Big Data Analytics



Academic Year 2022-23

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MapReduce: Data Processing using Programming



- MapReduce is a framework used for writing application that processes large data sets in a distributed manner with parallels algorithms.
- It is core component of the Apache Hadoop ecosystem
- MapReduce have main two function
 - 1. Map
 - 2. Reduce
- Both these Map & Reduce works only on key-value pairs

```
Ex. (Roll_No, Name)
(1, Sohan)
(2, Mohan)
```

In both function, input is Key-value pair and output also Key, Value pair

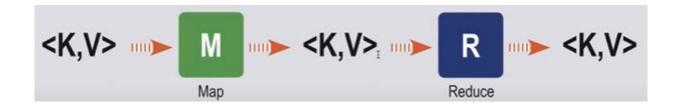
$$(k, v) -> Map -> (k, v) -> Reduce -> (k, v)$$

MapReduce, Why?



- MapReduce is a programming paradigm
- Traditional programming models work only when data is kept on a single machine and if data kept
 on multiple machine in a distributed manner than we require a new programming model to solve the
 problems.

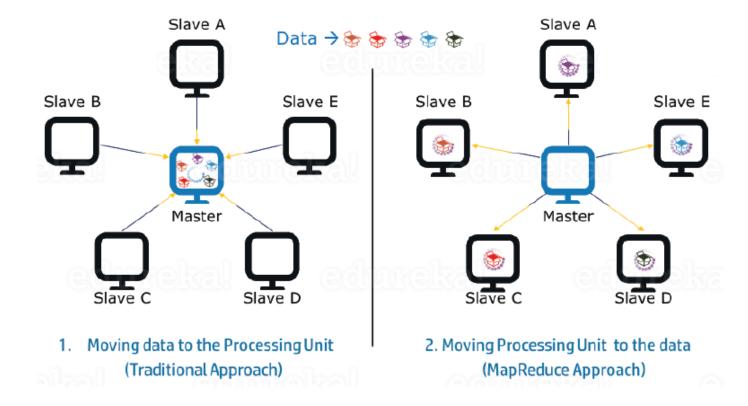
MapReduce is a computing paradigm for processing data that resides on many machines.



MapReduce: Working Mechanism

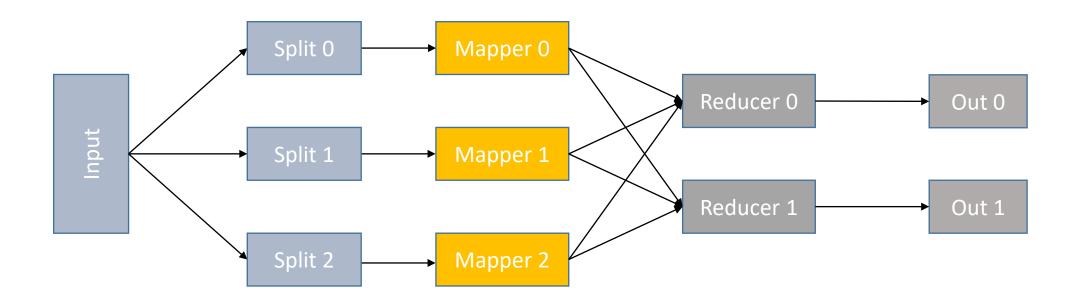


- Hadoop works on principle of data Locality, i.e., the data is processed where it is kept.
- In MapReduce code is move to data, data is not coming towards code, it makes fast processing.
- Data Block 128 MB but code around 20 KB.



MapReduce: Work Flow





MapReduce: Word Count





Car Bike Car Motor
Car Bike Motor
Car Bike Car
Bike Car
Bike Car

word_file.txt (300 MB)

