

minion executable documentation

Ian Gent, Chris Jefferson, Ian Miguel, Neil Moore, Karen Petrie, Andrea Rendl

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You are viewing documentation for minion. The same documentation is available from a minion executable by typing `minion help` at the command line. We intend that the command line help system be the main source of documentation for the system.

Each of the entries below concerns a different aspect of the system, and the entries are arranged hierarchically. For example to view information about the set of available constraints as a whole view “constraints” and to view specific information about the alldiff constraint view “constraints alldiff”.

A good place to start would be viewing the “input example” entry which exhibits a complete example of a minion input file.

Usage: `minion [switches] [minion input file]`

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1 constraints

Description

Minion supports many constraints and these are regularly being improved and added to. In some cases multiple implementations of the same constraints are provided and we would appreciate additional feedback on their relative merits in your problem.

Minion does not support nesting of constraints, however this can be achieved by auxiliary variables and reification.

Variables can be replaced by constants. You can find out more on expressions for variables, vectors, etc. in the section on variables.

References

`help variables`

2 constraints abs

Description

The constraint

`abs(x,y)`

makes sure that $x=|y|$, i.e. x is the absolute value of y .

Reference

`help constraints abs`

3 constraints alldiff

Description

Forces the input vector of variables to take distinct values.

Example

Suppose the input file had the following vector of variables defined:

`DISCRETE myVec[9] {1..9}`

To ensure that each variable takes a different value include the following constraint:

`alldiff(myVec)`

Notes

Enforces the same level of consistency as a clique of not equals constraints.

References

See

```
help constraints gacalldiff
```

for the same constraint that enforces GAC.

4 constraints difference

Description

The constraint

```
difference(x,y,z)
```

ensures that $z=|x-y|$ in any solution.

Notes

This constraint can be expressed in a much longer form, this form both avoids requiring an extra variable, and also gets better propagation. It gets bounds consistency.

5 constraints diseq

Description

Constrain two variables to take different values.

Notes

Achieves arc consistency.

Example

```
diseq(v0,v1)
```

6 constraints div

Description

The constraint

```
div(x,y,z)
```

ensures that $\text{floor}(x/y)=z$.

Notes

This constraint is only available for positive domains x , y and z .

References

help constraints modulo

7 constraints element

Description

The constraint

```
element(vec, i, e)
```

specifies that, in any solution, $\text{vec}[i] = e$ and i is in the range $[0 \dots |\text{vec}|-1]$.

Notes

Warning: This constraint is not confluent. Depending on the order the propagators are called in Minion, the number of search nodes may vary when using element. To avoid this problem, use `watchelement` instead. More details below.

The level of propagation enforced by this constraint is not named, however it works as follows. For constraint $\text{vec}[i]=e$:

- After i is assigned, ensures that $\min(\text{vec}[i]) = \min(e)$ and $\max(\text{vec}[i]) = \max(e)$.
- When e is assigned, removes idx from the domain of i whenever e is not an element of the domain of $\text{vec}[\text{idx}]$.
- When $\text{m}[\text{idx}]$ is assigned, removes idx from i when $\text{m}[\text{idx}]$ is not in the domain of e .

This level of consistency is designed to avoid the propagator having to scan through vec , except when e is assigned. It does a quantity of cheap propagation and may work well in practise on certain problems.

Element is not confluent, which may cause the number of search nodes to vary depending on the order in which constraints are listed in the input file, or the order they are called in Minion. For example, the following input causes Minion to search 41 nodes.

```
MINION 3
**VARIABLES**
DISCRETE x[5] {1..5}
**CONSTRAINTS**
element([x[0],x[1],x[2]], x[3], x[4])
alldiff([x])
**EOF**
```

However if the two constraints are swapped over, Minion explores 29 nodes. As a rule of thumb, to get a lower node count, move element constraints to the end of the list.

References

See the entry

```
constraints watchelement
```

for details of an identical constraint that enforces generalised arc consistency.

8 constraints element_one

Description

The constraint `element one` is identical to `element`, except that the vector is indexed from 1 rather than from 0.

References

See

```
help constraints element
```

for details of the `element` constraint which is almost identical to this one.

9 constraints eq

Description

Constrain two variables to take equal values.

Example

```
eq(x0,x1)
```

Notes

Achieves bounds consistency.

