TROPICAL SUPPORT VECTOR MACHINES AND MEAN PAYOFF GAMES

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ABSTRACT. We introduce a pseudo-polynomial algorithm for tropical support vector machines (SVMs), using mean payoff games for binary and multiclass classification. We provide a simple numerical criterion for tropical separability. In the separable case, we construct a separating tropical hyperplane, with margin guarantees. In the non-separable case, we use the same algorithm to provide a suitable classifier; it also measures the distance to tropical separability and allows to geometrically transform the data to reduce this distance to zero. We then extend our algorithm to the multiclass case. Finally, we study a feature map method to make the data separable in higher dimensions and propose a heuristic for sparsely adding relevant dimensions to better separate the data.

1. Introduction

Motivation and context. Support Vector Machines (SVMs), introduced by Boser et al. (1992) [4], are used in data science to determine separating hyperplanes for binary classification tasks.

Tropical geometry recently found applications in machine learning: tropical linear regression (Akian et al., 2021) [2], analysis of deep neural networks and parametric statistical models (Maragos et al., 2021) [8] and principal component analysis (Page et al., 2019) [9].

First considered by Gärtner and Jaggi (2006) [5], *Tropical Support Vector Machines* classify data points with a tropical hyperplane and using the tropical metric. Tropical SVMs find natural applications in phylogenomics (Tang et al., 2020) [10].

Initially formulated as linear programs [5], tropical hyperplanes can also be obtained by tropicalizing a classical hyperplane using $Maslov\ dequantization$ [6] or Viro's method [11]. Yet, redundant features can cause difficulties in the performance of L^2 -norm classical SVMs (Zhang and Zhou, 2010) [12]. Moreover, tropicalization relies on exponentialization and is subject to arithmetic errors, while not optimizing the margin directly in the initial space.

Tropical linear regression has been shown by Akian et al. (2021) [2] as equivalent to solving mean payoff games. We consider this parallel for Tropical Support Vector Machines.

The tropical linear classification problem. The tropical max-plus semifield \mathbb{R}_{max} is the set of real numbers, completed by $-\infty$ and equipped with the addition $a \oplus b = \max(a, b)$ and the multiplication $a \odot b = a + b$. A tropical hyperplane in the d-dimensional tropical vector space \mathbb{R}^d_{max} is a set of vectors of the form

(1.1)
$$\mathcal{H}_a = \left\{ x \in \mathbb{R}^d_{\max}, \quad \max_{1 \le i \le d} (a_i + x_i) \text{ is achieved at least twice} \right\}.$$

Such a hyperplane is parameterized by the vector $a = (a_1, \ldots, a_d) \in \mathbb{R}^d_{\max}$, which is required to be non-identically $-\infty$. It hence partitions the space depending on which coordinate reaches its maximum, laying the ground for d-class classification. By definition, tropical hyperplanes are invariant by translation along the constant vector $(1, \ldots, 1)$.

In this paper, we address the following tropical analogue of support vector machines. Given $n \in \{1, \ldots, d\}$ classes of d-dimensional data points X^1, \ldots, X^p , we look for a maximum-margin separating tropical hyperplane to build a classifier. The notion of margin depends on the metric: the canonical choice in tropical geometry is the (additive version of) Hilbert's projective metric. Its restriction to \mathbb{R}^d is induced by the so-called Hilbert's seminorm or Hopf oscillation

$$||x||_H \coloneqq \max_{1 \le i \le d} x_i - \min_{1 \le i \le d} x_i.$$

It is a projective metric, in the sense that the distance between two points is zero if and only if these two points differ by an additive constant. When dealing with hard-margin support vector machines (SVMs), we also have to ensure that every point lands in the right sector. Each class k belongs to some sectors

 $I^k \subset \{1,\ldots,d\}$, with each sector assigned to only one class. We then define the tropical parameterized hyperplane of configuration $\sigma = \{I^1,\ldots,I^n\}$, where I^k and I^ℓ are disjoint for $k \neq \ell$, as:

$$\mathcal{H} \coloneqq \left\{ x \in \mathbb{R}^d_{\max}, \quad \exists k \neq \ell, \quad (x-a) \, \text{reaches its max coordinate in } I^k \, \text{and} \, I^\ell \right\}.$$

We note $\mathcal{S}^k \subset \mathbb{R}^d_{\max}$ the *sector* classified as class k, where the maximum coordinate is reached in I^k . \mathcal{H} is said to *separate* point clouds $(X^k)_{1 \leq k \leq p}$ with a margin of ν when for every point x^k of class k:

- (1) x^k is on the right sector of the hyperplane, i.e its maximum coordinate is reached in I^k
- (2) x^k is at distance at least ν from any other sector \mathcal{S}^{ℓ} , $l \neq k$:

$$d_H(\mathcal{S}^{\ell}, x^k) \coloneqq \max(x^k - a) - \max_{i \in I^{\ell}} (x^k - a)_i \ge \nu.$$

In Figure 1.1a, we separate two classes of 3-dimensional points from a toy tropical dataset using a tropical hyperplane. The two lower sectors are assigned to one class, the upper sector corresponds to the other class. Its margin is maximal and is represented by the Hilbert ball in gray. The figure is represented in the projective space $\mathbb{P}\left(\mathbb{R}^3_{\max}\right)$, i.e. the quotient of the set of non-identically $-\infty$ vectors of \mathbb{R}^3_{\max} by the equivalence relation which identifies tropically proportional values.

