

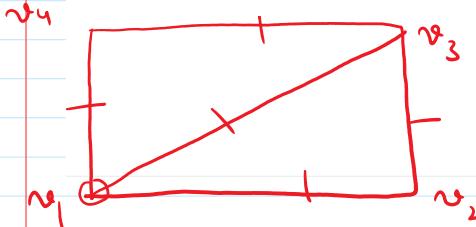
Theorems: Undirected Graphs

Theorem 1

The Handshaking theorem:

$$2e = \sum_{v \in V} \deg(v)$$

(why?) Every edge connects 2 vertices



Verify Hand-shaking for the above graph

$$\deg(v_1) = 2, \deg(v_2) = 2, \deg(v_3) = 2, \deg(v_4) = 2$$

$$\deg(v_1) + \deg(v_2) + \deg(v_3) + \deg(v_4) = 2 + 2 + 2 + 2 = 8$$

$$= 2 \times 4$$

$$= 2 \times \text{no. of edges}$$

$$\sum_{i=1}^n \deg(v_i) = 2e$$

$$= 2 \times 4 = 8$$

$$= 2e$$

$$= 2e$$

$$= 2 \times 4 = 8$$

Theorems: Undirected Graphs

Theorem 2:

An undirected graph has even number of vertices with odd degree

Proof V is the set of even degree vertex and V_2 refers to odd degree vertex

$$2e = \sum_{v \in V} \deg(v) = \sum_{u \in V_1} \deg(u) + \sum_{v \in V_2} \deg(v)$$

$\Rightarrow \deg(v)$ is even for $v \in V_1$.

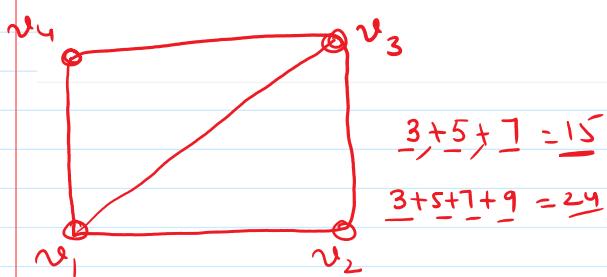
\Rightarrow The first term on the right-hand side of the last inequality is even.

\Rightarrow The sum of the last two terms on the right-hand side of the last inequality is even since it is sum of two even numbers.

Hence $\deg(v)$ is even for all $v \in V_1$.

$$\Rightarrow \text{second term } \sum_{v \in V_2} \deg(v) \text{ is even}$$

$$2+4+6 = \underline{\underline{\text{even}}}$$



$$V_1 = \{v_1, v_3\} \checkmark$$

$$V_2 = \{v_2, v_4\} \checkmark$$

$$3+5+7 = \underline{\underline{15}}$$

$$3+5+7+9 = \underline{\underline{24}}$$

$$\sum_{v_i \in \text{odd}} \deg(v_i) + \sum_{v_i \in \text{even}} \deg(v_i) = 2e$$

$$\sum_{v_i \in \text{odd}} \deg(v_i) + \text{even} = \underline{\underline{\text{even}}}$$

$$\sum_{v_i \in \text{odd}} \deg(v_i) = \underline{\underline{\text{even}}} - \underline{\underline{\text{even}}} = \underline{\underline{\text{even}}}$$

This situation is possible only when no of odd degree vertices are even.

Is it possible to have a graph with the following degree of vertices: $\frac{1}{\text{odd}}, 2, \frac{3}{\text{odd}}, 4, \frac{5}{\text{odd}}$.

This is not possible.

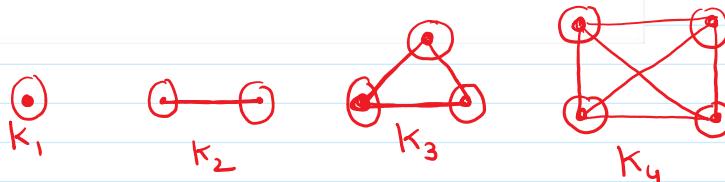
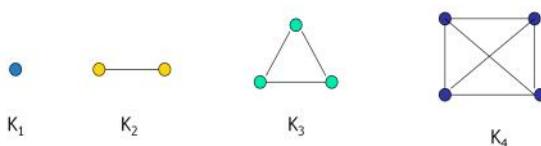
~~if v....., odd - , odd , odd~~

This is not possible.

