

STACK & APPLICATIONS

Unit-II: Stacks and Queues

Stacks: The Stack: Definition, operations, implementation using arrays, linked list and **Stack applications:** Infix to postfix expression conversion, Evaluation of Postfix expressions, balancing the symbols. **Queue:** definition, operations, implementation using arrays, linked list & its Applications. **Circular queue:** definition & its operations, implementation, **Dequeue:** definition & its types, implementation.

Learning Material

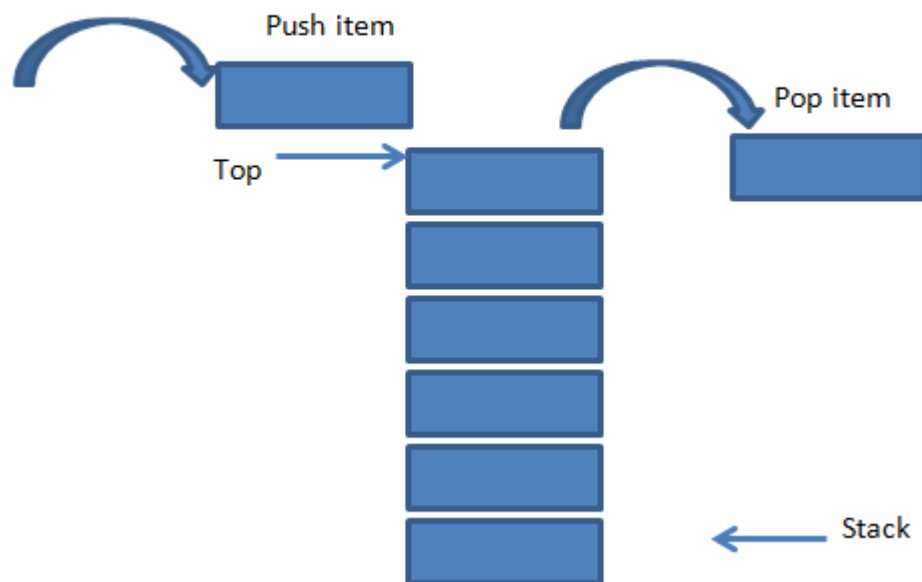
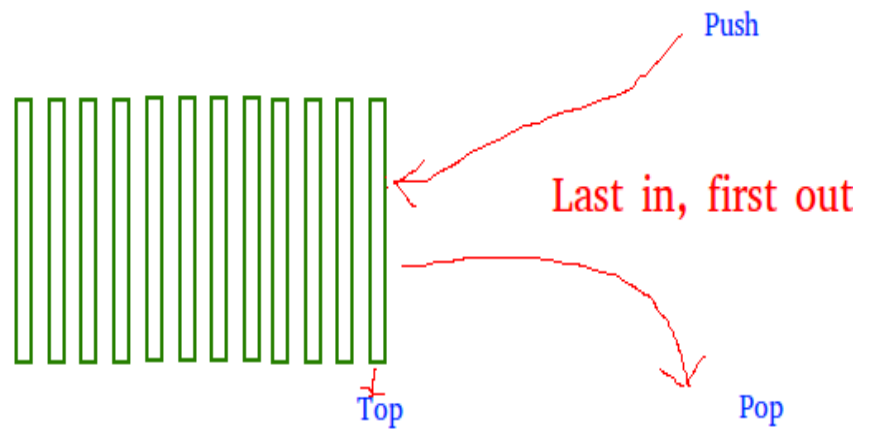
- ☐ Stack is a linear data structure.
- ☐ Stack can be defined as a collection of homogeneous elements, where insertion and deletion operations takes place at only one end called **TOP**.
- ☐ The insertion operation is termed as PUSH and deletion operation is termed as POP operation.
- ☐ The PUSH and POP operations are performed at TOP of the stack.
- ☐ An element in a stack is termed as ITEM.
- ☐ The maximum number of elements that stack can accommodate is termed as SIZE of the stack.
- ☐ Stack Pointer (SP) always points to the top element of the stack.
- ☐ Stack follows LIFO principle. i.e. **Last In First Out** i.e. the element which is inserted last into the stack will be deleted first from the stack.