Since our separator is translation-invariant by the constant vector, in order to separate any d-dimensional dataset, we plunge it into the $x_1 + \cdots + x_{d+1} = 0$ subspace of \mathbb{R}^{d+1}_{\max} . This amounts to artificially adding a new coordinate to each point, equal to the opposite of the sum of the initial coordinates. In the rare case where we are dealing with an intrinsically tropical dataset, for which this translation invariance makes sense, it is better to apply the tropical classifier on the initial features.

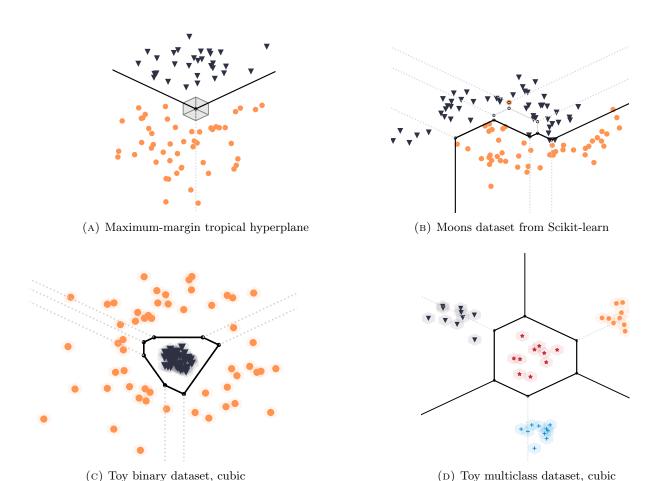


Figure 1.1. Tropical piecewise linear classifiers

Contribution. We formulate the tropical classification problem as a mean payoff game. We apply a Krasnoselskii-Mann iteration scheme to derive an ε -approximation of the tropical classifier in pseudopolynomial time, and we develop geometrical guarantees, e.g. on the margin in the separable case. Finally, we propose a feature map for fitting tropical polynomials while retaining margin information. Thanks to their evaluation speed and their piecewise linear hypersurfaces, we enable more complex classification tasks. All our methods are implemented in Python in the following repository: https: //gitlab.inria.fr/tropical/tropical-svm.

2. Shapley operators and mean payoff games

In this section, we have a finite set of points $X=(x_1,\ldots,x_p)\in\mathbb{R}_{\max}^{d\times p}$. We define the tropical span of X as the set of tropical linear combinations of these points:

$$\mathrm{Span}(X) := \left\{ \max_{1 \le i \le p} (x_i + \lambda_i), \quad \lambda \in \mathbb{R}^p \right\}.$$

As this is also the smallest tropically convex set containing these points, separating several point clouds amounts to separating their tropical spans. Given a tropically convex, compact and nonempty subset $\mathcal V$ of \mathbb{R}^d_{\max} , we define the projection of x in \mathbb{R}^d_{\max} by:

$$P_{\mathcal{V}}(x) \coloneqq \max\{y \in \mathcal{V}, \quad y \le x\}.$$

When V is the tropical span of X, we will also note this projection as P_X . Tropical projections can be expressed as a mean payoff game: for $i \in \{1, ..., p\}$ and x in \mathbb{R}^d_{\max} , [7]

$$P_X(x)_i = \max_{1 \le j \le p} \left\{ X_{ij} + \min_{1 \le k \le d} (-X_{kj} + x_k) \right\}.$$

Here, we recognize the payoff of a zero-sum deterministic game with perfect information. There are two players, "Max" and "Min" (the maximizer and the minimizer), who alternate their actions. Starting on node i, Player Max chooses to move to node j, and receives X_{ij} from Player Min. Similarly, Player Min in turn chooses a node k and has to pay $-X_{kj}$ to Player Max.

We now recall some tools from Perron-Frobenius theory, in relation with mean payoff games. We refer the reader to [1] for more information.

We call Shapley operator a map $T: \mathbb{R}^d_{\max} \to \mathbb{R}^d_{\max}$ that is order-preserving and additively homogenous, i.e. $T(\alpha+x) = \alpha + T(x)$ for all α in \mathbb{R}_{\max} and x in \mathbb{R}^d_{\max} . T is said to be diagonal free when $T_i(x)$ is independent of x_i for all $i \in \{1, \ldots, n\}$, i.e. when for all x, y in \mathbb{R}_{\max} which coincide on all coordinates except the *i*-th, we have $T_i(x) = T_i(y)$.

