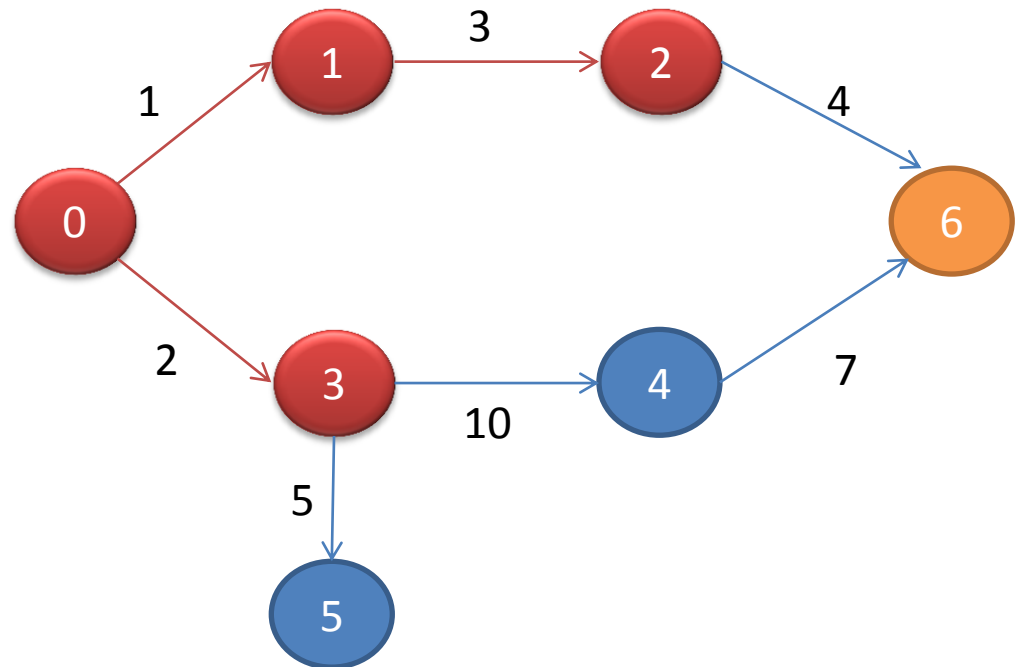
- Iteration 1: Node 0
- Iteration 2 : Node 1
- Iteration 3 : Node 3



# A-star Algorithm On a Simple Graph

- **Node picked by algorithm from open list:**

- Iteration 1: Node 0
- Iteration 2 : Node 1
- Iteration 3 : Node 3
- Iteration 4 : Node 2

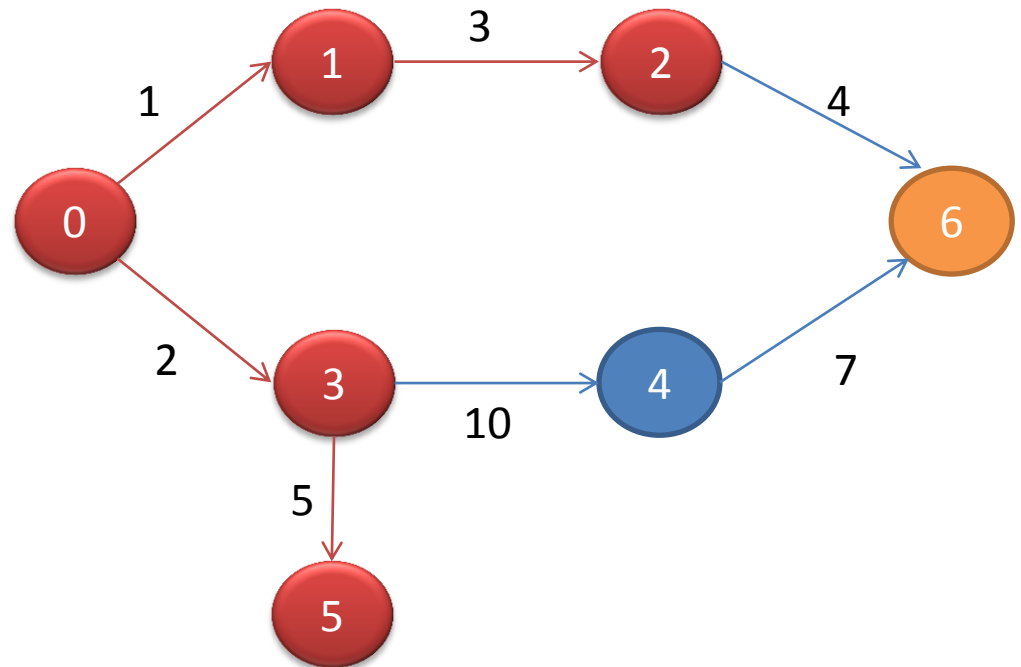




# A-star Algorithm On a Simple Graph

- **Node picked by algorithm from open list:**

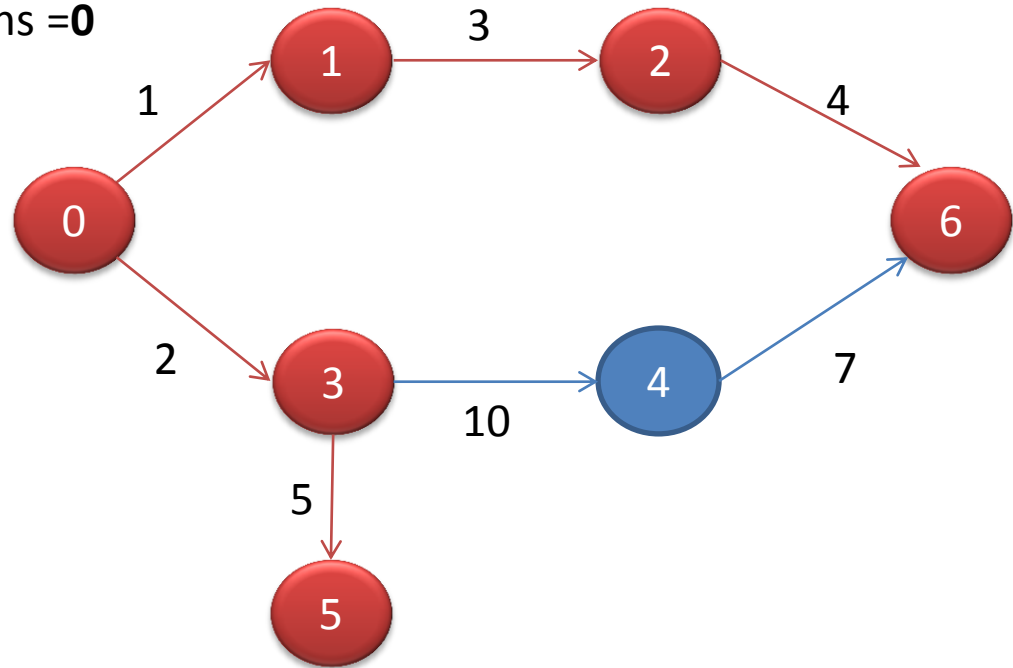
- Iteration 1: Node 0
- Iteration 2 : Node 1
- Iteration 3 : Node 3
- Iteration 4 : Node 2
- Iteration 5 : Node 5



# A-star Algorithm On a Simple Graph

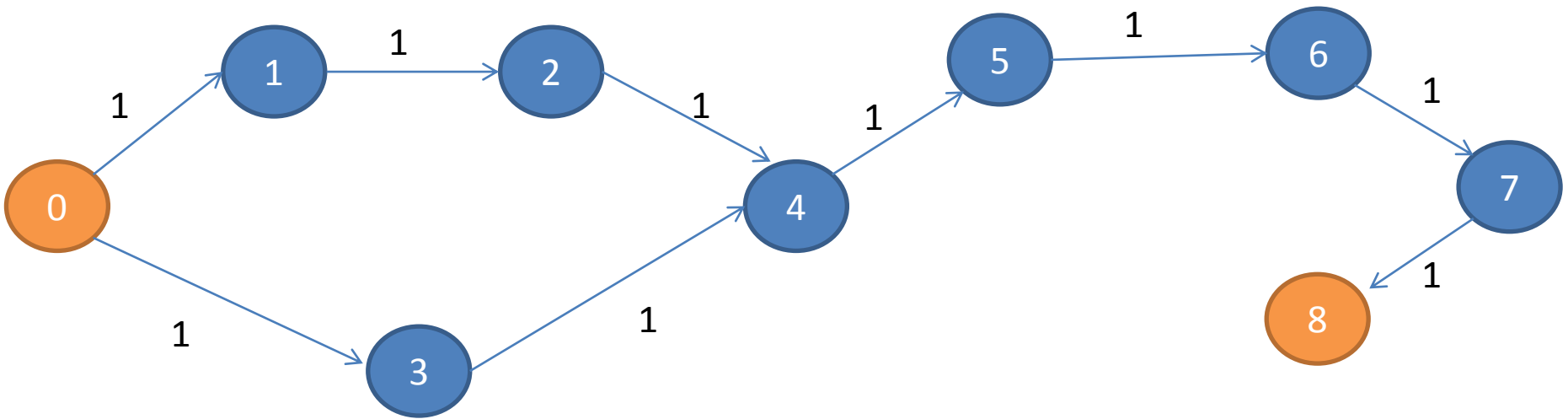
## Final Output

- The optimal path is: **0 1 2 6**
- Optimal path cost is **8**
- Number of iterations taken by the A\* Algo =**5**
- Number of Parent Pointer Redirections =**0**



# Parent Pointer Redirection

- We shall use the following graph & heuristic, to see a case where parent pointer redirection takes place.



