

Synthesizing Complementary Circuits Automatically

An Automatic Approach of
Complementary Circuits Synthesizing

ABSTRACT

One of the most difficult jobs in designing communication and multimedia chips, is to design and verify complex complementary circuit pair (E, E^{-1}) , in which circuit E transforms information into a format that is suitable for transmission and storage, while E 's complementary circuit E^{-1} recovers this information.

In order to ease this job, we propose a novel two-step approach to synthesize complementary circuit E^{-1} from E fully automatically. First, we assume the circuit E satisfies parameterized complementary assumption, which means its ~~parameter values~~ **parameter values** can be recovered from its output under some parameter setting. We check ~~this assumption~~ **values** with SAT solver and find out proper value for these parameters. Second, with ~~parameters value~~ **parameters value** and SAT instance obtained in the first step, we build the complementary circuit E^{-1} with an efficient satisfying assignments enumeration technique that is specially designed for circuits with lots of XOR gates.

To illustrate its usefulness and efficiency, we run our algorithm on several complex encoders from industrial projects, including PCIE and 10G ethernet, and successfully generate correct complementary circuits for them.

Categories and Subject Descriptors

B.5.2 [REGISTER-TRANSFER-LEVEL IMPLEMENTATION]: Design Aids—*Automatic synthesis*; B.4.4 [INPUT/OUTPUT AND DATA COMMUNICATIONS]: Performance Analysis and Design Aids—*Formal models*

General Terms

Algorithms, Design, Theory, Verification

Keywords

Synthesis, Complementary Circuit, Satisfying Assignments Enumeration

1. INTRODUCTION

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Communication and multimedia electronic applications are major driving forces of semiconductor industry. Many leading edge communication protocols and media formats, even still in non-standardized draft status, are implemented in chips and pushed to market, to maximize the chance of being accepted by consumer and becoming the de facto standards. Two such well known stories are the 802.11n wireless standard competition [1], and the disk format war between HD and blue ray [2]. In such highly competitive markets, designing correct chip is a great challenge.

One of the most difficult jobs in designing communication and multimedia chips, is to design and verify complex complementary circuit pair (E, E^{-1}) , in which circuit E transforms information into a format that is suitable for transmission and storage, while E 's complementary circuit E^{-1} recovers this information. To ease this job, we propose a novel approach to synthesize E^{-1} from E fully automatically in two steps.

1. In the first step, we assume that, for E under some parameters valuation, its input alphabet sequence can be uniquely determined by its output alphabet sequence. We call this assumption **parameterized complementary assumption**. We use a SAT solver to check this assumption and find out proper value for some parameters that make this assumption hold.
2. In the second step, with the SAT instance and ~~parameters value~~ **parameters value** obtained in the first step, we build a circuit E^{-1} with ~~an efficient~~ **an efficient** satisfying assignments enumeration **technology** (abbreviated as **ALLSAT**), which is specially designed for communication and arithmetic circuit with lots of XOR gates.

We implement our algorithm on zchaff [3], and run it on several complex encoder circuits from industrial projects, including PCIE and 10G Ethernet. We can build complementary circuits for all of them within 3000 seconds.

The contribution of this paper is twofold: 1) We propose the first approach to decide if it's possible to recover input sequence of an circuit E from its output sequence. 2) We propose an efficient ALLSAT algorithm for XOR in-

the high frequency and reliability
and compression ratio we are
pursuing highly affect the depth
of pipelines and complexity of
encoding mechanism.
(原句逻辑关系有误)

values

parameter values

All-SAT?

the SAT instance of circuit E

tensive circuits, to build complementary circuit E^{-1} from circuit E 's SAT instance.

The remainder of this paper is as follows. Section 2 presents background information. Section 3 presents how to check parameterized values under an assumption, and how to find a proper value of its parameters. Section 4 presents how to build complementary circuit. Section 5 presents experimental results of our approach. Section 6 presents related works. Section 7 concludes with a note on future work.

2. PRELIMINARIES

2.1 Propositional Satisfiability Problem

For a boolean formula F over variable set V , the **Propositional Satisfiability Problem** (abbreviated as **SAT**) is to find an **satisfying assignment** $A : V \rightarrow \{0, 1\}$, such that F can be evaluated to 1.

If such a satisfying assignment exists, then F is called a **satisfiable formula**, or else it is called **unsatisfiable formula**.