Tropical projections are Shapley operators. Given a Shapley operator T, we also define

$$\mathcal{S}(T) := \left\{ x \in \mathbb{R}^d_{\max}, \quad x \le T(x) \right\}.$$

More specifically, any tropically convex set V is equal to $S(P_{V})$. Indeed, as for any $x \in \mathbb{R}^{d}_{\max}$, $P_{V}(x) \leq x$, having $x \in \mathcal{S}(P_{\mathcal{V}})$ implies $x = P_{\mathcal{V}}(x)$, meaning x is in \mathcal{V} .

Operator P can be slightly tweaked to make it diagonal-free – but the resulting operator isn't a projector anymore. In the modified game, the opponent is prevented from replying eye-for-an-eye:

$$\left[P_X^{\mathrm{DF}}(x)\right]_i \coloneqq \max_{1 \le j \le p} \left\{ X_{ij} + \min_{k \ne i} (-X_{kj} + x_k) \right\}.$$

Noticeably, $S(P_X^{\mathrm{DF}}) = S(P_X)$. P and P^{DF} can be evaluated in time $\mathcal{O}(dp)$ as they involve computing the two greatest values of $(-X_{kj} + x_k)_{1 \leq k \leq d}$ for $j \in \{1, \ldots, p\}$.

3. Binary Tropical Support Vector Machines

3.1. Hard-Margin SVM. In this section, we introduce a practical criterion for the existence of a separating tropical hyperplane, and if it exists, we construct one with optimal margin. We also place ourselves in the binary setting, where we seek to separate two convex hulls \mathcal{V}^+ and \mathcal{V}^- . We assume we have two Shapley operators T^+ and T^- such that $\mathcal{S}(T^{\pm}) = \mathcal{V}^{\pm}$, which is for instance the case when separating finite point clouds: corresponding operators are projections $P_{\mathcal{V}^{\pm}}$. We define:

$$T = \min(T^+, T^-),$$

which is also a Shapley operator. We denote by \bot the vector of \mathbb{R}^d_{\max} identically equal to $-\infty$. The spectral radius of T is, by definition:

$$\rho(T) = \sup \left\{ \lambda \in \mathbb{R}_{\max}, \quad \exists u \in \mathbb{R}^d_{\max} \setminus \{\bot\}, \quad T(u) = \lambda + u \right\}.$$

Note that S(T) is equal to the intersection of $S(T^+)$ and $S(T^-)$, the convex hulls we are seeking to separate. From Allamigeon et al. [3], we know that V^+ and V^- being disjoint is equivalent to $\rho(T) < 0$. Given a the eigenvector corresponding to $\rho(T)$, we define the sectors:

$$I^{\pm} := \left\{ i \in \{1, \dots, d\}, \quad T^{\pm}(a)_i > \rho(T) + a_i \right\},$$

and the corresponding configuration $\sigma = \{I^+, I^-\}.$

Proposition 1. If $\rho(T) < 0$, the tropical hyperplane of apex a and configuration σ , given the sectors defined above, separates \mathcal{V}^+ and \mathcal{V}^- with a margin of at least $-\rho(T)$. Moreover, this margin is exact when T^{\pm} are of the form

$$\sup_{v \in \mathcal{V}^{\pm}} \left(v_i + \min(-v + x) \right),\,$$

which is in particular the case when separating finite point clouds.

Proof. For instance, let $i \in I^-$. As T^+ is non-expansive, for $v^+ \in \mathcal{V}^+$:

$$v_i^+ \le T^+(v^+)_i = (T^+(v^+) - T^+(a))_i + T^+(a)_i,$$

hence

$$v_i^+ \le \max(v^+ - a) + T^+(a)_i$$
.

As $i \in I^-$, $T^+(a)_i \le \rho(T) + a_i$, and

$$v_i^+ - a_i \le \max(v^+ - a) + \rho(T).$$

In particular, $v_i^+ - a_i < \max(v^+ - a)$ and any element of \mathcal{V}^+ cannot belong to any of the sectors in I^- with respect to \mathcal{H}_a , from which the sectors are well-defined. Finally,

$$d_H(S^-, v^+) = \max(v^+ - a) - \max(v^+ - a)_{I^-} > -\rho(T),$$

hence the margin is at least $-\rho(T)$ which is positive in the separable case, as shown in [3].

Finally, when T^{\pm} are of the aforementioned form, for $i \in I^-$ and $\varepsilon > 0$, we can find $v \in \mathcal{V}^+$ such that

$$T^+(a)_i - \varepsilon \le v_i - \max(v - a) \le T^+(a)_i$$

giving us

$$\rho(T) - \varepsilon < v_i - a_i - \max(v - a) < \rho(T).$$

Maximizing over all $i \in I^-$ gives that v is at most at distance $-\rho(T) + \varepsilon$ of \mathcal{S}^- , hence the margin. \square

Proposition 2. TBD

Proof. TBD \Box

Remark 3. To save on calculations, especially in large dimensions, we don't evaluate the coordinates associated with empty sectors. In fact, this is equivalent to defining corresponding apex coordinates to $+\inf$, which only increases the margin. For example, for $x^+ \in X^+$, and i a sector not assigned to the positive class,

$$d_H(S^-, x^+) = \max(x^+ - a) - (x_i^+ - a_i)$$

is an increasing quantity in a_i , as the max is not reached in the i coordinate. This reasoning also applies to the multiclass case.