Start Node: 0

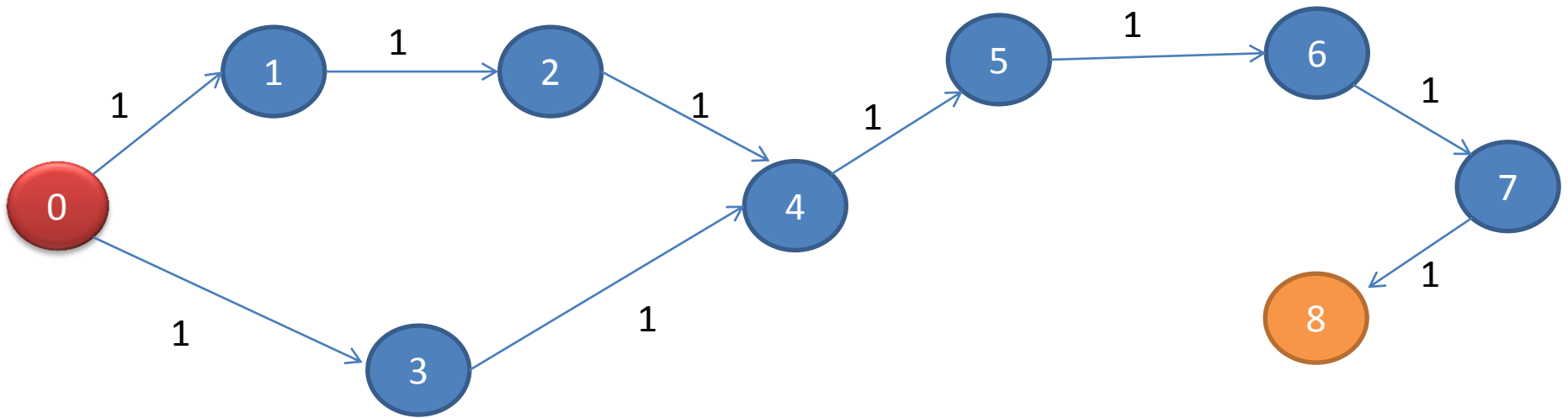
Goal Node: 8

$h(i) = 5$  for  $i=3$

0 otherwise

# Parent Pointer Redirection

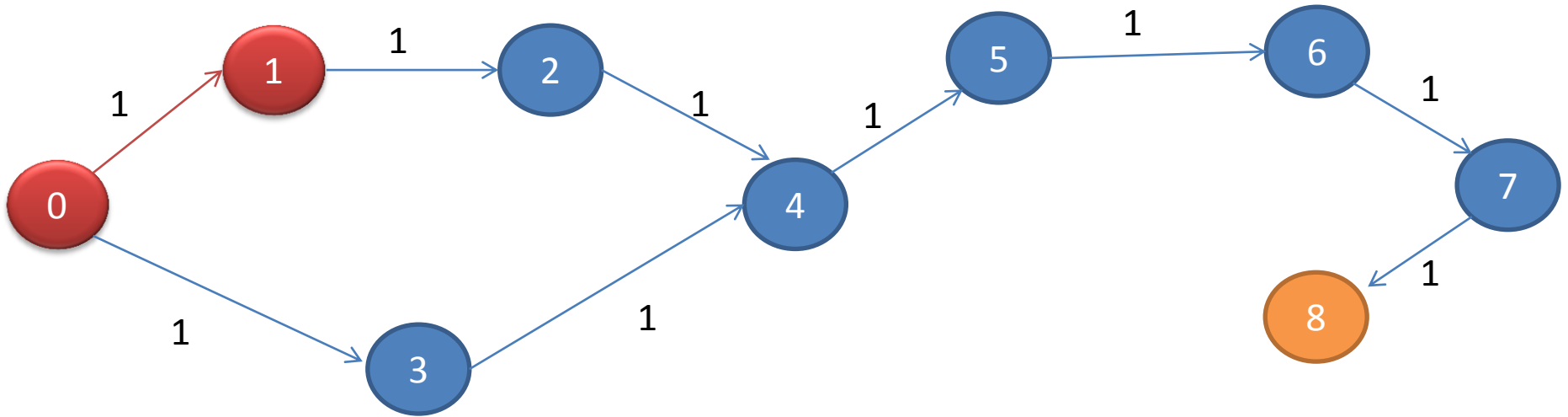
- Iteration 1:**



**The node chosen to be expanded from the open list: 0**

# Parent Pointer Redirection

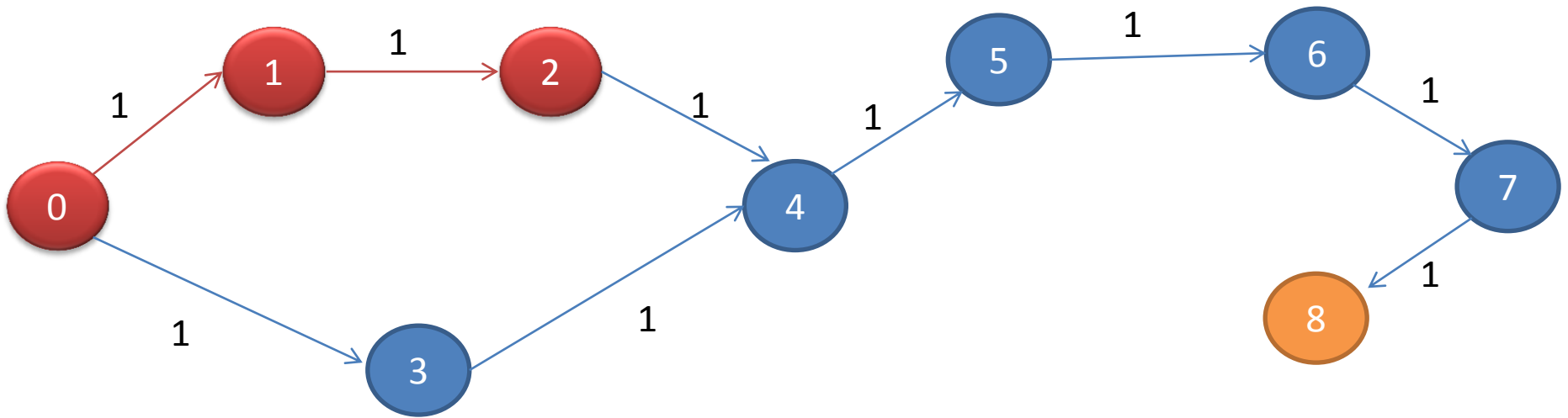
- Iteration 2:**



**The node chosen to be expanded from the open list: 1**

# Parent Pointer Redirection

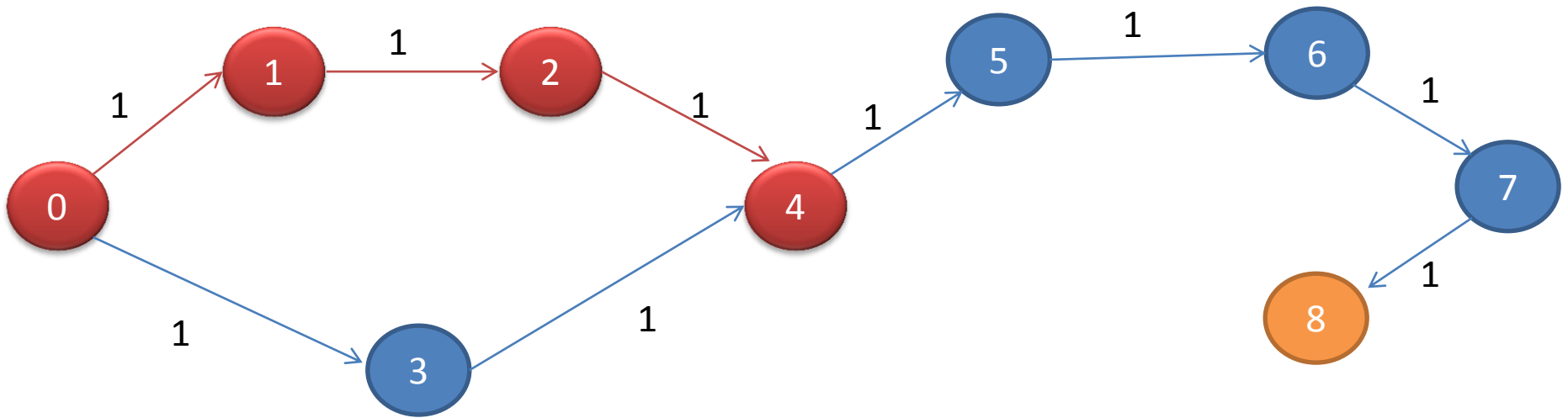
- Iteration 3:**



**The node chosen to be expanded from the open list: 2**

# Parent Pointer Redirection

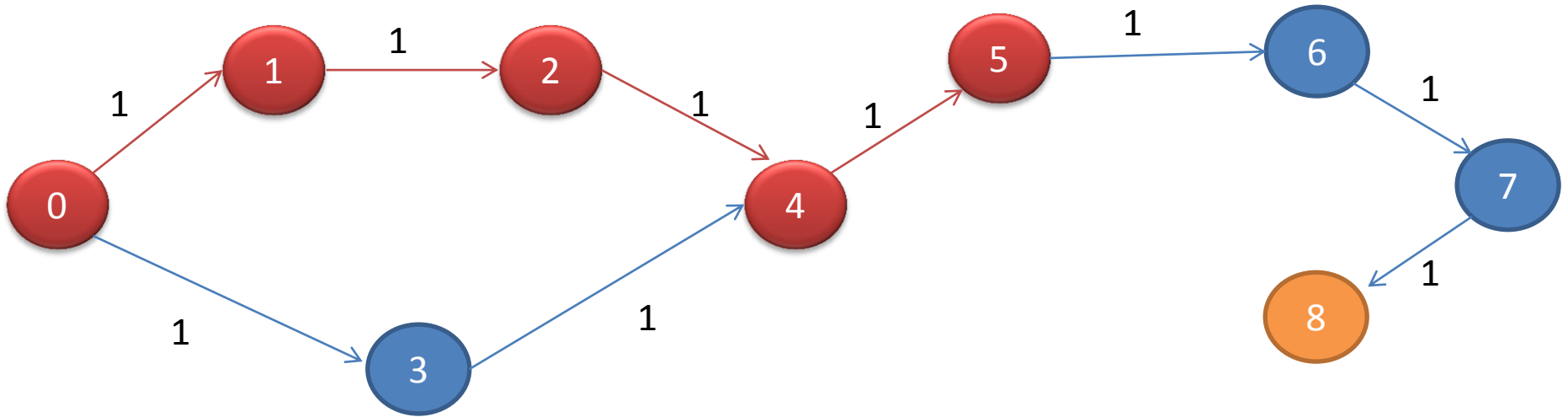
- Iteration 4:



The node chosen to be expanded from the open list: 4

# Parent Pointer Redirection

- Iteration 5:

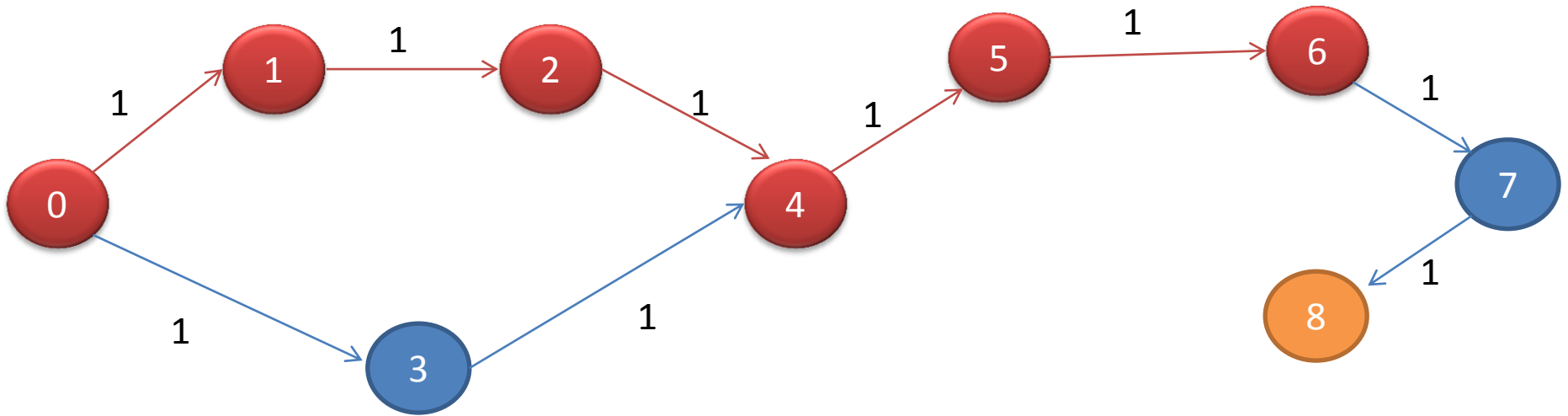


The node chosen to be expanded from the open list: 5



# Parent Pointer Redirection

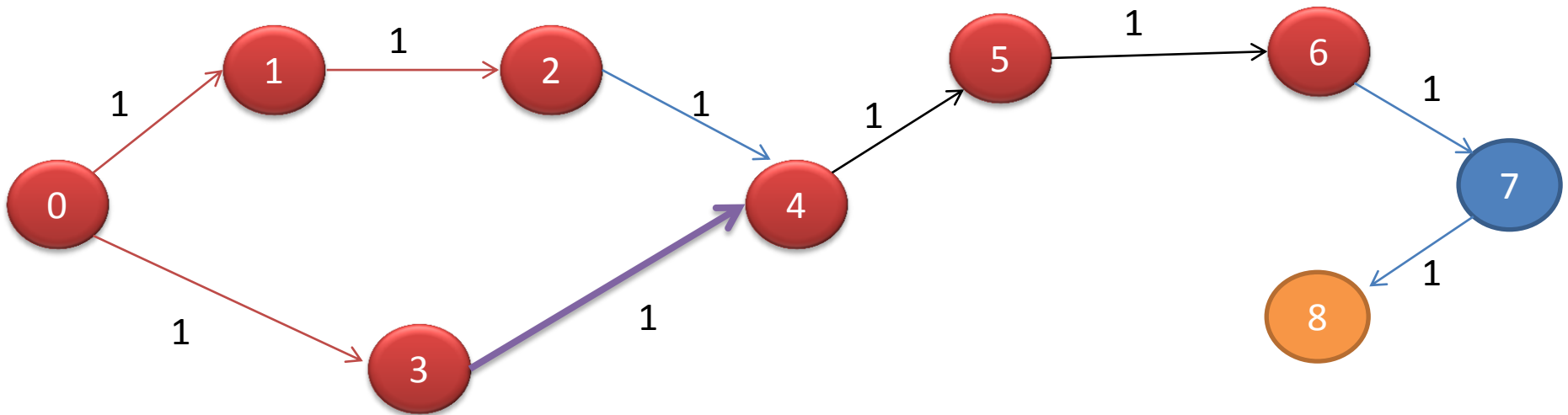
- Iteration 6:



The node chosen to be expanded from the open list: 6

# Parent Pointer Redirection

- Iteration 7:



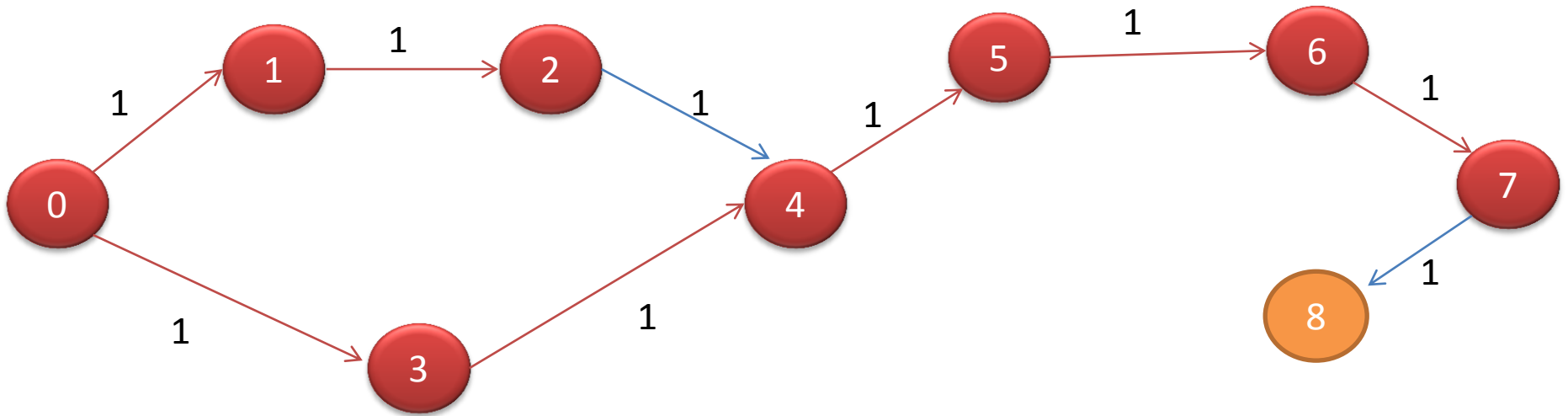
The node chosen to be expanded from the open list: 3