Simple graphs – special cases

- **Complete graph:** K_n is the simple graph that contains exactly one edge between each pair of distinct vertices.

Representation Example: K_1, K_2, K_3, K_4



In K_n graph

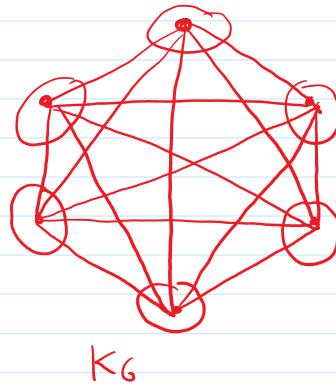
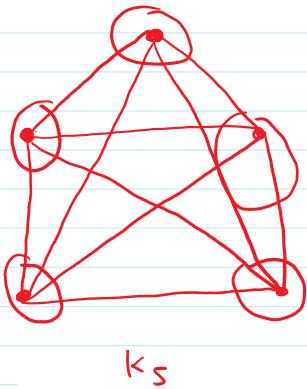
degree of each vertex = $n-1$

In K_{100} graph

what is degree of each vertex?

- (A) 98 (B) 99 (C) 100 (D) 101

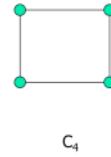
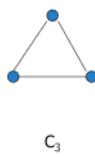
Q Draw K_5 and K_6 graph.



Simple graphs – special cases

- **Cycle:** C_n , $n \geq 3$ consists of n vertices $v_1, v_2, v_3 \dots v_n$ and edges $\{v_1, v_2\}, \{v_2, v_3\}, \{v_3, v_4\} \dots \{v_{n-1}, v_n\}, \{v_n, v_1\}$

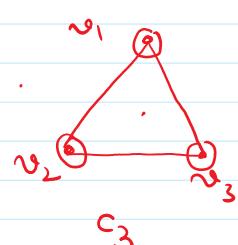
Representation Example: C_3, C_4



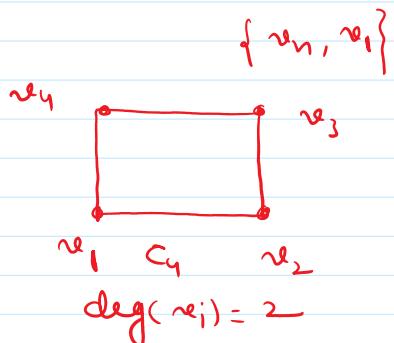
Draw C_5, C_6, C_7 & C_8 cycle.

Cycle C_n , $n \geq 3$

$\{v_1, v_2\}, \{v_2, v_3\}, \dots, \{v_{n-1}, v_n\}, \{v_n, v_1\}$



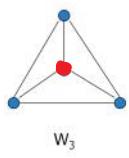
$\deg(v_i) = 2$



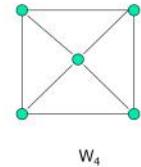
Simple graphs – special cases

- Wheels:** W_n , obtained by adding additional vertex to C_n and connecting all vertices to this new vertex by new edges.

Representation Example: W_3, W_4

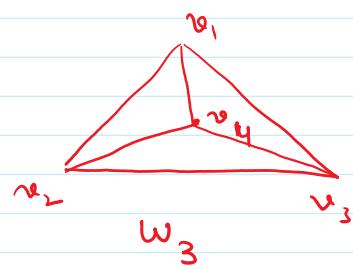


W_3

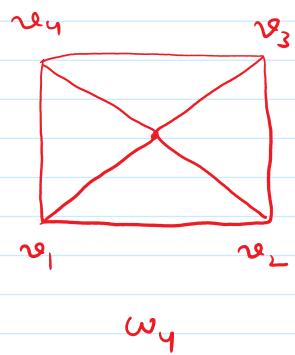


W_4

Wheel: $\underline{w_m}$



w_3



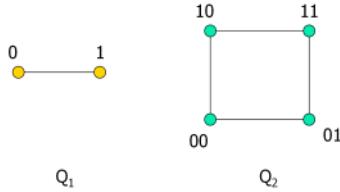
w_4

Draw w_5, w_6, w_7, w_8 .