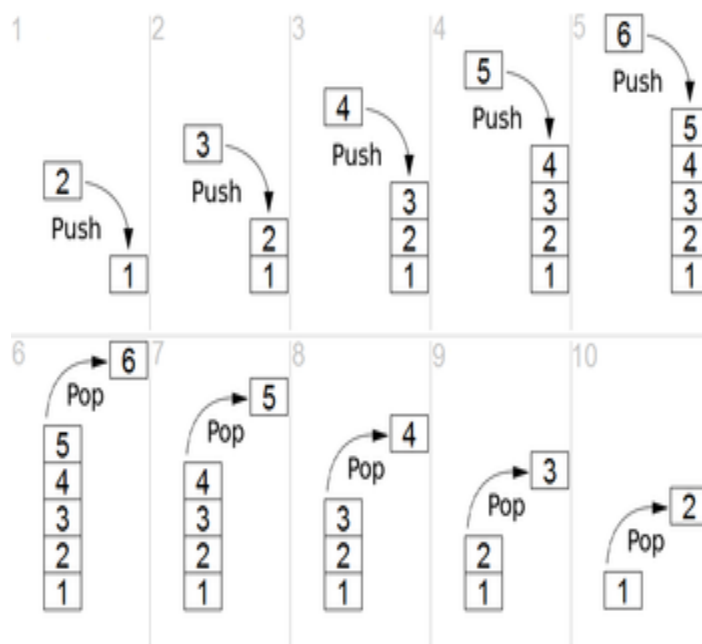
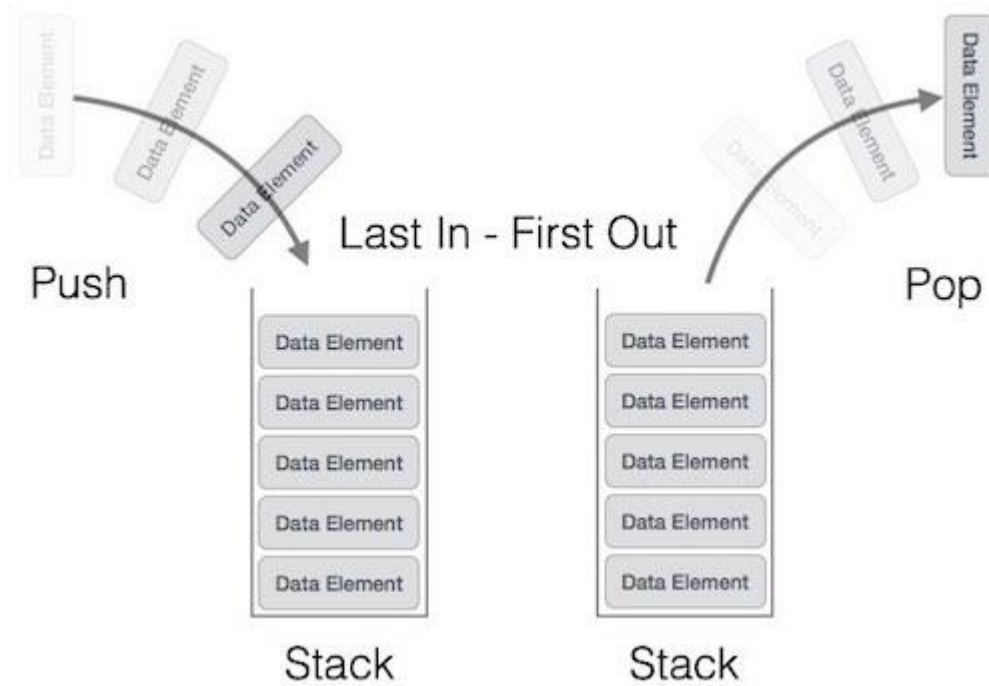
Diagram of a stack



Stack

Insertion and Deletion
happen on same end





Stack Operations:- We can perform mainly 3 operations on stack.

1. Push :- Inserting an element into the stack
2. Pop:- Deleting an element from the stack
3. Peep:- Displaying top most element of the stack

Representation of stack

There are two ways of representation of a stack.

1. Array representation of a stack.
2. Linked List representation of a stack.

1. Array representation of a stack.

First we have to allocate memory for array.

Starting from the first location of the memory block, items of the stack can be stored in sequential fashion.

Stack Initial Condition:

Initial condition of stack implemented using arrays is **top=-1**

Stack Full condition or overflow condition

Trying to PUSH an item into full stack is known as Stack overflow.

Stack overflow condition is $\text{top} == \text{ARRAYSIZE} - 1$

Stack underflow or Empty condition

Trying to POP an item from empty stack is known as Stack underflow. Stack underflow condition is $\text{top} == -1$

Algorithm Stack_PUSH(item)

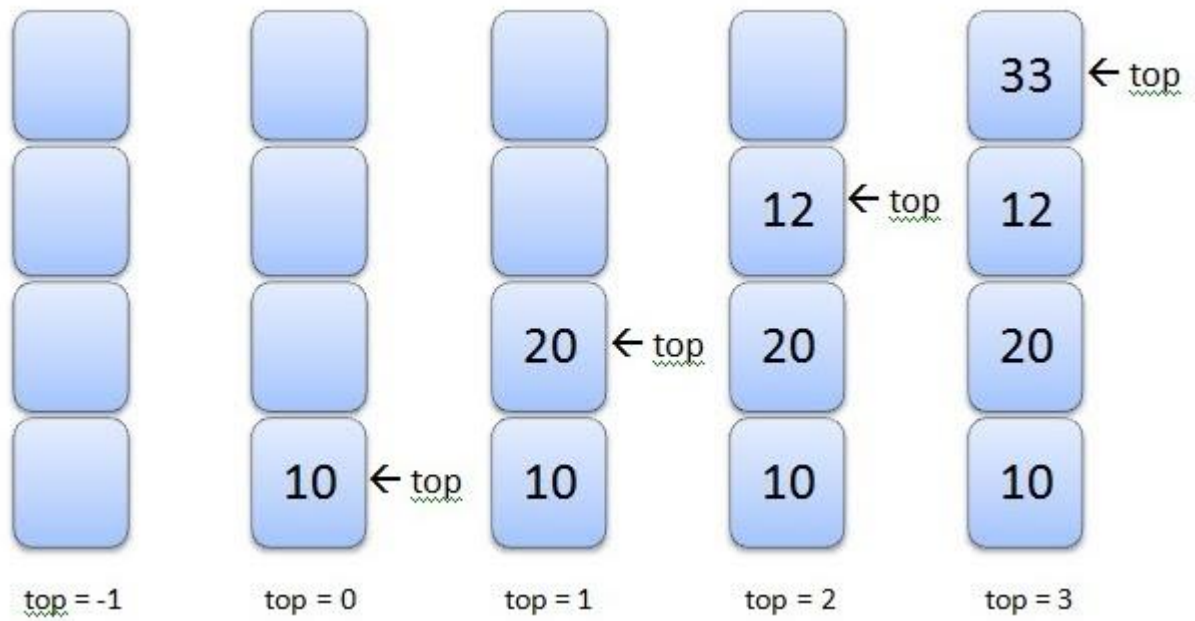
Input: item is new item to push into stack