A computer program that decides the existence of such satisfying assignment is called **SAT solver**.

For an assignment $A : U \rightarrow \{0, 1\}$, if $U \subset V$, then A is a **partial assignment**; if $U \equiv V$, then A is a **total assignment**.

For an assignment $A : U \rightarrow \{0, 1\}$, and $W \subset U$, $A|_W : W \rightarrow \{0, 1\}$ is the **projection** of A on W . Its definition is, for $v \in W$, $A|_W(v) = A(v)$. Intuitively, $A|_W$ is obtained from A by removing all variables $v \notin W$.

For an assignment $A : U \rightarrow \{0, 1\}$, and $u \notin U$, and $b \equiv 0$ or 1, $A|^{u \leftarrow b}$ is the **extension** of A on u , its definition is:

$$A|^{u \leftarrow b}(v) = \begin{cases} A(v) & v \in U \\ b & v = u \end{cases}$$

Intuitively, $A|^{u \leftarrow b}$ is obtained from A by adding assignment of u .

Normally, SAT solver takes a formula F to be represented in **CNF** form. F can be represented as $F = \bigwedge_{cl \in CL} cl$ is a conjunction of its clauses set CL , and **clause** $cl = \bigvee_{l \in Lit} l$ is a disjunction of its literals set Lit , and **literal** is a variable v or its negation $\neg v$. Lit is also called **SAT instance**.

For a satisfying assignment A of formula F , its **blocking clause** is:

$$bcl_A = \bigwedge_{v \in V} (A(v) \vee \neg A(v)) \quad (1)$$

It is obvious that A isn't satisfying assignment of $F \wedge bcl_A$. So bcl_A can be inserted into SAT solver to prevent A from becoming satisfying assignment again.

2.2 Satisfying Assignments Enumeration

Technologies that enumerate all satisfying assignments of a formula are called **ALLSAT**. It is obvious that we can enumerate all total satisfying assignments A by repeatedly calling a SAT solver, and adding blocking clause bcl_A of satisfying assignment A into SAT solver, until no more new satisfying assignments can be found.

But for a formula with n variables, there may be 2^n satisfying assignments to be enumerated. Thus, this approach is impractical for large n .

with a magnificent value

In order to reduce the number of satisfying assignments to be enumerated, we need **satisfying assignments minimization** technology to remove irrelevant variable's assignments from satisfying assignment A , such that A can cover more total satisfying assignments. For example, for OR gate $z \leftarrow u \vee v$, its total satisfying assignments that can make $z \equiv 1$ are $\{u \leftarrow 1, v \leftarrow 0\}, \{u \leftarrow 1, v \leftarrow 1\}$ and $\{u \leftarrow 0, v \leftarrow 1\}$, they contain 6 assignments to individual variables. It's obvious that the first two assignments can be merged into $\{u \leftarrow 1\}$, in which assignment to v is irrelevant. After two assignments can be merged into one, the number of assignments to u is removed. These two new merged partial assignments contain only two assignments to individual variables, which are much more succinct than previous three total assignments.

Formally, assume F is the formula over boolean variables set V , and obj is one object variable that should always be 1, and A is a satisfying assignment of $F \wedge obj$. If $F \wedge \neg obj \wedge A|_{V-\{v\}}$ is unsatisfiable, then $A|_{V-\{v\}}$ can't make obj to be 0, so obj must still be 1. Thus, we can merge $A|_{V-\{v\}}|^{v \leftarrow 0}$ and $A|_{V-\{v\}}|^{v \leftarrow 1}$ to remove v from A , a succinct partial satisfying assignment $A|_{V-\{v\}}$.

All existing ALLSAT approaches [4, 5, 6, 7, 8, 9, 10, 24] share this idea of satisfying assignments minimization.

On the other hand, for XOR gate $z \leftarrow u \oplus v$, its total satisfying assignments that can make $z \equiv 1$ are $\{u \leftarrow 1, v \leftarrow 0\}$ and $\{u \leftarrow 0, v \leftarrow 1\}$, they can't be merged. Unfortunately, XOR gates are widely used in almost all communication circuits, including but not limited to scrambler and descrambler, CRC generator and checker, pseudo random test pattern generator and checker.

An extreme example is a n -bits comparator that compares two n -bits variables. There are 2^n total satisfying assignments for this comparator, none of them can be merged with each other.

