Haskell or functional purity and laziness

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Preliminaries

```
{-# LANGUAGE BangPatterns #-}
module Intro where
import Test.QuickCheck
import Test.HUnit
import Text.Show.Functions
import Control.DeepSeq
```

Purity

Purity

Pure computations always result in the same value given the same inputs, meaning they can not perform side effects such as mutation of assignables or I/O.

Purity

```
var impure = function(x, y) {
  console.log("foobar");
  return x + y;
}
var pure = function(x, y) {
  return x + y;
}
```

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Haskell functions are curried by default

```
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(+) :: Num a => a -> a -> a

Prelude> :t uncurry (+)
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uncurry (+) :: Num c => (c, c) -> c
```

 Juxtaposition is function application and has highest fixity. Use \$ for least fixity.

```
Prelude> (*) 5 $ 3 + 2
25
Prelude> (*) 5 3 + 2
```

► Anonymous functions can be defined via the syntax

$$\xspace x -> x + 1$$

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► Function definition

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add1 x = x + 1
add1' = (+ 1)
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Ranges and list comprehensions

```
twoToFour = [2..4]
twoToFour' = [x + 1 | x <- [1..3]]</pre>
```

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Ranges and list comprehensions

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twoToFour = [2..4]
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```

▶ HUnit tests

```
map_test = map add1 [1,2,3] ~=? twoToFour
map_test' = map add1' [1,2,3] ~=? twoToFour'
```

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▶ Pattern matching allows one to define functions by (exhaustive) cases

```
doubleVision :: [a] -> [a]
doubleVision (x:xs) = x : x : doubleVision xs
doubleVision [] = []
vision_test = doubleVision [1,2,3] ~=? [1,1,2,2,3,3]
```

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```

► Guards enable more flexible checking in pattern matching

```
abs' x
| x >= 0 = x
| x < 0 = -x</pre>
```

Interlude - QuickCheck

Quickcheck property tests
abs_check :: Integer -> Bool
abs_check x = abs' x >= 0

Interlude - QuickCheck

```
Prelude> :load haskell.lhs
[1 of 1] Compiling Intro
Ok, modules loaded: Intro.
*Intro> quickCheck mymap_check
+++ OK, passed 100 tests.
Bool
abs_check :: Integer -> Bool
abs_check :: abs' x >= 0

Prelude> :load haskell.lhs
( haskell.lhs, interpreted )
```

► Let bindings

```
isPalindrome x = let rev = reverse x in
  rev == x
palindrome_test = isPalindrome "abba" ~=? True
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Where bindings

Recursion and Higher Order Functions

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Higher order functions

Here's how you implement map

```
myMap :: (a -> b) -> [a] -> [b]
myMap f (x:xs) = f x : myMap f xs
myMap f [] = []

myMap_check :: [Integer] -> Bool
myMap_check x = myMap (+1) x == map (+1) x
```

Counting to infinity

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```
infty = [1..]
```

► Forcing the issue?

```
finiteSlice = take 100 infty
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Forcing the issue?

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Don't try to reduce me to normal form..

```
take' n x = x `deepseq` take n x
loooop :: [Integer]
loooop = take' 10 [1..]
```

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```

Space leaks, or "Oh my god, it's full of stars.." foldl (+) 0 [1..10000000]

► This is pretty fast import Data.List foldl' (+) 0 [1..10000000]

Defining a strict(ish) fold using !
foldl'' f start xs = run start xs where
run !acc (x:xs) = run (f acc x) xs
run !acc [] = acc

Typeclasses and (Inductive) Data Types

Typeclasses

▶ Problem - Ad hoc polymorphism:

```
isEqual :: ? a -> a -> Bool
isEqual x y = ?
```

Typeclasses

Problem - Ad hoc polymorphism:

```
isEqual :: ? a -> a -> Bool
isEqual x y = ?
```

Solution:

```
class Equality a where
   eq :: a -> a -> Bool

isEqual :: (Equality a) => a -> a -> Bool
isEqual x y = x `eq` y
```

Algebraic Data Types

Defining new data types as sums of products

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▶ Defining new data types as sums of products

Using them

Defining instances

Custom instances

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Custom instances

instance Equality Boolean where

```
eq BTrue BTrue = True
    eq BFalse BFalse = True
    eq _ _ = False
  instance_tests = [BTrue `eq` BTrue ~=? True,
                    BTrue 'eq' BFalse ~=? False]
Nesting
  instance Equality a => Equality (ConsList a) where
    eq x y = isEq True x y where
      isEq True CNil CNil = True
      isEq True (Cons x xs) (Cons y ys) = isEq (x `eq` y) xs ys
      isEq _ _ = False
  instance_tests' = [Cons BTrue CNil `eq` Cons BTrue CNil ~=? True,
                     Cons BTrue CNil `eq` CNil ~=? False]
```

Records

Defining a carrot

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Accessors are auto generated

```
*Intro> :t color
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 Most obnoxious thing in all of Haskell: record names must be unique within a given module Functors, Idioms, Monads .. also Monoids

Motivation: You want to combine some things, like strings combineStrings xs = foldl (++) "" xs But suddenly you have to combine (Option)al strings and would prefer not to repeat yourself

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▶ Monoids to the rescue

```
class Monoid a where
  mempty :: a
  mappend :: a -> a -> a
  mconcat :: [a] -> a
  mconcat = foldr mappend mempty
```