Output: pushing new item into stack at top whenever stack is not full.

1. if($\text{top} == \text{ARRAYSIZE}-1$)
 - a) print(stack is full, not possible to perform push operation)
2. else
 - a) $\text{top}=\text{top}+1$
 - b) read item
 - c) $\text{s}[\text{top}]=\text{item}$

End

Stack_PUSH

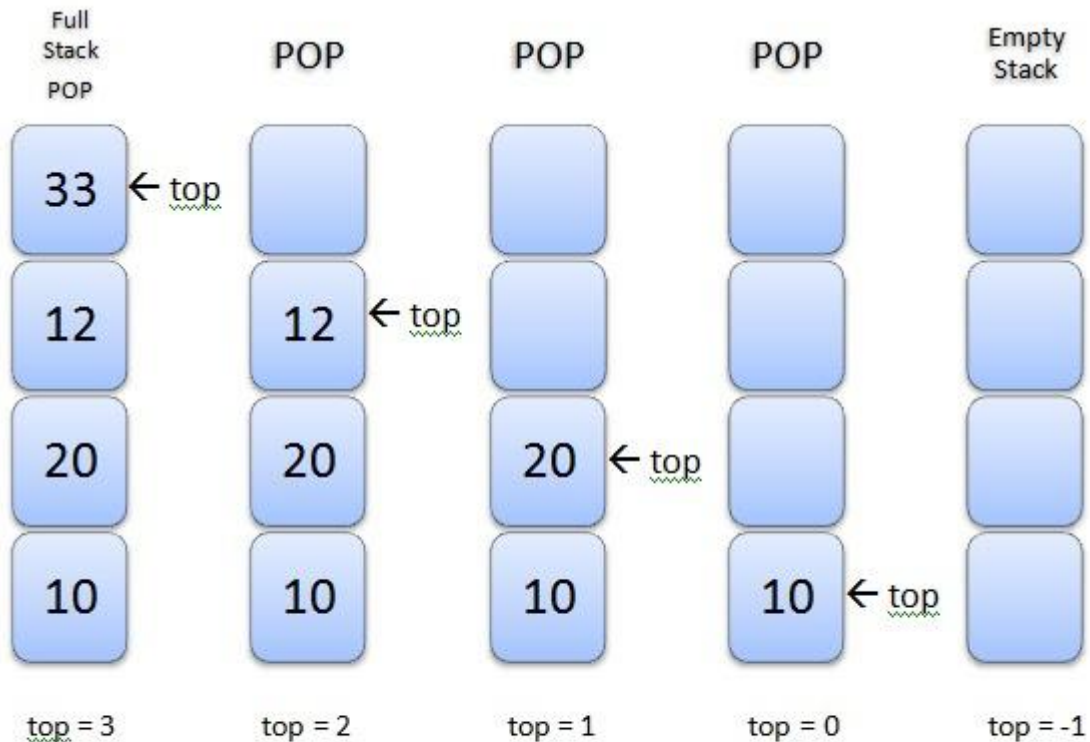


Algorithm Stack_POP()

Input: Stack with some elements.

Output: item deleted at top most end.

1. if($\text{top} == -1$)
 - a) print(stack is empty not possible to pop)
2. else
 - a) $\text{item} = \text{s}[\text{top}]$
 - b) $\text{top} = \text{top} - 1$
 - c) print(deleted item)



POP operation on Stack with the position of top pointer
 Each time an element is popped, then $top = top - 1$.

End Stack_POP

C Program to implement stack using arrays

```
#include<stdio.h>
#include<conio.h>
void push(int ele );
void pop( );
void peep( );
void display( );
int stack[10], top= -1, ele;
void main()
{
    int choi, ele;
    clrscr();
    while(1)
    {
        printf("\t\t Stack ADT Using Arrays\n");
        printf("\n1.Push \n2.Pop \n3.Peep \n4.Display \n5.Exit");
        printf("\nEnter your choice");
        scanf("%d",&choi);
        switch(choi)
        {
            case 1: printf("\nEnter the element to be inserted into the stack");
                     scanf("%d",&ele);
                     push(ele );
                     break;
            case 2: pop();
                     break;
            case 3: peep();
                     break;
            case 4: display( );
                     break;
            case 5: exit(0);
        }
    }
    getch();
}
```

```

}

void push( int ele)
{
    if( top==9)
        printf("\nStack is full");
    else
    {
        top++;
        stack[top]=ele;
    }
}

void pop( )
{
    if( top== -1)
        printf("\nStack is empty");
    else
    {
        ele=stack[top];
        printf("\nThe deleted element from the stack is %d",ele);
        top--;
    }
}

void peep( )
{
    if( top == -1)
        printf("\nStack is empty");
    else
    {
        ele=stack[top];
        printf("\nThe top most element of the stack is %d",ele);
    }
}

void display( )
{
    int i;
    if( top== -1)
        printf("\nStack is empty");
    else
    {
        printf("\nThe elements of the stack are:\n");
        for(i=top;i>=0;i--)
            printf("%d\n",stack[i]);
    }
}