Simple graphs – special cases

- N-cubes:** Q_n , vertices represented by $2n$ bit strings of length n . Two vertices are adjacent if and only if the bit strings that they represent differ by exactly one bit positions

Representation Example: Q_1, Q_2

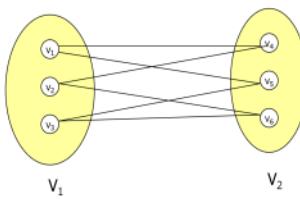


Bipartite graphs

- In a simple graph G , if V can be partitioned into two disjoint sets V_1 and V_2 such that every edge in the graph connects a vertex in V_1 and a vertex in V_2 (so that no edge in G connects either two vertices in V_1 or two vertices in V_2)

Application example: Representing Relations

Representation example: $V_1 = \{v_1, v_2, v_3\}$ and $V_2 = \{v_4, v_5, v_6\}$,



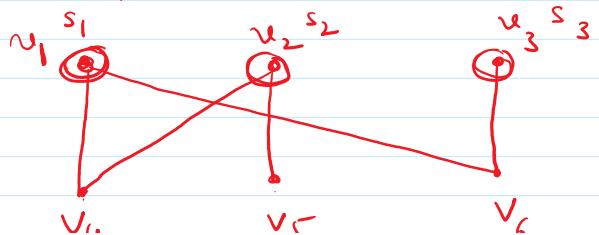
Bi-partite graph.

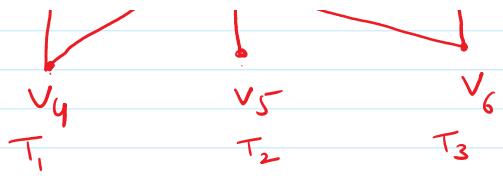
Two partitions of the vertex set

$$V = \{v_1, v_2, v_3, v_4, v_5, v_6\}$$

$$V_1 = \{v_1, v_2, v_3\} \quad V_1 \cap V_2 = \emptyset$$

$$V_2 = \{v_4, v_5, v_6\} \quad V_1 \cup V_2 = V$$

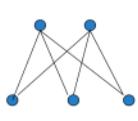




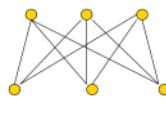
Complete Bipartite graphs

- $K_{m,n}$ is the graph that has its vertex set portioned into two subsets of m and n vertices, respectively. There is an edge between two vertices if and only if one vertex is in the first subset and the other vertex is in the second subset.

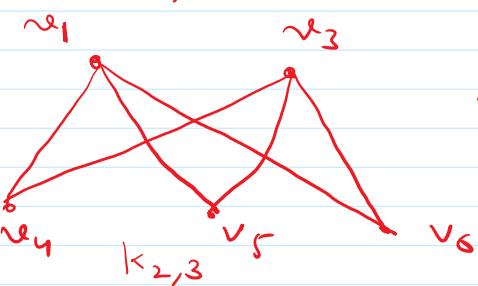
Representation example: $K_{2,3}, K_{3,3}$