► Instantiating..

```
instance Monoid a => Monoid (Option a) where
  mempty = None
  mappend (Some x) (Some y) = Some (x `mappend` y)
  mappend None x = x
  mappend x None = x

-- strings are lists of characters
instance Monoid [a] where
  mempty = []
  mappend = (++)
```

```
combineThings :: Monoid a => [a] -> a
combineThings xs = mconcat xs

combining_tests = [
   combineThings ["foo", "bar", "baz"] ~=?
   "foobarbaz",
   combineThings [Some "foo", None, Some "bar"] ~=?
   Some "foobar"]
```

► Classic problem:

```
getFromMaps x y =
  case lookup "foo" x of
   Just val -> case lookup "bar" y of
   Just val2 -> Just (val ++ val2)
   Nothing -> Nothing
  Nothing -> Nothing
-- ARGH
```

Classic problem:

```
getFromMaps x y =
   case lookup "foo" x of
   Just val -> case lookup "bar" y of
   Just val2 -> Just (val ++ val2)
   Nothing -> Nothing
   Nothing -> Nothing
-- ARGH
Or even just:
```

Nothing -> Nothing

Enter functors

```
doSomethingMaybe' x = fmap (++ "bar") x
class Functor f where
  fmap :: (a -> b) -> f a -> f b
instance Functor Option where
  fmap f (Some x) = Some (f x)
  fmap f None = None
```

Enter functors

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Option has kind * → *, which means it forms a concrete type of kind * only when applied to a type argment.

Idioms - Applicative Functors

▶ Queue Strauss' Sunrise - What if I could just write this getFromMaps' x y = (++) <\$> lookup "foo" x <*> lookup "bar" y

Idioms - Applicative Functors

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Applicative functors

```
class (Functor f) => Applicative f where
  pure :: a -> f a
  (<*>) :: f (a -> b) -> f a -> f b

instance Applicative Option where
  pure = Some
  None <*> _ = None
  Some f <*> v = fmap f v

(<$>) :: (Functor f) => (a -> b) -> f a -> f b
f <$> x = fmap f x
```

► Even though they have many, many other uses, let's look at I/O. Idea: type IO a = RealWorld -> (a, RealWorld)

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- ▶ Applicatives can't guarantee ordering of effects, we have to make sure that things are evaluated and bound in the right sequence.

Even though they have many, many other uses, let's look at I/O. Idea: type IO a = RealWorld -> (a, RealWorld)

- Applicatives can't guarantee ordering of effects, we have to make sure that things are evaluated and bound in the right sequence.
- The IO Monad allows one to write

```
getInputAndPrint :: IO ()
getInputAndPrint = do
  putStrLn "Give me some input"
  input <- getLine
  putStrLn input</pre>
```

making sure that the actions are performed in the specified sequence

▶ So how does this work? Let's start with a definition:

```
class (Applicative m) => Monad m where
  return :: a -> m a
  (>>=) :: m a -> (a -> m b) -> m b
  -- alternatively join :: m (m a) -> m a
```

▶ So how does this work? Let's start with a definition:

```
class (Applicative m) => Monad m where
  return :: a -> m a
  (>>=) :: m a -> (a -> m b) -> m b
-- alternatively join :: m (m a) -> m a
```

Option instance

```
instance Monad Option where
  return = Some
  Some x >>= f = f x
  None >>= f = None
```

Since this doesn't look so nice, introduce some sugar

```
getFromMaps''' x y = do
  val1 <- lookup "foo" x
  val2 <- lookup "bar" y
  return $ val1 ++ val2</pre>
```

fs <- [(+1), (+2)] return \$ fs x [2,3,3,4,4,5]

How to use it? getFromMaps', x y = lookup "foo" x >>= (\val1 -> (lookup "bar" y >>= (\val2 -> return \$ val1 ++ val2))) Since this doesn't look so nice, introduce some sugar getFromMaps'', x y = do val1 <- lookup "foo" x</pre> val2 <- lookup "bar" y return \$ val1 ++ val2 List is also a monad do $x \leftarrow [1,2,3]$

▶ That was easy, so let's move on to State..

```
newtype State' s a = State' { runState :: s -> (a, s) }
instance Monad (State' a) where
return x = State' $ \s -> (x, s)
oldState >>= f = State' $ \s ->
let (intermediateVal, intermediateState) = runState oldState s
in runState (f intermediateVal) intermediateState
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```

Utilities

```
put :: a -> State' a ()
put s = State' $ \_ -> ((), s)
get = State' $ \s -> (s, s)
```

► [1..] rewritten infty' = go 1 where go n = n : go (n + 1)

▶ [1..] rewritten

```
infty' = go 1 where
  go n = n : go (n + 1)
```

An infinite list of fibonacci numbers

```
fibs = 1 : 1 : rec fibs where
  rec (x:y:xs) = (x + y) : rec (y : xs)
cofibonacci n = head $ drop (n - 1) $ take n fibs
```

▶ [1..] rewritten

```
infty' = go 1 where
  go n = n : go (n + 1)
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▶ The blurred distinction between data and codata

```
Prelude> head $ 1 : undefined
1
```

Existential Types

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Join the Haskell meetup!: http://www.meetup.com/NY-Haskell/

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- Grab the slides and run make extract to produce plain Haskell code or runhaskell haskell.lhs to just run it
- ► To play around with it in ghci use

 Prelude>:load haskell.lhs