Comment (2024): this exam is from an older version of this course, taught at Brown. There are some differences of style and emphasis compared to Math 252 here. I've updated the notation in this file to notation we use in our course, but you may see some different notation in the solutions, since I used slightly different notation in this course.

- 1. Alice and Bob are using Elgamal encryption, with public parameters p=29, g=19. The summary table for Elgamal is on the back page of this packet. There is also a multiplication table for  $\mathbb{Z}/29\mathbb{Z}$ , so that you do not need to do the arithmetic by hand.
  - (a) [3 points] Alice chooses the private key a=9. What is her public key? Express your answer as an integer in  $\{0,1,2,\cdots,28\}$ .
  - (b) [3 points] Bob sends ciphertext (8, 25) to Alice. What is the plaintext? Express your answer as an integer in  $\{0, 1, 2, \dots, 28\}$ .
- 2. (a) [3 points] Suppose that  $a \mid m$ . Prove that the congruence  $ax \equiv ab \pmod{m}$  holds if and only if the congruence  $x \equiv b \pmod{\frac{m}{a}}$  holds (all variables are integers).
  - (b) [3 points] Suppose that gcd(a, m) = 1. Prove that the congruence  $ax \equiv ab \pmod{m}$  holds if and only if the congruence  $x \equiv b \pmod{m}$  holds.
- 3. [6 points] Write a program that reduces breaking Diffie-Hellman key exchange to breaking Elgamal encryption.

More precisely: suppose that Eve has written a function break\_elg(p,g,A,c1,c2) with the following behavior: if the arguments are as in Table 2.3 (back of the packet), then this function will return m. Make use of this function to write a function break\_dh(p,g,A,B), which accepts arguments as labeled in Table 2.2 and returns the corresponding shared secret.

You may use any functions that are built into Python (or any language you have written your homework in), plus the hypothetical function  $break_elg$ . You may also assume that you have already written a function  $ext_euclid(a,b)$ , with the following behavior: given two positive integers a, b, this function returns a list of three integers [u, v, d], where  $d = \gcd(a, b)$  and au + bv = d.

For full points, your program must require at most  $\mathcal{O}(\log p)$  arithmetic operations (not counting any operations needed to compute break\_elg).

- 4. Let p be a prime number, and g an element of  $(\mathbb{Z}/p\mathbb{Z})^{\times}$ .
  - (a) [3 points] Prove that if  $\operatorname{ord}_p g = 17$ , then  $p \equiv 1 \pmod{17}$ .
  - (b) [3 points] Prove conversely that if  $p \equiv 1 \pmod{17}$ , then there exists some element  $[g]_p$  of order 17.
- 5. [6 points] Suppose that Alice and Bob use the following variant of Elgamal. The parameters p, g are as in table 2.3, and Alice chooses a secret key a and public key A in the same manner as in table 2.3. However, instead of following table 2.3, Bob computes his ciphertext as follows: he chooses a random element k, and computes

$$c_1 \equiv A^k \pmod{p}$$
  
 $c_2 \equiv m \cdot q^k \pmod{p}$ ,

then sends  $(c_1, c_2)$  to Alice.

Explain how Alice can efficiently decipher messages, i.e. determine m from  $(c_1, c_2)$ . You will need to place a restriction on Alice's original choice of private key a in order for decryption to be possible; clearly state this restriction.

## Reference information. You may detach this sheet for easier use.

Public parameter creation											
A trusted party chooses and publishes a (large) prime p											
and an integer $g$ having large prime order in $\mathbb{F}_p^*$ .											
Private computations											
Alice	Bob										
Choose a secret integer a.	Choose a secret integer b.										
Compute $A \equiv g^a \pmod{p}$ .	Compute $B \equiv g^b \pmod{p}$ .										
Public exchange of values											
Alice sends $A$ to Bob $\longrightarrow$ $A$											
$B \leftarrow$ Bob sends $B$ to Alice											
Further private computations											
Alice	Bob										
Compute the number $B^a \pmod{p}$ .	Compute the number $A^b \pmod{p}$ .										
The shared secret value is $B^a \equiv$	$(g^b)^a \equiv g^{ab} \equiv (g^a)^b \equiv A^b \pmod{p}.$										

Table 2.2: Diffie–Hellman key exchange

Public paran	eter creation										
A trusted party chooses and publishes a large prime p											
and an element $g$ modulo	an element $g$ modulo $p$ of large (prime) order.  Alice Bob  Key creation  key $1 \le a \le p-1$ . $g^a \pmod{p}$ .  blic key $A$ .  Encryption  Choose plaintext $m$ .  Choose random element $k$ .										
Alice	Bob										
Key creation											
Choose private key $1 \le a \le p-1$ .											
Compute $A = g^a \pmod{p}$ .											
Publish the public key $A$ .											
Encryption											
	Choose plaintext $m$ .										
	Choose random element $k$ .										
	Use Alice's public key A										
	to compute $c_1 = g^k \pmod{p}$										
	and $c_2 = mA^k \pmod{p}$ .										
	Send ciphertext $(c_1, c_2)$ to Alice.										
Decryption											
Compute $(c_1^a)^{-1} \cdot c_2 \pmod{p}$ .											
This quantity is equal to $m$ .											