1 Parent Pointers Redirected: 3 → 4

g values changed for 3 nodes: 4 5 6

# Parent Pointer Redirection

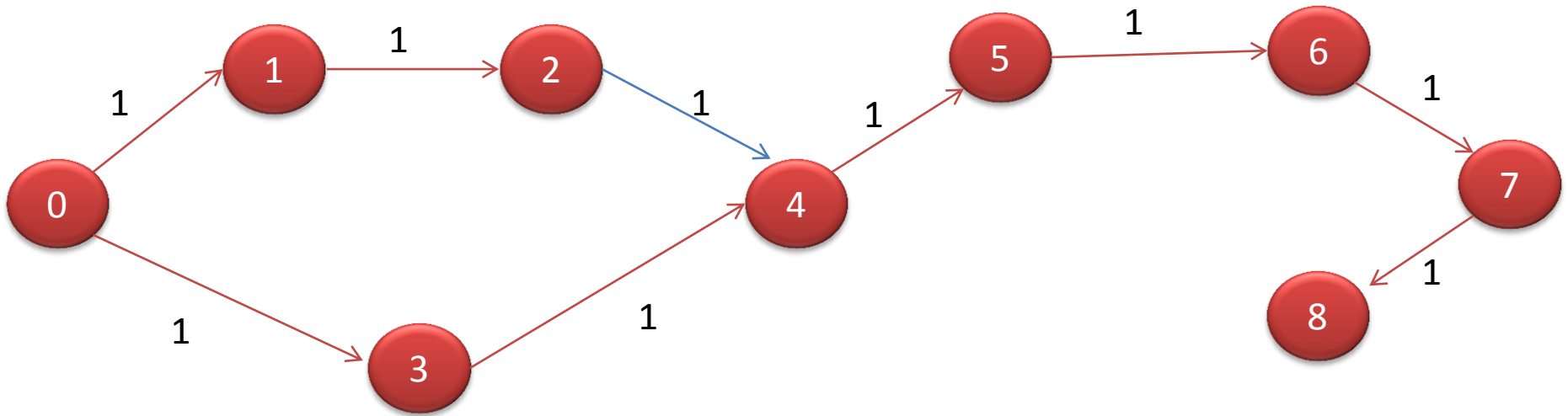
- Iteration 8:



The node chosen to be expanded from the open list: 7

# Parent Pointer Redirection

- **Final Output:**



- The optimal path is: 0 3 4 5 6 7 8
- Optimal path cost is 6
- Number of iterations taken by the A\* Algo = 8
- Number of Parent Pointer Redirections = 3

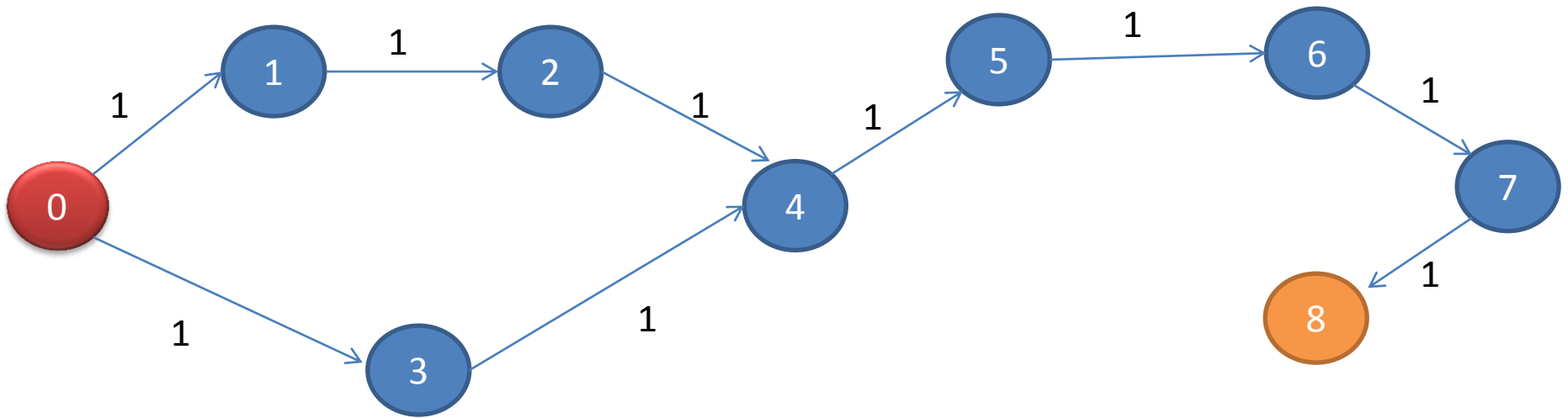
# Better Heuristic Performs Better

- We now change the heuristic of previous case as follows:

– $h(0) = 6$	$h(5) = 3$
– $h(1) = 6$	$h(6) = 2$
– $h(2) = 5$	$h(7) = 1$
– $h(3) = 5$	$h(8) = 0$
– $h(4) = 4$	
- Clearly above heuristic is better than the previous heuristic which was  $h(i) = \begin{matrix} 5 & \text{for } i=3 \\ 0 & \text{otherwise} \end{matrix}$

# Better Heuristic Performs Better

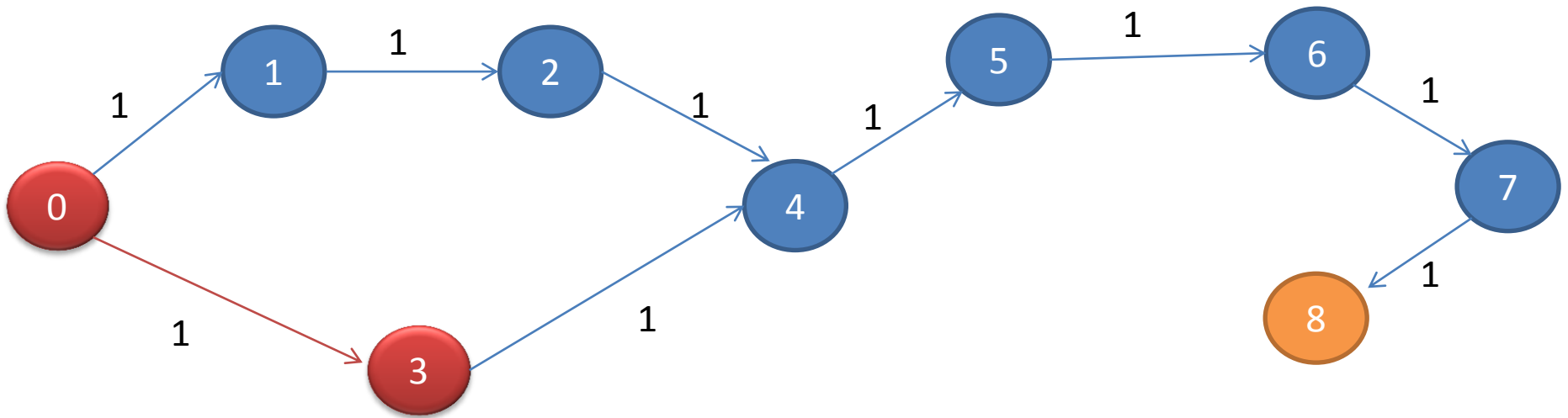
- Iteration 1:



The node chosen to be expanded from the open list: 0

# Better Heuristic Performs Better

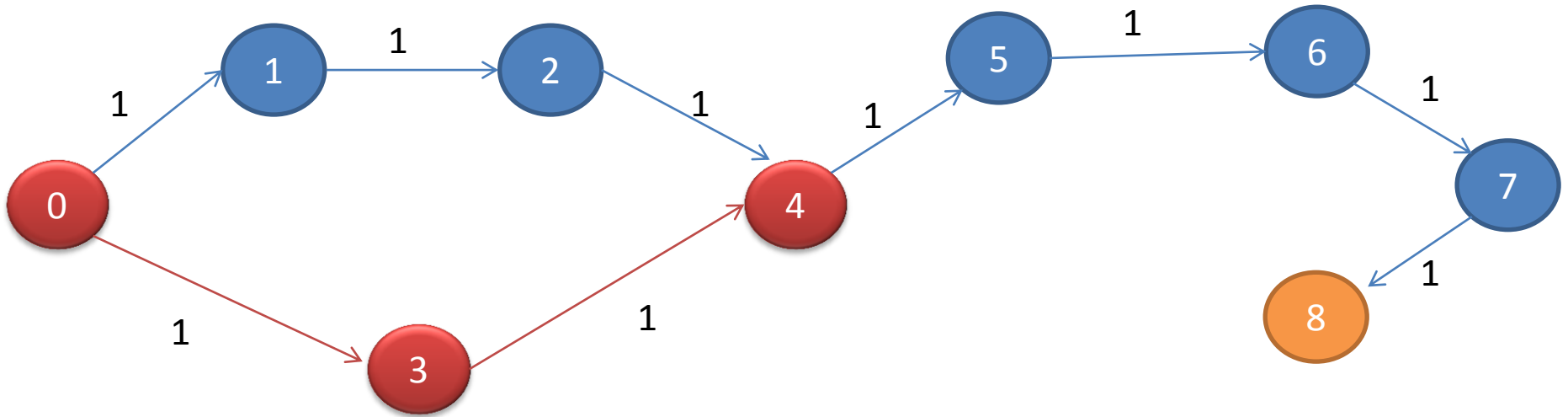
- Iteration 2:**



**The node chosen to be expanded from the open list: 3**

# Better Heuristic Performs Better

- Iteration 3:

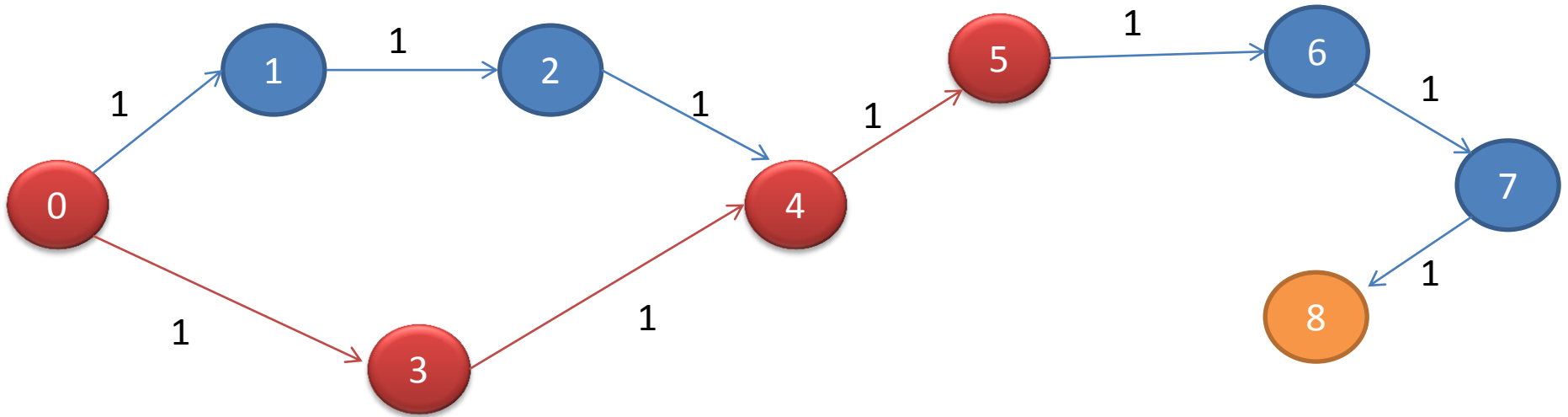


The node chosen to be expanded from the open list: 4



# Better Heuristic Performs Better

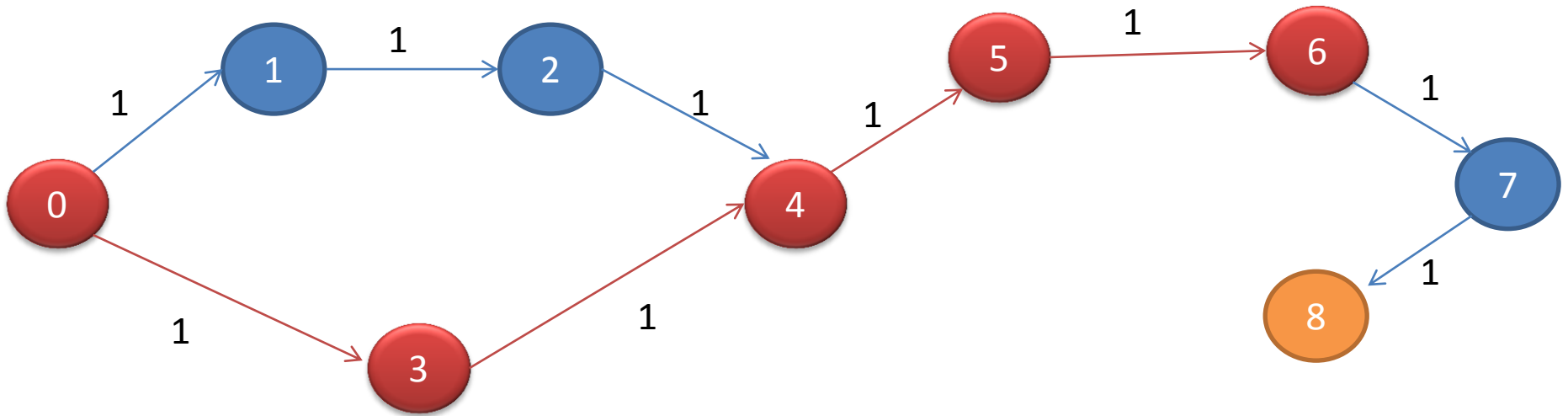
- Iteration 4:



