Meltdown and Spectre Samples

Written in Assembly

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1 Introduction

1.1 Overview

TBD

1.2 Conventions

1.2.1 Introduction

In this section we define some convention that are specific for this document.

1.2.2 Data Sections

The data is divided into three parts: read-only data, initialized data and uninitialized data. Code chunks with this type of data will all have defined sufficies.

Definition 1 Read-only data is data that is not modified during program execution. The suffix for read-only data is **-rodata**.

Definition 2 Initialized data is data that is changeable during program execution. The data is already initialized with data when the program starts. The suffix for initialized data is **-idata**.

Definition 3 Uninitialized data is data that is changeable during program execution. The data is not initialized. The suffix for uninitialized data is **-udata**.

1.3 Nasm

2.1 Introduction

TBD

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2.2 Detect Cache Access Time

2.2.1 High Resolution Timer

First we need a high resolution timer to determine the cache access time. For this we use the time stamp counter. The time stamp counter is monotonically incrementing. When reading the time stamp counter (with rdtsc) the result is delivered back in the registers EDX and EAX forming a 64bit value. The time stamp counter is not an absolute value but a relative value, meaning that you cannot (easily) calculate from the time stamp counter to some time units (e.g. ns). But this is no problem as we only want to measure relative times.

To retrieve a 64bit value for the time we shift the value in EDX 32 bits to the left and add the value of EAX to this.

$$\langle tsc\text{-}64bit \ 7 \rangle \equiv$$
 (8a)

rdtsc

shl RDX,32

add RAX,RDX

2.2.2 Cache Access Time Routine

Next we need a routine that calculates the cache access time for us.

First we have to ensure in this routine that the speculative execution of the processor does not interfere with our time measurement. For this we use the instruction lfence which ensures that all previous reads are done before executing the next instructions.

Next we access a memory location with the address RDI by loading this into RCX and measure the time before and after the access.

The command lfence before reading the time stamp counter is needed because we have to ensure that all reads before the time measurements are done.

At last we calculate the relative time needed to access the memory location. In theory we should see a difference whether the memory location is accessed before or not.

Parameters

RDI the address of the memory which is loaded either from the cache or from memory

Return

RAX the relative time of the cache access

```
\langle calculate\text{-}cache\text{-}access\text{-}time \ 8a \rangle \equiv
8a
                                                                                                               (11f 18c 21a 23c 31b)
            _calccachetime:
                     lfence
            \langle tsc-64bit 7 \rangle
                     mov
                                       R8, RAX
                     mov
                                       RCX, [RDI]
                     lfence
            \langle tsc-64bit 7 \rangle
                     sub
                                       RAX,R8
                     ret
         Defines:
```

2.3 Measure Cache Access Time

_calccachetime, used in chunks 10b, 11b, and 15b.

2.3.1 **Setup**

To measure the cache timing we create a standalone program that shows us the time for a cached and for an uncached memory access.

First we need some area in memory with data which we can later read from. This data area goes into the area .bss which contains uninitialized data. We align the data at a page boundary and reserve one pages for our data.

```
\langle data\text{-}udata \ 8b \rangle \equiv
                                                                                              (11f 18c 21a 23c 31b)
8b
                  alignb
                                        pagesize
                  data:
                                        resb pagesize
        Defines:
          data, used in chunks 9, 11b, 14b, 17a, 18a, 22c, 23b, and 27.
        Uses pagesize 5.
          From time to time we need a small scratch area so we define an area with 32 bytes.
        \langle scratch\text{-}udata \ 8c \rangle \equiv
                                                                                              (11f 18c 21a 23c 31b)
8c
                  scratch:
                                        resb 32
        Defines:
          scratch, used in chunks 10f, 11d, 17, 18a, 23b, 28, and 29d.
```

The program begins with the label _start.

9a ⟨cachetiming-program 9a⟩≡
_start:
(11f) 9f⊳

Uses _start 5.

Now we start with initialising the data area with some random data. For this we load RDI with the address of the data area.

9b $\langle init\text{-}random\text{-}data \text{ 9b} \rangle \equiv$ (9f 14a 21b 24b) 9c \triangleright

mov RDI, data

Uses data 8b.

9c

9e

Next we load the number of bytes to fill into RSI. For this we load the pagesize into RSI.

 $\langle init\text{-}random\text{-}data 9b\rangle + \equiv$ (9f 14a 21b 24b) \triangleleft 9b 9d \triangleright

mov RSI,pagesize

Uses pagesize 5.

At last we load EDX with some random seed. For this we use rdtsc and only use the lower 32 bit of the value.

9d $\langle init\text{-}random\text{-}data \text{ 9b} \rangle + \equiv$ (9f 14a 21b 24b) \triangleleft 9c 9e \triangleright rdtsc mov EDX,EAX

Now we call _xorshift to fill the data area.

 $\langle init\text{-random-data 9b}\rangle + \equiv$ (9f 14a 21b 24b) \triangleleft 9d

call _xorshift

Uses _xorshift 36a.

Now we add this data initialization to our program.

9f $\langle cachetiming\text{-}program 9a \rangle + \equiv$ (11f) $\triangleleft 9a 9g \triangleright$ $\langle init\text{-}random\text{-}data 9b \rangle$

2.3.2 Measure Time

Uses data 8b.

Now that we have setup our data area we can now cache data from the first page by loading it into a register which also loads this into the cache.

For this we load RDI with the address of the data area.

9g $\langle cachetiming-program 9a \rangle + \equiv$ (11f) \triangleleft 9f $10a \triangleright$

mov RDI,data

Before we load the data into a register now we will clear the cache lines with the given address. For this we use the instruction clflush. After flushing the cache line we ensure (with 1fence) that all reads from memory are finished before we load the data into a register again (and filling the cache).

```
\langle cachetiming-program 9a \rangle + \equiv
10a
                                                                                                   (11f) ⊲9g 10b⊳
                  clflush
                                  [RDI]
                  lfence
                                 RCX, [RDI]
                  mov
```

Now we can determine the time that is needed to load this data once again. We do not need to load RDI again because it has not changed.

```
10b
         \langle cachetiming-program 9a \rangle + \equiv
                                                                                                     (11f) ⊲ 10a 10e⊳
                   call
                                   _calccachetime
```

Uses _calccachetime 8a.

Now we have the relative cache access time in register RAX. We store this value to the stack and print out an explaining text.

For this we define the text to print.

```
\langle cachetiming-rodata \ 10c \rangle \equiv
10c
                                                                                                               (11f) 11c⊳
            \langle common-rodata \ 10d \rangle
                    scached:
                                           db "Cached Access Time: ",0x00
```

scached, used in chunk 10e.

Additionally we define some helper data, in this case line feed (LF).

```
\langle common-rodata \ 10d \rangle \equiv
10d
                                                                                                          (10c 18c 21a 23c 31b)
                     slf:
                                              db 0x0a
```

Defines:

Defines:

slf, used in chunks 11, 17, 18a, 23b, and 30.

