

# Fuzzing

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20210209

# Outline

Introduction

Principles of fuzzing

The new generation

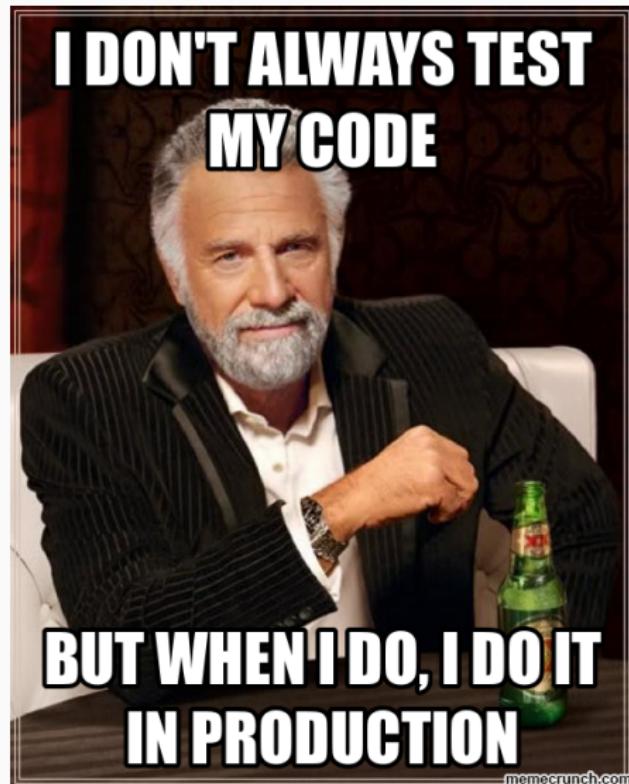
# Introduction

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# Fuzzing : your code is buggier than mine



It's about testing



# What is fuzzing ?

At its core, fuzzing **is** random testing.

# A short (selective) pre-history

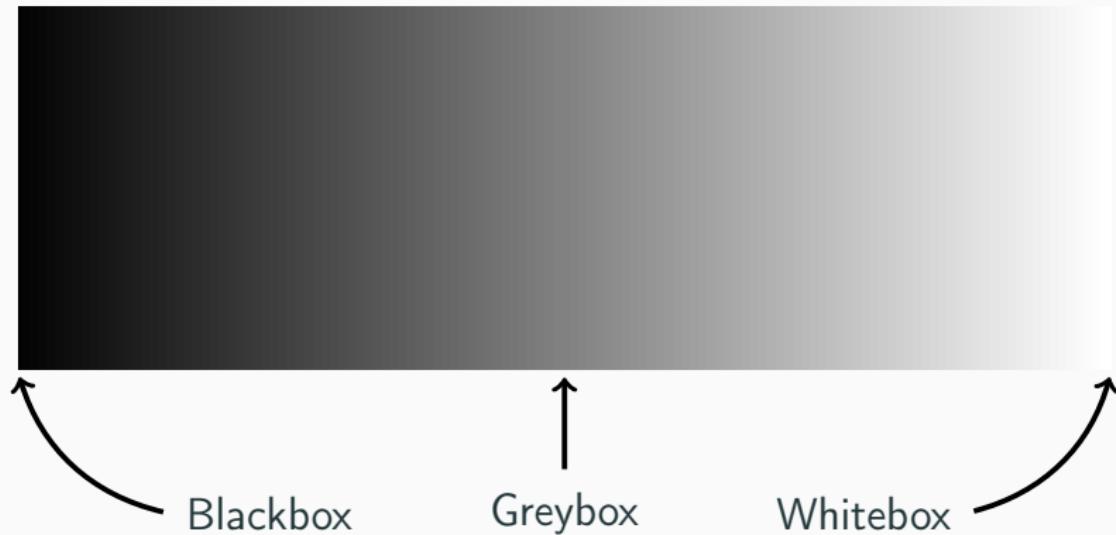
- 1981** Random testing is a cost-effective alternative to systematic testing techniques (Duran & Natos)
- 1983** "The Monkey" (Capps)
- 1988** Birth of the term "fuzzing" (Miller)

# The initial assignment

*The goal of this project is to evaluate the robustness of various UNIX utility programs, given an unpredictable input stream. [...] First, you will build a fuzz generator. This is a program that will output a random character stream. Second, you will take the fuzz generator and use it to attack as many UNIX utilities as possible, with the goal of trying to break them.*

1/3 of Unix utilities crashed, hung or failed upon fuzzing inputs.

# Shades of fuzzing



# Blackbox

## Key property

A blackbox fuzzer is unaware of the program structure

Because of that, blackbox fuzzers are necessarily limited. They thus can be considered **dumb**.

# Whitebox

## Key distinction

A whitebox fuzzer leverages program analysis to reach its targets:

Usual targets are:

1. code coverage;
2. program location.

This is usually synonym with "Dynamic Symbolic Execution".