To build a practical algorithm, we still need to compute the spectral radius of T. As it amounts to solving mean payoff games, we can apply the Krasnoselskii-Mann iteration scheme described in Algorithm 1.

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Algorithm 1 Krasnoselskii-Mann iteration scheme for finding Shapley eigenpairs

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1: input: A Shapley operator T and a relative convergence criterion \eta > 0.

2: x \coloneqq \mathbf{1}_d \in \mathbb{R}_{\max}^d

3: repeat

4: z \coloneqq (x + T(x))/2

5: x \coloneqq z - \max(z) \cdot \mathbf{1}_d

6: until relative change of x is lower than \eta

7: a \coloneqq x - \overline{x}; \lambda \coloneqq 2 \max(z)

8: return (a, \lambda)
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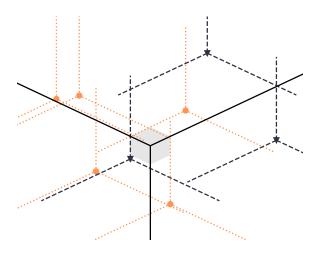


Figure 3.1. Eigenvalue is inner radius of intersection

3.2. Geometric approach for soft-margin SVM. We seek to classify convex hulls \mathcal{V}^+ and \mathcal{V}^- . Assuming that the data overlap, we want to have a sense of their non-separability. Using the operator T we previously defined, $\rho(T)$ is the (positive) inner radius of $\mathcal{S}(T)$, i.e. the supremum of the radii of the Hilbert balls contained in the intersection between \mathcal{V}^+ and \mathcal{V}^- , as shown by Allamigeon et al. (2018) [3] and as seen in Figure 3.1.

The next Proposition shows us that this measure can be interpreted geometrically, in the case where V^+ and V^- are generated by point clouds X^+ and X^- :

Proposition 4. By moving some points from X^+ and X^- by a distance of at most $\rho(T)$, we can nullify the interior of the intersection of the tropical spans.

More specifically, we note $a \in \mathbb{R}^d_{\max}$ the eigenvector corresponding to $\rho(T)$, and \mathcal{H}_a the (unparameterized) tropical hyperplane as defined by equation 1.1. We project all points of X^+ and X^- whose distance to \mathcal{H}_a is less than $\rho(T)$, onto \mathcal{H}_a . Then the intersection of new convex hulls is of empty interior.

Proof. The distance between any point x in \mathbb{R}^d_{\max} and the tropical hyperplane of apex a is:

$$d_H(x, \mathcal{H}_a) = \max(x - a) - \max_2(x - a),$$

where \max_2 stands for the second maximal value of any vector.

Considering X the cloud consisting of all points (regardless of sign), and for $x_j \in X$, we note s_j its sector and d_j the second argmax of $(x_j - a)$. For each point $x_j = X_{j-1} \in X$ at distance less than $\rho(T)$ of \mathcal{H}_a , the transformation described above consists in setting

$$W_{kj} := \begin{cases} X_{kj} & \text{if } k \neq s_j \\ X_{s_j j} - d_H(x_j, \mathcal{H}_a) & \text{at } s_j \end{cases}$$

so that w_j is projected on the hyperplane. As T^{\pm} is non-expansive and diagonal-free, let's remark that for $x^{\pm} \in V^{\pm}$:

$$x_i^{\pm} \le T^{\pm}(x^{\pm})_i = (T^{\pm}(x^{\pm}) - T^{\pm}(a))_i + T^{\pm}(a)_i,$$

hence

(3.1)
$$x_i^{\pm} \le \max(x^{\pm} - a)_{\neq i} + T^{\pm}(a)_i.$$

Let $1 \le i \le d$. If $x_j \in X$ is not in the *i*-th sector, then for k different from the sector of x_j , by definition:

$$(w_j - a)_k \le (x_j - a)_{d_i} \le (w_j - a)_{s_j},$$

hence

$$W_{ij} - \max_{k \neq i} (W_{kj} - a_k) \le a_i.$$

Otherwise, x_i is in the *i*-th sector and:

$$\max_{k \neq i} (W_{kj} - a_k) = X_{d_j j} - a_{d_j},$$

thus

$$W_{ij} - \max_{\neq i} (w_j - a) = (w_j - a)_{s_j} - (x_j - a)_{d_j} + a_{s_j} \ge a_{s_j} = a_i,$$

with equality iff $d_H(x_j, \mathcal{H}_a) \leq \rho(T)$.