Now we can store RAX and print the text.

```
\langle cachetiming\text{-}program 9a \rangle + \equiv
10e
                                                                                                        (11f) ⊲ 10b 10f⊳
                    push
                                    RAX
                                    RDI, scached
                    mov
                    call
                                    _print
```

Uses _print 38a and scached 10c.

We now restore the value and print the measured time to stdout.

```
10f
        \langle cachetiming-program 9a \rangle + \equiv
                                                                                                (11f) ⊲ 10e 11a⊳
                  pop
                  mov
                                RSI, scratch
                                _printdu64bit
                  call
```

Uses _printdu64bit 39a and scratch 8c.

At last we append a LF to the output.

11a $\langle cachetiming\text{-}program \text{ 9a} \rangle + \equiv$ (11f) \triangleleft 10f 11b> mov RSI,slf mov RDI,1 call _nprint

Uses _nprint 37b and slf 10d.

Now we do the same with an uncached value. The difference is that we do not load the value before.

11b $\langle cachetiming\text{-}program \ 9a \rangle + \equiv$ (11f) \triangleleft 11a 11d \triangleright mov RDI,data clflush [RDI] lfence call _calccachetime

Uses _calccachetime 8a and data 8b.

Now we have the time of the uncached data access in RAX and can print it out with some explaining text.

11c $\langle cachetiming\text{-}rodata \ 10c \rangle + \equiv$ (11f) $\triangleleft 10c$ suncached: db "Uncached Access Time: ",0x00

Defines:

suncached, used in chunk 11d.

```
\langle cachetiming-program 9a \rangle + \equiv
11d
                                                                                      (11f) ⊲11b 11e⊳
                             RAX
                push
                mov
                             RDI, suncached
                call
                             _print
                             RDI
                pop
                             RSI, scratch
                mov
                call
                             _printdu64bit
                             RSI,slf
                mov
                             RDI,1
                mov
```

Uses _nprint 37b, _print 38a, _printdu64bit 39a, scratch 8c, slf 10d, and suncached 11c.

At last we exit the program.

_nprint

call

```
11e \langle cachetiming\text{-}program 9a \rangle + \equiv (11f) \langle exitProgram 35b \rangle
```

Now we can put everything together and have our **cachetiming** program that we can now execute.

```
11f \langle cachetiming.asm \ 11f \rangle \equiv \langle preamble \ 5 \rangle
```

```
section .rodata \langle cachetiming\text{-}rodata \ 10c \rangle
section .bss \langle data\text{-}udata \ 8b \rangle
\langle scratch\text{-}udata \ 8c \rangle
section .text \langle cachetiming\text{-}program \ 9a \rangle
\langle calculate\text{-}cache\text{-}access\text{-}time \ 8a \rangle
\langle xorshift\text{-}prng \ 36a \rangle
\langle utilities \ 35a \rangle
```

The program source is placed in asm/. With make in the folder we can create an executable which is moved to bin/. There we can execute this program.

```
$ bin/cachetiming
Cached Access Time: 72
Uncached Access Time: 372
$
```

2.4 Read Byte via Cache Access Time

2.4.1 Introduction

We have seen that we can determine if the content of a memory address is in the cache or not (see 2.3 Measure Cache Access Time).

So next we try to read a single byte from the memory by only detecting the cache access time.

2.4.2 Clear Cache for Measurement

Before we can determine the cache access times we need to clear the cache. We define a subroutine for this.

Parameters

```
RDI the address of the probe memory

RSI the step size in the probe memory
```

```
⟨clearcache 13a⟩≡
                                                                                       (18c 21a 23c 31b)
13a
          _clearcache:
                mov
                              RCX,256
                cld
          .nextflush:
                clflush
                              [RDI]
                              RDI, RSI
                add
                loop
                              .nextflush
                lfence
                ret
        Defines:
          _clearcache, used in chunks 13b, 22a, and 25c.
          Now we add this to our program.
13b
        \langle cacheread by te-program \ 13b \rangle \equiv
                                                                                        (18c 21a) 14a⊳
                              RDI, probe
                mov
                              RSI, pagesize
                mov
                call
                              _clearcache
        Uses \_clearcache 13a, pagesize 5, and probe 13c.
```

2.4.3 Indexed Array Access

To read the value of a byte via the cache we use the byte to index into a probe array and then determine the cache access times of this probe array.

For this we will first create a probe array.

```
13c \langle probe\text{-}udata \ 13c \rangle \equiv (18c 21a 23c 31b) alignb pagesize probe times 256 resb pagesize

Defines: probe, used in chunks 13, 14c, 16c, 22c, and 27a.
Uses pagesize 5.
```

Next we will fill this **probe** array with some random data (similar to the chunks for data 9b, 9c, 9d and 9e).

```
13d ⟨init-random-probe 13d⟩≡ (14a 21b 24b)

mov RDI,probe

mov RSI,pagesize

shl RSI,8

rdtsc

mov EDX,EAX

call _xorshift
```

Uses $\mbox{\tt_xorshift}\ 36a,\ \mbox{\tt pagesize}\ 5,\ and\ \mbox{\tt probe}\ 13c.$

Now we add the initialization of the data and probe area to the program.

```
14a
         \langle cachereadbyte-program \ 13b \rangle + \equiv
                                                                                            (18c 21a) ⊲13b 14b⊳
           _start:
           ⟨init-random-data 9b⟩
           \langle init\text{-}random\text{-}probe 13d \rangle
         Uses _start 5.
           Now we can read a byte from data into AL.
14b
         \langle cachereadbyte-program \ 13b \rangle + \equiv
                                                                                            (18c 21a) ⊲14a 14c⊳
                  mov
                                 RDI, data
                                 RAX, RAX
                  xor
                                 AL, [RDI]
                  mov
         Uses data 8b.
            We use the value in RAX to access the probe array.
14c
         \langle cachereadbyte-program \ 13b \rangle + \equiv
                                                                                            (18c 21a) ⊲14b 16c⊳
                  mov
                                 RDX, pagesize
                  mul
                                 RDX
                  mov
                                 RSI, probe
                                 AL, [RSI+RAX]
                  mov
```

Now we read the datum back via the cache access times. For this we create subroutines.

2.4.4 Read a Byte from the Cache

Uses pagesize 5 and probe 13c.

First we create a subroutine to read the cache access timings for the probe area.