The node chosen to be expanded from the open list: 5

# Better Heuristic Performs Better

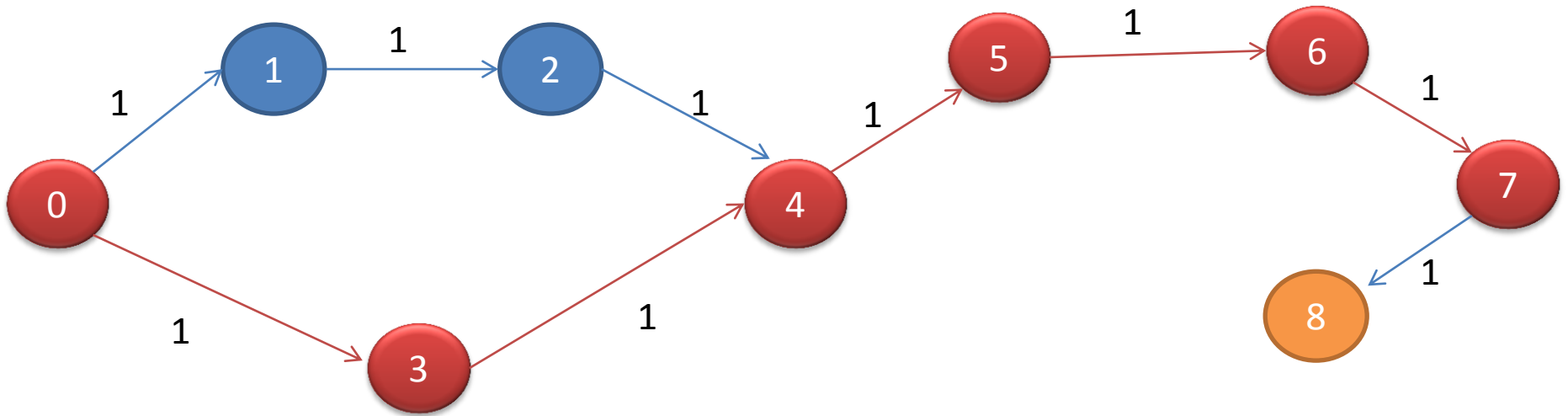
- Iteration 5:



The node chosen to be expanded from the open list: 6

# Better Heuristic Performs Better

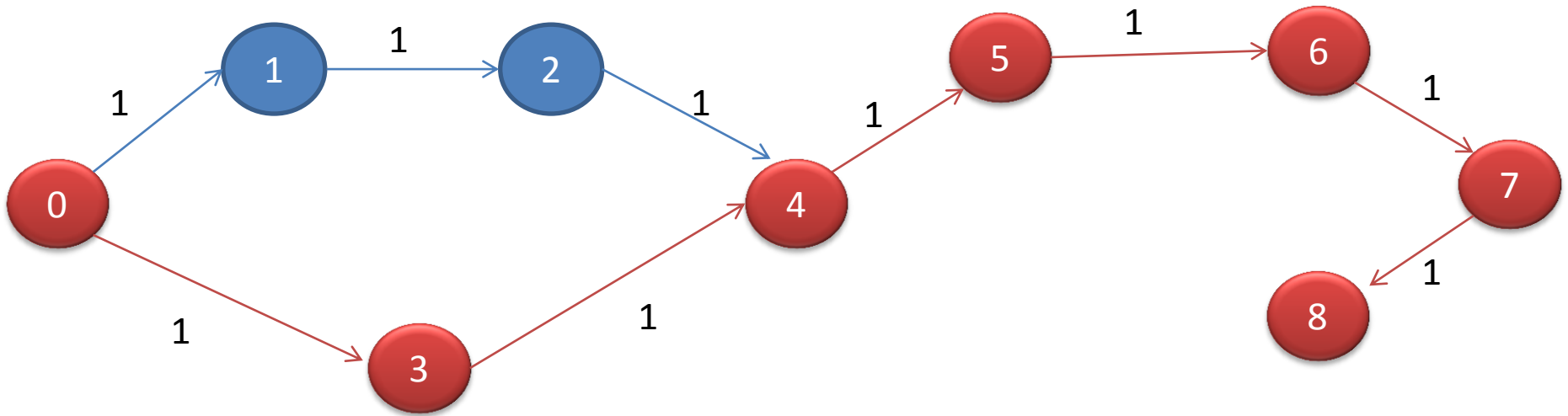
- Iteration 6:



The node chosen to be expanded from the open list: 7

# Better Heuristic Performs Better

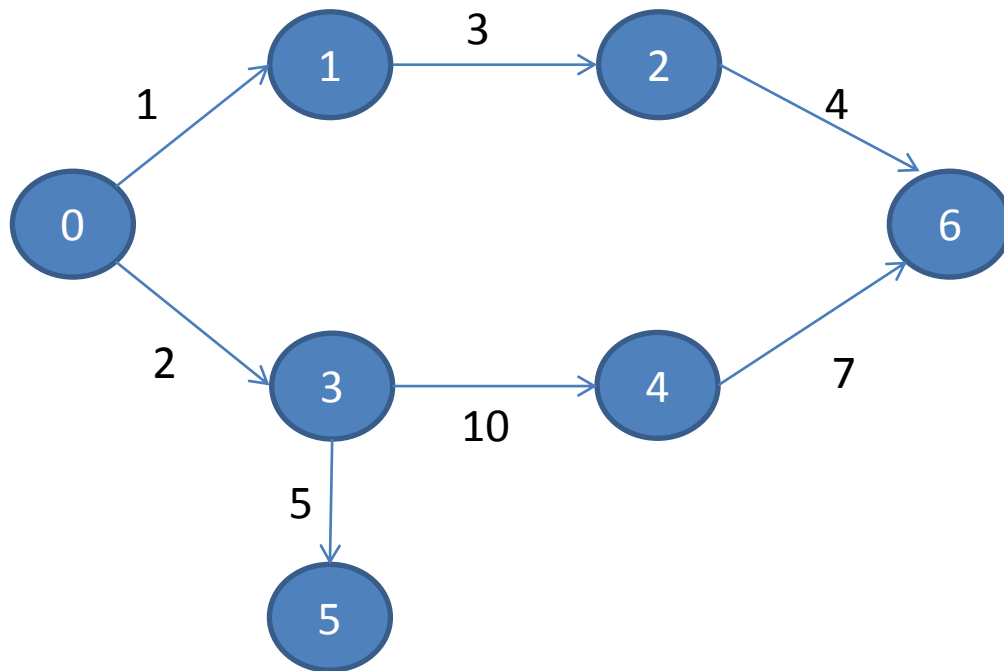
- **Final Output:**



- The optimal path is: **0 3 4 5 6 7 8**
- Optimal path cost is **6**
- Number of iterations taken by the A\* Algo = **6** (as opposed to 8 previously)
- Number of Parent Pointer Redirections = **0**
- So we see that with a better heuristic A\* algorithms converges faster

# $h$ greater than $h^*$

- We now verify that “if  $h(n) > h^*(n)$ ”, for all  $n$ ,  $A^*$  may find the goal faster, but may discover a suboptimal path.
- We consider the graph used initially and run the  $A^*$  algorithm on it using a heuristic such that  $h > h^*(n)$  for all  $n$