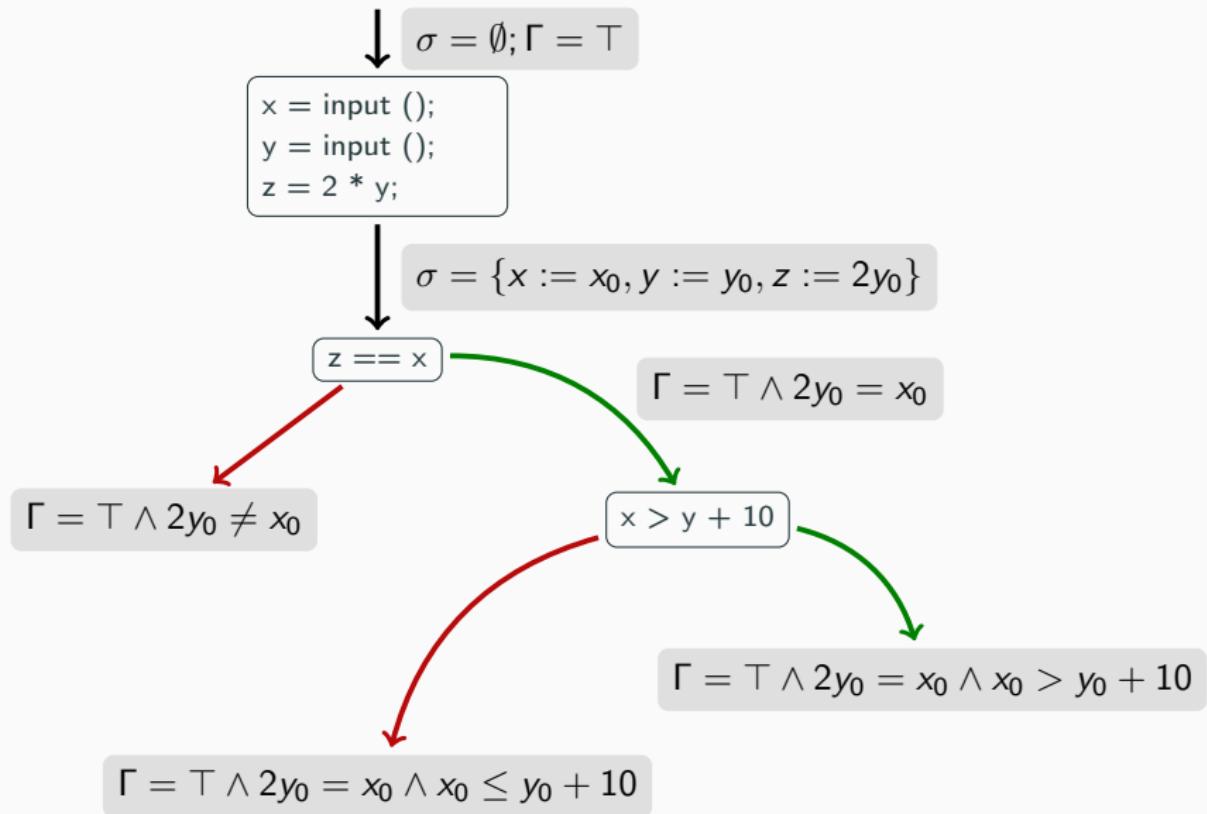
Whitebox fuzzers are **smart**.

# Coverage criteria

```
1 void
2 f(int a, int b, int * x)
3 {
4     if (a > 1 && b == 0)
5         *x = *x / a;
6     if (a == 2 || *x > 1)
7         *x = *x + 1;
8 }
```

- Instructions (I)
- All decisions (D)
- All simple conditions (C)
- All conditions / decisions (DC)
- All combinations of conditions (MC)
- All paths (P)

# Wait ? Whitebox fuzzers ?



# Greybox

Greybox fuzzers favor leveraging **instrumentations** instead of program analysis.

# Principles of fuzzing

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# Stages

1. Preprocess
2. Scheduling
3. Input Generation
4. Input Evaluation
5. Configuration Updating
6. Continue

# General algorithm (Manes et al. 2019 )

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Input:  $\mathbb{C}$ ,  $t_{\text{limit}}$

Output:  $\mathbb{B}$  // a finite set of bugs

- 1  $\mathbb{B} \leftarrow \emptyset$
  - 2  $\mathbb{C} \leftarrow \text{Preprocess}(\mathbb{C})$
  - 3 while  $t_{\text{elapsed}} < t_{\text{limit}} \wedge \text{Continue}(\mathbb{C})$  do
  - 4      $\text{conf} \leftarrow \text{Schedule}(\mathbb{C}, t_{\text{elapsed}}, t_{\text{limit}})$
  - 5      $\text{tcs} \leftarrow \text{InputGen}(\text{conf})$   
      //  $O_{\text{bug}}$  is embedded in a fuzzer
  - 6      $\mathbb{B}', \text{execinfos} \leftarrow \text{InputEval}(\text{conf}, \text{tcs}, O_{\text{bug}})$
  - 7      $\mathbb{C} \leftarrow \text{ConfUpdate}(\mathbb{C}, \text{conf}, \text{execinfos})$
  - 8      $\mathbb{B} \leftarrow \mathbb{B} \cup \mathbb{B}'$
  - 9 return  $\mathbb{B}$
-

# Stages (summarized) i

## Preprocess

- User supplies set of fuzz configurations as input, gets a potentially-modified set of fuzz configurations.
- May perform a variety of actions
  - insert instrumentation code to Program Under Tests (PUT),
  - measure execution speed of seed files
  - etc.

# Stages (summarized) ii

## Schedule

- Takes in :
  - current set of fuzz configurations
  - current time, and
  - a timeout
- Selects a fuzz configuration to be used for the current fuzz iteration.

# Stages (summarized) iii

## Input Generation

- Takes in a fuzz configuration
- Returns a set of concrete test cases.
- Some fuzzers use a seed in for generating test cases, while others use a model or grammar as a parameter.

# Stages (summarized) iv

## Input Evaluation

- Takes in :
  - a fuzz configuration
  - a set of test cases, and
  - a bug oracle.
- Executes PUT on test cases and checks if executions violate the security policy using the bug oracle.
- Outputs set of bugs found  $\mathbb{B}'$  and information about each of the fuzz runs.

# Stages (summarized) v

## Configuration Update

- Takes in:
  - a set of fuzz configurations,
  - current configuration, and
  - the information of each fuzz runs.
- May update the set of fuzz configurations.

For example, many grey-box fuzzers reduce the number of fuzz configurations based on fuzz runs.

# Stages (summarized) vi

## Continue

- Takes a set of fuzz configurations as input and outputs a boolean indicating whether a next fuzz iteration should happen or not.
- This models white-box fuzzers that can terminate when there are no more paths to discover.

# Preprocess: Instrumentation

## Goal

Gather execution feedback during runs

## Types of instrumentation

**Static** usually at compile time, sometimes rewriting the binary (e.g., afl-gcc)

**Dynamic** more costly but can instrument dlls (e.g., afl-qemu)

# Preprocess: Seed Selection

## Goal

Find minimal set that maximizes a coverage metric

## Examples of metrics

**AFL** branch coverage with logarithmic counters on each branch

**Honggfuzz** # executed instructions, branches, unique basic blocks

# Scheduling Problem

## Goal

Pick the configuration that is the most likely to lead to the most favorable outcome.