Parameters

```
RDI the address of the probe memory

RSI the step size in the probe memory

RDX an area to keep the detected cache access times (256 * 8 bytes)

14d \( \langle \text{readcachetiming 14d} \rangle \subseteq \text{ (18c 21a 23c 31b) 15ab} \)
\[ \text{_readcachetiming:} \( \langle \text{enterstackframe 35c} \rangle \)

Defines:
\[ \text{_readcachetiming, used in chunks 16c, 22c, and 26b.} \]
```

Now we create space on the stack to keep the variables. Next we save the parameters to the stack space created.

```
15a ⟨readcachetiming 14d⟩+≡ (18c 21a 23c 31b) ⊲14d 15b⊳

sub RSP,32

mov [RBP-8],RDI

mov [RBP-16],RSI

mov [RBP-24],RDX
```

Now we can start detecting the cache access times.

```
\langle readcachetiming 14d \rangle + \equiv
15b
                                                                        (18c 21a 23c 31b) ⊲15a 15c⊳
                mov
                            RCX,256
          .nextcacheread:
                mov
                             [RBP-32], RCX
                             _calccachetime
                call
                            RDX, [RBP-24]
                mov
                mov
                             [RDX], RAX
                            RDX,8
                add
                             [RBP-24],RDX
                mov
                            RDI, [RBP-8]
                mov
                add
                            RDI, [RBP-16]
                mov
                             [RBP-8], RDI
                            RCX, [RBP-32]
                mov
```

 $\begin{array}{c} \text{loop} & . \\ \text{Uses _calccachetime } 8a. \end{array}$

At the end we clean up the stack again and return to the caller.

.nextcacheread

```
15c ⟨readcachetiming 14d⟩+≡ (18c 21a 23c 31b) ⊲15b ⟨leavestackframe 35d⟩
```

After we determined all cache access times we can now find the lowest access time and with this the possible byte. We return two results from this subroutine, in AL the byte with the lowest cache access time and in AH the count of the lowest cache access time. Only if AH is 1 then the value in AL is valid.

Parameters

RDI the area with the detected cache access times (256 * 8 bytes)

Return

AL the read byte (in AL) with the lowest cache access time

AH the number of bytes read with the lowest cache access time

```
16a
       \langle analyze cachemintiming 16a \rangle \equiv
                                                                                     (18c 31b)
         _analyzecachetiming:
               push
                           RDI
                           R8,0xffffffffffffff
               mov
               xor
                           R9, R9
                           RCX, RCX
               xor
                           RSI, RDI
               mov
         .nexttry:
               lodsq
               cmp
                           RAX, R8
                           .nohit
               ja
                           R8,RAX
               mov
                           R9,RCX
               mov
         .nohit:
               inc
                           RCX
                           RCX,256
               cmp
               jb
                           .nexttry
               xor
                           RCX, RCX
                           RSI
               pop
         .nextcount:
               lodsq
                           RAX,R8
               cmp
                           .nomin
               ja
                           R10
               inc
         .nomin:
               inc
                           RCX
               cmp
                           RCX,256
                           .nextcount
               jb
                           RAX,R10
               mov
                           RAX,8
               shl
                           AL,R9b
               mov
               ret
```

2.4.5 The Whole Program to Read a Byte from Cache

Before we can start using our new subroutine **_readcachetiming** we need to define a data area for the cache access times.

```
16b \langle timings\text{-}udata \text{ 16b} \rangle \equiv (18c 21a 23c 31b) timings resq 256
```

Now we have all subroutines together we now can start implementing the main program and output the byte read.

```
16c \langle cachereadbyte-program \ 13b \rangle + \equiv (18c 21a) \langle 14c \ 17b \rangle
```

```
mov RDI,probe
mov RSI,pagesize
mov RDX,timings
call _readcachetiming
mov RDI,timings
call _analyzecachetiming
Uses _readcachetiming 14d, pagesize 5, and probe 13c.
```

Now we define a string to output for the read byte and the expected byte.

```
17a \langle cachereadbyte\text{-}rodata\ 17a \rangle \equiv (18c 21a 23c) sreadbyte: db "Byte read via cache access: ",0x00 ssountbyte: db "Count of bytes with min timing: ",0x00 sexpectedbyte: db "Expected byte from data: ",0x00 Uses data 8b.
```

We save the value from RAX (only AL is interesting to us) to the stack and print out the text.

```
17b \langle cachereadbyte\text{-}program \ 13b \rangle + \equiv (18c 21a) \triangleleft 16c 17c \triangleright push RAX mov RDI, sreadbyte call _print
```

Uses _print 38a.

Now we print the read byte and end the line with a LF.

```
\langle cachereadbyte-program \ 13b \rangle + \equiv
                                                                                   (18c 21a) ⊲17b 17d⊳
17c
                 pop
                              RDI
                              RDI
                 push
                              RDI, Oxff
                 and
                              RSI, scratch
                 mov
                              _printh8bit
                 call
                              RDI,1
                 mov
                              RSI,slf
                 mov
                 call
                              _nprint
```

Uses _nprint 37b, _printh8bit 41b, scratch 8c, and slf 10d.

Next we print (for information) the number of bytes read with the minimum cache access timing.

```
17d ⟨cachereadbyte-program 13b⟩+≡ (18c 21a) ⊲17c 18a⊳

mov RDI,ssountbyte

call _print

pop RDI

shr RDI,8

and RDI,0xff

mov RSI,scratch
```

```
call _printdu64bit
mov RDI,1
mov RSI,slf
call _nprint
```

Uses $_nprint 37b$, $_print 38a$, $_printdu64bit 39a$, scratch 8c, and slf 10d.

Now we read the byte from the original data array and print this also.

```
18a
        \langle cacheread by te-program \ 13b \rangle + \equiv
                                                                                 (18c 21a) ⊲17d 18b⊳
                mov
                             RDI, sexpected byte
                             _print
                call
                             RSI, data
                mov
                             RAX, RAX
                xor
                             AL, [RSI]
                mov
                             RDI, RAX
                mov
                mov
                             RSI, scratch
                             _printh8bit
                call
                             RDI,1
                mov
                             RSI,slf
                mov
                call
                             _nprint
```

Uses _nprint 37b, _print 38a, _printh8bit 41b, data 8b, scratch 8c, and slf 10d.

At last we exit the program.

```
18b \langle cachereadbyte\text{-}program \ 13b \rangle + \equiv  (18c 21a) \langle 18a \rangle + \equiv  (18c 21a) \langle 18a \rangle + \equiv
```

Now we put all together to get the program cachereadbyte that we can execute.

```
18c \langle cachereadbyte.asm \ 18c \rangle \equiv \langle preamble \ 5 \rangle

section .rodata \langle common\text{-}rodata \ 10d \rangle \langle cachereadbyte\text{-}rodata \ 17a \rangle

section .bss \langle data\text{-}udata \ 8b \rangle \langle probe\text{-}udata \ 13c \rangle \langle scratch\text{-}udata \ 8c \rangle \langle timings\text{-}udata \ 16b \rangle

section .text \langle cachereadbyte\text{-}program \ 13b \rangle
```

 $\langle clearcache 13a \rangle$

```
\langle calculate\text{-}cache\text{-}access\text{-}time 8a \rangle
\langle readcachetiming 14d \rangle
\langle analyzecachemintiming 16a \rangle
\langle xorshift\text{-}prng 36a \rangle
\langle utilities 35a \rangle
```

2.4.6 Improve Cache Access Time Analysis

As we can see – when running the program cachereadbyte – the result is not always as clear as it could be. Simply getting the lowest cache access time seems not to be enough. Sample outputs of the program are

\$ bin/cachereadbyte

```
Byte read via cache access:
                                 2b
Count of bytes with min timing: 1
Expected byte from data:
                                 2h
$ bin/cachereadbyte
Byte read via cache access:
                                 ff
Count of bytes with min timing: 11
Expected byte from data:
                                 b3
$ bin/cachereadbyte
Byte read via cache access:
                                 2f
Count of bytes with min timing: 1
Expected byte from data:
```

So we have to improve our cache time detection routine. We will change the implementation of the chunk 16a to define a threshold that is a little bit above the min access time and run the cache detection routine multiple times if no clear result is returned.