Start Node: 0

Goal Node: 6

$h(1) = 20$

$h(1) = 30$

$h(1) = 40$

$h(1) = 10$

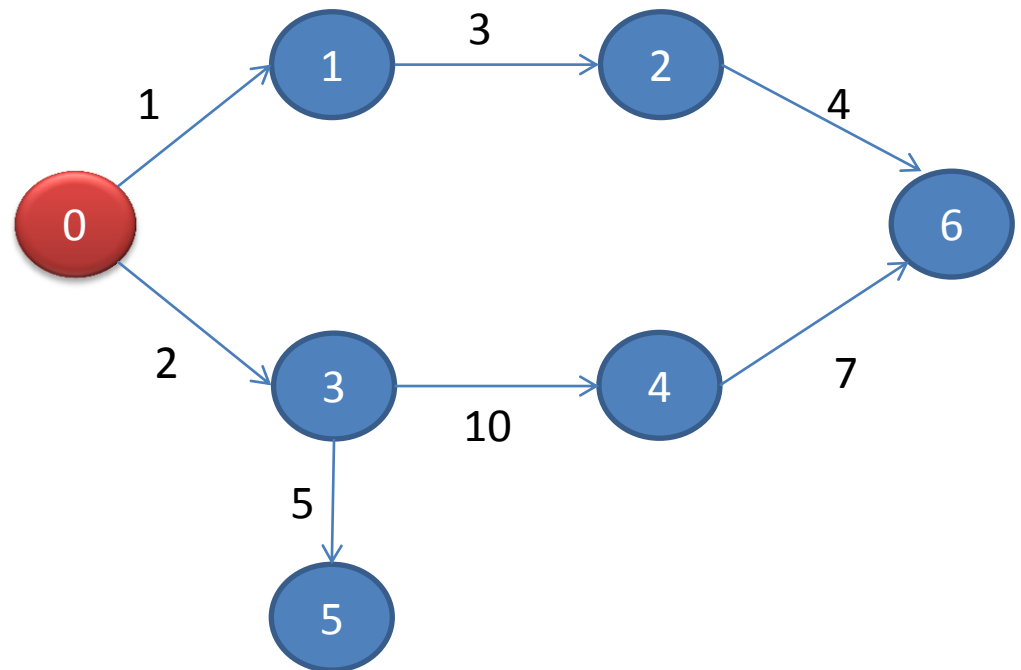
$h(1) = 10$

$h(1) = 10$

$h(1) = 0$

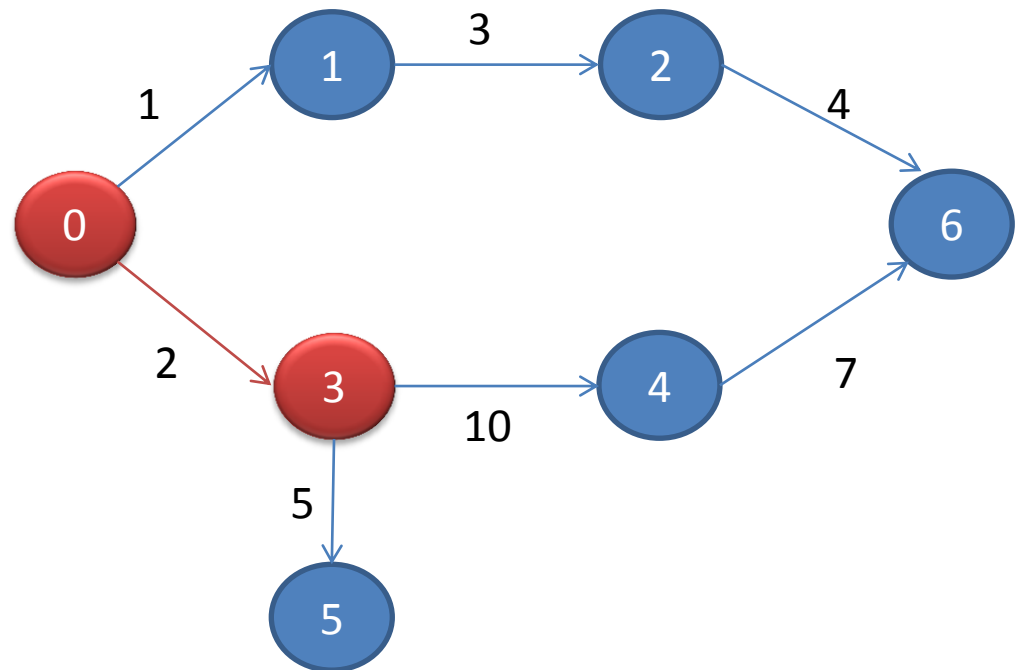
# $h$ greater than $h^*$

- **Node picked by algorithm from open list:**
  - Iteration 1: Node 0



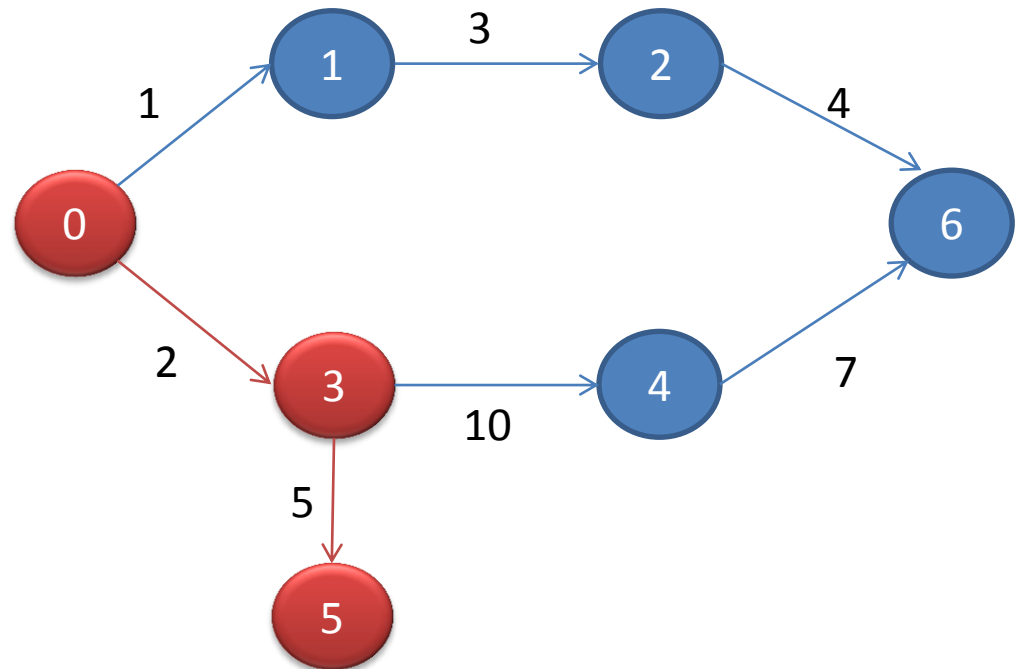
# $h$ greater than $h^*$

- **Node picked by algorithm from open list:**
  - Iteration 1: Node 0
  - Iteration 2: Node 3



# $h$ greater than $h^*$

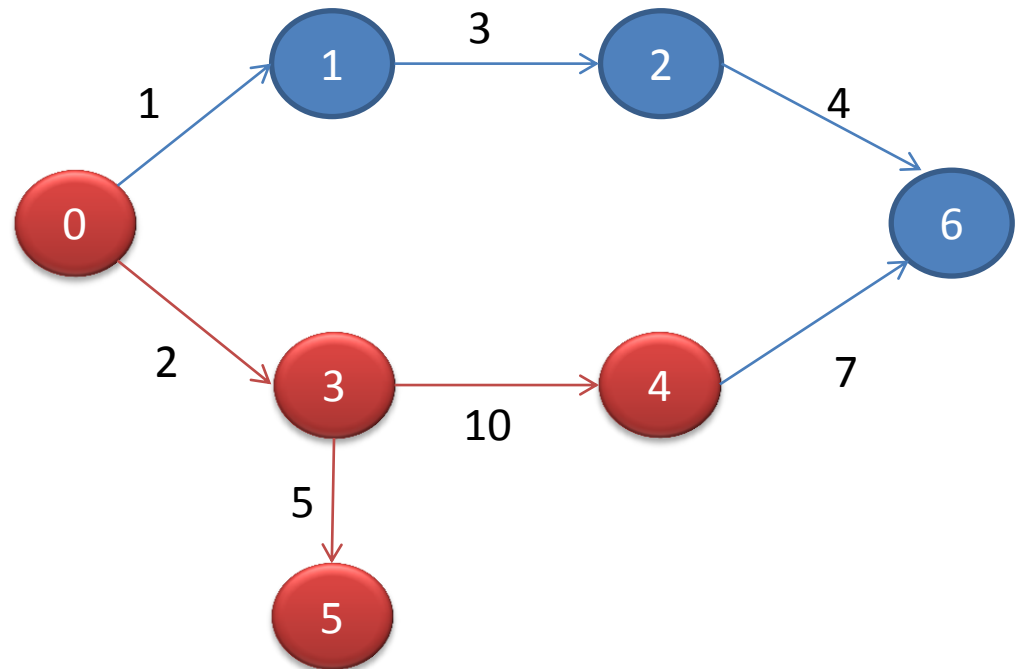
- **Node picked by algorithm from open list:**
  - Iteration 1: Node 0
  - Iteration 2: Node 3
  - Iteration 3: Node 5





# $h$ greater than $h^*$

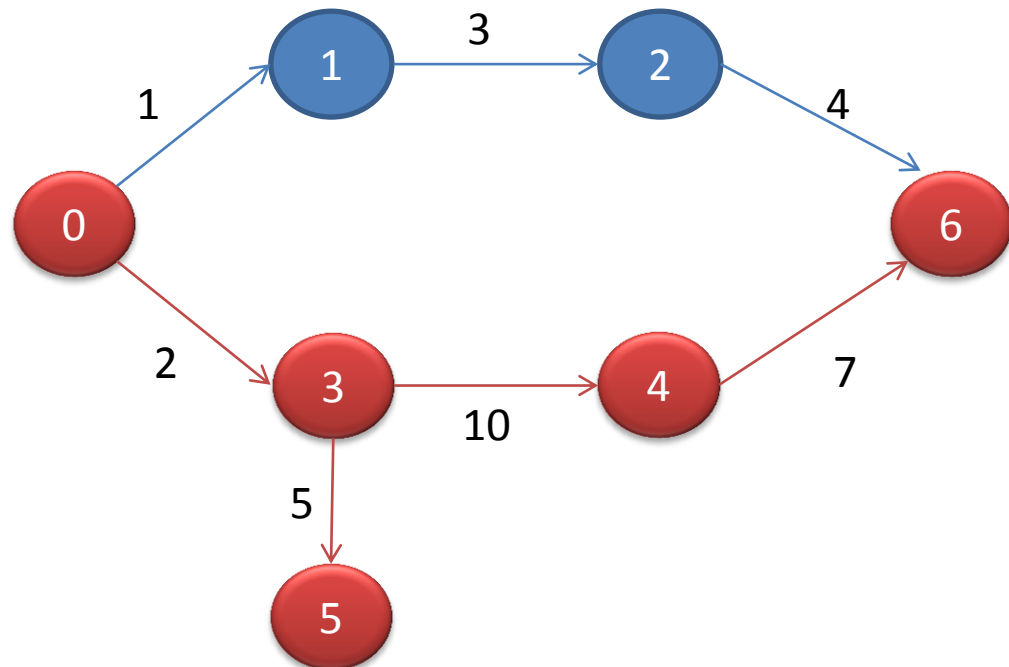
- **Node picked by algorithm from open list:**
  - Iteration 1: Node 0
  - Iteration 2: Node 3
  - Iteration 3: Node 5
  - Iteration 4: Node 4



# $h$ greater than $h^*$

- **Final Output:**

- Number of iterations taken by the A\* Algo = 4
- Number of Parent Pointer Redirections = 0
- Found path is: 0 3 4 6
- Found path cost is 19
- The discovered path is a suboptimal path.



# Monotone Restriction

- We now change the heuristic of second case as follows:

- $h(0) = 6$

- $h(5) = 3$

- $h(1) = 6$

- $h(6) = 2$

- $h(2) = 5$

- $h(7) = 1$

- $h(3) = 5$

- $h(8) = 0$

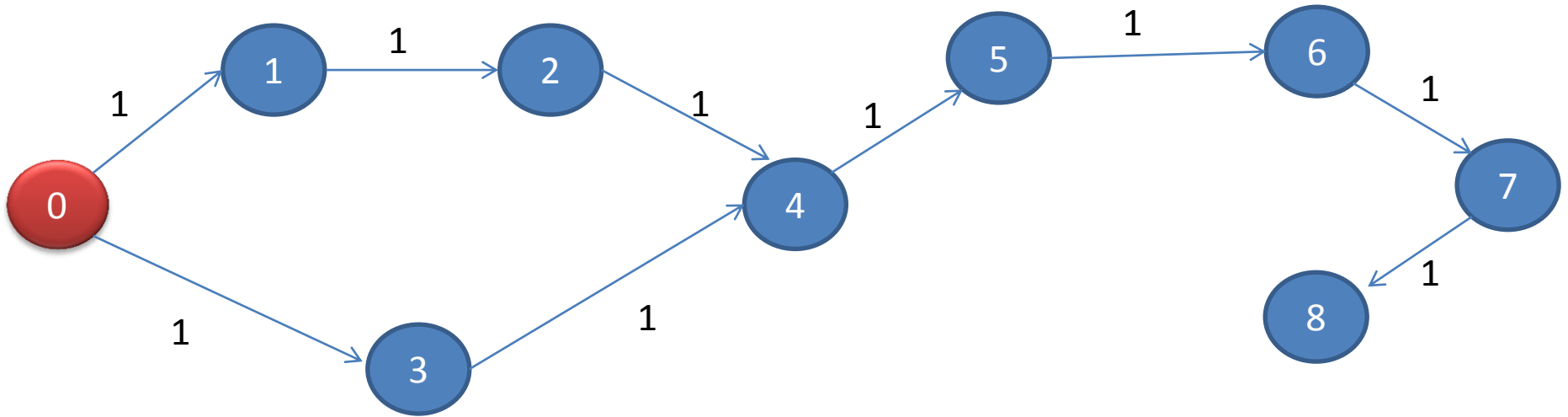
- $h(4) = 4$

- Note that in the above case MR is satisfied i.e.

$$h(\text{parent}) \leq h(\text{child}) + C(\text{parent}, \text{child})$$

# Monotone Restriction

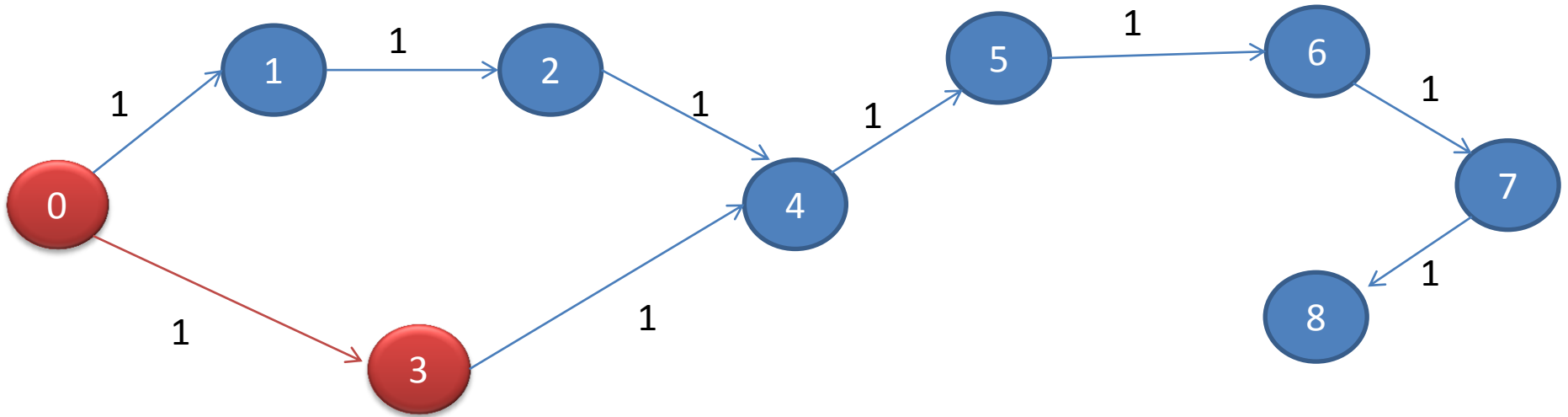
- Iteration 1:



The node chosen to be expanded from the open list: 0

# Monotone Restriction

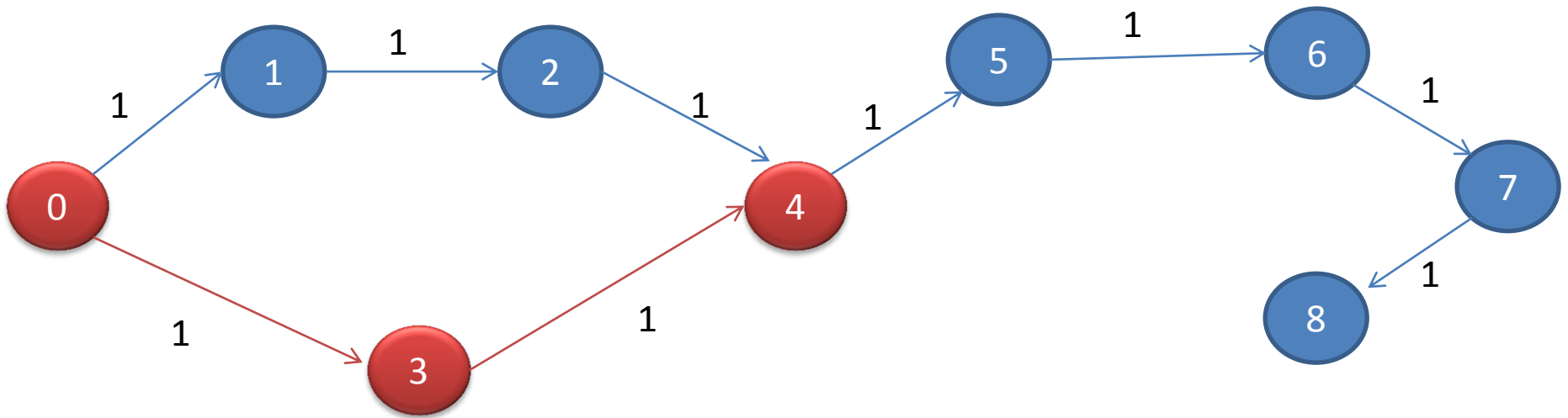
- Iteration 2:



The node chosen to be expanded from the open list: 3

# Monotone Restriction

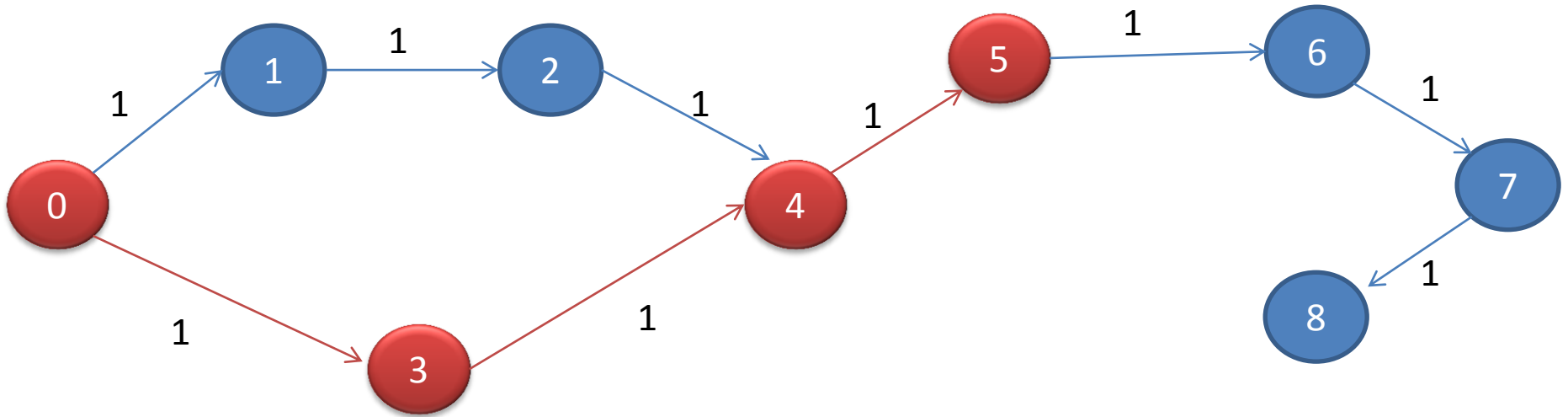
- Iteration 3:



The node chosen to be expanded from the open list: 4

# Monotone Restriction

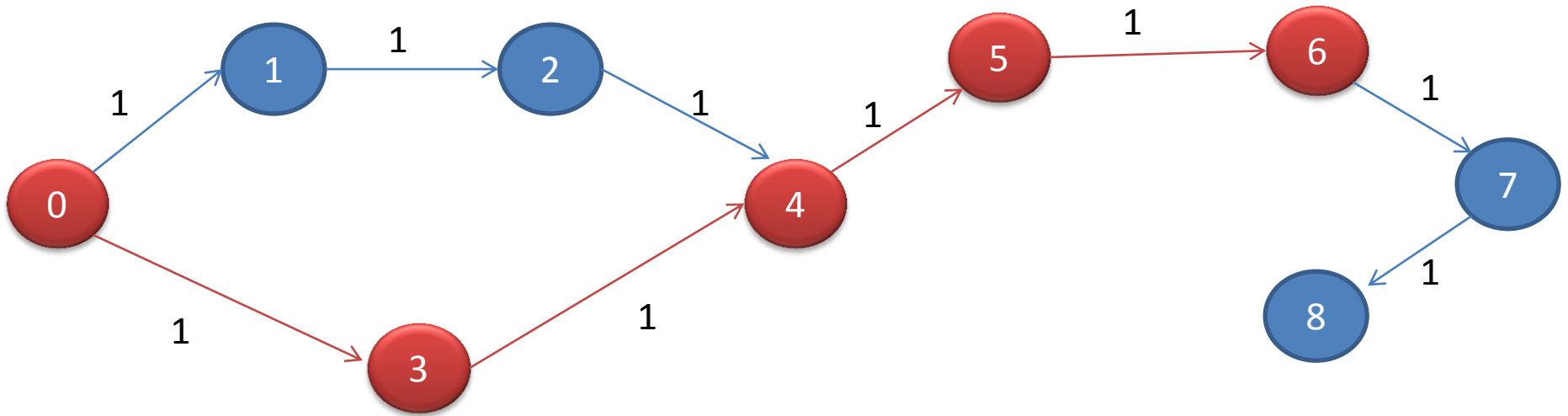
- Iteration 4:



The node chosen to be expanded from the open list: 5

# Monotone Restriction

- Iteration 5:

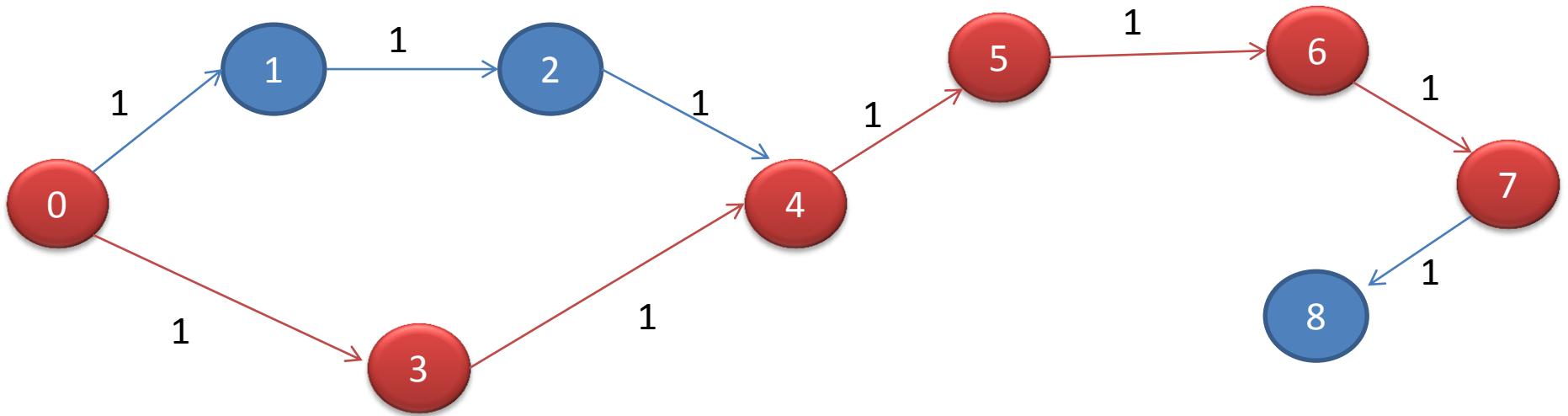


The node chosen to be expanded from the open list: 6



# Monotone Restriction

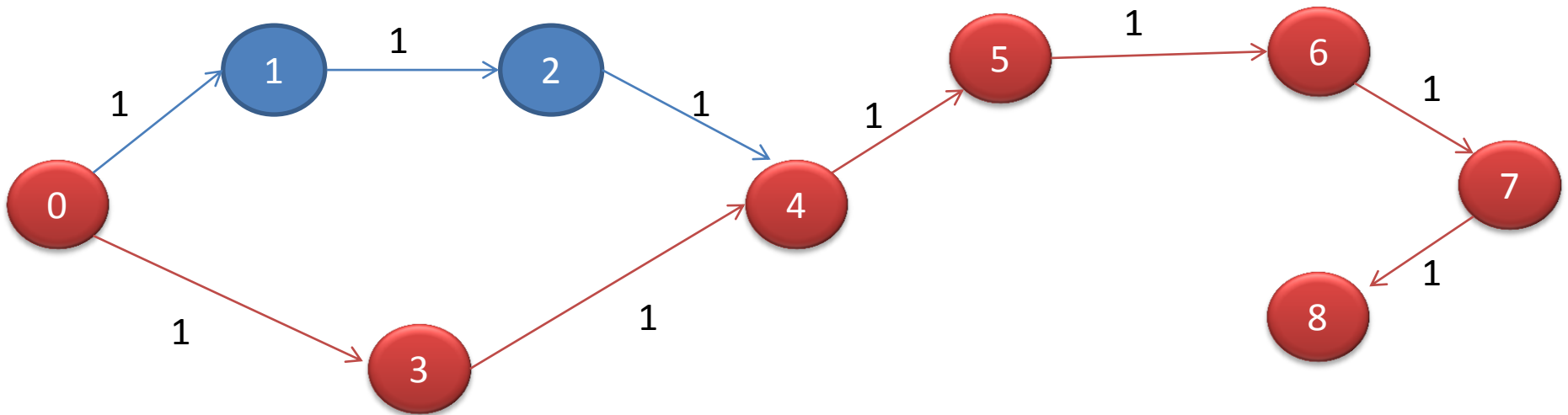
- Iteration 6:



The node chosen to be expanded from the open list: 7

# Monotone Restriction

- **Final Output:**



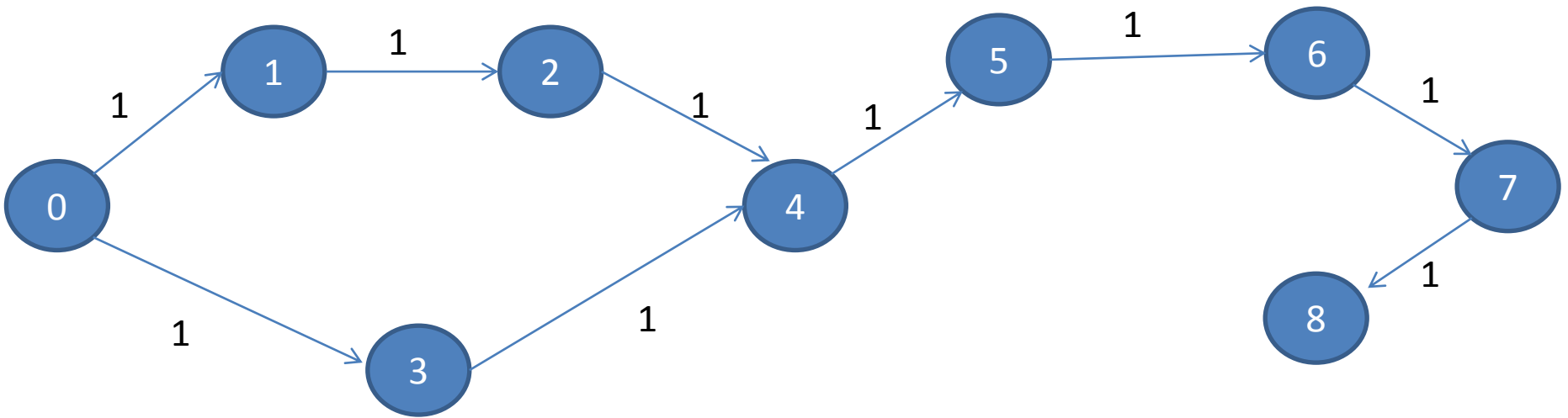
- The optimal path is: **0 3 4 5 6 7 8**
- Optimal path cost is **6**
- Number of iterations taken by the A\* Algo = **6**
- Number of Parent Pointer Redirections = **0**
- **So we see that with a when Monotone Restriction is satisfied, parent pointer redirection is not needed**

# Bidirectional Search Algorithm

- Carry out the Forward and Backward search parallelly.
- For both Forward Search & Backward Search, maintain separate Open Lists and Closed Lists.
- Push the Start Node in Open List of Forward Search and Goal Node in Open List of Backward Search.
- In each iteration of A\* Algorithm,
  - Carry out Forward Search.
  - Carry out Backward Search.
  - Check the intersection of CL of Forward Search & CL of Backward Search.
  - If intersection is not NULL, Stop!
  - Otherwise carry out another iteration of A\* Search
- The path discovered is given by :
  - Following the parent pointers from Start Node to Intersection Node in Forward Search
  - Reversing the direction of Parent Pointers from intersection node to Start node of Backward Search

# Bidirectional Search

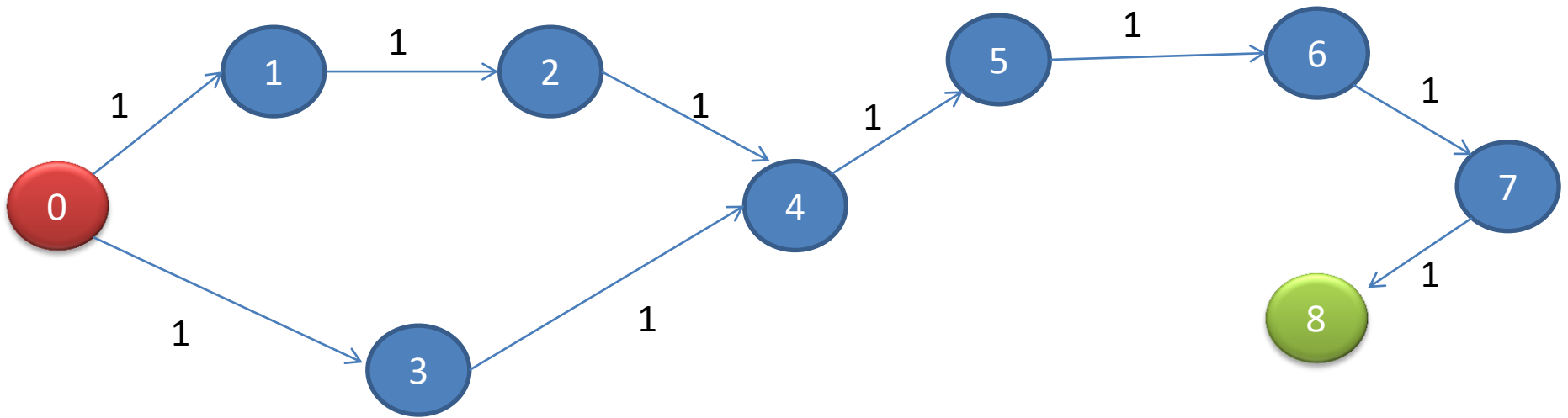
- We shall use the following graph & heuristic, to do a bidirectional search.