Suppose by symmetry that $T(a)_i = T^+(a)_i = \rho(T) + a_i$. We also have $T^-(a)_i \ge \rho(T) + a_i$. Then, using the proof of Theorem 22 in [2], we know that there is j^+ , j^- in $\{1,\ldots,p\}$ such that $x_{j^+} \in X^+$ and $x_{j^-} \in X^-$ are in sector i, with x_{j^+} being at distance $\rho(T)$ from \mathcal{H}_a and x_{j^-} at distance greater than $\rho(T)$. Therefore, $W_{ij^+} - \max_{\neq i} (w_{j^+} - a) = a_i$ and $W_{ij^-} - \max_{\neq i} (w_{j^-} - a) \ge a_i$. Moreover, for any j such that $x_j \in X^+$ is in sector i, equation 3.1 gives $d_H(x_j, \mathcal{H}_a) \le \rho(T)$.

Let Q^{\pm} be the diagonal-free projections over transformed point clouds, and $Q = \inf(Q^+, Q^-)$. We've just shown that $Q(a)_i = Q^+(a)_i = a_i$, and finally Q(a) = a.

Thus, the spectral radius can be interpreted as a bound on the distance under which some points have to be moved to separate the interiors of convex hulls.

As a is the center of the inner ball of $\mathcal{V}^+ \cap \mathcal{V}^-$, we also have a purely geometric heuristic for determining a hyperplane in the non-separable case: we simply take the hyperplane \mathcal{H} of apex a and then assign the sectors based on dominant population, for instance.

4. Multiclass classification

In this section, we consider convex hulls of n point classes, noted $(\mathcal{V}^k)_{1 \leq k \leq n}$ and described by Shapley operators T^k . We give a simple criterion for these classes to be separable by a tropical hyperplane, meaning that there exists a tropical signed hyperplane such that each \mathcal{V}^k belong to sectors of $I^k \subset \{1,\ldots,d\}$ with $(I^k)_{1\leq k\leq n}$ being pairwaise disjoint. We then characterize the margin as another spectral radius in this case.

We now consider the Shapley operator

$$T \coloneqq \max_{1 \le k < l \le n} \min(T^k, T^l),$$

which can be rewritten, for any x in \mathbb{R}^d_{\max} and i in $\{1,\ldots,d\}$:

$$T(x)_i := \max_2 \{ T^k(x)_i, 1 \le k \le n \}.$$

Morally, we can think of S(T) as the union of pairwise intersections of data classes, although the max operator does not strictly speaking correspond to a union. We also remark that when n = 2, T is the same as previously defined.

Let $(a, \rho(T))$ be the eigenpair of T approximated by the Kranoselskii-Mann algorithm, verifying $T(a) = \rho(T) + a$. We define the sectors:

$$I^k := \{1 \le i \le d, \quad T^k(a)_i > \rho(T) + a_i\}.$$

Given the definition of T, it is clear that these sets are disjoint, and we have the following result:

Proposition 5. If $\rho(T) \leq 0$, the tropical parameterized hyperplane of apex a and configuration σ , given the configuration $\sigma = \{I^k\}_{1 \leq k \leq n}$, separates all point clouds with a margin of at least $-\rho(T)$.

Proof. We consider class $k \in \{1, ..., n\}$ and sector S^{ℓ} with $\ell \neq k$. For $j \in I^{\ell}$, we have $T^{k}(a)_{i} \leq a_{i} + \rho(T)$ and therefore can use the same reasoning as in the proof of Proposition 1 to obtain that for all v^{k} in \mathcal{V}^{k} ,

$$d_H(\mathcal{S}^\ell, v^k) = \max(v^k - a) - \max_{i \in I^\ell} (v^k - a)_i \ge -\rho(T),$$

and v^k cannot belong to sector \mathcal{S}^{ℓ} , from which, as again, the sectors are well-defined.

In practice, our method tends to minimize the minimum margin between two classes. For example, if two classes are very close, and the third is far away, the method won't bother separating the latter from the others as much as possible.

But don't assume that the margin is reached, even between two classes. Indeed, let $i \in \{1, ..., d\}$. As $T(a) = \rho(T) + a$, From the expression of T, noting k and ℓ respectively the second argmax and argmax of the family $(T^j(a)_i)_{1 \le j \le n}$, we have

$$T^k(a)_i = \lambda + a_i$$
 and $T^\ell(a)_i > \lambda + a_i$,

meaning $i \in \mathcal{I}^{\ell}$. When separating finite point clouds, T^k is of the form

$$\sup_{v \in \mathcal{V}^k} \left(v_i + \min(-v + x) \right).$$

Hence, adapting the ideas of Proposition 1, for $\varepsilon > 0$, we can find $v \in \mathcal{V}^k$ such that

$$\max(v^k - a) - (v_i^k - a_i) \le -\rho(T) + \varepsilon,$$

which means that

$$d_H\left(\mathcal{V}^k, \mathcal{S}^\ell\right) \le \max(v^k - a) - \max_{j \in I_l} (v^k - a)_j \le \max(v^k - a) - (v_i^k - a_i) \le -\rho(T) + \varepsilon.$$

So, for every sector (assigned to a certain class), there's a point of a class not assigned to that sector within margin distance. In general, nothing says that this situation will be symmetrical between 2 sectors in particular, which would achieve the margin between them. In simple, general cases, more complex relationships can be observed.

