# Second Homework Report

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A full report on the second assigned Homework for the course "Fondamenti di Comunicazioni e Internet", Sapienza University of Rome.

Our working group is **Group 13**.

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# 1 Description

#### The goal

of this project was to simulate, using the network simulator ns-3, the network seen in Figure 1. Tasked with this network, albeit not complex, we naturally opted for dividing it into smaller sub-networks, planning to link them at a later stage. We decided to use the following alphabetical identifiers for each sub-net:

- **Network A**: comprised of Nodes 10-16 and an Access Point (Node 9).
- **Network B**: comprised of Nodes 6, 7, 8 and a router (Node 5).
- **Network C**: comprised of Nodes 0, 1 and a router (Node 2).
- **Network D**: comprised of Node 3 and a router (Node 4).

We will refer to the links as **L1** (the network that links Network D and Network C), **L2** (Network A and Network D), and to **L3** (Network B and Network D). Our network had a total of four applications installed on different nodes, exchanging data via:

• A TCP file delivery of 1221MB, starting at 0.72s, from Node 7 (Net B) to Server 0 (Net C).



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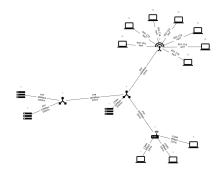


Figure 1: Our network topology.

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- A TCP file delivery of 1264MB, starting at 3.93s, from Node 16 (Net A) to Server 1 (Net C).
- A TCP file delivery of 1754MB, starting at 3.87s, from Node 13 (Net A) to Server 0 (Net C).
- A UDP Echo Application, sending packets of 1964B, with a periodicity of 20ms, and the maximum number of packets being 250. The sender was Client 15 (Network A) and the receiver Server 3 (in Network D).

# 2 Network topology and delays

#### 2.1 Known Topologies

We identified Network A as having a *star topology*, and the same goes for Network B, C. Note that a **single-layer tree topology** is closely related to a star topology, so it makes sense that Network C could be ambiguous as to whether it denotes a single layer tree topology rather than a star topology; while the distinction is more evident in Network A, since it has a larger number or spokes.

Finally, Network D, along with Networks L1-3, exhibits a simple point-to-point (or linear) topology.

#### 2.2 Packets Path

We analyzed the traffic of each TCP application, highlighting the path taken by each packet, which was coherent with what we expected before the analysis. To exclude the UDP Echo Application, we used the "tcp" filter in Wireshark.

Table 1: TCP file delivery of 1221MB

NODE ID	SENDER IP ADDRESS	RECEIVER IP ADDRESS
Node 7 (start)	192.168.2.2	192.168.3.1
Node 5 (interface 2)	//	//
Node 4 (interface 2)	//	//
Node 2 (interface 0)	//	//
Node 0 (end)	//	192.168.3.1

Table 2: TCP file delivery of 1264MB

NODE ID	SENDER IP ADDRESS	RECEIVER IP ADDRESS
Node 16 (start)	192.168.1.7	192.168.3.5
Node 5 (interface 2)	//	//
Node 4 (interface 2)	//	//
Node 2 (interface 0)	//	//
Node 1 (end)	//	192.168.3.5

Table 3: TCP file delivery of 1754MB

NODE ID	SENDER IP ADDRESS	RECEIVER IP ADDRESS
Node 13 (start)	192.168.1.8	192.168.3.1
Node 5 (interface 2)	//	//
Node 4 (interface 2)	//	//
Node 2 (interface 0)	//	//
Node 0 $(end)$	//	192.168.3.1

#### 2.3 Bottlenecks: Countermeasures and Solutions

From Fig. 1, it is apparent that the following are indeed bottlenecks:

- PTP link (Node 0 Node 2)
- PTP link (Node 1 Node 2)

- CSMA link (Node 6 Node 5)
- CSMA link (Node 7 Node 5)
- CSMA link (Node 8 Node 5)
- CSMA link (Node 3 Node 4)
- Each 802.11g link in Network A

There are multiple strategies to eliminate bottlenecks in a network, e.g. increasing the link's bandwidth, design priority packets and introduce network prioritization, optimizing routing algorithms to mitigate delays, improving network monitoring in order to act swiftly to intervene in case of malfunctioning, and to implement effective CA, CD protocols.