- finding most unique bugs
- maximizing coverage

## Exploration vs Exploitation

Scheduling balances:

- time gathering more information on configurations (**exploration**) and
- time fuzzing configurations believed to lead to favorable outcomes (**exploitation**).

# Scheduling: Blackbox

Blackbox fuzzers can only use fuzz outcomes (time spent, bugs found).

## Examples

- Favoring the configuration with a high success rate ( $\# \text{bugs} / \# \text{runs}$ )
- Prefer faster configurations (collection of information on them is quicker) & fix the time per configurations selection instead the number of runs (avoid spending a lot of time in slow configuration).

## Scheduling: Greybox i

### AFL

- Maintain a population of configurations with a fitness metric, apply some degree of genetic transformation (mutation, recombination)
- On a control-flow edge, AFL considers fastest and smallest inputs as "fitter"
- Fix number of run per selection.

## Scheduling: Greybox ii

### AFLFast

- On a control-flow edge, favor the one that has been chosen least
- On tie, favor the one that exercises path that has been selected the least
- The fuzzing time follows a power schedule

# Input generation: Generation-based

## Predefined model

- Network protocols
- EBNF
- System calls (types and number of arguments)

## Inferred

- Synthesize the grammar of the parser
- Capture packets and infer network protocols
- Observe I/O behavior and infer state machine
- Machine learning

## Input generation: Mutation-based

- Use initial seeds providing a structure of a valid input (file, network packets, . . . )
- Mutate portions of previous inputs to generate new *mostly valid* test cases.

# Bit-flip

## Basics

Flip:

- a fixed-number of bits or
- a random number of bits

## Mutation ratio

Determines the number of bit flips for a single generated input.

A good mutation ratio can be inferred through program analysis.

# Arithmetic mutation

## Basics

Perform an arithmetic operation on a sequence of bytes seen as an integer.

## Example

AFL selects 4-bytes values  $i$ , generates a random integer  $r$  and applies  $i + r$  or  $i - r$ .

The range of  $r$  is tunable.

# Block-based mutation

## Block

Arbitrary sequence of bytes

## Mutations

1. Insert a new block at a random position
2. Delete a randomly selected block
3. Replace a randomly selected block
4. Permutes the order of block sequences
5. Resize a seed by appending a block
6. Take a random block from another seed to insert/replace

# Dictionary-based mutation

Some fuzzers use "magic" values like 0 or  $-1$ , or format strings for mutation.

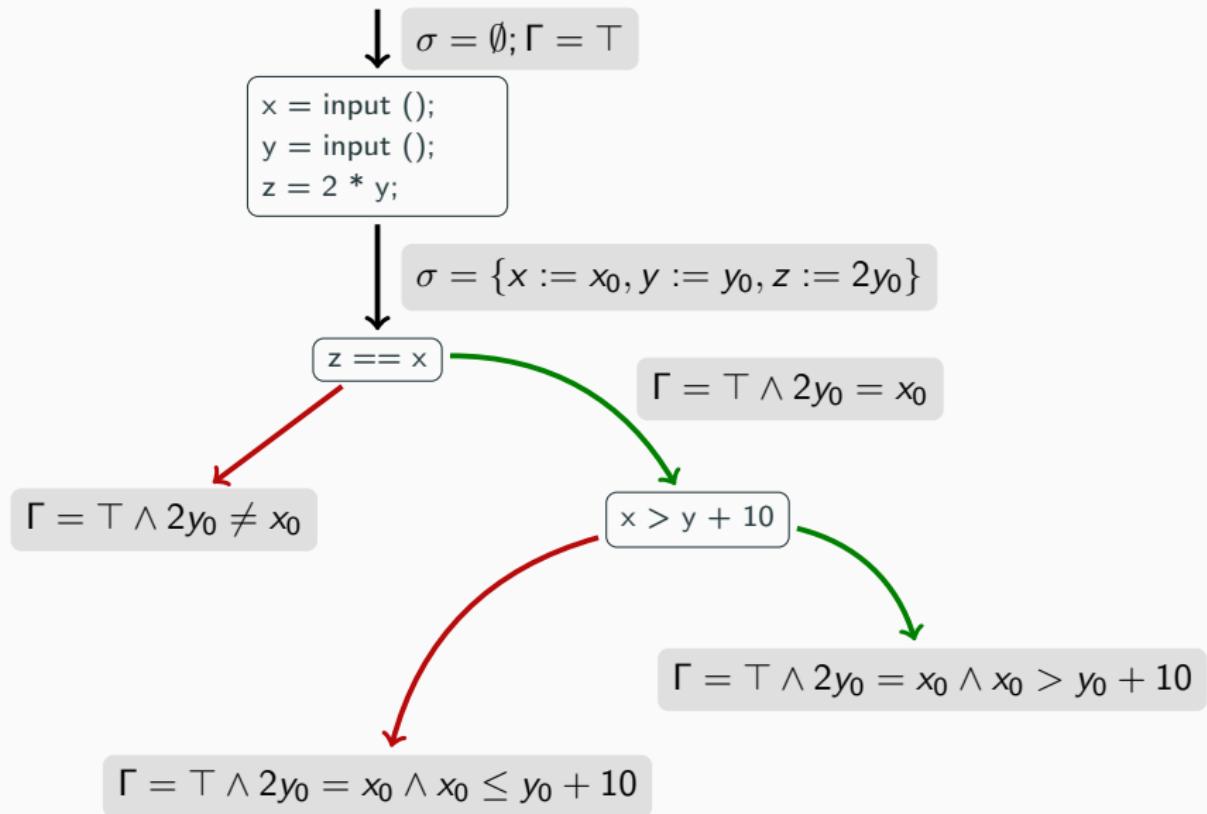
## Examples

**AFL** Uses 0, 1,  $-1$  for integer mutations

**Radamsa** unicode strings

**GPF** `%x`, `%s` for string mutations

# What about whitebox fuzzing?



# Input Evaluation: Bugs

## Default policy

A program execution terminated by a fatal signal is a violation

This is enough for memory vulnerabilities since they usually trigger segmentation faults.

