

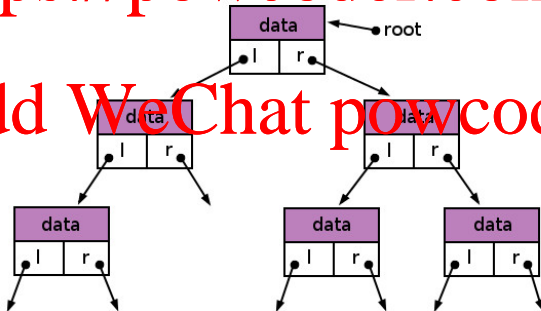
Dr Timothy Kimber

## Assignment Project Exam Help

January 2018

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# Dynamic Data Structures

Having efficient data structures is crucial for successful algorithms.

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- The problems seen so far involved fixed length lists
- In most languages we have a simple way to implement this efficiently — arrays
- Our algorithms assumed some sort of array type was available

Other problems require dynamic data structures such as

- Lists
- Stacks and Queues
- Sets and Dictionaries

These are designed to hold variable, essentially unlimited amounts of data.

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# Ordered Data Structures

A *list* is an ordered collection of {nodes, items, elements}.

- The key property of a list is the ordering of the nodes
- A list might support operations such as

`push` adds an element to the end of the list

`pop` removes the last element of the list

`shift` removes the first element of the list

`unshift` adds an element to the front of the list

`insert` adds an element at a given position

`remove` removes the element at a given position

`iterate` returns the items in order

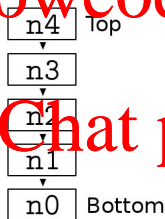
- Plus sorting, searching, copying, joining, splitting ...
- The most appropriate implementation depends on which operations are needed.

# Stacks

A *stack* is a **last-in first-out** (LIFO) list.

- Stacks support only
  - push** for adding elements
  - pop** for removing elements

- Stacks are usually pictured as a vertical (stacked!) structure



- Stacks support recursive algorithms including fundamental operations such as calling subprocedures and evaluating arithmetic expressions

## Stacks

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## Question

How would you implement a stack?

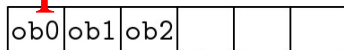
- Must be able to add “unlimited” objects
- Push and Pop must implement LIFO behaviour

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# Stack Implementation

Could use array or linked-list as storage model

- Both very simple implementation
- Array has fixed capacity
- Linked list has higher overheads



- Can make array **dynamic** (variable size)
- Integer `sp` points to the top of the stack
- Update `sp` within Push and Pop

## Dynamic Array-Based Stack

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```
Push(input: stack s[0, ..., k-1], object x)  
    if s.sp == k // array is full  
        s' = new stack of size 2k  
        for i = 0 to k-1  
            s'[i] = s[i]  
        s = s'  
    s[s.sp] = x  
    s.sp = s.sp + 1
```

- push increases the capacity of the stack if it is full

## Dynamic Array-Based Stack

Pop(input: stack s[0..k-1])