Since this method optimizes the minimum margin, it works best on previously balanced data. If necessary, it's still possible to use the binary setting of the previous section with one-vs-all approaches, in order to maximize all margins 2 by 2.

Remark 6. According to Proposition 2, we extend the sector assignment policy as follows: the sector corresponding to the maximum coordinate i is assigned to the class with the greatest number of points at distance $\max(0, -\rho(T))$ from it. Unreached sectors are removed for more sparsity, which preserves the margin, as explained in Remark 3

Conclusion. We therefore introduce Algorithm 2 to solve the tropical linear classification problem.

Algorithm 2 Determining tropical linear classifiers by solving mean payoff games

- 1: **input**: Data classes X^1, \ldots, X^n corresponding to labels $1, \ldots, n$, and a relative convergence criterion η for the iteration scheme.
- 2: $T := \max_2 (T^k)_{1 \le k \le n}$
- 3: $(a, \lambda) := \text{Krasnoselskii-Mann}(T, \eta)$
- 4: **assign** sectors σ to classes based on Remark 6
- 5: **return** classifier H_a^{σ} and margin $\max(0, -\lambda)$

In the separable case, we have margin information, which is exact in the binary setting (as it is reached by the two classes, see Proposition 1 and 5). In the non-separable case, by Proposition 4, the overlap zone is contained around the (non-parameterized) hyperplane H_a in a tubular envelope of radius λ , the (positive) eigenvalue approximated by the algorithm.

Let's analyze the complexity of Algorithm 2. Step 3 typically involves a few dozen iterations, in which n tropical projections have to be calculated each time, for a total cost in $\mathcal{O}(dp)$, where p is the total number of points in the dataset. We therefore have a pseudo-polynomial algorithm for calculating an ε -approximation of the classifier described above.

This gain in efficiency is achieved at the cost that tropical SVMs are ultimately very sensitive to outliers, which count as much as any point in the margin calculation.

Why not tropicalizing classical SVMs? Tropical hyperplanes can be obtained as the limit of classical hyperplanes seen on log-log paper, thanks to what we call *Maslov dequantization* [6] or Viro's method [11]. A natural way of doing Tropical SVM is therefore to take our data to the power $\beta > 0$, apply a classical SVM (which would maximize the margin in the new space), then go back to the logarithm. As β tends towards infinity, we are approaching some limiting tropical hyperplane. However, this method is a source of significant numerical error, especially when β is large. Moreover, it does not guarantee a good margin in the starting space, as illustrated in Figure 4.1, where the hypersurface in yellow tends towards the suboptimal tropical one in red.

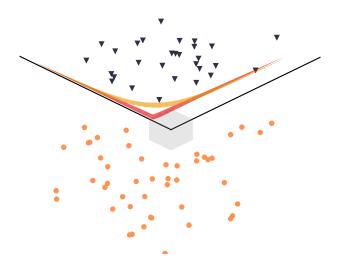


FIGURE 4.1. Tropicalized hyperplane tends towards a suboptimal tropical hyperplane

5. Tropical Polynomials as Piecewise Linear Classifiers

In this section, we extend our previous approach to piecewise linear fitting, using tropical polynomials. A tropical polynomial function over \mathbb{R}^d_{\max} is defined by

$$f(x) = \bigoplus_{\alpha \in \mathbb{Z}^d} f_{\alpha} x^{\alpha} = \max_{\alpha \in \mathbb{Z}^d} (f_{\alpha} + \langle x, \alpha \rangle),$$

where $f_{\alpha} = -\infty$ except for finitely many values of α , and with $\langle x, \alpha \rangle = \sum_{i} x_{i} \cdot \alpha_{i}$ under the convention $(-\infty) \cdot 0 = 0$. The tropical hypersurface associated with f is the set of points from \mathbb{R}^{d}_{\max} where the maximum in f is attained twice. Tropical hypersurfaces are piecewise linear and divide \mathbb{R}^{d}_{\max} between sectors corresponding to the maximal monomial: they can perform more complex classification.

If $\mathcal{A} \subset \mathbb{Z}^d$ is a set of vectors, we define the *Veronese embedding* of x as:

$$\operatorname{ver}_{\mathcal{A}}(x) \coloneqq (\langle x, \alpha \rangle)_{\alpha \in \mathcal{A}} \in \mathbb{R}_{\max}^{\mathcal{A}},$$

allowing us to map our point space into a larger space, made up of various integer combinations of features. To take a fairly exhaustive subset of monomials, we consider the integer combinations defined by the integer points of the dilated simplex. Noting $s \in \mathbb{N}^*$ a scale parameter, we define

$$\mathcal{A}_s := \{ x \in \mathbb{N}^d, \quad x_1 + \dots + x_d = s \},$$

which grows polynomially with s. Note that we choose homogeneous monomials to preserve the translation invariance property in the base space. The Veronese embedding defined by \mathcal{A}_s is invertible and if \mathcal{H}_a is a tropical hyperplane of apex a in $\mathbb{R}^{\mathcal{A}}_{\max}$, $\operatorname{ver}_{\mathcal{A}_s}^{-1}(\mathcal{H}_a)$ is a tropical polynomial hypersurface in the initial space, corresponding to the polynomial

$$f_{-a,\mathcal{A}_s}(x) = \bigoplus_{\alpha \in \mathcal{A}_s} (-a_\alpha) x^\alpha$$

Hence, $\operatorname{ver}_{\mathcal{A}_s}$ is a feature map enabling us to perform piecewise linear classification in \mathbb{R}^d_{\max} . This brings us close to the framework of neural networks with piecewise linear activations, e.g. relaxed linear units [13]. \mathcal{A}_1 corresponds to tropical hyperplanes: we are generalizing the approach previously developed.