$\underline{K}_{2,3}$

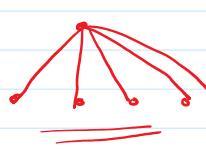


$\underline{K}_{3,3}$



$\underline{K}_{2,3}$

$\underline{K}_{1,5}$

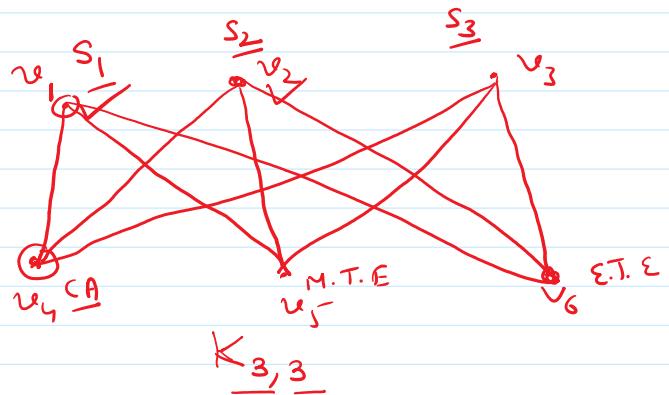


$$V = \{v_1, v_2, v_3, v_4, v_5, v_6\}$$

$$V_1 = \{v_1, v_2, v_3\} \checkmark$$

$$V_2 = \{v_4, v_5, v_6\} \checkmark$$

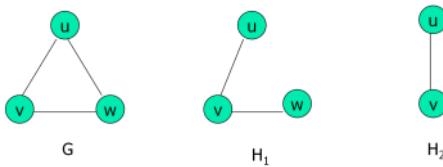
$$V_1 \cap V_2 = \emptyset, \quad V_1 \cup V_2 = V$$



$\underline{K}_{3,3}$

Subgraphs

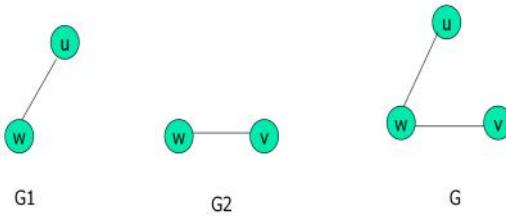
- A subgraph of a graph $G = (V, E)$ is a graph $H = (V', E')$ where V' is a subset of V and E' is a subset of E
- Application example: solving sub-problems within a graph
Representation example: $V = \{u, v, w\}$, $E = (\{u, v\}, \{v, w\}, \{w, u\})$, H_1 , H_2



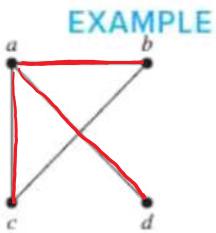
Subgraphs

- $G = G_1 \cup G_2$ wherein $E = E_1 \cup E_2$ and $V = V_1 \cup V_2$, G , G_1 and G_2 are simple graphs of G

Representation example: $V_1 = \{u, w\}$, $E_1 = \{\{u, w\}\}$, $V_2 = \{w, v\}$, $E_2 = \{\{w, v\}\}$, $V = \{u, v, w\}$, $E = \{\{u, w\}, \{w, v\}\}$



Adjacency matrix



Use an adjacency matrix to represent the graph shown in Figure .

Solution: We order the vertices as a, b, c, d . The matrix representing this graph is

$$\begin{bmatrix} 0 & 1 & 1 & 1 \\ 1 & 0 & 1 & 0 \\ 1 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 \end{bmatrix}.$$

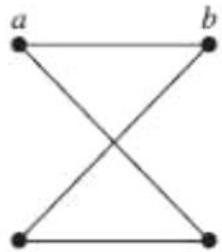
A simple graph.

$$A = \begin{bmatrix} 0 & 1 & 1 & 1 \\ 1 & 0 & 1 & 0 \\ 1 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 \end{bmatrix}$$

	a	b	c	d
a	0	1	1	1
b	1	0	1	0
c	1	1	0	0
d	1	0	0	0

EXAMPLE

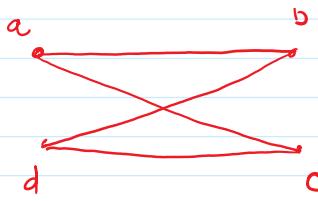
Draw a graph with the adjacency matrix



$$\begin{array}{ccccc} & a & b & c & d \\ a & \left[\begin{array}{cccc} 0 & 1 & 1 & 0 \\ 1 & 0 & 0 & 1 \\ 1 & 0 & 0 & 1 \\ 0 & 1 & 1 & 0 \end{array} \right] & \checkmark & & \end{array}$$

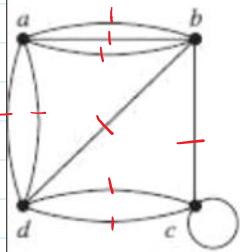
with respect to the ordering of vertices a, b, c, d .