```
if s.sp < k/2 // array is less than half full
```

```
    s' = new stack of size k/2
```

```
    for i = 0 to k/2 - 1
```

```
        s'[i] = s[i]
```

```
    s = s'
```

```
s.sp = s.sp - 1
```

```
return s[s.sp]
```

- pop decreases the capacity if it is too big
- a full implementation should have a minimum size



# Performance of Push

## Question

Given a stack containing  $N$  objects, what is the worst case time complexity of **push**?

- Assume: time to insert (copy, add) one object to array is  $c$
- Assume: initial capacity is 4

The time taken for pushing objects is:

- $c, c, c, c, c + 4c, c, \dots$

So:

- Worst time to push a single object is  $Nc$
- So  $T(N) = O(N)$
- Want to reflect fact that most pushes are not  $O(N)$

## Performance of Push

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Revised Question  
Given an empty stack, what is the worst case time to push  $N$  objects?

- Assume: initial capacity is 4
- Assume: time to insert (copy and add) one object to array is  $c$

The time taken for the each push is still:

- $c, c, 4c, c+4c, c, \dots$
- For  $N$  pushes the worst single push is  $4c$
- $T(N) = O(N^2)$

However, this is a big overestimate

## Performance of Push

The time for  $N$  pushes is at most

$$T(N) = Nc + (4c + 8c + \dots + (N/2)c + Nc)$$

where

- the first  $Nc$  is the cost of writing  $N$  elements
- the rest is for copying to new arrays

The time for copying is  $2Nc - 4c$  (see notes on solving series) so

- $T(N) \leq 3Nc - 4c$
- $T(N) = O(N)$

# Amortisation

The time for  $N$  pushes is  $\mathcal{O}(N)$ , so:

- A single push is *effectively* a constant time operation
- More correctly: push is **amortised**  $\Theta(1)$
- NOT the same as  $\Theta(1)$

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## Amortisation

- Related to accountancy method used to defer large costs
- **Amortised analysis** considers a sequence of operations
- Cost of individual ops is “amortised” across the sequence
- Unlike accountancy, must never be in debt

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## Amortised Analysis

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Rather than calculating cost of full sequence of  $V$  steps we can

- Pick a **representative** subsequence
- Subsequence is some “cycle” that repeats
- Pick an amortised cost for operations
- Show that paying amortised cost covers all costs (never in debt)

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### Exercise

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Find a representative cycle (subsequence) of pushes into the stack and show that the amortised cost of  $3c$  covers all costs.

## Amortised Analysis

- Start **after** any expensive push
- Up to and including next expensive one

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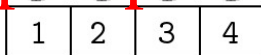