The initial metric can be easily linked to the metric in the augmented space: for $x \in \mathbb{R}^d_{\max}$, we have

$$\left[\frac{d}{s+d} + \frac{sd^2}{(s+d)^2}\right]^{-1} \leq \frac{\|\mathrm{ver}(x)\|_{H,\mathbb{R}^{A_s}_{\max}}}{\|x\|_{H,\mathbb{R}^d_{\max}}} \leq s.$$

Indeed,

$$\|\operatorname{ver}(x)\|_{H,\mathbb{R}^{\mathcal{A}_s}_{\max}} = \max_{\alpha \in \mathcal{A}_s} \sum_{i} \alpha_i x_i - \min_{\alpha \in \mathcal{A}^s} \sum_{i} \alpha_i x_i$$
$$\leq s(\max x - \min x) = s\|x\|_{H,\mathbb{R}^d_{\max}}.$$

The rest of the proof is in Appendix A. We deduce the following result:

Proposition 7. If data is separable in the augmented space, noting $(a, \rho(T))$ the eigenpair of T applied to \mathbb{R}^s_{\max} , data is separable in the initial space by the tropical hypersurface of f_{-a,\mathcal{A}^s} with a margin of at least $-\rho(T)/s$.

In Figures 1.1b, 1.1c and 1.1d, we visualize the decision boundaries of the piecewise linear classifier using monomials from A_3 . Dotted lines in light gray indicate merged sectors after assignment, and margin lower estimation is represented by the Hilbert balls around each point. These methods obviously work in higher dimensions, but we can only visualize them in lower ones.

Complexity. The learning speed depends essentially on the dimension of the feature space, i.e. the starting dimension and the number of monomials considered. Although our iteration scheme is efficient, given that

$$|\mathcal{A}^s| = \binom{s+d-1}{s},$$

the complexity grows very quickly, opening the way to feature selection questions for directly learning sparse polynomials.

On the other hand, when the tropical polynomial is trained, since only active monomials are kept, inference is in $\mathcal{O}(rd)$ only, where r is the number of monomials considered.

Feature selection. A heuristic for adaptive monomial choice might be to sample pairs of points of different classes, and make perpendicular slopes admissible. In many cases, this gives better results, as we can see in Figures 5.1c and 5.1d, in which we allow homogeneous positive real monomials.

6. Concluding remarks

We have formulated a pseudo-polynomial algorithm to determine a classifying tropical hyperplane, with geometrical guarantees. It would be interesting to conduct further work on monomial selection heuristics, to better separate complex datasets in the real world.

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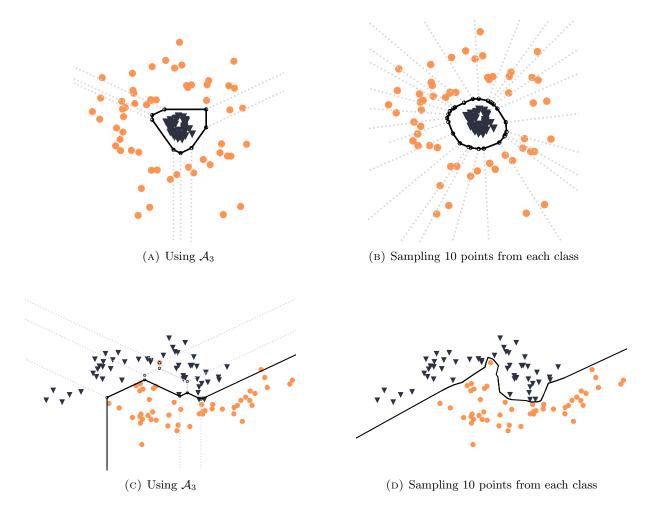


FIGURE 5.1. Feature engineering might yield better results

- [12] Li Zhang and Weida Zhou. On the sparseness of 1-norm support vector machines. Neural Networks, 23(3):373–385, 2010
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APPENDIX A. NORM EQUIVALENCE BETWEEN VERONESE AND STARTING SPACES

We note V the (invertible) Veronese embedding, $Y \in \mathbb{R}^d_{\max}$ from the initial space, and y. For the reciprocal side, using the lemma

$$||Y||_H = 2 \inf_{\alpha \in \mathbb{R}} ||Y + \alpha e||_{\infty},$$

we only need to show that

$$||Y||_{\infty} \le C' ||VY||_{\infty}$$

using some constant C'. Denoting y = VY, we have $V^Ty = V^TVY$, where V^TV is definite positive. Hence, $Y = (V^TV)^{-1}V^Ty$, and

$$||Y||_{\infty} = ||(V^T V)^{-1} V^T y||_{\infty} \le ||(V^T V)^{-1} V^T ||_{\text{op},\infty} ||V Y||_{\infty}$$

hence we define $C' = \|(V^T V)^{-1} V^T\|_{\text{op},\infty}$.

