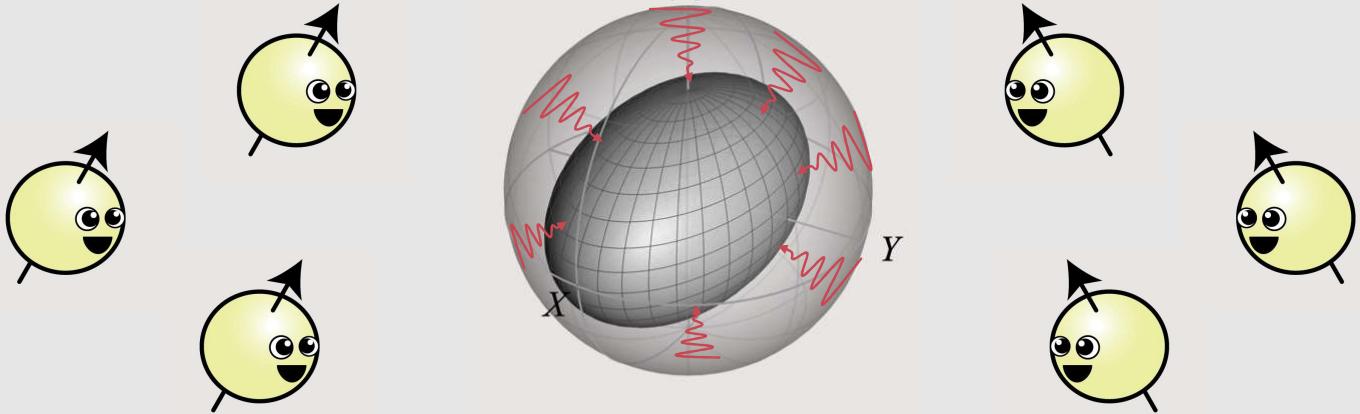


Introduction to Quantum Noise

IBM Quantum

Qiskit Global Summer School: Quantum Simulations



Zlatko K. Minev

IBM Quantum

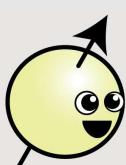


@zlatko_minev



zlatko-minev.com

What do I need to know before
I run quantum simulation on a
real, noisy quantum processor?





Chapter 1: Hello World!

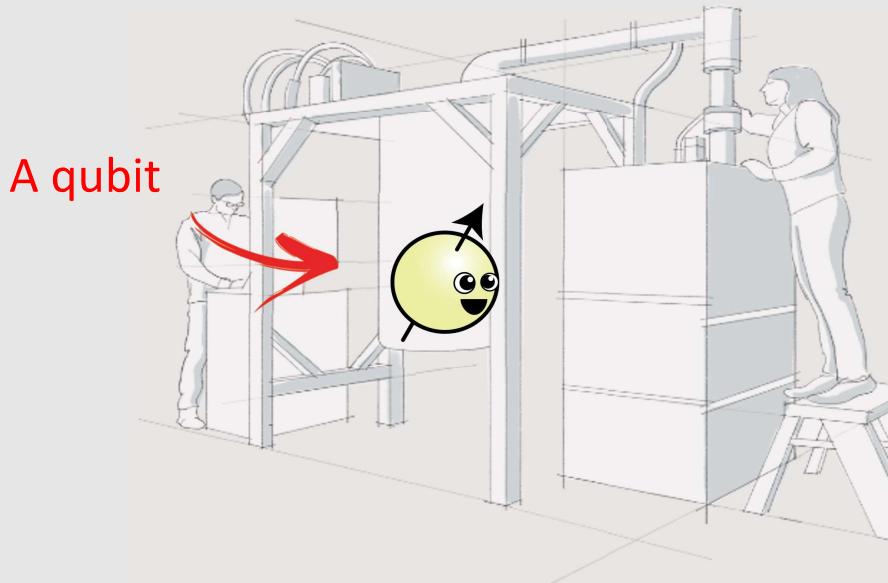
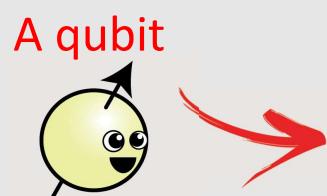
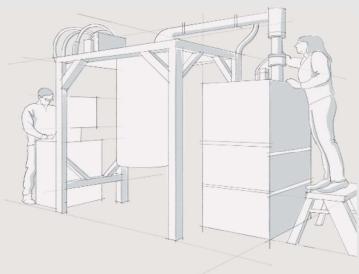


image: qiskit

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Hello World! building blocks



$|1\rangle$

$|0\rangle$

Computational
basis states

image: qiskit

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$$X |0\rangle = |1\rangle$$
$$X |1\rangle = |0\rangle$$

Zlatko Minev, IBM Quantum (5)

$$X |0\rangle = |1\rangle$$
$$X |1\rangle = |0\rangle$$

Zlatko Minev, IBM Quantum (6)



Hello World! qubit flipper quantum circuits



refresher:

Zlatko Minev, IBM Quantum (7)



Hello World! “debugger” step through

$$X |0\rangle = |1\rangle$$

$$X |1\rangle = |0\rangle$$

Zlatko Minev, IBM Quantum (8)



Hello World! “debugger” step through

$$X |0\rangle = |1\rangle$$

$$X |1\rangle = |0\rangle$$

Zlatko Minev, IBM Quantum (9)



Hello World! “debugger” step through

$$X |0\rangle = |1\rangle$$

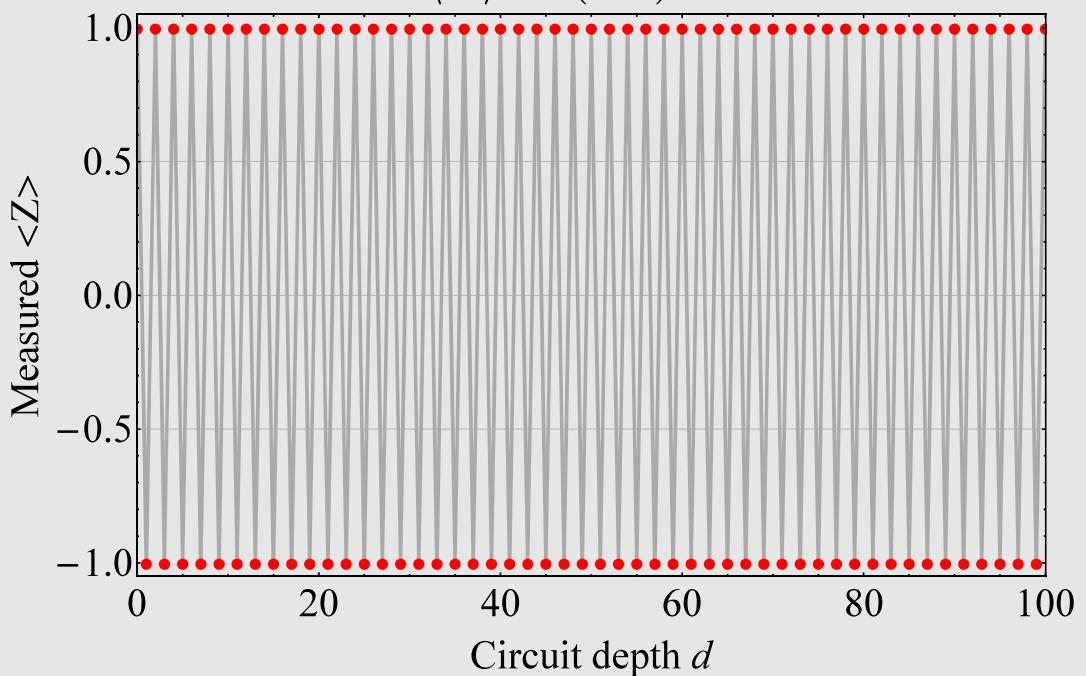
$$X |1\rangle = |0\rangle$$

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Hello World! Ideal expectation results

$$\langle Z \rangle = (-1)^d$$

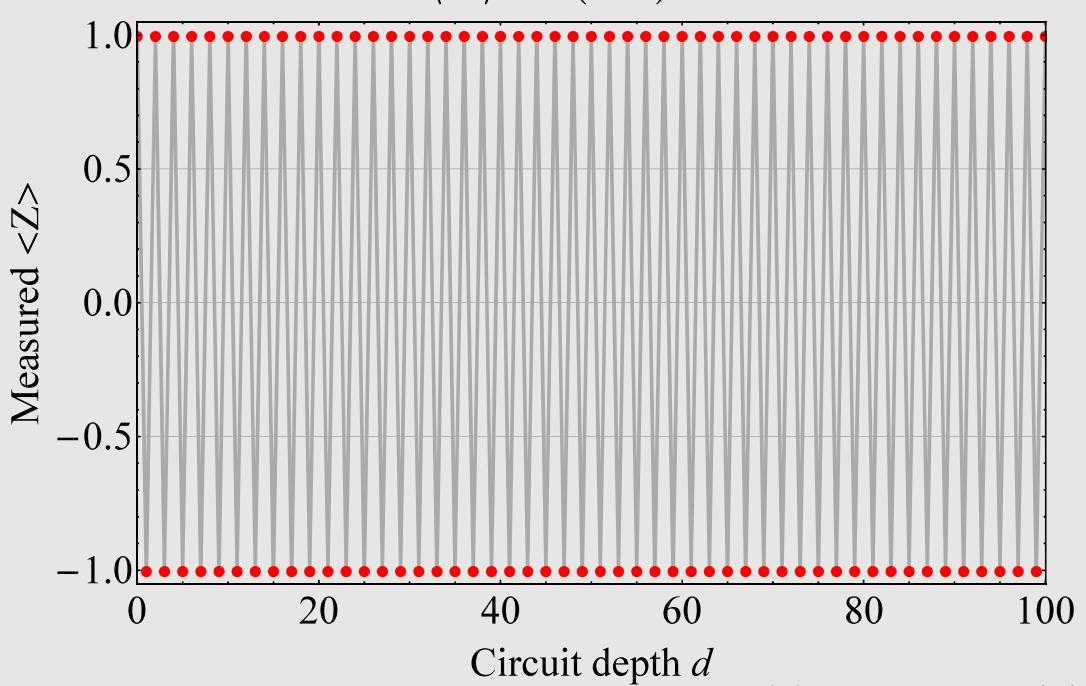


Zlatko Minev, IBM Quantum (11)



Hello World! Ideal expectation results

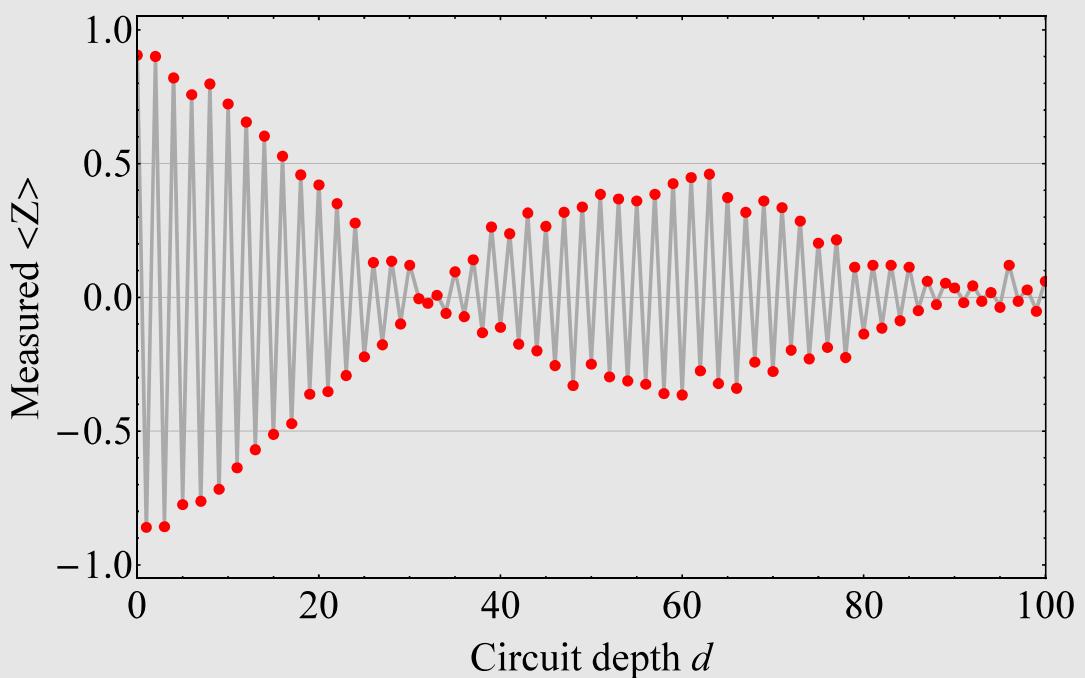
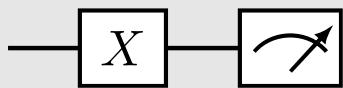
$$\langle Z \rangle = (-1)^d$$



Zlatko Minev, IBM Quantum (12)



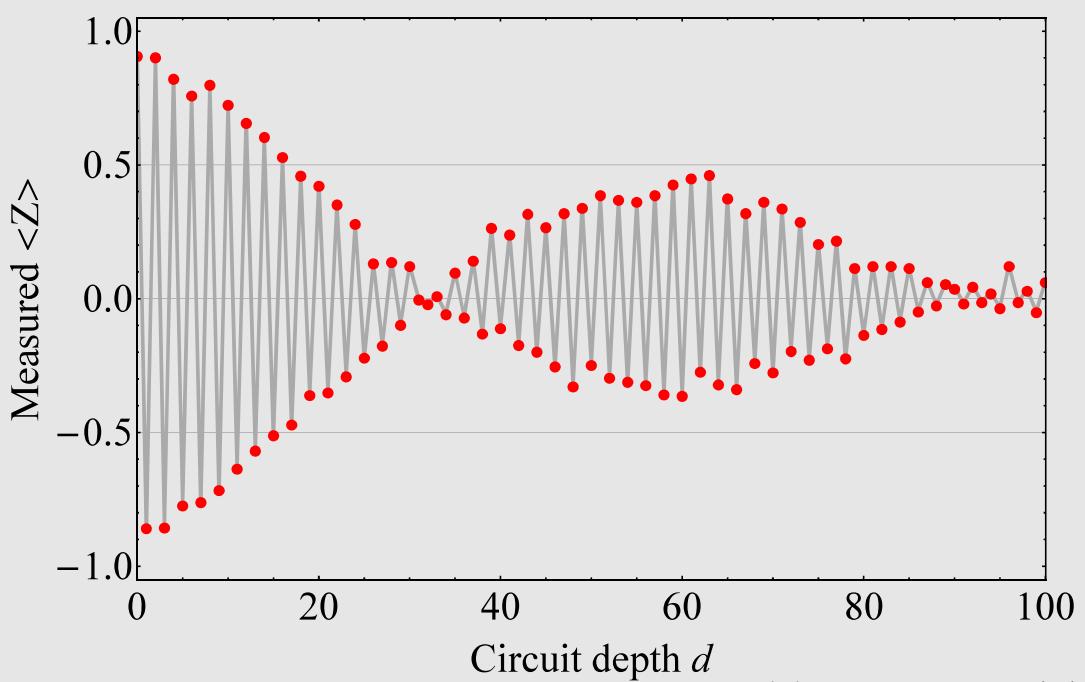
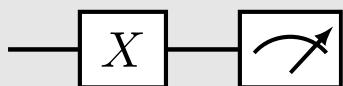
Hello World! Real expectation results



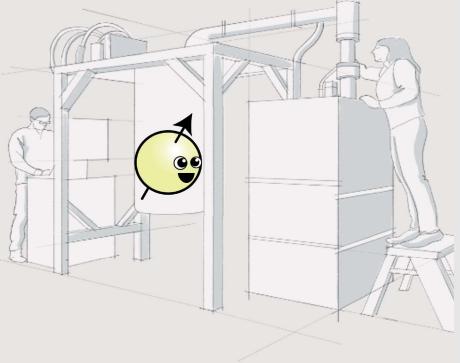
Zlatko Minev, IBM Quantum (13)



Real & noisy quantum processors: Why study noise?



Zlatko Minev, IBM Quantum (14)

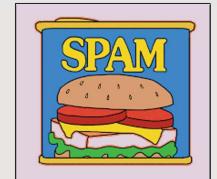
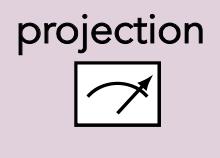
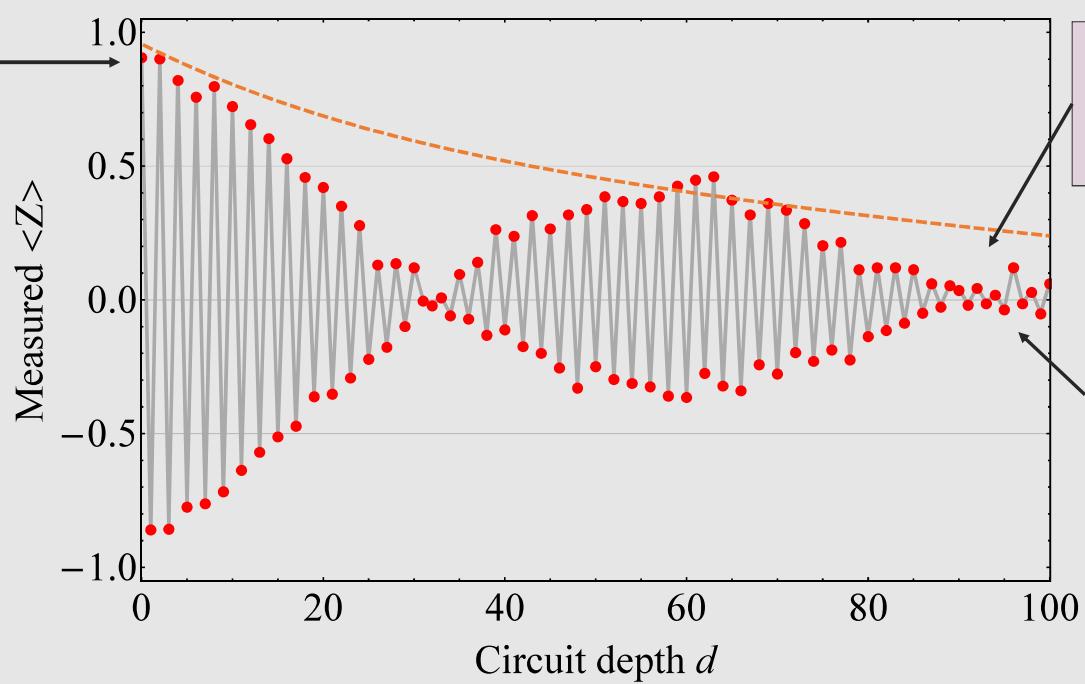
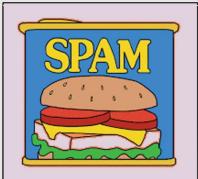


“Quantum phenomena
do *not* occur in a Hilbert space,
they occur in a laboratory.”