- Cheap pushes each put  $2c$  in the bank
- Have enough to cover extra costs when next expensive push occurs
- (If you started with a copy you are immediately over budget)

# Amortised Analysis

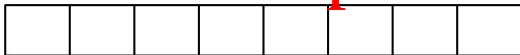
Argument only works because array is initially empty and size is **doubled**

- Say we have  $N$  objects on stack after a copy
- Before next copy we always push  $N$  more
- This is how cost is covered

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- Multiplying by **any factor** will do - will affect amortisation constant

## Queues

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- The **earliest** one added (FIFO Queue)
- The one with highest **priority** (Priority Queue)

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## Questions

- How could you implement a priority queue (PQ)?
- Given a PQ following your design that contains  $N$  objects, what would be the worst case time to add a new object? (Each object has a key attribute that determines its priority.)

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# Priority Queue Design

If we maintain a total ordering of the queue:

- It will take  $O(N)$  time to add new elements
- Can search a sorted array quickly but have to shift existing objects
- Finding position in a linked list is  $O(N)$

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9, 8, 7, 6, 3  
←  
remove items

Max Priority Queue

← 6, 9, 7, 3, 8  
add items

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Do not actually need total ordering.

- Queue does not support indexed access
- Just want to find object with highest priority

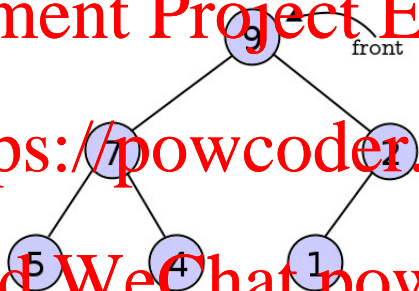
# Priority Queue Design

Solution is to divide and conquer the data

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- **Key Property:** Maintain order within each branch
- Highest (or lowest) key will be at the root
- Behaves like lots of mini queues

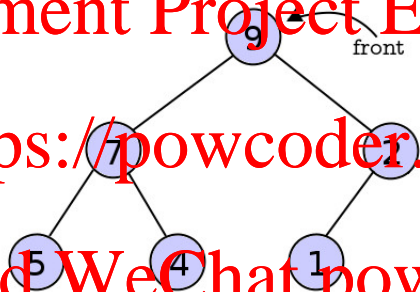
## Priority Queue Design

Solution is to divide and conquer the data

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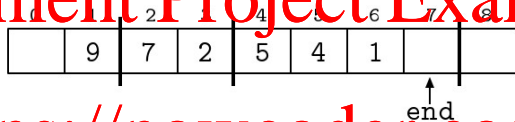
### Question

A new object could go in any branch. (Do you agree?) So, where should it go? Why?

## Heap: a Tree in an Array

We want to know where the “end” of the tree is:

- Build a tree within an array



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- Track end using “stack pointer”
- Navigate by indices
- Leaving a [0] blank means:

- parent of  $a[n]$  is  $a[n/2]$
- children of  $a[n]$  are  $a[2*n]$  and  $a[2*n+1]$

### Exercise

How should a new object be added to a **max** binary heap? (i.e. the greatest key should be at the root).

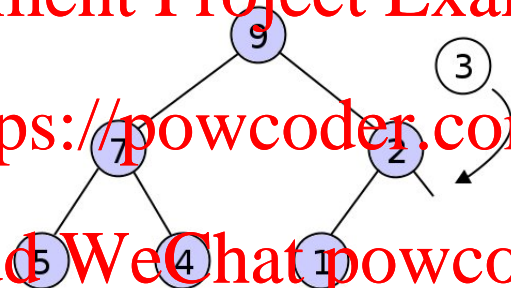
# Binary Heap Operations

Adding an object to the heap

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- Restore the “shape”
- Then restore the order