EXAMPLE

Use an adjacency matrix to represent the pseudograph shown.



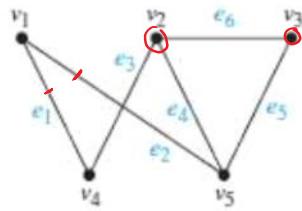
Solution: The adjacency matrix using the ordering of vertices a, b, c, d is

$$\begin{bmatrix} 0 & 3 & 0 & 2 \\ 3 & 0 & 1 & 1 \\ 0 & 1 & 1 & 2 \\ 2 & 1 & 2 & 0 \end{bmatrix}.$$

	a	b	c	d
a	0	3	0	2
b	3	0	1	1
c	0	1	1	2
d	2	1	2	0

EXAMPLE Represent the graph shown in Figure with an incidence matrix.

Solution: The incidence matrix is



$$\begin{array}{c} e_1 \quad e_2 \quad e_3 \quad e_4 \quad e_5 \quad e_6 \\ \hline v_1 & 1 & 1 & 0 & 0 & 0 & 0 \\ v_2 & 0 & 0 & 1 & 1 & 0 & 1 \\ v_3 & 0 & 0 & 0 & 0 & 1 & 1 \\ v_4 & 1 & 0 & 1 & 0 & 0 & 0 \\ v_5 & 0 & 1 & 0 & 1 & 1 & 0 \end{array}$$

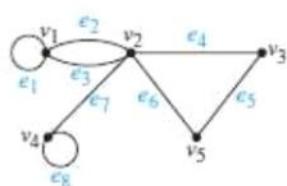
Incidence matrix
may or may not be
a square matrix.

FIGURE An undirected graph.

	e_1	e_2	e_3	e_4	e_5	e_6
v_1	1	1	0	0	0	0
v_2	0	0	1	1	0	1
v_3	0	0	0	0	1	1
v_4	1	0	1	0	0	0
v_5	0	1	0	1	1	0

EXAMPLE Represent the pseudograph shown in Figure using an incidence matrix.

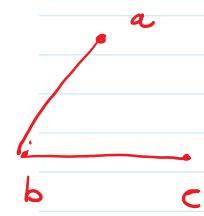
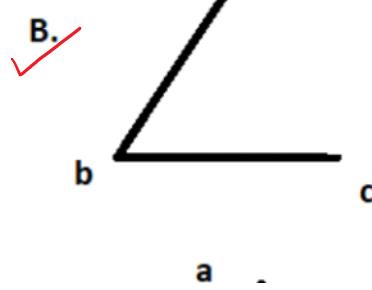
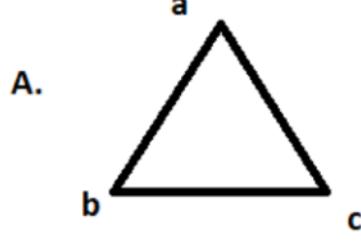
Solution: The incidence matrix for this graph is



$$\begin{matrix} & e_1 & e_2 & e_3 & e_4 & e_5 & e_6 & e_7 & e_8 \\ v_1 & 1 & 1 & 1 & 0 & 0 & 0 & 0 & 0 \\ v_2 & 0 & 1 & 1 & 1 & 0 & 1 & 1 & 0 \\ v_3 & 0 & 0 & 0 & 1 & 1 & 0 & 0 & 0 \\ v_4 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 1 \\ v_5 & 0 & 0 & 0 & 0 & 1 & 1 & 0 & 0 \end{matrix}$$

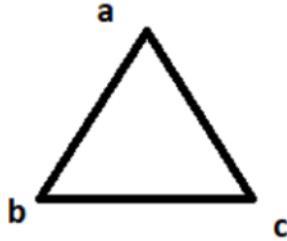
A pseudograph.

Quiz 2 : Which of the following graph has the adjacency matrix $\begin{bmatrix} a & b & c \\ a & 0 & 1 & 0 \\ b & 1 & 0 & 1 \\ c & 0 & 1 & 0 \end{bmatrix}$?