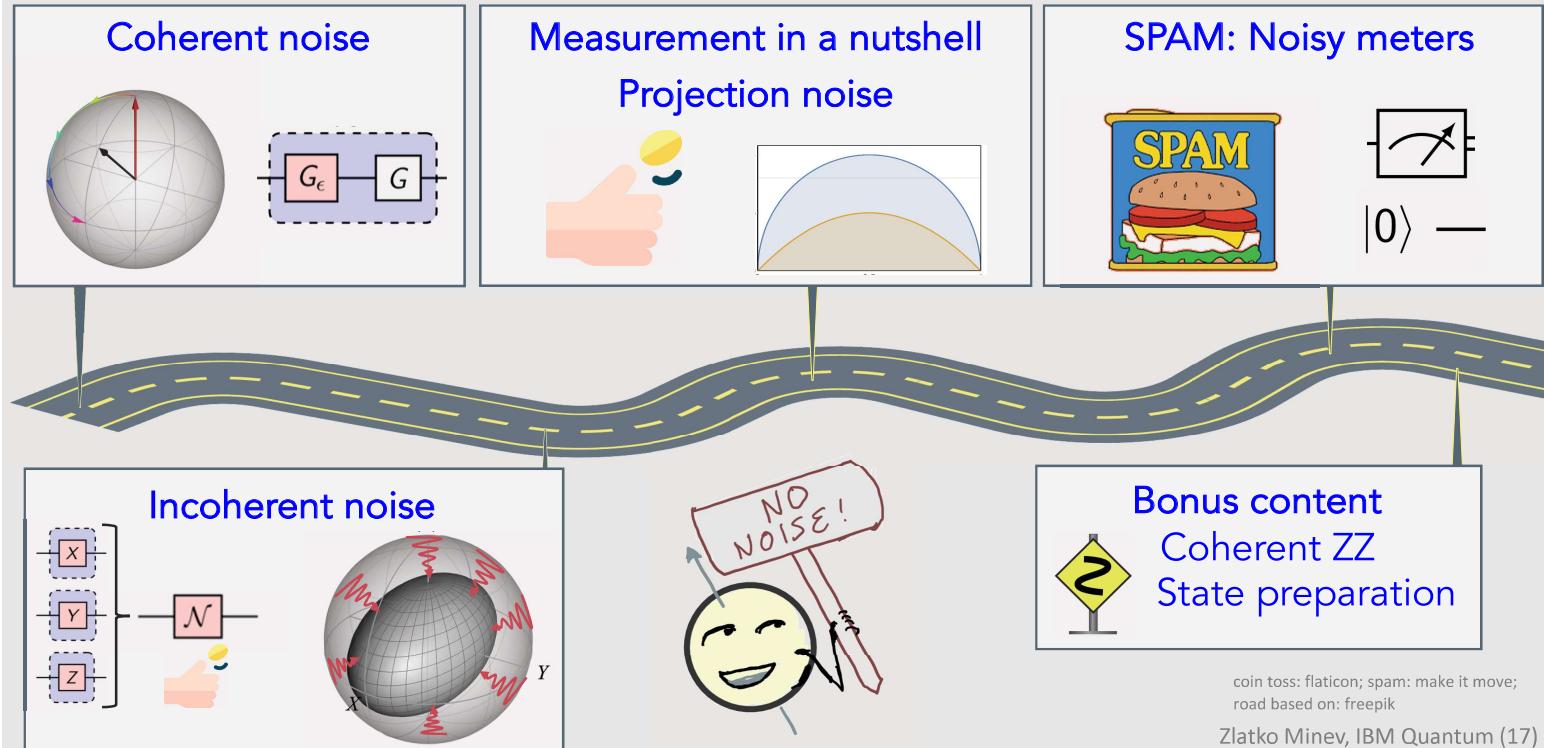
Asher Peres



Elements of noise



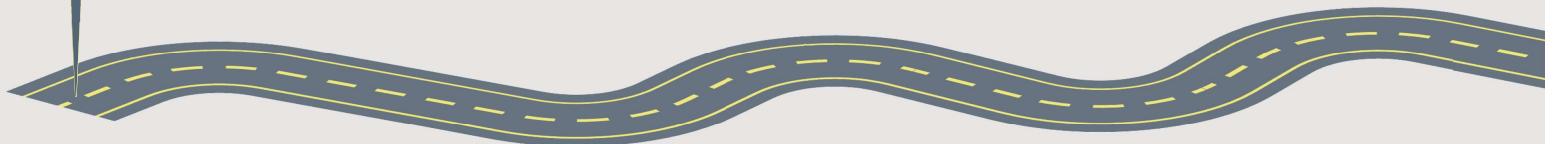
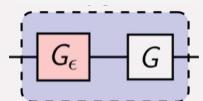
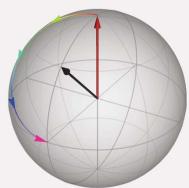
The road ahead



The destination is the journey

Chapter 2: Coherent noise

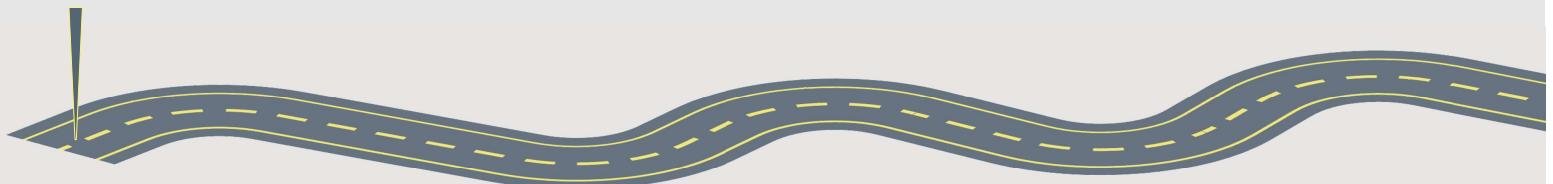
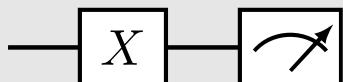
Coherent



road based on: freepik

Zlatko Minev, IBM Quantum (19)

Return to the Hello World example

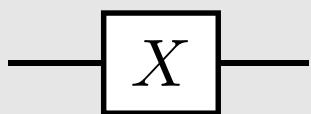


road based on: freepik

Zlatko Minev, IBM Quantum (20)



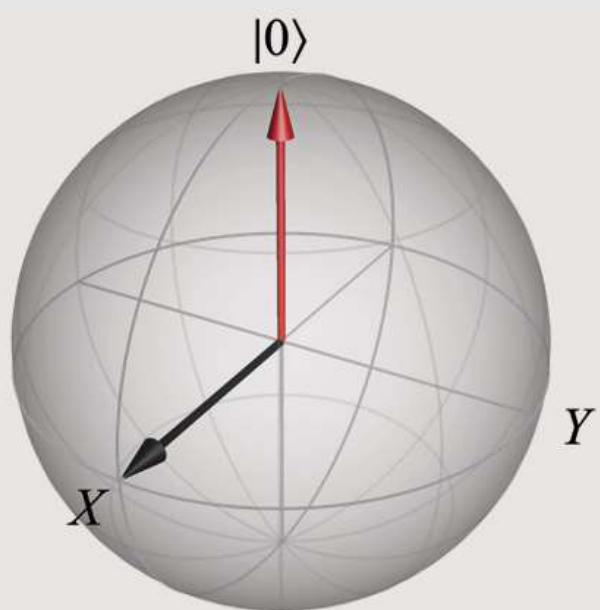
Origin of our X gate: time evolution



Refresher:

Zlatko Minev, IBM Quantum (21)

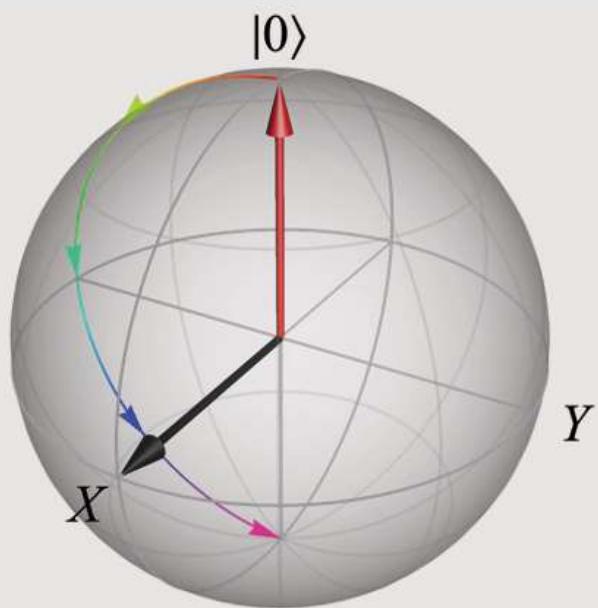
Visualize: Bloch sphere



Zlatko Minev, IBM Quantum (22)

Visualize: Evolution on the Bloch sphere

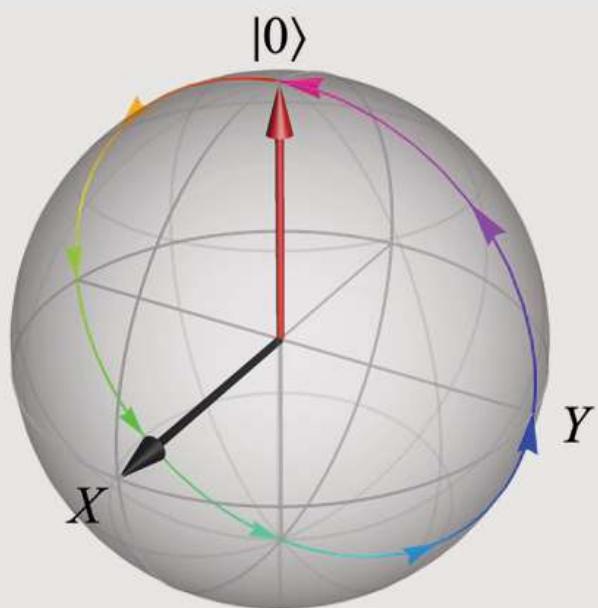
$$R_X(\theta) = \exp\left(-\frac{i\theta}{2}X\right)$$



Zlatko Minev, IBM Quantum (23)

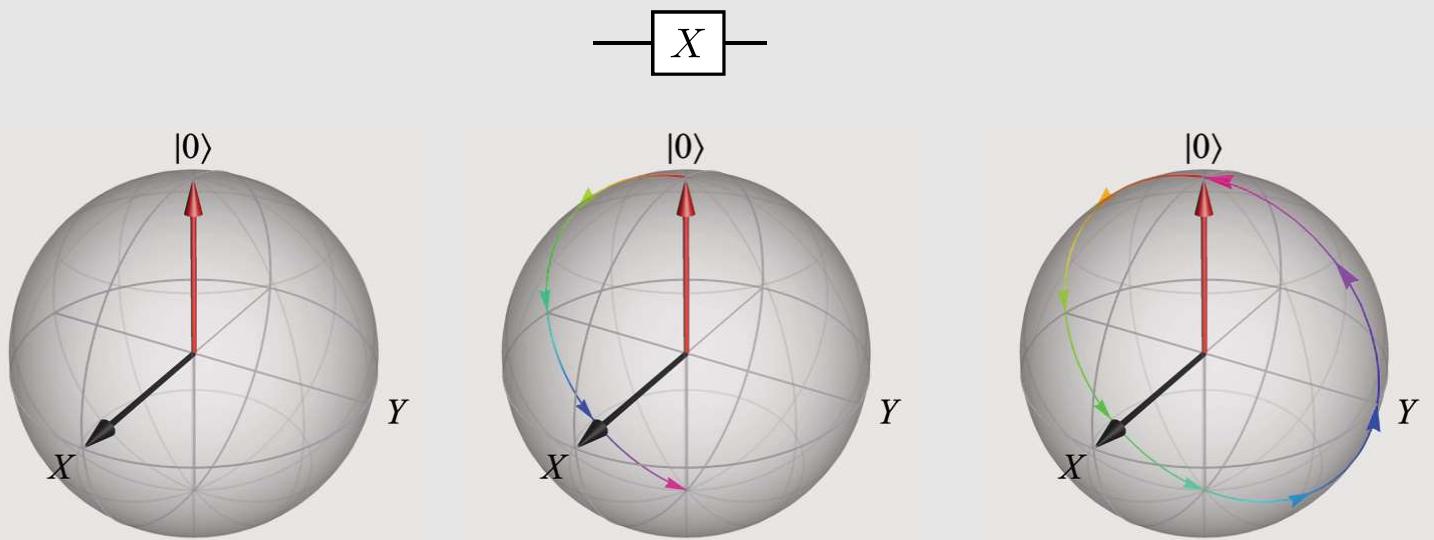
Visualize: Evolution on the Bloch sphere

$$R_X(\theta) = \exp\left(-\frac{i\theta}{2}X\right)$$



Zlatko Minev, IBM Quantum (24)

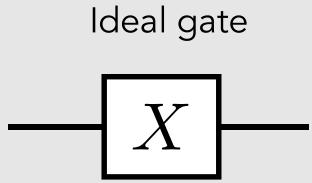
Evolution on the Bloch sphere



Zlatko Minev, IBM Quantum (25)



Miscalibrated gate

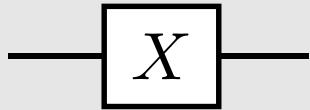


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Miscalibrated gate

Ideal gate

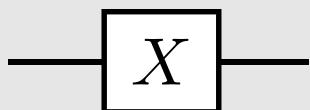


Zlatko Minev, IBM Quantum (27)



Noisy gate decomposition

Ideal gate



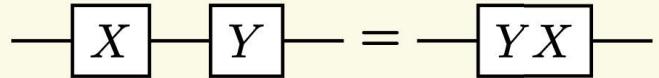
Zlatko Minev, IBM Quantum (28)



Careful

⚠ Common pitfall

The order in which gates appear in a schematic is the reverse of how they appear in the algebra.



Zlatko Minev, IBM Quantum (29)



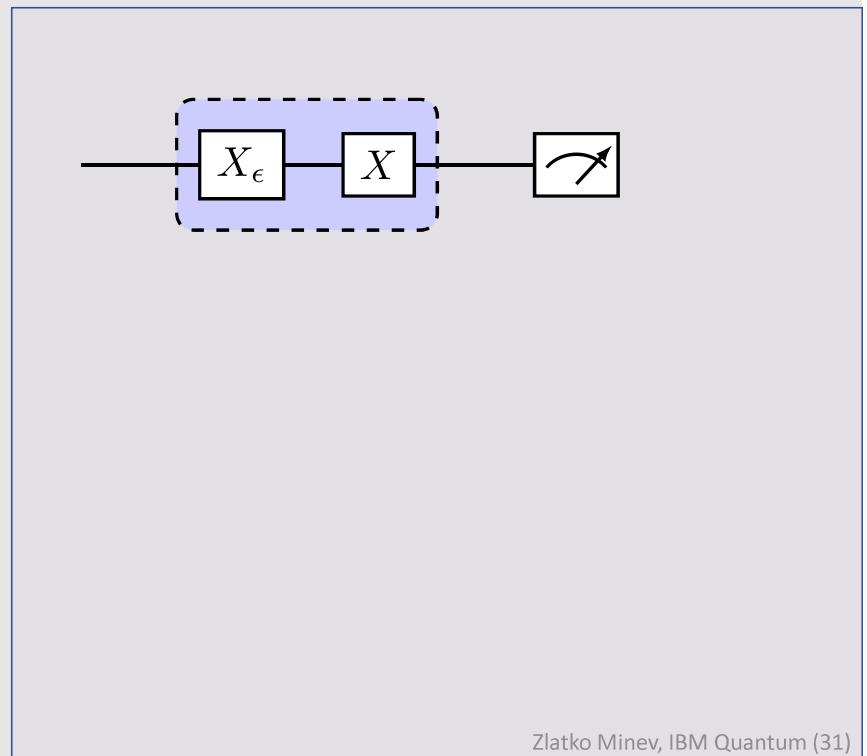
Using a noisy gate in a quantum circuit

$$X = R_X(\pi)$$

Zlatko Minev, IBM Quantum (30)



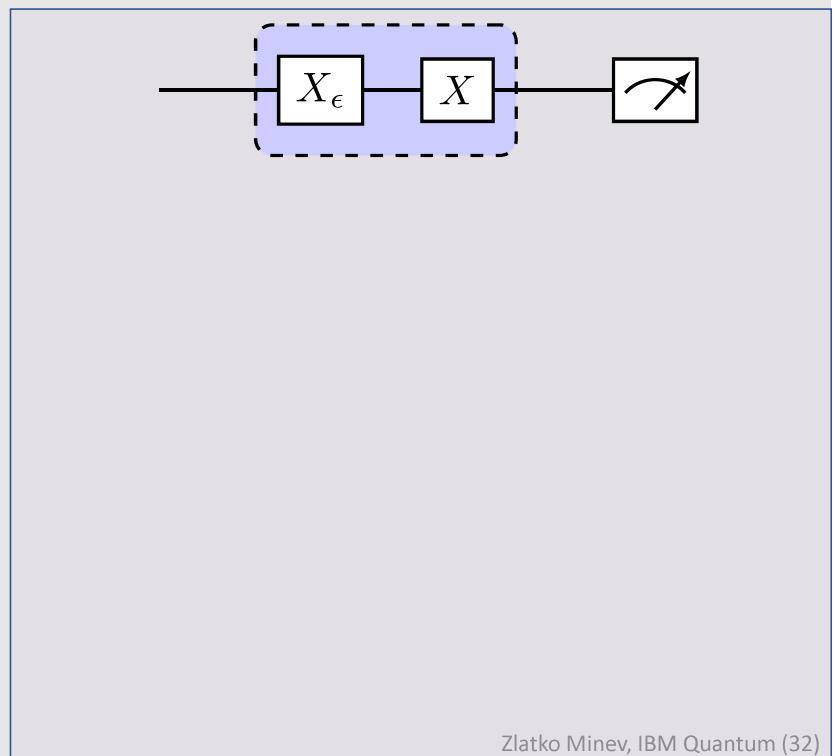
Using a noisy gate in a quantum circuit: final state