First start with the subrotuine to analyze the cache access timing. We define a threshold 25~% above the minimum cache access time.

First we search for the minimum cache access time.

Parameters

RDI the area with the detected cache access times (256 * 8 bytes)

Return

AL the first byte (in AL) with a cache access time below the threshold

```
AΗ
                    the number of bytes read with a cache access time below the threshold
20a
        \langle analyze caches impthrestiming 20a \rangle \equiv
                                                                                     (21a 23c) 20b⊳
          _analyzecachetiming:
                            RDI
                push
                            R8,0xffffffffffffff
                mov
                            RCX,RCX
                xor
                mov
                            RSI,RDI
          .nextmin:
                lodsq
                cmp
                            RAX,R8
                             .nonewmin
                ja
                mov
                            R8,RAX
          .nonewmin:
                inc
                            RCX
                cmp
                            RCX,256
                jb
                             .nextmin
          Now we have the minimum cache access time in R8. Next we will add \frac{1}{4} to this to
       have our threshold.
        \langle analyze caches impthrestiming 20a \rangle + \equiv
20b
                                                                               (21a 23c) ⊲20a 20c⊳
                            RAX,R8
                mov
                            RAX,4
                shr
                add
                            R8, RAX
          Now we scan the cache access times a second time and take all values below the
        threshold into account.
        \langle analyze caches impthrestiming 20a \rangle + \equiv
20c
                                                                                     (21a 23c) ⊲20b
                            RSI
                pop
                xor
                            RCX, RCX
                xor
                            R9, R9
          .nextbyte:
                lodsq
                            RAX,R8
                cmp
                ja
                             .nonewbyte
                inc
                            R9
                            R10,RCX
                mov
          .nonewbyte:
                            RCX
                inc
                            RCX,256
                cmp
                             .nextbyte
                jb
```

mov shl

mov ret RAX,R9

RAX,8 AL,R10b Now we put all together to get the program cachereadbyte2 that we can execute.

```
\langle cachereadbyte2.asm\ 21a \rangle \equiv
   \langle preamble 5 \rangle
   section .rodata
   \langle common-rodata \ 10d \rangle
   \langle cachereadbyte-rodata 17a \rangle
   section .bss
   \langle data-udata \ 8b \rangle
   \langle probe-udata \ 13c \rangle
   ⟨scratch-udata 8c⟩
   \langle timings-udata \ 16b \rangle
   section .text
   \langle cacheread by te-program 13b \rangle
   ⟨clearcache 13a⟩
   \langle calculate\text{-}cache\text{-}access\text{-}time \ 8a \rangle
   \langle readcachetiming 14d \rangle
   \langle analyze caches impthrestiming 20a \rangle
   \langle xorshift\text{-}prng 36a \rangle
   \langle utilities 35a \rangle
```

21a

Now when we only find a single hit then the possibility that the byte from the cache timing is the original byte is much higher.

Next we will create a program that tries to read the value from the cache until we have a single result.

First we initialize our data and probe areas.

```
21b ⟨cachereadbyte3-program 21b⟩≡
_start:
⟨init-random-data 9b⟩
⟨init-random-probe 13d⟩
Uses _start 5.
```

Next we create a subroutine that clears the cache and reads in a byte via the probe array.

Parameters

```
RDI
                    the address of the byte to read
       RSI
                    the address of the probe memory
       RDX
                    the step size in the probe memory
        \langle readbyte2cache 22a \rangle \equiv
22a
                                                                                           (23c) 22b⊳
          _readbyte2cache:
                push
                             RDI
                push
                             RSI
                push
                             RDX
                mov
                             RDI, RSI
                             RSI, RDX
                mov
                call
                             _clearcache
       Defines:
          _readbyte2cache, used in chunk 22c.
        Uses _clearcache 13a.
```

Next we can add the read of the byte and caching the data from the probe array.

```
\langle readbyte2cache\ 22a\rangle + \equiv
22b
                                                                                                   (23c) ⊲22a
                                RDX
                  pop
                                RSI
                  pop
                                RDI
                  pop
                                RAX, RAX
                  xor
                                AL, [RDI]
                  mov
                                RDX
                  mul
                                AL, [RSI+RAX]
                  mov
```

Now we add the call to this subroutine to our program and determine the byte by analyzing the cache access times.

```
22c
       \langle cachereadbyte3-program 21b \rangle + \equiv
                                                                                 (23c) ⊲21b 23a⊳
          .startreadcache:
                            RDI, data
               mov
                            RSI, probe
               mov
                            RDX, pagesize
               mov
               call
                            _readbyte2cache
                            RDI, probe
               mov
                            RSI, pagesize
               mov
                           RDX, timings
               mov
                            _readcachetiming
               call
               mov
                           RDI, timings
               call
                            _analyzecachetiming
```

Uses $_$ readbyte2cache 22a, $_$ readcachetiming 14d, data 8b, pagesize 5, and probe 13c.

Now we check if the read byte was a single byte, else we will do this again.

```
\langle cachereadbyte3-program 21b \rangle + \equiv
23a
                                                                                             (23c) ⊲22c 23b⊳
                  cmp
                                AH,1
                                .startreadcache
                  ja
           Now we print out our result.
23b
         \langle cachereadbyte3-program 21b \rangle + \equiv
                                                                                                    (23c) ⊲23a
                  push
                                RAX
                  mov
                                RDI, sreadbyte
                  call
                                _print
                                RDI
                  pop
                                RDI, Oxff
                  and
                                RSI, scratch
                  mov
                  call
                                _printh8bit
                  mov
                                RDI,1
                                RSI,slf
                  mov
                  call
                                _nprint
                                RDI, sexpected byte
                  mov
                  call
                                _print
                                RSI, data
                  mov
                                RAX, RAX
                  xor
                                AL, [RSI]
                  mov
                                RDI, RAX
                  mov
                                RSI, scratch
                  mov
                  call
                                _printh8bit
                                RDI,1
                  mov
                                RSI,slf
                  mov
                                _nprint
                  call
           \langle exitProgram 35b \rangle
         Uses _nprint 37b, _print 38a, _printh8bit 41b, data 8b, scratch 8c, and slf 10d.
         \langle cachereadbyte3.asm\ 23c \rangle \equiv
23c
           \langle preamble 5 \rangle
           section .rodata
           ⟨common-rodata 10d⟩
           ⟨cachereadbyte-rodata 17a⟩
           section .bss
           ⟨data-udata 8b⟩
           \langle probe-udata \ 13c \rangle
           \langle scratch\text{-}udata \ 8c \rangle
           \langle timings-udata \ 16b \rangle
```

```
section .text
\(\alpha cachereadbyte3-program 21b\rangle\)
\(\alpha readbyte2cache 22a\rangle\)
\(\alpha clearcache 13a\rangle\)
\(\alpha calculate-cache-access-time 8a\rangle\)
\(\alpha readcachetiming 14d\rangle\)
\(\alpha analyzecachesimpthrestiming 20a\rangle\)
\(\alpha vorshift-prng 36a\rangle\)
\(\alpha utilities 35a\rangle\)
```

Even if this program is not perfect because it is not reliable all the time it is reliable enough to demonstrate the next steps.

