

Recursion

"We achieve more when we chase the dream instead of the competition."

Simon Sinek

Good Evening



Today's content

- Function call tracing
- Recursion
- How to write a recursive code?
- Tracing the recursive code
- TC / SC of recursive code { Next class }

Why recursion?

- Used in Mergesort / Quicksort
- Binary Tree / BST / Tries
- Dynamic programming
- Back tracking
- Graphs

Function call tracing

```
int add(n,m){  
    return n+m;  
}
```

```
int mul(x,y){  
    return x*y;  
}
```

```
int sub(x,y){  
    return x-y;  
}
```

Last In First Out

Stack

```
main()  
x=10, y=20  
print(sub(mul(add(x,y),30),75));  
}  
    ↳ sub(mul(add(x,y),30),75); = 825  
        ↳ mul(mul(add(x,y),30),30) = 900  
            ↳ add(x,y) = 30  
                10, 20
```

main()

odd(x,y) : 30
mul(add(x,y),30) : 900
sub(mul(add(x,y),30),75)
print(825) x=10, y=10

Obs 1 → Whenever we add a function, it gets added at top

Obs 2 → Whenever a function returns something, it gets wiped out of the memory (At top)

Recursion \rightarrow A function calling itself

↳ Repetitive code

↳ Solving a problem using a smaller instance of
problem

Subproblem

Q1. Sum of n natural number

$$\text{sum}(4) = 1 + 2 + 3 + 4$$

$$\text{sum}(5) = \underbrace{1 + 2 + 3 + 4}_{\text{sum}(4)} + 5$$

a subproblem

$$\text{sum}(n) = \underbrace{1 + 2 + 3 + \dots + (n-1)}_{\text{sum}(n-1)} + n$$

$$\text{sum}(n) = \text{sum}(n-1) + n$$

Steps to write Recursive code

01. Assumption/
Expectation → Fix what our function supposed to do/
what is expected from the function

02. Main logic → Solving assumption using subproblem.

03. Base condition → Last valid input for which the
recursion has to stop

01. Sum of n natural number $n > 0$

// Assumption → Given n , calculate & return the
sum of 1 to n no.

```
int sum(n)
{
    if (n == 1) return 1;
    return sum(n-1) + n
}
```

Tracing

main()

print (sum(4));
10

3

```
int sum(N) n=4
if (N==1) { return 1; }
return (sum(N-1)+N);
}
```

$$6+4$$

```
int sum(N) n=3
if (N==1) { return 1; }
return (sum(N-1)+N);
}
```

$$3+3$$

return 6

Output = 10

```
int sum(N) n=2
if (N==1) { return 1; }
return (sum(N-1)+N);
}
```

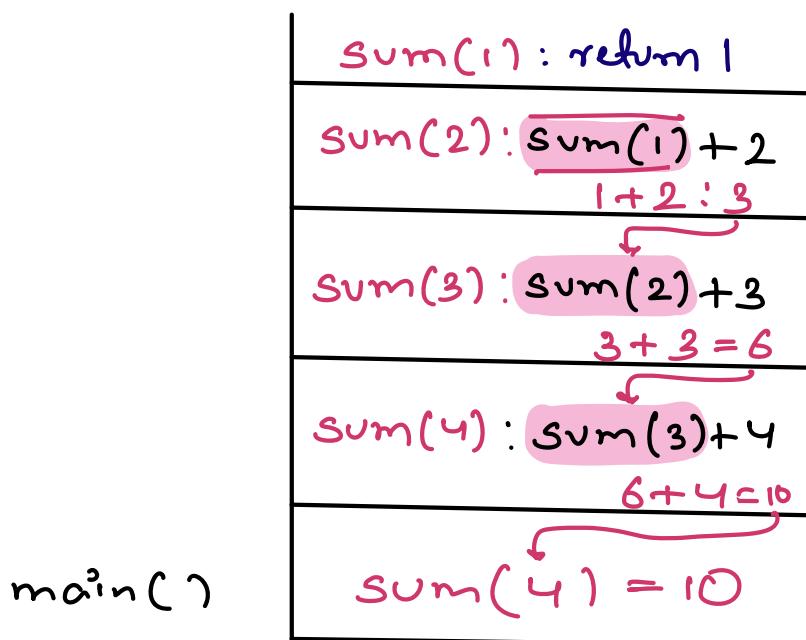
$$1+2$$

return 3

```
int sum(N) n=1
if (N==1) { return 1; }
return (sum(N-1)+N);
}
```

return 1

Stack Tracing



without base condition

Recursion will never stop

(a) TLE X

(b) Memory limit exceeded / Stack overflow

Note :- Limit to function calls $\approx 10^5$

Note :- If ever getting the stack overflow, check for base condition

Base condition → Must be present at the top

Q2. Factorial of a number ($n > 0$)

$$\text{fact}(3) = 3 * 2 * 1 = 6$$

$$\text{fact}(4) = 4 * 3 * 2 * 1 = 24$$

$$\text{fact}(n) = n * (n-1) * (n-2) * \dots * 1$$

$$\text{fact}(4) = 4 * \text{fact}(3)$$

// Assumption → Given n, calculate & return $n!$

int fact(n)

if ($n == 1$) return 1;

return fact($n-1$) * n;