Reference

```
help constraints minuseq
```

10 constraints gacalldiff

Description

Forces the input vector of variables to take distinct values.

Example

Suppose the input file had the following vector of variables defined:

```
DISCRETE myVec[9] {1..9}
```

To ensure that each variable takes a different value include the following constraint:

```
gacalldiff(myVec)
```

Notes

This constraint enforces generalized arc consistency.

11 constraints gcc

Description

The Generalized Cardinality Constraint (GCC) constrains the number of each value that a set of variables can take.

```
gccweak([primary variables], [values of interest], [capacity variables])
```

For each value of interest, there must be a capacity variable, which specifies the number of occurrences of the value in the primary variables.

This constraint is new, and its syntax and implementation are not finalised.

Description

The Generalized Cardinality Constraint (GCC) constrains the number of each value that a set of variables can take.

```
gcc([primary variables], [values of interest], [capacity variables])
```

For each value of interest, there must be a capacity variable, which specifies the number of occurrences of the value in the primary variables.

This constraint is new, and its syntax and implementation are not finalised.

Example

Suppose the input file had the following vectors of variables defined:

```
DISCRETE myVec[9] {1..9}
BOUND cap[9] {0..2}
```

The following constraint would restrict the occurrence of values 1..9 in myVec to be at most 2 each initially, and finally equal to the values of the cap vector.

```
gcc(myVec, [1,2,3,4,5,6,7,8,9], cap)
```

Notes

This constraint enforces a hybrid consistency. It reads the bounds of the capacity variables, then enforces GAC over the primary variables only. Then the bounds of the capacity variables are updated using flow algorithms similar to those proposed by Quimper et al, Improved Algorithms for the Global Cardinality Constraint.

This constraint provides stronger propagation to the capacity variables than the gccweak constraint.

12 constraints gccweak

Example

Suppose the input file had the following vectors of variables defined:

```
DISCRETE myVec[9] {1..9}
BOUND cap[9] {0..2}
```

The following constraint would restrict the occurrence of values 1..9 in myVec to be at most 2 each initially, and finally equal to the values of the cap vector.

```
gccweak(myVec, [1,2,3,4,5,6,7,8,9], cap)
```

Notes

This constraint enforces a hybrid consistency. It reads the bounds of the capacity variables, then enforces GAC over the primary variables only. Then the bounds of the capacity variables are updated by counting values in the domains of the primary variables.

The consistency over the capacity variables is weaker than the gcc constraint, hence the name gccweak.

13 constraints hamming

Description

The constraint

```
hamming(X,Y,c)
```

ensures that the hamming distance between X and Y is at least c. That is, that the size of the set $\{i \mid X[i] \neq Y[i]\}$ is greater than or equal to c.

14 constraints ineq

Description

The constraint

```
ineq(x, y, k)
```

ensures that

```
x <= y + k
```

in any solution.

Notes

Minion has no strict inequality ($<$) constraints. However $x < y$ can be achieved by

```
ineq(x, y, -1)
```

15 constraints lexleq

Description

The constraint

```
lexleq(vec0, vec1)
```

takes two vectors `vec0` and `vec1` of the same length and ensures that `vec0` is lexicographically less than or equal to `vec1` in any solution.

Notes

This constraints achieves GAC.

References

See also

```
help constraints lexless
```

for a similar constraint with strict lexicographic inequality.

16 constraints lexless

Description

The constraint

```
lexless(vec0, vec1)
```

takes two vectors `vec0` and `vec1` of the same length and ensures that `vec0` is lexicographically less than `vec1` in any solution.

Notes

This constraint maintains GAC.

References

See also

```
help constraints lexleq
```

for a similar constraint with non-strict lexicographic inequality.

17 constraints litsumgeq

Description

The constraint `litsumgeq(vec1, vec2, c)` ensures that there exists at least `c` distinct indices `i` such that `vec1[i] = vec2[i]`.

Notes

A SAT clause `{x,y,z}` can be created using:

```
litsumgeq([x,y,z],[1,1,1],1)
```

Note also that this constraint is more efficient for smaller values of `c`. For large values consider using `watchsumleq`.

Reifiability

This constraint is not reifiable.

References

See also

```
help constraints watchsumleq  
help constraints watchsumgeq
```

18 constraints max

Description

The constraint

```
max(vec, x)
```

ensures that `x` is equal to the maximum value of any variable in `vec`.

