

BIT MANIPULATION

2

The first step to achieving your goal, is to take a moment to respect your goal. Know what it means to you to achieve it.

DWAYNE JOHNSON



Good
Evening

Today's content

A Good news &

a bad news today

01. Power of left shift

(a) Set i^{th} bit

(b) Flip i^{th} bit

(c) Check i^{th} bit

02. Check bit is set or not

03. Count set bits

04. Negative number

05. Ranges

06. Importance of constraints

Left Shift Operator

$$n = \begin{smallmatrix} 5 & 4 & 3 & 2 & 1 & 0 \\ 1 & 0 & 1 & 0 & 1 & 0 \end{smallmatrix} \Rightarrow 2^5 + 2^3 + 2^1 = 32 + 8 + 2 = 42$$

$$n \ll 1 = \begin{smallmatrix} 5 & 4 & 3 & 2 & 1 & 0 \\ x & 0 & 1 & 0 & 1 & 0 \end{smallmatrix} \rightarrow 2^4 + 2^2 = 16 + 4 = 20$$

Quiz

$$7 \ll 3$$

$$a \ll n = a * 2^n$$

$$\begin{array}{r} \begin{smallmatrix} 31 & \dots & 5 & 4 & 3 & 2 & 1 & 0 \\ 0 & 0 & 0 & \dots & 0 & 0 & 0 & 1 & 1 & 1 \end{smallmatrix} \\ \underbrace{\hspace{1cm}}_{x} \end{array} \ll 3 \quad \Rightarrow \quad \dots 0 \ 0 \ 1 \ 1 \ 1 \ 0 \ 0 \ 0 = 2^5 + 2^4 + 2^3 = 56$$

Quiz

$$15 \ll \boxed{2} \Rightarrow 15 * 2^2$$

$$= 15 * 4 = 60$$

$$\begin{array}{r} \begin{smallmatrix} 0 & 0 & 0 & 0 & 0 & \dots & 0 & 1 & 1 & 1 \end{smallmatrix} \\ \underbrace{\hspace{1cm}}_{x} \end{array} \quad \begin{array}{r} \begin{smallmatrix} 0 & 0 & 0 & \dots & 0 & 1 & 1 & 1 & 0 & 0 \end{smallmatrix} \\ \underbrace{\hspace{1cm}}_{x} \end{array}$$

$$* \quad 1 \ll 0 = 1 \quad (\text{No shifting})$$

$$* \quad 1 \ll 1 = 0 \ 0 \ 0 \ 0 \ 0 \dots 0 \ 0 \ 0 \ 1 \ll 1$$

$$\begin{array}{r} \begin{smallmatrix} 0 & 0 & 0 & 0 & 0 & \dots & 0 & 0 & 0 & 1 & 0 \end{smallmatrix} \\ \underbrace{\hspace{1cm}}_{x} \end{array} \rightarrow 2^1 = 2$$

$$* \quad 1 \ll 2 = 4$$

$$* \quad 1 \ll 3 = 1 * 2^3 = 8$$

$$2^n = 1 << n$$

One thing ↗

$$\begin{aligned} 2 &<< (-1) \\ &= \underline{\underline{2 >> 1}} \\ &\text{return } 0 \end{aligned}$$

* Power of left shift

* $45 = (101101)_2$

01. Left shift with OR

$$\begin{array}{r} \text{OR} \\ 45 = 101101 \\ 1 << 2 = 000100 \\ \hline 101101 \end{array}$$



Left shift with
OR operator
will set the
particular bit

$$\begin{array}{r} \text{OR} \\ 45 = 101101 \\ 1 << 4 = 010000 \\ \hline 111101 \end{array}$$

$$\begin{array}{r} \text{OR} \\ 45 = 101101 \\ 1 << 1 = 000010 \\ \hline 101111 \end{array}$$