- **Heuristic**  $h(i) = 5$  for  $i=3$
- $0$  otherwise
- **Forward Pass:** Start Node: 0    Goal Node: 8
- **Reverse Pass:** Start Node: 8    Goal Node: 0
- We chose a node to be expanded from the Open List in forward direction and similarly in backward direction.

# Bidirectional Search

- Iteration 1:

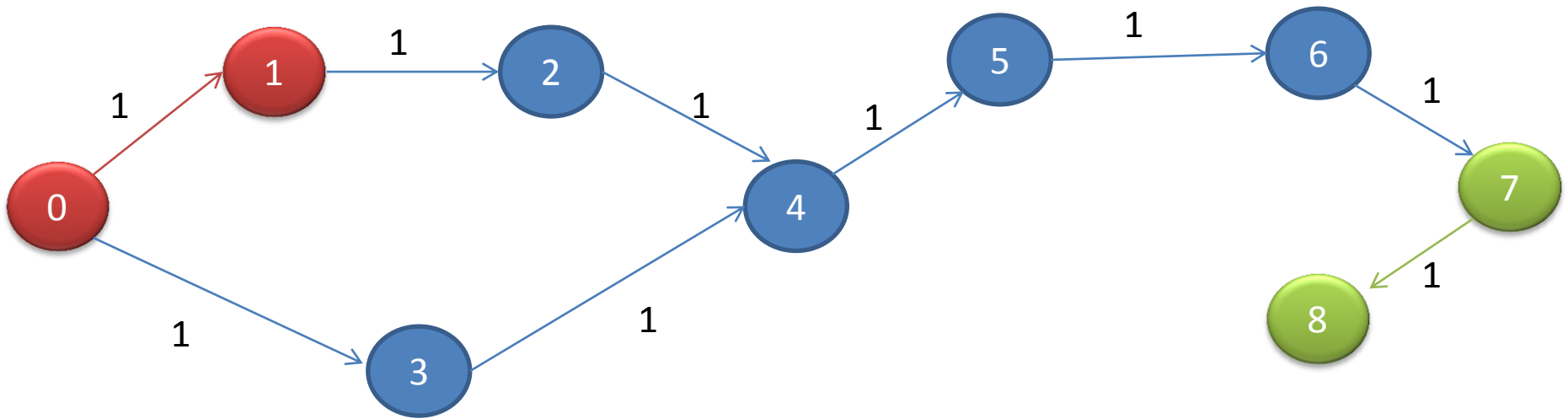


The node chosen to be expanded from the open list in FORWARD direction: 0

The node chosen to be expanded from the open list in BACKWARD direction: 8

# Bidirectional Search

- Iteration 2:**

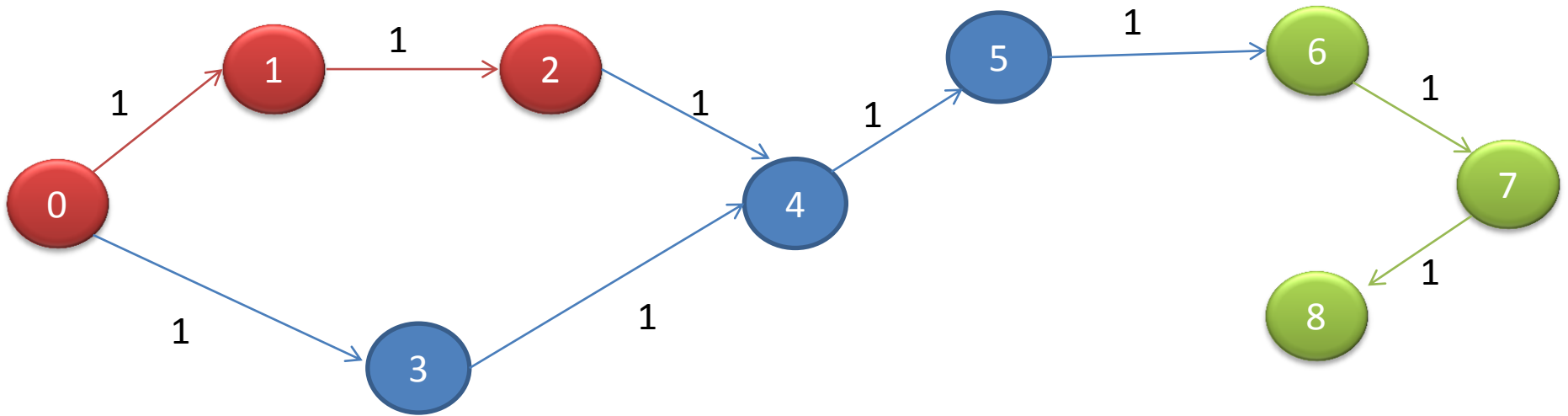


**The node chosen to be expanded from the open list in FORWARD direction: 1**

**The node chosen to be expanded from the open list in BACKWARD direction: 7**

# Bidirectional Search

- Iteration 3:**

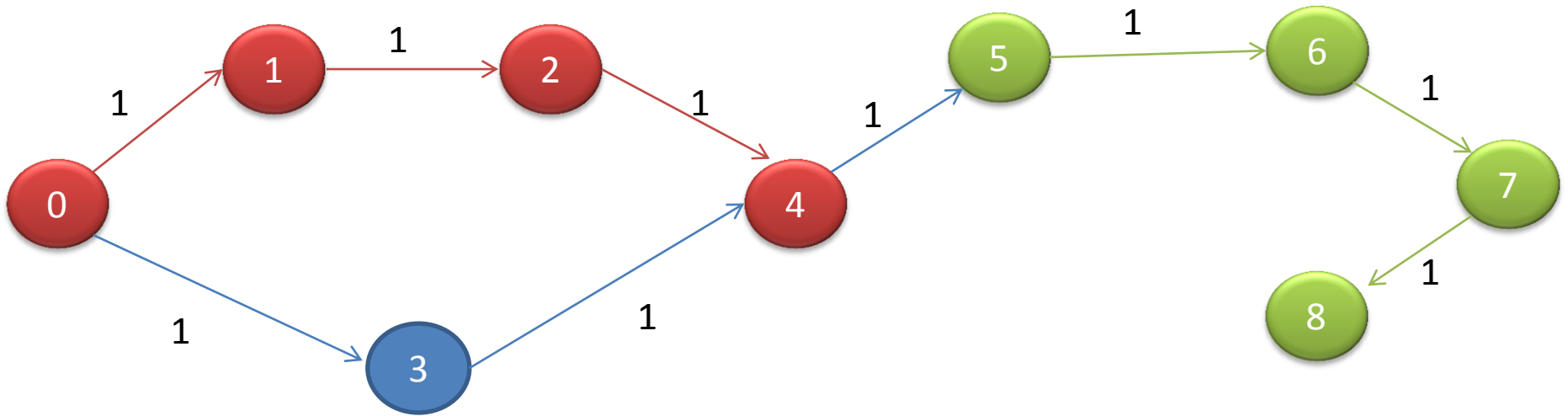


**The node chosen to be expanded from the open list in FORWARD direction: 2**

**The node chosen to be expanded from the open list in BACKWARD direction: 6**

# Bidirectional Search

- Iteration 4:



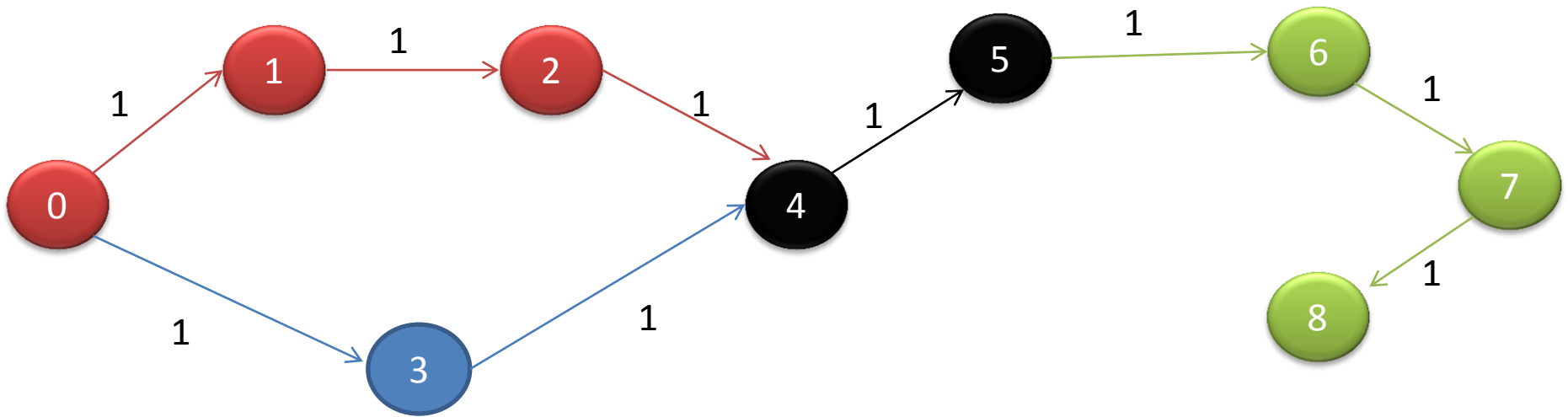
The node chosen to be expanded from the open list in FORWARD direction: 4

The node chosen to be expanded from the open list in BACKWARD direction: 5



# Bidirectional Search

- **Stop Condition Reached:**



- The search stops when the two searches meet at a common node.
- **The Algorithm takes only 4 iterations to converge.**
- **However the path found is NOT an optimal path.**
- **Cost of discovered path = 7**
- **Cost of optimal path = 6 (via 0-3-4-5-6-7-8)**

# Missionaries and Cannibals

- Constraints
  - The boat can carry at most 2 people
  - On no bank should the cannibals outnumber the missionaries
- State :  $\langle \#M, \#C, P \rangle$ 
  - $\#M$  = Number of missionaries on bank  $L$
  - $\#C$  = Number of cannibals on bank  $L$
  - $P$  = Position of the boat
- *Start State* =  $\langle 3, 3, L \rangle$
- *Goal State* =  $\langle 0, 0, R \rangle$
- Operations
  - $M2$  = Two missionaries take boat
  - $M1$  = One missionary takes boat
  - $C2$  = Two cannibals take boat
  - $C1$  = One cannibal takes boat
  - $MC$  = One missionary and one cannibal takes boat

# Missionaries and Cannibals

- **Heuristic,  $h(n) = 2n-1$** 
  - **where  $n$ =number of people on left bank**
  - **if  $n = 0$  ,  $h = 0$  as we have reached the goal state**
- Nodes chosen to be expanded from the open list: 0 1 3 2 4 5 6 7 8 9 10 11 12
- Optimal Path discovered:  
0 1 4 5 6 7 8 9 10 11 12 14
- Optimal Path cost = 11
- Number of iterations taken by A\* Algorithm = 13
- Number of parent pointer redirections = 0
- Note that this heuristic is **Neither Admissible Nor Monotone**
- Consider state  $S = \langle 2,2,L \rangle$ 
  - $h^*(S) = 5$  (optimal path :  $\langle 2,2,L \rangle - \langle 0,2,R \rangle - \langle 0,3,L \rangle - \langle 0,1,R \rangle - \langle 0,2,L \rangle - \langle 0,0,R \rangle$ )
  - $h(S) = 7$  ( $2*4-1 = 7$ )      Hence Not Admissible
- Also for  $P = \langle 2,2,L \rangle$        $C = \langle 0,2,L \rangle$ 
  - $h(P) = 7$        $h(C) = 3$
  - $C(P,C) = 1$       Hence Not Monotone