Zlatko Minev, IBM Quantum (31)



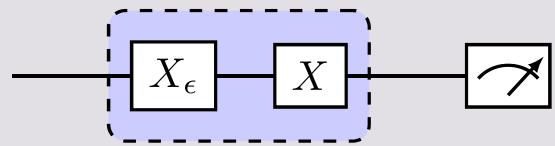
Ideal vs. noisy observable



Zlatko Minev, IBM Quantum (32)



Compare to full experiment



Zlatko Minev, IBM Quantum (33)

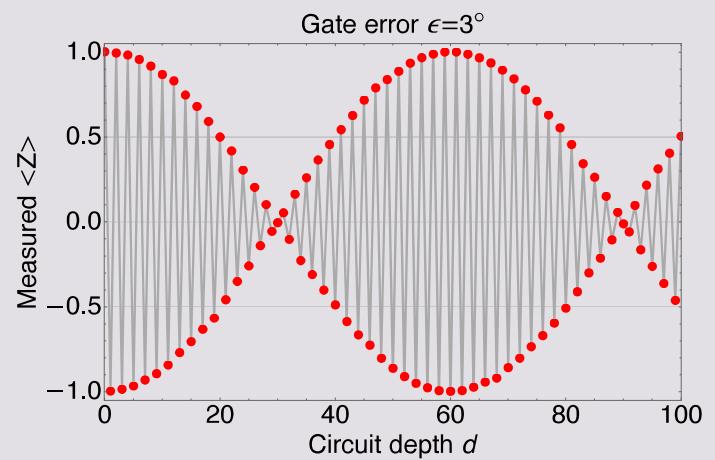
$$\langle \tilde{\psi}_f | Z | \tilde{\psi}_f \rangle = \cos(d\pi + d\epsilon)$$

Zlatko Minev, IBM Quantum (34)



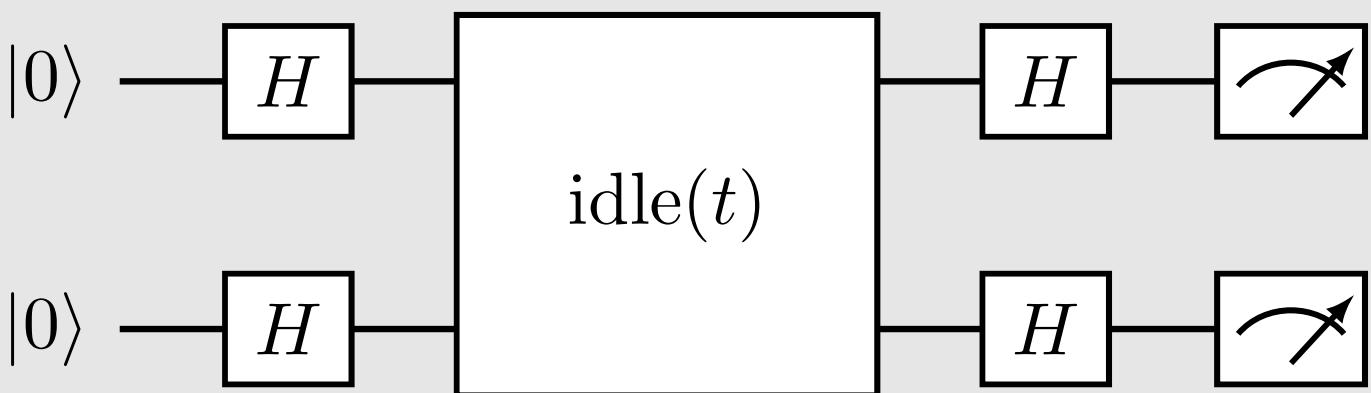
Coherent errors: brief summary

- are ubiquitous
- can be described by unitary operations
- do not loose quantum information
- data: can create oscillations in the data
- data: do not yield exponential decays
- have a quadratic impact on algorithmic accuracy (worst-case error)



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Bonus content: two-qubit coherent ZZ error



Zlatko Minev, IBM Quantum (36)



Questions

Answer these multiple-choice questions
in the chat; for example, type “1a 2b.”

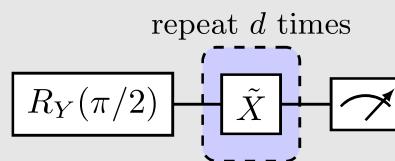
1. Coherent noise can be caused by

- a) loss of energy of the qubit
- b) miscalibration, such as over-rotation
- c) wanted coupling to neighboring qubit

2. Coherent noise can be really bad because

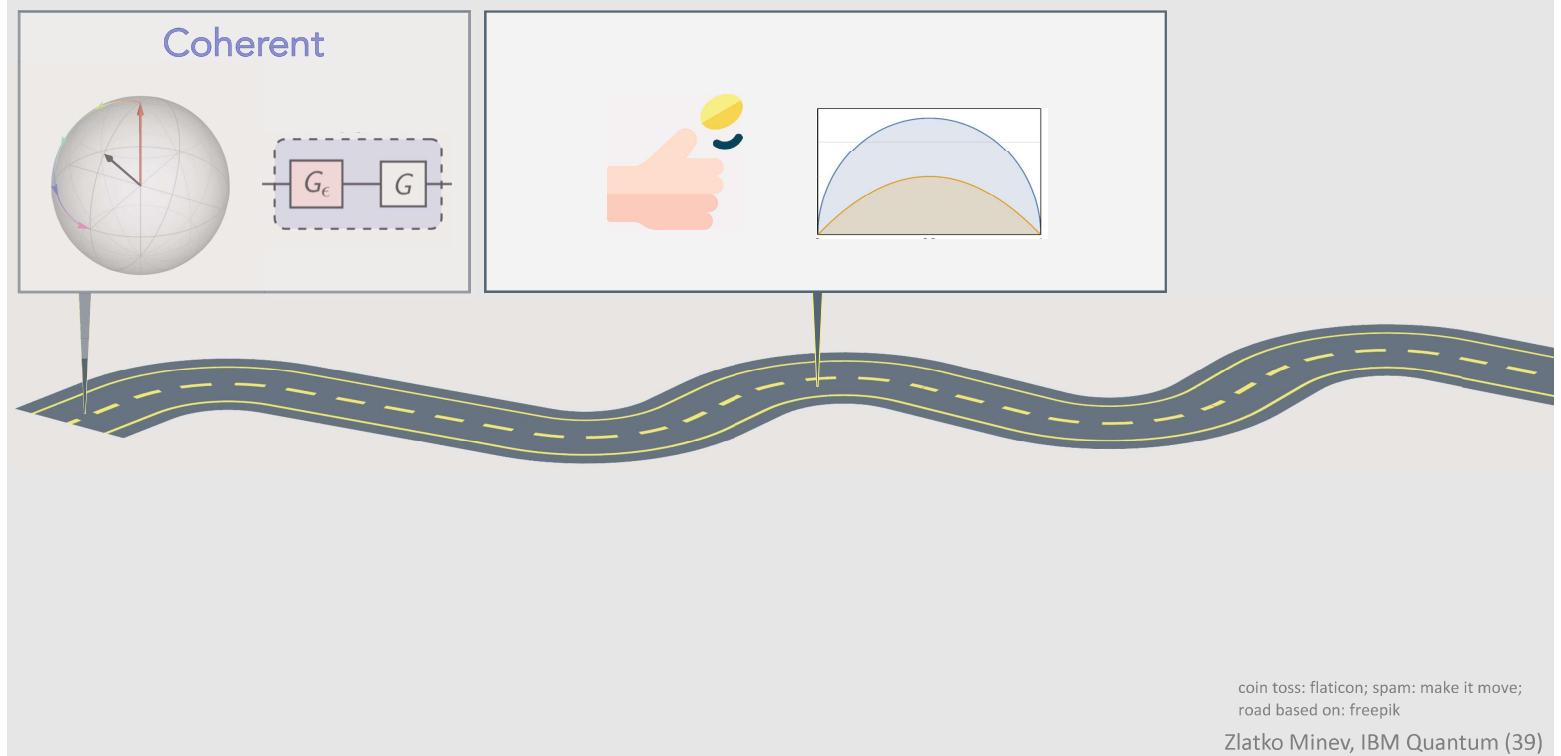
- a) it results in loss of information
- b) you cannot undo it
- c) the worst-case error often grows quadratically

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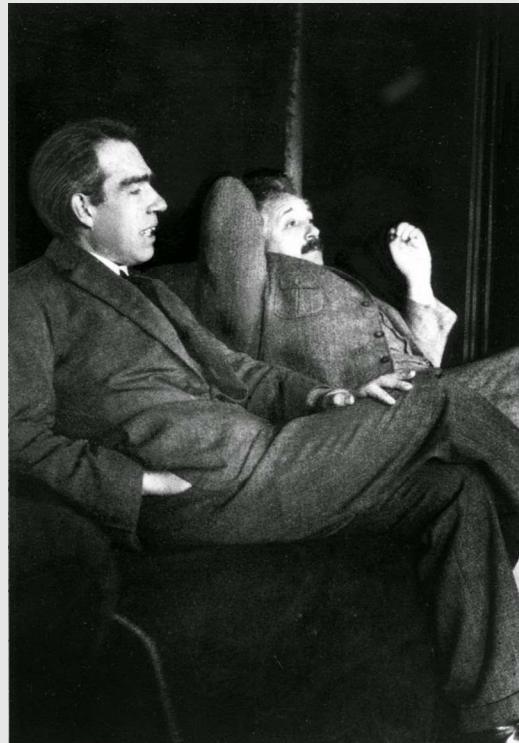


Zlatko Minev, IBM Quantum (38)

Chapter 3: Measurement theory in a nutshell + projection noise



Quantum Measurement



Bohr and Einstein (1925)
Source: Wikimedia

Zlatko Minev, IBM Quantum (40)



Measurement apparatus: general

The diagram illustrates the Schrödinger's cat thought experiment. On the left, two states of a cat are shown: a vertical superposition state $\frac{1}{\sqrt{2}} |\text{alive}\rangle + |\text{dead}\rangle$ and a horizontal state $|\psi\rangle$. A horizontal arrow points from the superposition state towards a central box. The box contains a curved arrow pointing from left to right, representing a measurement process. Two parallel horizontal lines extend from the right side of the box, followed by an equals sign and the letter M , indicating the measurement operator M .

Quantum

Classical

cat image: docencia

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Measurement basis

Computational basis

$$|0\rangle, |1\rangle$$

$$|\psi\rangle \xrightarrow{\hspace{1cm}} \text{[Quantum Circuit with two qubits and three gates]} \xrightarrow{\hspace{1cm}} M$$

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Resolution of the identity

$$|\psi\rangle \xrightarrow{\quad} M$$

| outcome | outcome probability |
|---------|--|
| $M = 0$ | $: p(M = 0) = \langle 0 \psi\rangle ^2$ |
| $M = 1$ | $: p(M = 1) = \langle 1 \psi\rangle ^2$ |

Zlatko Minev, IBM Quantum (43)



Resolution of the identity

$$|\psi\rangle \xrightarrow{\quad} M$$

| outcome | outcome probability |
|---------|--|
| $M = 0$ | $: p(M = 0) = \langle 0 \psi\rangle ^2$ |
| $M = 1$ | $: p(M = 1) = \langle 1 \psi\rangle ^2$ |

$$\begin{aligned} & p(M = 0) + p(M = 1) \\ &= |\langle 0|\psi\rangle|^2 + |\langle 1|\psi\rangle|^2 \\ &= \langle\psi|0\rangle\langle 0|\psi\rangle + \langle\psi|1\rangle\langle 1|\psi\rangle \\ &= \langle\psi|(|0\rangle\langle 0| + |1\rangle\langle 1|)|\psi\rangle \\ &= \langle\psi|\hat{I}|\psi\rangle \\ &= ||\psi\rangle|^2 \\ &= 1 \end{aligned}$$

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Projects and overlaps



| | |
|---------|---|
| outcome | outcome probability |
| $M = 0$ | $: p(M = 0) = \langle 0 \psi \rangle ^2$ |
| $M = 1$ | $: p(M = 1) = \langle 1 \psi \rangle ^2$ |

$$p(M = 0) + p(M = 1)$$

$$= \langle \psi | (|0\rangle\langle 0| + |1\rangle\langle 1|) |\psi \rangle$$

projectors

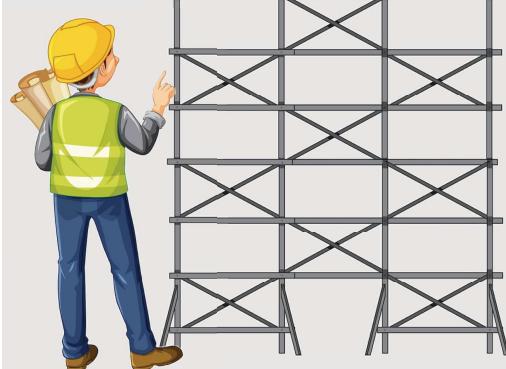
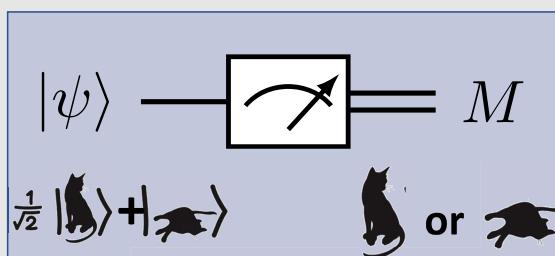
$$\hat{\Pi}_0 = |0\rangle\langle 0| \quad \hat{\Pi}_1 = |1\rangle\langle 1|$$

inner-product in operator space (overlap between states)

$$p(m) = \langle \hat{\Pi}_m, |\psi \rangle \rangle = \text{Tr} (\hat{\Pi}_m |\psi \rangle \langle \psi |)$$

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Measurement theory 101 summary: Scaffolding in terms of measurement operators



Measurement aspect

Measurement outcome M

Example

$M = 0$

Set of measurement outcomes Σ

$\Sigma = \{0, 1\}$

Measurement operator Π_M

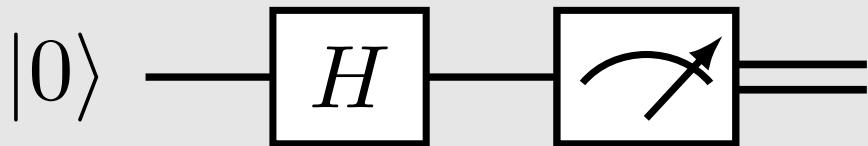
$\Pi_M = |0\rangle\langle 0|$

Probability of measuring $M = m$

$$p(M=m) = \langle \Pi_m, \psi \rangle = |\langle \Pi_m | \psi \rangle|^2$$



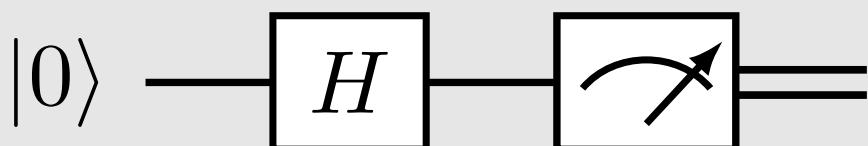
Putting it to use: example ideal circuit



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Example ideal circuit



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Mean

Statistics

$$\begin{aligned}
 \mathbb{E}[M] &= \sum_m m P(m) = 0 P(m=0) + 1 P(m=1) \\
 &\stackrel{\text{classical}}{\approx} P(M=1) \\
 &= p \\
 &= \sum_m m \langle \hat{M}_m \rangle \\
 &= \sum_m m \langle |m\rangle \langle m| \rangle \\
 &= \langle \sum_m m |m\rangle \langle m| \rangle \\
 &= \langle \hat{M} \rangle
 \end{aligned}$$

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$$\begin{aligned}
 \mathbb{V}[M] &= \mathbb{E}[M^2] - \underbrace{\mathbb{E}[M]^2}_{\text{Quantum}} = \langle \hat{M}^2 \rangle - \langle \hat{M} \rangle^2 \\
 \mathbb{E}[M^2] &= \sum_m m^2 P(m) \\
 &= \sum_m m^2 \langle |m\rangle \langle m| \rangle \\
 &= \langle \sum_m m^2 |m\rangle \langle m| \rangle \\
 &= \langle \hat{M}^2 \rangle \\
 &= \cancel{p^2(1-p)} + 1^2 p \\
 &= p \\
 \mathbb{V}[M] &= p - p^2 \\
 &\approx p(1-p) &= 0 & \text{if } p=0 \\
 &= \sigma_M^2 &= 1 & \text{if } p=1
 \end{aligned}$$

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Different observables

Different observables

$$\text{For example } \hat{M} = 0|0\rangle\langle 0| + 1|1\rangle\langle 1| = \frac{1}{2}(\hat{I} + \hat{Z}) = \begin{pmatrix} 0 & 0 \\ 0 & 1 \end{pmatrix}$$

$$\hat{M} = (1|0\rangle\langle 0| - 1|1\rangle\langle 1|) = \hat{Z} = \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix}$$