```

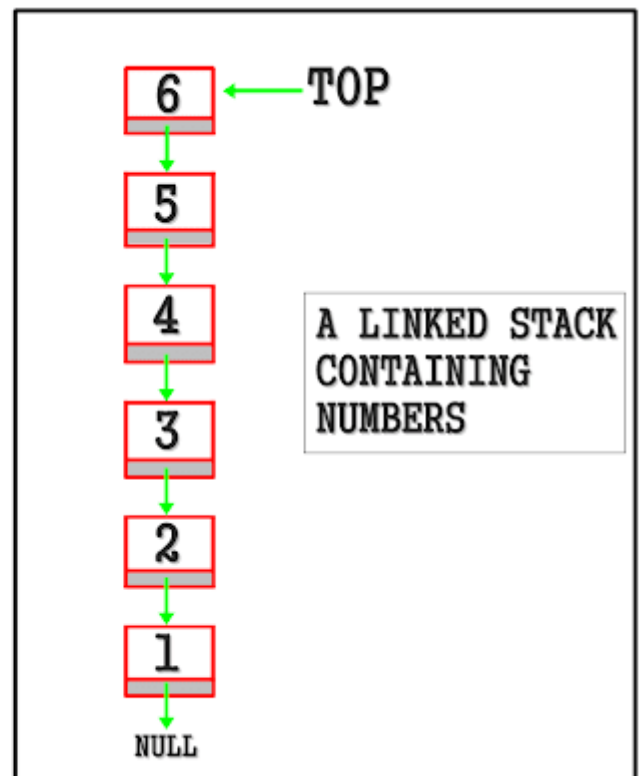
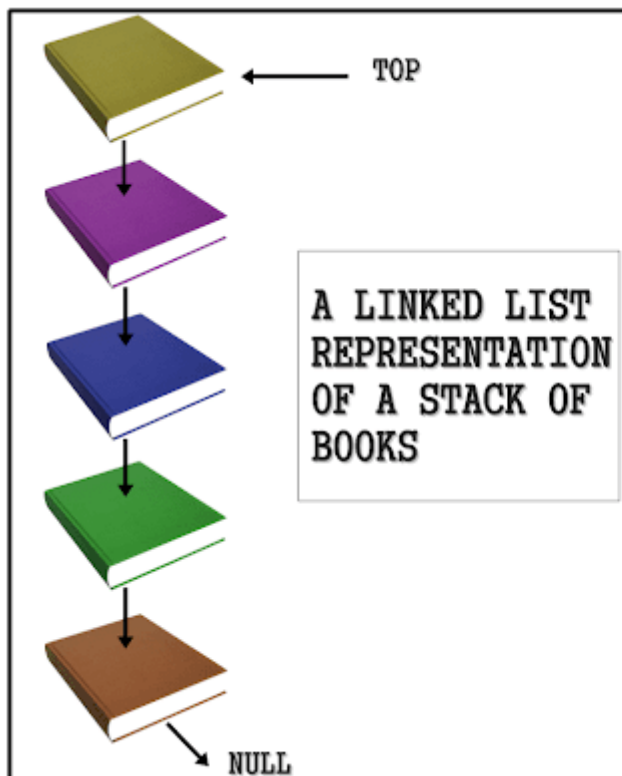
2. Linked List representation of a stack

- The array representation of stack allows only fixed size of stack. i.e. static memory allocation only.
- To overcome the static memory allocation problem, linked list representation of stack is preferred.
- In linked list representation of stack, each node has two parts. One is data field is for

the item and link field points to next node.

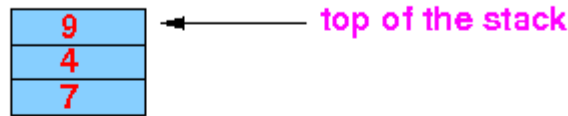
- Empty stack condition is
 $\text{top} = \text{NULL}$
- Full condition is not applicable for Linked List representation of stack. Because here memory is dynamically allocated.
- In linked List representation of stack, top pointer always points to top most node only.
i.e. first node in the list.