2.5 Read Array via Cache Access Time

2.5.1 Introduction

Now we have read a byte via the cache access times. Now it is time to read a complete memory area.

2.5.2 **Setup**

For this we use the data defined before and read in the complete area. For this we need additionally a memory area that holds the read data.

```
\langle readback\text{-}udata 24a \rangle \equiv
                                                                                                              (31b)
24a
                   alignb
                                         pagesize
                   readbackdata
                                        resb pagesize
         Defines:
           readbackdata, used in chunk 27.
         Uses pagesize 5.
           First we initialize the data and probe areas in our program with some random data.
24b
         \langle cacheread\text{-}program 24b \rangle \equiv
                                                                                                        (31b) 27a⊳
           ⟨init-random-data 9b⟩
           \langle init-random-probe 13d\rangle
         Uses _start 5.
```

Next we will define a subroutine that reads the data area and writes the results of the cache read into readbackdata.

Parameters

Uses _clearcache 13a.

25a

```
RDI
            the address of the data memory
            the size of the data memory
RSI
            the address of the probe memory
RDX
RCX
            the step size in the probe memory (the probe area needs to be at least
            256 * RCX bytes in size)
R8
            the address of the readback area (must be at least the same size as the data
            area)
R9
            the address of the the area to keep the timing data (at least 256 * 8 bytes)
\langle readarea \ 25a \rangle \equiv
                                                                                 (31b) 25b ⊳
  _readarea:
Defines:
  _readarea, used in chunk 27a.
```

Now we create some place on the stack and store the parameters on it. We reserve an extra place at [RBP-56] for a counter into the data memory.

```
\langle readarea 25a \rangle + \equiv
25b
                                                                                          (31b) ⊲25a 25c⊳
          \langle enterstack frame 35c \rangle
                               RSP,56
                 sub
                               [RBP-8],RDI
                 mov
                               [RBP-16],RSI
                 mov
                 mov
                               [RBP-24],RDX
                               [RBP-32], RCX
                 mov
                 mov
                               [RBP-40],R8
                               [RBP-48],R9
                 mov
                               RAX, RAX
                 xor
                               [RBP-56], RAX
                 mov
```

First we have to clear the cache before we can measure any cache access times.

```
25c ⟨readarea 25a⟩+≡
.startread:
mov RDI,[RBP-24]
mov RSI,[RBP-32]
call _clearcache (31b) ⊲25b 26a⊳
```

Now we can load the byte from the memory and cache the according value from the probe memory.

```
\langle readarea \ 25a \rangle + \equiv
26a
                                                                                         (31b) ⊲25c 26b⊳
                 mov
                               RSI, [RBP-8]
                 add
                               RSI, [RBP-56]
                               RAX, RAX
                 xor
                               AL, [RSI]
                 mov
                               RDX, [RBP-32]
                 mov
                              RDX
                 mul
                              RSI, [RBP-24]
                 mov
                               AL, [RSI+RAX]
                 mov
```

Now that we have filled our cache we can determine the cache access times.

```
26b ⟨readarea 25a⟩+≡
mov RDI,[RBP-24]
mov RSI,[RBP-32]
mov RDX,[RBP-48]
call _readcachetiming (31b) <26a 26c>
```

Uses _readcachetiming 14d.

Now we can analyze the cache access times.

```
26c \langle readarea\ 25a \rangle + \equiv (31b) \triangleleft 26b 26d \triangleright mov RDI,[RBP-48] call _analyzecachetiming
```

If we have more than 1 hit then we retry the reading of the byte.

```
26d \langle readarea\ 25a \rangle + \equiv (31b) \triangleleft 26c\ 26e \triangleright cmp AH,1 ja .startread
```

Now that we found a byte we store it in the resulting memory area.

```
\langle readarea 25a \rangle + \equiv
26e
                                                                                       (31b) ⊲26d 26f⊳
                              RDI, [RBP-40]
                mov
                mov
                              RCX, [RBP-56]
                              RDI, RCX
                add
                              [RDI],AL
                mov
                              RCX
                inc
                              [RBP-56],RCX
                mov
                              RCX, [RBP-16]
                 cmp
                 jb
                              .startread
```

Now we clean up the stack frame and return to the caller.

```
26f \langle readarea\ 25a \rangle + \equiv (31b) \triangleleft 26e \langle leavestackframe\ 35d \rangle ret
```

Now we can add this to our program.

```
\langle cacheread\text{-}program 24b \rangle + \equiv
27a
                                                                                           (31b) ⊲24b 27b⊳
                  mov
                               RDI, data
                               RSI, pagesize
                  mov
                               RDX, probe
                  mov
                               RCX, pagesize
                  mov
                               R8, readbackdata
                  mov
                               R9, timings
                  mov
                  call
                                _readarea
        Uses _readarea 25a, data 8b, pagesize 5, probe 13c, and readbackdata 24a.
           TBD
27b
        \langle cacheread\text{-}program 24b \rangle + \equiv
                                                                                            (31b) ⊲27a 27c⊳
                  mov
                               RDI.data
                               RSI, readbackdata
                  mov
                               _printcompare
                  call
        Uses _printcompare 31a, data 8b, and readbackdata 24a.
           At the end we end the program.
         \langle cacheread\text{-}program 24b \rangle + \equiv
27c
                                                                                                  (31b) ⊲27b
           \langle exitProgram 35b \rangle
           TBD
27d
        \langle cacheread\text{-}rodata \ 27d \rangle \equiv
                                                                                                        (31b)
                  sbgred:
                                      db 0x1b,"[1;41m",0x00
                                      db 0x1b,"[0m",0x00
                  sresetstyle:
                  sseparator:
                                      db "- ",0x00
                                      db " "
                  sblank:
                                      db "
                  semptybyte:
                                                ",0x00
        Defines:
           sbgred, used in chunk 29d.
           sblank, used in chunks 28e and 29d.
           sresetstyle, used in chunks 29d and 30.
           sseparator, used in chunk 29b.
```

First we define a subroutine which prints out up to 16 bytes each side by side on the screen. If two bytes in the arrays are different then the value at the right side (from the second array) will be printed with read background.

Parameters

RDI the address of the first array

RSI the address of the second array

RDX number of bytes to print (up to 16). If the value is above 16 then nothing is printed out

```
28a \langle print\text{-}comparision 16 \ 28a \rangle \equiv (31b) 28c \triangleright _printcompare16:
```

Defines:

_printcompare16, used in chunk 31a.

Additionally we need some scratch area for the printing.