References

See

```
help constraints min
```

for the opposite constraint.

19 constraints min

Description

The constraint

```
min(vec, x)
```

ensures that x is equal to the minimum value of any variable in vec .

References

See

```
help constraints max
```

for the opposite constraint.

20 constraints minuseq

Description

Constraint

```
minuseq(x,y)
```

ensures that $x=-y$.

Reference

```
help constraints eq
```

21 constraints modulo

Description

The constraint

```
modulo(x,y,z)
```

ensures that $x\%y=z$ i.e. z is the remainder of dividing x by y .

Notes

This constraint is only available for positive domains x , y and z .

References

```
help constraints div
```

22 constraints occurrence

Description

The constraint

```
occurrence(vec, elem, count)
```

ensures that there are count occurrences of the value elem in the vector vec.

Notes

elem must be a constant, not a variable.

References

```
help constraints occurrenceleq
help constraints occurrencegeq
```

23 constraints occurrencegeq

Description

The constraint

```
occurrencegeq(vec, elem, count)
```

ensures that there are AT LEAST count occurrences of the value elem in the vector vec.

Notes

elem and count must be constants

References

```
help constraints occurrence
help constraints occurrenceleq
```

24 constraints occurrenceleq

Description

The constraint

```
occurrenceleq(vec, elem, count)
```

ensures that there are AT MOST count occurrences of the value elem in the vector vec.

Notes

elem and count must be constants

References

`help constraints occurrence`
`help constraints occurrencegeq`

25 constraints pow

Description

The constraint

`pow(x,y,z)`

ensures that $x^y=z$.

Notes

This constraint is only available for positive domains x , y and z .

26 constraints product

Description

The constraint

`product(x,y,z)`

ensures that $z=xy$ in any solution.

Notes

This constraint can be used for (and, in fact, has a specialised implementation for) achieving boolean AND, i.e. $x \& y=z$ can be modelled as

`product(x,y,z)`

The general constraint achieves bounds generalised arc consistency for positive numbers.

27 constraints reification

Description

Reification is provided in two forms: `reify` and `reifyimply`.

`reify(constraint, r)` where r is a 0/1 var

ensures that r is set to 1 if and only if `constraint` is satisfied. That is, if r is 0 the constraint must NOT be satisfied; and if r is 1 it must be satisfied as normal. Conversely, if the constraint is satisfied then r must be 1, and if not then r must be 0.

`reifyimply(constraint, r)`

only checks that if `r` is set to 1 then constraint must be satisfied. If `r` is not 1, constraint may be either satisfied or unsatisfied. Furthermore `r` is never set by propagation, only by search; that is, satisfaction of constraint does not affect the value of `r`.

Notes

ALMOST ALL constraints are reifiable. Individual constraint entries mention if the constraint is NOT reifiable.

ALL constraints are reifyimplyable.

28 constraints reify

References

See
 `help constraints reification`

29 constraints reifyimply

References

See
 `help constraints reification`

30 constraints sumgeq

Description

The constraint

`sumgeq(vec, c)`

ensures that `sum(vec) >= c`.

31 constraints sumleq

Description

The constraint

`sumleq(vec, c)`

ensures that `sum(vec) <= c`.

32 constraints table

Description

An extensional constraint that enforces GAC. The constraint is specified via a list of tuples.

Example

To specify a constraint over 3 variables that allows assignments (0,0,0), (1,0,0), (0,1,0) or (0,0,1) do the following.

1) Add a tuplelist to the ****TUPLELIST**** section, e.g.:

```
**TUPLELIST**
myext 4 3
0 0 0
1 0 0
0 1 0
0 0 1
```

N.B. the number 4 is the number of tuples in the constraint, the number 3 is the -arity.

2) Add a table constraint to the ****CONSTRAINTS**** section, e.g.:

```
**CONSTRAINTS**
table(myvec, myext)
```

and now the variables of myvec will satisfy the constraint myext.

Example

The constraints extension can also be specified in the constraint definition, e.g.:

```
table(myvec, {<0,0,0>,<1,0,0>,<0,1,0>,<0,0,1>})
```

References

help input tuplelist

33 constraints watched-and

Description

The constraint

```
watched-and({C1,...,Cn})
```

ensures that the constraints C1,...,Cn are all true.