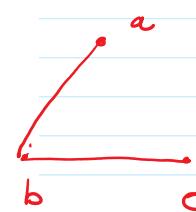
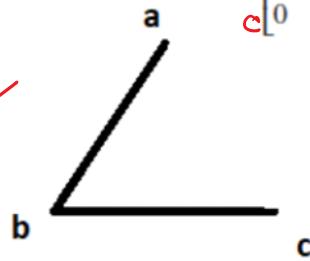


Quiz 2 : Which of the following graph has the adjacency matrix $\begin{bmatrix} 0 & 1 & 0 \\ 1 & 0 & 1 \\ 0 & 1 & 0 \end{bmatrix}$?

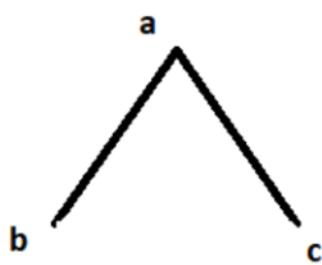
A.



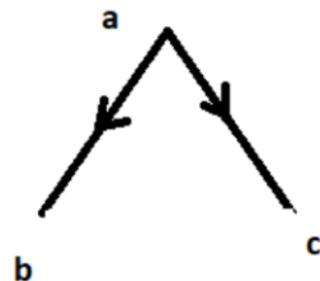
B.



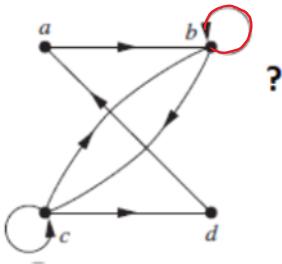
C.



D.



Quiz 3 : Which of the following adjacency matrix represents



A. $\begin{pmatrix} 1 & 1 & 1 & 0 \\ 0 & 1 & 0 & 1 \\ 0 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 \end{pmatrix} \times$

C. $\begin{pmatrix} 0 & 1 & 0 & 0 \\ 0 & 1 & 1 & 0 \\ 0 & 1 & 1 & 1 \\ 1 & 0 & 0 & 0 \end{pmatrix} \checkmark$

B. $\begin{pmatrix} 1 & 1 & 1 & 0 \\ 0 & 1 & 0 & 1 \\ 0 & 1 & 0 & 0 \\ 1 & 1 & 1 & 1 \end{pmatrix} \times$

D. $\begin{pmatrix} 0 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 1 & 1 & 1 \\ 1 & 0 & 0 & 1 \end{pmatrix} \times$

	a	b	c	d
a	0	1	0	0
b	0	1	1	0
c				
d				

$$\alpha_n + \alpha_{n-1} + \alpha_{n-2} + \alpha_{n-3} = 0$$

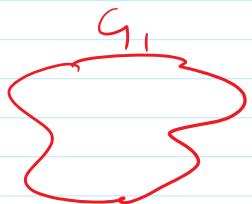
$$\alpha^3 + \alpha^2 + \alpha + 1 = 0$$