## Problem

The default policy will not detect corruptions leading to valid addresses.

To this effect, code **sanitizers** are needed, i.e., a form of runtime monitor that will trigger a fatal signal on property violation.

# Input Evaluation: Triage

## Deduplication

Prune test cases triggering the same bugs as another one.

Implemented by stack backtrace hashing or coverage-based deduplication (AFL).

## Prioritization

Rank or group violating test cases according to severity and uniqueness, aka determines a form of **exploitability**.

`!exploitable`: 4-rank scales (yes, probably, unknown, not likely).

# Input Evaluation: Triage ii

## Test case minimization

Identify the part of the test case that is necessary to trigger the violation.

Produce a smaller test case (e.g., AFL sets bytes to 0 and shorten the test case).

This is not specific to fuzzers and thus dedicated techniques have been developed and can be reused here.

# Configuration Updating

## Evolving the seed pool

Select the fittest.

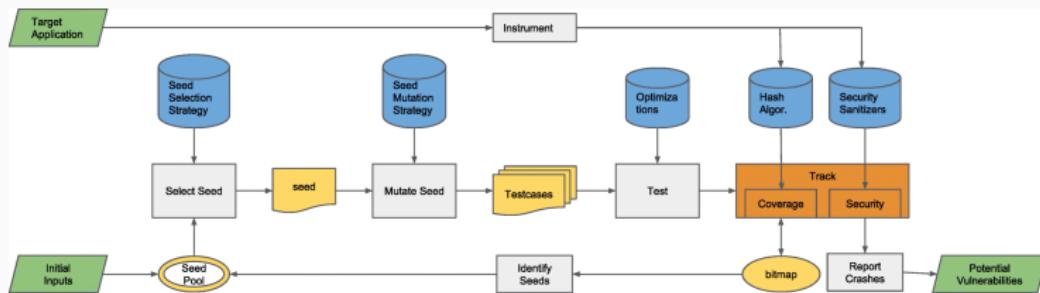
## Example

- Use on node/branch coverage: a newly discovered branch/node is sign of fitness !
- AFL also takes into account the number of times a branch has been taken
- Angora also adds calling context
- Steelix checks input offsets affect the progress in comparison instructions

## A concrete example : AFL



# The AFL loop



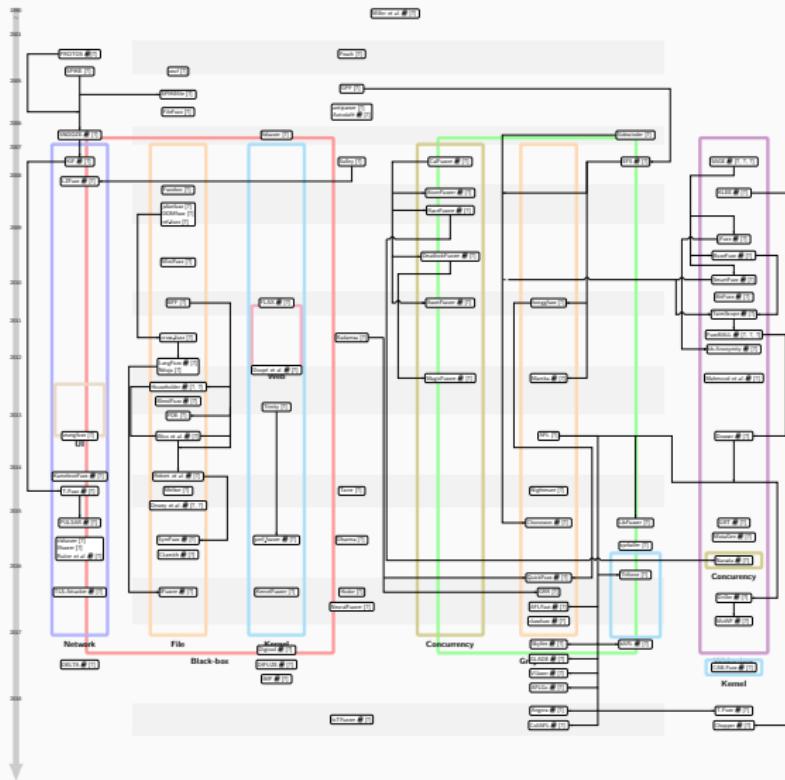
# Efficiency

<a href="#">clang / lld</a>	<a href="#">nasm</a>	<a href="#">ccls</a>
<a href="#">musl</a>	<a href="#">procmail</a>	<a href="#">fontconfig</a>
<a href="#">polkit</a>	<a href="#">Qt</a>	<a href="#">wavepack</a>
<a href="#">redis / lua-omgpack</a>	<a href="#">tigl</a>	<a href="#">privoxy</a>
<a href="#">perl</a>	<a href="#">libmpc</a>	<a href="#">radare2</a>
<a href="#">Slackware</a>	<a href="#">lwkern [reported by author]</a>	<a href="#">X.Org</a>
<a href="#">esifprobe</a>	<a href="#">jhead</a>	<a href="#">caproto</a>
<a href="#">XenFS C</a>	<a href="#">instacam</a>	<a href="#">djvulibre</a>
<a href="#">elixir</a>	<a href="#">Linux tools</a>	<a href="#">Knot DNS</a>
<a href="#">curl</a>	<a href="#">wpa_supplicant</a>	<a href="#">libde205 [reported by author]</a>
<a href="#">dnsmasq</a>	<a href="#">libpbig</a>	<a href="#">lame</a>
<a href="#">Ibmwmf</a>	<a href="#">usdecode</a>	<a href="#">MuPDF</a>
<a href="#">imlib2</a>	<a href="#">libraw</a>	<a href="#">libbson</a>
<a href="#">libress</a>	<a href="#">yarn</a>	<a href="#">W3C tidy-xmlfi</a>
<a href="#">VLC</a>	<a href="#">FreeBSD syscons</a>	<a href="#">John the Ripper</a>
<a href="#">screen</a>	<a href="#">trinix</a>	<a href="#">mosh</a>
<a href="#">UPX</a>	<a href="#">indent</a>	<a href="#">openjpeg</a>
<a href="#">MMX</a>	<a href="#">OpenMPPT</a>	<a href="#">nv</a>
<a href="#">dhcpd</a>	<a href="#">Mozilla NSS</a>	<a href="#">Nettle</a>
<a href="#">mbedtls TLS</a>	<a href="#">Linux netlink</a>	<a href="#">Linux eCryptfs</a>
<a href="#">Linux xfs</a>	<a href="#">botan</a>	<a href="#">expat</a>
<a href="#">Adobe Reader</a>	<a href="#">libav</a>	<a href="#">libical</a>