$N \mid (1 \ll i)$ → set the i^{th} bit

02. Left shift with XOR

$$\begin{array}{r} 0 \\ 1 \\ \hline 0 & = 0 \\ 0 & = 1 \end{array}$$

XOR

$$\begin{array}{r} 45 = 101101 \\ 1 \ll 2 = 000100 \\ \hline 101001 \end{array}$$

XOR

$$\begin{array}{r} 45 = 101101 \\ 1 \ll 4 = 010000 \\ \hline 111101 \end{array}$$

XOR

$$\begin{array}{r} 45 = 101101 \\ 1 \ll 1 = 000010 \\ \hline 101110 \end{array}$$

sane same puppy shame
Toggling the bits
with the help of
XOR & left shift

$N \wedge (1 \ll i) \rightarrow$ toggle the i^{th} bit

03. AND with left shift operator

AND

$$\begin{array}{r} 45 = 101101 \\ 1 \ll 2 = 000100 \\ \hline 000100 \end{array}$$

AND

$$\begin{array}{r} 45 = 101101 \\ 1 \ll 4 = 010000 \\ \hline 000000 \end{array}$$

AND

$$\begin{array}{r} 45 = 101101 \\ 1 \ll 1 = 000010 \\ \hline 000000 \end{array}$$

$N \& (1 \ll i)$

→ i^+ bit is unset

→ $1 \ll i \rightarrow i^+$ bit is set

Q1. Unset the i^+ bit of a given number if it is set, otherwise do nothing

$$n = 45 \quad \left\{ \begin{array}{c} \text{s } 4 \ 3 \ 2 \ 1 \ 0 \\ | 101101 \end{array} \right. \longrightarrow 101\textcolor{brown}{0}01$$

$$n = 45 \quad \left\{ \begin{array}{c} \text{s } 4 \ 3 \ 2 \ 1 \ 0 \\ | 101101 \end{array} \right. \longrightarrow 101101$$

Idea \rightarrow Check if the i^+ bit is set,
if yes, then toggle it

```
if (checkbit(n, i) == true)
|
|   n = n ^ (1 << i);
|
}
else {
|
|   // do nothing
|
3
```

Q How to check if i^+ bit is set or unset?

01 AND

$$N \& (1 \ll i) == (1 \ll i) \rightarrow i^+ \text{ bit is set}$$

02 OR

$$N \mid 1 \rightarrow N = \text{last bit is set}$$
$$N \mid 1 > N \rightarrow N = \text{last bit is unset}$$

$$N \mid (1 \ll i) == N \rightarrow i^+ \text{ bit is set}$$

$$N \mid (1 \ll i) > N \rightarrow i^+ \text{ bit is unset}$$

$$\begin{array}{r} N = 1010 \\ 1 \ll 2 = 0100 \\ \hline \text{Ans} = 1110 \end{array}$$

$$\begin{array}{r} N = 1011 \\ 1 \ll 3 = 1000 \\ \hline \text{Ans} = 1011 \end{array}$$