Notes

pointless, bearing in mind that a CSP is simply a conjunction of constraints already! However sometimes it may be necessary to use a conjunction as a child of another constraint, for example in a reification:

```
reify(watched-and({...}),r)
```

References

See also

help constraints watched-or

34 constraints watched-or

Description

The constraint

```
watched-or({C1,...,Cn})
```

ensures that at least one of the constraints C_1, \dots, C_n is true.

References

See also

help constraints watched-and

35 constraints watchelement

Description

The constraint

```
watchelement(vec, i, e)
```

specifies that, in any solution, $\text{vec}[i] = e$ and i is in the range $[0 \dots |\text{vec}|-1]$.

Notes

Enforces generalised arc consistency.

References

See entry

help constraints element

for details of an identical constraint that enforces a lower level of consistency.

36 constraints watchelement_one

Description

This constraint is identical to `watchelement`, except the vector is indexed from 1 rather than from 0.

References

See entry

`help constraints watchelement`

for details of `watchelement` which `watchelement_one` is based on.

37 constraints watchless

Description

The constraint `watchless(x,y)` ensures that x is less than y .

References

See also

`help constraints ineq`

38 constraints watchsumgeq

Description

The constraint `watchsumgeq(vec, c)` ensures that $\text{sum}(\text{vec}) \geq c$.

Notes

For this constraint, small values of c are more efficient.

Equivalent to `litsumgeq(vec, [1,...,1], c)`, but faster.

This constraint works on 0/1 variables only.

Reifiability

This constraint is not reifiable.

References

See also

`help constraints watchsumleq`

`help constraints litsumgeq`

39 constraints watchsumleq

Description

The constraint `watchsumleq(vec, c)` ensures that $\text{sum}(\text{vec}) \leq c$.

Notes

Equivalent to `litsumgeq([vec1,...,vecn], [0,...,0], n-c)` but faster.

This constraint works on binary variables only.

For this constraint, large values of `c` are more efficient.

References

See also

`help constraints watchsumgeq`
`help constraints litsumgeq`

40 constraints watchvecneq

Description

The constraint

`watchvecneq(A, B)`

ensures that `A` and `B` are not the same vector, i.e., there exists some index `i` such that `A[i] != B[i]`.

41 constraints weightedsumgeq

Description

The constraint

`weightedsumgeq(constantVec, varVec, total)`

ensures that $\text{constantVec} \cdot \text{varVec} \geq \text{total}$, where $\text{constantVec} \cdot \text{varVec}$ is the scalar dot product of `constantVec` and `varVec`.

References

`help constraints weightedsumleq`
`help constraints sumleq`
`help constraints sumgeq`

42 constraints weightedsumleq

Description

The constraint

```
weightedsumleq(constantVec, varVec, total)
```

ensures that $\text{constantVec} \cdot \text{varVec} \leq \text{total}$, where $\text{constantVec} \cdot \text{varVec}$ is the scalar dot product of constantVec and varVec .

References

```
help constraints weightedsumgeq
help constraints sumleq
help constraints sumgeq
```

43 constraints w-inrange

Description

The constraint `w-inrange(x, [a,b])` ensures that $a \leq x \leq b$.

References

See also

```
help constraints w-notinrange
```

44 constraints w-inset

Description

The constraint `w-inset(x, [a1,...,an])` ensures that x belongs to the set $\{a1, \dots, an\}$.

References

See also

```
help constraints w-notinset
```

45 constraints w-literal

Description

The constraint `w-literal(x, a)` ensures that $x=a$.

References

See also

```
help constraints w-notliteral
```

46 constraints w-notinrange

Description

The constraint `w-notinrange(x, [a,b])` ensures that $x < a$ or $b < x$.

References

See also

`help constraints w-inrange`

47 constraints w-notinset

Description

The constraint `w-notinset(x, [a1,...,an])` ensures that x does not belong to the set $\{a1, \dots, an\}$.

References

See also

`help constraints w-inset`

48 constraints w-notliteral

Description

The constraint `w-notliteral(x, a)` ensures that $x \neq a$.