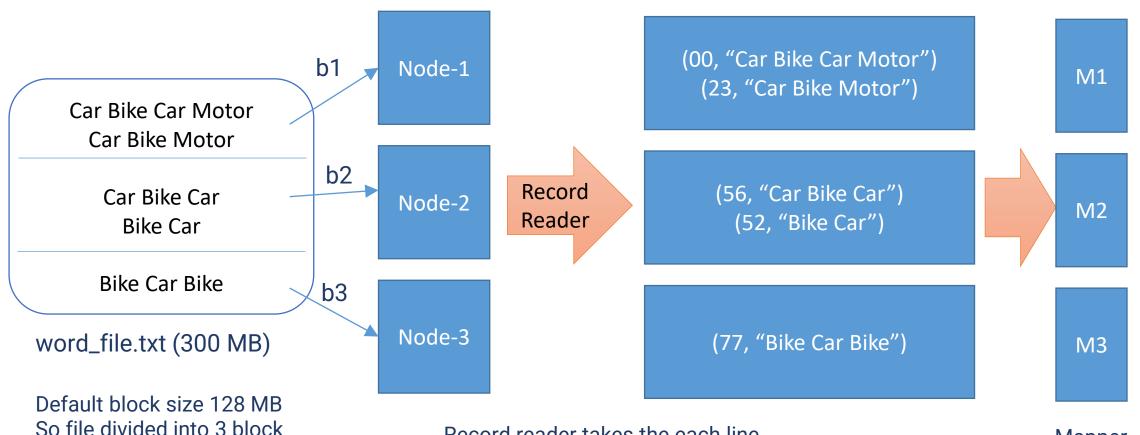
Output

'Car': 7, 'Bike': 6, 'Motor': 2

Expected output

MapReduce: Word Count





Record reader takes the each line as input, add dummy key and return key-value pairs as output.

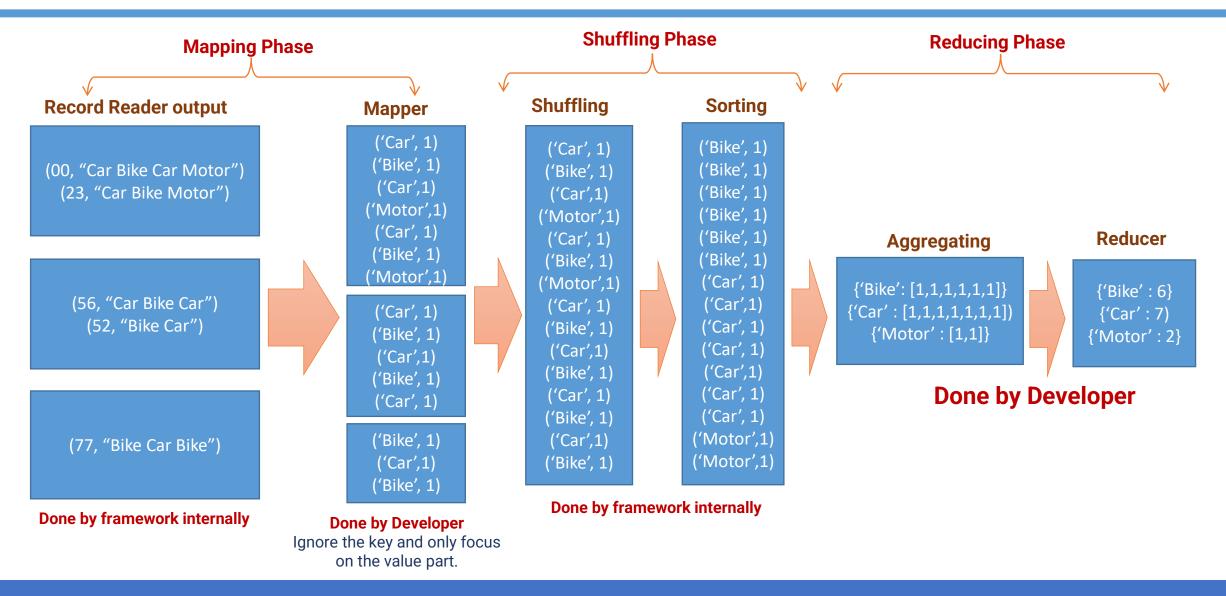
It is done by framework internally.

Mapper Program /logic has

to be written by the developer.

MapReduce: Word Count





MapReduce with Combiner: Work Flow

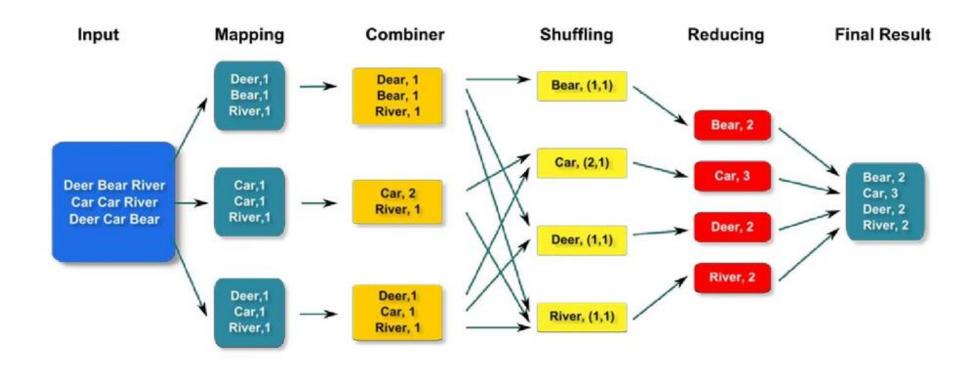


- Combiner is also known as "Mini-Reducer" that synopsizes the Mapper output record with the same
 Key before passing to the Reducer.
- On a huge dataset when we run a MapReduce job. So Mapper creates large chunks of intermediate
 data. Then the framework passes this intermediate data on the Reducer for further handling. This
 leads to huge network congestion. The Hadoop framework offers a function known as Combiner
 that plays a key role in reducing network congestion.
- The main job of Combiner a "Mini-Reducer is to handle the output data from the Mapper, before passing it to Reducer. It works after the mapper and before the Reducer. Its usage is optional.