#### 2.4 Average Throughput

Since we had to ignore UDP Echo Application traffic, we chose to analyze the throughput on Network L1 (using, again, the "tcp" filter). To isolate the packets sent up until t, we filtered packets based on their number, hence the filter used was frame.number >=0 && frame.number <=N, where N represents the packet number at the requested time t. At t=2.0s, the calculated throughput is 0.100464Mbps, while at t=5.0s, the calculated throughput amounted to 0.269821Mbps. This is in line with the I/O Graph provided by Wireshark (Fig. 2), since the traffic at t=2.0s appears to be much less compared to the traffic at t=5.0s. Note that we ignored the ACKs, since we wanted to focus on analyzing the data transmission to get a more accurate metric.

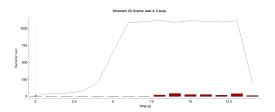


Figure 2: I/O Graph for Node 4, interface 2

# 3 Wireless Network, Infrastructure Mode

#### 3.1 AP Behavior and Beacon Frame

We expected the Access Point (AP) to broadcast its Beacon frame as soon as the simulation started, followed by the

Association Request (AReq) and its consequential ACK and Association Response (ARes), sent respectively by a Wi-Fi Station and the AP. We observed this exact behavior. Furthermore, the Beacon Frame provides important details:

- It tells each Station the AP's SSID, to which the AReq must be sent to.
- The DS flag indicates whether the AP is receiving or transmitting, respectively, it set to 01 or 10.
- The Payload is the Wireless Management, which includes important information such as whether Ad Hoc or Infrastructure Mode is used, and the available transmission rates.
- A timestamp to synchronize the Wi-Fi Stations.

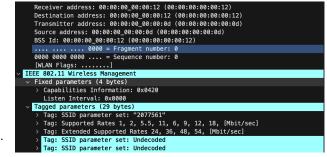


Figure 3: Screenshot of a Beacon frame Capture

#### 3.2 Collisions

When applying the wlan.fc.retry==1 filter, to highlight every re-transmission at Node 9 (interface 1), many packets appear to be colliding (Fig. 4), e.g. the probable behavior of each Station after receiving the Beacon frame is that each one sends back the AReq frame at the same time, thereby causing collisions. Also, it's worth to note that the total number of packets that match the filter is 718. The cause for this high collision rate is that the RTS/CTS mechanism was not applied in this capture.

Figure 4: Screenshot of the wlan.fc.retry==1 filter applied to the Node 9 - 1 capture, without RTS/CTS

#### 3.3 Overhead (No RTS/CTS)

To calculate the total Overhead of our Echo Application, we used two filters (see Fig. (6)), namely:

• wlan.addr == 00:00:00:00:00:0d || udp.

The 00:00:00:00:00:00 address refers to the sender Node in our application, and filters the Echo Reply. We then calculated the **overhead of each packet** and multiplied it by the number of packets that matched the filter, which resulted in a total of 18982B.

## 3.4 RTS/CTS

As expected, the use of RTS/CTS eliminates the risk of collisions, which is due to the inherent mechanism of RTS/CTS. By enabling it, a Station that wants to send data first asks for permission to the AP, which replies with a broadcast message telling the requesting Station that it can start sending its frames. By broadcasting the CTS frame, each Station that has not sent an RTS knows that the AP is busy. Note that the total number of re-transmissions is now 47, and they could be caused by any other factor, e.g. congestion.

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Figure 5: Screenshot of the wlan.fc.retry==1 filter applied to the Node 9 - 1 capture, with RTS/CTS

## 3.5 Overhead with RTS/CTS

Before tackling the calculation of the overhead, we should ask ourselves the question "would it be any different?" to which the **answer is obviously yes**. By enabling the RTS/CTS mechanism, we unavoidably generate more traffic, since the nodes now have to exchange additional frames. To identify which frames to consider,

we used the  $wlan.addr == 00:00:00:00:00:0d \mid | udp$  filter to isolate the packets we were interested in analyzing, and added each individual overhead:

$$(14*490) + (20*245) + (64*243) = 27312B$$
 (1)

In conclusion, the total overhead with RTS/CTS is equal to 27312B > 18982B.

Figure 6: Screenshot of the wlan.addr ==  $00:00:00:00:00:00:0d \parallel udp$  filter applied to the Node 15 capture, with RTS/CTS