# The new generation

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# Genealogy (Manes et al. 2019)



# Driller (2016)

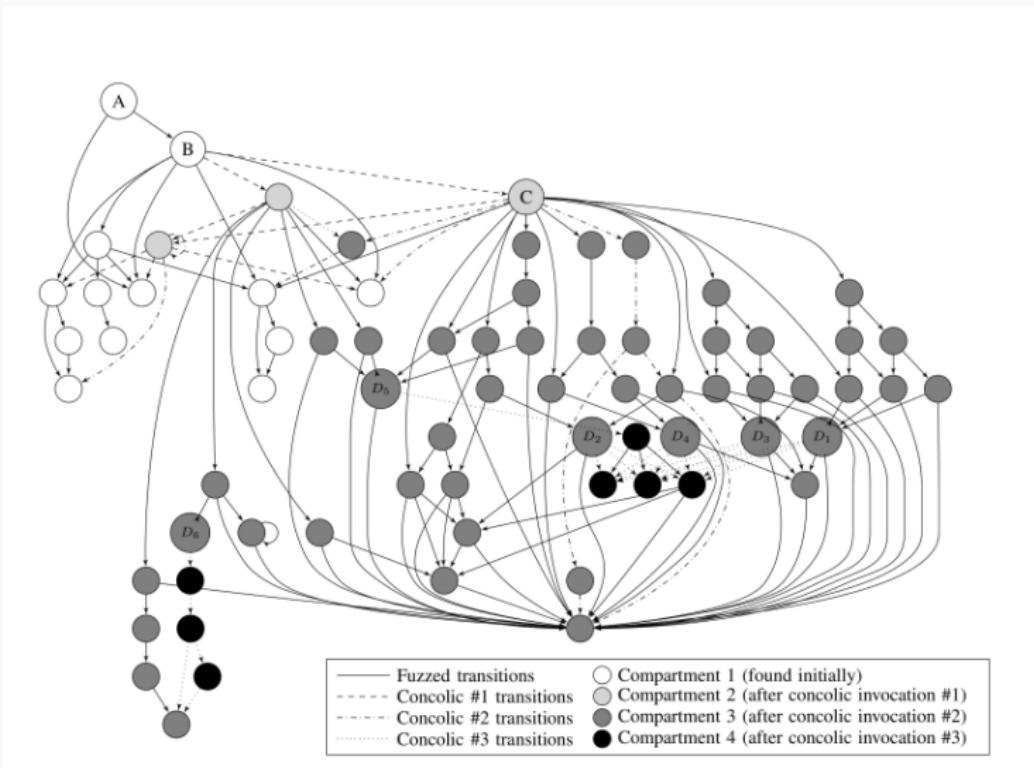
## Goal

Combination of SE & greybox fuzzing

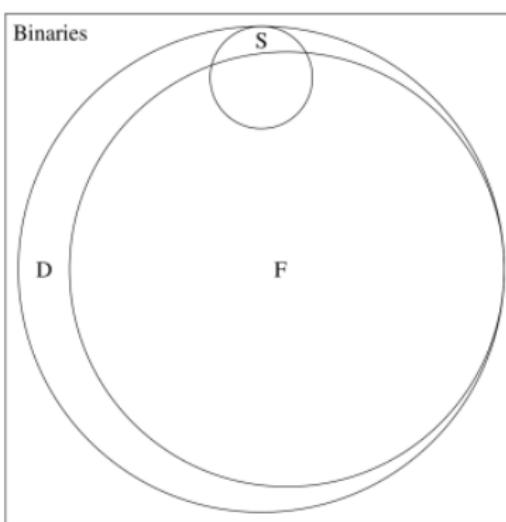
## Key insight

When the fuzzer finds "hard conditions", launch SE to solve them, then let the fuzzer continue with this new input.

# Driller : Inner workings



# Driller : Results



Method	Crashes Found
Fuzzing	68
Fuzzing $\cap$ Driller	68
Fuzzing $\cap$ Symbolic	13
Symbolic	16
Symbolic $\cap$ Driller	16
Driller	77

# Vuzzer (2017) i

## Proposal

Configurations are added only upon discovering a new non-error handling block (statically determined).

The configuration fitness is the weighted sum of the log of the frequency over exercised blocks.

## Key insight

Error-handling blocks lower the chance of vulnerabilities.

## Consequence

VUzzer prefers normal blocks that are rare according to CFG random walks.