MapReduce with Combiner: Work Flow



Combiner - Local Reduce



MapReduce with Combiner



Advantages

- Use of combiner decreases the time taken for data transfer between mapper and reducer.
- Combiner increases the overall performance of the reducer.
- It reduces the amount of data that the reducer has to process.

Disadvantages

- In the native filesystem, when Hadoop stores
 the key-value pairs and runs the combiner
 later this will result in expensive disk IO.
- MapReduce jobs can't rely on the combiner execution as there is no guarantee in its execution.

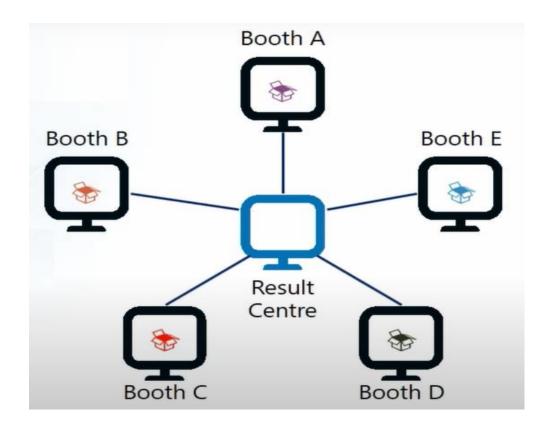
MapReduce Use Case: Election Vote Counting



Votes is stored at different Booths and Result centre has the details of all the booths

Counting: Traditional Approach

- Votes are moved to result centre for counting.
- Moving all the votes to centre is costly.
- Results centre is over burdened.
- Counting takes time.



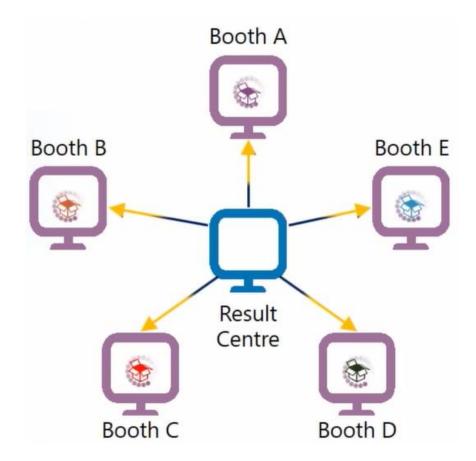
MapReduce Use Case: Election Vote Counting



Votes is stored at different Booths and Result centre has the details of all the booths

Counting: MapReduce Approach

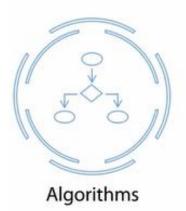
- Votes are counted at individual booths.
- Booth-wise results are sent to the result centre.
- Final result is declared easily and quickly using this way.



Real-Time Uses of MapReduce



Real-Time Uses of MapReduce











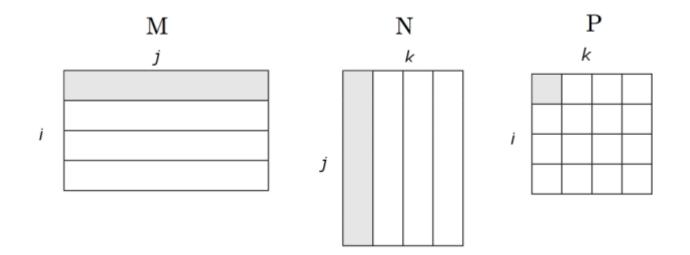




Algorithm uses MapReduce : Matrix Vector Multiplication



Matrix-vector and matrix-matrix calculations fit nicely into the MapReduce style of computing.