$$\underline{\underline{\text{Ans} > N}}$$

$$\text{Ans} == N$$

03

XOR operator $\rightarrow \{ \text{TODO} \}$

* Right shift operator = $a \gg x = \frac{a}{2^x}$

$$01. \quad 10 \gg 2 \Rightarrow \frac{10}{2^2} = \frac{10}{4} = 2$$

$000001010 \gg 2$

00000010

\downarrow

2

$$02. \quad 8 \gg 3 = \frac{8}{2^3} = \frac{8}{8} = 1$$

$$03. \quad 8 \gg 4 = \frac{8}{2^4} = 0$$

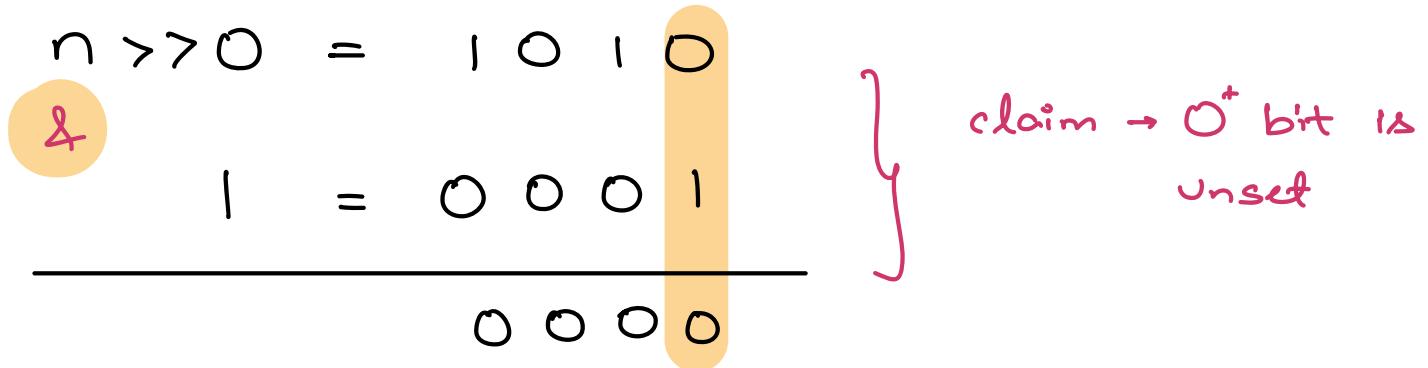
$$04. \quad 1 \gg 2 = 0$$

$N \& 1$

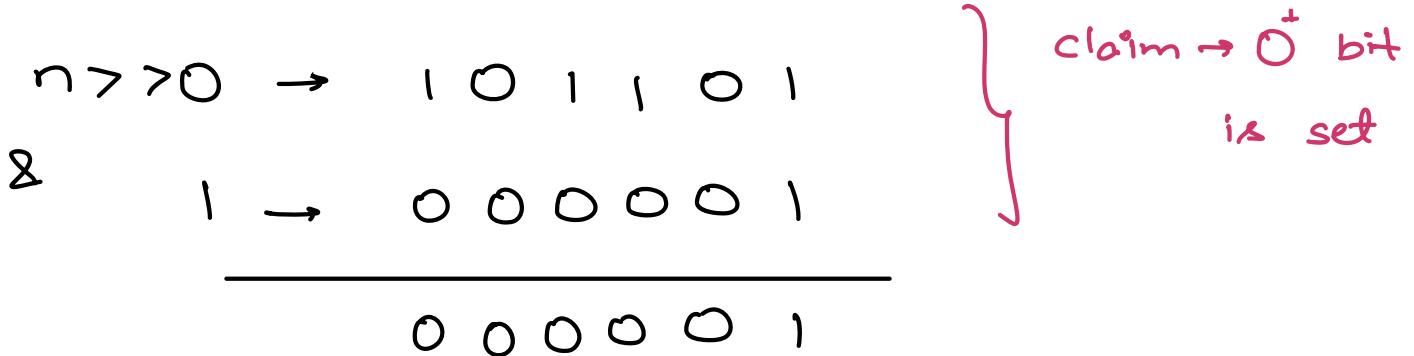
0 {last bit is unset}

1 {last bit is set}

$$\Rightarrow n = 1010$$

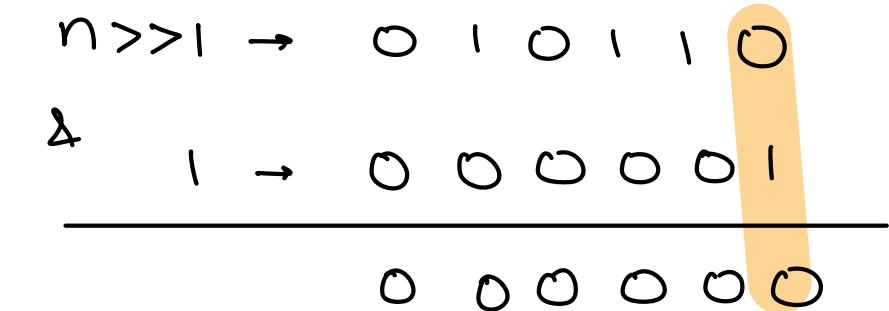


$$n = 45 \rightarrow 101101$$



$$n \rightarrow 101101$$

$$n >> 1 \rightarrow 010110$$



$n >> 2$ 001011
 & 1 0000001
 ——————
 000001

Check if i^{th} bit is set or not

$(n >> i) \& 1$ → 0 (i^{th} bit is unset)
 → 1 (i^{th} bit is set)

Q Count the no. of set bits in a no.

$n = 45$ 101101 Ans = 4

$n = 8$ 001000 Ans = 1