Thus, enumerating satisfying assignments for XOR intensive circuits is a major difficulty of state-of-the-art ALLSAT approaches, we will solve this problem in section 4.

3. CHECKING PARAMETERIZED COMPLEMENTARY ASSUMPTION

In this section, we will introduce how to check whether input sequence of circuit E can be recovered from its output sequence.

3.1 Parameterized Complementary Assumption

Our circuit cares about the input and output sequence of circuits E , so **Mealy finite state machine**[25] is a suitable model for us.

Definition 1. **Mealy finite state machine** is a 6-tuple $M = (S, s_0, I, O, T, G)$, consisting of the following

1. A finite set of state S
2. An initial state $s_0 \in S$
3. A finite set of input alphabet I
4. A finite set of output alphabet O
5. A state transition function $T : S \times I \rightarrow S$
6. An output function $G : S \times I \rightarrow O$

The circuit E can be modeled by such a Mealy finite state machine. The relationship between its output sequence $o \in O^\omega$ and input sequence $i \in I^\omega$ is shown in figure ???. This relationship is built by unfolding the transition function T and output function G by d times, as shown in formula (2).

(2)

In order to recover $i \in I^\omega$ from $o \in O^\omega$, we must know how to compute i_n for every n , that is, to find a function f^{-1} that can compute i_n from $o \in O^\omega$.

But due to the limited memory of realistic computers, we can't take the infinite length sequence $o \in O^\omega$ as input to f^{-1} , we can only use a finite length sub-sequence of o . This sub-sequence has two parameters, its length l and its delay d compared to i_n , as shown in following figure ??.

Thus, $f^{-1} : O^l \rightarrow I$ is now a boolean function that takes finite length sequence $\langle o_{n+d-l}, \dots, o_{n+d-1} \rangle$ as input, and computes i_n .

For a particular pair of d and l , f^{-1} exists if the following assumption holds:

Definition 2. Parameterized Complementary Assumption: For any valuation of sequence $\langle o_{n+d-l}, \dots, o_{n+d-1} \rangle$, assume there exists no more than one valuation of i_n that can make formula (2) satisfiable.

This assumption holds if and only if following formula (3) is unsatisfiable.

While the first line of formula 3 is the (3)

In formula (3), the first line is same as formula (2), the second line is a copy of formula (2), except that its variables are all renamed by appending a prime. These two lines mean two individual unfolding of circuit E . The third line constraints their output sequences $\langle o_{n+d-l}, \dots, o_{n+d-1} \rangle$ and $\langle o'_{n+d-l}, \dots, o'_{n+d-1} \rangle$ to be the same, and the fourth line constraints that their input alphabet i_n and i'_n are different.

For a particular pair of d and l , checking formula (3) may return two results:

1. **Satisfiable.** This means, for a $\langle o_{n+d-l}, \dots, o_{n+d-1} \rangle$, there exist two different i_n and i'_n that can both make formula (2) satisfiable. This violates definition 2, so no f^{-1} exists for this pair of d and l . We should continue

2. **Unsatisfiable.** This means parameterized complementary assumption is satisfied, $f^{-1} : O^l \rightarrow I$ exists for this pair of d and l . We will build f^{-1} with formula (2) in section 4.

3.2 Ruling out Invalid Input Alphabets with on the basis of

Most communication protocols and systems have some restrictions on valid pattern of input alphabet. Assume this restriction is expressed as an assertion predicate $R : I \rightarrow \{0, 1\}$, in which $R(i_n) \equiv 0$ means that i_n is an invalid input alphabet. Invalid input alphabets will be mapped

some predefined error output alphabet, that is, for $i_n, i'_n \in \{i_n | R(i_n) \equiv 0\}$, they will both be mapped to the same error output alphabet $e \in O$. This will prevent our approach from distinguishing i_n from i'_n .

Such restrictions are often documented clearly in specification of communication protocols, so we chose to employ an assertion based mechanism, such that the user can code these restrictions R into their script or source code.

Thus, formula (2) and (3) should be changed into following formula (4) and (5), in which bold formulas are used to account for predicate R .

rewritten as following Eqn.4 and 5

(4)

(5)

3.3 Approximating Reachable State Set with Prefix Sequence

In last subsection, we have constrained the valid pattern of i_m . But s_n in figure ?? still hasn't been constrained yet. This s_n may be outside of reachable state set of circuit E , which may make checking parameterized complementary assumption fail unnecessary.

