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377 Operating Systems

Process Abstraction and API





THE CRUX OF THE PROBLEM: HOW TO PROVIDE THE ILLUSION OF MANY CPUS? Although there are only a few physical CPUs available, how can the OS provide the illusion of a nearly-endless supply of said CPUs?



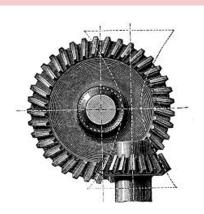
Virtualizing the CPU

Time Sharing

- The OS can promote the illusion that many virtual CPUs exist.
- o The OS runs one process, stops it, then runs another, and so on.

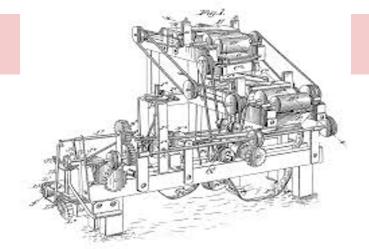


- Low-level machinery and high-level intelligence.
- Mechanisms: low-level methods implementing a needed piece of machinery.
 - Context Switch
- Policies: algorithms for making decisions within the O.
 - Which program should the OS run?
 - Scheduling Policies
 - Historical information, workload knowledge, and performance metrics





The Abstraction: A Process



A Process

• An instance of a running program.

At any moment in time, we can summarize a process by taking an inventory of the different pieces of the system it <u>accesses</u> or <u>effects</u> during execution.

Machine State

- What parts of the machine are important to the execution of the program.
- What can a program read or update when it is running?
 - **Memory:** data & instructions, also known as its <u>address space</u>.
 - **Registers:** program counter (PC), stack pointer, frame pointer, etc.
 - Persistent Storage: I/O information such as open files.

- Create
- Destroy
- Wait
- Miscellaneous Control
- Status

- Create: An operating system must include some method to create new processes. When you type a command into the shell, or double-click on an application icon, the OS is invoked to create a new process to run the program you have indicated.
- Destroy
- Wait
- Miscellaneous Control
- Status

- Create
- Destroy: As there is an interface for process creation, systems also provide an interface to destroy processes forcefully. Of course, many processes will run and just exit by themselves when complete; when they don't, however, the user may wish to kill them, and thus an interface to halt a runaway process is quite useful.
- Wait
- Miscellaneous Control
- Status

- Create
- Destroy
- Wait: Sometimes it is useful to wait for a process to stop running; thus some kind of waiting interface is often provided.
- Miscellaneous Control
- Status

- Create
- Destroy
- Wait
- Miscellaneous Control: Other than killing or waiting for a process, there are sometimes other controls that are possible. For example, most operating systems provide some kind of method to suspend a process (stop it from running for a while) and then resume it (continue it running).
- Status

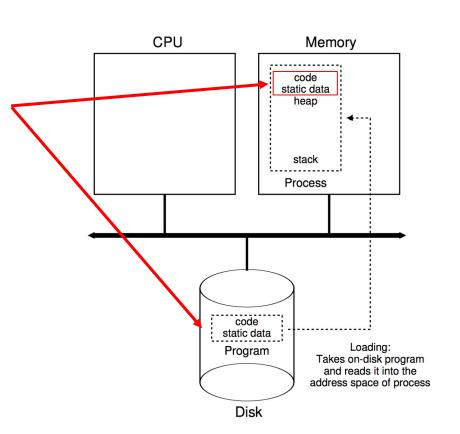
- Create
- Destroy
- Wait
- Miscellaneous Control
- Status: There are usually interfaces to get some status information about a process as well, such as how long it has run for, or what state it is in.

Process Creation: Loading

The first thing the OS must do is load its code and any static data into memory - the <u>address space</u> of the process.

Early operating systems performed this task <u>eagerly</u>, whole program.

A modern OS performs this <u>lazily</u>, by loading pieces of code and data as they are needed.

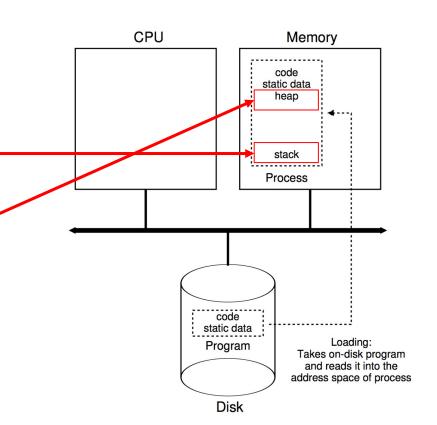


Process Creation: Memory Allocation

The next step is to allocate important memory regions for the process.

Stack: for storing local variables and function parameters and return values.

Heap: for dynamically allocated data.

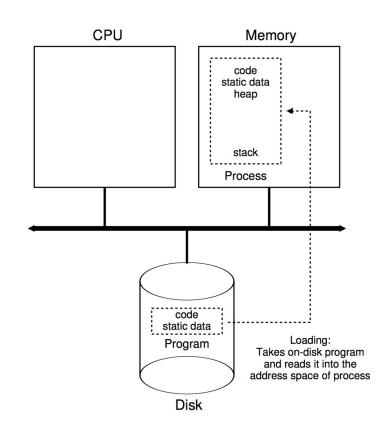


Process Creation: Initialization

The OS will also do some other <u>initialization</u> <u>tasks</u>, particularly as related to input/output (I/O).