Idea 1 → Check for every bit

Set
 ——————
 $c = c + 1$
 unset
 ——————
 // do nothing

```

count = 0

for (i=0 ; i<32 ; i++)
    if (checkbit(n, i) == true)
        count += 1

```

return count;

Not a
Generalised
code

$n=2 \Rightarrow 00000\dots\dots 00010$

$\begin{matrix} 31 & 30 & 29 & \dots & 4 & 3 & 2 & 1 & 0 \end{matrix}$

Issues with the code

01 Unnecessary iterations

02. Hard code by taking 32

$\left\{ \begin{array}{l} \text{int } i : 0 \rightarrow 31 \\ \text{long } i : 0 \rightarrow 63 \\ \text{byte } i : 0 \rightarrow 7 \end{array} \right.$

Idea 2

count = 0

$n = \dots 1 0 1 0$	\circ
$n \gg 1 = 0 1 0 1$	$x \quad $
$n \gg 2 = 0 0 1 0$	$x \quad 0$
$n \gg 3 = 0 0 0 1$	$x \quad $
$n \gg 4 = \underline{\underline{0 0 0 0}}$	$\underline{\underline{2}}$

↳ if $n == 0$ return count

```
count = 0
while (n > 0)
    if (n & 1 == 1) {
        count = count + 1;
    }
    n = n >> 1; →  $\frac{n}{2}$ 
}
return count;
```