3

$n \geq 0$

int fact(n)

if ($n == 0$) return 1

return fact($n-1$) * n

3

$$0! = 1$$

Fibonacci numbers ($n \geq 0$)

# Input	0	1	2	3	4	5	6	7	8
Fib()	0	1	1	2	3	5	8	13	21

$\text{fib}(n) = \text{sum of two previous fib no.}$

$$\text{fib}(7) = \text{fib}(6) + \text{fib}(5)$$

$$\text{fib}(8) = \text{fib}(7) + \text{fib}(6)$$

$$\text{fib}(n) = \text{fib}(n-1) + \text{fib}(n-2)$$

// Assumption → Calculate & return n^{th} fib number

```
int fib0(n){
```

```
    if (n ≤ 1) return n;
```

```
}
```

```
if (n == 0) return 0;
if (n == 1) return 1;
```

```
    return fib(n-1) + fib(n-2)
```

Base Case → Last valid input for which we want our main logic to stop

$$\text{fib}(0) = \text{fib}(0-1) + \text{fib}(0-2)$$

$$\text{fib}(1) = \text{fib}(1-1) + \text{fib}(1-2)$$

$$\text{fib}(2) = \text{fib}(2-1) + \text{fib}(2-2)$$

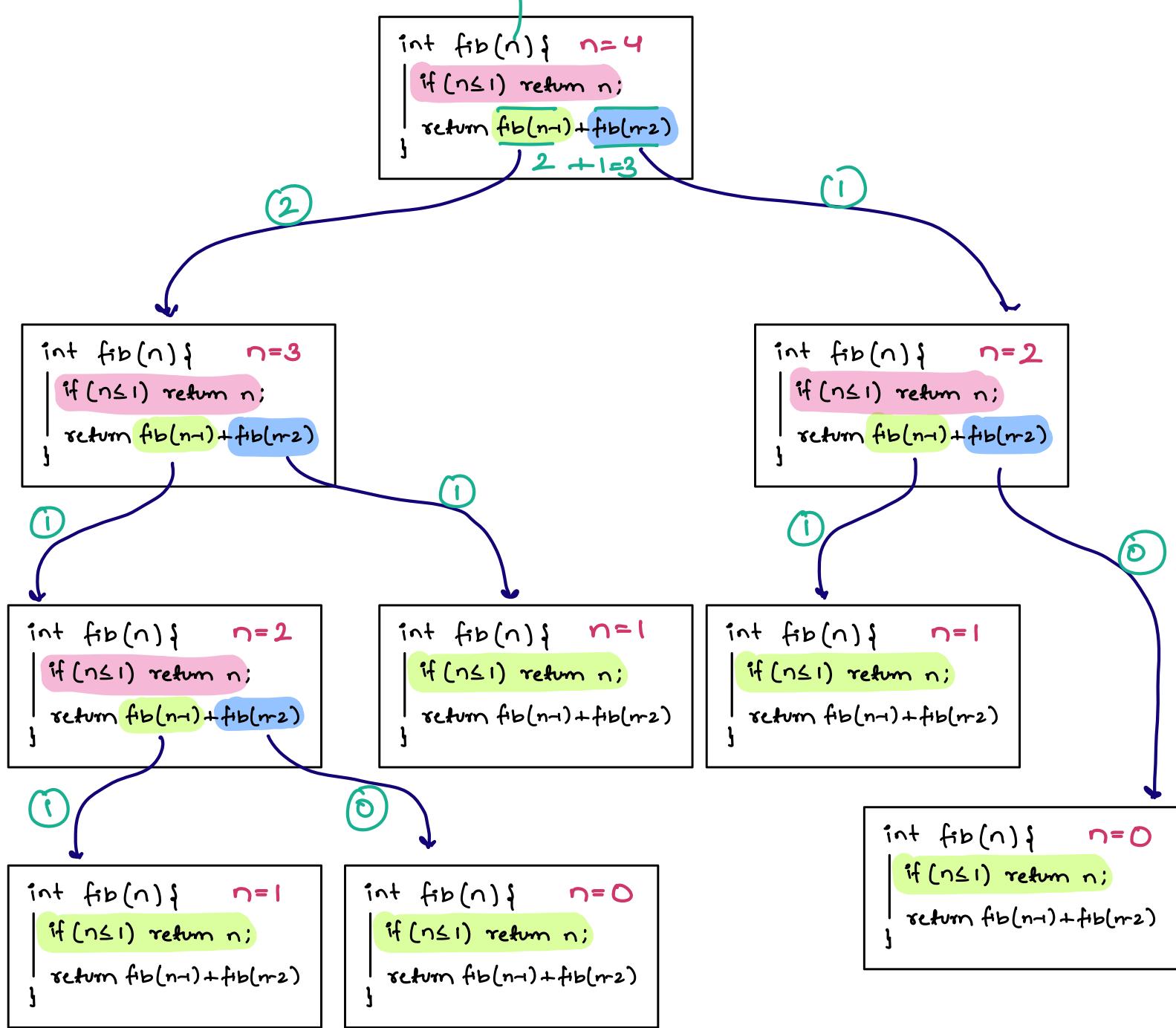
} for 0 & 1
main logic is invalid, make them as base condition