For $\hat{M} = \hat{Z}$

$$\Sigma = \{-1, +1\}$$

Pauli Observables

$\hat{M} = \hat{X}$

$$\Sigma = \{+, -\}$$

$$\hat{\Pi}_+ = |+\rangle\langle +|$$

$$\hat{\Pi}_- = |- \rangle\langle -|$$

$$\hat{X} = \sum_{m=1}^{\infty} m |m\rangle\langle m|$$

MEG

$$= |+1\rangle\langle +1| - |-1\rangle\langle -1|$$

$$\hat{X}|+\rangle = +1|\times\rangle$$

$$\hat{X}|-\rangle = -1|\times\rangle$$

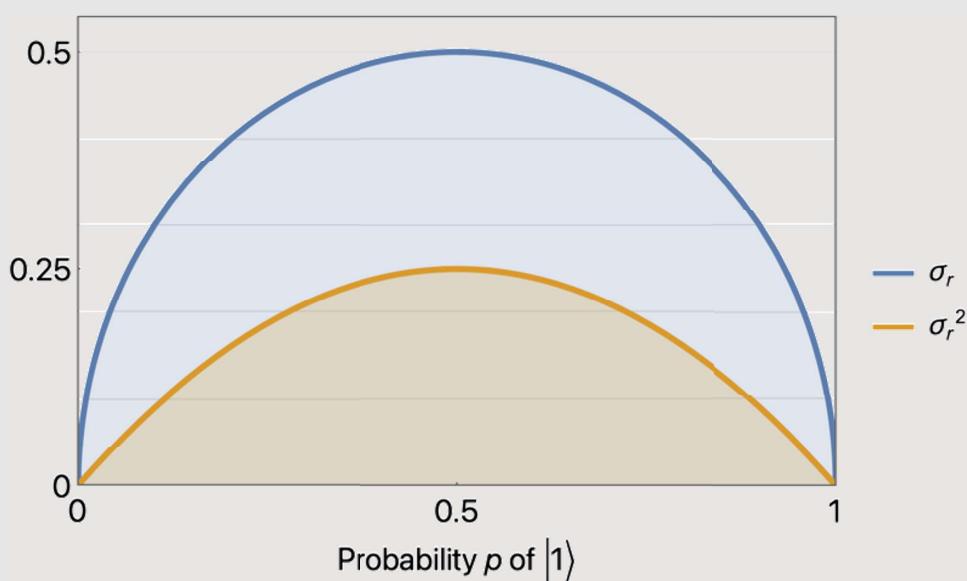
$$|+\rangle = \frac{1}{\sqrt{2}}(|1\rangle)$$

$$|- \rangle = \frac{1}{\sqrt{2}}(|-1\rangle)$$

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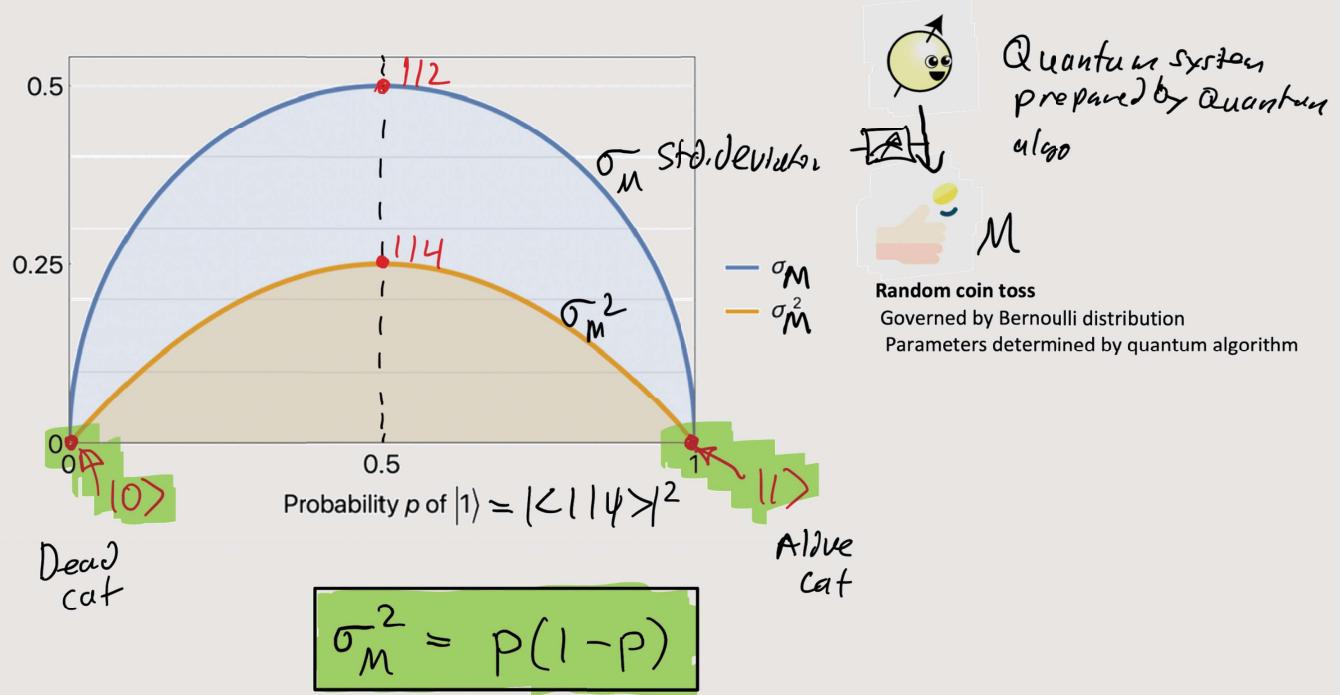
C1 C2 C3 C4 C5

Variance of the random classical variable vs. probability to obtain 1



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Variance of the random classical variable vs. probability to obtain 1



Zlatko Minev, IBM Quantum (54)

Shots, shots, shots

shot 1: $|0\rangle \xrightarrow{H} \xrightarrow{\text{meter}} m_1 = 1$

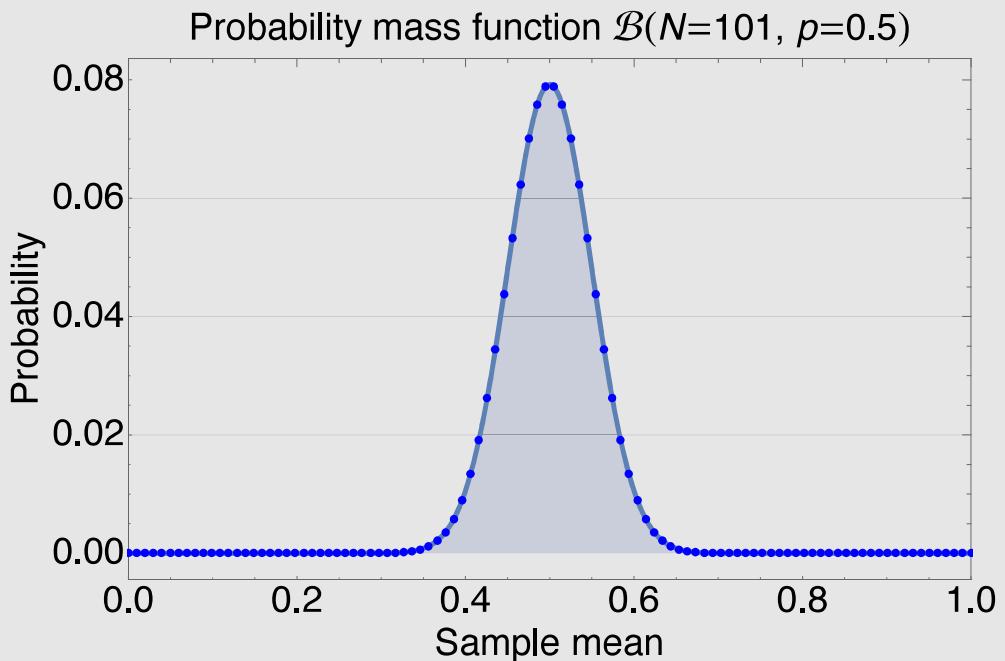
shot 2: $|0\rangle \xrightarrow{H} \xrightarrow{\text{meter}} m_2 = 0$

shot 3: $|0\rangle \xrightarrow{H} \xrightarrow{\text{meter}} m_4 = 1$

shot 4: $|0\rangle \xrightarrow{H} \xrightarrow{\text{meter}} m_3 = 0$

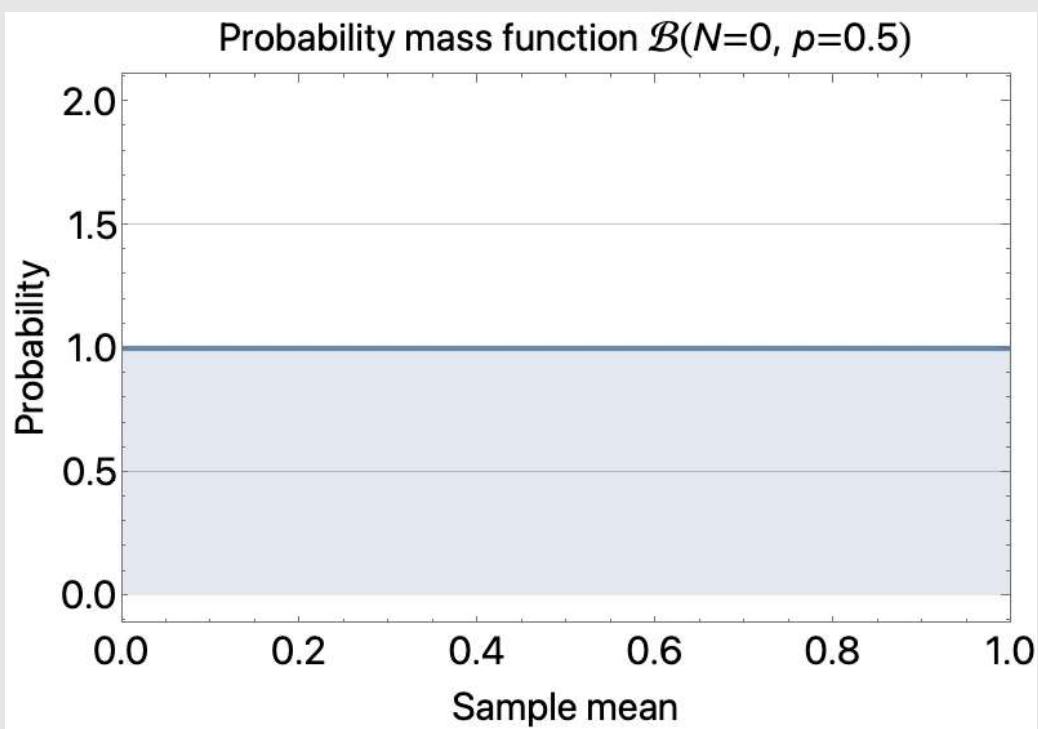
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Sampled output distribution



Zlatko Minev, IBM Quantum (56)

Animation of convergence of shots expectation value and mean



Zlatko Minev, IBM Quantum (57)

Concentration inequalities and tail bounds

*Making a list,
checking it twice,
going to see
which inequality
is nice!*

*Markov? Hoeffding?
Jensen? Chebyshev?
Chernoff?*

1. Probability (Technical note 11.9 v0.6)

1A. Concentration inequalities and tail bounds

Unless otherwise specified, all variables are real \mathbb{R} . Inequalities come as one-sided $\Pr(\dots \leq \dots)$ and two-sided $\Pr(|\dots| \leq \dots)$. Notation: X is a random variable, $\mu := \mathbb{E}[X]$, $\sigma^2 := \text{Var}[X]$, $S_n := X_1 + \dots + X_n$.

| Inequality | Conditions | Common form | Notes / Alternate form |
|--------------------------------|--|---|---|
| Markov ¹ | Non-negative $X \geq 0$ | $\Pr[X \geq a] \leq \frac{\mathbb{E}[X]}{a}$ | $\forall a > 0 \quad \Pr[X \geq k\mathbb{E}[X]] \leq \frac{1}{k} \quad k > 1$ [3, Sec. 5.1][6, Thm 1.13] |
| extension | + non-negative, strictly increasing func Φ : $\Phi(X) \geq \Phi(a) \leq \frac{\mathbb{E}[\Phi(X)]}{\Phi(a)}$ | $\Pr[X \geq a] = \Pr[\Phi(X) \geq \Phi(a)] \leq \frac{\mathbb{E}[\Phi(X)]}{\Phi(a)}$ | $\forall a > 0$ Wiki |
| Reverse Markov | upper-bounded by U : $\max X = U$ (can be positive) | $\Pr[X \leq a] \leq \frac{U - \mathbb{E}[X]}{U - a}$ | $\forall a > 0$ [1, Sec. 3.1] |
| Chebyshev ² | Finite mean and variance: $\mathbb{E}[X], \text{Var}[X]$ finite | $\Pr[X - \mathbb{E}[X] \geq a] \leq \frac{\sigma^2}{a^2}$ | $\Pr[X - \mathbb{E}[X] \geq a \cdot n] \leq \frac{1}{a^2}$ [1, Sec. 3.2] $\forall a > 0, \sigma^2 = \text{Var}[X]$ [3, Sec. 5.1][2, Thm 18.11] |
| Cantelli | Improved Chebyshev (same; but one-sided) | $\Pr[X - \mathbb{E}[X] \geq a] \leq \frac{\sigma^2}{a^2}$ | $\forall a > 0, \sigma^2 = \text{Var}[X]$ Wiki |
| Chernoff ³ | Generic | $\Pr[X \geq a] = \Pr[e^{tX} \geq e^{ta}]$ | $\forall t > 0, a \in \mathbb{R}$ [1, Sec. 3.3] |
| Jensen | $f: \mathbb{R} \rightarrow \mathbb{R}; f$ convex | $f(\mathbb{E}[X]) \leq \mathbb{E}[f(X)]$ | [3, Prob. 5.3][6, Thm 1.14] |
| Hoeffding's lemma | $\mathbb{E}[X] = \mu$ $a \leq X \leq b$ | $\Pr[e^{\lambda X}] \leq e^{\lambda \mu} e^{\frac{\lambda^2(b-a)^2}{8}}$ | $\lambda \in \mathbb{R}$ [1, Sec. 3.4] |
| Sum of random variables | | | |
| Chernoff-Hoeffding (one-sided) | n independent random vars X_1, \dots, X_n indep $S_n = X_1 + \dots + X_n$ $X_i \in [a_i, b_i] \quad \forall i$ | $\Pr[S_n - \mathbb{E}[S_n] \geq t] \leq \exp\left(\frac{-2t^2n^2}{\sum_{i=1}^n (b_i - a_i)^2}\right)$ | [1, Sec. 3.5] |
| (two-sided) ⁴ | (same as above) | $\Pr[S_n - \mathbb{E}[S_n] > t] \leq 2 \exp\left(\frac{-2t^2n^2}{\sum_{i=1}^n (b_i - a_i)^2}\right)$ | $\forall t \in (0, \frac{1}{2})$ [5, Thm.1.1] |
| (two-sided iid) | same plus iid, range, mean μ for each | $\Pr\left[\left \frac{S_n}{n} - \mu\right \geq \epsilon\right] \leq 2 \exp\left(-2n\epsilon^2\right)$ | $\forall \epsilon > 0$ [6, Thm 1.16] |
| Thm 1.3 | n independent random vars X_1, \dots, X_n indep $S_n = X_1 + \dots + X_n$ | $\Pr[S_n - \mathbb{E}[S_n] > \epsilon] \leq 2 \exp\left(\frac{-\epsilon^2}{\frac{4}{n} \sum_{i=1}^n \text{Var}[X_i]}\right)$ | $\epsilon \in (0, 2\sqrt{\text{Var}[S_n]} / (\max_i X_i - \mathbb{E}[X_i]))$ [5, Thm. 1.3] |
| Azuma | Weak law of large numbers n independent iid random vars X_1, \dots, X_n indep $\mathbb{E}[X_i] = \mu$ iid (same) | $\lim_{n \rightarrow \infty} \Pr\left[\left \frac{1}{n} S_n - \mu\right \geq \epsilon\right] = 0$ | $\forall \epsilon > 0$ [3, Sec. 5.2][6, Thm 1.15] |
| Advanced | | | |
| Bennett | n independent zero mean X_1, \dots, X_n indep $\mathbb{E}[X_i] = 0$ iid | $\Pr[S_n > \epsilon] \leq \exp\left(-n\sigma^2 h\left(\frac{\epsilon}{n\sigma^2}\right)\right)$ | $\sigma^2 := \frac{1}{n} \sum_{i=1}^n \text{Var}[X_i], \forall \epsilon > 0,$ $h(a) := (1+a) \log(1+a) - a \text{ for } a \geq 0$ [1, 4.1] |
| Bernstein | (same) | $\Pr[S_n > \epsilon] \leq \exp\left(\frac{-n\epsilon^2}{2(\sigma^2 + \epsilon/3)}\right)$ | (same) [1, 4.2] |
| Efron-Stein | scalar func of vars $f: \mathbb{R}^n \rightarrow \mathbb{R}$ X_1, \dots, X_n indep $w/ values in set \chi$ | $\text{Var}[Z] \leq \sum_{i=1}^n \mathbb{E}\left[(Z - \mathbb{E}_i[Z])^2\right]$ | $Z := g(X_1, \dots, X_n)$ $\mathbb{E}_i[Z] := \mathbb{E}[Z X_1, \dots, X_{i-1}, X_{i+1}, \dots, X_n]$ [1, 4.3] |
| McDiarmid's | scalar func of vars $f: \mathbb{R}^n \rightarrow \mathbb{R}$ X_1, \dots, X_n indep $w/ values in set \chi$ | $\Pr[f(X_1, \dots, X_n) - \mathbb{E}[f(X_1, \dots, X_n)] \geq \epsilon] \leq \exp\left(\frac{-2\epsilon^2}{\sum_{i=1}^n \epsilon_i^2}\right)$ | condition: c -bounded difference property $\forall i$ $0 < c_i < 1$ $ f(X_1, \dots, X_i, \dots, X_n) - f(X_1, \dots, X'_i, \dots, X_n) \leq c_i$ [1, 4.4] |

¹Markov's inequality bounds the first moment of random variable. Use it when a constant probability bound is sufficient [1, Sec. 3.3].

²Chebyshev is derived from Markov. It bounds the second moment. If σ is unknown, we can use the bounds of $X \in [a, b]$.

³Chernoff bound is used to bound the tails of the distribution for a sum of independent random variables. By far the most useful tool in randomized algorithms [1, Sec. 3.3].

⁴This probability can be interpreted as the level of significance α (probability of making an error) for a confidence interval around the mean of size 2δ . Therefore, we require at least $\log(2\alpha)/2\delta^2$ samples to acquire $1 - \alpha$ confidence interval $\mathbb{E}[X] \pm t$.