V is a $\ell_{s,d} \times n$ matrix, where $\ell_{s,d}$ is the cardinal of \mathcal{A}^s , i.e. the number of homogenous polynomials of degree d with n variables:

$$\ell_{s,d} = \binom{s+d-1}{d-1}.$$

Rows of V are points from the external face of the dilated simplex. As the latter is invariant by permutating axes, we have

$$V^T V = (V_{\cdot i}^T V_{\cdot j})_{ij} = \begin{pmatrix} \alpha + \beta & \beta \\ & \ddots & \\ \beta & \alpha + \beta \end{pmatrix} = \alpha I_{d \times d} + \beta \mathbf{1}_{d \times d},$$

where $\alpha + \beta = ||V_{\cdot 1}||^2$ and $\beta = \langle V_{\cdot 1}, V_{\cdot 2} \rangle$. We find that:

$$\alpha = \binom{s+d}{s-1} \quad \text{and} \quad \beta = \binom{s+d-1}{s-2},$$
 and as $\mathbf{1}_{d\times d}$ is of rank 1, V^TV is easy to invert, and

$$(V^T V)^{-1} = \alpha^{-1} \left(I - \beta \alpha^{-1} (1 + d\beta \alpha^{-1})^{-1} \mathbf{1}_{d \times d} \right)$$

As $\mathbf{1}_{d \times d} V^T = s \mathbf{1}_{d \times \ell_{s,d}}$,

$$M \coloneqq (V^T V)^{-1} V^T = \alpha^{-1} \left(I - B \mathbf{1}_{d \times \ell_{s,d}} \right),\,$$

where

$$B = s \cdot \beta \alpha^{-1} (1 + d\beta \alpha^{-1})^{-1} = \frac{s - 1}{d + 1}.$$

is a $d \times \ell_{s,d}$ matrix whose lines are permutations of each other. Hence, their L^1 norms are equal and:

$$||M||_{\text{op},\infty} = \alpha^{-1} \sum_{i=1}^{d} |V_{i1} - B| = \alpha^{-1} \sum_{k=0}^{s} \ell_{s-k,d-1} |k - B|.$$

Splitting this sum with respect to the sign of k-B yields, noting b=|B|:

$$||M||_{\text{op},\infty} = \alpha^{-1} \left(\sum_{k=0}^{b} \ell_{s-k,d-1}(B-k) + \sum_{k=b+1}^{s} \ell_{s-k,d-1}(k-B) \right)$$

$$\leq \alpha^{-1} \left(\sum_{k=0}^{b} \binom{s-k+d-2}{d-2} B + \sum_{k=b+1}^{s} \binom{s-k+d-2}{d-2} (s-B) \right)$$

Choosing p+1 items among q+1 ones amounts to first choosing the maximum, and then the n remaining elements below it. Hence

$$\sum_{k=p}^{q} \binom{k}{p} = \binom{q+1}{p+1},$$

and

$$\sum_{b+1 \leq k \leq s} \binom{s-k+d-2}{d-2} = \binom{s-b+d-2}{d-1},$$

and similarly

$$\sum_{0 \leq k \leq b} \binom{s-k+d-2}{d-2} = \ell_{s,d} - \binom{s+d-b-2}{d-1}.$$

Hence

$$||M||_{\text{op},\infty} \le \alpha^{-1} \left[B\ell_{s,d} + (s-2B) \binom{s-b+d-2}{d-1} \right],$$

but

$$B\alpha^{-1}\ell_{s,d} = \frac{(s-1)d}{s(s+d)}, \qquad s-2B = \frac{s(d-1)+2}{d+1},$$

and

$$\frac{\binom{s-b+d-2}{d-1}}{\binom{s+d}{d+1}} = d(d+1)\frac{(s-b+d-2)(s-b+d-3)\cdots(s-b)}{(s+d)(s+d-1)\cdots(s+1)s}$$

$$= \frac{d(d+1)}{(s+d-1)(s+d)}\frac{s-b}{s}\frac{s+1-b}{s+1}\cdots\frac{s+d-2-b}{s+d-2}$$

$$\leq \frac{d(d+1)}{(s+d-1)(s+d)},$$

combinatori possible)

as $b \leq s$. Finally, for $s \geq 2$:

$$\|(V^T V)^{-1} V^T\|_{\infty, \text{op}} \le \frac{(s-1)d}{s(s+d)} + \frac{d(s(d-1)+2)}{(s+d-1)(s+d)}$$
$$\le \frac{d}{s+d} + \frac{sd^2}{(s+d)^2}$$