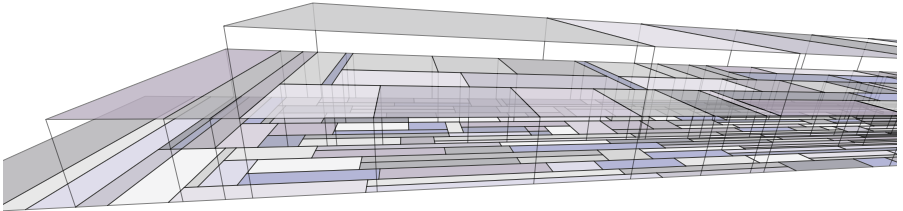




Aalto-yliopisto

*Lectio praecursoria, October 22, 2010*

# Accessing Multiversion Data in Database Transactions



Tuukka Haapasalo  
tuukka.haapasalo@tkk.fi

# Background on databases

# Background on databases

## What are databases?

### Databases

- Data storage for programs
- Examples:
  - Addresses in an address book
  - Bank account information
  - Calendar events
  - ...
  - Images, videos, music



# Background on databases

What are databases used for?

## Information retrieval

- Quick access to information
- Analogies: document archives, libraries
  - Books ordered by the author (in a library)
  - Easy to locate a certain book
  - Easy to locate all books written by a given author

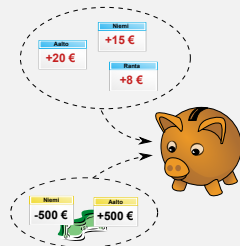


# Background on databases

## How are databases used?

### Transactions

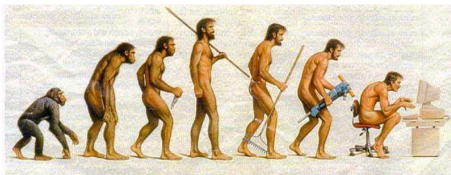
- Data are modified and accessed in transactions
- Atomicity
- Multiple concurrent updates
  - State remains consistent
  - Index structure remains intact



# Multiversion databases

# Multiversion databases

## Difference to traditional databases?



## Multiversion databases

- Store the evolution of data
- What information was stored before?
- Examples:
  - What documents the archive consisted of when it was created?
  - Who were the users of the system on the June 6th, 2010?
  - What was the balance of Mr. X's bank account a month ago?

# Modeling versioned data



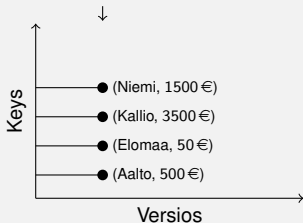
# Modeling versioned data

## Data model

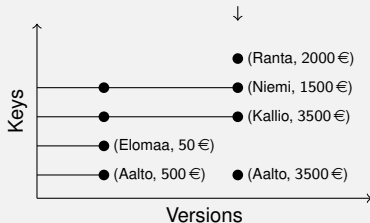
### Multiversion databases

Dissertation p. 12

Any change creates a new version (state) from the records of the index:



(a) Before modifications



(b) After modifications

Queries can target previous versions in addition to the current version.

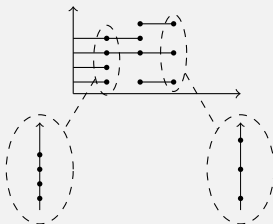
# Modeling versioned data

## Efficient queries

### Optimality

Dissertation p. 35

A multiversion database index is **optimal**, if querying a **version  $v$**  is as efficient as in a **single-version** database index that only indexes the data items of the given **version  $v$** .



# Modeling versioned data

## Efficient multiversion indexes

### ■ **Time-split $B^+$ -tree (TSB-tree); Lomet and Salzberg [4]**

- The first efficient multiversion index (1989)
- Not optimal
- Dissertation p. 55

### ■ **Multiversion $B^+$ -tree (MVBT); Becker et al. [1, 2]**

- Second efficient multiversion index (1993–1996)
- Optimal
- Each update creates a new version
- Dissertation p. 61

### ■ **Multiversion access structure (MVAS); Varman and Verma [5]**

- Third efficient multiversion index (1997)
- Optimal according to a different (not so strict) definition
- Each update creates a new version
- Dissertation p. 69