Table 2.3: Elgamal key creation, encryption, and decryption

## Multiplication table modulo 29

	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26	27	28
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
1	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26	27	28
2	0	2	4	6	8	10	12	14	16	18	20	22	24	26	28	1	3	5	7	9	11	13	15	17	19	21	23	25	27
3	0	3	6	9	12	15	18	21	24	27	1	4	7	10	13	16	19	22	25	28	2	5	8	11	14	17	20	23	26
4	0	4	8	12	16	20	24	28	3	7	11	15	19	23	27	2	6	10	14	18	22	26	1	5	9	13	17	21	25
5	0	5	10	15	20	25	1	6	11	16	21	26	2	7	12	17	22	27	3	8	13	18	23	28	4	9	14	19	24
6	0	6	12	18	24	1	7	13	19	25	2	8	14	20	26	3	9	15	21	27	4	10	16	22	28	5	11	17	23
7	0	7	14	21	28	6	13	20	27	5	12	19	26	4	11	18	25	3	10	17	24	2	9	16	23	1	8	15	22
8	0	8	16	24	3	11	19	27	6	14	22	1	9	17	25	4	12	20	28	7	15	23	2	10	18	26	5	13	21
9	0	9	18	27	7	16	25	5	14	23	3	12	21	1	10	19	28	8	17	26	6	15	24	4	13	22	2	11	20
10	0	10	20	1	11	21	2	12	22	3	13	23	4	14	24	5	15	25	6	16	26	7	17	27	8	18	28	9	19
11	0	11	22	4	15	26	8	19	1	12	23	5	16	27	9	20	2	13	24	6	17	28	10	21	3	14	25	7	18
12	0	12	24	7	19	2	14	26	9	21	4	16	28	11	23	6	18	1	13	25	8	20	3	15	27	10	22	5	17
13	0	13	26	10	23	7	20	4	17	1	14	27	11	24	8	21	5	18	2	15	28	12	25	9	22	6	19	3	16
14	0	14	28	13	27	12	26	11	25	10	24	9	23	8	22	7	21	6	20	5	19	4	18	3	17	2	16	1	15
15	0	15	1	16	2	17	3	18	4	19	5	20	6	21	7	22	8	23	9	24	10	25	11	26	12	27	13	28	14
16	0	16	3	19	6	22	9	25	12	28	15	2	18	5	21	8	24	11	27	14	1	17	4	20	7	23	10	26	13
17	0	17	5	22	10	27	15	3	20	8	25	13	1	18	6	23	11	28	16	4	21	9	26	14	2	19	7	24	12
18	0	18	7	25	14	3	21	10	28	17	6	24	13	2	20	9	27	16	5	23	12	1	19	8	26	15	4	22	11
19	0	19	9	28	18	8	27	17	7	26	16	6	25	15	5	24	14	4	23	13	3	22	12	2	21	11	1	20	10
20	0	20	11	2	22	13	4	24	15	6	26	17	8	28	19	10	1	21	12	3	23	14	5	25	16	7	27	18	9
21	0	21	13	5	26	18	10	2	23	15	7	28	20	12	4	25	17	9	1	22	14	6	27	19	11	3	24	16	8
22	0	22	15	8	1	23	16	9	2	24	17	10	3	25	18	11	4	26	19	12	5	27	20	13	6	28	21	14	7
23	0	23	17	11	5	28	22	16	10	4	27	21	15	9	3	26	20	14	8	2	25	19	13	7	1	24	18	12	6
24	0	24	19	14	9	4	28	23	18	13	8	3	27	22	17	12	7	2	26	21	16	11	6	1	25	20	15	10	5
25	0	25	21	17	13	9	5	1	26	22	18	14	10	6	2	27	23	19	15	11	7	3	28	24	20	16	12	8	4
26	0	26	23	20	17	14	11	8	5	2	28	25	22	19	16	13	10	7	4	1	27	24	21	18	15	12	9	6	3
27	0	27	25	23	21	19	17	15	13	11	9	7	5	3	1	28	26	24	22	20	18	16	14	12	10	8	6	4	2
28	0	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1