Linked list representation of stack:-

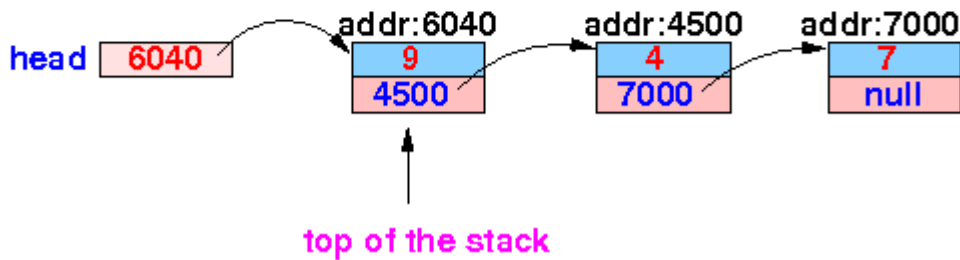


Representing a stack with a list:

Stack:



List:



C program to implement stacks using linked list:-

```
#include<stdio.h>
#include<stdlib.h>
struct node
{
    int info;
    struct node *next;
};
struct node *top=NULL;
void push(int ele);
void pop();
void display();
void main()
{
    int ele,choice;
    while(1)
    {
        clrscr();
        printf("\t\t Stack ADT using Linked List \n");
        printf("1.push\n");
        printf("2.pop\n");
        printf("3.display\n");
        printf("4.exit\n");
        printf("enter choice\n");
        scanf("%d",&choice);
        switch(choice)
        {
            case 1:printf("\n enetr ele ");
                    scanf("%d",&ele);
                    push(ele);
```

```

                break;
            case 2:pop();
                break;
            case 3:display();
                break;
            case 4:exit(0);
        }
    }
}

void push(int ele)
{
    struct node *temp;
    temp=(struct node*)malloc(sizeof(struct node));
    temp->info=ele;
    if(top==NULL)
    {
        temp->next=NULL;
        top=temp;
    }
    else
    {
        temp->next=top;
        top=temp;
    }
}

void pop()
{
    struct node *temp;
    if(top==NULL)
        printf("\n stack is empty");
    else if(top->next==NULL)
    {
        temp=top;
        top=NULL;
        free(temp);
    }
    else
    {
        temp=top;
        top=top->next;
        free(temp);
    }
    getch();
}

void display()
{
    struct node *p;

```

```

    if(top==NULL)
    printf("\n stack is empty");
    else
    {
    p=top;
    while(p!=NULL)
    {
        printf("\n elements are %d",p->info);
        p=p->next;
    }
    }
    getch();
}

```

Arithmetic expressions or Algebraic expressions:-

An expression is a combination of operands and operators.

Eg. $c = a + b$

In the above expression a, b, c are operands and +, = are called as operators.

We have 3 notations for the expressions.

- i. Infix notation
- ii. Prefix notation
- iii. Postfix notation

Infix notation: Here operator is present between two

operands. eg. $a + b$

The format for Infix notation as follows

$\langle \text{operand} \rangle \langle \text{operator} \rangle \langle \text{operand} \rangle$

Prefix notation: Here operator is present before two operands.

eg. $+ a b$

The format for Prefix notation as follows

$\langle \text{operator} \rangle \langle \text{operand} \rangle \langle \text{operand} \rangle$

Postfix notation: Here operator is present after two operands.

eg. a b +

The format for Postfix notation as follows

<operand> <operand> <operator>

Applications of stack

1. Infix to postfix conversion
2. Evaluation of postfix expression
3. Balancing symbols or delimiter matching

1. Infix to postfix conversion

While conversion of infix expression to postfix expression, we must follow the precedence (priority) of the operators.

Operator	priority
(0
+ -	1
* / %	2
^ or \$	3

To convert an infix expression to postfix expression, we can use one stack.

Within the stack, we place only operators and left parenthesis only. So stack used in conversion of infix expression to postfix expression is called as operator stack.

Algorithm Conversion of infix to postfix

Input: Infix expression.

Output: Postfix expression.

1. Perform the following steps while reading of infix expression is not over

- a) if symbol is left parenthesis then push symbol into stack.
- b) if symbol is operand then add symbol to post fix expression.
- c) if symbol is operator then check stack is empty or not.
 - i) if stack is empty then push the operator into stack.
 - ii) if stack is not empty then check priority of the operators.

(I) if priority of current operator > priority of operator present at top of stack then push operator into stack.

(II) else if priority of operator present at top of stack >= priority of current operator then pop the operator present at top of stack and add popped operator to postfix expression (go to step I)

d) if symbol is right parenthesis then pop every element form stack up corresponding left parenthesis and add the popped elements to postfix expression.

2. After completion of reading infix expression, if stack not empty then pop all the items from stack and then add to post fix expression.

End conversion of infix to postfix

Contin.....

Conversion of Infix to Postfix

Example to Convert Infix to Postfix using stack

$$a + (b * c)$$

Read character	Stack	Output
a	Empty	a
+	+	a
(+(a
b	+(ab
*	+(*	ab
c	+(*	abc
)	+	abc*
	9CM305.25	abc*+ 19

Infix to Postfix using stack



- Example $A*B+C$ become $AB*C+$

	current symbol	operator stack	postfix string
1	A		A
2	*	*	A
3	B	*	AB
4	+	+	AB * {pop and print the '*' before pushing the '+'}
5	C	+	AB * C
6			AB * C +

The infix expression is

$$(P/(Q-R)*S+T)$$

Symbol	Stack	Expression
((—
P	(P
/	(/	P
((/(P
Q	(/(PQ
-	(/(-	PQ
R	(/(-	PQR
)	(/	PQR-
*	(*	PQR-/
S	(*	PQR-/S
+	(+	PQR-/S*
T	(+	PQR-/S*T
)		PQR-/ST*+

So, the postfix expression is $PQR-/ST+*$.