Assume circuit E can be modeled by Mealy state machine $M_E = (S, s_0, I, O, T, G)$. Its reachable state set with assertion predicate R is

rewrite Eqn.4 and 5 as Eqn. 7 and 8

(6)

Thus, to rule out unreachable s_n , we need to change formula (4) and (5) into formula (7) and (8) below:

(7)

On the first extreme occasion, (8)

Now we have two extreme cases:

1. One extreme case is formula (4) and (5) with error ratio in checking complementary assumption.
2. The other extreme case is formula (7) and (8), which doesn't affected by unreachable state set, but with very high computation complexity in computing reachable state set RS_E .

So can we find a tradeoff between these two extremes? That should be a method with both acceptable computation complexity and low risk of unnecessary fail in checking parameterized complementary assumption.

To achieve this, we approximate RS_E with a prefix state transition sequence of length p :

create the function

Obviously, Eqn.9

It is obvious that formula (9) is very similar to (6), except that (9) doesn't consider initial state s_0 . So RS_E and all RS_E^p form a total order relation :

Eqn.9

Eqn.6

that

$RS_E \subseteq \dots \subseteq RS_E^p \subseteq \dots \subseteq RS_E^q \subseteq \dots$ where $p > q$

So now, in addition to parameter d and l , we have the third parameter p to be searched. In order to account for RS_E^p , we need to change formula (4) and (5) to following formula (10) and (11), by changing starting index of m from

rewrite Eqn.4 and 5
as Eqn.10 and 11

through sets of values

Now put it altogether, with formula (10) and (11), we iterate through all valuations of d, l and p from smaller one to larger one, until we find one valuation of d, l and p that makes formula (11) unsatisfiable, then that valuation and F_E in formula (10) will be used in section 4 to build complementary circuit E^{-1} .

4. BUILDING COMPLEMENTARY CIRCUIT WITH ALLSAT ALGORITHM DESIGN FOR XOR INTENSIVE CIRCUITS

If we find proper value for d, l at iteration 3, we can now build the complementary circuit boolean function $f^{-1} : O^l \rightarrow I$ in this section.

4.1 Algorithm Framework

According to section 3.3, $f^{-1} : O^l \rightarrow I$ can be built from formula (10) by enumerating satisfying assignments of $\langle o_{n+d-l}, \dots, o_{n+d-1} \rangle$ and i_n .

Assume the set of all total satisfying assignments of formula (10) is $\{A_1, \dots, A_t\}$, then f^{-1} can be defined as:

$$f^{-1}(\langle o_{n+d-l}, \dots, o_{n+d-1} \rangle) = \begin{cases} A_1(i_n) & \text{if } \bigwedge_{m=n+d-l}^{n+d-1} o_m \equiv A_1(o_m) \\ \dots & \dots \\ A_t(i_n) & \text{if } \bigwedge_{m=n+d-l}^{n+d-1} o_m \equiv A_t(o_m) \end{cases}$$

But this naive approach suffers from state space explosion problem. For O^l that contains m boolean variables, there may be 2^m satisfying assignments, which make it impossible to build f^{-1} for large m .

There exists many much more efficient approaches to enumerating satisfying assignments of SAT instance [4, 5, 6, 7, 8, 9, 10, 24]. Their basic idea has been introduced in subsection 2.2.

But they are still not efficient. The essential reason that leads to this is the wide usage of XOR gates in communication and arithmetic circuits. As explained in subsection 2.2, satisfying assignments of XOR can't be minimized.

So we invent a novel approach that is specially designed for XOR intensive circuits, as shown below:

ALGORITHM 1. Synthesizing Complementary variable sets

1. assume I_{var} and O_{var} are respectively boolean variables set that represent i_n and $\langle o_{n+d-l}, \dots, o_{n+d-1} \rangle$ and F_E is defined in formula (10)
2. $G_{XOR} \leftarrow \{\}$

parameters

```
3. foreach  $v \in I_{var}$  {
4.    $SA_v \leftarrow \{\}$ 
5.   while ( $F_E \wedge v \equiv 1$  is satisfiable) {
6.     Assume  $A$  is a satisfying assignment
7.      $A_{BFL} \leftarrow BFL(F_E, A, v)$ 
8.      $\{A_{XOR}, G\} \leftarrow XORMIN(F_E, A_{BFL}, v)$ 
9.      $SA_v \leftarrow SA_v \cup \{A_{XOR}\}$ 
10.     $G_{XOR} \leftarrow G_{XOR} \cup G$ 
11.     $F_E \leftarrow F_E \wedge bcl_{SA_{XOR}}$ 
12.     $F_E \leftarrow F_E \wedge \bigwedge_{(x_1 \leftarrow v_1 \oplus v_2) \in G} \{x_1 \equiv v_1 \oplus v_2\}$ 
13.  }
```