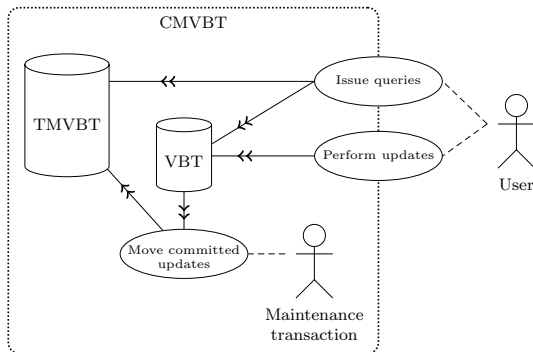
# Our research

### Contributions

- 1 Transactions to the MVBT index: the *transactional MVBT* (TMVBT)
  - Only a single updating transaction at a time
  - As efficient as the MVBT
  - Dissertation p. 75
- 2 The *concurrent MVBT* (CMVBT) for concurrent updating transactions
  - $\text{CMVBT} = \text{TMVBT} + \text{VBT}$
  - VBT = a versioned  $B^+$ -tree
  - Dissertation p. 111
- 3 Experimental evaluation
  - CMVBT is as efficient as the TSB-tree in the general situation
  - CMVBT is more efficient for key-range queries
  - Dissertation p. 137

### Concurrent multiversion $B^+$ -tree [3]

Dissertation p. 113



# Experimental evaluation

# Experimental evaluation

## Tested index structures

### Indexes we have evaluated

- **CMVBT** index
- **TSB-tree**

### Also implemented

- **TMVBT** index (one transaction at a time)
- **VBTree** index (alone)



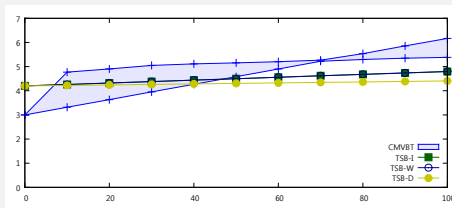
# Experimental evaluation

## Queries and updates

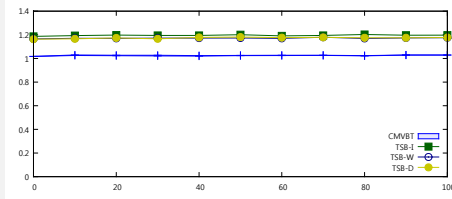
### Queries and updates, short transactions

Dissertation p. 147

Page fixes



Page reads from disk



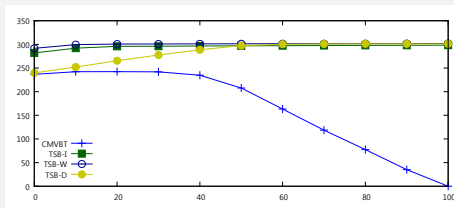
# Experimental evaluation

## Key-range queries

### Key-range queries

Dissertation p. 151

Page reads from disk (almost identical to the number of page fixes in this test):



# Summary

## Index structures

- **TMVBT** = transactional, optimal MVBT
- **CMVBT** = **TMVBT** + **VB**T
- Multiple updating transactions can operate on the **CMVBT** concurrently

## Kokeelliset tulokset

- **CMVBT** is as efficient as the **TSB-tree** in the general situation
- **CMVBT** is more efficient than the **TSB-tree** in key-range queries
- **CMVBT** takes 10–60 % more space than the **TSB-tree**

- [1] B. Becker, S. Gschwind, T. Ohler, B. Seeger, and P. Widmayer. On optimal multiversion access structures. In *Proceedings of the 3rd International Symposium on Advances in Spatial Databases*, pages 123–141, 1993.
- [2] B. Becker, S. Gschwind, T. Ohler, B. Seeger, and P. Widmayer. An asymptotically optimal multiversion B-tree. *The VLDB Journal*, 5(4):264–275, 1996.
- [3] T. Haapasalo, I. Jaluta, S. Sippu, and E. Soisalon-Soininen. Concurrent updating transactions on versioned data. In *Proceedings of the 2009 International Database Engineering and Applications Symposium*, pages 77–87, September 2009.
- [4] D. Lomet and B. Salzberg. Access methods for multiversion data. In *Proceedings of the 1989 ACM SIGMOD International Conference on Management of Data*, pages 315–324, 1989.
- [5] P. J. Varman and R. M. Verma. An efficient multiversion access structure. *IEEE Transactions on Knowledge and Data Engineering*, 9(3):391–409, 1997.