References

See also

`help constraints w-literal`

49 input

Description

Minion expects to be provided with the name of an input file as an argument. This file contains a specification of the CSP to be solved as well as settings that the search process should use. The format is

```
Minion3Input := MINION 3
               <InputSection>+
               **EOF**

InputSection := <VariablesSection>
               | <SearchSection>
               | <ConstraintsSection>
```

| <TuplelistSection>

i.e. 'MINION 3' followed by any number of variable, search, constraints and tuplelists sections (can repeat) followed by '**EOF**', the end of file marker.

All text from a '#' character to the end of the line is ignored.

See the associated help entries below for information on each section.

Notes

You can give an input file via standard input by specifying '--' as the file name, this might help when minion is being used as a tool in a shell script or for compressed input, e.g.,

```
gunzip -c myinput.minion.gz | minion
```

50 input constraints

Description

The constraints section consists of any number of constraint declarations on separate lines.

```
ConstraintsSection ::= **CONSTRAINTS**  
                    <ConstraintDeclaration>*
```

Example

```
**CONSTRAINTS**  
eq(bool,0)  
alldiff(d)
```

References

See help entries for individual constraints under

```
help constraints
```

for details on constraint declarations.

51 input example

Example

Below is a complete minion input file with commentary, as an example.

```
MINION 3  
  
# While the variable section doesn't have to come first, you can't  
# really do anything until  
# You have one...
```

```

**VARIABLES**

# There are 4 type of variables
BOOL bool      # Boolean don't need a domain
BOUND b {1..3} # Bound vars need a domain given as a range
DISCRETE d {1..3} # So do discrete vars

#Note: Names are case sensitive!

# Internally, Bound variables are stored only as a lower and upper bound
# Whereas discrete variables allow any sub-domain

SPARSEBOUND s {1,3,6,7} # Sparse bound variables take a sorted list of values

# We can also declare matrices of variables!

DISCRETE q[3] {0..5} # This is a matrix with 3 variables: q[0],q[1] and q[2]
BOOL bm[2,2] # A 2d matrix, variables bm[0,0], bm[0,1], bm[1,0], bm[1,1]
BOOL bn[2,2,2,2] # You can have as many indices as you like!

#The search section is entirely optional
**SEARCH**

# Note that everything in SEARCH is optional, and can only be given at
# most once!

# If you don't give an explicit variable ordering, one is generated.
# These can take matrices in interesting ways like constraints, see below.
VARORDER [bool,b,d]

# If you don't give a value ordering, 'ascending' is used
#VALORDER [a,a,a,a]

# You can have one objective function, or none at all.
MAXIMISING bool
# MINIMISING x3

# both (MAX/MIN)IMISING and (MAX/MIN)IMIZING are accepted...

# Print statement takes a vector of things to print

PRINT [bool, q]

# You can also give:
# PRINT ALL (the default)
# PRINT NONE

# Declare constraints in this section!
**CONSTRAINTS**

# Constraints are defined in exactly the same way as in MINION input
formats 1 & 2
eq(bool, 0)
eq(b,d)

```



```

# To get a single variable from a matrix, just index it
eq(q[1],0)
eq(bn[0,1,1,1], bm[1,1])

# It's easy to get a row or column from a matrix. Just use _ in the
# indices you want
# to vary. Just giving a matrix gives all the variables in that matrix.

#The following shows how flattening occurs...

# [bm] == [ bm[_,_] ] == [ bm[0,0], bm[0,1], bm[1,0], bm[1,1] ]
# [ bm[_,1] ] = [ bm[0,1], bm[1,1] ]
# [ bn[1,_,0,_] ] = [ bn[1,0,0,0], b[1,0,0,1], b[1,1,0,0], b[1,1,0,1] ]

# You can string together a list of such expressions!

lexleq( [bn[1,_,0,_], bool, q[0]] , [b, bm, d] )

# One minor problem.. you must always put [ ] around any matrix expression, so
# lexleq(bm, bm) is invalid

lexleq( [bm], [bm] ) # This is OK!

# Can give tuplelists, which can have names!
# The input is: <name> <num_of_tuples> <tuple_length> <numbers...>
# The formatting can be about anything..

**TUPLELIST**

Fred 3 3
0 2 3
2 0 3
3 1 3

Bob 2 2 1 2 3 4

#No need to put everything in one section! All sections can be reopened..
**VARIABLES**

# You can even have empty sections.. if you want

**CONSTRAINTS**

#Specify tables by their names..

table([q], Fred)

# Can still list tuples explicitly in the constraint if you want at
# the moment.
# On the other hand, I might remove this altogether, as it's worse than giving
# Tuplelists

table([q],{ <0,2,3>,<2,0,3>,<3,1,3> })

#Must end with the **EOF** marker!

**EOF**

```

Any text down here is ignored, so you can write whatever you like (or nothing at all...)