这指代不清？

```
14. Building  $f_v^{-1} : \{0, 1\}^{O_{var}} \rightarrow \{0, 1\}^{I_{var}}$  for  $v \in I_{var}$ 
15. }
16. Building  $f^{-1} : \{0, 1\}^{O_{var}} \rightarrow \{0, 1\}^{I_{var}}$ 
```

At line 2, G_{XOR} is a set of XOR gates discovered by XORMIN on line 8. They will help to speedup the process of enumerating satisfying assignments of $F_E \wedge v \equiv 1$ on line 5. More detail will be given in subsection 4.3.

At line 4, SA_v is a set of satisfying assignments that can make $v \equiv 1$, they will be used in line 14 to build f_v^{-1} . At line 5, we iterate through all satisfying assignments A of $F_E \wedge v \equiv 1$, and minimize them in two steps.

1. BFL in line 7 is a modified BFL[6] that will be described in subsection 4.2.

2. XORMIN in line 8 will further minimize the result of BFL by discovering hidden XOR gates, and return the minimized assignment A_{XOR} and discovered XOR gate set G . This is one of our major contributions, and will be described in subsection 4.3.

At line 11, we will rule out the enumerated satisfying assignment A_{XOR} by adding its blocking clauses into F_E .

At line 12, we will add clauses into F_E for newly discovered XOR.

4.2 Minimizing Satisfying Assignments with Modified BFL algorithm

The basic idea of BFL has been introduced in subsection 2.2, so we present the modified BFL here directly.

ALGORITHM 2. $BFL(F_E, A, v)$

1. Assume O_{var} is boolean variables set that represents $\langle o_{n+d-l}, \dots, o_{n+d-1} \rangle$
2. foreach $u \in O_{var}$ {
3. if ($F_E \wedge \neg v \wedge A|_{O_{var}-\{u\}}$ is unsatisfiable) {
4. $A \leftarrow A|_{O_{var}-\{u\}}$
5. $O_{var} \leftarrow O_{var} - \{u\}$
6. }
7. }
8. return A

In line 2, we iterate through all output variable $u \in O_{var}$. In line 3, if that formula is unsatisfiable, then $A|_{O_{var}-\{u\}}$ is sufficient to lead to $v \equiv 1$. Thus, we can remove the assignment of u from A .

According to subsection 2.2, for circuits with lots of XOR gates, BFL can't minimize its satisfying assignments. Thus, algorithm 1 can't terminate in reasonable time for large XOR intensive circuits.

From

4.3 Minimizing Satisfying Assignments by Discovering XOR gates

Intuitively, assume circuit in figure ??a) has input variables set $V = \{v_1, v_2, v_3, \dots, v_n\}$, and it MAY contain a XOR $v' \leftarrow v_1 \oplus v_2$. We can avoid enumerating the assignments of v_1 and v_2 by first checking the existence of this XOR, and then, as shown in figure ??b); if this XOR actually exists, we add another XOR $x_1 \leftarrow v_1 \oplus v_2$ into this circuit, and enumerate assignments on $V \cup \{x_1\} - \{v_1, v_2\}$ instead of V .

Formally, to check existence of XOR $v' \leftarrow v_1 \oplus v_2$, for a satisfying assignment A_F of formula F_E , we define a new assignment A_{x_1} in formula (12), by first removing assignments of v_1 and v_2 from A_F , and then adding assignment of x_1 as result of XORing v_1 and v_2 into A_F .

With this A_{x_1} , existence of $v' \leftarrow v_1 \oplus v_2$ can be decided by checking unsatisfiability of following formula :

(12)

Unsatisfiable of formula (13) means that A_{x_1} can't make v to be 0, so A_{x_1} must be a satisfying assignment of $F_E \wedge v \wedge \{x_1 \equiv v_1 \oplus v_2\}$.