28b $\langle cacheread\text{-}udata 28b \rangle \equiv$

scratch: resb 64

Uses scratch 8c.

At the start of the subroutine we prepare a stack frame for further operations as we will need to save and restore the registers RDI, RSI, RDX and RCX multiple times. Additionally we store R12 to the stack to use this register as scratch register.

```
28c
        \langle print\text{-}comparision16 \ 28a \rangle + \equiv
                                                                                            (31b) ⊲28a 28d⊳
                 push
                                RBP
                                RBP, RSP
                 mov
                 sub
                                RSP,32
                                [RBP-8],RDI
                 mov
                                [RBP-16],RSI
                 mov
                 mov
                                [RBP-24],RDX
                 push
                                R12
```

Now we first start and check that no more than 16 bytes should be printed, otherwise we will end the subroutine immediately.

```
28d \langle print\text{-}comparision16 \ 28a \rangle + \equiv (31b) \triangleleft 28c 28e \triangleright cmp RDX,0x10 ja .done
```

Next we can start and handle the "left" side of the output. We output up to 16 bytes and then continue at .leftbytesdone (29a).

```
28e \langle print\text{-}comparision16 \ 28a \rangle + \equiv (31b) \triangleleft 28d 29a \triangleright xor RCX,RCX .nextbyteleft: cmp RCX,RDX
```

```
.leftbytesdone
jae
           AL, [RDI+RCX]
mov
           AH, AH
xor
           DI, AX
mov
           RSI, scratch
mov
           _printh8bit
call
mov
           RDI,1
           RSI, sblank
mov
call
           _nprint
```

mov

[RBP-32], RCX

```
RDI, [RBP-8]
mov
           RDX, [RBP-24]
mov
           RCX, [RBP-32]
mov
inc
           RCX
jmp
           .nextbyteleft
```

Uses _nprint 37b, _printh8bit 41b, sblank 27d, and scratch 8c.

Now we fill up the space so that the space of 16 bytes is occupied.

29a $\langle print\text{-}comparision16 28a \rangle + \equiv$ (31b) ⊲28e 29b⊳

.leftbytesdone:

```
RCX,0x10
cmp
jae
           .leftdone
mov
           RDI, semptybyte
call
           _print
inc
           RCX
           .leftbytesdone
jmp
```

.leftdone:

Uses _print 38a.

Next we print out the separator between the two compare block.

29b $\langle print\text{-}comparision16 \ 28a \rangle + \equiv$ (31b) ⊲29a 29c⊳

```
mov
           RDI, sseparator
call
           _print
```

Uses _print 38a and sseparator 27d.

To print the second half (for comparision) we restore the values of the parameters first.

 $\langle print\text{-}comparision16 \ 28a \rangle + \equiv$ 29c

(31b) ⊲29b 29d⊳

```
RDI, [RBP-8]
mov
           RSI, [RBP-16]
mov
mov
           RDX, [RBP-24]
```

Now we compare each byte with the original value first and then print it out. If the value differs from the original value we additionally mark the byte.

29d $\langle print\text{-}comparision16 \ 28a \rangle + \equiv$ xor

(31b) ⊲29c 30⊳

.nextbyteright:

```
[RBP-32], RCX
mov
           RCX, RDX
cmp
           .rightbytesdone
jae
           AL, [RSI+RCX]
mov
           AH, [RDI+RCX]
mov
           R12W, AX
mov
           AH, AL
cmp
```

RCX, RCX

```
.printplain
              jе
                          RDI, sbgred
              mov
              call
                          _print
        .printplain:
              xor
                          RDI,RDI
                          AX,R12W
              mov
                          AH, AH
              xor
                          DI,AX
              mov
                          RSI, scratch
              {\tt mov}
              call
                          _printh8bit
              mov
                          AX,R12W
                          AH,AL
              cmp
                          .printdone
              jе
                          RDI, sresetstyle
              mov
              call
                          _print
         .printdone:
              mov
                          RDI,1
              mov
                          RSI,sblank
                          _nprint
              call
                          RDI, [RBP-8]
              mov
              {\tt mov}
                          RSI, [RBP-16]
                          RDX, [RBP-24]
              mov
              mov
                          RCX, [RBP-32]
                          RCX
              inc
                          .nextbyteright
              jmp
        .rightbytesdone:
      Uses _nprint 37b, _print 38a, _printh8bit 41b, sbgred 27d, sblank 27d, scratch 8c,
        and sresetstyle 27d.
      \langle print\text{-}comparision16 28a \rangle + \equiv
30
                                                                                   (31b) ⊲29d
              cmp
                          RCX,0x10
                          .rightdone
              jae
                          RCX
              inc
              jmp
                          .rightbytesdone
         .rightdone:
         .done:
                          RDI, sresetstyle
              mov
              call
                          _print
              mov
                          RDI,1
                          RSI,slf
              mov
                          _nprint
              call
                          R12
              pop
                          RSP,RBP
              mov
              pop
                          RBP
```

```
ret
          Uses _nprint 37b, _print 38a, slf 10d, and sresetstyle 27d.
            TBD
          Parameters
         RDI
                         the address of the first array
         RSI
                         the address of the second array
         RDX
                         number of bytes to print. In each line 16 bytes from the first and 16 bytes
                         from the right side are printed
          \langle print\text{-}comparision 31a \rangle \equiv
                                                                                                                       (31b)
31a
            _printcompare:
                                    RDX,16
                    mov
                    call
                                    _printcompare16
                    ret
          Defines:
            _printcompare, used in chunk 27b.
          Uses _printcompare16 28a.
            TBD
          \langle cacheread.asm 31b \rangle \equiv
31b
            \langle preamble 5 \rangle
            section .rodata
            \langle common-rodata \ 10d \rangle
            \langle cacheread\text{-}rodata 27d \rangle
            section .bss
            \langle data\text{-}udata \text{ 8b} \rangle
            \langle probe-udata \ 13c \rangle
            ⟨readback-udata 24a⟩
            \langle timings-udata \ 16b \rangle
            \langle scratch\text{-}udata \ 8c \rangle
            section .text
            \langle cacheread\text{-}program 24b \rangle
```

⟨clearcache 13a⟩

 $\langle readcachetiming 14d \rangle$

⟨calculate-cache-access-time 8a⟩

```
\langle analyze cachemintiming 16a \rangle
\langle readarea 25a \rangle
\langle print\text{-}comparision 31a \rangle
\langle print\text{-}comparision16 28a \rangle
\langle xorshift\text{-}prng 36a \rangle
\langle utilities 35a \rangle
```

3 Signals

3.1 Basics

TBD

3.2 Detecting Signals

TBD

3.3 Handling Signals

TBD

4 Utilities

4.1 Introduction

```
TBD
```

```
35a \langle utilities 35a \rangle \equiv (11f 18c 21a 23c 31b) \langle nprint 37b \rangle \langle print 38a \rangle \langle printdu64bit 39a \rangle \langle printh8bit 41b \rangle
```

4.2 Common Chunks

4.2.1 Exit Program

This chunk ends the program with exit code 0.

```
35b \langle exitProgram \ 35b \rangle \equiv (11e 18b 23b 27c) xor RDI,RDI mov RAX,60 syscall
```

4.2.2 Stack Frame

A chunk to create a stack frame.

```
35c \langle enterstackframe \ 35c \rangle \equiv (14d 25b) push RBP mov RBP,RSP
```

A chunk to clean up the created stack frame.

```
35d \langle leavestackframe \ 35d \rangle \equiv (15c 26f)
mov RSP,RBP
pop RBP
```

4.3 Random Number Generator

To initialize the data a random number generator (RNG) is used. The sample programs use xorshift¹ as RNG.