52 input search

Description

Inside the search section one can specify

- variable orderings,
- value orderings,
- optimisation function, and
- details of how to print out solutions.

```
SearchSection ::= <VarValOrdering>*  
                <OptimisationFn>?  
                <PrintFormat>?
```

If no varval ordering is given then the variables are assigned in instantiation order and the values tried in ascending order.

If a variable order is given as a command line argument it will override anything specified in the input file.

Multiple variable orders can be given, each with an optional value ordering:

```
VarValOrdering ::= <VarOrder>  
                  <ValOrder>?
```

In each VarOrder an instantiation order is specified for a subset of variables. Variables can optionally be \"auxiliary variables\" (add \"AUX\" to the varorder) meaning that if there are several solutions to the problem differing only in the auxiliary variables, only one is reported by minion.

```
VarOrder ::= VARORDER AUX? <ORDER>? [ <varname>+ ]
```

where

```
<ORDER> ::= STATIC | SDF | SRF | LDF | ORIGINAL | WDEG | CONFLICT | DOMOVERWDEG
```

The value ordering allows the user to specify an instantiation order for the variables involved in the variable order, either ascending (a) or descending (d) for each. When no value ordering is specified, the default is to use ascending order for every search variable.

```
ValOrder ::= VALORDER [ (a|d)+ ]
```

To model an optimisation problem the user can specify to minimise or maximise a variable's value.

```
OptimisationFn ::= MAXIMISING <varname>  
                  | MINIMISING <varname>
```

Finally, the user can control some aspects of the way solutions are

printed. By default (no PrintFormat specified) all the variables are printed in declaration order. Alternatively a custom vector, or ALL variables, or no (NONE) variables can be printed. If a matrix or, more generally, a tensor is given instead of a vector, it is automatically flattened into a vector as described in 'help variables vectors'.

```
PrintFormat ::= PRINT <vector>
              | PRINT ALL
              | PRINT NONE
```

References

See also

```
switches -varorder
```

53 input tuplelist

Description

In a tuplelist section lists of allowed tuples for table constraints can be specified. This technique is preferable to specifying the tuples in the constraint declaration, since the tuplelists can be shared between constraints and named for readability.

The required format is

```
TuplelistSection ::= **TUPLELIST**
                   <Tuplelist>*
```

```
Tuplelist ::= <name> <num_tuples> <tuple_length> <numbers>+
```

Example

```
**TUPLELIST**
AtMostOne 4 3
0 0 0
0 0 1
0 1 0
1 0 0
```

References

help constraints table

54 input variables

Description

The variables section consists of any number of variable declarations on separate lines.

```
VariablesSection ::= **VARIABLES**
                   <VarDeclaration>*
```

Example

```
**VARIABLES**

BOOL bool                                #boolean var
BOUND b {1..3}                           #bounds var
SPARSEBOUND myvar {1,3,4,6,7,9,11}        #sparse bounds var
DISCRETE d[3] {1..3}                     #array of discrete vars
```

References

See the help section

```
help variables
```

for detailed information on variable declarations.

55 switches

Description

Minion supports a number of switches to augment default behaviour. To see more information on any switch, use the help system. The list below contains all available switches. For example to see help on -quiet type something similar to

```
minion help switches -quiet
```

replacing 'minion' by the name of the executable you're using.

56 switches -check

Description

Check solutions for correctness before printing them out.

Notes

This option is the default for DEBUG executables.

57 switches -dumptree

Description

Print out the branching decisions and variable states at each node.

58 switches -findallsols

Description

Find all solutions and count them. This option is ignored if the problem contains any minimising or maximising objective.

59 switches -fullprop

Description

Disable incremental propagation.

Notes

This should always slow down search while producing exactly the same search tree.

Only available in a DEBUG executable.

60 switches -nocheck

Description

Do not check solutions for correctness before printing them out.

Notes

This option is the default on non-DEBUG executables.

61 switches -nodelimit

Description

To stop search after N nodes, do

```
minion -nodelimit N myinput.minion
```

References

```
help switches -timelimit
help switches -sollimit
```

62 switches -noprintsols

Description

Do not print solutions.

