Introduction Computability & Complexity Polynomial time

Complexity Theory & Philosophy

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Chapter

Complexity Theory

Introduction

Computability & Complexity

Polynomial time

Turing test & Al

Time & Space

Some historical background

- ▶ 85 years ago Gödel introduced incompleteness theorems¹
- 5 years later Turing generalized this to limitations of computation (Halting)²
- ▶ When "computing machines" became reality, time and space requirements of algorithms were of interest ³

¹[Gödel, 1931]

²[Turing, 1936]

³[Hartmanis and Stearns, 1965]: On the comp. complexity of algorithms

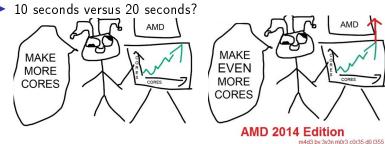
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 - Multiplication vs. prime factoring
 - Reading a book with 400 pages vs. reading all possible 400p books

► Computability Theory: Truncated HALTING problem

Is there an F-proof of S in n or fewer symbols?

- Decidable by exponential bruteforce algorithm
- ► If there was a polynomial (with sane constants and exponent) algorithm Hilbert's dream (almost) possible

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 - ▶ Is it possible to store the decimal expansion somewhere? Would 2^{2²¹²³¹} - 1 not be a known prime (if it is prime) because the universe is too small?

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 Would 2^{2²¹²³¹} 1 not be a known prime (if it is prime)
 because the universe is too small?
 - Maybe the existence of polynomial (n = number of digits) procedure to output the decimal digits of p would explain "knowing"

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Turing test & Al

Turing Test

Simulating humans

Can a computer think?

Time & Space

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The Turing Test I

Evaluator converses with machine that pretends to be human and a real human. If evaluator can't distinguish the two the machine passes the test.

- Additional requirements could be added, e.g.: vision, hearing, handwriting, drawing
- ► The test allows a maximum amount of information to be exchanged, e.g. 2³⁰ bits.

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Argument against "strong Al", presented by John Searle in 1980

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⇒ Passing the Turing test is not a question of computability (as e.g. Roger Penrose argued in *The Emperor's New Mind*) but of complexity

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Chinese Room Argument II

- Chinese room argument initially made to counter "Strong Al" claim
 - Chinese conversation
 - Person in huge library executes moves of the computer program
- ► Nobody would claim that this person understands Chinese, therefore the program cannot give a computer "mind" or "understanding" either
- ► Widely considered a bad argument
- ► Flood of refutations inspired Pat Hayes to quip:

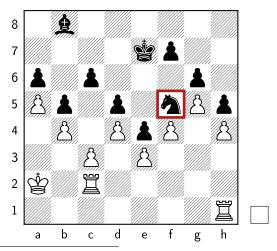
 The field of cognitive science should be redefined as "the ongoing research program of showing Searle's Chinese Room Argument to be false"

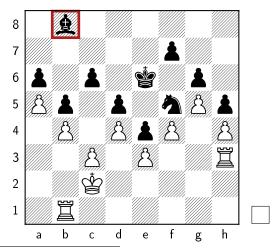
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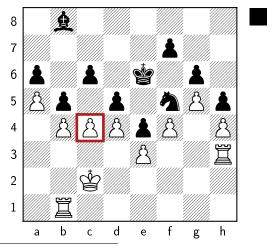
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- ► Can humans solve NP-complete problems efficiently?

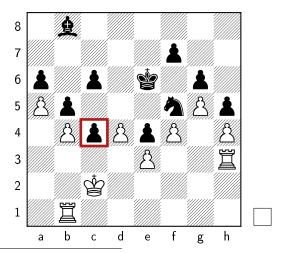
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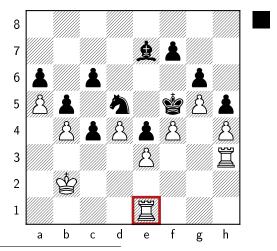
- ► Would a compact & efficient program passing the Turing test be intelligent? If yes, is this possible?
- Can humans solve NP-complete problems efficiently? No
- Humans are usually superior to machines at search problems with higher-level structure, semantics, global patterns or symmetries
 - Proving Fermat's Last Theorem (still took us a while)
 - Pigeonhole Principle (put 1001 objects in 1000 slots)
 - Determining if a chess position is "dead"



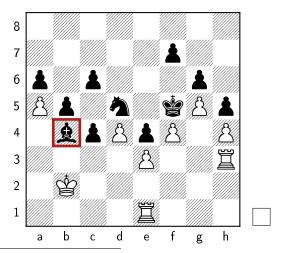




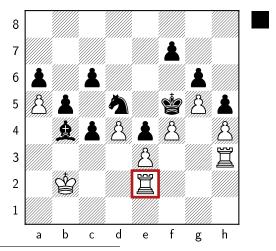




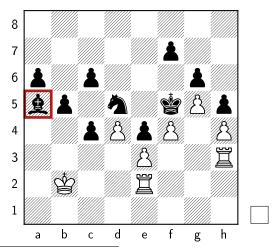
Rybka vs Hikaru Nakamura (ICC, 3+0, 2008)



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Can a computer think?

- Metaphysical question: When does intelligence or understanding arise? Most certainly not when outputting predefined sentences (current state of Siri & Cortana)
 - Underlying theme of the game The Talos Principle (2014)
- Practical question: Are there patterns that a computer can't recognize efficiently? If yes – why? (lower bounds on complexity of algorithms⁴)

⁴Currently most algorithms require exponential time for proofs using pigeon hole principle [Beame and Pitassi, 2001]

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Closed Timelike Curves

CTC Computation

References

Closed Timelike Curves I

- Possibility derived from general relativity by Gödel in 1949
- Since then some other more or less practical ways to achieve CTC discovered, e.g. Tipler cylinder & wormholes
- Famous example:

Time traveller goes back in time to kill his grandfather to prevent existence of time travellers parents. Then time traveller does not exist and can't kill his grandfather.

Possible solutions for grandfather paradox:

Closed Timelike Curves II

- Universe ensures consistency and he somehow fails to kill his grandfather
- ▶ Quantum mechanic solution: S maps quantum state before time travel to quantum state after time travel. Universe finds fixed point, e.g.: time traveller is born with p=1/2 and may or may not be born and kill his grandfather.⁵

⁵[Deutsch, 1991]

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- ► No logical paradox
- ▶ But one of computational complexity: Knowledge comes into existence without causal process (*Evolutionary Principle*)

Closed Timelike Curve Computation [Brun, 2003]

- ▶ Take any NP-complete problem p and encode all possible solutions s to binary representation $E: S \rightarrow \{0, 1, ..., |S-1|\}$
- Let f(E(s)) = 1 iff s is a valid solution for p (runs in polynomial time).
- ▶ Run the following program inside a CTC with it's own output:
 - ▶ If f(x) = 1 then output x
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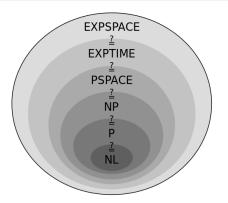
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- ► PSPACE would be the largest class solvable independent from classical or quantum computation

Complexity Classes



$$LSPACE \subseteq P \subseteq NP \subseteq PSPACE \subseteq EXPTIME \subseteq EXPSPACE$$

$$P \subseteq PSPACE$$

Image from Wikipedia

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