Thus, A_F and $A_F|_{V-\{v_1, v_2\}}|_{v_1 \leftarrow \neg A_F(v_1)}|_{v_2 \leftarrow \neg A_F(v_2)}$, that can't be merged by BFL, have now been merged into A_{x_1} with the help of a newly discovered XOR gate $x_1 \leftarrow v_1 \oplus v_2$. If we repeatedly merge assignments by checking unsatisfiability of formula (13), we can get a partial assignment of $F_E \wedge v \wedge \{x_1 \equiv \dots \oplus \dots\} \dots \wedge \{x_n \equiv \dots \oplus \dots\}$ in formula (14), which contains 2^n total assignments, that can't be merged by BFL.

(14)

With above discussion in mind, we describe XORMIN below:

ALGORITHM 3. **XORMIN**(F_E, A_F, v)

```

1.  $G = \{\}$ 
2. do {
3.    $G_{new} = \{\}$  // the set of newly discovered XOR
4.   foreach  $v_1, v_2 \in O_{var}$  {
5.     if(formula (13) is unsatisfiable){
6.        $G_{new} \leftarrow G_{new} \cup \{x_1 \leftarrow v_1 \oplus v_2\}$ 
7.        $A_F \leftarrow A_F|_{O_{var}-\{v_1, v_2\}}|_{x_1 \leftarrow A_F(v_1) \oplus A_F(v_2)}$ 
8.        $O_{var} \leftarrow O_{var} \cup \{x_1\} - \{v_1, v_2\}$ 
9.        $F_E \leftarrow F_E \wedge bcls_{A_F}$ 
10.       $F_E \leftarrow F_E \wedge \bigwedge_{\{x_1 \leftarrow v_1 \oplus v_2\} \in G_{new}} \{x_1 \equiv v_1 \oplus v_2\}$ 
11.    }
12.  }
13.   $G = G \cup G_{new}$ 
14. } while( $G_{new} \neq \{\}$ )
15. return  $\{A_F, G\}$ 

```

In line 1, G is a set that will be used to hold all XOR gates discovered by this algorithm.

In line 2, the do-while statement will repeatedly discover new XOR gates, until no more XOR gates can be discovered.

In line 4, foreach statement will enumerate each pair of $v_1, v_2 \in V$, and line 5 will test if there is a XOR gate between v_1 and v_2 .

4.4 Building f_v^{-1} and f^{-1}

The complementary function $f^{-1} : O^I \rightarrow I$ is the function that takes $\langle o_{n+d-l}, \dots, o_{n+d-1} \rangle$, and computes i . Assume i_n is represented by boolean variables set I_{var} , and $\langle o_{n+d-l}, \dots, o_{n+d-1} \rangle$ is represented by boolean variables set O_{var} .

Thus, f_v^{-1} main is $f_v^{-1} : \{0, 1\}^{O_{var}} \rightarrow \{0, 1\}^{I_{var}}$. Then building f_v^{-1} can be partitioned into multiple tasks, each task builds a boolean function $f_v^{-1} : \{0, 1\}^{O_{var}} \rightarrow \{0, 1\}$ for a $v \in I_{var}$.

To build f_v^{-1} , let's assume that SA_v is the set of satisfying assignments that makes v to take on 1. Then f_v^{-1} can be defined as :

$$f_v^{-1}(x) = \begin{cases} 1 & \exists A \in SA_v. st. x \equiv A(x) \\ 0 & else \end{cases}$$

5. EXPERIMENTAL RESULTS

5.1 Benchmarks

Our approach is the first one that can synthesize complementary circuits automatically, so we can't compare it with other's research results.

Temporal logic synthesis is a research topic that is somewhat close to us, but it can't be scaled to large circuits, and no commercial available IP cores are written in temporal logic. Thus, it's impossible to compare our result with temporal logic synthesis.

Table 1 shows some information of following benchmarks.

1. The first benchmark is a XGXS encoder that is compliant to clause 48 of IEEE-802.3ae 2002 standard [11].
2. The second benchmark is a XFI encoder that is compliant to clause 49 of the same IEEE standard.
3. The third benchmark is a 66 bit scrambler that is used to make a data sequence to have enough transitions between 0 and 1, such that it can run through high speed noisy serial transmission channel.
4. The fourth benchmark is a PCIE physical coding module.
5. The fifth benchmark is Ethernet module of Sun's OpenSparc T2 processor.