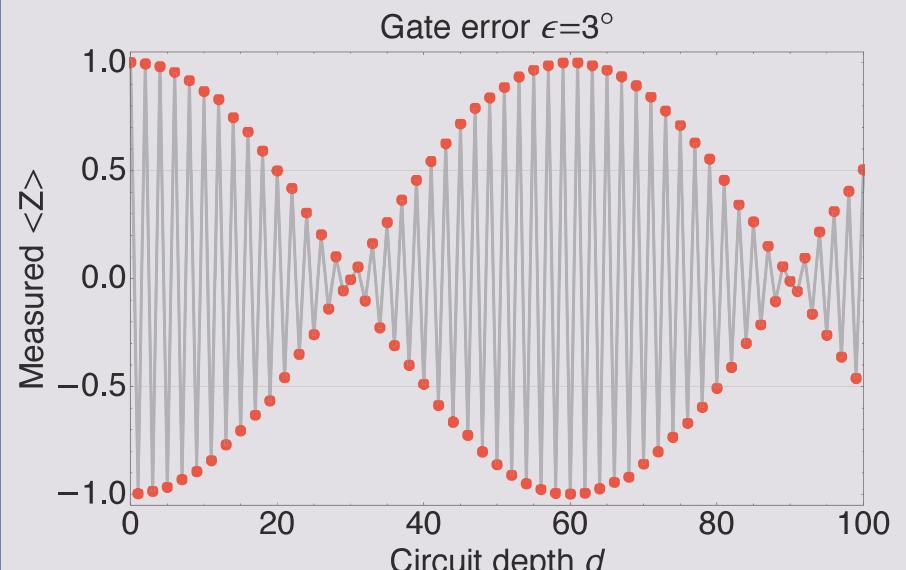
1 of 2 1. Probability (Technical note 11.9 v0.6)

Zlatko Minev, IBM Quantum (58)

<https://www.zlatko-minev.com/blog/inequalities>

Recall gate error result

$$\langle \tilde{\psi}_f | Z | \tilde{\psi}_f \rangle = \cos(d\pi + d\epsilon)$$



Zlatko Minev, IBM Quantum (59)

Projection & sampling noise

$$\langle \tilde{\psi}_f | Z | \tilde{\psi}_f \rangle = \cos(d\pi + d\epsilon)$$

Zlatko Minev, IBM Quantum (60)



Questions

Answer these multiple-choice questions
in the chat; for example, type “1a 2b.”

1. Projection noise is due to

- a) measurement apparatus that could be made more efficient
- b) classical limitations
- c) core nature of quantum physics

2. To reduce projection noise

- a) increase the number of sample
- b) you cannot undo it
- c) apply readout error mitigation

Zlatko Minev, IBM Quantum (61)



Dive deeper? Try the following



1. Calculate the following for a qubit

1. The expectation value of the sample variance for N shots of the observable $|1\rangle\langle 1|$.

where the sample mean is defined as

$$S = \frac{1}{N} \sum_{n=1}^N M_n$$

and the sample variance is defined as

2. Is the estimate biased?

3. The variance of V.

4. Can you find an expression for an unbiased estimate of the sample variance?

2. What about two qubits?

1. Can you find the projection operators for the observable ZZ?

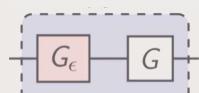
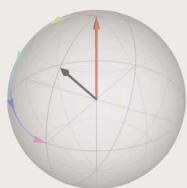
2. Find the probability distribution for the observables ZI, IZ, and ZZ for a general state.

3. If you take 10 shots and find all 10 outcomes to be 1, what is the probability the qubit is in the $|0\rangle$ state? (hint: it's not zero!)

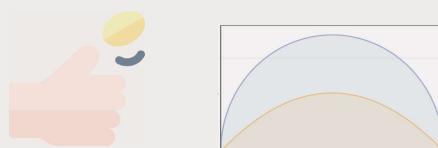
Zlatko Minev, IBM Quantum (62)

State preparation & measurement

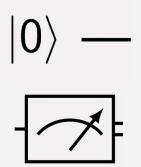
Coherent



Projection & measurement theory



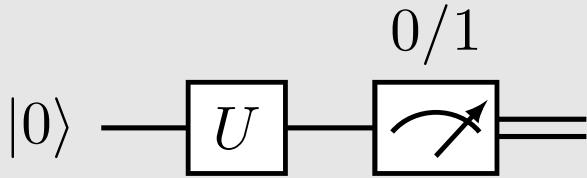
State preparation & measurement



coin toss: flaticon; spam: make it move;
road based on: freepik

Zlatko Minev, IBM Quantum (63)

Qubit example



Zlatko Minev, IBM Quantum (64)

Measurement error

Qiskit Global Summer School on Quantum Machine Learning
Zlatko K. Minev

$|0\rangle \xrightarrow{U} \begin{matrix} 0/1 \\ M \end{matrix} = \begin{cases} \text{outcome} & \text{probability} & \text{quantum} \\ \mathbb{P}(M=0) = \frac{1-p}{1-p} & = 1-p & = |\langle 0|\psi\rangle|^2 \\ \mathbb{P}(M=1) = p & = p & = |\langle 1|\psi\rangle|^2 \end{cases}$

$\hat{M} = 1|1\rangle\langle 1| = \frac{1}{2}(\hat{I} - \hat{\sigma}_z)$

M ideal classical outcome

\tilde{M} noisy classical outcome

$P_M = \begin{pmatrix} \mathbb{P}(M=0) \\ \mathbb{P}(M=1) \end{pmatrix} = \begin{pmatrix} 1-p \\ p \end{pmatrix}$

$P_{\tilde{M}} = \begin{pmatrix} \mathbb{P}(\tilde{M}=0) \\ \mathbb{P}(\tilde{M}=1) \end{pmatrix} = \begin{pmatrix} 1-\tilde{p} \\ \tilde{p} \end{pmatrix}$

$\begin{cases} \mathbb{P}(\tilde{M}=0) = \mathbb{P}(\tilde{M}=0|M=0)\mathbb{P}(M=0) + \mathbb{P}(\tilde{M}=0|M=1)\mathbb{P}(M=1) \\ \mathbb{P}(\tilde{M}=1) = \mathbb{P}(\tilde{M}=1|M=0)\mathbb{P}(M=0) + \mathbb{P}(\tilde{M}=1|M=1)\mathbb{P}(M=1) = \varepsilon(1-p) + (1-\varepsilon)p = \tilde{p} \end{cases}$

$\tilde{P}_M = A P_M$

$A = \begin{matrix} M=0 & M=1 \\ \tilde{M}=0 & \begin{pmatrix} \mathbb{P}(\tilde{M}=0|M=0) & \mathbb{P}(\tilde{M}=0|M=1) \\ \mathbb{P}(\tilde{M}=1|M=0) & \mathbb{P}(\tilde{M}=1|M=1) \end{pmatrix} \end{matrix}$

$= \begin{pmatrix} 1-\varepsilon & \varepsilon \\ \varepsilon & 1-\varepsilon \end{pmatrix} = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \text{ for an ideal measurement!}$

$\sum_n A_{nn} = 1 \text{ for any } n$

$\tilde{p} = \varepsilon(1-p) + (1-\varepsilon)p$

$= \varepsilon - p\varepsilon + p - \varepsilon p$

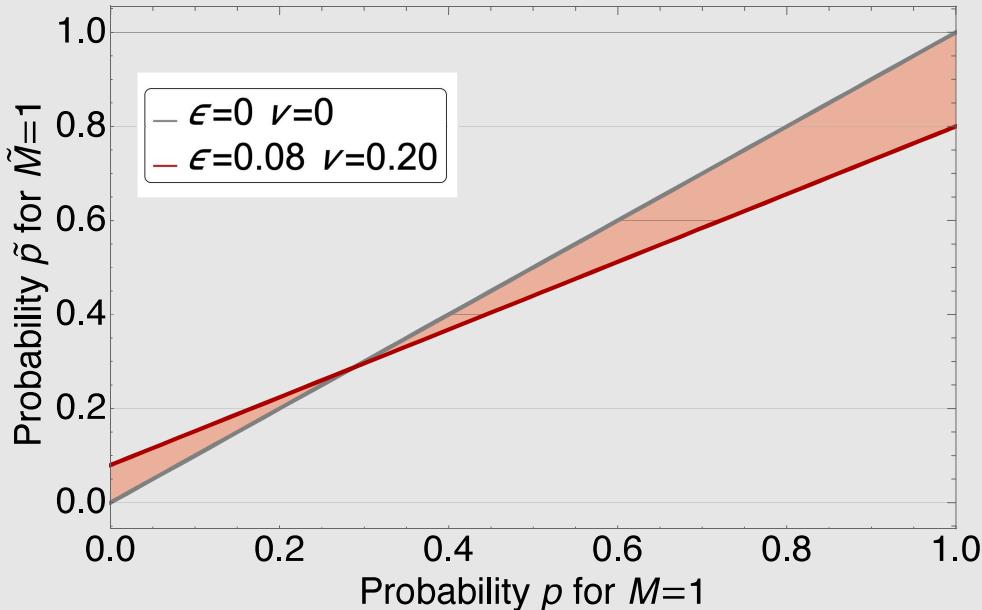
$= p + \varepsilon - (\varepsilon + p)$

Stochastic matrix

Zlatko Minev, IBM Quantum (65)

Readout error

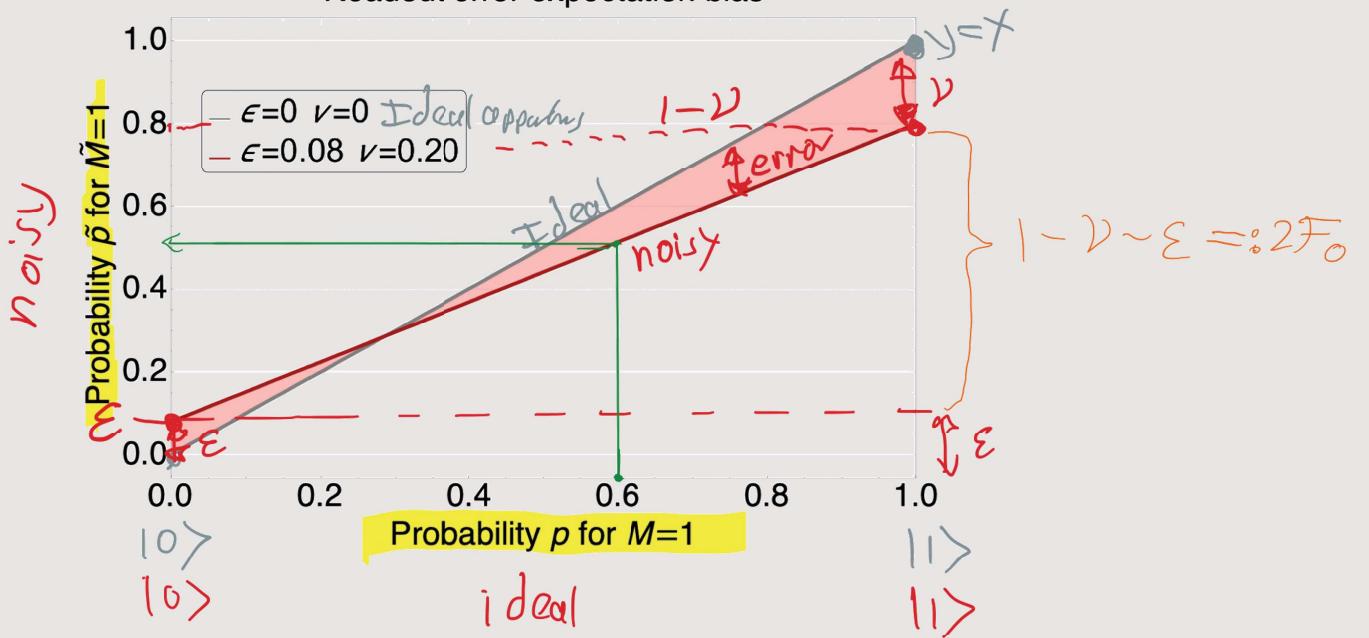
Readout error expectation bias



Zlatko Minev, IBM Quantum (66)

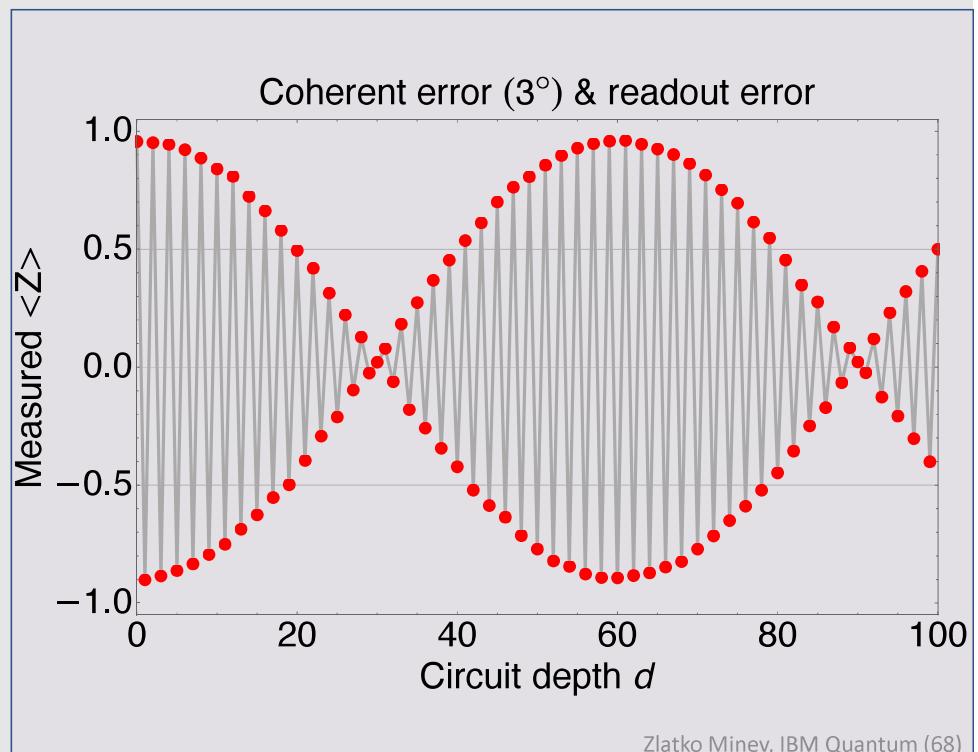
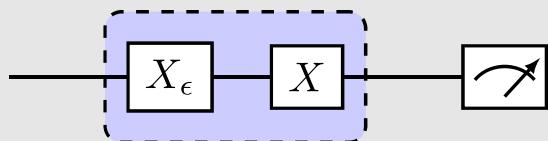
Assignment fidelity

Readout error expectation bias



Zlatko Minev, IBM Quantum (67)

Projection & sampling noise



Questions

Answer these multiple-choice questions
in the chat; for example, type “1a 2b.”

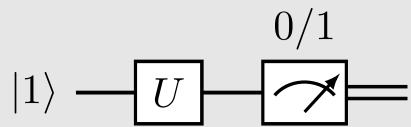
1. Readout error is due to

- a) measurement apparatus that could be made more efficient
- b) classical limitations
- c) core nature of quantum physics

2. To reduce readout error bias

- a) increase the number of sample
- b) you cannot undo it
- c) apply readout error mitigation

State prep



Zlatko Minev, IBM Quantum (70)

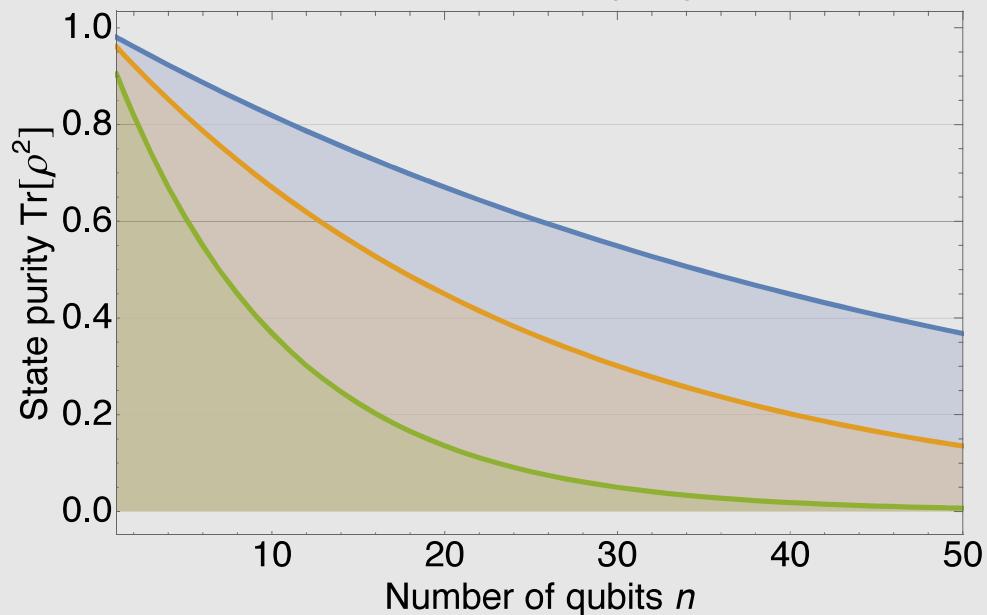
Multiple qubits

$$[(1 - p) |0\rangle\langle 0| + p |1\rangle\langle 1|]^{\otimes n} \equiv \text{---} U \text{---} = \text{---} \curvearrowright \text{---}$$