Matrix Data Model for MapReduce



Algorithm 1: The Map Function

- 1 for each element m_{ij} of M do
- produce (key, value) pairs as $((i, k), (M, j, m_{ij}))$, for k = 1, 2, 3, ... up to the number of columns of N
- 3 for each element n_{jk} of N do
- produce (key, value) pairs as $((i, k), (N, j, n_{jk}))$, for i = 1, 2, 3, ... up to the number of rows of M
- 5 return Set of (key, value) pairs that each key, (i,k), has a list with values (M, j, m_{ij}) and (N, j, n_{jk}) for all possible values of j

Algorithm 2: The Reduce Function

- 1 for each key (i,k) do
- sort values begin with M by j in $list_M$
- sort values begin with N by j in $list_N$
- 4 multiply m_{ij} and n_{jk} for j_{th} value of each list
- sum up $m_{ij} * n_{jk}$
- 6 return (i,k), $\sum_{j=1} m_{ij} * n_{jk}$



Matrix Data Model for MapReduce

$$M_{ij} = \begin{bmatrix} 10 & 20 \\ 30 & 40 \end{bmatrix} \qquad N_{jk} = \begin{bmatrix} 5 & 6 \\ 7 & 8 \end{bmatrix}$$

As input data files, we store matrix M and N on HDFS in following format

$$(M, i, j, m_{ij})$$
 and (N, i, j, m_{ij})
 $(M, 1, 1, 10)$, $(M, 1, 2, 20)$, $(M, 2, 1, 30)$
 $(N, 1, 1, 5)$, $(N, 1, 2, 6)$, $(N, 2, 1, 7)$

Note:

Most matrices are sparse so large amount of cells have value zero. When we represent matrices in this form, we do not need to keep entries for the cells that have values of zero to save large amount of disk space.



for each element m_{ij} of M do produce (key, value) pairs as $\left((i,k), \left(M,j,M_{ij}\right)\right)$ for k=1,2, ... Up to column of matrix N

Mapper code for Matrix M

k=#colums of N	i = row number of M	j = column number of M	formula (key, value) $\Big((i,k), ig(M,j,m_{ij}ig)\Big)$
k=1	i=1	j=1	((1,1), (M, 1, 10))
		j=2	((1,1), (M, 2, 20))
	i=2	j=1	((2,1), (M, 1, 30))
		j=2	((2,1), (M, 2, 40))
k=2	i=1	j=1	((1,2), (M, 1, 10))
		j=2	((1,2), (M, 2, 20))
	i=2	j=1	((2,2), (M, 1, 30))
		j=2	((2,2), (M, 2, 40))

for each element n_{jk} of N do produce (key, value) pairs $\operatorname{as} \left((i,k), \left(N,j,N_{jk} \right) \right)$ for i=1,2, .. Up to row of matrix M

Mapper code for Matrix N

i=#colums of M	j = row number of N	k = column number of N	formula (key, value) $\left((i,k), \left(N, j, n_{ij}\right)\right)$
i=1	j=1	k=1	((1,1), (N, 1, 5))
		k=2	((1,2), (N, 1, 6))
	j=2	k=1	((1,1), (N, 2, 7))
		k=2	((1,2), (N, 2, 8))
i=2	j=1	k=1	((2,1), (N, 1, 5))
		k=2	((2,2), (N, 1, 6))
	j=2	k=1	((2,1), (N, 2, 7))
		k=2	((2,2), (N, 2, 8))



Return Set of (key, value) pairs that same key has a list with values M & N



Algorithm 2: The Reduce Function

```
1 for each key (i,k) do

2 sort values begin with M by j in list_M

3 sort values begin with N by j in list_N

4 multiply m_{ij} and n_{jk} for j_{th} value of each list

5 sum up m_{ij} * n_{jk}

6 return (i,k), \sum_{j=1}^{n} m_{ij} * n_{jk}
```

Reduce Code:

Sort values in M and N list according to J and perform multiplication.

$$((1,1), 10*5+20*7) => ((1,1), 190)$$

 $((1,2), 10*6+20*8) => ((1,2), 220)$
 $((2,1), 30*5+40*7) => ((2,1), 430)$
 $((2,2), 30*6+40*8) => ((2,2), 500)$

Result:

$$M_{ij} = \begin{bmatrix} 10 & 20 \\ 30 & 40 \end{bmatrix}$$
 $N_{jk} = \begin{bmatrix} 5 & 6 \\ 7 & 8 \end{bmatrix}$ $M \times N = \begin{bmatrix} 190 & 220 \\ 430 & 500 \end{bmatrix}$

Failure in MapReduce



Failures are norm in commodity hardware

Worker failure

- Detect failure via periodic heartbeats
- Re-execute in-progress map/reduce tasks

• Master failure

Single point of failure; Resume from Execution Log

Robust

 Google's experience: lost 1600 of 1800 machines once!, but finished fine



Learn Fundamentals & Enjoy Engineering





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