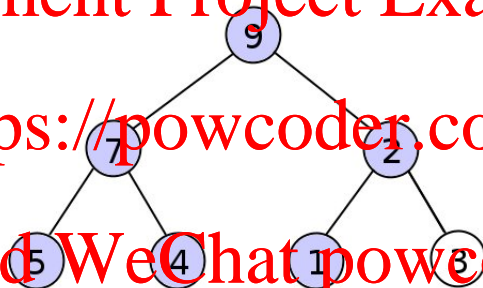
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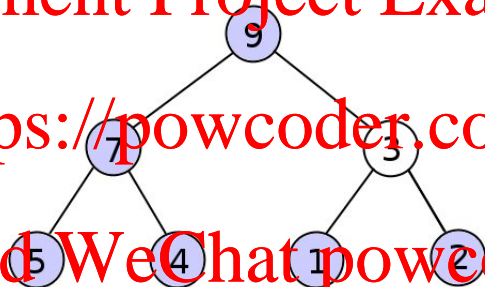
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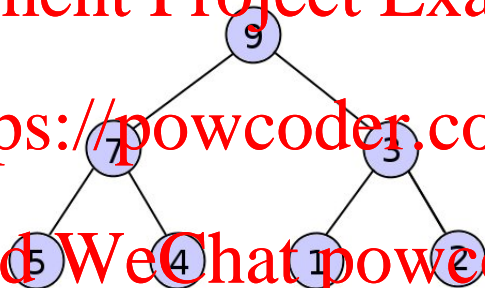
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Adding an object to the heap

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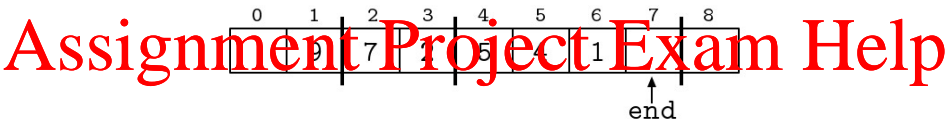
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- Restore the “shape”
- Then restore the order



## Heap: a Tree in an Array



- <https://powcoder.com>
- Track end using “stack pointer”
  - Leaving  $a[0]$  blank means:

- parent of  $a[i]$  is  $a[\lfloor i/2 \rfloor]$
  - children of  $a[i]$  are  $a[2*i]$  and  $a[2*i+1]$
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### Exercise

How should the object with the greatest key be removed from a **max** binary heap?

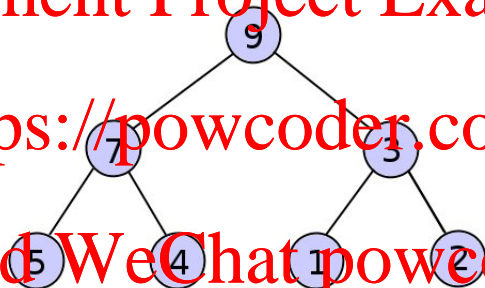
# Binary Heap Operations

To remove the object at the root

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- Restore the “shape”
- Then restore the order

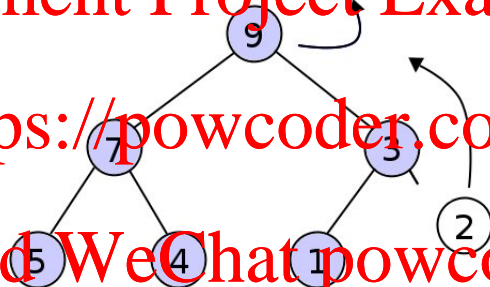
# Binary Heap Operations

To remove the object at the root

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- Restore the “shape”
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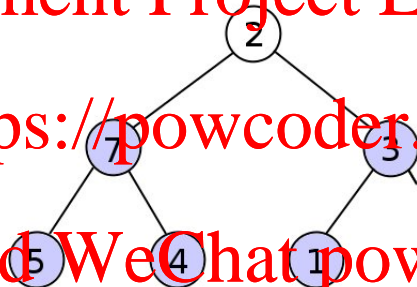
## Binary Heap Operations

To remove the object at the root

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- Restore the “shape”
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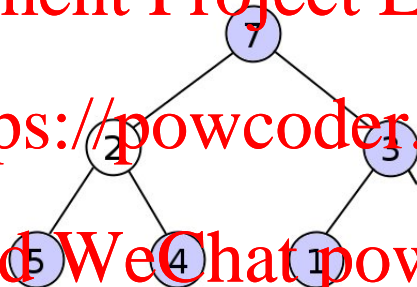
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To remove the object at the root

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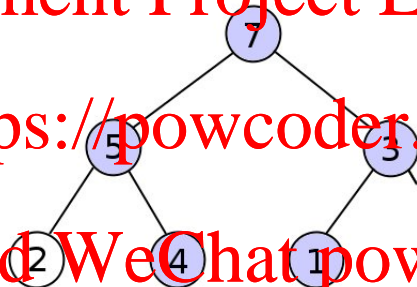
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To remove the object at the root

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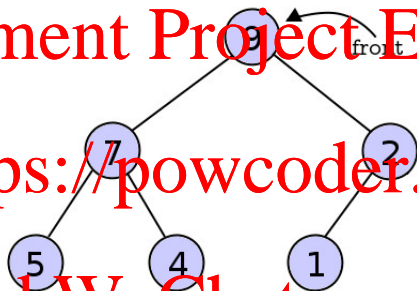
- Restore the “shape”
- Then restore the order

## Binary Heap Performance

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### Question

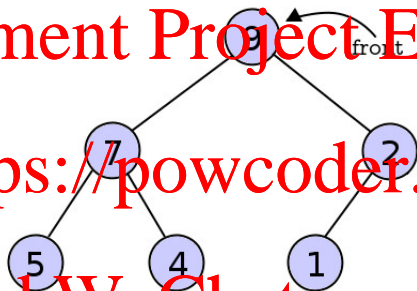
Given a heap containing  $N$  objects, what is the time complexity for adding or removing one object?

# Binary Heap Performance

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Both operations are  $O(\log_2 N)$

- Height of the heap is  $\Theta(\log_2 N)$
- Each operation confined to one branch



# Heapsort

Heaps also provide us with the **Heapsort** algorithm (JWJ Williams, 1964)

Heapsort (given a list  $L$ )

- Create an empty heap  $H$
- Remove each element of  $L$  and add it to  $H$
- Remove each element of  $H$  and add it to  $L$
- HALT

- What could be simpler?
- Performance is again  $\Theta(N \log_2 N)$
- Can also be implemented **in place** by setting up list and heap partitions within a single array