For example, in UNIX systems, each process by default has three open file descriptors, for <u>standard input</u>, <u>output</u>, and <u>error</u>

these descriptors let programs easily read input from the terminal as well as print output to the screen.

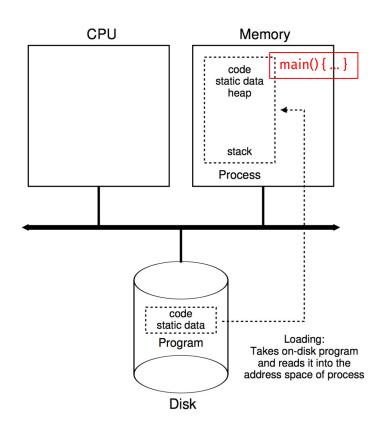


Process Ready

The OS has now set the stage for program execution. It thus has one last task: to start the program running at the entry point.

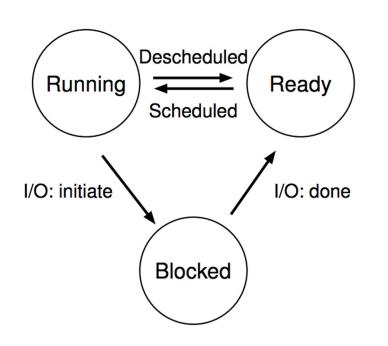
Historically, a special function called main().

By "jumping" to the main() routine, the OS transfers control of the CPU to the newly-created process, and thus the program begins its execution.



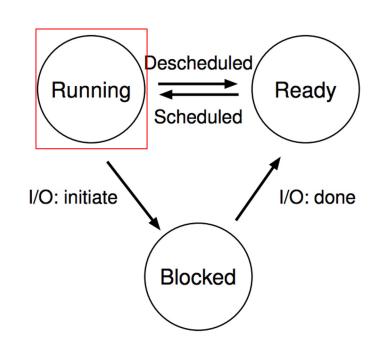
In a simplified view a process can be in one of three states:

- Running
- Ready
- Blocked



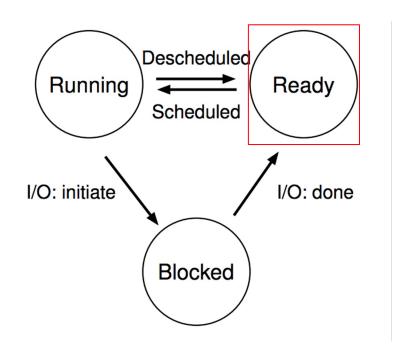
In a simplified view a process can be in one of three states:

- Running: In the running state, a process is running on a processor. This means it is executing instructions.
- Ready
- Blocked



In a simplified view a process can be in one of three states:

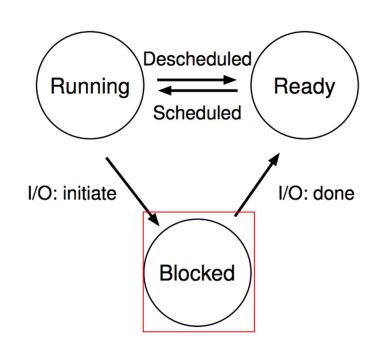
- Running
- Ready: In the ready state, a process is ready to run but for some reason the OS has chosen not to run it at this given moment.
- Blocked



In a simplified view a process can be in one of three states:

- Running
- Ready
- Blocked: In the blocked state, a process has performed some kind of operation that makes it not ready to run until some other event takes place.

Example: when a process initiates an I/O request to a disk, it becomes blocked and thus some other process can use the processor.



Time	$\mathbf{Process}_0$	$\mathbf{Process}_1$	Notes
1	Running	Ready	
2	Running	Ready	
3	Running	Ready	
4	Running	Ready	Process ₀ now done
5	_	Running	
6	_	Running	
7	_	Running	
8	_	Running	Process ₁ now done

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Data Structures

Operating Systems are rich in data structures!

To provide an abstraction for a processes we require a data structure!

Process Control Block: C structure that contains information about each process.

Process List: Keeps track of all processes (PCBs) in the running system.

Example: Xv6* OS Proc Structure

struct context {

```
int eip;
  int esp;
  int ebx;
  int ecx;
  int edx;
  int esi;
  int edi;
  int ebp;
// the different states a process can be in
enum proc_state { UNUSED, EMBRYO, SLEEPING,
                  RUNNABLE, RUNNING, ZOMBIE };
```

// to stop and subsequently restart a process

// the registers xv6 will save and restore * Xv6 is a UNIX-like OS from MIT.

Example: Xv6 OS Proc Structure

```
// the information xv6 tracks about each process
// including its register context and state
struct proc {
                            // Start of process memory
 char *mem;
                           // Size of process memory
 uint sz;
 char *kstack;
                         // Bottom of kernel stack
                            // for this process
 enum proc_state state;  // Process state
 int pid;
                         // Process ID
 struct proc *parent; // Parent process
                       // If non-zero, sleeping on chan
 void *chan;
 int killed;
                            // If non-zero, have been killed
 struct file *ofile[NOFILE]; // Open files
 struct inode *cwd; // Current directory
 struct context; // Switch here to run process
 struct trapframe *tf; // Trap frame for the
                            // current interrupt
```

CRUX: HOW TO CREATE AND CONTROL PROCESSES

What interfaces should the OS present for process creation and control? How should these interfaces be designed to enable ease of use as well as utility?

The End