First we clear the direction flag to ensure that we are incrementing the data pointer RDI.

Next we move the number of values to be generated to RCX (which is a counter in x86 processors) and divide it by 4 (because we use a 32bit RNG). Additionally we move the seed to EAX.

Parameters

```
RDI
                     the address of the memory which is to be filled with random numbers
        RSI
                     the number of bytes that are filled with random numbers. This must be a
                     multiple of 4
        EDX
                     the seed of the RNG
        \langle xorshift\text{-}prng \ 36a \rangle \equiv
36a
                                                                              (11f 18c 21a 23c 31b) 36b⊳
          _xorshift:
                 cld
                 mov
                              RCX, RSI
                 shr
                              RCX,2
                              EAX, EDX
                 mov
        Defines:
          _xorshift, used in chunks 9e and 13d.
          Now we can generate the next 32bit random number.
36b
        \langle xorshift\text{-}prng 36a \rangle + \equiv
                                                                        (11f 18c 21a 23c 31b) ⊲36a 37a⊳
           .next_random:
                 mov
                              EBX, EAX
                 shl
                              EAX,13
                              EAX, EBX
                 xor
                              EBX, EAX
                 mov
                              EAX,17
                 shr
                              EAX, EBX
                 xor
                              EBX, EAX
                 mov
                              EAX,5
                 shl
                              EAX, EBX
```

xor

¹https://en.wikipedia.org/wiki/Xorshift

Because we want to generate multiple random numbers we store the value of EAX to [RDI] and loop for the next random number.

```
37a \langle xorshift\text{-}prng \ 36a \rangle + \equiv (11f 18c 21a 23c 31b) \triangleleft 36b stosd loop .next_random ret
```

4.4 Printing Strings

4.4.1 Printing Strings with Length

The routine _nprint prints a string with the given length to stdout.

the address to the bytes to print to stdout

We move the number of bytes to print to RDX which is the 3rd parameter to the systemcall. Next we move the address of the bytes to print to RSI which is the 2nd parameter to the systemcall. The 1st argument (in RDI) to the systemcall is the file descriptor (1 is stdout). Additionally the number of the systemcall (1) is passed in RAX. The systemcall (syscall) now prints RDX bytes from [RSI] to the file descriptor RDI.

At the end we return to the caller.

Parameters

RSI

```
RDI the number of bytes to print to stdout
```

```
37b ⟨nprint 37b⟩≡
_nprint:

mov RDX,RDI

mov RDI,1

mov RAX,1

syscall
```

Defines:

_nprint, used in chunks 11, 17, 18a, 23b, 28-30, 38d, 41a, and 42a.

4.4.2 Printing C-Strings

The routine _print prints a null-terminated string to stdout.

First we clear the direction flag to increment the address in RDI while scanning the data.

Next we start with clearing AL (setting it to null) and saving the address of the string to RSI. We're using RSI because we later need the address to calculate the length of the string.

Parameters

RDI the address to the null-terminated bytes to print to stdout

```
38a ⟨print 38a⟩≡
_print:
cld
xor AL,AL
mov RSI,RDI
```

Defines:

_print, used in chunks 10e, 11d, 17, 18a, 23b, 29, and 30.

Next we search for the terminating null ('\0') character. For this we use the instruction scasb (scan string byte) which compares the byte at the address [RDI] with the value in AL and sets the flags accordingly. When the byte at [RDI] is not the value of AL the next instruction (jne) jumps to the given label (.next_char in this case).

scasb additionally increments RDI so that we go through the string until '\0' is found.

```
38b ⟨print 38a⟩+≡ (35a) ⊲38a 38c⊳
.next_char:
scasb
jne .next_char
```

After we have found the string termination we calculate the number of bytes that the string has. In RSI we now have the starting address of the bytes to print and in RDI we have the end address of the bytes to print. After that we calculate the number of bytes to print.

38c
$$\langle print \ 38a \rangle + \equiv$$
 (35a) \triangleleft 38b 38d ▷ sub RDI,RSI

Now we have the address of the string in RDI and the length of the string in RSI which are the 1st and 2nd argument in the call of _nprint.

```
38d ⟨print 38a⟩+≡ (35a) ⊲38c
call _nprint
ret
Uses _nprint 37b.
```

4.5 Printing Numbers

4.5.1 Printing a Decimal 64bit Unsigned Integer

The routine _printdu64bit prints a given 64bit integer as unsigned decimal number to stdout.

To print a decimal number we have to divide the number by 10 and get the remainder for printing (from right to left). For this we move the divisor to a register and the dividend to RAX. We have to use RAX because this is the only register we can use for division.

Additionally we need the address of the scratch area in RDI for storing the result. We also save the address of the scratch area to R8 for later use.

To increment the address during the processing we clear the direction flag.

Parameters

RDI the number number to print to stdout

RSI the address of a scratch area with a size of at least 20 bytes

```
39a ⟨printdu64bit 39a⟩≡
_printdu64bit:

mov RAX,RDI

mov RDI,RSI

mov R8,RDI

mov RCX,10

cld
```

Defines:

_printdu64bit, used in chunks 10f, 11d, and 17d.

Now we define a label to jump back when we see that there are still more digits to print. Then we test RAX for 0 and end the processing of the digits.

```
39b ⟨printdu64bit 39a⟩+≡ (35a) ⊲39a 39c⊳
.next:
cmp RAX,0
je .done
```

Next we divide RAX by RCX. For this we have to clear RDX because this is the higher value of the dividend. The result is then placed into RAX and the remainder into RDX.

```
39c \langle printdu64bit 39a\rangle + \equiv (35a) \triangleleft 39d \triangleright xor RDX,RDX div RCX
```

We now exchange the result and the remainder because we now need the remainder in RAX (or AL) for further processing. Now we can add the ASCII character '0' to AL and have the correct ASCII value in AL. Now we can store the ASCII character to the scratch area.

```
39d \langle printdu64bit 39a\rangle + \equiv (35a) \triangleleft 39c 40a\triangleright xchg RDX,RAX add AL,'0' stosb
```

Now we restore RAX (which we saved to RDX) to go into the next round.

```
40a \langle printdu64bit 39a \rangle + \equiv (35a) \triangleleft 39d \triangleleft 40b \triangleright mov RAX,RDX jmp .next
```

Now that we have all the numbers as ASCII characters we are nearly done. We now have to reverse the number in memory because the number saved at the lowest address is the digit with the least significance.