$$\alpha_n - \alpha_{n-(k-1)} - \alpha_{n-k} = 0$$

Its order = $n - (n-k) = k$

$$\alpha^K = (\alpha - 1)^k = 0$$

Isomorphic Graphs

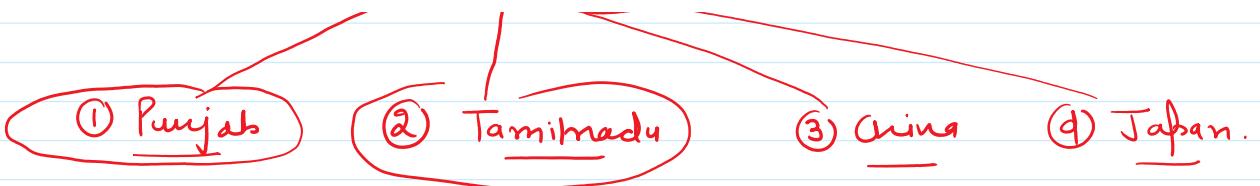


G1

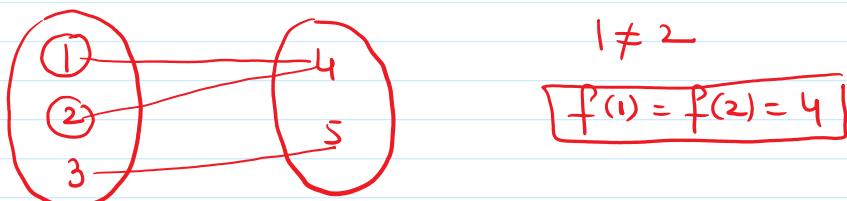


G2

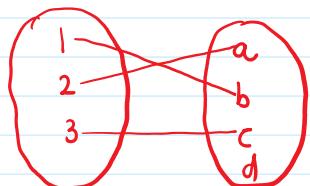
object
Apple



One-one mapping: Let X and Y be two non-empty sets. Then a function $f: X \rightarrow Y$ is called 1-1 function if $x_1, x_2 \in D_f$ s.t. $x_1 \neq x_2 \Rightarrow f(x_1) \neq f(x_2)$



This function is not a 1-1 function / injective function.

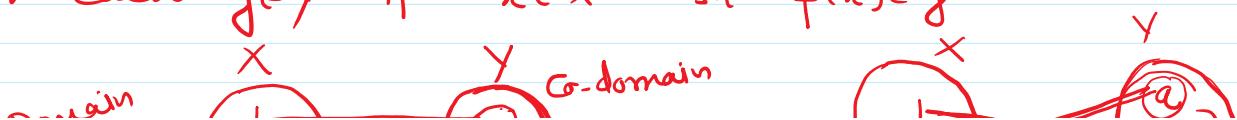


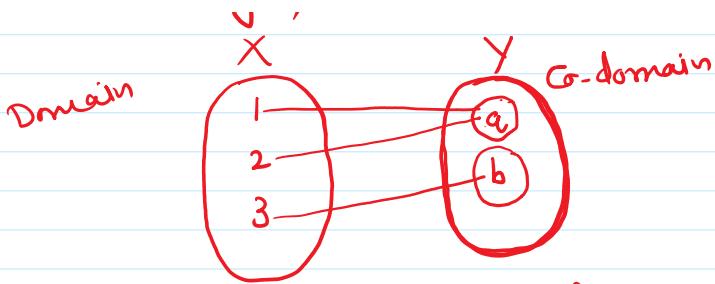
This function is a 1-1 function.

Onto function: Let X and Y be two non-empty sets.

Then a function $f: X \rightarrow Y$ is called an onto function.

If for each $y \in Y$ if $x \in X$ s.t. $f(x) = y$

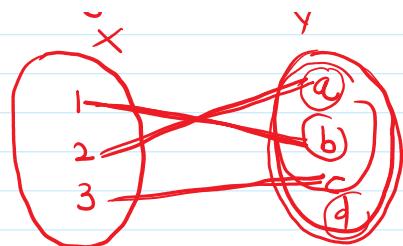




This function is an onto function.

$$\text{Range} = \{a, b\}, \quad \text{Codomain} = \{a, b\}$$

$\text{Range} = \text{Codomain}$



This function is not onto function.

$$\text{Range} = \{a, b, c\}$$

$$\text{Codomain} = \{a, b, c, d\}$$

$\text{Range} \subseteq \text{Codomain}$

Bijective function: A function $f: X \rightarrow X$ is called bijective function if it is both 1-1 and onto.

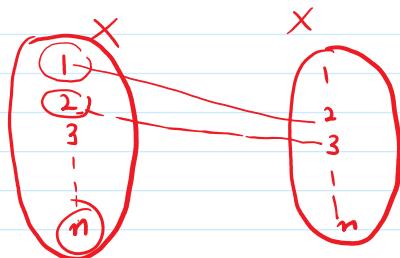
Total no. of bijective functions:

If X contains n elements then total no. of bijections

$$= n!$$

$$= n \times (n-1) \times (n-2) \times \dots \times 1$$

$$= n!$$



How many bijections are possible for a set X having two elements $\{a, b\}$. $2! = 2 \times 1 = 2$