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Thermal state

Thermal state purity



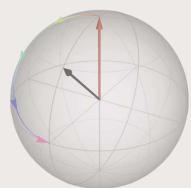
Individual qubit population

- $p=0.01$
- $p=0.02$
- $p=0.05$

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Incoherent noise

Coherent



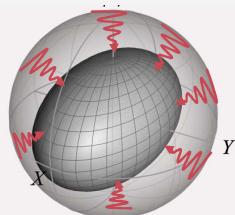
Projection & measurement theory



State preparation & measurement



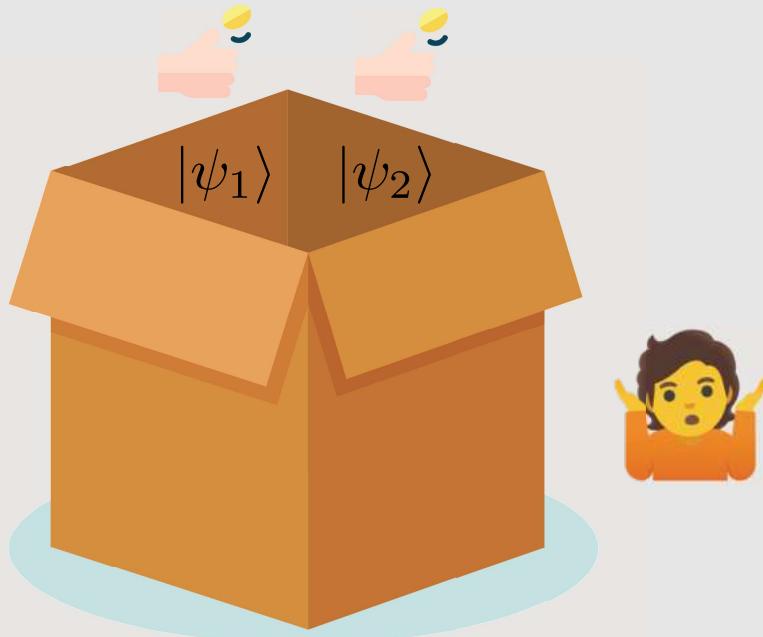
Incoherent noise



coin toss: flaticon; spam: make it move;
road based on: freepik

Zlatko Minev, IBM Quantum (73)

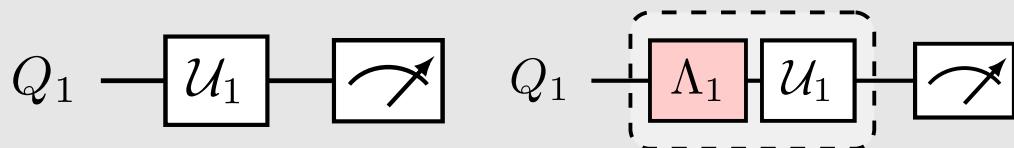
Review: mixed state (density matrix)



box; freepik. person: facebook

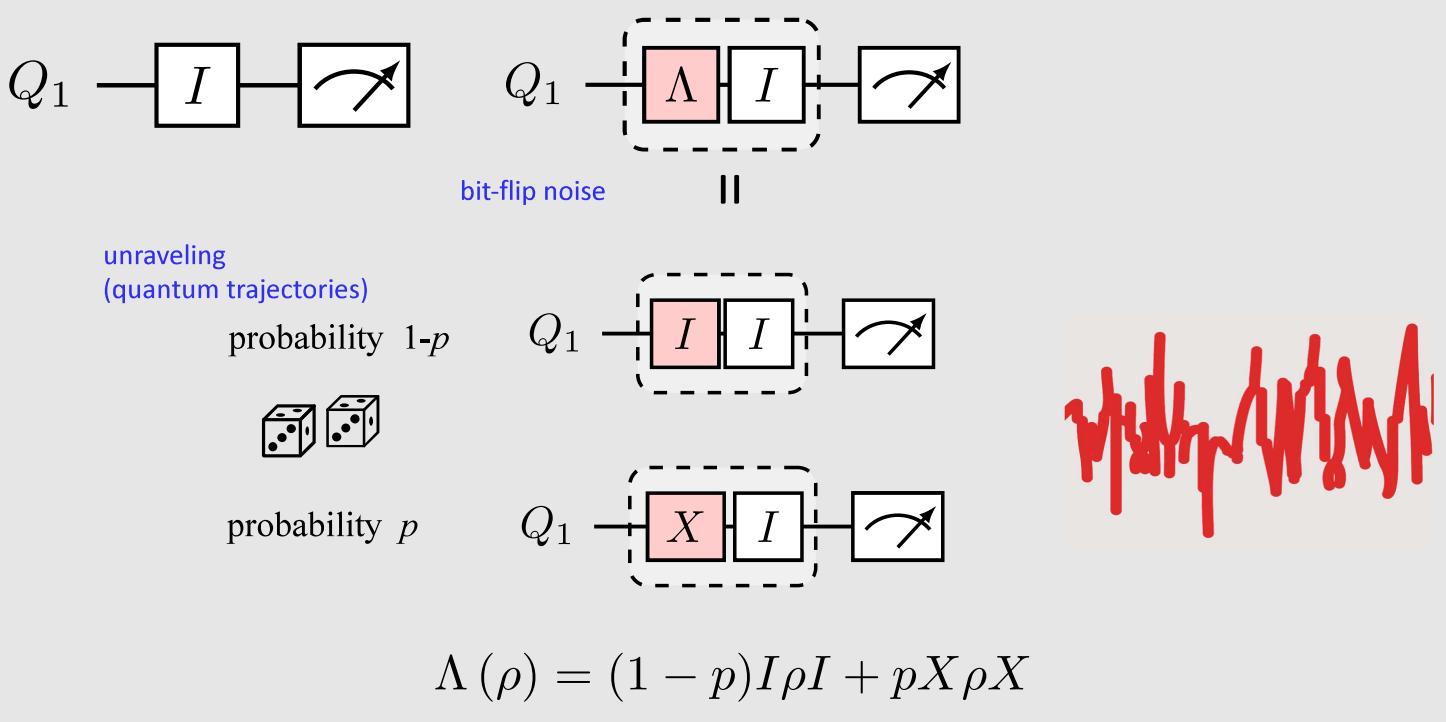
Zlatko Minev, IBM Quantum (74)

Toy model



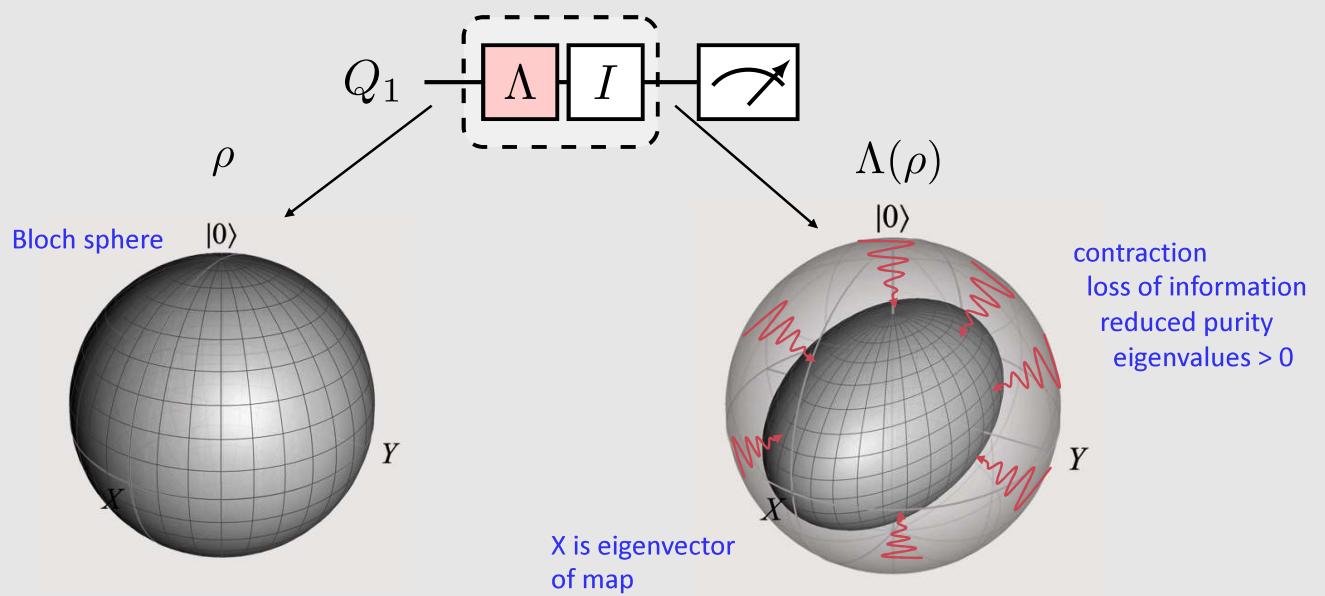
Zlatko Minev, IBM Quantum (75)

Toy model: noise unraveling into quantum trajectories



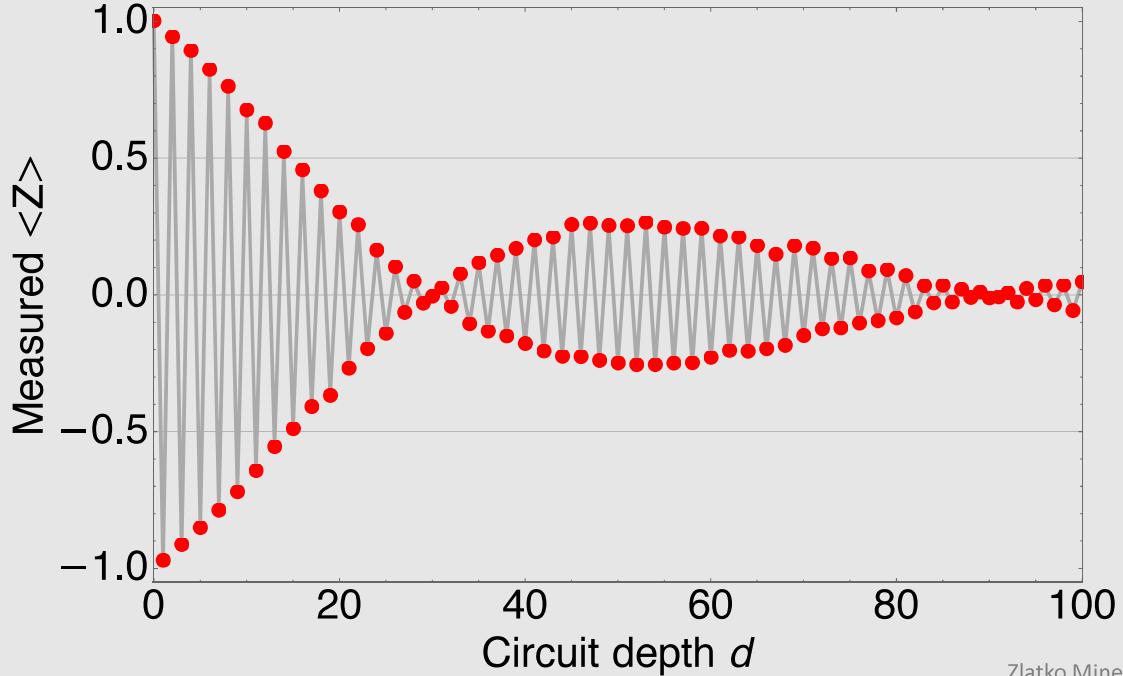
Zlatko Minev, IBM Quantum (76)

Toy model: noise unraveling into quantum trajectories

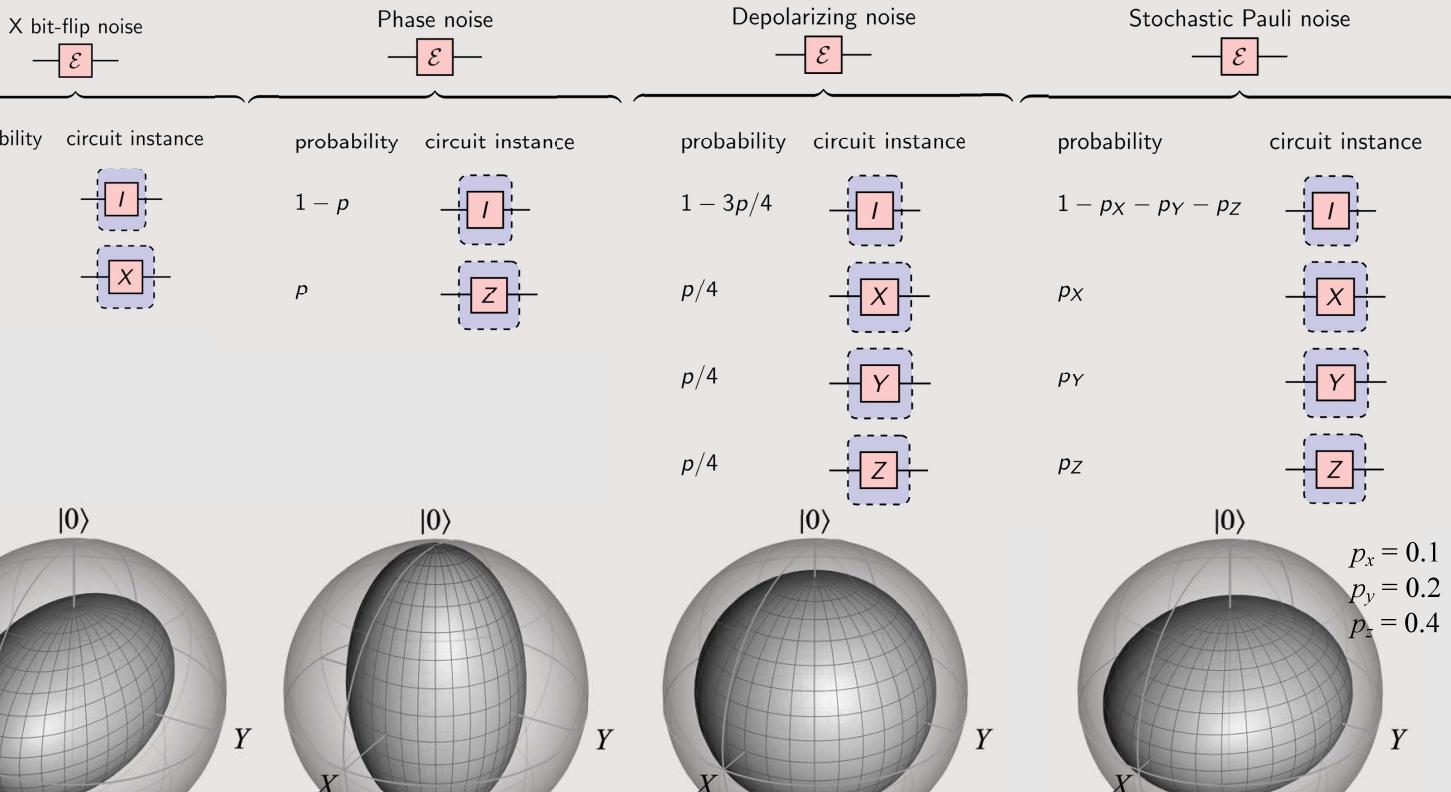


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Coherent error (3°) & incoherent error $p=0.012$

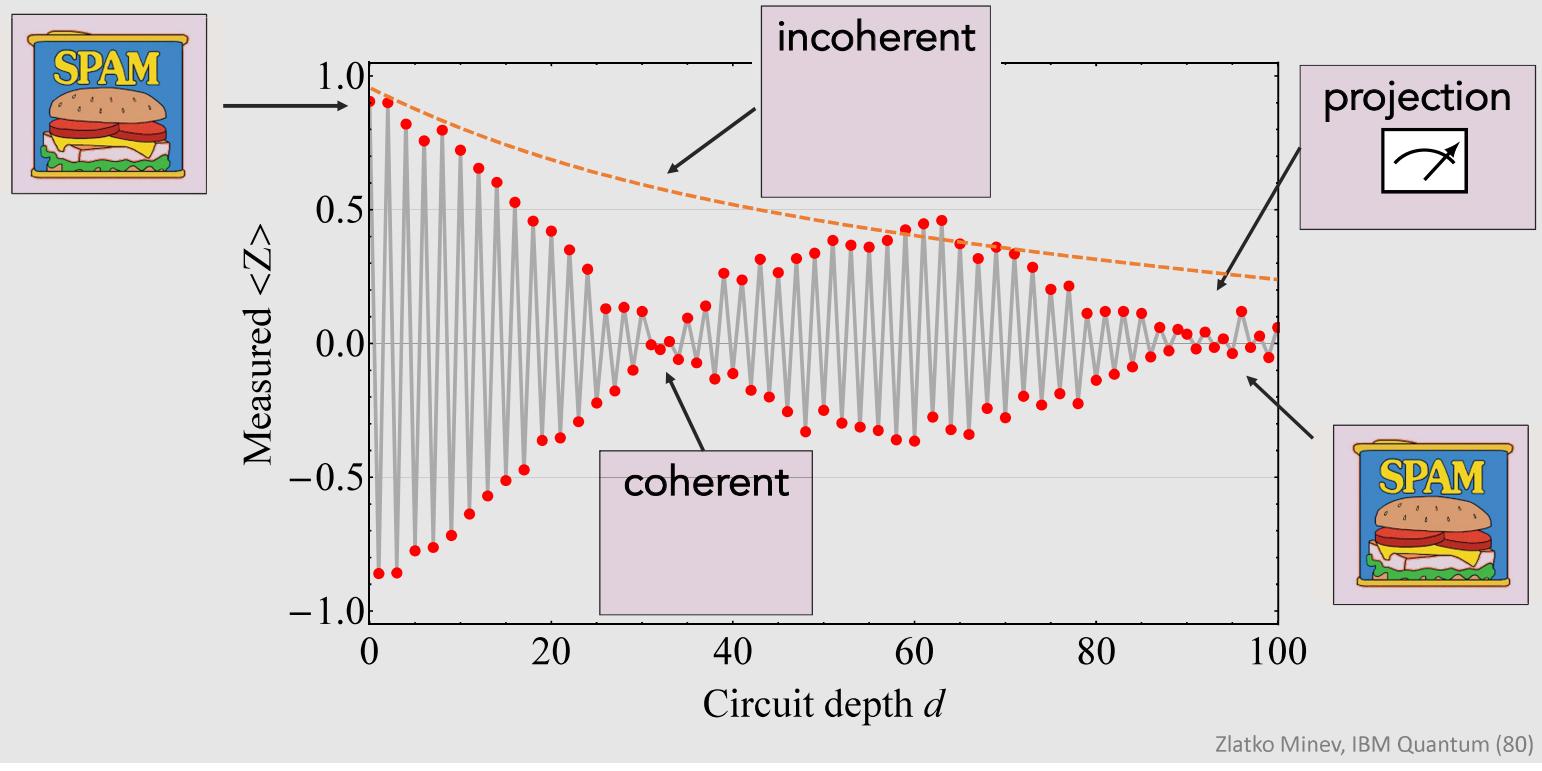


Zlatko Minev, IBM Quantum (78)