We now start with checking if we have written any character. If not then we write the ASCII character '0' into the memory. We use the instruction stosb for this to adjust the address in RDI at the same time.

```
40b ⟨printdu64bit 39a⟩+≡ (35a) ⊲40a 40c⊳

.done:

cmp RDI,RSI

jne .printout

mov AL,'0'

stosb
.printout:
```

Next we calculate the number of digits that the number has. For this we move the address of the last digit to RDX and subtract the start of the scratch area from this. Next we adjust RDI because it points to the first address after the number.

```
40c \langle printdu64bit 39a \rangle + \equiv (35a) \triangleleft 40b 40d \triangleright mov RDX,RDI sub RDX,RSI dec RDI
```

We now have RSI with the address of the start of the number and RDI with the address of the end. We now have to exchange the digits from the front and the end to get the right number. For this we increment RSI and decrement RDI after each exchange and when the addresses pass each other we are done.

```
40d
        \langle printdu64bit 39a \rangle + \equiv
                                                                                           (35a) ⊲40c 41a⊳
           .reverse:
                               AL, [RSI]
                 mov
                               AH, [RDI]
                 mov
                               [RSI],AH
                 mov
                               [RDI],AL
                 mov
                 dec
                               RDI
                               RSI
                 inc
                               RSI, RDI
                 cmp
                               .reverse
                 jb
```

Now we restore the address of the scratch area to RSI and move the number of digits (which we stored in RDX) to RDI and can the call _nprint to print the number.

```
41a ⟨printdu64bit 39a⟩+≡ (35a) ⊲40d

mov RSI,R8

mov RDI,RDX

call _nprint

ret

Uses _nprint 37b.
```

4.5.2 Printing a Hexadecimal 8bit Integer

The routine _printh8bit prints a given 8bit integer as hexadecimal number to stdout. To print a hexadecimal number we mask a nibble (4bit) and have the number to print. First we clear the register RAX and move the number to AX for further processing and clear the higher 8bit (AH). Additionally we move it to R8 for later restore.

Additionally we need the address of the scratch area in RDI for storing the result.

To increment the address during the processing we clear the direction flag.

Parameters

DI the number number to print to stdout. Only the lower 8bit are used.

RSI the address of a scratch area with a size of at least 2 bytes

```
41b ⟨printh8bit 41b⟩≡
_printh8bit:

xor RAX,RAX
mov AX,DI
xor AH,AH
mov R8,RAX
mov RDI,RSI
cld (35a) 41c⊳
```

Defines:

_printh8bit, used in chunks 17c, 18a, 23b, 28e, and 29d.

Now we mask the higher 4 bit of AL by shifting it 4 bits to the right and mask out all but the lower 4 bit. Next we call the internal method printh8bit.printh4bit to print out this nibble.

```
41c \langle printh8bit\ 41b \rangle + \equiv (35a) \triangleleft 41b \triangleleft 42a \triangleright shr AL,4 and AL,0x0f call .printh4bit
```

Next we restore the number and print out the lower 4 bits.

```
42a
        \langle printh8bit 41b \rangle + \equiv
                                                                                                  (35a) ⊲41c
                 mov
                               RAX,R8
                 and
                               AL, 0x0f
                 call
                               .printh4bit
                               RDI,2
                 mov
                 call
                               _nprint
                 ret
           ⟨printh8bit.printh4bit 42b⟩
        Uses _nprint 37b.
```

Now we define the internal method to print a hexadecimal digit.

First we test if the digit is above or equal to 10. In this case we have to print out a character between 'a' and 'f' else we print out a decimal digit (between '0' and '9').

Parameters (internal)

AL the lower 4 bit contain the hexadecimal digit print to stdout

RDI the address of a scratch area

```
42b \langle printh8bit.printh4bit 42b\rangle \equiv .printh4bit: cmp AL,10
```

cmp AL,10 jae .printa2f

Defines:

printh8bit.printh4bit, never used.

Now we add '0' to get the code for the digit between '0' and '9'.

```
42c \langle printh8bit.printh4bit.42b\rangle + \equiv (42a) \triangleleft 42b 42d\triangleright add AL,'0' jmp .printout
```

Else we print a digit between 'a' and 'f'. We first subtract 10 because the value in AL is now between 10 and 15.

```
42d \langle printh8bit.printh4bit 42b \rangle + \equiv (42a) \triangleleft 42c 42e \triangleright .printa2f: sub AL,10 add AL,'a'
```

Now we store the character into the storage area.

```
42e \langle printh8bit.printh4bit.42b \rangle + \equiv (42a) \triangleleft 42d .printout: stosb ret
```

A Glossary

 ${\bf x86}\,$ x86 denotes a microprocessor architecture based on the 8086/8088 36

B Acronyms

 ${\sf ASCII}$ American Standard Code for Information Interchange 39, 40

LF line feed 10, 11, 17

 ${\sf RNG}$ random number generator 36

C x86-Instructions

```
clflush Flush Cache Line, introduced with Intel<sup>®</sup> Pentium<sup>®</sup> 4 10

lfence Load Fence, introduced with Intel<sup>®</sup> Pentium<sup>®</sup> 4 7, 10

rdtsc Read Time Stamp Counter, introduced with Intel<sup>®</sup> Pentium<sup>®</sup> 7, 9
```

D Code Chunks

```
\langle analyze cachemintiming 16a \rangle
\langle analyze caches impthrestiming 20a \rangle
\langle cacheread\text{-}program 24b \rangle
\langle cacheread\text{-}rodata 27d \rangle
\langle cacheread\text{-}udata 28b \rangle
\langle cacheread.asm 31b \rangle
\langle cachereadbyte-program 13b \rangle
\langle cachereadbyte-rodata 17a \rangle
\langle cachereadbyte.asm \ 18c \rangle
\langle cacheread by te 2. asm 21a \rangle
\langle cachereadbyte3-program 21b \rangle
\langle cacheread by te 3. asm 23c \rangle
\langle cachetiming-program 9a \rangle
\langle cachetiming-rodata \ 10c \rangle
\langle cachetiming.asm 11f \rangle
\langle calculate\text{-}cache\text{-}access\text{-}time \ 8a \rangle
⟨clearcache 13a⟩
\langle common-rodata \ 10d \rangle
\langle data\text{-}udata \text{ 8b} \rangle
\langle enterstack frame 35c \rangle
\langle exitProgram 35b \rangle
\langle init\text{-}random\text{-}data 9b \rangle
\langle init\text{-}random\text{-}probe \ 13d \rangle
\langle leavestackframe 35d \rangle
\langle license 74 \rangle
\langle nprint 37b \rangle
\langle preamble 5 \rangle
\langle print 38a \rangle
\langle print\text{-}comparision 31a \rangle
\langle print\text{-}comparision16 28a \rangle
\langle printdu64bit 39a \rangle
⟨printh8bit 41b⟩
⟨printh8bit.printh4bit 42b⟩
\langle probe-udata \ 13c \rangle
⟨readarea 25a⟩
⟨readback-udata 24a⟩
⟨readbyte2cache 22a⟩
```

D Code Chunks

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