63 switches -preprocess

This switch allows the user to choose what level of preprocess is applied to their model before search commences.

The choices are:

- GAC
- generalised arc consistency (default)
- all propagators are run to a fixed point
- if some propagators enforce less than GAC then the model will not necessarily be fully GAC at the outset

- SACBounds
- singleton arc consistency on the bounds of each variable
- AC can be achieved when any variable lower or upper bound is a singleton in its own domain

- SAC
- singleton arc consistency
- AC can be achieved in the model if any value is a singleton in its own domain

- SSACBounds
- singleton singleton bounds arc consistency
- SAC can be achieved in the model when domains are replaced by either the singleton containing their upper bound, or the singleton containing their lower bound

- SSAC
- singleton singleton arc consistency
- SAC can be achieved when any value is a singleton in its own domain

These are listed in order of roughly how long they take to achieve. Preprocessing is a one off cost at the start of search. The success of higher levels of preprocessing is problem specific; SAC preprocesses may take a long time to complete, but may reduce search time enough to justify the cost.

Example

To enforce SAC before search:

```
minion -preprocess SAC myinputfile.minion
```

References

help switches -X-prop-node

64 switches -printsols

Description

Print solutions.

65 switches -printsolonly

Description

Print only solutions and a summary at the end.

66 switches -quiet

Description

Do not print parser progress.

References

help switches -verbose

67 switches -randomiseorder

Description

Randomises the ordering of the decision variables. If the input file specifies an ordering it will randomly permute this. If no ordering is specified a random permutation of all the variables is used.

68 switches -randomseed

Description

Set the pseudorandom seed to N. This allows 'random' behaviour to be repeated in different runs of minion.

69 switches -redump

Description

Print the minion input instance file to standard out. No search is carried out when this switch is used.

70 switches -sollimit

Description

To stop search after N solutions have been found, do

```
minion -sollimit N myinput.minion
```

References

help switches -odelimit

help switches -timelimit

71 switches -solsout

Description

Append all solutions to a named file.
Each solution is placed on a line, with no extra formatting.

Example

To add the solutions of myproblem.minion to mysols.txt do

```
minion -solsout mysols.txt myproblem.minion
```

72 switches -tableout

Description

Append a line of data about the current run of minion to a named file.
This data includes minion version information, arguments to the executable, build and solve time statistics, etc. See the file itself for a precise schema of the supplied information.

Example

To add statistics about solving myproblem.minion to mystats.txt do

```
minion -tableout mystats.txt myproblem.minion
```

73 switches -timelimit

Description

To stop search after N seconds, do

```
minion -timelimit N myinput.minion
```

References

```
help switches -nodelimit  
help switches -sollimit
```

74 switches -varorder

Description

Enable a particular variable ordering for the search process. This flag is experimental and minion's default ordering might be faster.

The available orders are:

```
- sdf - smallest domain first, break ties lexicographically
```


- sdf-random - sdf, but break ties randomly
- srf - smallest ratio first, chooses unassigned variable with smallest percentage of its initial values remaining, break ties lexicographically
- srf-random - srf, but break ties randomly
- ldf - largest domain first, break ties lexicographically
- ldf-random - ldf, but break ties randomly
- random - random variable ordering
- static - lexicographical ordering

75 switches -verbose

Description

Print parser progress.

References

help switches -quiet

76 switches -X-prop-node

Description

Allows the user to choose the level of consistency to be enforced during search.

See entry 'help switches -preprocess' for details of the available levels of consistency.

Example

To enforce SSAC during search:

```
minion -X-prop-node SSAC input.minion
```

References

help switches -preprocess

77 variables

General

Minion supports 4 different variable types, namely

- 0/1 variables,
- bounds variables,
- sparse bounds variables, and
- discrete variables.

Sub-dividing the variable types in this manner affords the greatest opportunity for optimisation. In general, we recommend thinking of the variable types as a hierarchy, where 1 (0/1 variables) is the most efficient type, and 4 (Discrete variables) is the least. The user should use the variable which is the highest in the hierarchy, yet encompasses enough information to provide a full model for the problem they are attempting to solve.

Minion also supports use of constants in place of variables, and constant vectors in place of vectors of variables. Using constants will be at least as efficient as using variables when the variable has a singleton domain.