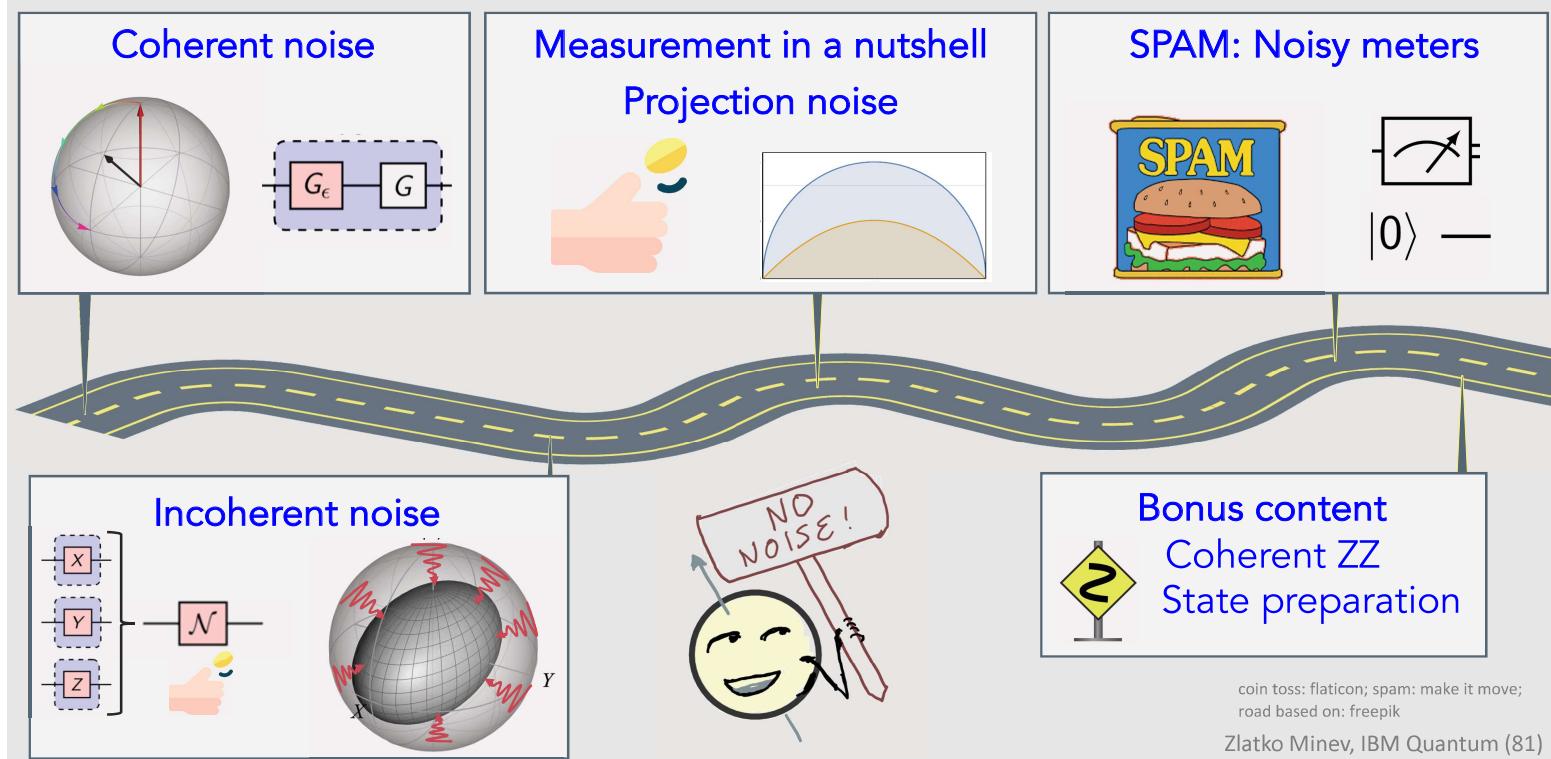


Zlatko Minev, IBM Quantum (79)

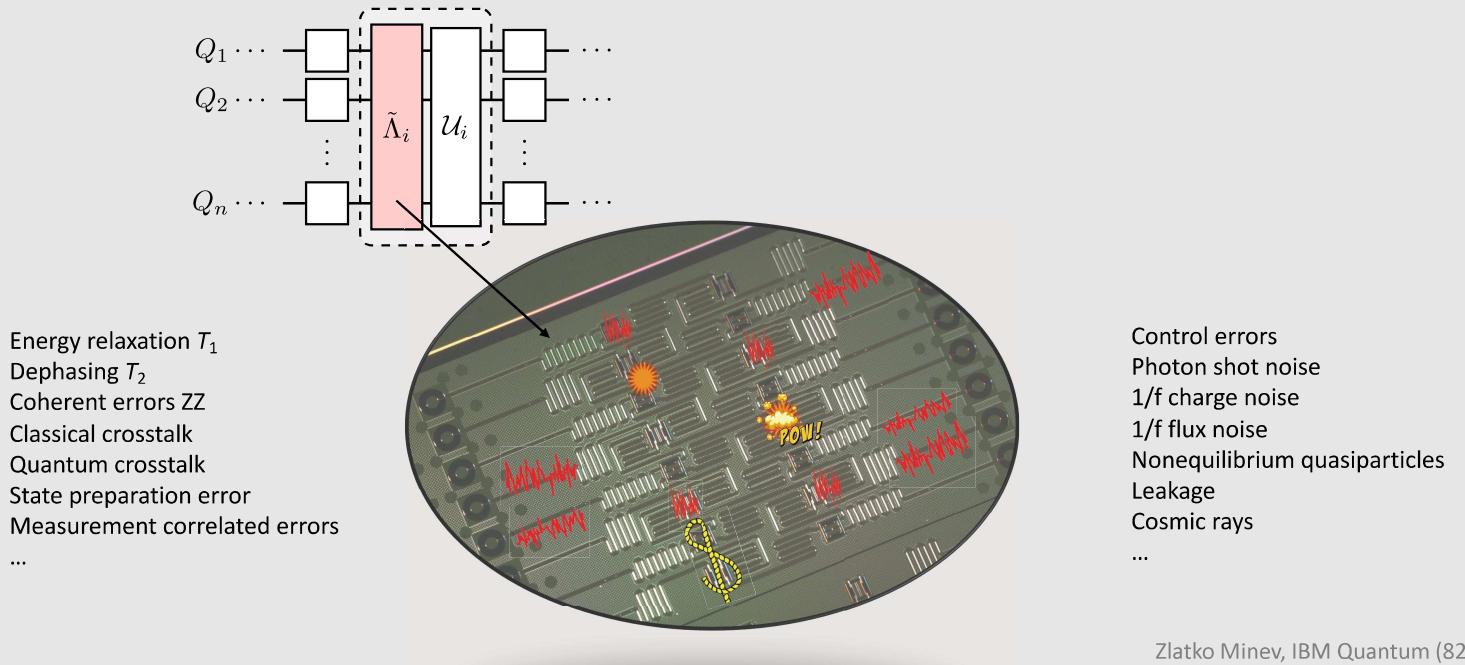
Elements of noise



Our journey so far



Outlook: Is it possible to learn & correct the noise with accuracy, efficiency, and scalability?



Error mitigation and error correction

Error mitigation: working with what you have

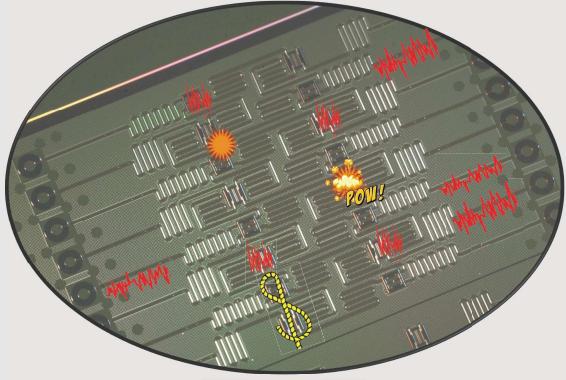
- **benefit** suppress errors on classical results (expectation values)
- **q-cost** no extra qubits or hardware resources needed
- **c-cost** trades classical resources (post-processing) for lower error
- **limitation** bad asymptotic scaling: high number of samples & circuits

Error correction: protecting quantum information

- **benefit** suppress & correct errors to arbitrarily small level
- **q-cost** very large qubit and hardware overhead
- **c-cost** decoding and encoding can be classically costly
- **challenge** requires fault-tolerant operations and readout

Latest: Qiskit Quantum Seminar YouTube Series

Ways to learn more
TODO



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Introduction to Quantum Noise

Lab work with Qiskit

Run experiments on real devices

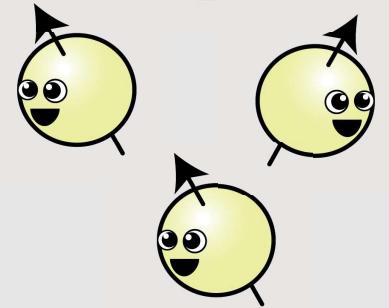
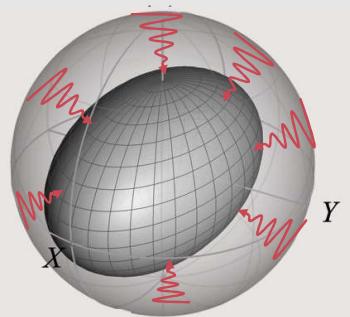


Check out references, problems given in
the lecture, dangerous bends

Stay in touch

Thank you!

Zlatko K. Minev



@zlatko_minev



zlatko-minev.com

IBM Quantum

The important thing is not to stop questioning.
Curiosity has its own reason for existence.

One cannot help but be in awe when he
contemplates the mysteries of eternity, of life, of the
marvelous structure of reality.

It is enough if one tries merely to comprehend a
little of this mystery each day.

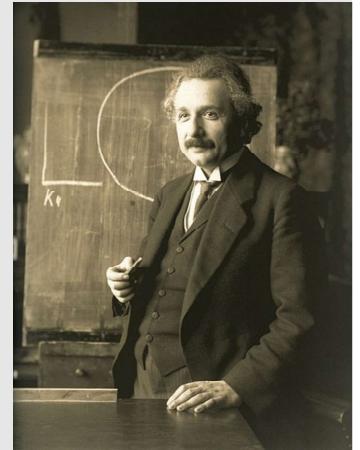


Photo: F. Schmutzler

Albert Einstein



@zlatko_minev

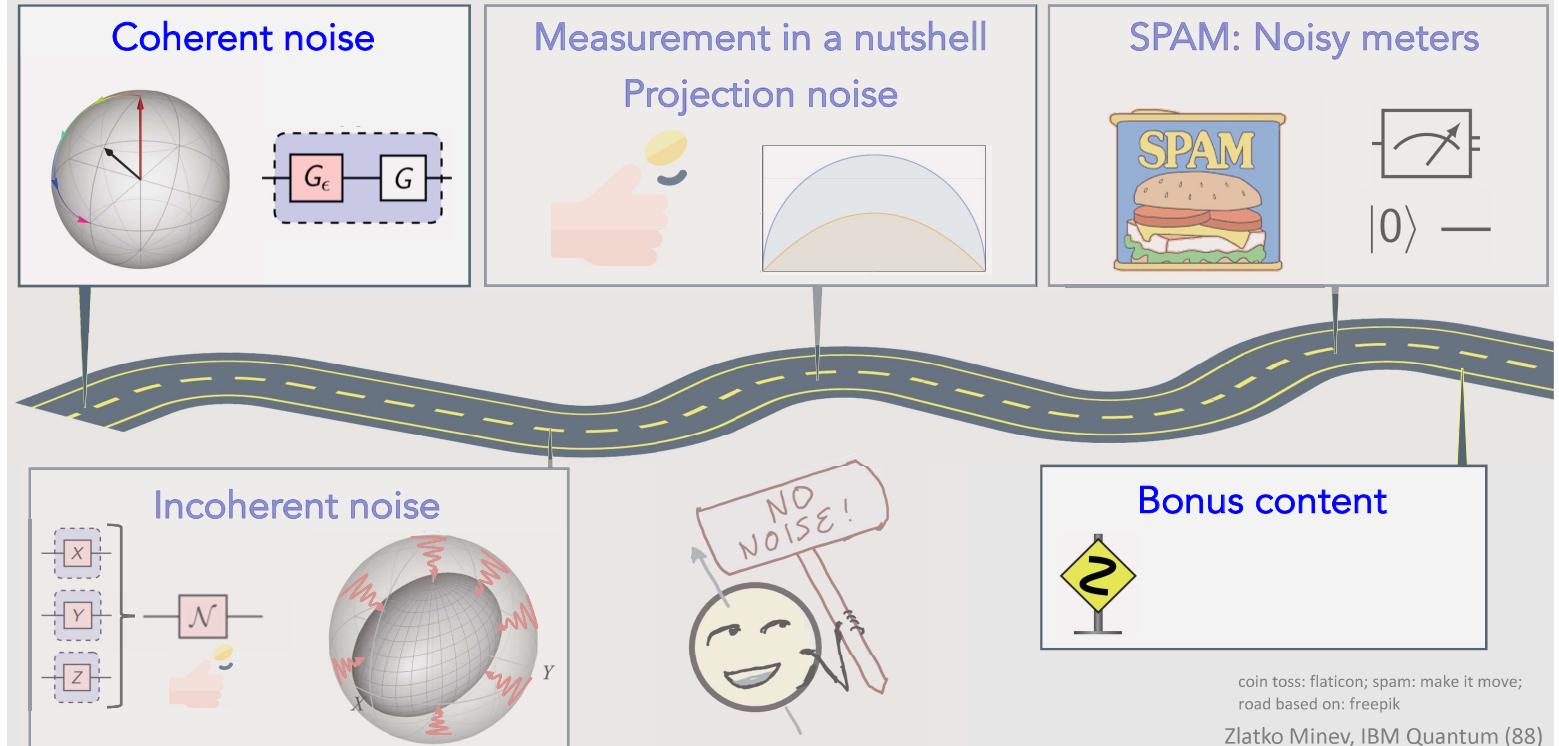


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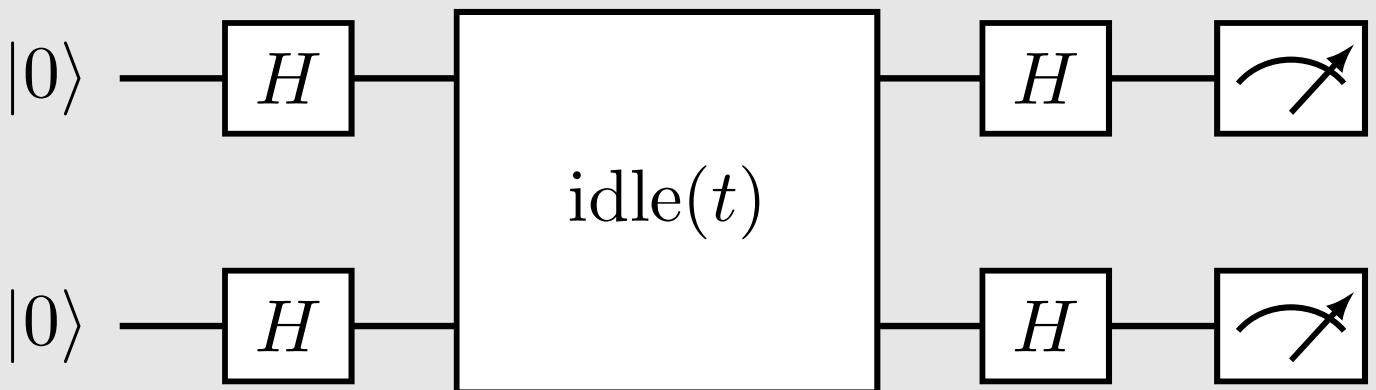
IBM Quantum

Bonus content

Bonus content

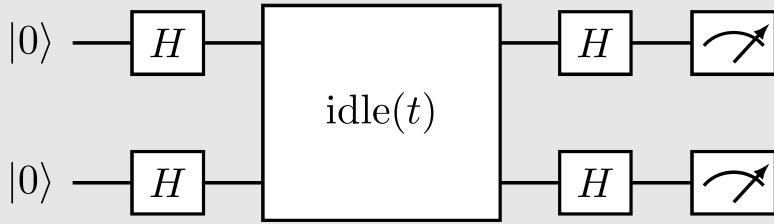


Bonus content: two-qubit coherent ZZ error





Bonus content: two-qubit ZZ error



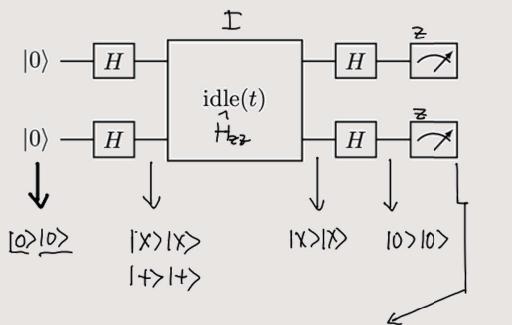
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Introduction to quantum noise

Coherent errors

Qiskit Global Summer School on Quantum Machine Learning

Zlatko K. Minev



$$\begin{aligned}
 &\langle Z|Z\rangle = + \\
 &\langle Z|Z\rangle = +1 \\
 &\langle Z|Z\rangle = \langle 0|Z|0\rangle \underbrace{\langle Z|0\rangle}_{=0} \langle 0|0\rangle \\
 &= \langle 0|Z|0\rangle \langle 0|Z|0\rangle \\
 &\approx (+) (+) \\
 &\approx +1
 \end{aligned}$$