# Vuzzer (2017) ii

**Table 1:** Vuzzer vs AFL: #crashes

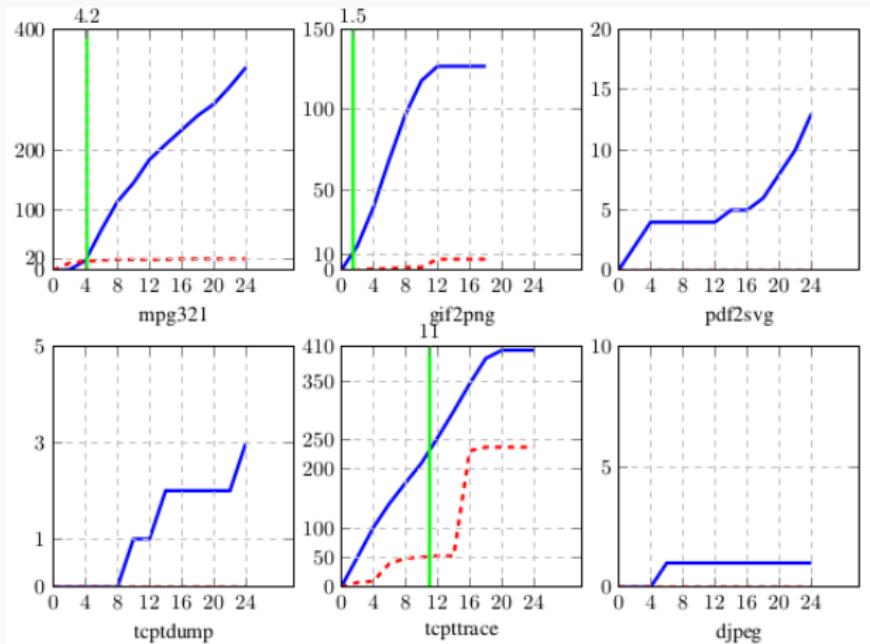
Software	Vuzzer	AFL
mpg321	337	19
gif2png/libpng	127	7
pdf2svg/libpoppler	13	0
tcdpump/libpcap	3	0
tcptrace/libpcap	403	238
djpeg/libjpeg	1	0

# Vuzzer (2017) iii

**Table 2:** Vuzzer vs AFL: #inputs

Software	Vuzzer	AFL
mpg321	23.6k	883k
gif2png/libpng	43.2k	1.84m
pdf2svg/libpoppler	5k	923k
tcdpump/libpcap	77.8k	2.89m
tcptrace/libpcap	40k	3.29m
djpeg/libjpeg	90k	35.9m

# Vuzzer (2017) iv



# Vuzzer (2017) v

Table 3: New bugs discovered

Program	Bug type	Fixed ?	Reported
tcpdump	Oob read	✓	✗
mpg321	Oob read	✗	✓
mpg321	Double free	✗	✓
pdf2svg	Null ptr deref	✓	✗
pdf2svg	Abort	✓	✗
pdf2svg	Assert failure	✓	✗
tcptrace	Oob read	✗	✓
gif2png	Oob read	✗	✓

# AFLGo (2017) i

## Goal

Generating inputs with the objective of reaching a set of program locations.

Key idea : Directed Greybox Fuzzing

## Applications

- Patch testing
- Crash reproduction
- Static analysis report verification
- Information flow detection

# AFLGo (2017) ii

## Measure

- Distance between a function & a set of target functions  
= harmonic mean
- Harmonic mean can distinguish between a node that is close to one target and further from another and one that is equidistant from both (average mean may be equal).

## Scheduling

Key insight :: simulated annealing

Use more energy to fuzz seeds closer to the targets.

Enter exploitation after given exploration time has elapsed.

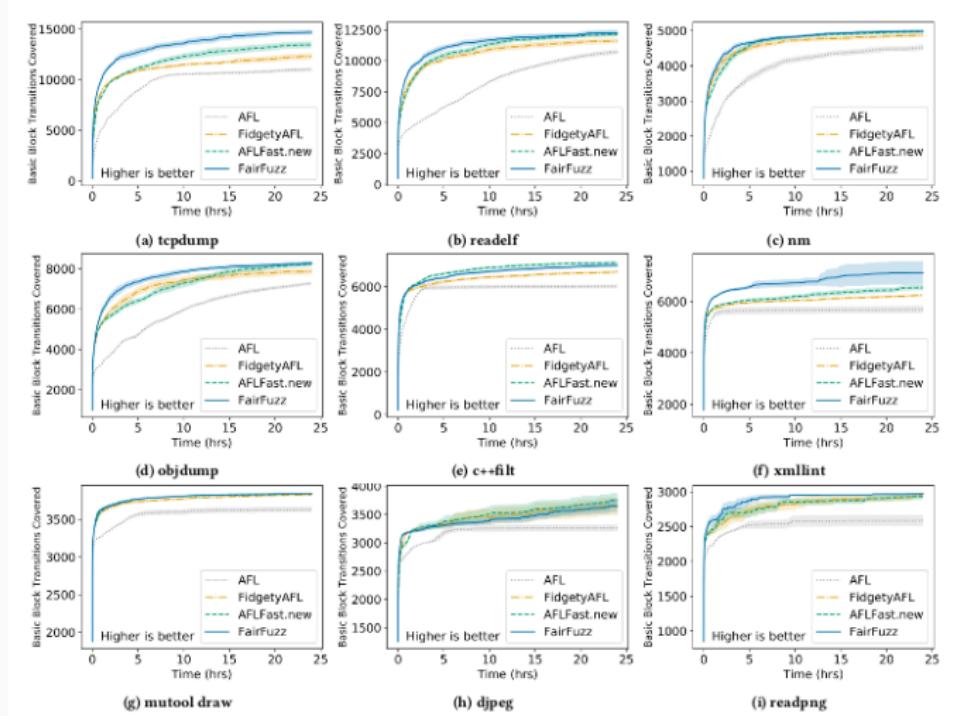
## Goal

- Achieve better branch coverage for AFL
- No extra instrumentation

## Key steps

1. Identify *rare* branches
2. New mutation technique to increase probability of hitting rare branches.

# FairFuzz (2018) ii



# Angora (2018)



# Angora (2018) i

## Goal

Increase branch coverage by solving path constraints *without symbolic execution*

## Key ingredients

- Context-sensitive branch coverage
- Byte-level taint tracking
- Gradient descent-based search
- Type & shape inference