main() {

 print(fib(4)) —→ 3

3

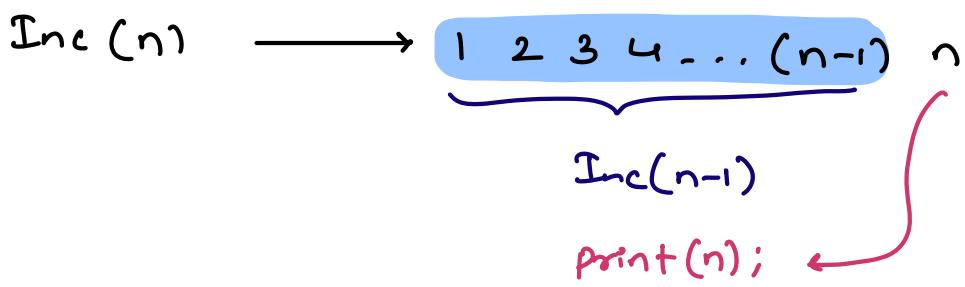
Tracing = 4



Q3. Given n , print all the numbers from 1 to N in increasing order ($n > 0$)

Inc (3) \rightarrow 1 2 3

Inc (4) \rightarrow 1 2 3 4
Inc (3) 4



// Assumption → Given n , print all the no. from 1 to n

```
void Inc (n):
    if (n==1) { point(1);
                 return; }
    Inc(n-1)
    print(n);
    }
```

Note:- Always return from the base condition

Tracing -

```
void Inc (n) n=3
if (n==1) { point(1);
             return; }
Inc(n-1)
print(n);
}
```

```
void Inc (n) n=2
if (n==1) { point(1);
             return; }
Inc(n-1)
print(n);
}
```

void Inc (n) n=1

if ($n == 1$) { print(1);
return; }

Inc(n-1)

print(n);

3

Output
1 2 3

Dec(4) → 4 3 2 1

Dec(n) → n n-1 n-2 n-3 ... 1

TODO

05. Given a ch[], check if the given ch[] is palindrome from indexes s to e (s < e)

ch[] = { a m a z o n } s = 2 e = 4 false

ch[] = { n a y a n } s = 0 e = 4 true

`ch[] = { good dad } s=4 c=6 True`
 0 1 2 3 4 5 6

`ch[] = { good dad } s=2
 0 1 2 3 4 5 6 c=5`

$\text{ch}[] = \{ \underline{s} \quad \underline{\text{stt}} - \underline{e-1} \quad \underline{e} \}$

`ch[] = { m a d a m } s=0 e=4 }` True

`ch[] = { m P c m }` $s=0$ $e=3$ } False

boolean ispalindrome (char [], s, e)

if ($s > e$) return True;

if ($ch[s] == ch[e]$ && $isplain(ch, s+1, c-1) == T$)

return True:

3

return false;

main()

ch = {noon}

s=0

e=3

print (ispali(ch,0,3))

3

True

```
boolean ispalindrome(ch,s,e){  
    if (s>e) return true;  
  
    if (ch[s]==ch[e] &&  
        ispalindrome(ch,s+1,e-1));  
        return true;  
    }  
    return false  
}
```

s=0

e=3

3

True

```
boolean ispalindrome(ch,s,e){  
    if (s>e) return true;  
  
    if (ch[s]==ch[e] &&  
        ispalindrome(ch,s+1,e-1));  
        return true;  
    }  
    return false  
}
```

s=1

e=2

True

boolean ispalindrome(ch,s,e){

if (s>e) return true;

if (ch[s]==ch[e] &&

ispalindrome(ch,s+1,e-1));

return true

} return false

s=2

e=1