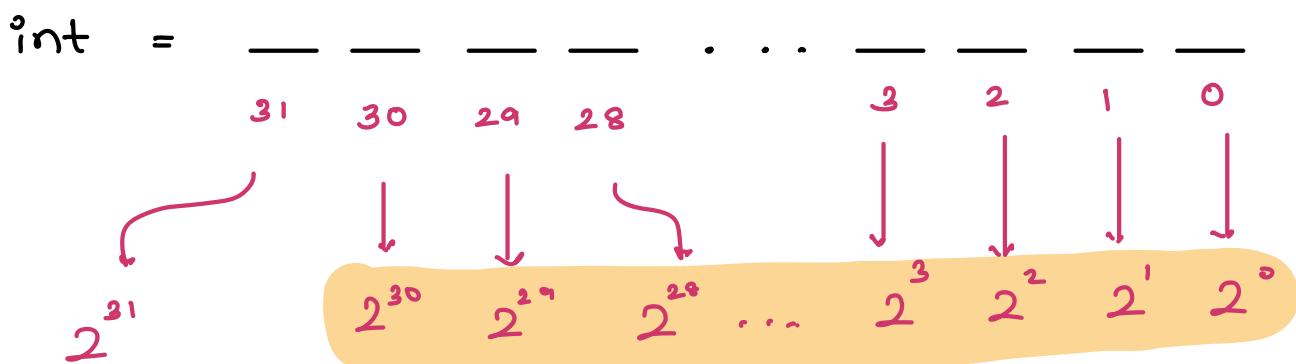
TC: $O(\log_2 n)$

SC: $O(1)$

Break time : 10:20 → 10:30 pm

* Negative numbers

$$(-45)_{10} \longrightarrow (?)_2$$



$$= 2^{30} + 2^{29} + 2^{28} + \dots + 2^3 + 2^2 + 2^1 + 2^0$$

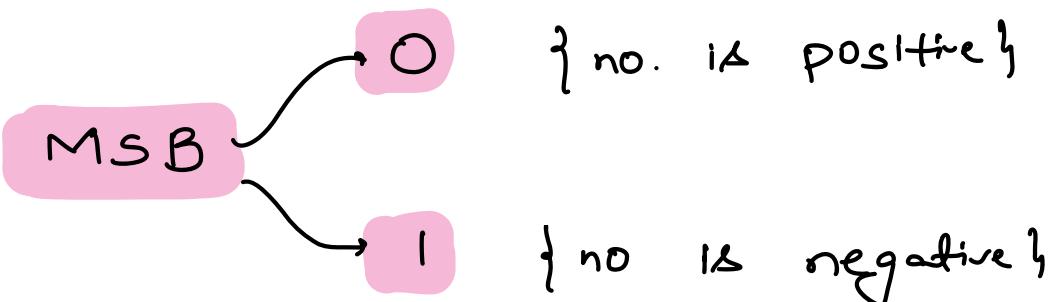
$\underbrace{\hspace{10em}}$

$$2^0 + 2^1 + 2^2 + \dots + 2^{30}$$

$$\frac{a(r^n - 1)}{r - 1} \Rightarrow \frac{2^0 * (2^{31} - 1)}{2 - 1} = 2^{31} - 1$$

$$\underbrace{2^{31}}_{>} > 2^{30} + 2^{29} + 2^{28} + \dots + 2^2 + 2^1 + 2^0$$

MSB → Most significant bit



$$(-45)_{10} \rightarrow (\quad)_2$$

Assumption → int → 8 bits

$$(45)_{10} = \underline{0} \underline{0} \underline{1} \underline{0} \underline{1} \underline{1} \underline{0} \underline{1}$$

01. 1's compliment \Rightarrow Toggle each & every bit

02. 2's compliment \Rightarrow 1's compliment + 1

$$(45) = \underline{0} \underline{0} \underline{1} \underline{0} \underline{1} \underline{1} \underline{0} \underline{1}$$

$$1's \text{ compliment} = \underline{1} \underline{1} \underline{0} \underline{1} \underline{0} \underline{0} \underline{1} \underline{0}$$

$$+ 1 = \underline{0} \underline{0} \underline{0} \underline{0} \underline{0} \underline{0} \underline{0} \underline{1}$$

$$2's \text{ compliment} = \underline{\underline{1}} \underline{\underline{1}} \underline{\underline{0}} \underline{\underline{1}} \underline{\underline{0}} \underline{\underline{0}} \underline{\underline{1}} \underline{\underline{1}}$$

7 6 5 4 3 2 1 0

$$-2^7 + 2^6 + 2^4 + 2^1 + 2^0$$

$$\Rightarrow -128 + 64 + 16 + 2 + 1$$

$$= -128 + 83$$

$$\Rightarrow -45$$

\oplus sum = sum / 2

carry = sum / 2

$$02. (-12)_{10} \rightarrow ()_2$$

$$12 = \underline{0} \underline{0} \underline{0} \underline{0} \underline{1} \underline{1} \underline{0} \underline{0}$$

i's compliment = $\underline{1} \underline{1} \underline{1} \underline{1} \underline{0} \underline{0} \underline{1} \underline{1}$

$$+ 1 \quad \underline{\underline{0} \underline{0} \underline{0} \underline{0} \underline{0} \underline{0} \underline{0} \underline{1}}$$