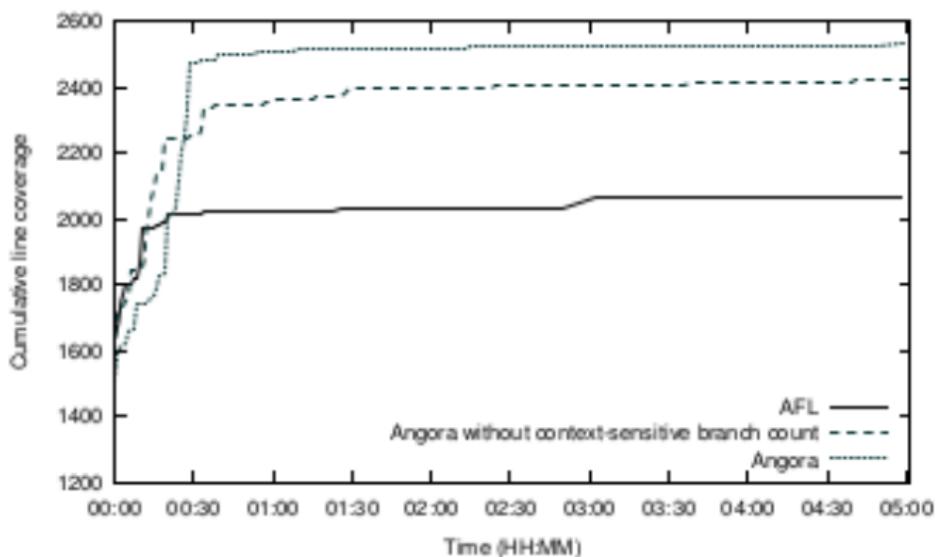
# Angora (2018) ii

**Table 4:** Results

Program	Listed bugs	Angora	AFL	Vuzzer	Steelix
uniq	28	29	9	27	7
base64	44	48	0	17	43
md5sum	57	57	1	✖	28
who	2136	1541	1	50	194

Omitted from table : SES, FUZZER

# Angora (2018) iii



# Angora (2018) iv

**Table 5:** Angora vs AFL line coverage

Program	AFL	Angora
file-5.32	2070	2534
jhead-3.00	347	789
xmlwf(expat)-2.2.5	1980	2025
djpeg(ijg)-v9b	5401	5509
readpng(libpng)-1.6.34	1592	1799
nm-2.29	6372	7721
objdump-2.29	3448	6216
size-2.29	2839	4832

# Angora (2018) v

**Table 6:** Angora vs AFL branch coverage

Program	AFL	Angora
file-5.32	1462	1899
jhead-3.00	218	789
xmlwf(expat)-2.2.5	2905	3158
djpeg(ijg)-v9b	1677	1782
readpng(libpng)-1.6.34	872	1007
nm-2.29	4105	4693
objdump-2.29	2071	3393
size-2.29	1792	2727

# Angora (2018) vi

**Table 7:** Angora vs AFL unique crashes

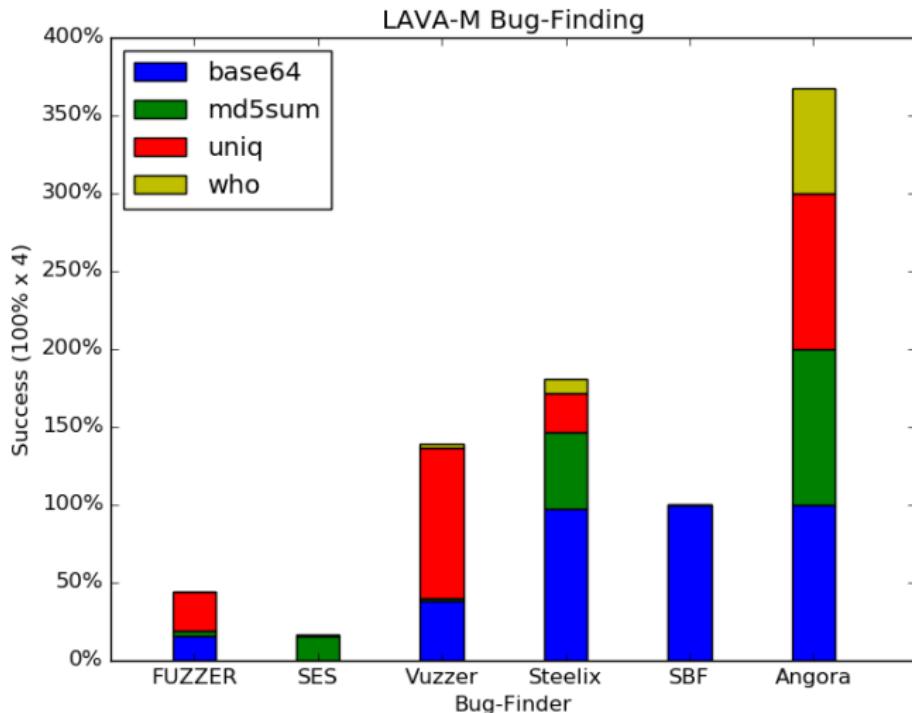
Program	AFL	Angora
file-5.32	0	6
jhead-3.00	19	52
xmlwf(expat)-2.2.5	0	0
djpeg(ijg)-v9b	0	0
readpng(libpng)-1.6.34	0	0
nm-2.29	12	29
objdump-2.29	4	48
size-2.29	6	48

