3. Solutions

ASI36

2018

1 Basics (basics.c)

Assuming this program has been compiled to an executable named p.

1.1 Question 1

This program may have both expected and unexpected behaviors The expected behaviors are:

- Printing "Usage: p num1 num2" and exiting due to not enough arguments.
 This can be triggered with the following input p 1
- 2. Printing "You lose"

Whenever x == 0, this will happen. If you only look at the syntactic level, this should be always, since there is no assignment to x. This is however not the case as we will see below.

p 1 2 will exhibit this behavior

Other behaviors that are "kind of" unexpected are the following:

1. Printing "You win".

To print this message, one needs to overwrite the value of the local x with something other than 0. Looking at the assembly (objdump -d -M intel p), we have the following initialization sequence:

```
1 | 59d: c6 45 e3 00 mov BYTE PTR [ebp-0x1d],0x0 ; x
2 | 5a1: c7 45 db 00 00 00 mov DWORD PTR [ebp-0x25],0x0 ; t[0-3]
3 | 5a8: c7 45 df 00 00 00 mov DWORD PTR [ebp-0x21],0x0 ; t[4-7]
```

So t ends at ebp - 0x21 + 0x4 and x is located at ebp - 0x1d. So there is a no gap between the end of t and x. If we can write 9 bytes from the start of t, we might rewrite x as well.

The number of writes is controlled through argv[2]; what we write by argv[1]. We want to write something other than 0 (say 1).

For example, p 1 8 does that (p 11 2222 as well).

2. Looping forever

3. crash

The two behaviors below come from the same problem. It is also possible to overwrite i. On my machine it is at ebp - 0x1c, right above x on the stack — you can locate by putting a value at initialization in it if you wish.

If you run p 1 25 (whatever value greater than 9 instead of 25 and whatever value in [0..8] instead of 1 does it). You will loop forever. You can observe it in gdb

```
1 break main
2 watch i
3 continue
```

You will see the value of i loop until 9 then come back to 2 since t[10] points to i and rewrites it with 1 in our case. Then it is incremented back to 2.

Now if the value you put instead of 1 does not make the index i come back to a range between [0..8] then you can either win – if you exit the loop with a value other than 0 for x, or provoke a segmentation fault if you jump above x and continue overwriting after i.

The former is achieved for example by p 11 12, the latter p 14 68. You may even print you win then provoke a segmentation fault: p 14 50 does that.

1.2 Question 2

All the expected behaviors are still observable but not the others.

There are two reasons:

- 1. The order of the locals has been changed: the buffer is now above x and i and thus cannot change them anymore, e.g., to print "You win" or loop.
- 2. Only the potential of 3 remains but is replaced by the message whenever you overwrite the bounds of the stack frame.

```
1 *** stack smashing detected ***: <unknown> terminated
```

The initialization is now:

```
1 61c:
               8b 50 04
                                            mov
                                                   edx, DWORD PTR [eax+0x4]
2 61f:
               89 55 d4
                                                   DWORD PTR [ebp-0x2c],edx
                                            mov
               65 8b 0d 14 00 00 00
                                                   ecx,DWORD PTR gs:0x14
3 622:
                                                                           ; canary
                                            mov
4 629:
               89 4d f4
                                                   DWORD PTR [ebp-0xc],ecx
                                            mov
5 62c:
              31 c9
                                            xor
                                                   ecx,ecx
               c6 45 e7 00
6 62e:
                                            mov
                                                   BYTE PTR [ebp-0x19],0x0;x
  632:
              c7 45 ec 00 00 00 00
                                                   DWORD PTR [ebp-0x14],0x0 ; t[0-3]
                                            mov
8 639:
               c7 45 f0 00 00 00 00
                                                   DWORD PTR [ebp-0x10], 0x0; t[4-7]
  640:
               c6 45 ec 30
                                                   BYTE PTR [ebp-0x14],0x30
              c7 45 e8 01 00 00 00
10 644:
                                            mov
                                                   DWORD PTR [ebp-0x18], 0x1; i
```

2 Take the heap (h.c)

The program has a potential vulnerability on the heap, since p and p3 are both dynamically allocated. f seems to correctly check against strcpy manipulation but the real problem lies the handling of p3.

Now let's try to examine what happens right before the scanf. Any entry triggering free(p) suffices. For example AAAABBBBCCCCDDDD as argv[1] is long enough.

If you add a break at the scanf and enter the two lines above, you will see that p and p3 points to the same address: malloc has reused the freed space. It means that system(p) will execute whatever you enter. So now if you enter, say fortune or sh, you will execute this program.

3 Format-string exploitation (fmt.c)

This exercise is explained in the book "Hacking: the Art of Software Exploitation" in the relevant section about format string exploitation.