5.2 Experimental Results

Table 2 shows the run-time of checking parameterized complementary assumption on these circuits, and the discovered proper values of parameters.

Table 3 shows the results of generated verilog description of complementary circuits. With BFL only, complementary circuits of the three most complex circuits: XFI, PCIE

and T2 ethernet, can't be built within 10,000 seconds. While with **BFL** and **XORMIN**, we can finally build all complementary circuits successfully within 3000 seconds.

These generated complementary circuits are all been verified by dynamic simulations.

6. RELATED WORKS

6.1 Temporal Logic Synthesis

Automatic synthesis of program from logic specification is first identified as Church's problem at 1962[12]. Some early researches [13, 14] solve this problem by reducing it to checking emptiness of two automata. After invention of SAT at early 1980s, this problem has been considered again [15, 16]. But at 1980, A. Pnueli and P. Rosner[17] pointed out complexity of LTL synthesis is double exponential in the size of the formula.

This high complexity drives researchers turning their focus to find smaller but still useful subset logic, such that synthesis problem can be solved with lower complexity.

One line of research [18, 19, 20] focus on the so-called generalized reactive formulas of the form: $(\Box \Diamond p_1 \wedge \dots \Box \Diamond p_m) \rightarrow (\Box \Diamond q_1 \wedge \dots \Box \Diamond q_n)$. Complexity of solving synthesis problem for such formula is $O(N^3)$.

The other line of research focus on finding efficient symbolic algorithm [23] for expensive symbolic determination algorithm [21] on a useful formula subset, or just avoiding it[22].

6.2 Satisfying Assignments Enumeration

There are two research directions for satisfying assignments enumeration. One is bottom up, which removes irrelevant variables from total assignments. The other is top down, which adds relevant variable's assignments into an empty set, to form a smaller and smaller assignment set.

The first bottom up approach is proposed by K. L. McMillan [5]. He constructs an alternative implication graph in SAT solver, which records the reasoning relation that leads to the assignment of a particular object variable *obj*. All variables outside this graph can be ruled out from the total assignment. Kavita Ravi et al.[6] and P. P. Chauhan et al.[10] remove those variables that can make *obj* $\equiv 0$ satisfiable one by one. Shen et al.[8] and HoonSang Jin et al.[4, 7] use a conflict analysis based approach, to remove multiple irrelevant variables in one SAT run. Orna Grumberg et al.[9] separates the variables set into important subset and non-important subset. Variables in important subset have higher decision priority than non-important ones. Thus, the important subset form a search tree, each leaf is another search tree for non-important set.

Cofactoring [24] is the only top down approach, which starts from an empty set of assignments, and add relevant variable's assignments one by one.

6.3 AND-XOR Logic Synthesis

Classical logic synthesis works on AND-OR network. Its kernel is two-level logic minimization, which tries to find a small sum-of-products expression for boolean function *f*. It is obvious that such two-level logic minimization algorithms are similar to satisfying assignments enumeration described in previous subsection, except that they don't work on SAT solvers.

Three most well known two-level logic minimization algorithms are Quine-McCluskey[26], Scherzo[27], and Espresso-

II[28].

Just like state-of-the-art ALLSAT that can't deal with XOR-intensive circuits efficiently, classical logic synthesis also has the same problem with many researches propose synthesis algorithms that target XOR-intensive circuit.

One research direction focus on extending classical two-level AND-OR minimization to describe D-XOR network [29, 30]. These works normally express circuits with most general ESOP (exclusive sum of product) expressions. But very large-scaled complexity of these approaches prevents them from handling large circuits.

Another line of research relies on Reed-Muller expansion[31], one of its most used variant is Fixed Polarity Reed-Muller Form (FPRM) given by Davio and Deschamps[32], in which a variable can have either positive or negative polarity. Some related works that rely on FPRM are [33, 34, 35].

7. CONCLUSIONS AND FUTURE WORKS

In this paper, we propose the first fully automatic approach that synthesizes complementary circuits for communication applications. According to experimental results, our approach can synthesize correct complementary circuits for many very complex circuits, including but not limited to PCIE and Ethernet.

One possible future work is to automatically generate assertions that rule out invalid input data patterns, such that the users can be free from the burden of inspecting documentation and writing assertions.

Another possible future work is to deal with circuits with memory array and multiple clocks, such that more complex communication mechanism, such as data link layer and transaction layer, can be deal with by our approach.

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