2's compliment = $\underline{\underline{1} \underline{1} \underline{1} \underline{1} \underline{0} \underline{1} \underline{0} \underline{0}}$

$$\text{Ans} = -2^7 + 2^6 + 2^5 + 2^4 + 2^2$$

$$= -128 + 64 + 32 + 16 + 4$$

$$= -128 + 116$$

$$\Rightarrow -12$$

* Ranges

01. Byte = 8 bits

$$\text{Min num} = 10000000 = -128 \\ = -2^7$$

$$\text{mask num} = 01111111 = 2^7 - 1$$

-128 to 127

02 int → 32 bits

$$\min \text{ num} = 1\ 00000\dots 000000 \rightarrow -2^{31}$$

31 30 29 - - - .. 4 3 2 1 0

$$\max \text{ num} = 011111\dots111111 = \overset{31}{2} - 1$$

$$\text{Range} = -2^{31} \text{ to } 2^{31} - 1$$

$$\text{Range} = -2 * 10^9 \text{ to } 2 * 10^9$$

03. long = 64 bits

$$\min \text{ num} = \frac{1}{6^3} \quad \underline{0} \quad \underline{0} \quad \underline{0} \quad \underline{0} \quad \dots \quad \underline{0} \quad \underline{0} \quad \underline{0} = -2^{63}$$

$$\max \text{ num} = \underline{0} \quad | \quad | \quad | \quad | \quad | \quad - \quad - \quad | \quad | \quad | \quad | \quad | = 2^{63} - 1$$

$$\text{Range} = -2^{63} \text{ to } 2^{63} - 1$$

Q Given two integers a & b. Return a * b

constraints $\begin{cases} a \leq 2 * 10^9 \\ b \leq 2 * 10^9 \end{cases}$

01 int ans = a * b \rightarrow incorrect

return ans overflow

02 long ans = a * b \rightarrow incorrect

return ans; Overflow

$\overbrace{\text{int} * \text{int}}$

int \rightarrow garbage value

03. long ans = long(a * b) \rightarrow incorrect

overflow

$\overbrace{\text{int} * \text{int}}$

int \rightarrow (garbage value)

04. $\text{long ans} = \text{long}(a) * b$ → Correct

$\text{long} \neq \text{int}$
long
=====

05 $\text{long ans} = a;$
 $\text{ans} *= b;$
 return ans

} Correct

_____ & _____ & _____ & _____ & _____

Doubts

Signed integer - MSB
0 positive
1 negative

unsigned integer -

⇒ Unique elements when ele are repeating

$$\boxed{a^a = 0}$$

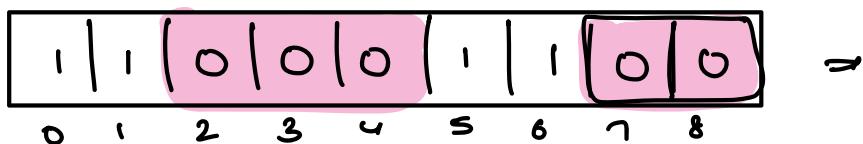
$$a^b^c = a^c^b \\ = a^b^c$$

= array = 2 3 4 5 3 5 2

⇒ $2^3^4^5^3^5^2 \rightarrow 4$

- $2^2^3^3^5^5^4 \rightarrow 4$

* Subarray with OR = 1



Subarray whose OR = 1

$$\text{Total no. of subarrays} = \frac{n(n+1)}{2} \Rightarrow \frac{9 \times 10}{2} = \underline{\underline{45}}$$

Obs → Even if single 1 is present, the OR of complete subarray is 1.

* count the no. of subarrays only having zeros

Ans = Total no. of subarrays - subarrays only having zero

count = 0

for (i=0; i<n; i++)

 if (arr[i] == 0)

 zeroes ++;

 else {

 count = count + (zeroes * ($\frac{zeroes+1}{2}$))

 zeroes = 0

}

3

count = count + (zeroes * ($\frac{zeroes+1}{2}$))

$$Ans = \frac{n * (n+1)}{2} - count$$