See the entry on vectors for information on how vectors, matrices and, more generally, tensors are handled in minion input. See also the alias entry for information on how to multiply name variables for convenience.

78 variables 01

Description

01 variables are used very commonly for logical expressions, and for encoding the characteristic functions of sets and relations. Note that wherever a 01 variable can appear, the negation of that variable can also appear. A boolean variable x 's negation is identified by $!x$.

Example

Declaration of a 01 variable called `bool` in input file:

```
BOOL bool
```

Use of this variable in a constraint:

```
eq(bool, 0) #variable bool equals 0
```

79 variables alias

Description

Specifying an alias is a way to give a variable another name. Aliases appear in the ****VARIABLES**** section of an input file. It is best described using some examples:

```
ALIAS c = a
```

```
ALIAS c[2,2] = [[myvar,b[2]], [b[1],anothervar]]
```

80 variables bounds

Description

Bounds variables, where only the upper and lower bounds of the domain are maintained. These domains must be continuous ranges of integers i.e. holes cannot be put in the domains of the variables.

Example

Declaration of a bound variable called myvar with domain between 1 and 7 in input file:

```
BOUND myvar {1..7}
```

Use of this variable in a constraint:

```
eq(myvar, 4) #variable myvar equals 4
```

81 variables constants

Description

Minion supports the use of constants anywhere where a variable can be used. For example, in a constraint as a replacement for a single variable, or a vector of constants as a replacement for a vector of variables.

Examples

Use of a constant:

```
eq(x,1)
```

Use of a constant vector:

```
element([10,9,8,7,6,5,4,3,2,1],idx,e)
```

82 variables discrete

Description

In discrete variables, the domain ranges between the specified lower and upper bounds, but during search any domain value may be pruned, i.e., propagation and search may punch arbitrary holes in the domain.

Example

Declaration of a discrete variable x with domain {1,2,3,4} in input file:

```
DISCRETE x {1..4}
```

Use of this variable in a constraint:

```
eq(x, 2) #variable x equals 2
```

83 variables sparsebounds

Description

In sparse bounds variables the domain is composed of discrete values (e.g. {1, 5, 36, 92}), but only the upper and lower bounds of the domain may be updated during search. Although the domain of these variables is not a continuous range, any holes in the domains must be there at time of specification, as they can not be added during the solving process.

Notes

Declaration of a sparse bounds variable called myvar containing values {1,3,4,6,7,9,11} in input file:

```
SPARSEBOUND myvar {1,3,4,6,7,9,11}
```

Use of this variable in a constraint:
eq(myvar, 3) #myvar equals 3

84 variables vectors

Description

Vectors, matrices and tensors can be declared in minion input. Matrices and tensors are for convenience, as constraints do not take these as input; they must first undergo a flattening process to convert them to a vector before use. Additional commas at the end of vectors are ignored (see example below).

Examples

A vector of 0/1 variables:

```
BOOL myvec[5]
```

A matrix of discrete variables:

```
DISCRETE sudoku[9,9] {1..9}
```

A 3D tensor of 0/1s:

```
BOOL mycube[3,3,2]
```

One can create a vector from scalars and elements of vectors, etc.:

```
alldiff([x,y,myvec[1],mymatrix[3,4]])
```

When a matrix or tensor is constrained, it is treated as a vector whose entries have been strung out into a vector in index order with the rightmost index changing most quickly, e.g.

```
alldiff(sudoku)
```

is equivalent to

```
alldiff([sudoku[0,0],...,sudoku[0,8],...,sudoku[8,0],...,sudoku[8,8]])
```

Furthermore, with indices filled selectively and the remainder filled with underscores (_) the flattening applies only to the underscore indices:

```
alldiff(sudoku[4,_])
```

is equivalent to

```
alldiff([sudoku[4,0],...,sudoku[4,8]])
```

Lastly, one can optionally add square brackets ([]) around an expression to be flattened to make it look more like a vector:

```
alldiff([sudoku[4,_]])
```

is equivalent to

```
alldiff(sudoku[4,_])
```

Example

Additional hanging commas at the end of array are ignored, e.g.

```
lexleq([A,B,C,],[D,E,F,])
```

is equivalent to

```
lexleq([A,B,C],[D,E,F])
```

This feature is provided to make it easier to computer-generate input files.