Infix Expression : A+B*(C^D-E)				
Token	Action	Result	Stack	Notes
A	Add A to the result	A		
+	Push + to stack	A	+	
B	Add B to the result	AB	+	
*	Push * to stack	AB	* +	* has higher precedence than +
(Push (to stack	AB	(* +	
C	Add C to the result	ABC	(* +	
^	Push ^ to stack	ABC	^ (* +	
D	Add D to the result	ABCD	^ (* +	
-	Pop ^ from stack and add to result	ABCD^	(* +	- has lower precedence than ^
	Push - to stack	ABCD^	- (* +	
E	Add E to the result	ABCD^E	- (* +	
)	Pop - from stack and add to result	ABCD^E-	(* +	Do process until (is popped from stack
	Pop (from stack	ABCD^E-	* +	
	Pop * from stack and add to result	ABCD^E-*	+	Given expression is iterated, do Process till stack is not Empty, It will give the final result
	Pop + from stack and add to result	ABCD^E-*+		
Postfix Expression : ABCD^E-*+				

	RPN	Stack	Input Expression		RPN	Stack	Input Expression
①			$A+(B*(C-D)/E)$	⑨	ABC	(* (+	$D)/E)$
②	A		$+(B*(C-D)/E)$	⑩	ABCD	(* (+	$) / E)$
③	A	+	$(B*(C-D)/E)$	⑪	ABCD-	* (+	$/ E)$
④	A	(+	$B*(C-D)/E)$	⑫	ABCD-*	/ (+	$E)$
⑤	AB	(+	$*(C-D)/E)$	⑬	ABCD-*E	/ (+	$)$
⑥	AB	* (+	$(C-D)/E)$	⑭	ABCD-*E/	+	
⑦	AB	(* (+	$C-D)/E)$	⑮	ABCD-*E/+		
⑧	ABC	(* (+	$-D)/E)$				

Consider the following arithmetic infix expression P

$$P = A + (B / C - (D * E ^ F) + G) * H$$

Character scanned	Stack	Postfix Expression (Q)
A	(A
+	(+	A
((+ (A
B	(+ (AB
/	(+ (/	AB
C	(+ (/	ABC
-	(+ (-	ABC /
((+ (- (ABC /
D	(+ (- (ABC / D
*	(+ (- (*	ABC / D
E	(+ (- (*	ABC / DE
^	(+ (- (* ^	ABC / DE
F	(+ (- (* ^	ABC / DEF
)	(+ (-	ABC / DEF ^ *
+	(+ (+	ABC / DEF ^ * -
G	(+ (+	ABC / DEF ^ * - G
)	(+	ABC / DEF ^ * - G +
*	(+ *	ABC / DEF ^ * - G +
H	(+ *	ABC / DEF ^ * - G + H
)		ABC / DEF ^ * - G + H * +

Write a Turbo C program to convert the Infix expression to Postfix Expression using Stacks implementation of Arrays.

EXAMPLE 10.6. Convert $X: A + (B * C - (D / E \uparrow F) * G) * H$ into postfix form showing stack status after every step in tabular form.

Solution.

Symbol Scanned	Stack	Expression Y
1. A	(A
2. +	(+	A
3. ((+ (A
4. B	(+ (A B
5. *	(+ (*	A B
6. C	(+ (*	A B C
7. -	(+ (-	A B C *
8. ((+ (- (A B C *
9. D	(+ (- (A B C * D
10. /	(+ (- (/	A B C * D
11. E	(+ (- (/	A B C * D E
12. ↑	(+ (- (/ ↑	A B C * D E
13. F	(+ (- (/ ↑	A B C * D E F
14.)	(+ (-	A B C * D E F ↑ /
15. *	(+ (- *	A B C * D E F ↑ /
16. G	(+ (- *	A B C * D E F ↑ / G
17.)	(+	A B C * D E F ↑ / G * -
18. *	(+ *	A B C * D E F ↑ / G * -
19. H	(+ *	A B C * D E F ↑ / G * - H
20.)		A B C * D E F ↑ / G * - H * +

Infix Expression: **A+ (B*C-(D/E^F)*G)*H**, where ^ is an exponential operator

Symbol	Scanned	STACK	Postfix Expression	Description
1.		(Start
2.	A	(A	
3.	+	(+	A	
4.	((+(A	
5.	B	(+(AB	
6.	*	(+(*	AB	
7.	C	(+(*	ABC	
8.	-	(+(-	ABC*	'*' is at higher precedence than '-'
9.	((+(-(ABC*	
10.	D	(+(-(ABC*D	
11.	/	(+(-(/	ABC*D	
12.	E	(+(-(/	ABC*DE	
13.	^	(+(-(/^	ABC*DE	
14.	F	(+(-(/^	ABC*DEF	
15.)	(+(-	ABC*DEF^/	Pop from top on Stack , that's why '^' Come first
16.	*	(+(-*	ABC*DEF^/	
17.	G	(+(-*	ABC*DEF^/G	
18.)	(+	ABC*DEF^/G*-	Pop from top on Stack , that's why '^' Come first
19.	*	(+*	ABC*DEF^/G*-	
20.	H	(+*	ABC*DEF^/G*-H	
21.)	Empty	ABC*DEF^/G*-H*+	END