# Missionaries and Cannibals

- **Heuristic,  $h(n) = 2n+1$** 
  - **where  $n$ =number of people on left bank**
  - **if  $n = 0$  ,  $h = 0$  as we have reached the goal state**
- Nodes chosen to be expanded from the open list: 0 1 3 2 4 5 6 7 8 9 10 11 12
- Optimal Path discovered:  
0 1 4 5 6 7 8 9 10 11 12 14
- Optimal Path cost = 11
- Number of iterations taken by A\* Algorithm = 13
- Number of parent pointer redirections = 0
- Note that this heuristic is **Neither Admissible Nor Monotone**
- Consider state  $S = \langle 2,2,L \rangle$ 
  - $h^*(S) = 5$  (optimal path :  $\langle 2,2,L \rangle - \langle 0,2,R \rangle - \langle 0,3,L \rangle - \langle 0,1,R \rangle - \langle 0,2,L \rangle - \langle 0,0,R \rangle$ )
  - $h(S) = 9$  ( $2*4+1 = 9$ )      Hence Not Admissible
- Also for  $P = \langle 2,2,L \rangle$   $C = \langle 0,2,L \rangle$ 
  - $h(P) = 9$        $h(C) = 5$
  - $C(P,C) = 1$       Hence Not Monotone

# Missionaries and Cannibals

- **Heuristic,  $h(n) = n-1$** 
  - **where  $n$ =number of people on left bank**
  - **if  $n = 0$  ,  $h = 0$  as we have reached the goal state**
- Nodes chosen to be expanded from the open list: 0 1 3 2 4 5 6 7 8 9 10 11 12 13
- Optimal Path discovered:  
0 1 4 5 6 7 8 9 10 11 12 14
- Optimal Path cost = 11
- Number of iterations taken by A\* Algorithm = 13
- Number of parent pointer redirections = 0
- Note that this heuristic is **Admissible but Not Monotone**
- Consider  $P = \langle 2, 2, L \rangle$   $C = \langle 0, 2, L \rangle$ 
  - $h(P) = 3$                        $h(C) = 1$
  - $C(P, C) = 1$                       Hence Not Monotone

# Missionaries and Cannibals

- **Heuristic,  $h(n) = n/2$** 
  - **where  $n$ =number of people on left bank**
  - **if  $n = 0$  ,  $h = 0$  as we have reached the goal state**
- Nodes chosen to be expanded from the open list: 0 3 4  
1 2 5 6 7 8 9 10 11 12 13
- Optimal Path discovered:  
0 3 4 5 6 7 8 9 10 11 12 14
- Optimal Path cost = 11
- Number of iterations taken by A\* Algorithm = 14
- Number of parent pointer redirections = 0
- This heuristic is Admissible as well as Monotone.

# 8-Puzzle Problem

2		3
1	8	5
4	7	6

Start State

1	2	3
4	5	6
7	8	

Goal State

- Tile movement represented as the movement of the blank space.
- Operators:
  - L : Blank moves left
  - R : Blank moves right
  - U : Blank moves up
  - D : Blank moves down
  - $C(L) = C(R) = C(U) = C(D) = 1$

# 8-Puzzle Problem

- Heuristic,  $h(n)$  = sum of Manhattan distances of tiles from their destined position

$h(n) = \text{X-dist} + \text{Y-dist from Goal State}$

- The optimal path is: 0-1- 4-10-22-39-69-119
- Optimal path cost is 7
- Number of nodes expanded from Open List =7
- Number of Parent Pointer Redirections =0



# 8-Puzzle Problem

- Heuristic,  $h(n)$  = no. of tiles displaced from their destined position.
- The optimal path is: 0 1 4 10 22 39 69 119
- Optimal path cost is 7
- Number of nodes expanded from Open List = 8
- Number of Parent Pointer Redirections = 0
- Here also we see that ***Manhattan Distance***, which is a better heuristic than ***No. of Displaced Tiles*** heuristic performs better (converges faster).

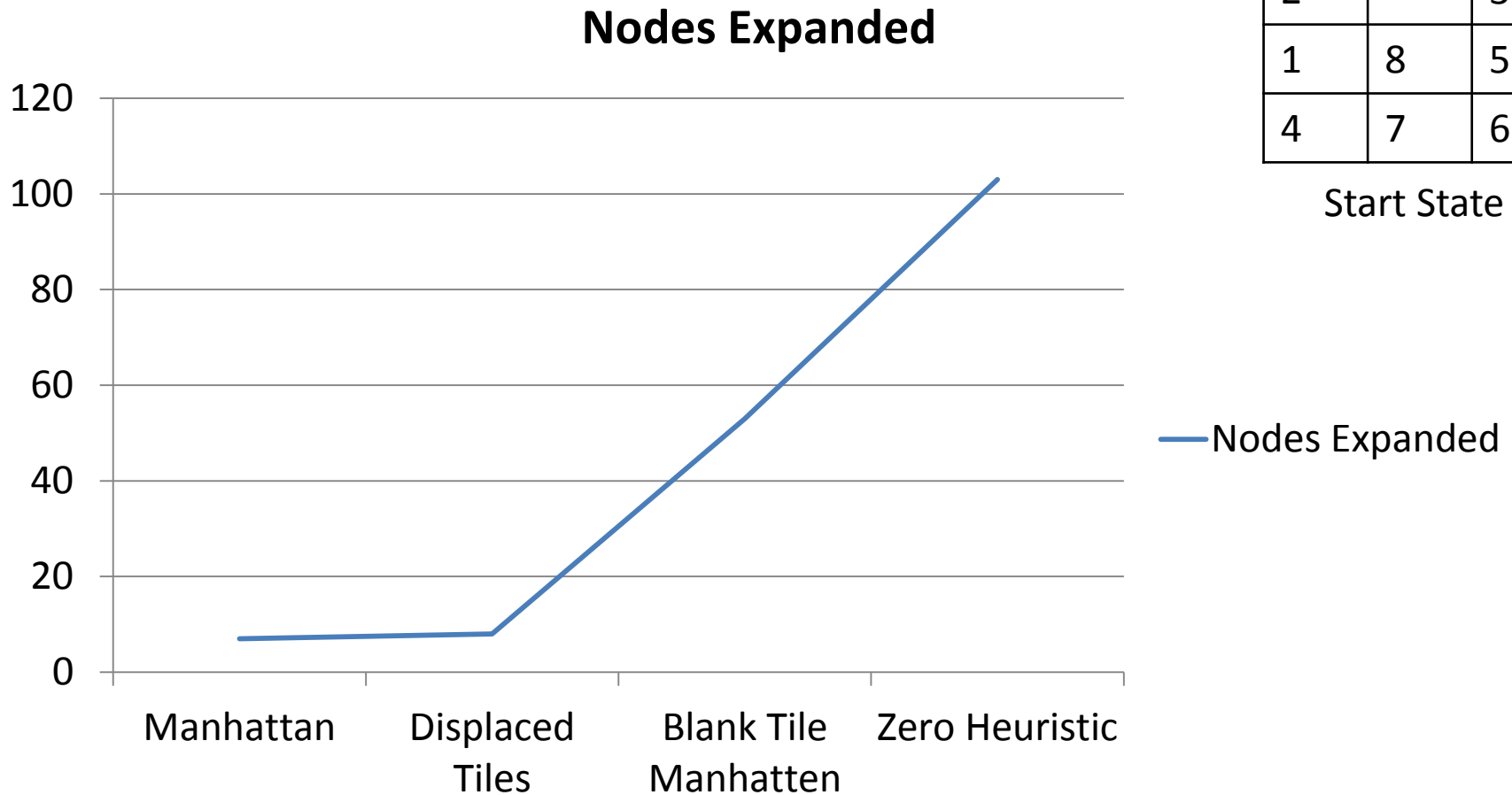
# 8-Puzzle Problem(our own heuristic)

- Heuristic,  $h(n)$  = Manhattan distances of only the blank tile
- The optimal path is:0-1- 4-10-22-39-69-119
- Optimal path cost is 7
- Number of nodes expanded from Open List =**53**
- Number of Parent Pointer Redirections =0
- Now we observe that this heuristic expands a lot higher number of nodes from the Open List as this is a very poor heuristic.

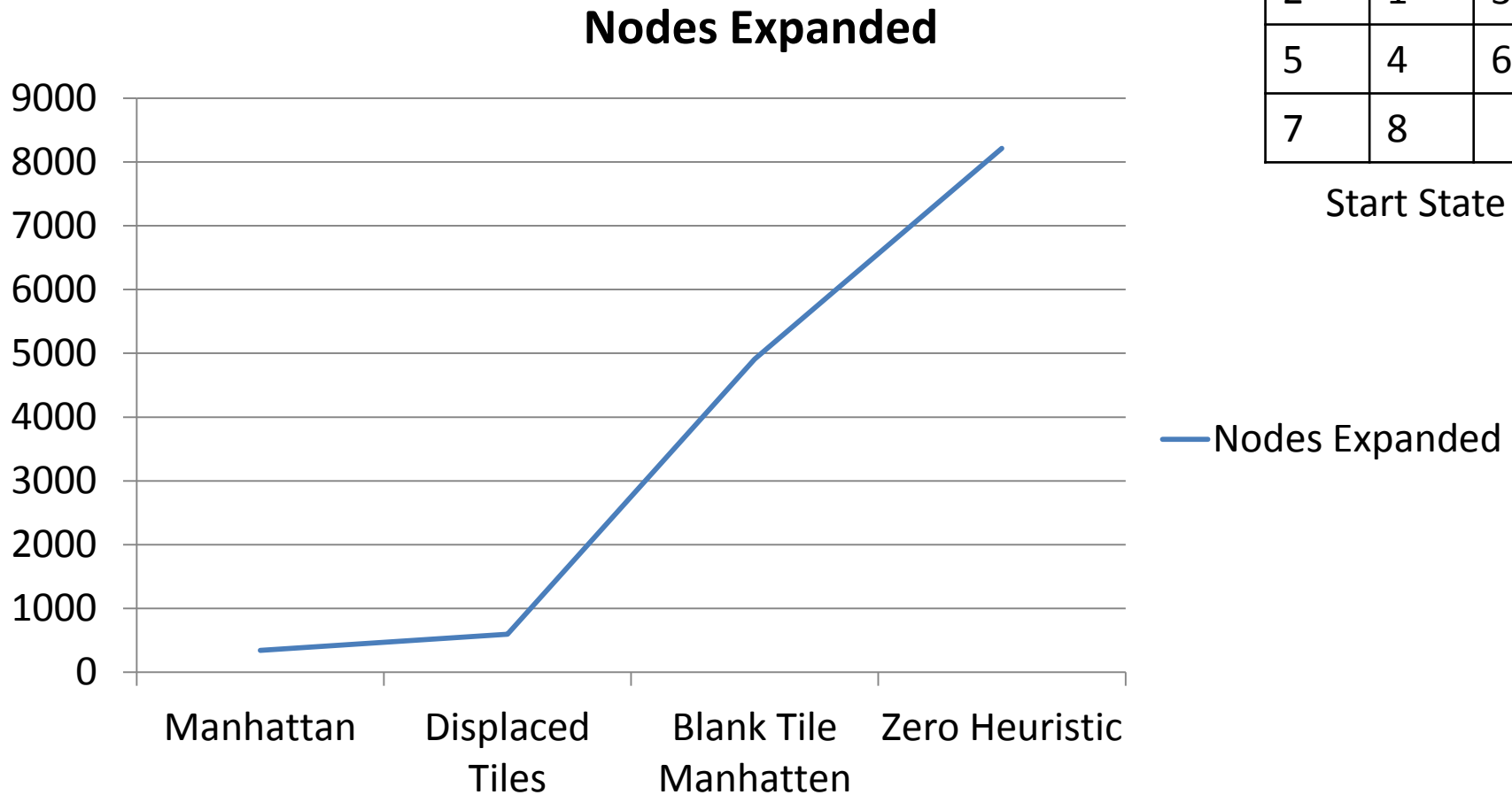
# 8-Puzzle Problem(our own heuristic)

- Heuristic,  $h(n) = 0$  for all nodes
- The optimal path is:0-1- 4-10-22-39-69-119
- Optimal path cost is 7
- Number of nodes expanded from Open List =103
- Number of Parent Pointer Redirections =0
- Now we observe that this heuristic expands even a lot higher number of nodes from the Open List as this is a very poor heuristic.

# Comparing Heuristics (8-Puzzle)



# Comparing Heuristics (8-Puzzle)



# 8-Puzzle Non Reachability

Start State

2	1	3
5	4	7
6	8	

Goal State

1	2	3
4	5	6
7	8	

- We now have a start state from which we can not reach to the goal state.
- However, we can figure the non-reachability before running the A\* Algorithm by using the following strategy:
  - I. Write the start state puzzle in row major form i.e. in above case start state can be written as [2 1 3 5 4 7 6 8].
  - II. Count the no. of inversions in this array.
  - III. Repeat the above steps (I) & (II) for goal state. Since Goal State is fixed, row major form is [1 2 3 4 5 6 7 8] & its no of inversions are zero.
  - IV. Now we Use the following rule:

“Start State with even no of inversions can reach a Goal State with even no. of inversions and Start State with odd no of inversions can reach a Goal State with odd no. of inversions”
  - V. Since our goal state contains 0 (even) inversions, so **our start state must have even no. of inversions.**