# Achievements summary

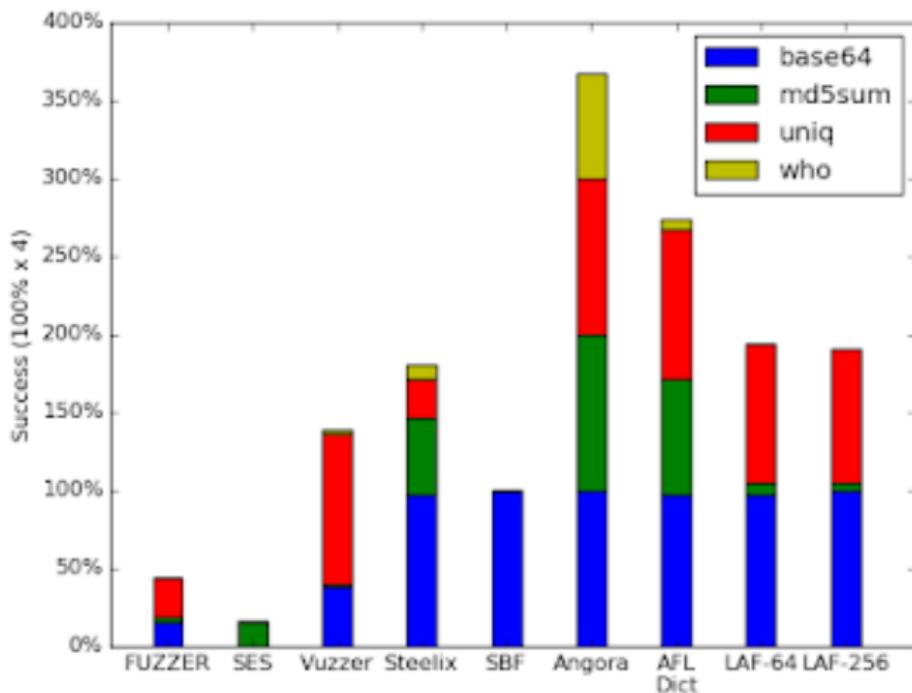
Ideas have been explored

- Better branch coverage (deeper, broader)
- Directed targeting
- Lightweight dynamic constraint solving
- Combination with other analyses:
  - Symbolic execution
  - Static analyses

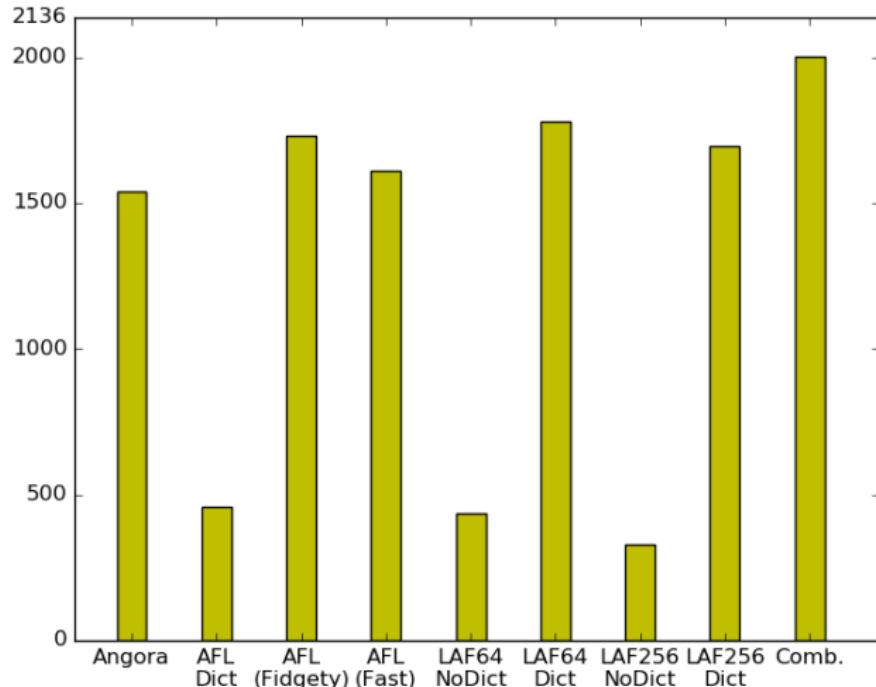
# Standard benchmarks: an emerging concern



# Baseline



# Who's who



# More goodness to come . . .

Fuzzing is a **very active** research area

Check <https://wcventure.github.io/FuzzingPaper/>

## New developments in 2019-2020

- Hawkeye (CCS '18)
- DigFuzz (NDSS '19)
- MemFuzz (ICSE '19)
- Eclipser (ICSE '19)
- Matryoshka (CCS '19)
- SAVIOR (S&P '20)
- ...

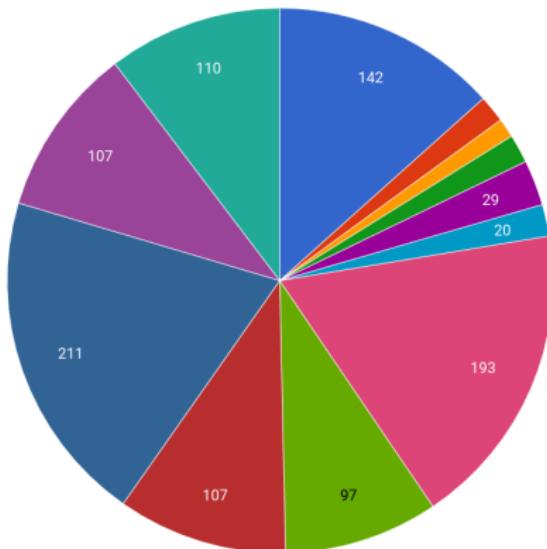
# FAAS

Recent years have seen initiatives from "MAANG" members

**Google OSS-Fuzz**

**Microsoft Project Springfield**

# Types of bugs (OSS-Fuzz 2016)



- heap buffer overflows
- global buffer overflows
- stack buffer overflows
- use after frees
- uninitialized memory
- stack overflows
- timeouts
- ooms
- leaks
- ubsan
- unknown crashes
- other (e.g. assertions)

# Anti-fuzzing (Antifuzz, 2019) i

## Attacking fuzzer's assumptions

Fuzzers depends on some implicit support:

- Coverage feedback
- Crash detection
- Speedy loop
- (sometimes) Solvable constraints

## Coverage feedback countermeasure

Add fake function calls depending on input. This add many new paths, all of them "interesting"

# Anti-fuzzing (Antifuzz, 2019) ii

## Crash detection countermeasure

Use common anti-debugging technique.

Fake crash catching

## Execution speed countermeasure

Check validity of inputs : this induces a slowdown that is enough to delay fuzzers.

## Thwart constraint solving

Instead of checking conditionals against a value  $v$ , check against  $\text{sha256}(v)$ .

# Evaluating anti-fuzzing

Fuzzers do not find bugs in LAVA-M anymore

Code coverage is reduced by  $\geq 90\%$

Performance overhead is negligible

# Questions ?



<https://rbonichon.github.io/teaching/2021/asi36/>