C Program to convert infix to postfix expression:-

```
#define SIZE 50          /* Size of Stack */
#include <ctype.h>
char s[SIZE];
int top=-1;             /* Global declarations */
void push(char elem)
{
    /* Function for PUSH operation */
    s[++top]=elem;
}

char pop()
{
    /* Function for POP operation */
```

```

        return(s[top--]);
    }

int pr(char elem)
{
    /* Function for precedence of operators*/
    switch(elem)
    {
        case '#': return 0;
        case '(': return 1;
        case '+':
        case '-': return 2;
        case '*':
        case '/': return 3;
    }
}

```

```

void main()
{
    /* Main Program */
    char infx[50],pofx[50],ch,elem;
    int i=0,k=0;
    printf("\n\nRead the Infix Expression ? ");
    scanf("%s",infx);
    push('#');
    while( (ch=infx[i++]) != '\0')
    {
        if( ch == '(')
            push(ch);
        else if(isalnum(ch))
            pofx[k++]=ch;
        else if( ch == ')')
        {
            while( s[top] != '(')
                pofx[k++]=pop();
            elem=pop(); /* Remove ( */
        }
        else
        {
            /* Operator */
            while( pr(s[top]) >= pr(ch) )
                pofx[k++]=pop();
        }
    }
}

```

```

        push(ch);
    }
} //end of while
while( s[top] != '#') /* Pop from stack till empty */
    pofx[k++]=pop();
pofx[k]='\0'; /* Make pofx as valid string */
printf("\n\nGiven Infix Expn: %s Postfix Expn:
%s\n",infx,pofx);
}

```

2.Evaluation of postfix expression

- To evaluate a postfix expression we use one stack.
- For Evaluation of postfix expression, in the stack we can store only operand. So stack used in Evaluation of postfix expression is called as operand stack.

Algorithm PostfixExpressionEvaluation

Input: Postfix expression

Output: Result of Expression

1. Repeat the following steps while reading the postfix expression.
 - a) if the read symbol is operand, then push the symbol into stack.
 - b) if the read symbol is operator then pop the top most two items of the stack and apply the operator on them, and then push back the result to the stack.
2. Finally stack has only one item, after completion of reading the postfix expression. That item is the result of expression.

End PostfixExpressionEvaluation

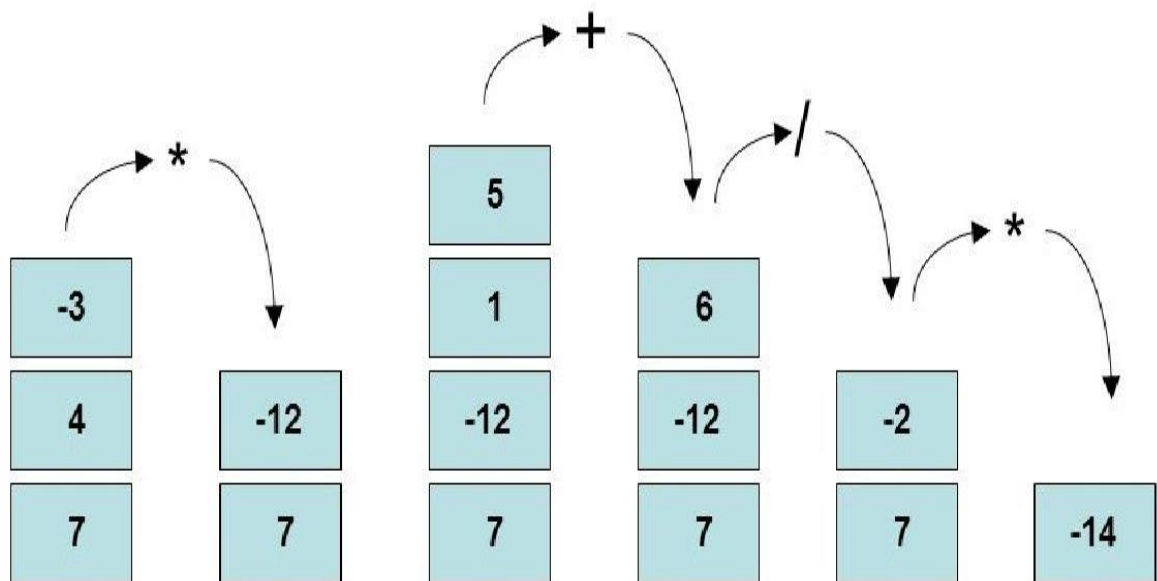
Example: $4325^*-+=4+3-2*5=-3$

Symbol	opnd1	opnd2	value	opndstack
4				4
3				4, 3
2				4, 3, 2
5				4, 3, 2, 5
*	2	5	10	4, 3
				4, 3, 10
-	3	10	-7	4
				4, -7
+	4	-7	-3	
				-3

result

Evaluating Postfix Expressions

- Expression = 7 4 -3 * 1 5 + / *



Evaluation of Postfix Expression

Postfix \rightarrow $a\ b\ c\ * +\ d\ -$ Let, $a = 4, b = 3, c = 2, d = 5$

Postfix \rightarrow $4\ 3\ 2\ * + 5\ -$

Operator/Operand	Action	Stack
4	Push	4
3	Push	4, 3
2	Push	4, 3, 2
*	Pop (2, 3) and $3*2 = 6$ then Push 6	4, 6
+	Pop (6, 4) and $4+6 = 10$ then Push 10	10
5	Push	10, 5
-	Pop (5, 10) and $10-5 = 5$ then Push 5	5