Hadamard gate

$$H = \begin{pmatrix} |0\rangle & |1\rangle \\ \frac{1}{\sqrt{2}} & \frac{1}{\sqrt{2}} \\ \frac{1}{\sqrt{2}} & \frac{-1}{\sqrt{2}} \end{pmatrix}$$

$$\begin{cases} H|0\rangle = |+x\rangle = \frac{1}{\sqrt{2}}(1) \\ H|1\rangle = |-x\rangle = \frac{1}{\sqrt{2}}(1) \end{cases}$$

$$|+x\rangle = +|+x\rangle$$

$$|-x\rangle = -|-x\rangle$$

$$|+x\rangle := |+\rangle$$

$$|-x\rangle := |-\rangle$$

$$\begin{cases} Z|0\rangle = +|0\rangle \\ Z|1\rangle = -|1\rangle \end{cases}$$

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NOTES

ZZ Interaction

$$\hat{H} = \frac{1}{2} \hbar \omega \hat{Z}\hat{Z}$$

$$\hat{U}(t) = \exp(-i\frac{\hbar\omega}{2} \hat{Z}\hat{Z}t)$$

$$= \exp(-i\frac{\omega t}{2})$$

$$= \cos(\frac{\omega t}{2}) \hat{I} - i\sin(\frac{\omega t}{2}) \hat{Z}\hat{Z}$$

$$= \hat{R}_{\frac{\omega t}{2}} (\theta = \omega t)$$

$$|0\rangle|0\rangle \xrightarrow{HH} |+\rangle|+\rangle \xrightarrow{\text{idle}} R_{ZZ(\theta)}|+\rangle|+\rangle = \cos\frac{\theta}{2}|+\rangle|+\rangle - i\sin\frac{\theta}{2}|-\rangle|-\rangle \xrightarrow{HH} \cos\frac{\theta}{2}|0\rangle|0\rangle - i\sin\frac{\theta}{2}|1\rangle|1\rangle$$

$$R_{ZZ(\theta)}|+\rangle|+\rangle = \cos\frac{\theta}{2}|+\rangle|+\rangle - i\sin\frac{\theta}{2}(|-\rangle|-\rangle)$$

$$Z|+\rangle = \frac{1}{\sqrt{2}}(|0\rangle + |1\rangle) = \frac{1}{\sqrt{2}}(|+\rangle - |-\rangle)$$

$$Z|+\rangle = |-\rangle$$

$$Z|-\rangle = |+\rangle$$

$$\langle ZI \rangle = \cos \omega t$$

$$\langle IZ \rangle = \cos \omega t$$

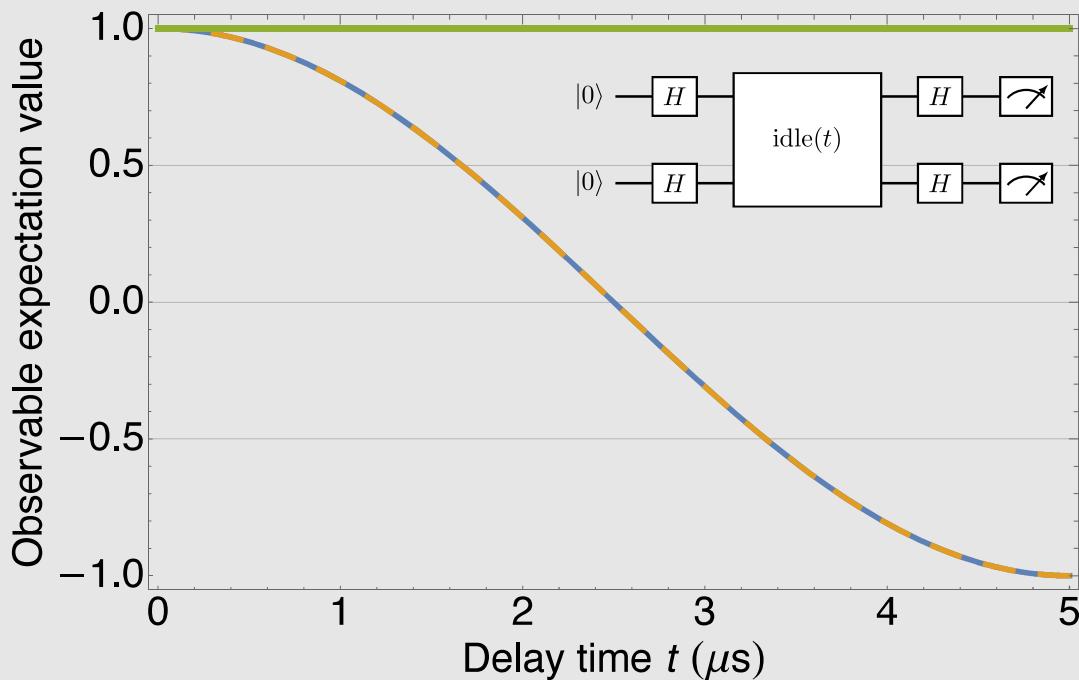
$$\langle ZZ \rangle = 1$$

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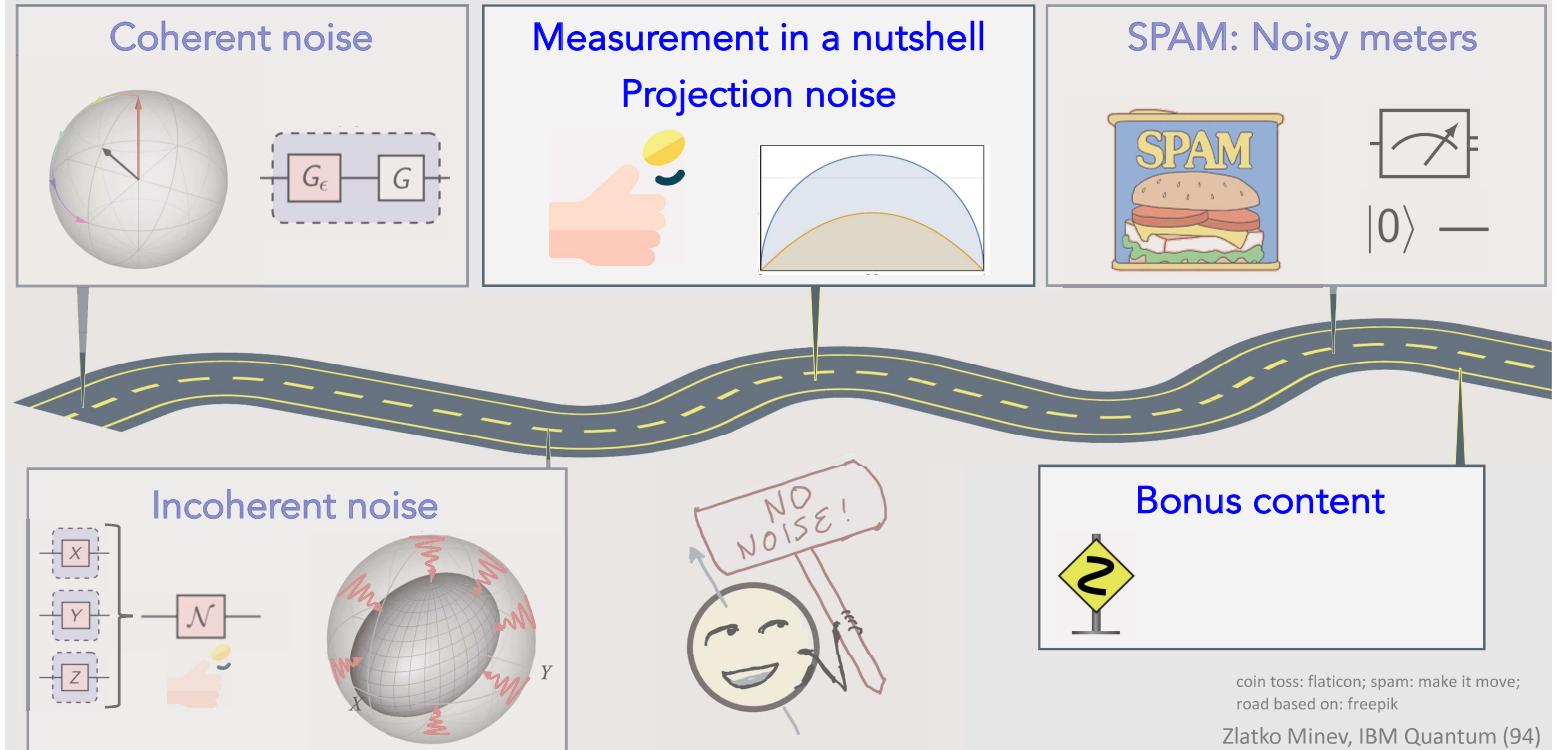
Bonus content: two-qubit ZZ error

Gate error $\omega = 10$ kHz



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Bonus content



A matrix

Bonus section content:

Reconstruct A matrix

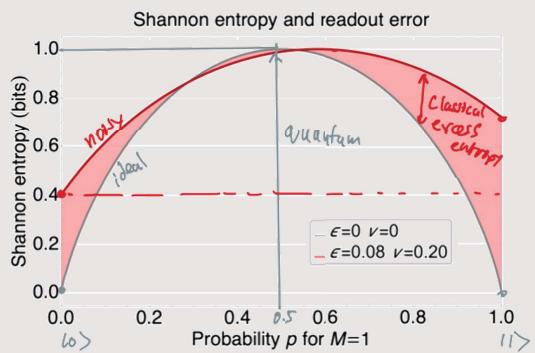
$$|0\rangle \xrightarrow{\text{A}} (M=0) \xrightarrow{\text{A}} \tilde{M} \quad p=0 \quad \tilde{p}=\epsilon$$

$$|0\rangle \xrightarrow{\text{A}} (M=1) \xrightarrow{\text{A}} \tilde{M} \leftarrow w_{\text{data\&tgt}}^{\text{best}} \quad p=1 \quad \tilde{p}=1-\epsilon$$

Noise mitigation

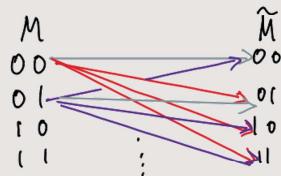
we know A
measure \tilde{p} , \tilde{P}_M noisy
find p , P_M ideal
 $\dim A = 2^n \times 2^n \quad n = \text{number}$

$$\begin{aligned} \tilde{P}_M &= A P_M \\ P_M &= A^{-1} \tilde{P}_M \\ p &= \end{aligned}$$



Laser System

$\dim A = 2^n \times 2^n$



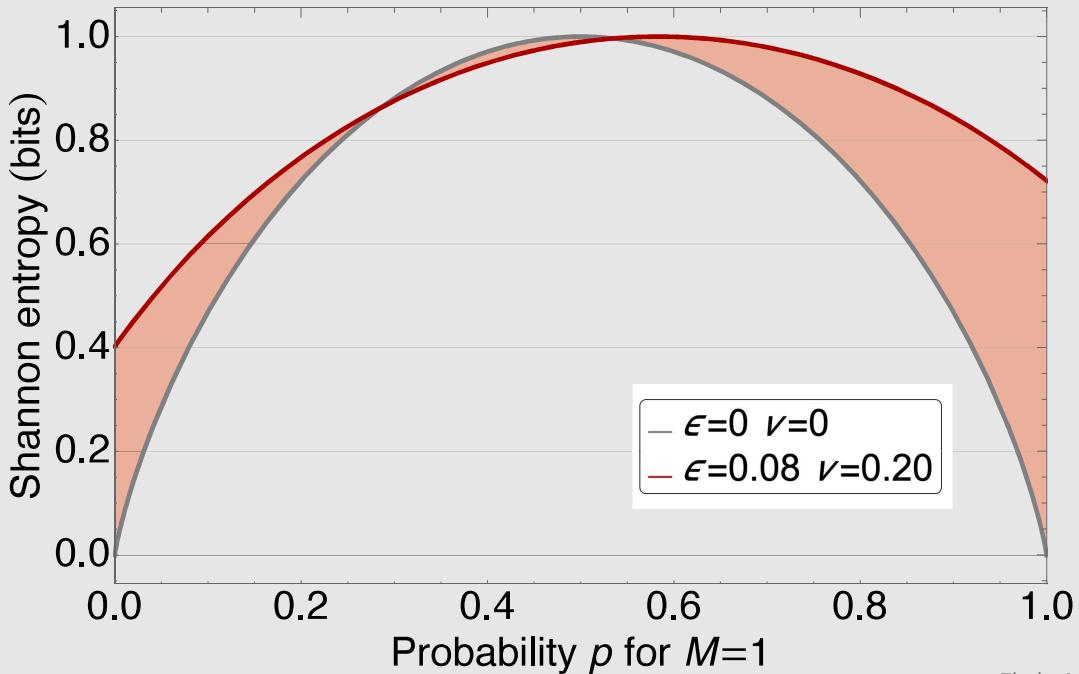
Shannon Entropy

$$H(A) = -\sum_m p_m \log_2 p_m = - (1-p) \log_2 (1-p) - p \log_2 p$$

Binary entropy

Entropy

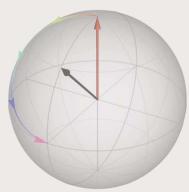
Shannon entropy and readout error



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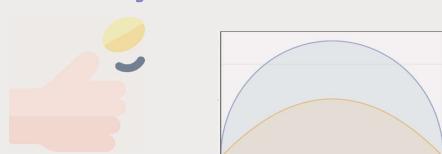
Bonus content

Coherent noise

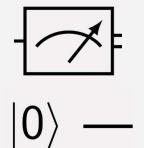


Measurement in a nutshell

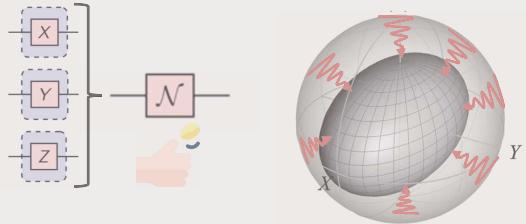
Projection noise



SPAM: Noisy meters



Incoherent noise



Bonus content



coin toss: flaticon; spam: make it move;
road based on: freepik

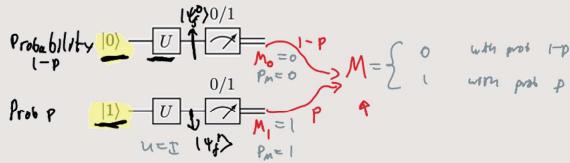
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State prep noise

Introduction to quantum noise

State preparation noise

Qiskit Global Summer School on Quantum Machine Learning
Zlatko K. Minev



Density operator

$$\rho = (1-p) |0\rangle\langle 0| + p |1\rangle\langle 1|$$

$$= \text{Pno} \rho_{\text{ideal}} + \text{Perr} \rho_{\text{error}}$$

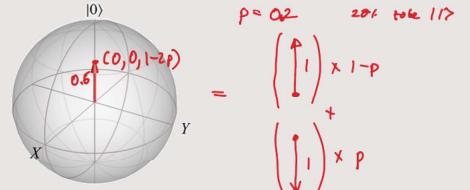
$$\begin{aligned} \rho_f &= \text{Pno} \rho_{\text{ideal}} + \text{Perr} \rho_{\text{error}} \\ &= (1-p) |1\rangle\langle 1| + p |1\rangle\langle 1| \\ &= (1-p) U |0\rangle\langle 0| U^\dagger + p U |1\rangle\langle 1| U^\dagger \\ &= U \left[(1-p) |0\rangle\langle 0| + p |1\rangle\langle 1| \right] U^\dagger \\ &= U \rho U^\dagger \end{aligned}$$

$$\rho_v = U_p U_v^\dagger$$

$$\begin{aligned} (U|\psi\rangle)^+ &= (|\psi_f\rangle)^+ \\ &\quad \langle \psi_f | = \langle \psi | U^\dagger \end{aligned}$$

$$\begin{aligned} \rho_v &= U_p U_v^\dagger \\ |0>\langle 0| &= \begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix} \\ |1>\langle 1| &= \begin{pmatrix} 0 & 0 \\ 0 & 1 \end{pmatrix} \\ \rho &= \begin{pmatrix} 1-p & 0 \\ 0 & p \end{pmatrix} = \frac{1}{2} (\hat{I} + (1-2p)\hat{Z}) \\ \langle X \rangle &= \text{Tr}[\Sigma X \rho] = 0 \\ \langle Y \rangle &= \text{Tr}[\Sigma Y \rho] = 0 \\ \langle Z \rangle &= \text{Tr}[\Sigma Z \rho] = 1-2p \end{aligned}$$

$$\text{Tr}[\rho_{AB}] = 2S_{AB} \quad \text{for } \alpha, \beta \in \{X, Y, Z\}$$



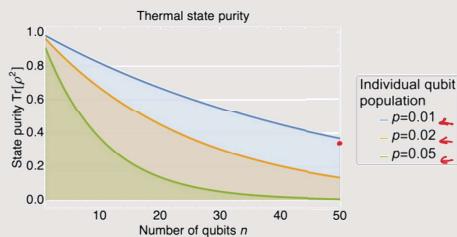
$$\begin{aligned} \text{Tr}(\rho^2) &= \langle X \rangle^2 + \langle Y \rangle^2 + \langle Z \rangle^2 = 0^2 + 0^2 + (1-2p)^2 \\ &= 1 - 4p + 4p^2 \end{aligned}$$

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State prep noise

Scaling to larger number of qubits

$$\begin{aligned} [(1-p)|0\rangle\langle 0| + p|1\rangle\langle 1|]^{\otimes n} &\equiv U \otimes \dots \otimes U \\ \rho_0 &= \underbrace{[(1-p)|0\rangle\langle 0| + p|1\rangle\langle 1|]}_{(1-p)^n |000\dots 0\rangle\langle 000\dots 0|}^{\otimes n} \\ \rho_0 &= \begin{cases} 1 & \text{if } p=0 \text{ (ideal case)} \\ (1-p)^n \approx 1-np+\mathcal{O}(p^2) & \text{for } p>0 \end{cases} \\ \langle Z_k \rangle &= \langle |1| |1| |1| \dots \rangle = 1-2p \xrightarrow{\text{much easier}} \\ \langle Z Z Z \dots Z \rangle &= \text{Tr}[\Sigma Z^{\otimes n} \rho_0] = (1-2p)^n \\ \rho_0 &= \prod_{k=1}^n \frac{1}{2} (\hat{I}_k + (1-2p)\hat{Z}_k) \\ &= \frac{1}{2^n} (1-2p)^n \otimes (1-2p) \otimes \dots \\ &= \frac{1}{2^n} \sum (1-2p)^n Z^{\otimes k} + \dots \\ \text{Tr}[\rho_0^2] &= \prod_{k=1}^n \text{Tr}(\rho_0^2) = (1-2p)^{2n} \end{aligned}$$



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The End!

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