Evaluating Arithmetic Postfix Expression



PostFix Expression 2 3 4 + * 5 *

Move	Token	Stack
1	2	2
2	3	2 3
3	4	2 3 4
4	+	2 7 (3+4=7)
5	*	1 4 (2*7=14)
6	5	14 5
7	*	70 (14*5=70)

Evaluate below postfix expression 234+*6-

Step No.	Value of i	Operation	Stack
1	2	Push 2 in stack	2
2	3	Push 3 in stack	3 2
3	4	Push 4 in stack	4 3 2
4	+	Pop 2 elements from stack and perform addition operation. And push result back to stack. i.e. 4+3 = 7	7 2
5	*	Pop 2 elements from stack and perform multiplication operation. And push result back to stack. i.e. 7 * 2 = 14	14
6	6	Push 6 in stack	6 14
7	-	Pop 2 elements from stack and perform subtraction operation. And push result back to stack. i.e. 14 - 6 = 8	8
8		Pop result from stack and display	

C Program to evaluate postfix expression:-

```
#include<stdio.h>
#include<conio.h>
#include<ctype.h>
int st[100],top=-1;
int cal(char post[]);
void push_item(int);
int pop_item();
void main()
{
    char in[50];
    int result;
    clrscr();
    printf("\n \t enter the postfix expression");
    gets(in);
    result=cal(in);
    printf("\n result=%d",result);
    getch();
}
void push_item(int it)
{
    if(top==99)
        printf("stack is overflow\n");
    else
    {
        top++;
        st[top]=it;
    }
}
int pop_item()
{
    int it;
    if(top== -1)
        printf("stack underflow\n");
    else
    {
        return (st[top--]);
    }
}
int cal(char post[])
{
    int m,n,x,y,j=0,len;
    len=strlen(post);
    while(j<len)
    {
        if(isdigit(post[j]))
        {
            x=post[j]-'0';
            push_item(x);
        }
        else
        {
            m=pop_item();
            n=pop_item();
            switch(post[j])
            {
                case '+':x=m+n;
                    break;
                case '-':x=m-n;
                    break;
                case '*':x=m*n;
                    break;
                case '/':x=m/n;
                    break;
            }
            push_item(x);
        }
    }
}
```



```

        j++;
    }
    if(top>0)
    {
        printf("no of operands are more than operators");
        exit(0);
    }
    else
    {
        y=pop_item();
        return(y);
    }
}

```

3.Balancing Symbols or Delimiter matching :-

The objective of this application is to check the symbols such as parenthesis () , braces { } , brackets [] are matched or not.

Thus every left parenthesis, brace and bracket must have its right counterpart.

Algorithm for Balancing Symbols or Delimiter matching :-

1. Make an empty stack.
2. Read an expression from left to right.
3. If the reading character is opening symbol, then push it into stack.
4. If the reading character is closing symbol and if the stack is empty, then report as unbalanced expression.
5. If the reading character is closing symbol and if the stack is not empty, then pop the stack.
6. If the symbol popped is not the corresponding opening symbol, then report as unbalanced expression.
7. After processing the entire expression and if the stack is not empty then report as unbalanced expression.
8. After processing the entire expression and if the stack is empty then report as balanced expression.

C program to implement balancing symbols or delimiter matching:-

```

#include<stdio.h>
#include<conio.h>
#include<ctype.h>
void push(char);
char pop();
char stack[20];
int top=-1;
main()
{
    char expr[20],ch;

```

```

int i;
clrscr();
printf("\nEnter an expression\n");
gets(expr);
for(i=0;expr[i]!='\0';i++)
{
    ch=expr[i];
    if(ch=='(' || ch=='{' || ch=='[')
        push(ch);
    else if(ch==')')
    {
        if(top== -1)
        {
            printf("\nUnbalanced expression");
            exit();
        }
        else if( (ch=pop())!='(')
        {
            printf("\nUnbalanced expression");
            exit();
        }
    }
    else if(ch=='}')
    {
        if(top== -1)
        {
            printf("\nUnbalanced expression");
            exit();
        }
        else if( (ch=pop())!='{')
        {
            printf("\nUnbalanced expression");
            exit();
        }
    }
    else if(ch==']')
    {
        if(top== -1)
        {
            printf("\nUnbalanced expression");
            exit();
        }
        else if( (ch=pop())!='[')
        {
            printf("\nUnbalanced expression");
            exit();
        }
    }
}
if(top== -1)
    printf("\nBalanced expression");
getch();
}

```

```
void push(char x)
{
    top++;
    stack[top]=x;
}
int pop()
{
    return stack[top--];
}
```